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# A LOGICAL FRAMEWORK FOR INTEGRATING SOFTWARE MODELS VIA REFINEMENT

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# BACKGROUND

- ▶ Formal software engineering is a set of mathematically grounded techniques for the specification, development and verification of software and hardware systems.
- ▶ A formal specification is the exact definition in mathematical notation of what the system is required to do (and not do).



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# EVENT B

- ▶ The Event B formal specification language is used in the verification of safety critical systems



- ▶ Event B models are an instance of the specification



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# PROBLEM

- ▶ Different formalisms do not integrate well e.g. Event B models the specification it does nothing for the implementation and its proofs are not easily transferable to other formalisms



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# SOLUTION

- ▶ Establish a theoretical framework within which refinement steps, and their associated proof obligations, can be shared between different formalisms
- ▶ Hypothesis: the theory of institutions can provide this framework and, we will construct an institution based specification of the Event B formalism



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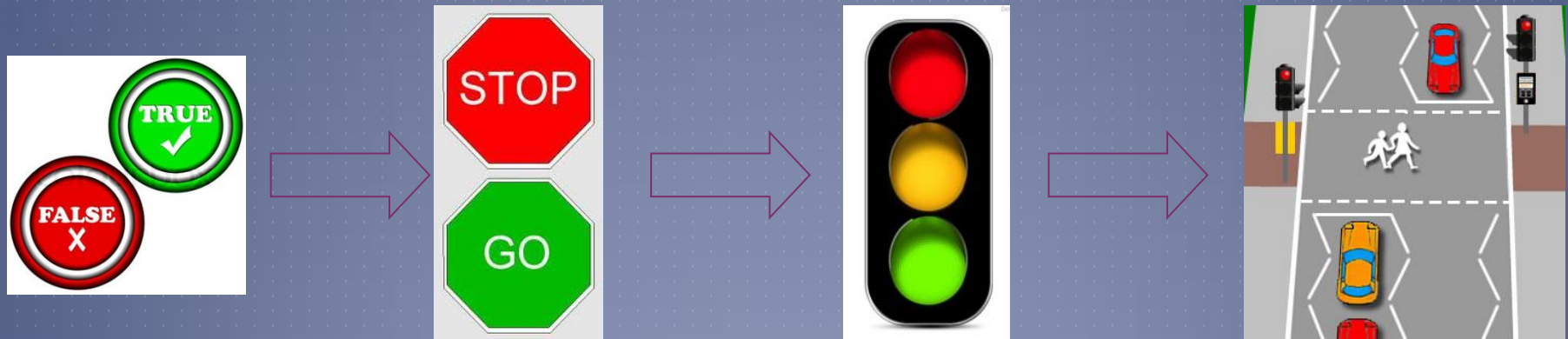
# RESEARCH QUESTIONS



1. Can the theory of institutions ensure the accuracy of the translation between Event-B and other specification formalisms?
2. Can this theory allow us to investigate proof obligations generated by Event-B in different formalisms?

# REFINEMENT

- ▶ Refinement provides a way for us to model software at different levels of abstraction



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# SOCIAL NETWORK

```

MACHINE
  mac1
SEES
  ctx1
VARIABLES
  person
  rawcontent
  content
  owner
INVARIANTS
  inv1 : person  $\subseteq$  PERSON
  inv2 : rawcontent  $\subseteq$  RAWCONTENT
  inv3 : content  $\in$  rawcontent  $\leftrightarrow$  person
  inv4 : owner  $\in$  rawcontent  $\rightarrow$  person
EVENTS
  INITIALISATION  $\triangleq$ 
    STATUS
  ordinary
BEGIN
  act1 : person :=  $\emptyset$ 
  act2 : rawcontent :=  $\emptyset$ 
  act3 : content :=  $\emptyset$ 
  act4 : owner :=  $\emptyset$ 
END

  transmit  $\triangleq$ 
    STATUS
  ordinary
ANY
  rc
  pe
WHERE
  grd1 : rc  $\in$  rawcontent
  grd2 : pe  $\in$  person
  grd3 : rc  $\mapsto$  pe  $\notin$  content
THEN
  act1 : content := content  $\cup$  {rc  $\mapsto$  pe}
END
END
  
```



```

MACHINE
  mac2
REFINES
  mac1
SEES
  ctx1
  ctx2
VARIABLES
  person
  rawcontent
  content
  owner
  visible
  viewpermission
INVARIANTS
  inv1 : visible  $\in$  rawcontent  $\leftrightarrow$  person
  inv2 : viewpermission  $\in$  person  $\leftrightarrow$  person
EVENTS
  INITIALISATION  $\triangleq$ 
    extended
    STATUS
  ordinary
BEGIN
  act1 : person :=  $\emptyset$ 
  act2 : rawcontent :=  $\emptyset$ 
  act3 : content :=  $\emptyset$ 
  act4 : owner :=  $\emptyset$ 
  act5 : visible :=  $\emptyset$ 
  act6 : viewpermission :=  $\emptyset$ 
END

  transmit  $\triangleq$ 
    STATUS
  ordinary
REFINES
  transmit
ANY
  rc
  pe
WHERE
  grd1 : rc  $\in$  rawcontent
  grd2 : pe  $\in$  person
  grd3 : rc  $\mapsto$  pe  $\notin$  content
THEN
  act1 : visible := visible  $\cup$  {rc  $\mapsto$  pe}
  act2 : viewpermission := viewpermission  $\cup$  {owner(rc)  $\mapsto$  pe}
END
  
```



# REFINEMENT CALCULUS

- ▶ Refinement calculus is a notation and a set of rules for deriving programs from their specifications
- ▶ Refinement calculi are an extension of Dijkstra's language of guarded commands and both specification and implementation occur within the same formalism
- ▶ There are three main theories of refinement:
  1. Carroll Morgan
  2. Ralph-Johan Back
  3. Joseph Morris

# MORGAN VS BACK VS MORRIS

- ▶ The definition of what constitutes refinement appears to be the same in all calculi
- ▶ The rules, however, are slightly different: Morgan is the only one to use miracles
- ▶ Back's refinement calculus is much more theoretical than that of Morgan using lattice and category theory as its underlying mathematical basis
- ▶ Morris extended Back's refinement calculus to include the notion of prescription
- ▶ Since the meaning of what is a valid refinement stays the same then regardless of how it is carried out we should always be able to refine a given specification to an implementation that is semantically consistent across all calculi.

# IS THIS REFINEMENT?

- ▶ Regular expression  $\rightarrow$  NFA  $\rightarrow$  DFA  $\rightarrow$  min state DFA
- ▶ Context free grammar  $\rightarrow$  LR Parser
- ▶ Parsing in general
- ▶  $\alpha$ - conversion
- ▶  $\beta$ - reduction
- ▶ A class extending another class
- ▶ A class and an interface it implements
- ▶ An interface and another interface it extends
- ▶ A generic class/ interface and one of its instantiations
- ▶ A class and an instance of the class
- ▶ Liskov Substitution
- ▶ Refactoring
- ▶ UML with OCL  $\rightarrow$  C#/with contracts
- ▶ Is refinement a consequence relation á la Tarski?

# GENERAL THEORY OF REFINEMENT

## - REEVES AND STREADER 2008

- ▶ The general model takes as primitive:
  1. A set of entities: the specifications and implementations we wish to develop by refinement
  2. A set of contexts: the environment with which the entities interact
  3. A user formalised by defining the set of observations that can be made when an entity is executed in a given context
- ▶ The general definition of refinement is parameterised by a set  $\Xi$  of possible contexts and a function  $O$  which determines what can be observed
- ▶ The concrete entity **C** is a refinement of an abstract entity **A** when no user of **A** could observe if they were given **C** in place of **A**.

# DEFINITION

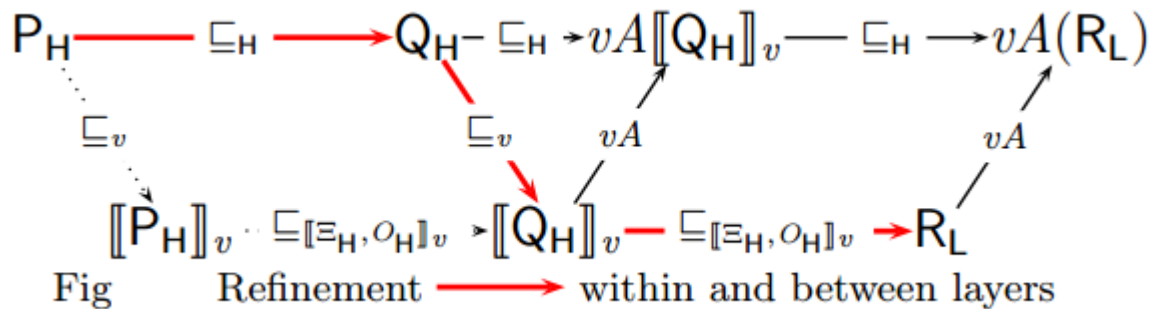
- ▶ Let  $\Xi$  be a set of contexts each of which entities **C** and **A** can communicate privately with, and  $O$  be a function which returns a set of traces, each trace being what a user observes of an execution then:

$$A \sqsubseteq_{\Xi, O} C \triangleq \forall x \in \Xi. O([C]_x) \subseteq O([A]_x)$$

- ▶ Since general refinement has contexts  $\Xi$  as a parameter, by changing  $\Xi$  we are able to model different types of interaction
- ▶ This definition of refinement can be further specialised for refinement of specific cases

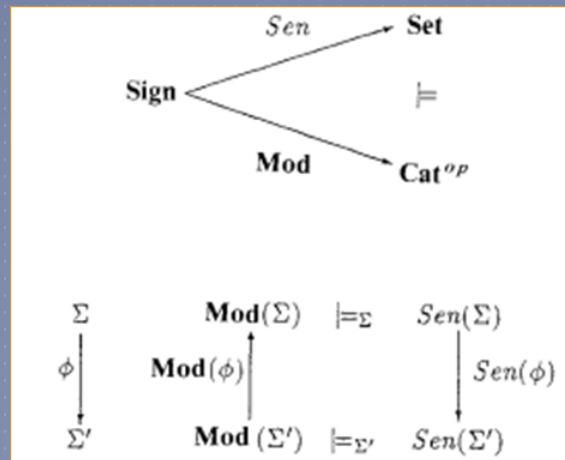
# VERTICAL REFINEMENT

- ▶ We can view each special model of refinement as a layer in the grand scheme of things each encompassing a set of entities and a refinement relation
- ▶ Mathematically our vertical refinement is a Galois connection between the layers.
- ▶ This allows us to interpret high level entities as low level entities using a semantic mapping, however, these low level entities cannot interact with the high level ones so the contexts must also be refined



# CATEGORY THEORY / INSTITUTIONS

- ▶ Category Theory is a special branch of Mathematics that allows us not only to describe objects but also to investigate the relationships between them
- ▶ Institutions are an application of category theory that allow us to relate the syntactic and semantic structures of different formal languages



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# $\Pi$ - INSTITUTIONS

- ▶ Alternative to institution – replacing the notions of model and satisfaction by Tarski's consequence operator
- ▶ Definition:
  - ▶ A  $\pi$ -institution is a triple  $(\text{Sign}, \varphi, \{Cn_{\Sigma}\}_{\Sigma:\text{Sign}})$  consisting of
    1. A category  $\text{Sign}$  (of signatures)
    2. A functor  $\varphi:\text{Sign} \rightarrow \text{Set}$  (set of formulae over each signature)
    3. For each object  $\Sigma$  of  $\text{Sign}$ , a consequence operator  $Cn_{\Sigma}$  defined in the power set of  $\varphi(\Sigma)$  satisfying for each  $A, B \subseteq \varphi(\Sigma)$  and  $\mu: \Sigma \rightarrow \Sigma'$ 

$(\text{RQ1}) A \subseteq Cn_{\Sigma}(A)$	(Extensiveness)
$(\text{RQ2}) Cn_{\Sigma}(Cn_{\Sigma}(A)) = Cn_{\Sigma}(A)$	(Idempotence)
$(\text{RQ3}) Cn_{\Sigma}(A) = \bigcup_{B \subseteq A, B \text{ finite}} Cn_{\Sigma}(B)$	(Compactness)
$(\text{RQ4}) \varphi(\mu)(Cn_{\Sigma}(A)) \subseteq Cn_{\Sigma'}(\varphi(\mu)(A))$	(Structurality)



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# CONCLUSION

- ▶ Work to date:
  - ▶ Denotational Semantics
  - ▶ Communicating Sequential Processes (Hoare)
  - ▶ Category Theory/Institutions/ $\pi$ -institutions
  - ▶ Refinement: Morgan, Back, Morris, general refinement
  - ▶ Questions: Is this refinement?
  - ▶ Tarski Consequence and refinement
- ▶ Work for next semester:
  - ▶ Reading course and project in category theory
  - ▶ Developing an institution for Event B
  - ▶ Developing Event B case studies



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ANY  
QUESTIONS  
?



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