

## Exercises for Lecture 2

Proof theory for modal and non-classical logics

June 2023

Rules of **NK**

$$\begin{array}{c} \text{init} \frac{}{\Gamma\{p, \bar{p}\}} \quad \vee \frac{\Gamma\{A, B\}}{\Gamma\{A \vee B\}} \quad \wedge \frac{\Gamma\{A\} \quad \Gamma\{B\}}{\Gamma\{A \wedge B\}} \\[10pt] \Box \frac{\Gamma\{[A]\}}{\Gamma\{\Box A\}} \quad \Diamond \frac{\Gamma\{\Diamond A, [A, \Delta]\}}{\Gamma\{\Diamond A, [\Delta]\}} \end{array}$$

Rules of **NKX**, for  $X$  45-closed and  $X \subseteq \{D, T, B, 4, 5\}$

$$\begin{array}{c} \text{d} \frac{\Gamma\{\Diamond A, [A]\}}{\Gamma\{\Diamond A\}} \quad \text{t} \frac{\Gamma\{\Diamond A, A\}}{\Gamma\{\Diamond A\}} \quad \text{b} \frac{\Gamma\{[\Delta, \Diamond A], A\}}{\Gamma\{[\Delta, \Diamond A]\}} \\[10pt] \text{4} \frac{\Gamma\{\Diamond A, [\Diamond A, \Delta]\}}{\Gamma\{\Diamond A, [\Delta]\}} \quad \text{5} \frac{\Gamma\{\Diamond A\}\{\Diamond A\}}{\Gamma\{\Diamond A\}\{\emptyset\}} \text{depth}(\Gamma\{\}\{\emptyset\}) > 0 \end{array}$$

Calculi **NKX**, for  $X$  45-closed and  $X \subseteq \{D, T, B, 4, 5\}$ , are defined by adding to **NK** the rules corresponding to the letters in  $X$ . For instance, **NKDB** is **NK** + d + b. We denote by  $\vdash_{\mathbf{NKX}}$  derivability of sequent  $\Delta$  in the nested calculus **NKX**.

**Exercise 1.** Prove the following:

- a)  $\vdash_{\mathbf{NK}} \Diamond(A \vee B) \supset (\Diamond A \vee \Diamond B)$
- b)  $\vdash_{\mathbf{NK}} \Diamond(A \supset B) \supset \Box(A \supset B)$
- c)  $\vdash_{\mathbf{NKT}} \Box p \supset p$
- d)  $\vdash_{\mathbf{NKB}} p \supset \Box \Diamond p$
- e)  $\vdash_{\mathbf{NK4}} \Box p \supset \Box \Box p$

f)  $\vdash_{\mathbf{NK5}} \Diamond p \supset \Box \Diamond p$

**Exercise 2.** (optional)

Have a look at the countermodel construction from a failed proof search detailed in the following paper (subsection "Completeness", Theorem 3):

Brnnler, Deep sequent systems for modal logic. Arch. Math. Logic 48, 551577 (2009). <https://link.springer.com/article/10.1007/s00153-009-0137-3>

Construct a Kripke model  $\mathcal{M}$  containing a world  $x$  such that  $\mathcal{M}, x \not\models p \supset \Box \Diamond p$  from a failed branch of a proof search tree for the formula in  $\mathbf{NK}$ , applying the strategy from the paper. Hint: the strategy can be simplified, as proof search in  $\mathbf{NK}$  always terminates.