## Exercises for Lecture 2

Proof theory for modal and non-classical logics

June 2023

Rules of NK

$$\begin{array}{ccc} \operatorname{init} \frac{}{\Gamma\{p,\bar{p}\}} & \vee_{\mathsf{R}} \frac{\Gamma\{A,B\}}{\Gamma\{A\vee B\}} & \wedge_{\mathsf{R}} \frac{\Gamma\{A\} & \Gamma\{B\}}{\Gamma\{A\wedge B\}} \\ & & & \\ & \square_{\mathsf{R}} \frac{\Gamma\{[A]\}}{\Gamma\{\square A\}} & & \diamond_{\mathsf{R}} \frac{\Gamma\{\diamondsuit A,[A,\Delta]\}}{\Gamma\{\diamondsuit A,[\Delta]\}} \end{array}$$

Rules of **NKX**, for *X* 45-closed and  $X \subseteq \{D, T, B, 4, 5\}$ 

$$\mathrm{d} \frac{\Gamma\{\diamondsuit A, [A]\}}{\Gamma\{\diamondsuit A\}} \qquad \mathrm{t} \frac{\Gamma\{\diamondsuit A, A\}}{\Gamma\{\diamondsuit A\}} \qquad \mathrm{b} \frac{\Gamma\{[\Delta, \diamondsuit A], A\}}{\Gamma\{[\Delta, \diamondsuit A]\}}$$
 
$$4\frac{\Gamma\{\diamondsuit A, [A, \Delta]\}}{\Gamma\{\diamondsuit A, [\Delta]\}} \qquad 5\frac{\Gamma\{\diamondsuit A\}\{\diamondsuit A\}}{\Gamma\{\diamondsuit A\}\{\emptyset\}} \ depth(\Gamma\{\}\{\emptyset\}) > 0$$

Calculi **NKX**, for *X* 45-closed and  $X \subseteq \{D, T, B, 4, 5\}$ , are defined by adding to **NK** the rules corresponding to the letters in *X*. For instance, **NKDB** is **NK** + d + b. We denote by  $\vdash_{\mathbf{NKX}}$  derivability of sequent  $\Delta$  in the nested calculus **NKX**.

**Exercise 1.** Prove the following:

$$a) \vdash_{\mathbf{NK}} \Diamond (A \lor B) \supset (\Diamond A \lor \Diamond B)$$

$$b) \vdash_{\mathbf{NK}} \Diamond (A \supset B) \supset \Box (A \supset B)$$

- $c) \vdash_{\mathbf{NKT}} \Box p \supset p$
- $d) \vdash_{\mathbf{NKB}} p \supset \Box \Diamond p$
- $e) \vdash_{\mathbf{NK4}} \Box p \supset \Box \Box p$

 $f) \vdash_{\mathbf{NK5}} \Diamond p \supset \Box \Diamond p$ 

## Exercise 2. (optional)

Have a look at the countermodel construction from a failed proof search detailed in the following paper (subsection "Completeness", Theorem 3):

Brünnler, Deep sequent systems for modal logic. Arch. Math. Logic 48, 551–577 (2009). https://link.springer.com/article/10.1007/s00153-009-0137-3

Construct a Kripke model  $\mathcal{M}$  containing a world x such that  $\mathcal{M}, x \not\models p \supset \Box \Diamond p$  from a failed branch of a proof search tree for the formula in  $\mathbf{N}\mathbf{K}$ , applying the strategy from the paper. Hint: the strategy can be simplified, as proof search in  $\mathbf{N}\mathbf{K}$  always terminates.