

# Subtyping and Generics

## Polymorphism

- Two principal forms of polymorphism are subtyping and generics
- Their interactions consist in two main areas: bounds and variance

## Type Bounds

- Consider the method ***assertAllPos*** which:
  - \* Takes an ***IntSet***
  - \* Returns the ***IntSet*** itself if all its elements are positive
  - \* Throws an exception otherwise
- What is the best type you can give to ***assertAllPos***?
- One might want to express that ***assertAllPos*** takes Empty sets to Empty sets and NonEmpty sets to NonEmpty sets
- A way to express this is: ***def assertAllPos[S <: IntSet](r: S): S = ...***
- Here, “***<: IntSet***” is an upper bound of the type parameter S (this means that S can be instantiated only to types that conform to ***IntSet***)
- Generally, the notation:
  - \* ***S <: T*** means: S is a subtype of T
  - \* ***S >: T*** means: S is a supertype of T
- You can also use a lower bound for a type variable
- Example: ***[S >: NonEmpty]*** introduces a type parameter s that can range only over supertypes of ***NonEmpty***. So S could be one of ***NonEmpty***, ***IntSet***, ***AnyRef*** or ***Any***
- It is also possible to mix a lower bound with an upper bound, but the lower bound comes first
- For instance: ***[S >: NonEmpty <: IntSet]*** would restrict S any type on the interval between ***NonEmpty*** and ***IntSet***

## Covariance

- There’s another interaction between subtyping and type parameters we need to consider: Given ***NonEmpty <: IntSet*** is ***List[NonEmpty] <: List[IntSet]*** ?
- Intuitively, this makes sense: a list of non-empty sets is a special case of a list of arbitrary sets
- We call types for which this relationship holds covariant because their subtyping relationship varies with the type parameter

## Array Typing Problem

- Arrays in Java are covariant, but covariant typing causes problems
- To see why, consider the Java code below:

```
NonEmpty[] a = new NonEmpty[]{  
    new NonEmpty(1, new Empty(), new Empty())  
};  
IntSet[] b = a;  
b[0] = new Empty();  
NonEmpty s = a[0];
```

- By the third line, we don't have a NonEmpty element in the array anymore; the element here is now empty
- That means that in the fourth line we assign an Empty element to a variable of type NonEmpty, which is a violation of type soundness
- So in the third line, you would get a runtime error similar to a class cast exception

## The Liskov Substitution Principle

- The following principle, stated by Barbara Liskov, tells us when a type can be subtype of another:

***If  $A \leq B$ , then everything one can do with a value of type B, one should also be able to do with a value of type A.***

**Exercise:** *The problematic array example would be written as follows in Scala:*

***val a: Array[NonEmpty] = Array(NonEmpty(1, Empty(), Empty()))***

***val b: Array[IntSet] = a***

***b(0) = Empty()***

***val s: NonEmpty = a(0)***

*When you try out this example, what do you observe? – A type error in line two, because Arrays in Scala are not covariant (otherwise, the Liskov principle would be violated)*