# A Closer Look At Lists

Principles of Functional Programming

### Lists Recap

Lists are the core data structure we will work with over the next weeks.

```
Type: List[Fruit]
Construction:
   val fruits = List("Apple", "Orange", "Banana")
   val nums = 1 :: 2 :: Nil
Decomposition:
   fruits.head // "Apple"
   nums.tail // 2 :: Nil
   nums.isEmpty // false
   nums match
     case x :: y :: _ => x + y // 3
```

# List Methods (1)

#### Sublists and element access:

xs.length	The number of elements of xs.
xs.last	The list's last element, exception if xs is empty.
xs.init	A list consisting of all elements of xs except the
	last one, exception if xs is empty.
xs.take(n)	A list consisting of the first n elements of xs, or xs
	itself if it is shorter than n.
xs.drop(n)	The rest of the collection after taking n elements.
xs(n)	(or, written out, xs.apply(n)). The element of xs
	at index n.

## List Methods (2)

#### Creating new lists:

xs ++ ys The list consisting of all elements of xs followed

by all elements of ys.

xs.reverse The list containing the elements of xs in reversed

order.

xs.updated(n, x) The list containing the same elements as xs, except

at index n where it contains x.

#### Finding elements:

xs.indexOf(x) The index of the first element in xs equal to x, or

-1 if x does not appear in xs.

xs.contains(x) same as xs.indexOf(x) >= 0

The complexity of head is (small) constant time.

What is the complexity of last?

To find out, let's write a possible implementation of last as a stand-alone function.

```
def last[T](xs: List[T]): T = xs match
  case List() => throw Error("last of empty list")
  case List(x) =>
  case y :: ys =>
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So, last takes steps proportional to the length of the list xs.

```
def init[T](xs: List[T]): List[T] = xs match
  case List() => throw Error("init of empty list")
  case List(x) => ???
  case y :: ys => ???
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Answer: O(xs.length)
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  extension [T](xs: List[T])
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What is the complexity of reverse?
Answer: O(xs.length * xs.length)
Can we do better? (to be solved later).
```

Remove the n'th element of a list xs. If n is out of bounds, return xs itself.

```
def removeAt[T](n: Int, xs: List[T]) = ???
Usage example:
  removeAt(1, List('a', 'b', 'c', 'd')) > List(a, c, d)
```

## Exercise (Harder, Optional)

Flatten a list structure: