

Name: Marinna Ricketts-Uy

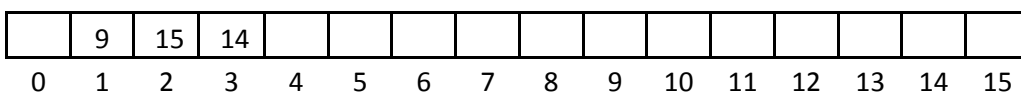
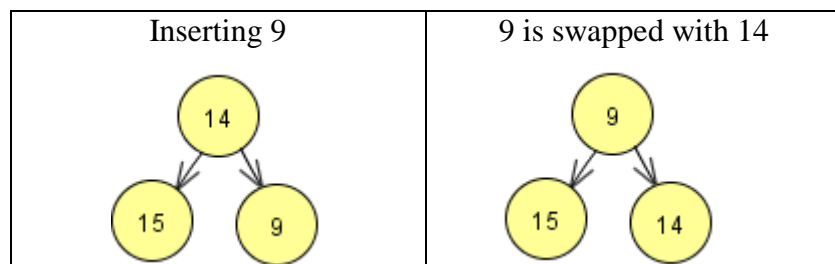
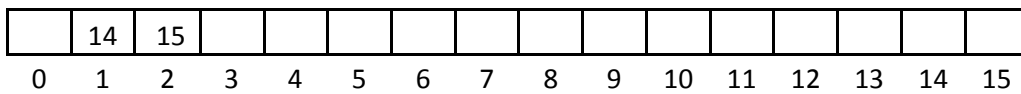
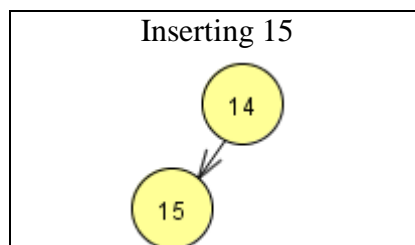
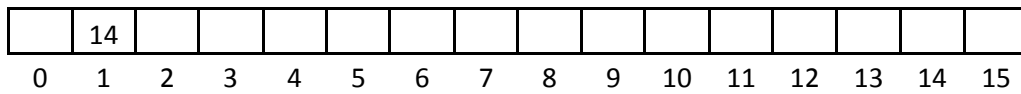
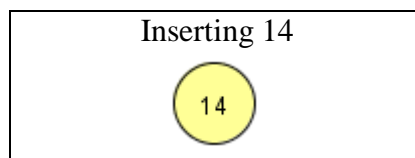
Section: 02 – Budhraj

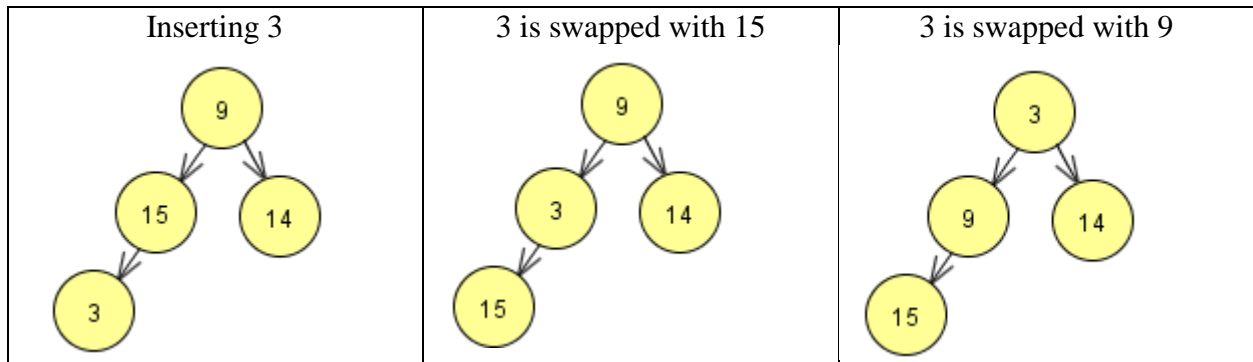
HW #: 4

Version: B

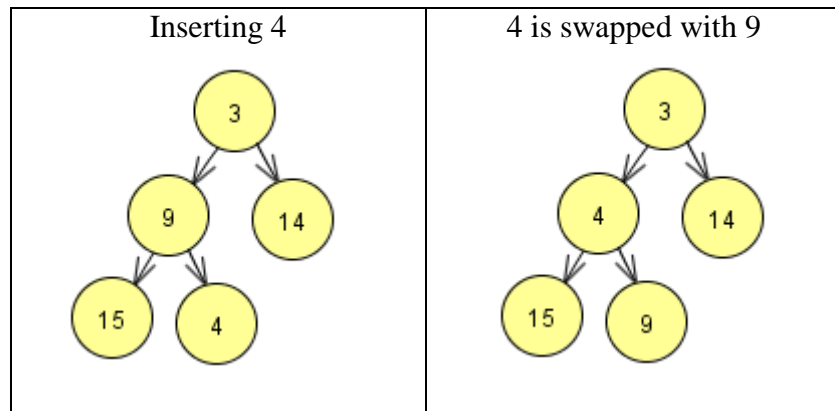
Username: pd12778

1. Inserting the following into a min heap: 14,15,9,3,4,6,8,12,10,13,7,11,1,2,5

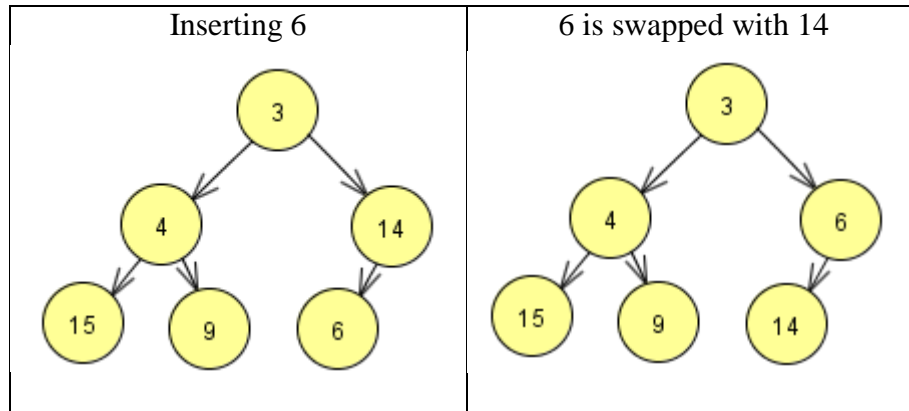




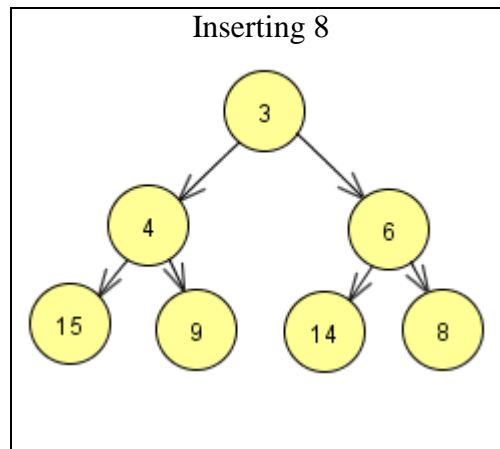
	3	9	14	15											
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



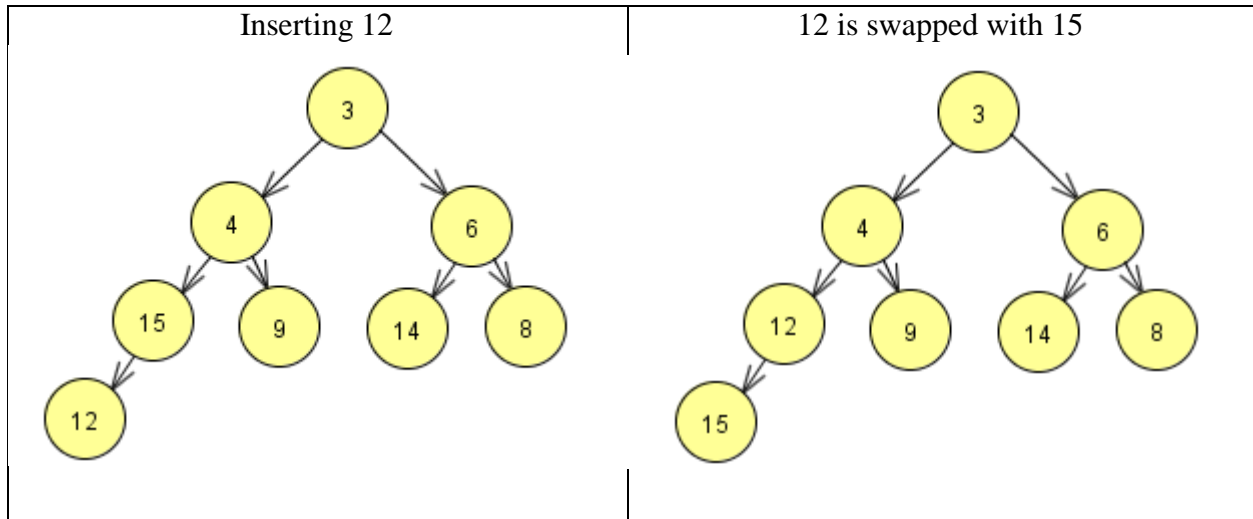
	3	4	14	15	9										
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



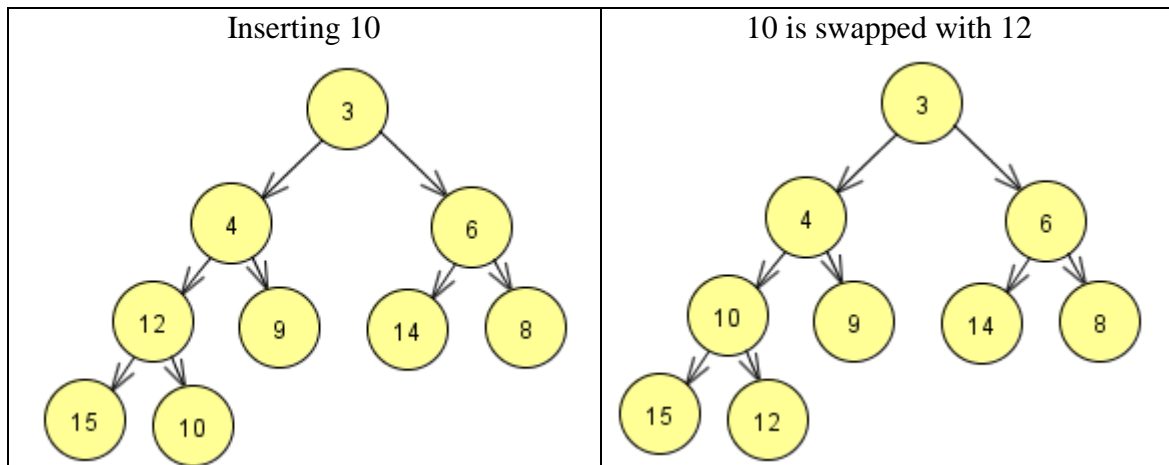
	3	4	6	15	9	14									
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



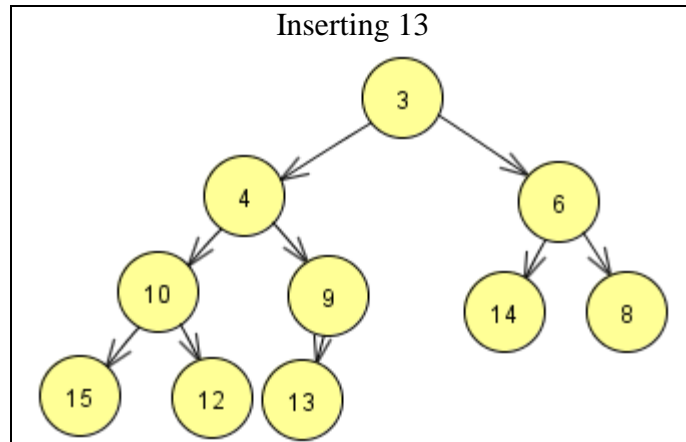
	3	4	6	15	9	14	8								
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



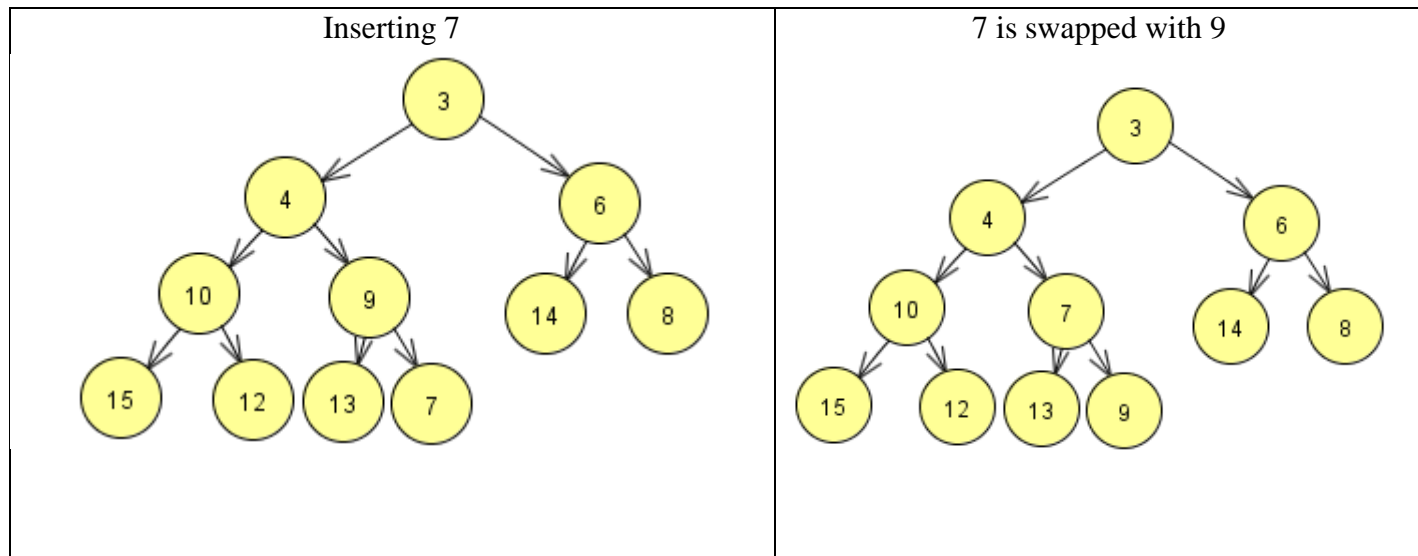
	3	4	6	12	9	14	8	15							
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



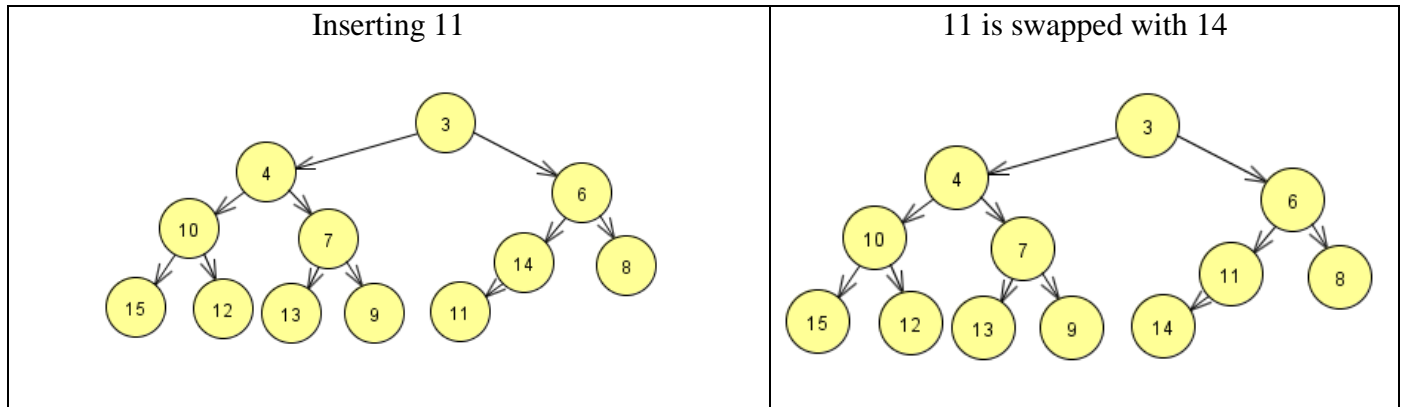
	3	4	6	10	9	14	8	15	12						
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



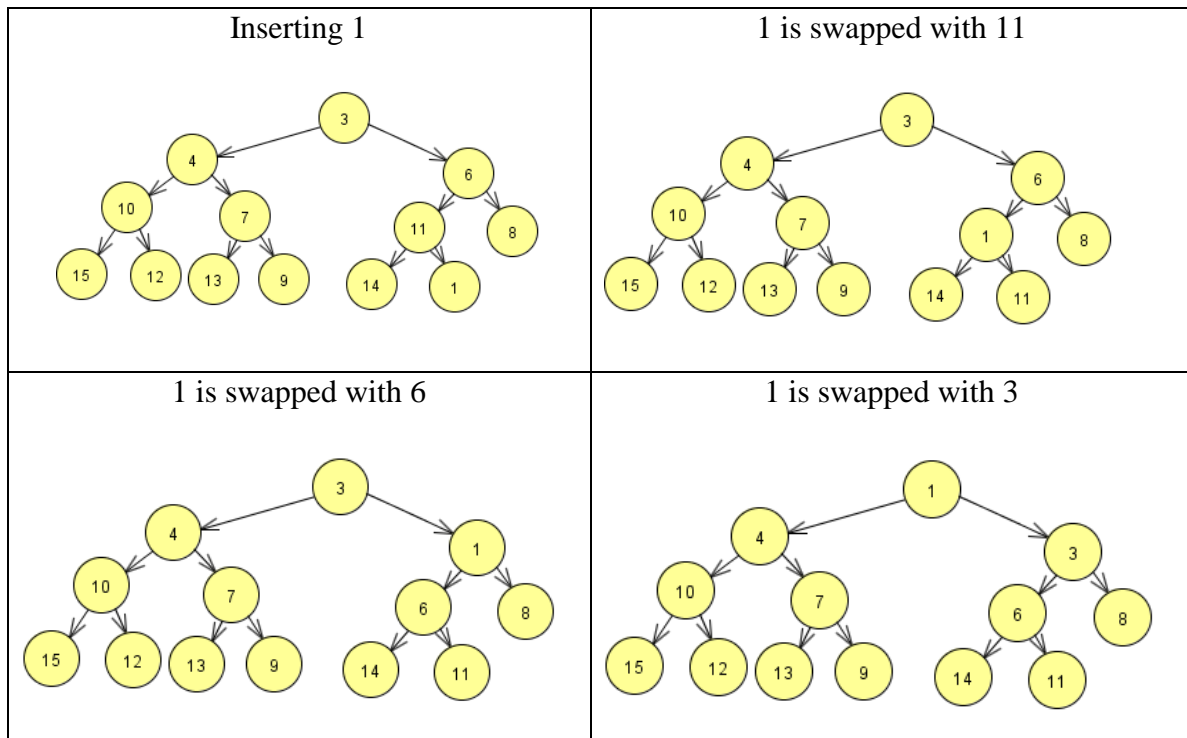
	3	4	6	10	9	14	8	15	12	13					
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



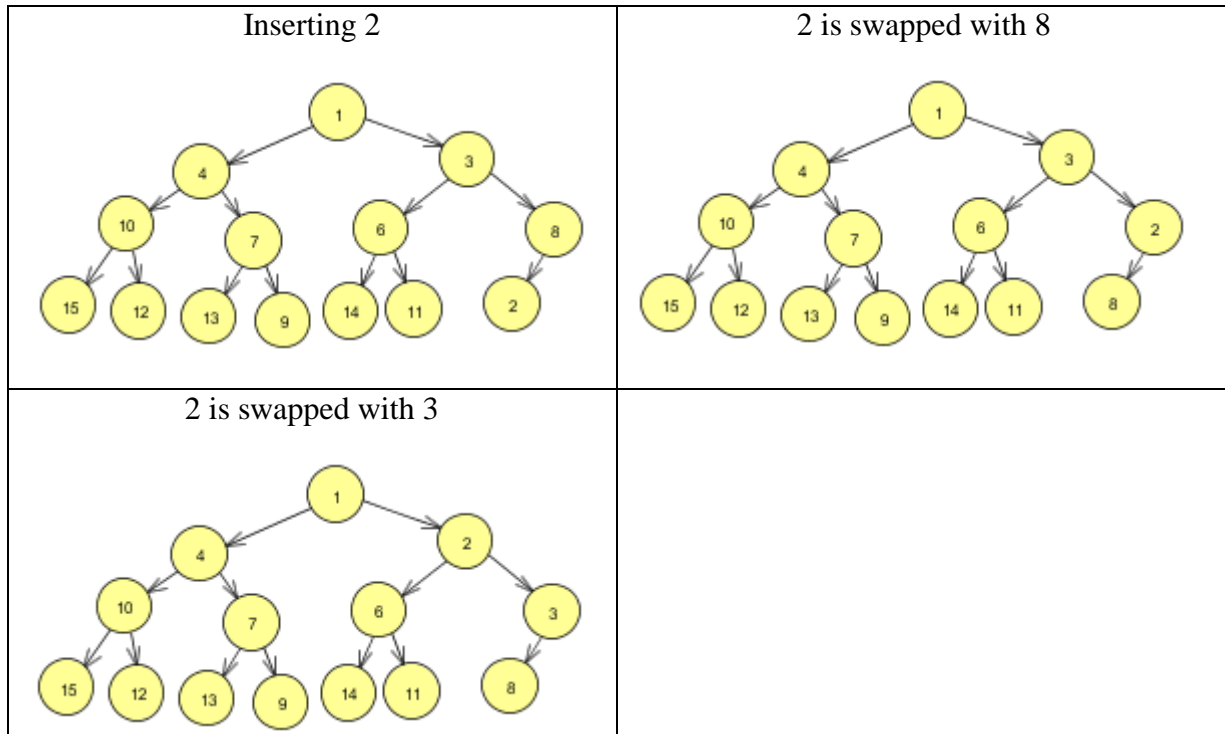
	3	4	6	10	7	14	8	15	12	13	9				
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



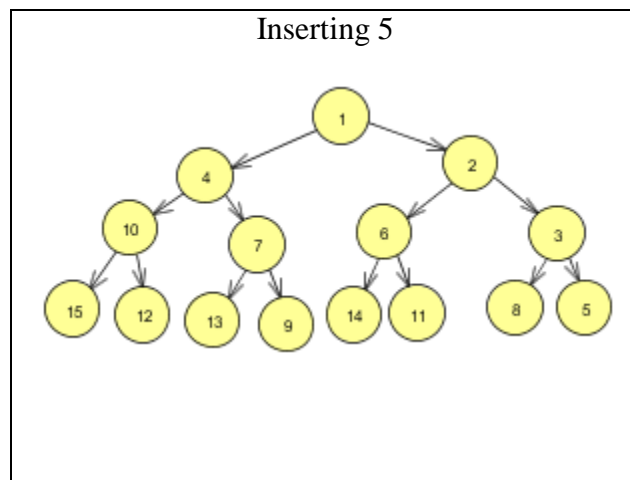
	3	4	6	10	7	11	8	15	12	13	9	14			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



	1	4	3	10	7	6	8	15	12	13	9	14	11		
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



	1	4	2	10	7	6	3	15	12	13	9	14	11	8	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15



	1	4	2	10	7	6	3	15	12	13	9	14	11	8	5
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

2.

a. Prove it must be at one of the leaves:

In order to prove that the maximum item in a min-heap must be one of the leaves, we can assume the opposite.

Proof by Contradiction:

We can assume that it is not a leaf. If this is the case, then the maximum item must have at least one child. If it has a child, and is the maximum item, then it must be greater than its child. However, then based on the property of a min-heap, this is not valid. The parent must be less than its children.

So, the maximum item must be one of the leaves in a min-heap.

b. Prove there are exactly $\text{ceiling}(N/2)$ leaves:

The last leaf in a min-heap is at the N^{th} index. The parent is at index $\text{floor}(N/2)$ and the leaves are indexed from $\text{floor}(N/2) + 1$ to N . So in order to calculate the number of leaves:

$$N - \text{floor}(N/2) = \text{ceiling}(N/2)$$

c. Prove every leaf must be examined to find it:

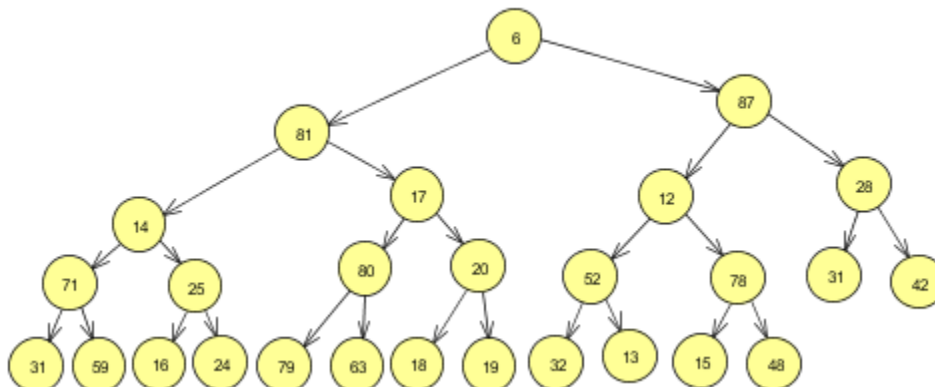
So let's assume the min-heap is a perfect tree of N nodes and the maximum item is M . The next step is to add $N + 1$ nodes that are all greater than M . The values added will be the leaves in the order in which each were inserted. Any of the values could be the maximum item in the min-heap.

So, the maximum item could be in any of the leaves, and every leaf must be examined in order to find which is the greatest amongst them all.

3.

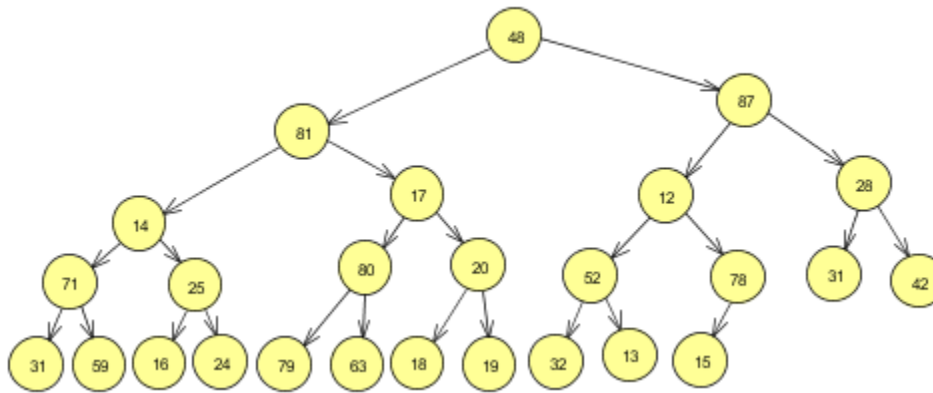
a. Figure 1 after two delete min operations:

Original:

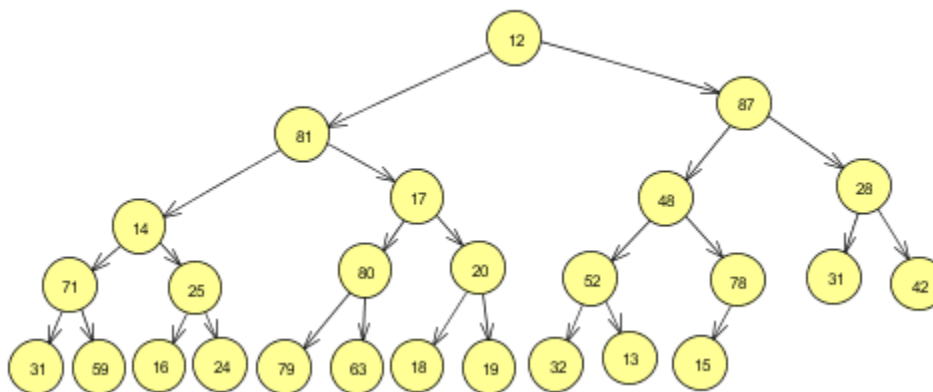


After one delete min operation:

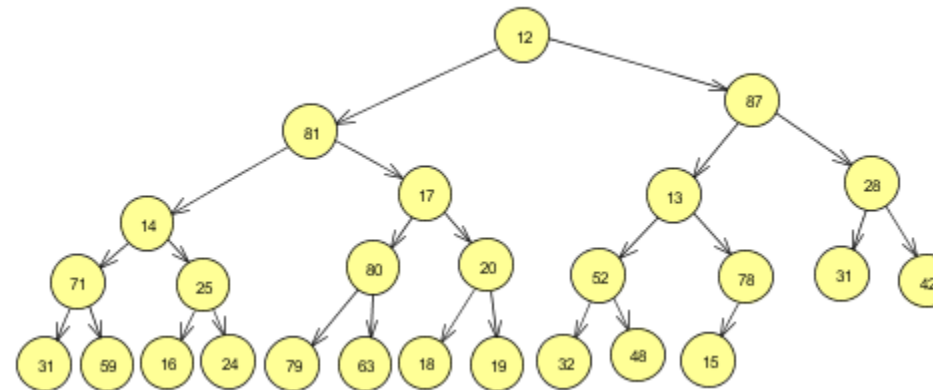
Step 1: Delete 6 and replace with last item in heap (48)



Step 2: Swap 48 with 12

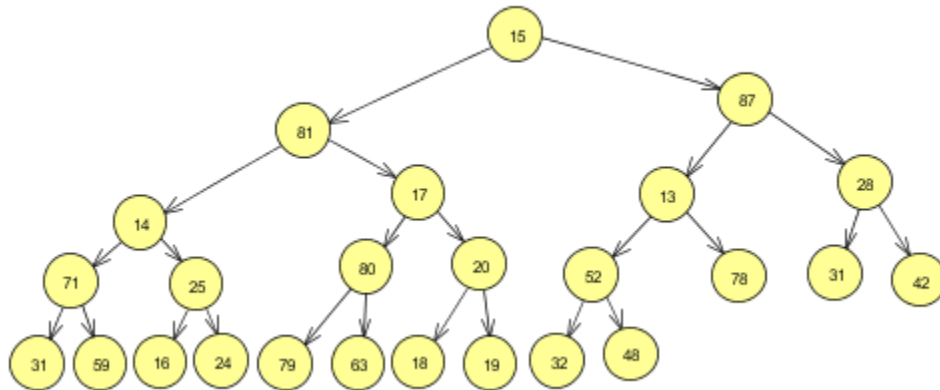


Step 3: Swap 48 with 13

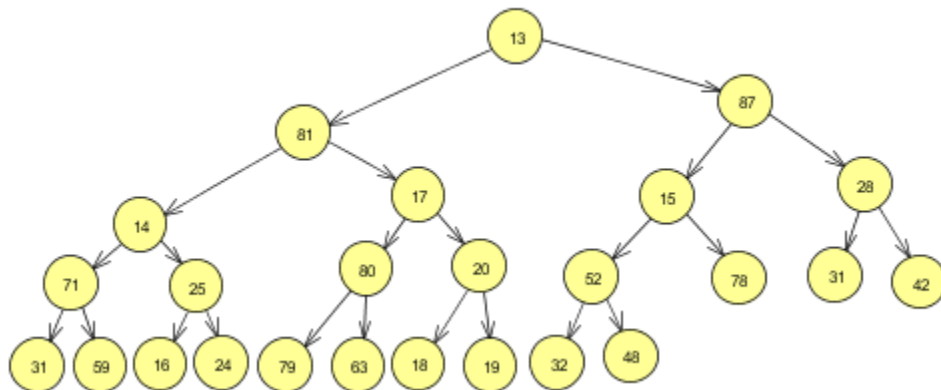


After another delete min operation:

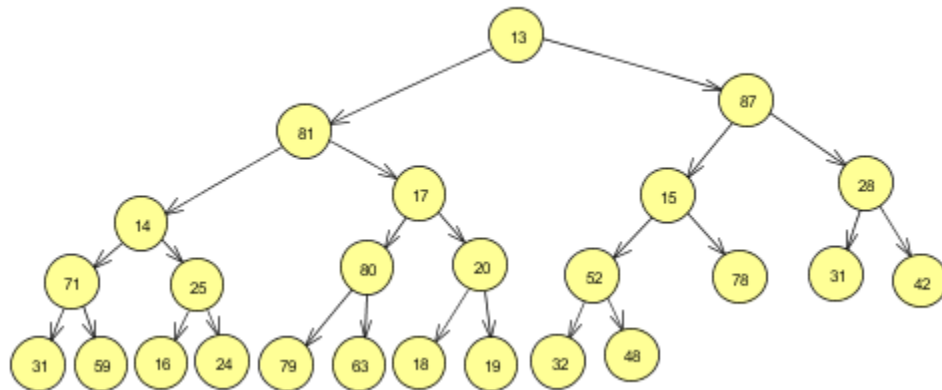
Step 1: Delete 12 and replace with last item in the heap (15)



Step 2: Swap 15 with 13

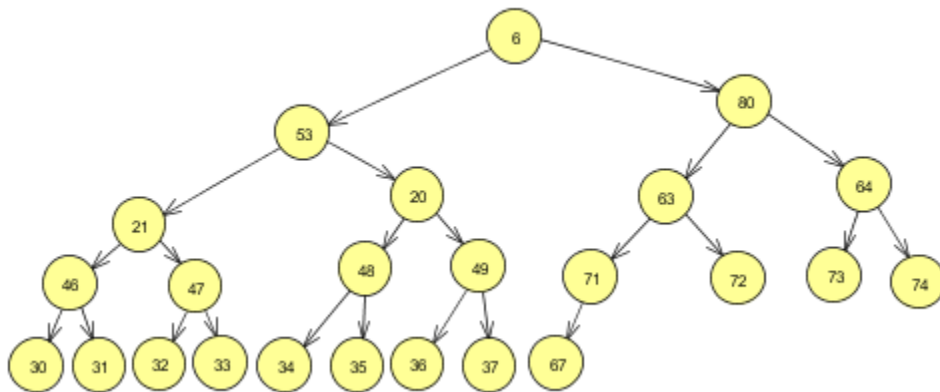


Final Answer (a):

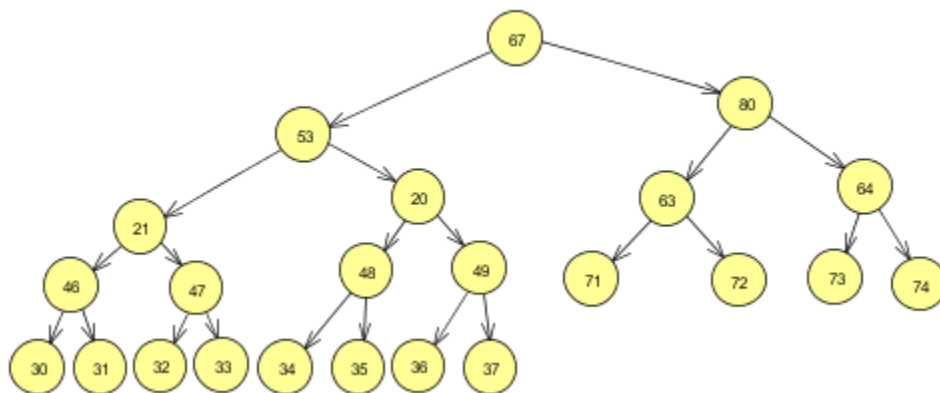


b. Figure 2 after one delete min operation:

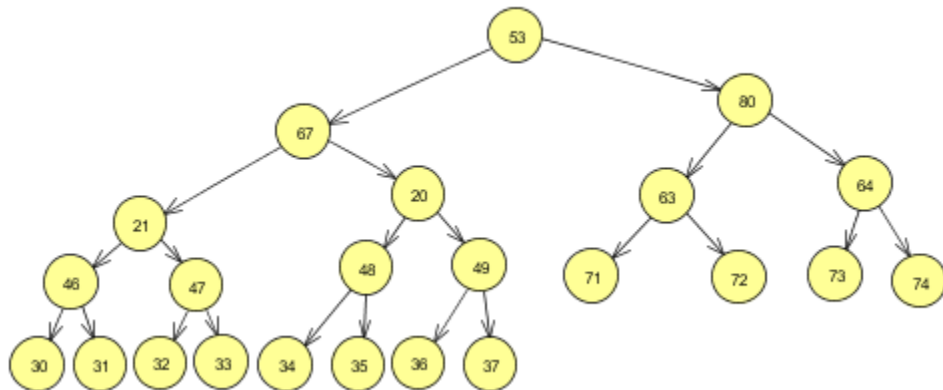
Original:



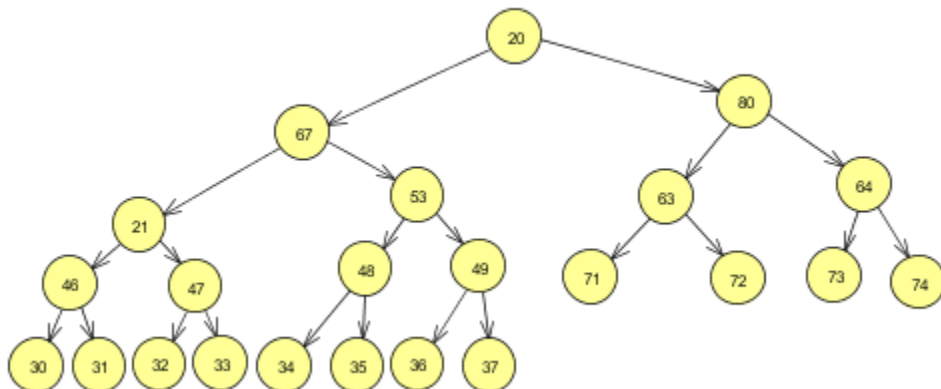
Step 1: Delete 6 and replace with last item in heap (67)



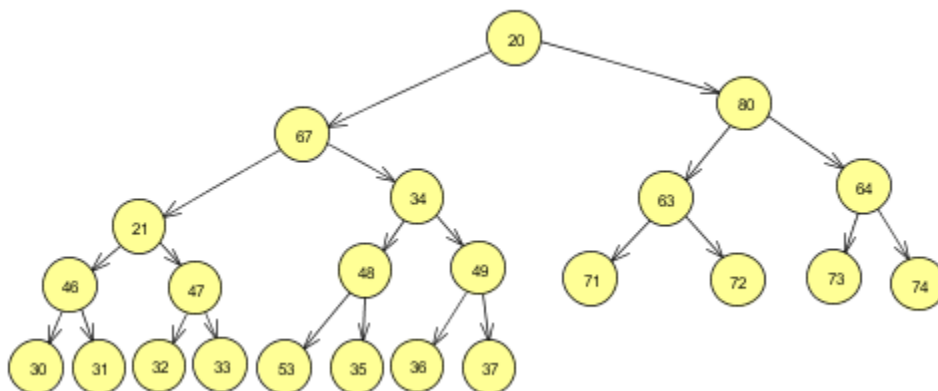
Step 2: Swap 67 with 53



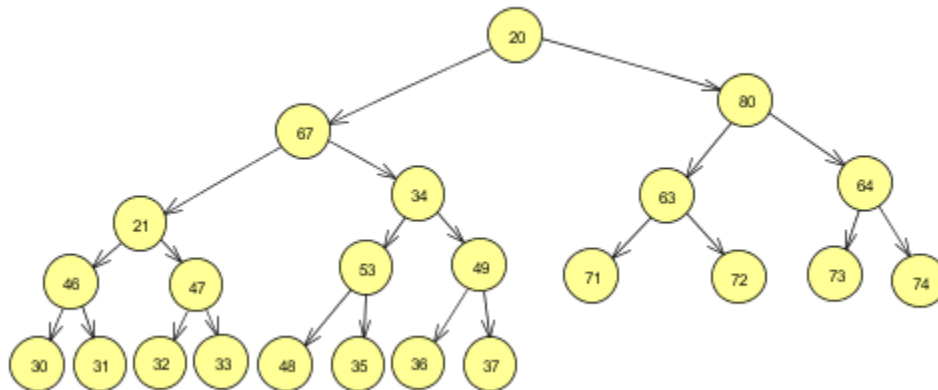
Step 3: Swap 53 with 20



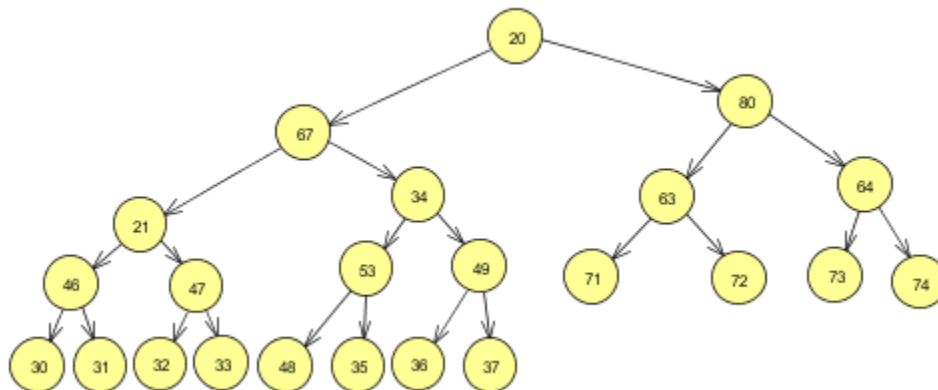
Step 4: Swap 53 with 34



Step 5: Swap 53 with 48

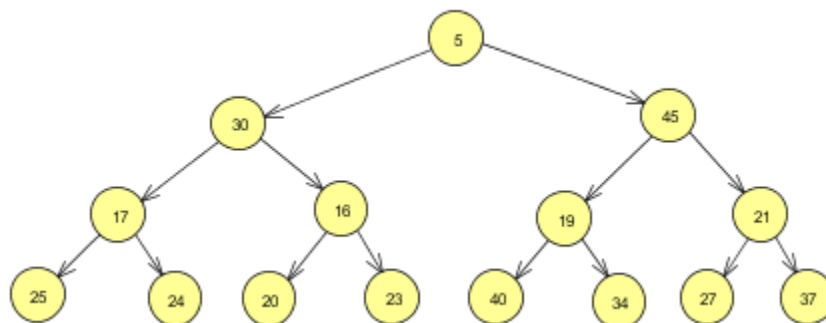


Final Answer (b):

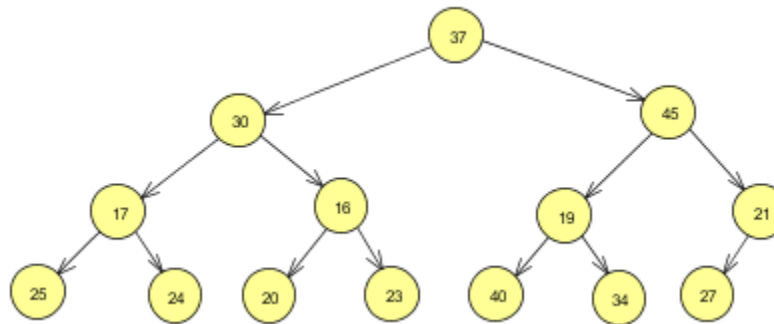


c. Figure 3 after one delete min operation:

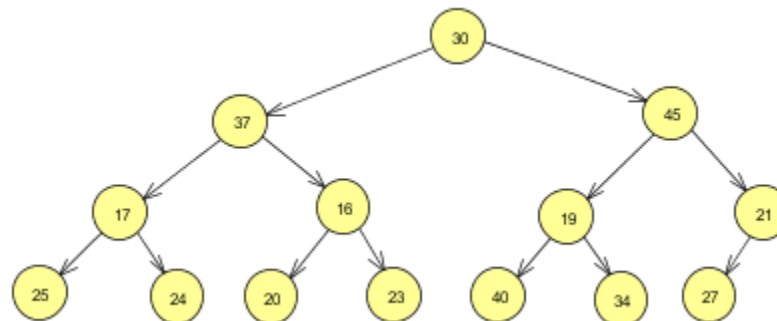
Original:



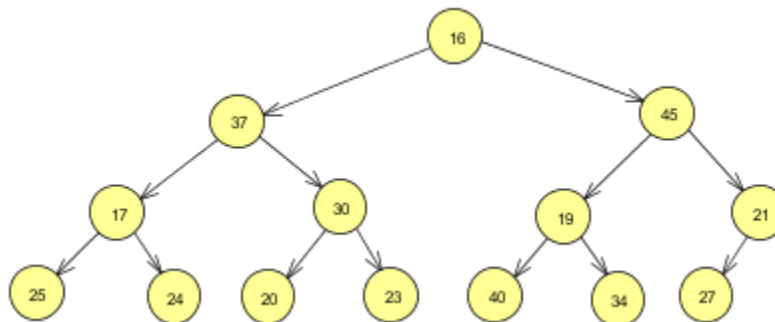
Step 1: Delete 5 and replace with last item in heap (37)



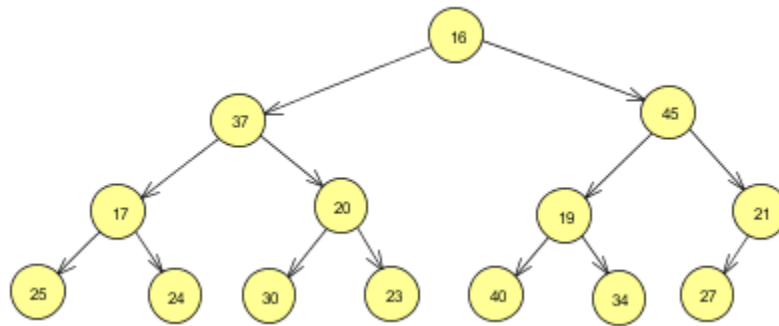
Step 2: Swap 37 with 30



Step 3: Swap 30 with 16



Step 4: Swap 30 with 20



Final Answer (c):

