

Course on Proof Theory - Lecture 2

# Exercises - Soundness, completeness and other metalogical results

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Gianluca Curzi, Marianna Girlando

University of Birmingham

Midlands Graduate School  
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## Exercises for Lecture 2

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1. Show that axiom HF5. is valid ( $x \notin \text{FV}(A)$ ):

$$(\forall x(A \rightarrow B)) \rightarrow (A \rightarrow \forall x B)$$

2. Show that the inference rule gen preserves validity:

$$\text{gen } \frac{A}{\forall x.A}$$

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3. For a set of formulas  $\Gamma$ , if  $\Gamma \not\models \perp$  then  $\Gamma$  is satisfiable. Why?
4. For a set of formulas  $\Gamma$ , show that the following are equivalent:
  - ▷ If  $A \notin \Gamma$ , then  $\Gamma \cup A$  is not consistent.
  - ▷ For all formulas  $A$ , either  $A \in \Gamma$  or  $A \notin \Gamma$
5. Show that the following forms of consistency are equivalent:
  - ▷  $\Gamma \not\models \perp$
  - ▷ There is no formula  $A$  such that  $\Gamma \vdash A$  and  $\Gamma \vdash \neg A$
  - ▷ There is a formula  $A$  such that  $\Gamma \not\models A$
6. In Lindenbaum Lemma, prove that every  $\Gamma_k$  in the construction is consistent.

7. **(Difficult)** A graph is a structure  $\mathfrak{G} = \langle V, E \rangle$  such that  $E \subseteq V \times V$ . For  $k \in \mathbb{N}$  we say that  $\mathfrak{G}$  is  $k$ -colourable if there is a function  $c : V \rightarrow \{1, \dots, k\}$  such that, whenever  $(u, v) \in E$  then  $c(u) \neq c(v)$ .

Using the compactness theorem, show that a (possibly infinite) graph is  $k$ -colourable iff every finite subgraph of it is  $k$ -colourable.