

USERS' GUIDE FOR ee9: AN ENGLISH ELECTRIC KDF9 EMULATOR

by BILL FINDLAY

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0: INTRODUCTION

This document is a guide for users of **ee9**, a program emulating the EE KDF9 computer. Readers not yet familiar with the KDF9 should consult the companion document, *The English Electric KDF9*; you will be referred to the 'Manual' for further details; this is the *KDF9 Programming Manual*, the nearest thing we have to a definitive reference work. **ee9** should be portable to any system with a basic POSIX API. It is written in Ada 2012 using the GNU Ada compiler, and is available for a range of computers and operating systems. For the characteristics of the latest version, see §6.

1: USING ee9 CONVENIENTLY

ee9 can be invoked directly, but it is *very* much easier to use the command files described in this section. The default setup runs them in the directory **Testing**, so **Testing** should be in your search path to make them accessible to the shell. That is what the immediately following descriptions assume. However, they may be made use of elsewhere so long as you create a compatible environment, for which see §3.1, Appendix 1, and the accompanying README file, *inter alia*.

ee9 requires KDF9 source code and data files to use the Latin-1 text encoding. When you create such a file, ensure that it is saved in the Latin-1 format and not, for example, as UTF-8. The latter would cause characters outside the ASCII subset to be saved as UTF-8 byte pairs, which are meaningless to KDF9 software. The terminal window background is assumed by default to be of a colour other than black or red. White best matches the look of Flexowriter printout. To work with a dark background, or other colours, see Appendix 1.

The directories, e.g. **Binary** and **Assembly**, etc, named in the following, are all subdirectories of **Testing**. They are defaults only; alternative locations for your files may be specified, for which see Appendix 1.

```
nine prog [ data | - ] [ mode | - ] [ miscellany | - ] [ tt [ tt ] ]
```

nine runs the KDF9 machine code program *prog*. The executable is taken from **Binary/prog**; the data for tape reader 1 (TR1), if specified, is taken from **Data/data.txt**, failing which from **Data/data**. A parameter may be given as '-' if it is not needed and a later parameter must be specified. All the optional parameters may be omitted, as they default wisely.

The allowable execution *mode* characters (see §2.1) are:

- **f**: for fast mode, the default (if **f** is explicitly specified, the real CPU time for the run is displayed on termination)
- **p**: for pause mode
- **t**: for trace mode
- **x**: for external trace mode

The allowable characters in the optional *miscellany* parameter are described in §5.2.

The parameters *tt* specify the character codes to be used by either one or two paper tape punches (the first *tt*), or paper tape readers (the second *tt*). *tt* is either one or two characters, each being either k/K or l/L, specifying KDF9 paper tape code or Latin-1 code respectively. If *tt* is one character, it refers to unit 0 of the type (i.e. TP0 or TR0). If *tt* is two characters, the first refers to unit 0 and the second to unit 1 of the type. The default code is Latin-1. See also §3.4 and Appendix 2.

```
crnine prog [ data | - ] [ mode | - ] [ miscellany | - ]
```

crnine is like **nine**, except that data, if specified, is made available via the punched card reader file, CR0.

```
nine_priv prog [ data | - ] [ mode | - ] [ miscellany | - ] [ tt [ tt ] ]
```

nine_priv is like **nine**, but the program is executed in the 'privileged program' mode. See §2.2.

tsdnine *prog* [*data* | -] [*mode* | -] [*miscellany* | -] [*tt* [*tt*]]

tsdnine runs a program with data supplied as TR1, under the control of the Time Sharing Director (TSD). The parameters are as for **nine**, except that a *data* file name without a **.txt** suffix is taken to be in KDF9 code. Before booting, **tsdnine** ensures that the magnetic tape files have valid labels. See Appendix 9 for a primer on operating the TSD.

A program run by **nine**, **crnine** or **nine_priv**, with **ee9** emulating the Director's API, is said to be running 'directly'.

Machine code programs for execution by **ee9** can be created by compiling KDF9 Usercode or Algol source programs. Examples can be found in the **Testing/Assembly**, **Testing/Kidsgrove** and **Testing/Whetstone** directories.

ucc *prog*

ucc compiles the KDF9 Usercode program **Assembly/prog.k3**, placing the object code in **Binary/prog**; a compilation listing is stored in **Assembly/prog-listing.txt**. For help with KDF9 Usercode see Appendix 13.

kalgol *prog*

kalgol compiles the KDF9 Algol program **Kidsgrove/prog.a60**, placing the object program in **Assembly/prog.k3**. The latter is then compiled by **ucc**. The executable left in **Binary** can be run using **nine**. The *mode* and *miscellany* flags for the *compiler* may be given in the environment variables **KIDSMODE** and **KIDSMISC**. By default **KIDSMISC=w** to suppress the compiler's voluminous Flexowriter output. For help with Kidsgrove Algol see Appendix 12.

kids *prog* [*data* | -] [*mode* | -] [*miscellany* | -] [*tt* [*tt*]]

kids compiles and runs **Kidsgrove/prog.a60** (using **kalgol** and then **nine**). See also **wich**, in Appendix 12.

walgol *prog* [*mode* | -] [*miscellany* | -]

walgol compiles and runs the program held in **Whetstone/prog.a60** along with any stream 20 input. The file must be in Flexowriter format. For more-convenient facilities see **whet**, following. For help with Whetstone Algol see Appendix 12.

whet *prog* [*mode* | -] [*miscellany* | -]

whet converts **Kidsgrove/prog.a60** from 'stropped' to Flexowriter format, then uses **walgol** to compile and run it. The result may need some amendment to pass compilation by **walgol**. In particular, unlike **kalgol**, **walgol** does not implement KDF9 'code' procedures. If the converted file already exists you are asked whether you want to overwrite it.

stone

stone runs the Whetstone Benchmark using **walgol** in authentic timing mode, taking about 7 minutes of real time!

pasc *prog* [*data* | -] [*mode* | -] [*miscellany* | -] [*tt* [*tt*]]

pasc compiles and runs **Pascal/prog.p** (using **PASKAL** and then **glance**). See Appendix 14.

2: EMULATION MODES

2.1: DIAGNOSTIC MODES

A KDF9 program is run, at option, in one of four **diagnostic modes**. These are:

- **fast** mode; in which **ee9** runs the program at maximum speed, with no execution tracing or interactive diagnostic facilities available
- **pause** mode, in which **ee9** single-shots the program, pausing to interact with the user after each instruction
- **trace** mode, in which **ee9** runs the program at speed with extensive retrospective tracing enabled
- **external** trace mode, in which **ee9** writes a summary of every traced instruction to an external file

More precisely, things work as follows.

In fast mode **ee9** interacts with the user only by providing informative messages, either because the KDF9 program has terminated, or to log significant events during the run (such as the allocation of an I/O device). All tracing overhead is avoided in fast mode.

In pause mode **ee9** uses console-window text I/O to interact with the user. After each instruction is executed a short summary of the machine state is displayed and a prompt asks the user how to continue. The user replies with an optional single letter (which may be given in upper case or lower case) followed by RETURN, selecting one of the following:

- **f**: execution proceeds in **fast** mode
- **p**: execution proceeds in **pause** mode
- **t**: execution proceeds in **trace** mode
- (nothing): execution proceeds in the current mode.

All retrospective tracing types described in §4 are available in pause mode, trace mode, and external trace mode; but the manner of execution depends on whether the current instruction execution lies within a set range of addresses, and within set instruction-count bounds. If so, instructions are added to their appropriate traces; and breakpoints and watchpoints are monitored. If not, execution proceeds as in fast mode (but at about a third of the speed).

2.2: RUN STATES

The run state specifies how the emulated KDF9 is to run the program:

- In **boot** mode the KDF9 reads a 9-word bootstrap routine from TR0, then jumps to word 0, in **Director** state.
- In **problem program** mode **ee9** reads into core, from TR0, a binary program prepared by a compiler. Its execution starts at word 0, in **program** state. **ee9** itself implements any OUTs (system calls) requested by the program, so that it is not necessary to have a Director running.
- In **privileged program** mode **ee9** reads and runs a binary program, exactly as in problem program mode, except that privileged instructions are allowed to execute successfully instead of failing.

2.3: BREAKPOINTS, FETCHPOINTS, STOREPOINTS, AND WATCHPOINTS

Certain addresses in core can be marked as breakpoints or as watchpoints, to force diagnostic interaction with the user. A **breakpoint** is set on an instruction word, and causes interaction after an instruction beginning in that word has been executed. A **fetchpoint** is set on a data word, and causes interaction after data has been fetched from that word. A **storepoint** is set on a data word, and causes interaction after data has been stored into that word. A **watchpoint** combines a fetchpoint and a storepoint on the same word.

2.4: AUTHENTIC TIMING MODE

At option, **ee9** can be made to insert timed pauses into its execution so that the elapsed time of a program run by **ee9** approximates the elapsed time of a run on the real KDF9 hardware. This may be instructive for younger users, who have never seen characters being output by a computer, one at a time, and with noticeable delays! See §5.2.

3: INPUTS AND OUTPUTS

3.1: EMULATED KDF9 I/O DEVICES

At the start of a run **ee9** looks in the current directory for files to represent the KDF9 peripherals. There is a fixed assignment for the console Flexowriter, which is associated with the user's interactive terminal window. Other devices are associated with files having names derived from the device type. Magnetic tape deck *d*, for example, is always associated with the file named MT*d*. The full list of these associations is as follows:

- a card punch is CP*d*
- a card reader is CR*d*
- a drum store is DR*d*
- a fixed disc store is FD*d*
- a graph plotter is GP*d*
- a line printer is LP*d*
- a KDF9 type 1081 magnetic tape deck is MT*d*
- a Standard Interface device is SI*d*
- an IBM-compatible seven-track tape deck is ST*d*
- a paper tape punch is TP*d*
- a paper tape reader is TR*d*

Every device file actually used in a run must exist, and be appropriately accessible, when **ee9** is invoked. A magnetic tape file, if not completely empty, must have a valid label block. (But see **RLT** in Appendix 11.) The assignment of devices to buffers (KDF9 DMA channels) is set by default, but can be changed by the user: see the K option in §5.1.

3.2: THE FLEXOWRITER CONSOLE TYPEWRITER

The terminal window is the means by which users, in their rôle as KDF9 operators, can mimic Flexowriter I/O. The Flexowriter is used to type-in responses to prompts output by problem programs or by Director. Repeatedly typing these responses quickly becomes tedious. If a file named FW0 exists, it is used as a source of “canned” responses. They are defined, with their identifying prompts, in FW0; and are picked up automatically by **ee9**. If a prompt spreads over more than one line, a KDF9 Line Shift can be represented in FW0 by a ‘®’, and a KDF9 Page Change by a ‘©’.

When a prompt is issued, **ee9** scans FW0, down from the last match found. If it finds a new match, it injects the given response into the Flexowriter input stream; but if it reaches the end of the file without finding a match, it returns control of the Flexowriter to the user's terminal window, so that a manual response can be given. If a prompt matches a line in FW0 that specifies a null response string (c.f. the second ‘OUT;’ in the following example) then **ee9** terminates the run.

For example, the Whetstone Algol compiler prompts ‘[q] OUT;’ to which a typical reply is ‘N. |’. If the Algol program compiles, it runs and prompts ‘[q] STREAM;’ to which a typical reply is ‘30. |’; but if the compilation fails the compiler loops back to its ‘[q] OUT;’ prompt, where the user will normally want to terminate the run so that the Algol source code can be amended. The following data in FW0 will achieve this without user intervention:

```
[q] OUT;N. |
[q] STREAM;30. |
[q] OUT;
```

For a second example, as the Time Sharing Director bootstraps it issues a series of requests for basic configuration parameters. The following data in FW0 supplies suitable responses without user intervention:

```
CORE MODULES;8. |
OUT 8 REEL NO;9. |
LEVELS;N. |
DATE D/M/Y;4/5/67. |
TIME ON 24-HOUR CLOCK®HOURS/MINS;1/23. |
```

This facility had a real equivalent: the Flexowriter incorporated an ‘edge-punched card’ reader. It read data (in paper tape code) from the edge of a non-standard punched card. Cards prepared with replies to prompts could be inserted into the reader and read at the maximum rate, thus speeding input and avoiding any delay due to typing errors by the operator.

Note that ee9 requires every Flexowriter input string to be terminated by a RETURN, even when a read-to-End Message instruction is being obeyed. In reality, KDF9 would end the transfer immediately at the End Message (EM) character, or when the required number of characters had been read; but data is not transferred to **ee9**’s input buffer until a RETURN is typed. A merely terminating RETURN is discarded from the input by **ee9**, and is not passed to the KDF9 program.

In response to CTRL-C, **ee9** outputs a prompt that lets the diagnostic mode be changed. Replying with RETURN causes a FLEX interrupt when running Director, evoking a ‘**TINT**’; (CTRL-C typed when **ee9** has failed to respond quits the run.)

Output to the Flexowriter was typed in red; input from the computer operators in black. This is simulated in **ee9** by using ANSI-standard SGR escape sequences to vary the displayed font style. Some Windows terminal utilities do not implement SGR escape sequences, so Flexowriter output styling can be disabled with a miscellany parameter or a settings file option.

3.3: READING MORE THAN ONE ROLL OF PAPER TAPE OR DECK OF CARDS

A means is provided to simulate the way in which KDF9’s computer operators could satisfy a program’s demand for data with several physically-separate rolls of paper tape or decks of cards, loaded into a reader in succession. If a program attempts to read from a tape reader, and the end of the associated file has been reached, **ee9** allows the user to specify a successor file to which the tape reader is re-attached. A prompt requesting more data is displayed.

- If the user responds ‘=’ then the reader is connected to the user’s terminal so that the data may be typed then and there.
- If the user responds ‘/’ then the pathname of a file supplying the next “roll of tape” may be typed. If the response is ‘@’ then the name of a file within the directory specified by KDF9_DATA may be typed; see Appendix 1. With both ‘/’ and ‘@’ a final ‘.txt’ may be omitted from the name. If the specified file is not found, or not readable, the prompt is repeated.
- The filename may be given on the same line as ‘/’ or ‘@’ if the latter is followed by a separator such as a space.
- If the user responds ‘q’ or ‘Q’ then execution terminates.
- If the user presses RETURN (only), or signals end-of-data (e.g. by typing ^D), the buffer is set to the abnormal state. Attempting another read without clearing the error will cause a Lock-In Violation.
- The foregoing also applies, *mutatis mutandis*, to punched card readers.
- See the **N** option in §5, about disabling this feature in non-interactive mode.

Examples (user’s responses in **green**):

```
1.
ee9: Type @ or / to name a file, = to type the data, ENTER key for EOF, Q or q to quit: =
ee9: Type the data for TR1:
3;|
2.
ee9: Type @ or / to name a file, = to type the data, ENTER key for EOF, Q or q to quit: /
ee9: Give the pathname of the file: next
The KDF9 operator must have made a mistake: next cannot be read.
3.
ee9: Type @ or / to name a file, = to type the data, ENTER key for EOF, Q or q to quit: / Data/next
4.
ee9: Type @ or / to name a file, = to type the data, ENTER key for EOF, Q or q to quit: @
ee9: Give the name of a file in Data/: next
5.
ee9: Type @ or / to name a file, = to type the data, ENTER key for EOF, Q or q to quit: q
...
Final State: Run stopped by the user: quit requested.
```

3.4: REPRESENTING THE KDF9 CHARACTER SETS

The KDF9 had its own, somewhat idiosyncratic, character sets. They are tabulated in Appendix 2a. For the convenience of the user, external data read and written by the paper tape readers and punches may be represented in either KDF9 paper tape code, or in the ISO Latin-1 character set with automatic conversion between Latin-1 and KDF9 codes. Several characters in the KDF9 paper tape set are absent from Latin-1, so a simple transliteration is used to represent them externally.

For the Flexowriter, and for Latin-1 I/O with tape punches and tape readers, Case Normal and Case Shift characters are generated on input, and interpreted on output. This means that when you are typing an input text, it is not necessary to type Case Normal and Case Shift characters, although there usually is no harm in doing so. When such a text is being read as the input stream for a two-shift device, an appropriate case-character is generated automatically by the emulator, if the Latin-1 character being read is not available in the input device’s current shift state. Two-shift output works in the complementary manner. Two-shift devices always start out in the Case Normal condition.

For example, the external Latin-1 string ‘Bill Findlay’ may be read into the KDF9 core store as the characters ‘BBILL ñFBINDLAY’, with ß denoting the Case Shift character and ñ denoting the Case Normal character. A KDF9 program that writes the characters ‘BBILL ñFBINDLAY’ to a two-shift device in Latin-1 mode will generate the Latin-1 string ‘Bill Findlay’ as its external representation.

Text-file input to **ee9** may use any of CR, LF, or a CRLF pair as the line terminator: **ee9** treats all three the same. Text-file output from **ee9** writes the line terminator most appropriate for the host’s operating system.

Non-graphic KDF9 characters also have Latin-1 external representations, to enable 1-to-1 inter-conversion between the internal and external data formats. Apart from the format effectors (Horizontal Tab, New Line, Form Feed), users should never need to type these characters, as they could not be typed on a Flexowriter.

In KDF9 mode, characters are transmitted using the KDF9 paper tape format, in which devices use an 8-bit byte externally, supplementing the 6 internal data bits with two ‘parity’ bits (see also the Manual, §17.4 and Appendix 2b).

KDF9 object programs are never given in Latin-1, but always natively, in the KDF9 paper tape code.

Characters displayed in tracing output are shown using their Case Normal Latin-1 transliterations, except as follows:

- the KDF9 Horizontal Tab character is depicted as `␣`
- the KDF9 Line Shift or LS character (Carriage **R**eturn) is depicted as `␣`
- the KDF9 Page **C**hange or PC character (Form Feed) is depicted as `␣`
- the KDF9 Filler character is depicted as `␣`

These depictions are also used by the card reader and card punch. Card images of less than 80 characters are padded with blanks to fill all 80 columns; any input line longer than a card is truncated to 80 characters. In ‘direct’ mode, card images may have up to 160 characters, notionally two per card column. The card punch writes files that are valid as input to the card reader and can be read to exactly reproduce the original data in core.

The representation of magnetic tapes is described in detail in Appendix 10.

Drums and disc drives are represented as plain POSIX files using the Latin-1 Case Normal transliteration. Sector boundaries are not explicitly represented: they are implicit in the data address within the POSIX file, occurring every 320 bytes for disc drives and every 1024 bytes for drums.

3.5: GRAPH PLOTTING

ee9 includes an emulation of the model 564 Calcomp graph plotter, as described in Appendix 6, §5, p.302 of the Manual (but see what I say about this in §6: Caveats, and in Appendix 3). There was provision on the KDF9 to switch a buffer manually between a tape punch and a graph plotter; in **ee9** this is done with **g** in the *miscellany* parameter, with a settings file option **G**, or with a settings file option **K**. When one or more of these is given, GP0 replaces TP1 on the shared buffer.

The KDF9 graph plotter takes commands that move the plotting position in steps of 0.005 inches; see Appendix 3 for the command codes. Steps are accumulated into vectors by **ee9**, and PostScript vector drawing commands are output to GP0, creating an Encapsulated PostScript (EPS) file. It is possible to fit the plotter with ‘pens’ having a variety of ink colour and ball-point tip size: see under **G** in §5.

4: LOGGING AND TRACING

Messages that record the progress of the emulation, and details of any errors that were detected, are written to the interactive console window, along with interactive diagnostics, and output produced for the KDF9 Flexowriter. A selection of these messages is also written to the file `KDF9_log.txt`. On completion of a run, the final machine state, any requested core store areas, and any retrospective traces may be written to the log file and to the console window.

It is possible to request the output of certain areas of the KDF9’s core store, in a variety of suitable formats. These printouts can be taken either before the start of execution; or on termination; or at both times, to allow comparisons.

In the **interrupt** trace, which is produced only in boot mode, interrupt requests are listed with the privilege state and priority of the interrupting device; the elapsed time of occurrence (in μ s); and the value of ICR, the Instruction Count Register, which is a count of the number of instructions executed so far. The interrupt trace also logs EXITD, return from interrupt, instructions.

The tracing of executed instructions is subject to instruction-count and address-range bounds. Instruction executions within those bounds are traced; those that fall outside the bounds are not.

In the peripheral I/O trace, the events shown are transfer initiations and terminations, busy-buffer and store-access lockouts, and I/O status test operations. Each is listed with the device name, Q-store parameter, privilege state (P-S for problem program slots and D for Director state) and priority of the transfer; the elapsed time of occurrence of the event; and the current instruction count.

The C part of the parameter used in a disc seek operation is logged in the format `DdPppSss`, where *d*, *pp* and *ss* are, respectively, the drive number, platter number and *seek area* number being addressed. In non-seek disc operations the C part is shown as `Sss`, where *ss* is the starting *sector* number for the transfer.

The C part of the parameter used in a drum transfer is logged in the format `DdTtSs` where *d*, *tt* and *s* are the drive number, track number and the starting sector number.

Transfer operations appear twice, once for the initiation (S) and once for the termination (E). Lockouts appear once, when they happen. A test operation gives the result of the test as a Boolean (Y or N) in the T field of a trace.

In the retrospective trace, instructions are listed in order, starting with the most recently executed. The trace includes the instruction itself, and its most relevant operand; ‘ND’ and ‘SD’, the Nest and SJNS Depths; ‘V’ and/or ‘T’ showing whether overflow and/or the test register is set; the CPU time of occurrence of the event; and the value the current instruction count. In the case of a store order, the traced operand is the value written to store. In the case of a fetch order, it is the value fetched. For a Q-store order, it is the content of the relevant Q register. For a conditional jump it is the determining value. For subroutine jump or exit, it is the relevant value in the SJNS. For a 1-syllable or 2-syllable ALU order, it is the value left in the top of the nest. And so on.

External trace mode is like retrospective mode, with additional output to the file `trace.txt`. This output has one line for each traced instruction. It contains: the instruction’s address; the value of ICR; the CPU time; the nest depth; the SJNS depth; ‘V’ and/or ‘T’ if overflow and/or the test register is set; the value in N1, if the nest is non-empty; and the disassembled instruction.

When tracing, and if requested, **ee9** will tally the number of traced executions of each type of KDF9 instruction. On termination of a run a histogram of dynamic instruction-type frequencies may be logged. A similar plot may be produced showing the frequency with which each instruction-containing word is executed.

At option, all tracing modes can compute a digital signature of the execution: a 48-bit cumulative hash, displayed in octal, of the contents of all the relevant KDF9 registers (nest, SJNS and Q stores) at the end of each traced instruction. Known values for this hash can be used to verify the proper operation of an implementation of **ee9**. (When the signature is enabled, the starting time-of-day is forced to midnight, to ensure a repeatable hash value.)

For examples of the output that can be obtained using these facilities, see Appendices 7A-7G.

5: THE MODE SETTINGS FILES AND THE MISCELLANY PARAMETER

The emulator has default settings for all of its options, but they may be over-ridden by settings specified in files that the emulator attempts to read as part of its initialization, and/or by specifying a miscellany parameter on the command line.

5.1: SETTINGS FILES

The file `settings_1.txt` applies to a first or sole program to be run, and `settings_2.txt` applies to a second program overlaid by it (e.g. the Whetstone Controller, overlaid after a successful compilation by the Translator). A setting specified by the command line over-rides a similar option specified in the `settings_1.txt` file.

The settings files contain a line for each option to be set, beginning with a letter that specifies the option concerned. This may be followed by one or two parameters. Numbers may be given in octal—preceded by a hash sign (`#`)—or in decimal; and this convention is also used systematically in output messages from the emulator. The options are presented here in upper case, but lower/mixed case is also accepted. The available options are as follows:

A MODERN_TIMES_MODE | AUTHENTIC_TIME_MODE

The **A** flag sets aspects of the AUTHENTICITY of execution. It takes one symbolic parameter. It is possible to set authentic elapsed timing (see §2.4), with the AUTHENTIC_TIME_MODE parameter; in MODERN_TIMES_MODE emulation proceeds at the full speed of the host computer.

B start [end]

This flag has either one or two parameters, which are instruction-word addresses. It sets a BREAKPOINT on every instruction word in the given range of addresses.

C l h

This flag is used to set two COUNT values, say *l* and *h*, that determine when tracing is done. No breakpoint or watchpoint fires, and no instruction is traced, unless $l \leq i \leq h$ is satisfied; where *i* is the current value of ICR. With suitable *l* and *h* values, tracing can be confined to a set time during execution (for example, the last few instruction executions before a program fails). The values *l* and *h* are given as unsigned decimal integers.

D FAST_MODE | TRACE_MODE | PAUSE_MODE | EXTERNAL_MODE

This flag sets the DIAGNOSTIC mode, specifying the type of tracing and the kind of logging that may be generated.

Fi start end [base] and Ii start end [base]

These flags have two mandatory parameters, which are word addresses. They request that the contents of that range of addresses be output in a specified interpretation, INITIALLY, or FINALLY (i.e. after the end of the run).

For both *Ii* and *Fi*, the interpretation is given by the string of letters *i*, each letter of which must be one of: **A**, for strings in ASCII/LATIN-1 code; **C**, for strings in card code; **L**, for strings in LINEPRINTER code; **N**, for strings in paper tape code, with case NORMAL shown; **S**, for strings in paper tape code, with case SHIFT shown; **T**, for strings in paper TAPE code, with shift characters actioned; **O**, for syllabic OCTAL/ ORDERS; **U**, for orders in an approximation to USERCODE format; and **W**, for data WORDS in octal, syllabic octal, line printer characters, Q store format, and signed decimal.

With **U**, **D** can also be given to get machine code addresses in DECIMAL instead of octal. A third parameters may be given. When printing the code of a Director it specifies an additional point of entry (e.g. for an overlay). For a problem program it specifies the address of V0 of the routine delimited by *start* and *end*, so that disassembly may present data accesses as V store references instead of absolute addresses. For examples of Usercode format, see Appendix 6.

The PIC, PID, POC and POD instructions for cards and paper tape permit the processing of data in arbitrary character codes. The **A** format for core-store printing is provided to facilitate the debugging of modern KDF9 programs that process data in ASCII/Latin-1, the native character set of **ee9**.

G [colour [tip size]]

This flag allows one or two optional symbolic parameters. The first, if given, sets the GRAPH PLOTTING pen colour from the list: BLACK (the default), BLUE, BROWN, CYAN, DARK_BLUE, DARK_CYAN, DARK_GREEN, DARK_GREY, DARK_MAGENTA, DARK_RED, GREEN, GREY, MAGENTA, RED, WHITE, YELLOW. If a colour is given, a second parameter may be given to set the pen tip size from the list: EXTRA_EXTRA_FINE (the default, 1 plotter step wide), EXTRA_FINE (2 steps wide), FINE (4 steps), MEDIUM (6), MEDIUM_BROAD (8), BROAD (10), EXTRA_BROAD (12).

In any case, the shared buffer is switched from TP1 to GP0.

Hh cutoff%

In tracing mode, **ee9** collects statistics on the execution of the KDF9 program. These may be displayed at the end of a run in the form of histograms. Two plots are available. If *h* contains P or p, an execution-frequency profile is output. If *h* contains T or t, then an instruction type frequency histogram is output. Both may be requested. *cutoff* is a number in the range 0.0 to 100.0, taken to be a percentage; the % is required to follow the number immediately. Histogram bins that account for less than that fraction of all executed instructions are suppressed. Nothing is output if the option is suppressed by the miscellany parameter.

K { buffer code }

This flag sets the KDF9 configuration by specifying the device types connected to the buffers. The parameters consist of one or more pairs of items, the first a *buffer* number (in the range 0 to 15) and the second a two-letter *code*, being the first two letters of the device name (e.g. CR or LP). GP may be specified to switch TP1's buffer to GP0. At most 1 DR unit, 1 FD unit, 2 TP, TR, CP, CR, LP, or SI units, at most 14 tape decks of either kind, and only one of either DR or FD can be configured. To leave a buffer with an absent device, specify the code AD; any attempt to use that buffer will fail. Every configuration must have FW on buffer 0 and TR on buffer 1. Any buffer not listed keeps its default device, which is that specified by:

K 0 FW 1 TR 2 TR 3 TP 4 TP 5 LP 6 CR 7 CP 8 MT 9 MT 10 MT 11 MT 12 MT 13 MT 14 FD 15 ST

The **K** flag is actioned only at the start of a run; any **K** options in the `settings_2.txt` file are ignored. N.B. The commands described in §1 may not work properly with a non-default configuration, especially if only one TR and no TP or LP is included.

L [*t*]

This flag gives a value, *t*, that sets an execution time **LIMIT** on how long the KDF9 program is allowed to execute before being terminated. The limit is specified in instruction executions, so the program is terminated if $ICR > t$ at the end of any instruction execution. The value *t* is given as an unsigned decimal integer. If the **L** flag is given without a parameter the time limit is taken to be the default time limit for non-interactive mode (see §6).

N [*t*]

The **N** flag has one optional parameter, with the same meaning at the *t* parameter of the **L** flag. It makes **ee9** run in **NON-INTERACTIVE** mode, suitable for invocation from a command script. In this mode it is not possible to supply responses to prompts, whether from the KDF9 program or from **ee9** itself; so if an interactive input is requested in non-interactive mode, **ee9** terminates with a suitable diagnostic message. If the **N** flag is given without a parameter, or on the command line, the time limit is taken to be the default time limit for non-interactive mode (see §6).

O

The **O** flag outputs the diagnostics (post-run dumps, histograms and traces) requested in `settings_1.txt` when a program terminates and nominates an overlaid successor by means of OUT 1 or OUT 2.

P *address value*

This flag is used to **POKE** the *value* (given in octal or decimal) into the *address*. The *address* takes the form of an octal or decimal number, immediately followed by one of the letters **W**, **U**, **L**, **S** or **C**. If **W** is specified, the *value* overwrites the whole word at the *address*. If **U** or **L** is given, the *value* overwrites the specified halfword. If **S** or **C** is given, it must be followed by the position within the word of the syllable or character, respectively, to be overwritten. If a part-word is specified, the *value* must be within its storable range.

Q

The **Q** flag outputs any specified pre-run dumps and then **quits**. It is useful for getting a disassembly of a program without running it.

R *a b*

This flag is used to set two addresses, say *a* and *b*, that delimit the **RANGE** of instructions where tracing is done. No breakpoint or watchpoint fires, and no instruction is traced, unless $a \leq i \leq b$ is satisfied; where *i* is the address of the word containing the instruction to be executed. With suitable *a* and *b* values, instruction tracing can be confined to the sequence of instructions that you are currently debugging.

S *start [end]*

This flag has either one or two parameters, which are data-word addresses. It sets a **STOREPOINT** on every word in the given range of addresses.

T **BOOT_MODE** | **PROGRAM_MODE** | **PRIVILEGED_MODE**

This flag is used to set the run state.

V{**ABDEFHIKPRSTWZ**}

The **V** flag is used to set the **VISIBILITY** of diagnostic output. It is immediately followed by a selection of the letters: **A** to suppress Director API messages, **E** to suppress confirmatory or warning messages—but not error messages—from **ee9**, **F** to suppress the **FINAL STATE** of the KDF9 at the end of a run, **H** for the **HISTOGRAMS**, **I** for the **INTERRUPT** trace, **P** for the **PERIPHERAL I/O** trace, **R** for the **RETRO** trace, **S** for the digital **SIGNATURE**, AND **W** for console Flexowriter output.

A trace is output if it is provided by the requested diagnostic mode, **and** its output is not suppressed.

The default is that all traces provided by the diagnostic mode are to be output, i.e. **not** suppressed.

The option **B** installs the **BSI** interface, **SI0**, on buffer 15/#17.

The option **C** enables the output of a **Core** image file in the event of abnormal termination.

The option **D** enables the output of any optional **Debugging** messages.

The option **K** installs the drum, **DR0**, on buffer 14/#16.

The option **T** enables execution with authentic **TIMING** (see §2.4).

Spaces and tabs after the **V** are ignored. **V** by itself restores the default options.

W *start [end]*

This flag has either one or two parameters, which are data-word addresses. It sets a **WATCHPOINT** (i.e., a **FETCHPOINT** **and** a **STOREPOINT**) on every word in the given range of addresses. Note that it is not possible to set a **FETCHPOINT** only.

X

The **X** flag over-rides other diagnostic options and runs the program in external tracing mode with only signature generation enabled.

Y

The **Y** flag sets up an operand symbol table entry to be used when disassembling orders during tracing and when printing Usercode format core dumps. **Y** by itself clears out the symbol table. See Appendix 6 for a full description of this feature.

/ | -

A line beginning with / or with - is taken to be comment and ignored.

5.2: THE MISCELLANY PARAMETER

The options permitted with the miscellany parameter on the command line are as follows:

abcdefghijklmnopqrstuvwxyz.0123456789-

The miscellany parameter is put into effect from left to right; it is case-insensitive. The letters **gnox** correspond to the option flags of the same name, with the defaults stated in §5.1 for their parameters. The letters **abcdefghijklmnopstvw** correspond to the options that can be given with the **V** flag, as defined in §5.1. A digit *d* requests an execution time limit of $(d+1) \times 100_000_000$ instructions; '.' specifies a time limit of 1_000_000. Option **y** specifies that the program is to be run as a 'bare' Director. Option **z** specifies suppression of all inessential output. A hyphen is ignored.

6: IMPLEMENTATION CHARACTERISTICS AND CAVEATS

CHARACTERISTICS

The defaults for the settable options in the present implementation of **ee9** are as follows:

- the default run mode is `PROGRAM_MODE`
- the default diagnostic mode is `FAST_MODE`
- the default diagnostic visibility generates all traces, the digital signature, and the histograms
- the default is to run interactively; that is, with non-interactive mode disabled
- the default elapsed time mode is `MODERN_TIMES_MODE`
- the default time limit allows for effectively unlimited execution
- the default time limit in non-interactive mode is 100 million instruction executions
- the default count bounds, *l* and *h*, are 0 and the time limit, respectively
- the default range bounds, *a* and *b*, are 0 and #17777 (8191), respectively
- no breakpoint, watchpoint, storepoint, or core dump is pre-set
- the shared buffer is switched by default to TP1, not GP0
- the default graph plotter pen colour is `BLACK`
- the default graph plotter pen tip is `EXTRA_EXTRA_FINE` (1 plotter step wide)

CAVEATS

Do not be discouraged by the length of this list; **ee9** successfully runs all distributed KDF9 programs, whether surviving, resurrected or newly-written.

- TSD OUTs 15 and 46 are not yet implemented.
- There are some bogus OUTs, with numbers distinct from those of any historically valid OUT, that are implemented by **ee9** in program mode for various reasons of convenience. They may be withdrawn in future releases.
- KDF9's nest-depth checking caused a NOUV interrupt **after** the maximum or minimum depth had been transgressed. Presently, **ee9** checks for all of these violations **before** the offending instruction is executed. This makes little difference in practice. KDF9 had 'imprecise' interrupts, which made recovery from a NOUV error impossible: Director could do no more than terminate (or perhaps restart) the offending program. The behaviour of **ee9** makes debugging a bit easier.
- There is some doubt as to the semantics of the various division instructions, particularly with respect to rounding, and their results on overflow and on division by zero (other than setting the overflow bit).
- There is some doubt as to the semantics of the FLOAT and FLOATD instructions when the scale factor operand in N1 is outside the range $-128 \leq N1 \leq +127$ specified by the Manual. EE issued a library subroutine, L77, to deal with such cases, so it seems that the behaviour of the hardware with out of range operands was not entirely satisfactory. **ee9** takes a safety first approach and uses the least significant 8 bits of N1 as the effective scale factor, ignoring the rest of N1. This is compatible with the programming of the `sqrt` function in the Kidsgrove Algol system.
- The FRB and TOB orders deliver results that differ from the KDF9 hardware in unusual cases involving radix values greater than 31, which are stated to be illegal in the Manual, and to malfunction in the hardware documentation.
- The all-zero syllable is documented as an invalid order, but strong evidence has come to light that the hardware actually treated it as a "no-op", like DUMMY, so **ee9** now does the same, identifying it in traces as DUMMY0.
- The Standard Interface (SI) device is present as a 'best guess' implementation, pending more information about it.
- The POE and POF orders, for TP and SI, produce NULs in lieu of paper tape runoff; this is suppressed in Latin-1 mode.
- It is assumed that the POF order for TP and SI is the same as the POE order, but that the count is multiplied by 8.
- The POK and POL orders are fully implemented for the CP and FD devices. The PMK and PML orders are fully implemented for 7-track Ampex tape decks. There is considerable doubt as to the effects of the PMG, PMK, PML, POK, and POL orders for all other devices.
- The CT order lets an ongoing transfer terminate normally, as if it were MANUAL; the hardware forced an immediate end.
- In Latin-1 mode the PIC and PID orders for FW, TR, TP and SI0 copy each input byte *verbatim* to the least significant 8 bits of a word; POC and POD for those devices copy the least significant 8 bits of each word *verbatim* to a byte of output.
- PID and POD transfers for FW, TR, TP and SI0 devices terminate after a word that has the 8-bit paper tape code for an EM, decimal 61, #75, in its least significant 6 bits.
- The POC and POD orders switch FW from writing to reading after a word that has the 8-bit paper tape code for a semicolon, decimal 60, #74, in its least significant 6 bits.
- **ee9** cannot fully implement the semantics of 5-hole paper tape. See Appendix 2C for more details.
- It is assumed that the graph plotter pen tip sizes and colours are the same as those of pens presently on sale.
- It is assumed that the plotter command codes listed in Appendix 5, §3, p.304 of the Manual are wrong, as there is an error whereby 11 plotting codes are claimed but only 9 are given. The codes used by **ee9** are given in Appendix 3 of this *Guide*. These have been confirmed by evidence from other computers of the era that had the same type of plotter.

- It is assumed that the POB and POD orders for the graph plotter have the same effect as POA and POC, respectively. The Manual says that the plotter responds to EM with a ‘peculiar’ pen movement, but does not end the transfer, and that invalid codes have ‘unpredictable’ effects. **ee9** in these cases implements an arbitrary, but safe and potentially useful, operation: it moves the plotting position to the origin of the plotter co-ordinate system.
- **ee9** supports 4 disc drives. It is assumed that the fixed-head area of a disc drive is platter 16, seek area 0.
- It is assumed that a drum has 40 tracks of 8 sectors. The drum storage supported by **ee9** includes 4 such drum units.
- Many other hypotheses have been made in emulating discs and drums; it remains to be seen whether they are justified.
- The emulation of OUT 8 by **ee9** has some differences from the authentic behaviour of TSD; see Appendix 5.
- When a program running directly uses OUT 8 to write a prompt to FW0 (a ‘query’ in KDF9 parlance, including a ‘;’ that turns the write into a read), **ee9** precedes the given data with ‘[q]’; but if the transfer is a simple output message **ee9** precedes it with ‘[m]’. This imitates the behaviour of the **non**-Time Sharing (‘standard’) Director. Also, such a query always terminates on reading EM, but Directors insist that the EM be preceded by a full stop.
- I/O lockouts in **ee9** remain in force for the whole duration of a transfer and are cleared *in toto* at the end; KDF9 cleared the lockout of each 32-word group as soon as the transfer crossed the boundary into the next group.
- There remain some issues related to the execution of problem programs under TSD control.
 - (a) The CPU time reported by TSD may differ by about 1% from that reported by running directly under **ee9**.
 - (b) TSD reports ludicrous execution times in the order of 800 minutes for some programs that run in a few seconds.
 - (c) The coordination of a FW0 transfer initiated directly by a problem program, with TSD’s typing-out of its own messages, is not satisfactory. Programmers were strongly discouraged from direct output to FW0 for this very reason.
- Failures due to programming error can occur during a run. Some are required by the KDF9 architecture to cause a LIV interrupt in problem program state, and **ee9** implements that. LIV interrupts are inhibited in Director state, but **ee9** does not allow execution to continue, as that might cause unpredictable failures. Other failures were not detected by the KDF9 hardware, but are apparent to **ee9**, which acts appropriately. Impossible I/O operations by Director end the run, because this indicates a serious malfunction. In boot mode they set the device abnormal and abandon the order; the onus is on the program to detect this and act accordingly—if it does not it may subsequently LIV. Failures when running a program directly end the run with a specific diagnostic. See Appendix 8 for a list of the various kinds of error message.

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REFERENCES AND FURTHER READING

Included with distributions of **ee9**:

- *The English Electric KDF9, The Hardware of the KDF9, The Software of the KDF9, The KDF9 and Benchmarking, The Purchase of a KDF9 in 1964, and The KDF9: a Bibliography*
- *ee9 Users’ Guide, ee9 Implementation Overview and ee9 Release History*
- *Kidsgrove Code Procedures, Kidsgrove Compilation Errors, Kidsgrove Compilation Limits, Kidsgrove Compiler Options, Kidsgrove IO Libraries, Kidsgrove Optional Output and Kidsgrove Runtime Errors*
- *README*
- *Whetstone Compilation Errors, Whetstone Controller-listing, Whetstone Runtime Errors and Whetstone Translator-listing*

Available at: www.findlayw.plus.com/KDF9:

- *KDF9 Programming Manual*, Publication 1002 mm ® 1001263; International Computers Ltd.; ‘the Manual’
- *KDF9 ALGOL programming*, Publication 1003 mm ® 1000565; English Electric–Leo Computers Ltd.

SEE ALSO

www.findlayw.plus.com/KDF9/#Emulator

which is updated occasionally with **ee9** news.

APPENDIX 1: **ee9** COMMAND FUNDAMENTALS AND USAGE OPTIONS

The emulator is invoked from the command line, thus:

```
ee9 { -ss | -dd | -m[m] | -TPtt | -TRtt } +program_file_name
```

where the parameters introduced by ‘-’ are optional and can be given in any number or order; *m* is a string that specifies a miscellany of options; *d* specifies a diagnostic execution mode; *s* is the initial CPU state for the KDF9 run; and *tt* specifies paper tape device character codes. Note that *program_file_name* is a path name; unlike the *prog* parameter of the **nine** shell command it is not by default made relative to **Testing/Binary**.

The allowable state flag characters *s* are:

- **b**: for **booting** into Director state, which is how operating systems such as the TSD are loaded and run
- **p**: for **problem** program state, the default, allowing user programs to be run without a Director
- **t**: for **privileged** program state, allowing programs to be run with OUTs serviced as in problem program state, but allowing privileged instructions; though inauthentic, this is useful for running ‘hardware test’ programs.

For the diagnostic flags *d*, see the **nine** command in §1. The miscellany parameter *m* is described in §5.2.

The parameters *tt* specify the character code to be used by either one or two paper tape punches (with **-TP**), or paper tape readers (with **-TR**). *tt* is either one or two characters, each being either **k/K** or **l/L**, specifying KDF9 paper tape code or Latin-1 code respectively. If *tt* is one character, it refers to unit 0 of the type (i.e. **TP0** or **TR0**). If *tt* is two characters, the first refers to unit 0 and the second to unit 1 of the type. See also §3.4 and Appendix 2. The current working directory in which **ee9** is invoked must contain files to represent the KDF9’s I/O devices. See §3. Invoking **ee9** is not trivial, so there are shell commands to do the donkey work for you. These are described in §1.

The **ee9** distribution contains a hierarchy of directories to locate its files systematically and conveniently for the casual user. **Testing**, and its **Adjuncts**, **Assembly**, **Binary**, **Data**, **Kidsgrove**, **Pascal**, **Whetstone** and **Workfiles** subdirectories contain materials for KDF9 programs to use. The **Adjuncts** directory contains components used to run the Algol and Pascal compilers and the Time Sharing Director. Temporary workfiles are created in, and deleted from, **Workfiles**. However, all of these locations are defaults only, and others can be substituted if that better meets the user’s needs. To specify alternative places shell environment variables may be exported, as follows.

ADJUNCTS	Adjuncts
DIRECTOR	Adjuncts/KKT40E007UPU
KDF9_USERCODE	Assembly
KDF9_BINARY	Binary
KDF9_DATA	Data
KIDSGROVE	Kidsgrove
PASCAL	Pascal
WHETSTONE	Whetstone
ALGOL_TEMP	Workfiles
environment variable	default path

E.g., to nominate a directory with pathname *d*, for stropped Algol source files (see Appendix 12), issue the command:

```
export KIDSGROVE=d
```

If a KDF9 program overlays itself by means of OUT 1 (see Appendix 4), the executable is taken from the directory specified by **KDF9_BINARY**. If a program interactively requests additional data and the user gives the ‘@’ response (see §3.3), the file is looked for in the directory specified by **KDF9_DATA**. **DIRECTOR** specifies the version of TSD used by **tsdnine**.

In each case the default path is used if the variable is unset (either undefined or containing the empty string).

The defaults conform to the distributed structure when working in **Testing**, where most executable binaries and shell files are normally held. Alternatively, source code and data files can be conveniently kept together in one directory with pathname *path*, independently of the **Testing** hierarchy, by setting its pathname in the current shell instance, thus (note the **.**):

- **workspace path**

ee9 colourizes Flexwriter output using ANSI terminal SGR escape codes. The codes used by the default ‘lite’ mode are for red and black text on a white background, but you can change the codes **ee9** uses by exporting their values in environment variables as follows, where **ee9** replaces the word “ESC” in the value with an ASCII Escape character.

RED_FONT	SGR code for ‘red’ FW output	ESC[31m	<i>unset</i>	NULL
BLACK_FONT	SGR code for ‘black’ FW input	ESC[37m	<i>unset</i>	NULL
UNDERLINE	SGR code for FW underlined text	ESC[4m	<i>unset</i>	NULL
PLAIN_FONT	SGR code for FW plain text	ESC[0m	<i>unset</i>	NULL
variable	defines	dark mode	lite mode	non-SGR terminal

The effects listed can be obtained for runs in the current shell instance with a command such as:

- **set_f_codes**

where *f* is one of **dark**, **lite** or **null**. The ‘dark’ mode gives red and white text on a black background. If you prefer another colour scheme, set the variables as you please.

APPENDIX 2A: KDF9 INTERNAL CHARACTER CODES AND THEIR TRANSCRIPTIONS

Line printer	SP		LF	FF			%	'
Punched Cards	SP	ISO:"	ISO:®	ISO:©	ISO:¬	ISO:#	%	'
Normal Case	SP	ISO:"	LF	FF	HT	ISO:#	CS ISO:β	CN ISO:ñ
Shift Case	SP	ISO:"	LF	FF	HT	ISO:#	CS ISO:β	CN ISO:ñ
Octal code	00	01	02	03	04	05	06	07

Line printer	:	=	()	£	*	,	/
Punched Cards	:	=	()	£	*	,	/
Normal Case	ISO:&	ISO:?	ISO:!	ISO:%	ISO:'	ISO:\$	ISO:~	/
Shift Case	ISO:&	ISO:?	ISO:!	ISO:%	ISO:'	ISO:\$	ISO:~	:
Octal code	10	11	12	13	14	15	16	17

Line printer	0	1	2	3	4	5	6	7
Punched Cards	0	1	2	3	4	5	6	7
Normal Case	0	1	2	3	4	5	6	7
Shift Case	↑ ISO: ^	[]	<	>	=	×	÷
Octal code	20	21	22	23	24	25	26	27

Line printer	8	9		₁₀ ISO:ø	;	+	-	.
Punched Cards	8	9	_ ISO: _	₁₀ ISO:ø	;	+	-	.
Normal Case	8	9	_ ISO: _	₁₀ ISO:ø	;	+	-	.
Shift Case	()	_ ISO: _	£	;	≠ ISO:±	*	,
Octal code	30	31	32	33	34	35	36	37

Line printer		A	B	C	D	E	F	G
Punched Cards	ISO:@	A	B	C	D	E	F	G
Normal Case	ISO:@	A	B	C	D	E	F	G
Shift Case	ISO:@	a	b	c	d	e	f	g
Octal code	40	41	42	43	44	45	46	47

Line printer	H	I	J	K	L	M	N	O
Punched Cards	H	I	J	K	L	M	N	O
Normal Case	H	I	J	K	L	M	N	O
Shift Case	h	i	j	k	l	m	n	o
Octal code	50	51	52	53	54	55	56	57

Line printer	P	Q	R	S	T	U	V	W
Punched Cards	P	Q	R	S	T	U	V	W
Normal Case	P	Q	R	S	T	U	V	W
Shift Case	p	q	r	s	t	u	v	w
Octal code	60	61	62	63	64	65	66	67

Line printer	X	Y	Z					
Punched Cards	X	Y	Z	ISO:{	ISO:}	→ ISO:	ISO:\	Ø
Normal Case	X	Y	Z	ISO:{	ISO:}	→ ISO:	ISO:\	see note
Shift Case	x	y	z	ISO:{	ISO:}	→ ISO:	ISO:\	see note
Octal code	70	71	72	73	74	75	76	77

NOTES

The transcription provides a Latin-1 representation for every KDF9 internal character code:

- Code 77 is represented by Ø on punched card devices and fast storage devices invariably, on two-shift devices for 'character mode' transfers only, and is otherwise suppressed on both input and output.
- An empty cell indicates a code that is completely suppressed by the line printer.
- SP is a blank space.
- LF is Line Feed, Line Shift (LS) in KDF9 terminology; it prints as ® in NEST displays, *etc.*
- FF is Form Feed, Page Change (PC) in KDF9 terminology; it prints as © in NEST displays, *etc.*
- HT is Horizontal Tab; it prints as ¬ in NEST displays, *etc.*
- CS is Case Shift.
- CN is Case Normal.
- 'y ISO:x' indicates that x is the ISO Latin-1 transcription of the non-Latin-1 KDF9 character y.
- 'ISO:x' indicates that x is the ISO Latin-1 external representation of a non-legible KDF9 character.
- Fast storage devices (e.g. magnetic tapes) always use the Normal Case representation.
- Except for 'character mode' output, β and ñ are acted upon by the Flexowriter, not transferred literally, so that output is presented in the correct case.

APPENDIX 2B: KDF9 PAPER TAPE CODES (WITH PARITY CHANNELS) AND THEIR TRANSCRIPTIONS

Line printer	SP		LF	FF		%	'	
Punched Cards	SP	ISO:"	ISO:®	ISO:©	ISO:¬	ISO:#	%	'
Normal Case	SP	ISO:"	LF	FF	HT	ISO:#	CS ISO:ß	CN ISO:ñ
Shift Case	SP	ISO:"	LF	FF	HT	ISO:#	CS ISO:ß	CN ISO:ñ
Paper tape	220	021	022	003	024	005	006	027

Line printer	:	=	()	£	*	,	/
Punched Cards	:	=	()	£	*	,	/
Normal Case	ISO:&	ISO:?	ISO:!	ISO:%	ISO:'	ISO:\$	ISO:~	/
Shift Case	ISO:&	ISO:?	ISO:!	ISO:%	ISO:'	ISO:\$	ISO:~	:
Paper tape	030	011	012	033	014	035	036	017

Line printer	0	1	2	3	4	5	6	7
Punched Cards	0	1	2	3	4	5	6	7
Normal Case	0	1	2	3	4	5	6	7
Shift Case	↑ ISO: ^			<	>	=	×	÷
Paper tape	060	041	042	063	044	065	066	047

Line printer	8	9		₁₀ ISO:º	;	+	-	.
Punched Cards	8	9	_ ISO: _	₁₀ ISO:º	;	+	-	.
Normal Case	8	9	_ ISO: _	₁₀ ISO:º	;	+	-	.
Shift Case	()	_ ISO: _	£	;	≠ ISO:±	*	,
Paper tape	050	071	072	053	074	055	056	077

Line printer		A	B	C	D	E	F	G
Punched Cards	ISO:@	A	B	C	D	E	F	G
Normal Case	ISO:@	A	B	C	D	E	F	G
Shift Case	ISO:@	a	b	c	d	e	f	g
Paper tape	120	101	102	123	104	125	126	107

Line printer	H	I	J	K	L	M	N	O
Punched Cards	H	I	J	K	L	M	N	O
Normal Case	H	I	J	K	L	M	N	O
Shift Case	h	i	j	k	l	m	n	o
Paper tape	110	131	132	113	134	115	116	137

Line printer	P	Q	R	S	T	U	V	W
Punched Cards	P	Q	R	S	T	U	V	W
Normal Case	P	Q	R	S	T	U	V	W
Shift Case	p	q	r	s	t	u	v	w
Paper tape	140	161	162	143	164	145	146	167

Line printer	X	Y	Z					
Punched Cards	X	Y	Z	ISO:{	ISO:}	→ ISO:	ISO:\	Ø
Normal Case	X	Y	Z	ISO:{	ISO:}	→ ISO:	ISO:\	see note
Shift Case	x	y	z	ISO:{	ISO:}	→ ISO:	ISO:\	see note
Paper tape	170	151	152	173	154	175	176	157

NOTE

- Internal code 77 is represented by Ø on punched card devices and fast storage devices invariably, on two-shift devices for 'character' mode transfers only, and is otherwise suppressed on both input and output.
- 'Character' mode transfers on two-shift devices in KDF9 mode use the full 8-bit paper tape code and transfers to End Message terminate after the 8-bit internal code #175 (which is the Latin-1 code for '}', not '|').

APPENDIX 2C: KDF9 5-HOLE (FERRANTI) PAPER TAPE CODES

Normal Case	ISO:®	A	B	C	D	E	F	G
KDF9 code	40	41	42	43	44	45	46	47
Letters Case	<i>fs</i>	A	B	C	D	E	F	G
Figures Case	<i>fs</i>	1	2	*	4	()	7
Ferranti code	00	01	02	03	04	05	06	07

Normal Case	H	I	J	K	L	M	N	O
KDF9 code	50	51	52	53	54	55	56	57
Letters Case	H	I	J	K	L	M	N	O
Figures Case	8	≠	=	-	ν	<i>lf</i>	<i>sp</i>	,
Ferranti code	10	11	12	13	14	15	16	17

Normal Case	P	Q	R	S	T	U	V	W
KDF9 code	60	61	62	63	64	65	66	67
Letters Case	P	Q	R	S	T	U	V	W
Figures Case	0	>	≥	3	→	5	6	/
Ferranti code	20	21	22	23	24	25	26	27

Normal Case	X	Y	Z	ISO:{	ISO:}	→ ISO:	ISO:\	see notes
KDF9 code	70	71	72	73	74	75	76	77
Letters Case	X	Y	Z	<i>ls</i>	.	?	£	<i>er</i>
Figures Case	x	9	+	<i>ls</i>	.	<i>n</i>	<i>cr</i>	<i>er</i>
Ferranti code	30	31	32	33	34	35	36	37

NOTES

The handling of 5-hole paper tape—in the teleprinter code of the Ferranti Pegasus, Sirius and Orion computers—is rather non-obvious. When a 5-hole tape was read by KDF9 the most significant bit of each input character was set to 1 by the hardware, adding #40 to the physical Ferranti code. On output to a 5-hole tape punch the hardware discarded that bit. This provided a 1-to-1 mapping and had the advantage that upper case letters got the same codes from 5-hole and 8-hole input tapes, namely the KDF9 native codes for those letters. **ee9** cannot implement this behaviour for direct transfers, because it does not always know whether a paper tape device is in 5-hole mode or 8-hole mode (for example, the type of the device attached to a buffer can be changed dynamically by operators' command to a Director) and it has no way to find out. Consistency demands that transfers that are notionally from 5-hole tape readers and to 5-hole tape punches work in exactly the same way as for 8-hole devices in **ee9**.

Things are very different when a program running under Director control writes to a 5-hole punch via OUT 8 streams 50-57. Most KDF9 characters with codes in the ranges #00-#37, and the letters A-Z, are output verbatim. Others are given special treatment. Director (not **ee9**) converts them in an attempt to preserve the appearance of a teleprinter transcription, inserting shift characters where necessary. Director also converts KDF9's LS character into a *fs cr lf* sequence. Clearly, this has the potential for muddle if the KDF9 data includes characters that are absent from the Ferranti teleprinter set.

The **mtp** and **a2b** utilities (see Appendices 10 and 11) have facilities for converting from Ferranti code to Latin-1. They assume that the Ferranti codes are those generated by OUT 8:

- Ferranti code 00, denoted *fs*, is the Figures Shift character. It switches printing to the Figures Case character set.
- Ferranti code 33, denoted *ls*, is the Letters Shift character. It switches printing to the Letters Case character set.
- Ferranti code 14 is the ν character in Figures Case, but L in Letters Case. The ν was treated by English Electric software as if it were the KDF9 10 character. It is rendered as 'ν' by **mtp** and **a2b**.
- Ferranti code 15, denoted *lf*, is the Line Feed character in Figures Case but M in Letters Case.
- Ferranti code 16, denoted *sp*, is the blank space character in Figures Case but N in Letters Case.
- Ferranti code 22 is the ≥ character in Figures Case, but R in Letters Case. It is rendered as ']' by **mtp** and **a2b**.
- Ferranti code 24 is a right-arrow character in Figures Case, but T in Letters Case. It looks like, but is not, a KDF9 EM. It is rendered as '»' by **mtp** and **a2b**.
- Ferranti code 35 is the *n* character in Figures Case, but ? in Letters Case. This Ferranti code works like EM for 5-track data, since it generated the KDF9 EM internal code (#75) when read by the hardware. In Figures Case it was treated by English Electric software as if it were an underlined EM. It is rendered as 'n' by **mtp** and **a2b**.
- The KDF9 LS character has the same CR/LF effect as the ASCII Newline, but in Ferranti code there are separate CR and LF characters, available only in Figures Case.
- Ferranti code 36, denoted *cr* above, is the Carriage Return character in Figures Case but £ in Letters Case.
- Ferranti code 37 was used to erase paper tape characters by punching holes in all five positions of the tape frame. It is suppressed on output by **mtp** and **a2b**, exactly as #77 is for devices working with KDF9 codes.

APPENDIX 3: KDF9 GRAPH PLOTTER CODES

Action	Decimal	Octal	Binary
None	0	#0	000000
Step paper back	1	#1	000001
Step paper forward	2	#2	000010
Step pen right	4	#4	000100
Step pen right, paper back	5	#5	000101
Step pen right, paper forward	6	#6	000110
Step pen left	8	#10	001000
Step pen left, paper back	9	#11	001001
Step pen left, paper forward	10	#12	001010
Lower pen	16	#20	010000
Raise pen	32	#40	100000

NOTES

• The command list in Appendix 6, §5, p.302 of the Manual has an obvious error whereby 11 plotting codes are claimed but only 9 are given, the last of them being inconsistent with the others. The present list provides the two missing codes and corrects the incorrect code for a step right and backwards. The codes listed here have been confirmed by the evidence of other computers of the era that used the same plotter; the bit positions and their interpretations correspond directly to the control bits in the Calcomp plotter's hardware interface. All other 6-bit character codes represent invalid plotter commands.

APPENDIX 4: TIME SHARING DIRECTOR OUTS (SYSTEM CALLS)

Director provides an API that is accessed by the OUT instruction with parameters in the NEST. See Table A.4.1.

TABLE A4.1: DIRECTOR OUT PARAMETERS

N1	N2	N3	Action
0†			Terminate this run normally.
1†	Program name	Terminate this run and overlay the program whose name is in N2-3.
2†	Time limit in seconds		Restart this program. Allow it the new time limit given in N2.
3†			Return the CPU time used in N1. The result is given as seconds to 23 integral places.
4‡	Tape label		Allocate the tape deck with the reel having the 1-word label in N2. Return its buffer number in N1.
5‡	Device type code		Allocate an I/O device of the type given in N2; see Table A4.2. Return its buffer number in N1.
6‡	Buffer number		Deallocate the I/O device with buffer number in N1. If it is a tape deck, unload any mounted tape reel. (Tape decks are treated the same by OUT 6 and OUT 7 in ee9.)
7‡	Buffer number		Deallocate an allocated tape deck with buffer number in N1. Do not unload any mounted tape reel.
8	Control word		Spool output or write to FW0. See Appendix 5.
9,			Return the time of day in N1. The result is given as seconds since midnight to 23 integral places.
10 ‡	Tape label	Allocate the tape deck with the reel having the 2-word label in N2-3. Return its buffer number in N1 and the 'Tape Serial Number' (TSN), as read from the label block, in N2.
11 *	Control word		Write core to allocated sectors of drum storage.
12 *	Control word		Read allocated sectors of drum storage into core.
13 *	Number of sectors		Allocate a number of drum sectors for exclusive use by this program run. Can be done once only per run.
14 *			Return the number of drum sectors available for allocation in N1, with the sign bit set if OUT 13 has already been executed.
16 **	Control word		Write to FW0, as if by OUT 8, but with a prefix of 'N9' . See Appendix 5.
17 ‡			Return the CPU Time in N1 and the Notional Elapsed Time in N2. The results are given as seconds to 23 integral places.
41 ††	Control word		Write to sector(s) in the currently selected set of disc platters.
42 ††	Control word		Read from sector(s) in the currently selected set of disc platters.
43 ††	0 or -1		Select the first (0) or the second (-1) set of disc platters for transfers.
44 ††	Number of platters		Allocate a number ≤ 8 of disc platters. Can be done only twice per run.
45 ††	0 or -1		Deallocate all platters in the first (0) or the second (-1) set.
47 ††			Check the last FD transfer: set TR if a parity error was flagged.

NOTES

† See the Manual, §26.2. 1.

- Obeying OUT with an empty NEST has the same effect as OUT 0.
- OUT 1 overlays the current program with one read from external media. Its name consists of 12 Case Normal characters, e.g. “KMW0301--UPU”.
- OUT 2 resumes the execution of the program with a new time limit, but without its previously allocated devices.

‡ See the Manual, §17.3.

- The label for OUT 4 consists of 8 Case Normal characters. e.g. “PRINTEND”.
- The label for OUT 10 consists of a plus sign followed by 15 Case Normal characters. e.g. “+KOUTPUT/0023004”.

* See the Manual, Appendix 6, §4.

- The control word for the drum I/O OUTs is of the form *Q starting sector / low core address / high core address*.
- The sectors allocated by OUT 13 are numbered from 0.

** OUT 16 writes to the console Flexowriter in a slightly different format from OUT 8.

The control word for OUT 16 takes the same form as for OUT 8, but only the FW0 stream is supported.

†† See the Manual, Appendix 6, §6.

- The control word for the disc I/O OUTs is of the form *Q starting sector / low core address / high core address*.
- The sectors allocated by OUT 44 are numbered from 0.

TABLE A4.2: DEVICE TYPE CODES

Type	FW	TP	TR	LP	CR	TP §	CP	GP	SI	PDP-8 *	MT †
Decimal	0	1	2	3	4	5	7	16	17	53	55
Octal	#00	#01	#02	#03	#04	#05	#07	#20	#21	#65	#67

NOTES

- Add 8 (#10) to codes 1-7 to reclaim a previously-relinquished device.

§ Type 5 is for a 5-channel tape punch, using Ferranti paper tape code.

* The PDP-8 front end for Eldon 2 and COTAN was attached to a magnetic tape buffer.

† A magnetic tape deck given type #67 is treated as having an unlabelled tape, which is not to be read before it is labelled.

APPENDIX 5: OUT 8

Director implements an output spooling system for paper tape punches and line printers. It also administers access to the console Flexowriter. Up to 32 virtual output streams are provided to each running problem program; see Table A5.1.

Stream output is written first to a designated magnetic tape. When that becomes full, or at the behest of the computer operator, it is rewound and the accumulated data for each stream is transferred from the tape to its associated device. A second magnetic tape may be designated to take output while this is happening, and the two tapes may switch rôles as often as necessary to transfer all the outputs to their intended destinations.

OUT 8 takes a parameter word in N2 that controls the spooling action. It is in Q store format. If the C-, I-, and M-parts are equal, and contain a stream number, that stream is closed. Otherwise the I-part is taken to be the starting address, a , of an output area and the M-part as its ending address, b .

If $b = a + 1$ and the area is laid out as follows:

- word $a+0$: *stream number*
- word $a+1$: control word in Q store format: 4095 / -1 / n

then a device-dependent ‘gap’ is written to the stream; for a tape punch stream this takes the form of runout (unpunched) tape, of length n characters, for $n / 10$ inches; applied to a line printer stream it produces a PC character.

Otherwise, and more usually, the words in address range a to b inclusive are written to the OUT 8 tape by a transfer-to-End-Message operation, the word at a having been overwritten by ‘red tape’ identifying the stream.

TABLE A5.1: OUT 8 STREAMS IMPLEMENTED BY DIRECTOR

stream number	Output Device
#00	Flexowriter
#10 .. 17	An 8-hole paper tape punch selected by the operator
#20 .. 27	Not used
#30 .. 37	The line printer
#40 .. 47	Not used
#50 .. 57	A 5-hole paper tape punch selected by the operator
#60 .. 67	Not used
#70 .. 77	The line printer or a paper tape punch, at the operator’s discretion

Special provision is made for queries, prompted responses on the Flexowriter console, which are implemented immediately, not deferred. In this case the parameter word in N2 must have a C-part of #100000. The output area must be laid out with a prompt (ending with a ‘;’ character) and be followed by an input area for the response. Various restrictions are imposed to preserve a tidy layout of the console’s typed log. In particular, word a is overwritten by Director with suitable format effectors. When running directly, an OUT 8 query always terminates on reading EM, but Director also requires the EM to be preceded by a period, thus ‘. |’.

In its problem program and privileged program modes, **ee9** emulates OUT 8 not by spooling, but by directly transferring the output to the intended destination device. Note that outputs to multiple streams intended for the same device are therefore interleaved by **ee9**, but are not interleaved by Director. In doing so **ee9** is somewhat more specific than Director in its provision of output streams, as indicated in Table A5.2. Streams 11, 13, 15 and 17 are directed to TP1 instead of TP0, as a small mitigation of the interleaving issue.

The 5-hole mode is not implemented by **ee9**; all paper tape output is in 8-hole mode.

Director requires a line of output to be written in a single OUT 8 call, the line ending in a LS, PC or EM character. **ee9** allows a line to be built up incrementally, in a series of calls to the emulated OUT 8. This behaviour was necessary to support an earlier version of the reconstructed I/O library in the Kidsgrove Algol run time environment.

So far these differences between **ee9** and Director have not been problematic.

TABLE A5.2: OUT 8 STREAMS IMPLEMENTED BY ee9

stream number	Output Device
#00	Flexowriter
#10, 12, 14, 16	TP0
#11, 13, 15, 17	TP1
#30 .. 37	LP0
#50 .. 57	TP1
#70 .. 77	LP0
others	Invalid

The KDF9 Algol I/O procedures work on notional I/O streams, identified by decimal numbers that are made to look like OUT 8 octal stream numbers. For example, output to Algol stream 30 goes to OUT 8 stream #30. Algol input uses the 20 range of stream numbers which OUT 8 does not use.

APPENDIX 6: DISASSEMBLED MACHINE CODE

U format core printing attempts to present the machine code and data in the designated area in a approximation to Usercode. Words deemed to be data are output in a variety of formats. Instructions are shown with octal and decimal operands, or with Usercode symbolic operands if a symbol table has been provided. A symbol table may be obtained from a **kal3** assembly listing using the ancillary program **extract_symbols**. It generates a file of Y options, which must be copied into the run's settings file if symbolic disassembly is wanted from that run. For convenience, **nine** creates the symbol table and concatenates it with **settings_3.txt** into **settings_1.txt**. Here is an example of the kind of output that can be expected. First the Usercode source text, with the Usercode symbolic names shown in boldface:

```
P102V1;
    V0=Q0/AW1/AW5;
    V1=P[7DC];
    Q13;      =W0;
    SET3;     =RC13;
1;  ZERO;    SHLD+16;SHL+32;   SHA-32;
    JSP100;   ZERO;    SHLD+6;   REV;    SHL+6;
    REV;      SHLD-6; OR;      REV;
    DC13;     J1C13NZ;
    ERASE;    =W4;
    SETB17;   OR;      =W3;
    SETB17;   OR;      =W2;
    V1;       =W5;
    SETB70;   =W1;
    V0;       SET8;    OUT;
    W0;       =Q13;
EXIT1;

P101V3;
    V0=B2037202020;
    V1=F1000;
    V2=B3600 0000 0000 0000;
    Q13;      =V3;
    STAND;    DUP;    ABSF;    V1;    xF;
    FIX;      SET47;   -;      =C13;   SHAC13;
    V3;       =Q13;
    JSP100;   ZERO;    SHLD+30; SHL+6;
    SHLD-30; ERASE;    V0; OR;
    REV;      J1<Z;
EXIT1;
1;  V2; OR;
EXIT1;
```

Here is its symbolic disassembly, again with the Usercode symbolic names shown in boldface:

```
VOP102, E#2700, E1472:
    #0000000667203341 = #000 #000 #006 #335 #006 #341 = " %W0%A" = Q #000000/ #003335/ #003341
    = 115148513 = 6.155297639320E-43 = 0.000000818179394 = Q 0/ 1757/ 1761
V1P102, E#2701, E1473:
    #7777777777777702 = #377 #377 #377 #377 #377 #302 = "00000000" = Q #177777/ #177777/ #177702
    = -62 = -1.918807060892E+28 = -0.0000000000000441 = Q -1/ -1/ -62
P102V1, E#2702, E1474:
    Q13; =W0; (#3334);
    SET3; =RC13;
E#2703/4 (1475/4):
    ZERO; SHLD+16; SHL+32; SHA-32;
    JSP100; (#2742); (LINK=#2704/5);
    ZERO; SHLD+6; REV; SHL+6; REV; SHLD-6; OR; REV; DC13;
    JE#2703/4C13NZ; (1475);
    ERASE; =W4; (#3340);
    SETB17; (15); OR; =W3; (#3337);
    SETB17; (15); OR; =W2; (#3336);
    V1P102; (#2701); =W5; (#3341);
    SETB70; (56); =W1; (#3335);
    VOP102; (#2700); SETB10; (8);
    OUT;
    W0; (#3334); =Q13;
    EXIT1;
    DUMMY0;

VOP101, E#2720, E1488:
    #00000002037202020 = #000 #000 #020 #175 #004 #020 = " 0.000" = Q #000000/ #010175/ #002020
    = 276628496 = 1.478725763829E-42 = 0.000001965563683 = Q 0/ 4221/ 1040
V1P101, E#2721, E1489:
    #2127640000000000 = #105 #175 #000 #000 #000 #000 = "17T" = Q #042575/ #000000/ #000000
    = 76403173228544 = 1.0000000000000E+3 = 0.542877197265625 = Q 17789/ 0/ 0
V2P101, E#2722, E1490:
    #3600000000000000 = #170 #000 #000 #000 #000 #000 = "-" = Q #074000/ #000000/ #000000
    = 131941395333120 = 0.0000000000000E+0 = 0.937500000000000 = Q 30720/ 0/ 0
V3P101, E#2723, E1491: zero
P101V3, E#2724, E1492:
    Q13; =V3P101; (#2723);
    STAND; DUP; ABSF; V1P101; (#2721); xF; FIX; SETB57; (47); -; =C13; SHAC13; V3P101; (#2723); =Q13;
    JSP100; (#2742); (LINK=#2730/2);
    ZERO; SHLD+30; SHL+6; SHLD-30; ERASE; VOP101; (#2720); OR; REV;
    JE#2734/0LTZ; (1500);
    EXIT1;
E#2734/0 (1500/0):
    V2P101; (#2722); OR;
    EXIT1;
```

The symbol table definitions available with the Y option setting are as follows:

Y#V *n* specifies the number, *n*, of the V stores of P0, the 'main program'
 Y#Y *n* specifies the number, *n*, of the Y stores
 YW *a* specifies the address, *a*, of W0
 YY *a* specifies the the address, *a*, of Y0
 YZ *a* specifies the the address, *a*, of Z0
 Y *y a* specifies the the address, *a*, of Yy0 (*y* stands for a capital letter; note the blank space after the Y)
 YP*p* @*a* specifies the start address, *a*, of routine P*p*
 YP*p*V*v* @*a* specifies the start address, *a*, of routine P*p* and gives the number, *v*, of its last V store

When a symbol table is supplied, operand references logged in traces are also shown in symbolic form.

An alternative to **nine** is the command **none** which creates no symbol table. Without a symbol table, **ee9** analyses the program's control flow and uses heuristics to determine whether a given word represents data or instructions. These do a good job, but cannot always be correct. Here is the non-symbolic disassembly of the same code:

```
E#2700, E1472:
#0000000667203341 = #000 #000 #006 #335 #006 #341 = " %W0%A" = Q #000000/ #003335/ #003341
= 115148513 = 6.155297639320E-43 = 0.000000818179394 = Q 0/ 1757/ 1761
E#2701, E1473:
#7777777777777702 = #377 #377 #377 #377 #377 #302 = "00000000" = Q #177777/ #177777/ #177702
= -62 = -1.918807060892E+28 = -0.0000000000000441 = Q -1/ -1/ -62
E#2702, E1474:
Q13; =E#3334;(1756);
SET3; =RC13;
E#2703/4, E1475/4:
ZERO; SHLD+16; SHL+32; SHA-32;
JSE#2742/0;(1506);(LINK=#2704/5);
ZERO; SHLD+6; REV; SHL+6; REV; SHLD-6; OR; REV; DC13;
JE#2703/4C13NZ;(1475);
ERASE; =E#3340;(1760);
SETB17;(15); OR; =E#3337;(1759);
SETB17;(15); OR; =E#3336;(1758);
E#2701;(1473); =E#3341;(1761);
SETB70;(56); =E#3335;(1757);
E#2700;(1472); SETB10;(8);
OUT;
E#3334;(1756); =Q13;
EXIT1;
DUMMY0;
E#2720, E1488:
#00000002037202020 = #000 #000 #020 #175 #004 #020 = " 0.000" = Q #000000/ #010175/ #002020
= 276628496 = 1.478725763829E-42 = 0.000001965563683 = Q 0/ 4221/ 1040
E#2721, E1489:
#2127640000000000 = #105 #175 #000 #000 #000 #000 = "17T" = Q #042575/ #000000/ #000000
= 76403173228544 = 1.0000000000000E+3 = 0.542877197265625 = Q 17789/ 0/ 0
E#2722, E1490:
#3600000000000000 = #170 #000 #000 #000 #000 #000 = "-" = Q #074000/ #000000/ #000000
= 131941395333120 = 0.0000000000000E+0 = 0.937500000000000 = Q 30720/ 0/ 0
E#2723, E1491: zero
E#2724, E1492:
Q13; =E#2723;(1491);
STAND; DUP; ABSF; E#2721;(1489); *F; FIX; SETB57;(47); -; =C13; SHAC13; E#2723;(1491); =Q13;
JSE#2742/0;(1506);(LINK=#2730/2);
ZERO; SHLD+30; SHL+6; SHLD-30; ERASE; E#2720;(1488); OR; REV;
JE#2734/0LTZ;(1500);
EXIT1;
E#2734, E1500:
E#2722;(1490); OR;
EXIT1;
```

Here is what the first of these routines looks like with the decimal address option in force:

```
E1472, E#2700:
#0000000667203341 = #000 #000 #006 #335 #006 #341 = " %W0%A" = Q #000000/ #003335/ #003341
= 115148513 = 6.155297639320E-43 = 0.000000818179394 = Q 0/ 1757/ 1761
E1473, E#2701:
#7777777777777702 = #377 #377 #377 #377 #377 #302 = "00000000" = Q #177777/ #177777/ #177702
= -62 = -1.918807060892E+28 = -0.0000000000000441 = Q -1/ -1/ -62
E1474, E#2702:
Q13; =E1756;(#3334);
SET3; =RC13;
E1475/4, E#2703/4:
ZERO; SHLD+16; SHL+32; SHA-32;
JSE1506/0;(#2742);(LINK=1476/5);
ZERO; SHLD+6; REV; SHL+6; REV; SHLD-6; OR; REV; DC13;
JE1475/4C13NZ;(#2703);
ERASE; =E1760;(#3340);
SET15;(#17); OR; =E1759;(#3337);
SET15;(#17); OR; =E1758;(#3336);
E1473;(#2701); =E1761;(#3341);
SET56;(#70); =E1757;(#3335);
E1472;(#2700); SET8;
OUT;
E1756;(#3334); =Q13;
EXIT1;
```

APPENDIX 7: TRACING

These examples give an indication of the types of diagnostic output that can be obtained.

APPENDIX 7A: THE RETROSPECTIVE TRACE

Under ND is given the NEST depth and under SD the depth of the SJNS. Under VTD are the states of the overflow and test registers and whether running in Director mode. For example:

Retrospective trace of all instructions.

	ND	SD	VTD	CPU TIME	ICR
Ended #23/4: OUT	3	0	V	1674616069	306451954
After #23/3: ZERO	4	0	V	1674616056	306451953
After #23/0: OUT	3	0	V	1674616054	306451952
After #22/3: SET6	5	0	V	1674616041	306451951
After #22/1: C13	4	0	V	1674616029	306451950
After #21/5: POEQ13	3	0	V	1674616024	306451949
After #136/5: EXIT2	3	0	V	1674616005	306451948
After #136/2: JE#137/2±; (95)	3	1	V	1674615986	306451947
After #135/5: SET29; (#35)	4	1	V	1674615981	306451946
After #135/2: JE#133/5LTZ; (91)	3	1	V	1674615977	306451945
After #135/1: DUP	4	1	V	1674615973	306451944
After #135/0: -	3	1	V	1674615971	306451943
After #134/3: SET32; (#40)	4	1	V	1674615970	306451942
After #134/1: SHLD+6	3	1	V	1674615959	306451941
After #134/0: ZERO	3	1	V	1674615956	306451940
After #133/5: ERASE	2	1	V	1674615954	306451939
After #133/4: DUP	3	1	V	1674615953	306451938
After #133/1: YOM7Q; (31488)	2	1	V	1674615951	306451937
After #132/5: POBQ14	1	1	V	1674615944	306451935
After #132/3: PIBQ15	1	1	V	1674615912	306451933
After #132/1: IM15 TO Q14	1	1	V	1674615880	306451931
After #131/5: =M15	1	1	V	1674615876	306451930
After #131/4: +	2	1	V	1674615866	306451929
After #131/2: I15	3	1	V	1674615865	306451928
After #130/5: SET125; (#175)	2	1	V	1674615859	306451927
After #130/3: =RM7	1	1	V	1674615855	306451926
After #130/2: ZERO	2	1	V	1674615852	306451925
After #130/0: =RM6	1	1	V	1674615850	306451924
After #20/5: JSE#130/0; (88)	2	1	V	1674615847	306451923
After #20/2: SETAYB1; (#72765)	2	0	V	1674615836	306451922
After #52/3: JE#20/2=Z; (16)	1	0	V	1674615832	306451921
After #52/1: SHL+24	2	0	V	1674615821	306451920
After #51/4: E0; (#0)	2	0	V	1674615817	306451919
After #51/3: ERASE	1	0	V	1674615803	306451918
After #51/2: ERASE	2	0	V	1674615802	306451917
After #51/0: =RM2	3	0	V	1674615801	306451916
After #50/5: DUP	4	0	V	1674615798	306451915
After #50/4: DUP	3	0	V	1674615796	306451914
After #50/3: ZERO	2	0	V	1674615794	306451913
After #160/1: EXIT1	1	0	V	1674615792	306451912
After #157/5: POAQ14	1	1	V	1674615779	306451911

(etc)

After earlier instructions, whose tracing is now lost.

APPENDIX 7B: THE EXTERNAL TRACE OF ALL INSTRUCTION TYPES

LOCATION	ICR	CPU TIME	ND	SD	VT	[N1]	INSTRUCTION
#00000/0	1	8	0	0			J#00012/0; (8)
#00012/0	2	12	1	0		#0000000000000002	SET2
#00012/3	3	19	2	0		#0000000000000005	SET5
#00013/0	4	32	1	0		#0000000000000002	OUT
#00013/3	5	35	0	0			=C15
...							
#00236/2	3439084	19706693	1	2	V	#0000000000000004	JS#00063/2; (51)
#00063/2	3439085	19706697	2	2	V	#0000000000001243	M2
#00063/4	3439086	19706699	3	2	V	#0000000000001243	DUP
#00063/5	3439087	19706703	2	2	V	#0000000000002506	+
#00064/0	3439088	19706705	1	2	V	#0000000000000004	=M3
#00064/2	3439089	19706709	2	2	V	#0000000000000001	M1
#00064/4	3439090	19706715	1	3	V	#0000000000000004	=LINK
#00065/0	3439091	19706721	2	3	V	#0000000000000000	YA1M3; (#4713)

(etc)

APPENDIX 7C: THE SINGLE-STEP TRACE

In pause mode **ee9** executes instructions one-by-one, interacting with the user at each step and presenting the values in the registers relevant to the instruction just obeyed:

At #01313/4 (715/4); ICR = 182908; EL = 9111955; the instruction was
#200:324:357, i.e. JSP40;(#1313)

JB: #01313/4; SJNS depth: 2

NEST:

N1:

#1766314631463146 = 69709037200998 = 1.999999999998E-1 = 0.4953124999999997
= Q #037546/ #063146/ #063146 = Q 16230/ 26214/ 26214
= #077 #146 #146 #146 #146 #146 = "/V9F9F9F"

Breakpoint: at #02357/0 (d:debug | f:ast | t:race | p:ause or q:uit)? p

At #02357/0 (1263/0); ICR = 182909; EL = 9111957; the instruction was
#042, i.e. DUP

NEST:

N1:

#1766314631463146 = 69709037200998 = 1.999999999998E-1 = 0.4953124999999997
= Q #037546/ #063146/ #063146 = Q 16230/ 26214/ 26214
= #077 #146 #146 #146 #146 #146 = "/V9F9F9F"

N2:

#1766314631463146 = 69709037200998 = 1.999999999998E-1 = 0.4953124999999997
= Q #037546/ #063146/ #063146 = Q 16230/ 26214/ 26214
= #077 #146 #146 #146 #146 #146 = "/V9F9F9F"

Breakpoint: at #02357/1 (d:debug | f:ast | t:race | p:ause or q:uit)?

At #02357/1 (1263/1); ICR = 182910; EL = 9111961; the instruction was
#203:104:366, i.e. JE#02366/3LEZ;(2366)

NEST:

N1:

#1766314631463146 = 69709037200998 = 1.999999999998E-1 = 0.4953124999999997
= Q #037546/ #063146/ #063146 = Q 16230/ 26214/ 26214
= #077 #146 #146 #146 #146 #146 = "/V9F9F9F"

Breakpoint: at #02357/4 (d:debug | f:ast | t:race | p:ause or q:uit)?

At #02357/4 (1263/4); ICR = 182911; EL = 9111967; the instruction was
#045, i.e. FIX

NEST:

N1:

#7777777777777776 = -2 = -6.189700196427E+26 = -0.0000000000000014
= Q #177777/ #177777/ #177776 = Q -1/ -1/ -2
= #377 #377 #377 #377 #377 #376 = "00000000\"

N2:

#3146314631463000 = 112589990684160 = 6.044629096666E+22 = 0.7999999999999272
= Q #063146/ #063146/ #063000 = Q 26214/ 26214/ 26112
= #146 #146 #146 #146 #146 #000 = "9F9F9F8 "

Breakpoint: at #02357/5 (f:ast | t:race | p:ause or q:uit)? q

APPENDIX 7D: THE PERIPHERAL I/O EVENT TRACE

Under CPL is given the TSD context (P, Q, R or S) if the transfer was initiated by a program, or D if initiated by Director; and the current priority level. Under the heading T the value of the T register is shown as Y (for 1) or N (for 0) in the case of an order, such as BUSY, that tests device status. S and E in the next column show the start and end of a transfer.

Retrospective trace of peripheral I/O events.

				CPL	T	EL. TIME	ICR
After #4310/3:	CLOQ7	Q #072500/	#004640/	#077337	D 0 N	2052470495	366874014
After #2520/4:	TWQ7	FW0 Q #000000/	#000472/	#000472	D 0 E	2052465840	366871442
After #2472/0:	BUSYQ0	FW0 Q #000000/	#000000/	#000000	D 0 Y	2051676802	366873352
After #1062/4:	CTQ6	TR1 Q #000002/	#177777/	#000002	D 0 N	2051672349	366872593
After #2520/4:	TWQ7	FW0 Q #000000/	#000472/	#000472	D 0 S	2051665840	366871441
After #2472/0:	BUSYQ0	FW0 Q #000000/	#000000/	#000000	D 0 N	2051665650	366871408
After #21/5:	PGAPQ13	TP0 Q #000004/	#000000/	#000454	P 0 E	2051660310	366870057
After #1534/4:	MANUALQ6	TP0 Q #000004/	#000001/	#000004	D 0 Y	2051660310	366870512
After #1255/2:	PARQ7	TP0 Q #000004/	#000000/	#000000	D 0 N	2051660106	366870474
After #1255/2:	PARQ7	TP0 buffer lockout			P 0	2051660106	366870473
After #623/5:	BUSYQ7	TP0 Q #000004/	#000000/	#000000	D 0 Y	2051654372	366870467
After #21/5:	PGAPQ13	TP0 Q #000004/	#000000/	#000454	P 0 S	2051652220	366870056
After #132/5:	PWEQ14	TP0 Q #000004/	#066440/	#066635	P 0 E	2051660025	366869806
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 E	2051651903	366869805
After #132/5:	PWEQ14	TP0 Q #000004/	#066440/	#066635	P 0 S	2051650935	366869805
After #157/5:	PWQ14	TP0 Q #000004/	#066440/	#066440	P 0 E	2051723445	366869784
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 S	2051650903	366869804
After #157/5:	PWQ14	TP0 Q #000004/	#066440/	#066440	P 0 S	2051650725	366869783
...							
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 E	373054471	59581364
After #132/5:	PWEQ14	TP0 Q #000004/	#066440/	#066635	P 0 E	373010462	59578367
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 S	372953471	59581363
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 E	372945668	59578366
After #132/5:	PWEQ14	TP0 Q #000004/	#066440/	#066635	P 0 S	372937742	59578366
After #132/5:	PWEQ14	TP0 Q #000004/	#066440/	#066635	P 0 E	373001334	59576600
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 S	372937668	59578365
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 E	372936540	59576599
After #132/5:	PWEQ14	TP0 Q #000004/	#066440/	#066635	P 0 S	372928614	59576599
After #20/0:	PGAPQ13	TP0 Q #000004/	#000000/	#000454	P 0 E	375655473	59576588
After #132/3:	PREQ15	TR1 Q #000002/	#066440/	#066635	P 0 S	372928540	59576598
After #20/0:	PGAPQ13	TP0 Q #000004/	#000000/	#000454	P 0 S	372928473	59576587

(etc)

After earlier interrupts, whose tracing is now lost.

Total time waiting for unoverlapped I/O to finish = 6635ms.

APPENDIX 7E: THE INTERRUPT TRACE

Under CPL is the context (P, Q, R or S) and priority level of the interrupted program. In addition, the trace shows control returning from Director to a problem program by means of the EXITD instruction, with the BA, NOL, and instruction address values for the resumed program.

Retrospective trace of interrupt requests.

				CPL	EL. TIME	ICR
Ended #575/4:	EDT	MT0 Q 8/#004335/	#004336	P 0	435959729	69288606
After #76/4:	EDT	FW0 Q 0/#005246/	#005250	P 0	435943366	69285714
After #416/3:	OUT	0		P 0	433847319	69079602
After #116/3:	EXITD	BA #004640 NOL 8191 @ #416/2		P 0	433847299	69079600
After #415/5:	OUT	8		P 0	433827181	69076031
...						
After #116/3:	EXITD	BA #004640 NOL 8191 @ #4430/0		P 0	429984577	68459560
After #4427/3:	CLOCK	1048580 KDF9 us		P 0	429982227	68459073
After #116/3:	EXITD	BA #004640 NOL 8191 @ #4430/0		P 0	428934534	68285222
After #4427/3:	CLOCK	1048584 KDF9 us		P 0	428925791	68283615
After #116/3:	EXITD	BA #004640 NOL 8191 @ #3153/4		P 0	427877727	68109705
After #60/5:	EDT	MT0 Q 8/#011657/	#011702	P 0	427877343	68109622
After #3153/4:	LOV	at #011656 (E2574)		P 0	427877149	68109587
After #116/3:	EXITD	BA #004640 NOL 8191 @ #2056/4		P 0	427877113	68109584
After #2056/1:	OUT	8		P 0	427858473	68106231
...						
After #116/3:	EXITD	BA #004640 NOL 8191 @ #3270/0		P 0	412965422	65759551
After #3267/3:	EDT	FW0 Q 0/#000472/	#000472	P 0	412956544	65757930
After #116/3:	EXITD	BA #004640 NOL 8191 @ #3153/4		P 0	412161464	65636232
After #60/5:	EDT	MT0 Q 8/#011657/	#011702	P 0	412161080	65636149
After #3153/4:	LOV	at #011656 (E2574)		P 0	412160886	65636114
After #116/3:	EXITD	BA #004640 NOL 8191 @ #2056/4		P 0	412160850	65636111
After #2472/0:	EDT	FW0 Q 0/#000472/	#000472	P 0	412148327	65633868
After #2056/1:	OUT	8		P 0	411606886	65632757
...						
After #116/3:	EXITD	BA #004640 NOL 8191 @ #2056/4		P 0	411352623	65591525
After #715/0:	EDT	MT0 Q 8/#004335/	#004336	P 0	411340926	65589427
After #2472/0:	EDT	FW0 Q 0/#000472/	#000472	P 0	411340010	65589283
After #712/5:	EDT	MT0 Q 8/#004335/	#004336	P 0	410541042	65587978
After #1027/5:	LOV	in #066440/#066635 for TR1		P 0	410532559	65586511
After #2472/0:	EDT	FW0 Q 0/#000472/	#000472	P 0	410531063	65586264
After #457/4:	EDT	MT0 Q 8/#004335/	#004336	P 0	410035534	65585510
After #713/5:	LOV	MT2 is busy		P 0	410026286	65583901
After #2472/0:	EDT	FW0 Q 0/#000472/	#000472	P 0	410021744	65583150
After #1045/3:	EDT	MT0 Q 8/#003021/	#003023	P 0	409224283	65581808
After #2472/0:	EDT	FW0 Q 0/#000472/	#000472	P 0	409212381	65579669
After #2056/1:	OUT	8		P 0	409007105	65578559
After #116/3:	EXITD	BA #004640 NOL 8191 @ #3157/0		P 0	408914735	65562404
After #2472/0:	EDT	FW0 Q 0/#000472/	#000472	P 0	408905390	65560684
After #3156/3:	OUT	17		P 0	408408682	65559504
After #116/3:	EXITD	BA #004640 NOL 8191 @ #0/0		P 0	408407665	65559338
After #2565/0:	EDT	TR1 Q 2/#004650/	#011447	P 0	408332663	65546631

(etc)

After earlier interrupts, whose tracing is now lost.

APPENDIX 7F: THE INSTRUCTION-TYPE AND INSTRUCTION-LOCUS FREQUENCY HISTOGRAMS

Either, both or neither of these plots may be obtained, at option. The examples show two different programs.

Histogram of the opcodes of 3438705 executed instructions with frequency $\geq 0.20\%$

004: ×F	54250	1.58%	:
012: PERM	68668	2.00%	:
024: FLOAT	10291	0.30%	:
027: NEG	54524	1.59%	:
033: NOT	14503	0.42%	:
034: ×D	37146	1.08%	:
036: −	66571	1.94%	:
037: SIGN	33410	0.97%	:
041: ZERO	41389	1.20%	:
042: DUP	136649	3.97%	:
043: DUPD	33299	0.97%	:
045: FIX	14447	0.42%	:
050: CONT	37951	1.10%	:
051: REVD	48162	1.40%	:
052: ERASE	52997	1.54%	:
056: +	141379	4.11%	:
060: /	17040	0.50%	:
062: /F	12042	0.35%	:
065: REV	119555	3.48%	:
066: CAB	17098	0.50%	:
074: +F	45208	1.31%	:
075: −F	49204	1.43%	:
102: MkMqQ	20663	0.60%	:
140: M+Iq	43817	1.27%	:
141: M−Iq	69939	2.03%	:
151: MqTOQk	15744	0.46%	:
161: SHA	51219	1.49%	:
164: SHL	15755	0.46%	:
166: SHLD	7628	0.22%	:
170: =[R]{Q C I M}q	250329	7.28%	:
171: {Q C I M}q	188054	5.47%	:
172: =+{Q C I M}q	23666	0.69%	:
173: LINK	40057	1.16%	:
174: =LINK	64503	1.88%	:
206: JrNEZ	9476	0.28%	:
213: Jr	62570	1.82%	:
215: JSr	74752	2.17%	:
217: EXIT	99197	2.88%	:
221: JrEQ	7125	0.21%	:
222: JrLTZ	12845	0.37%	:
224: JrGTZ	12795	0.37%	:
226: JrEQZ	44682	1.30%	:
260: JrCqNZ	32457	0.94%	:
300: EeMq	587012	17.07%	:
301: =EeMq	214536	6.24%	:
303: =EeMqQ	112303	3.27%	:
304: SET	201893	5.87%	:

Executions accounted for in the profile: 76.5%

Histogram of the loci of 306451954 executed instructions with frequency $\geq 0.50\%$

#63:	21524493	7.0%	:
#64:	21524493	7.0%	:
#65:	14349662	4.6%	:
#66:	7174831	2.3%	:
#70:	2544372	0.8%	:
#71:	1980726	0.6%	:
#72:	2592966	0.8%	:
#73:	3791307	1.2%	:
#74:	2939526	0.9%	:
#75:	28793967	9.4%	:
#76:	7174831	2.3%	:
#77:	11635855	3.8%	:
#100:	12165488	3.9%	:
#101:	12165488	3.9%	:
#102:	3041372	0.9%	:
#110:	3274426	1.0%	:
#111:	2971022	0.9%	:
#112:	4667583	1.5%	:
#113:	3073182	1.0%	:
#114:	6069284	1.9%	:
#115:	6069284	1.9%	:
#124:	1630169	0.5%	:
#125:	2220330	0.7%	:
#126:	2331633	0.7%	:
#127:	1554422	0.5%	:
#167:	2697092	0.8%	:
#171:	1992671	0.6%	:
#172:	1977597	0.6%	:
#173:	1977597	0.6%	:
#174:	2022819	0.6%	:
#204:	2697092	0.8%	:
#206:	2692318	0.8%	:
#207:	6742730	2.2%	:
#210:	9439822	3.0%	:
#211:	13485460	4.4%	:
#212:	6573495	2.1%	:
#213:	4552883	1.4%	:
#214:	3034642	0.9%	:
#215:	3034642	0.9%	:
#216:	4552423	1.4%	:
#217:	3022530	0.9%	:
#234:	3009498	0.9%	:
#235:	3022530	0.9%	:
#236:	4539949	1.4%	:
#240:	3022168	0.9%	:
#241:	3035562	0.9%	:
#242:	3371365	1.1%	:
#243:	2022819	0.6%	:
#244:	2058873	0.6%	:

Executions accounted for in the profile: 92.0%

APPENDIX 8: ERROR MESSAGES FROM **ee9**

The following headings introduce classes of error message. Each message may include more detailed diagnostics.

LIV INTERRUPT

NOUV INTERRUPT

RESET INTERRUPT

These are invalid operations that were detected by the KDF9 hardware. A RESET interrupt always connotes a failure, as do NOUV and LIV interrupts in problem programs. NOUV cannot happen in Director.

IMPOSSIBLE I/O OPERATION

IMPOSSIBLE I/O OPERATION IN DIRECTOR

LIV INTERRUPT IN DIRECTOR

These failures are caused by trying to execute an instruction that cannot be correctly interpreted. They include I/O orders with invalid parameters, requesting infeasible operations (such as trying to backspace a magnetic tape when it is already fully rewound), and miscellaneous other errors that **ee9** detects and fails. Although LIV could be signalled when running Director, it did not interrupt. However, it is always treated by **ee9** as failing, because it indicates a machine state that does not permit emulation to continue successfully.

FAILURE IN OUT

These failures are caused by OUT orders with invalid parameters, or that request impossible operations.

FEATURE NOT YET IMPLEMENTED

Presently, only unimplemented OUTs and unimplemented device types give rise to this error.

THE KDF9 OPERATOR HAS MADE A MISTAKE

These failures are caused by errors in the configuration of the emulated KDF9 that result from mistakes made in the invocation of **ee9**; e.g. by running it in the wrong working directory, so that some required file is absent.

INVALID PAPER TAPE FILE SUPPLIED

These failures result from the user not having set up an object program or data file correctly in KDF9 paper tape code.

APOLOGIES FOR THIS DISMAL FAILURE

These are failures of internal consistency checks made by **ee9**. They should never happen, but are handled defensively, with a view to debugging **ee9** itself.

OTHERS

The following is detected during diagnostic output at the end of a run. Should never happen. 8-)

Failure in finalize_ee9 ...

APPENDIX 9: OPERATING A KDF9 WITH THE TIME SHARING DIRECTOR

The following material draws heavily on the EE document ‘A1020 – Time Sharing Director – Mark I’, by A. Doust and M.R. Wetherfield, the authors of the Time Sharing Director program.

Commands from the computer operators to Director are known as ‘TINTs’, short for **T**ypewriter **I**nterrupts. They are initiated by pressing the interrupt button on the console Flexowriter. Director responds by typing **TINT**; and waiting for a command to be input. Each such command is terminated by typing ‘.→’. With ee9, interrupt by typing CTRL-C and answer the prompt with RETURN. For ‘→’, shown below, type ‘|’.

In the following, *dd* indicates a two-digit device (buffer) number; *tt* indicates a device type (see Table A4.2); *p* indicates a 12-character program name (also known as the PRN, or Program Reference Number); *n* is an octal number; *l* is a magnetic tape label; and *s* is a timesharing slot name (one of ‘P’, ‘Q’, ‘R’ or ‘S’).

A *s*.→

Terminate program in slot *s*, or terminate its loading.

A *s*+.→

Terminate program in slot *s* and type a basic diagnostic, showing the top of the SJNS and the NEST..

B *s n*.→

Read the octal integer *n* (up to 8 digits) and store its value into the less significant half of E0 of *s*. This can be used to convey options to the program at run time, separately from any tape or card data.

C *dd*.→

Load the magnetic tape on unit *dd*. Director reads the tape label, notes it, and positions the tape at the block immediately following, ready for allocation to a program that requests it by name, using OUT 4 or OUT 10.

D *dd*.→

Unload the magnetic tape on unit *dd*. The tape is rewound and disengaged from the deck, ready for removal.

E *dd*A.→ and E *dd*B.→

Nominate the magnetic tape loaded on unit *dd* for the input of A/B programs.

F .→

Forget (used if interrupt key pressed in error).

G.→

Type the peripheral unit list.

H.→

Type a list of ‘wanted’ tape labels. These are tapes that programs have requested, but have not yet been loaded.

I *s*0.→

Even restart for program *s*, i.e. make it execute the jump in the first halfword of word 0. This, and the odd restart, are for use when a program fails in a manner that has been anticipated, and can recover from.

I *s*1.→

Odd restart for program *s*, i.e. make it execute the jump in the second halfword of word 0.

L *dd/tt*.→

Change the peripheral device unit *dd* to type *tt*. If *tt* = 0, the unit is deleted from the free list.

M *x{s?}y*.→

Output a store print in syllabic octal. *x* and *y* are both octal integers. *y* words are output, starting at address *x*: if *y* is omitted, 1 is understood. If {*s?*} is P, Q, R or S then the base address of *s* is added to *x*. Any other separator (e.g. ‘/’) gives the absolute (Director) address *x*.

R *s*.→

Resume *s* (after suspension).

S *s*.→

Suspend *s*.

T *n/w*.→

Load a B-program and give it priority level *n*, with a store limit of *w* words. Absence of *w* implies that the program may want the whole store. If the / is replaced by P the paper tape reader from which its call tape is read is pre-allocated to the program. A-programs are loaded on the initiative of Director, not the human operators. When the resources that are needed become available, Director says:

n tt P dd s

n APIU dd s

The operator must have presented the program’s call tape on the A-program input unit *dd*. Director reads it in and continues:

* *date time*

n s P p or n s M p

depending on whether the program itself is to be read from paper tape or magnetic tape.

U.→

For each program in the machine, type its slot letter (P, Q, R or S), 12 character identifier, priority, base address, and running and elapsed times so far.

V n/n' .→

Put program in priority level n into priority level n' , and vice versa.

Z dd .→

Relabel the magnetic tape on unit dd . Director issues the query **TN/ID;** to which the operator replies **N/l.→**, e.g.:

TN/ID;N/+DIRECTORBINWORK.→

Director confirms the change, e.g.:

10L12 /Iden<+DIRECTORBINWORK>,TSN -00-0552

WHEN THINGS GO WRONG

If something goes wrong while reading in a program, Director types:

n CRNP Fx

where: CRNP means 'Can't Read New Program' and x is:

- 0: No paper tape reader is available for reading the A-block; this applies to B-programs only.
- 1: Parity failure in the A-block.
- 2: the A-block not in correct form.
- 3: C-block sum check failure.
- 4: Parity failure in the B-block.
- 5: E2 and E3 of the B-block are not the same as the first two words of the A-block.
- 6: Program's store requirement too large.
- 7: Incorrect fillers for a C-block.
- 8: Parity failure in a C-block.
- 9: the unit specified by the A-block is not available for reading B-blocks and C-blocks.

Program input can be terminated by giving a bogus A-block, thus inducing CRNP **F2**. B-program input can be terminated by the operator: if no reader is available, **TINT A s .→** causes CRNP **F0**, but after the A and B blocks have been read, and the program is waiting for core store, it causes CRNP **F6**.

In the event of a problem program failing, Director types:

$n s$ FAILS indicator [*SJNS, N1, N2 and NEST depth, where needed*]

REACT;

This gives the operator the opportunity to abandon, or to restart the run, using **A s .→**, **I $s0$.→**, or **I $s1$.→**. The *indicator* is taken from the following list of two or three digit codes:

- S** A Spurious interrupt, indicating a hardware (emulation) failure
- 00L** Lock-in violation
- 00N** Nest or SJNS over- or under-flow
- 00T** Time limit exceeded
- 01** Incorrect OUT number (it is still in N1)
- 02** Failure in OUT 5
- 03** Failure in OUT 6 or OUT 7
- 04** OUT 4 or OUT 10 requests an unavailable tape
- 05** OUT 1 obeyed with less than 3 items in nest
- 06** OUT 1 obeyed with incorrect program name
- 07** OUT 4 specifies incorrect identifier
- 11** OUT 10 specifies incorrect identifier
- 71** Closing OUT 8 stream without having written anything to it
- 72** Wrong terminator to output block in OUT 8
- 73** Incorrect character in OUT 8 block for the Flexowriter
- 74** Invalid stream number in OUT 8
- 75** No magnetic tape available for OUT 8
- 76** Parity failure in OUT 8 output
- 77** Invalid addresses for an OUT 8 block

The following failures do not give the option of a restart, i.e. **REACT**; is not typed:

- 10 CRNP failure in OUT 1
- 20 Incorrect program name in OUT 2
- 21 Time limit for OUT 2 absent or incorrect
- 22 OUT 2 obeyed with incorrect store limit in E1

NORMAL TERMINATION

When a program finishes without detected failure, the following message is typed out:

n p / / / e / run time / n.e.t. / elapsed time

where: *run time* is the CPU time used; *n.e.t.* is the ‘notional elapsed time’ which gives an estimate of the elapsed time would have been experienced had there been no delays due to multiprogramming; *elapsed time* is the ‘wall clock’ time since the program started; and *e* is the ‘ending number’, which gives the reason for termination, namely:

- 0: OUT 0
- 2: OUT 2
- 4: TINT A
- 5: CRNP in OUT 1
- 6: any failure in OUT 2
- 7: result of any response to **REACT**;

In the case of successful overlay by OUT 1, Director types the priority, slot and name of the new program:

n s OUT 1

n M p

‘DESCRIPTORS’

Messages of the form:

tt S dd s

indicate that unit *dd* is of type *tt* and has current status *S*. For units other than magnetic tape units, when the unit is unallocated *S* is the letter U. When the unit is pre-allocated to a program or allocated to a program, *S* is P or A respectively. *s* indicates the program to which the unit is (pre-)allocated. When a unit is allocated to Director (e.g. for TINT M), *S* is the letter D.

For magnetic tape units, these messages are only typed when the unit becomes loaded (L) or unloaded (U), except when the whole unit list is typed by TINT G). They then also include the Tape Serial Number (TSN) of the tape, e.g.:

10L12/Iden<PRINTEND>, TSN -20-0403

10L13/Ident<ZERO>, TSN -20-0284

When a list is typed by TINT G, the identifiers of tapes on loaded units only are typed, and the TSN is omitted, e.g.

10L12 PRINTEND

10L13 ZERO

When a unit is allocated or deallocated by a program, the descriptor is typed in the relevant column for P, Q, R or S. The only exceptions to this are the pre-allocation messages preceding A-program input, and messages concerning magnetic tape units. These are always typed in the Director column.

WAITING FOR ...

Messages of the form:

s WAITING FOR TYPE tt

mean that program *s* (which must in this case be a B-program) has attempted to allocate to itself a unit of type *tt*, but none is available. The operators have to decide whether or not to terminate program *s* (or another B-program), in order to free a device of type *tt*, or let it wait for one to be freed voluntarily.

APPENDIX 10: MAGNETIC TAPES IN **ee9**

Magnetic tapes are represented by Ada direct access files, using the `Ada.Direct_IO` library package, to allow selective overwriting of blocks. This is necessitated by the MWIPE and MGAP operations of the 1081-type tape deck. Tapes must be initialized (labelled) before any attempt is made to access them by a KDF9 problem program. (See **RLT** in Appendix 11.)

The “write permit ring”, which was a metal band inserted into the tape reel concentrically with the hub, depressed a switch to enable the deck’s write heads. It is modelled by the file’s access permission. Making the file read-only simulates an absent ring, allowing operations that do not change the contents of the tape, and failing those that do.

Each tape block, and each length of erased tape, is represented by one or more “slices”, a slice being a record in the direct access file. A slice has two components: a string and per-slice metadata. The string contains the Case Normal Latin-1 transliteration of all or part of a KDF9 data block. In the case of an erasure slice the string reserves space for possible future over-writing by data, but its contents are of no significance.

An emulated tape block is represented by a number of consecutive slices such that the total length of their contained strings is sufficient to encompass the data block or tape erasure. The maximum string size has been set so that blocks of 256 words (as used by POST), and short blocks containing OUT 8 print-line images, both have a space efficiency of better than 95%. The maximum block size (not slice size) is 32KW, or 256K characters, that being the maximum size of the core store.

No attempt is made to model either interblock gaps or the erased tape extending from the Beginning of Tape Window (BTW) to the start of the first block of data. However interblock gaps are taken into account when estimating the elapsed time of magnetic tape transfers.

Ampex 7-track tape files are recorded in exactly the same way as KDF9-native tapes, with the addition of single-slice blocks representing the two kinds of tape mark. They read into core as #170000000000 in the single input word. Tape marks do not appear in KDF9-native tapes.

The metadata in each record is as follows.

- Byte 0—a code indicating whether the slice represents a data block (code ‘D’), a length of tape erased by the MWIPE operation (code ‘W’), a length erased by the MGAP operation (code ‘G’), an even parity tape mark (code ‘e’), or an odd parity tape mark (code ‘o’).
- Byte 1—a code made up as follows:
 - (a) if the block is ‘LBM’ marked; i.e. if it was written by a MLWQq or MLWEQq instruction instead of the normal MWQq and MWEQq instructions, and so responds positively to the MLBQq order: +64;
 - (b) if this is the last slice of a multi-slice block: +8; and
 - (c) if this is the first slice of a multi-slice block: +1.

Byte 1 therefore has the following possible values (decimal = octal = Latin-1):

```
00 = 000 = NUL  => no flags
01 = 001 = SOH  => first slice of block
08 = 010 = BS   => last slice of block
09 = 011 = HT   => only slice of block (first and last)
64 = 100 = @    => LBM flag
65 = 101 = A    => first slice of block with LBM flag
72 = 110 = H    => last slice of block with LBM flag
73 = 111 = I    => only (first and last) slice of block with LBM flag
```

- Byte 2—the length of the string in this slice.

This encoding makes the contents of a magnetic tape somewhat legible using the POSIX **od** command, but **mtp** usually offers an a more useful view.

THE **mtp** UTILITY

Convenient views of a tape file are obtainable using the **mtp** program. The command takes the form:

```
mtp {MT|ST}{01234567}[T{PQRS}{1357}{01234567}]
```

The nominated magnetic tape file is read and an interpretation of its content is written to the standard output. Any error messages are written to the standard error output.

- To despool a stream from an OUT 8 tape written by TSD, append the slot letter and stream number, e.g. **mtp MT0P30**:

Despooling OUT 8 stream 30 for slot P, written on 23/09/69 by WHETSTONEUPU

```
N=   0 J=   0 K=   0 X1=+1.00000000 X2=-1.00000000 X3=-1.00000000 X4=-1.00000000
N= 120 J= 140 K= 120 X1=-0.06834220 X2=-0.46263766 X3=-0.72971839 X4=-1.12397907
N= 140 J= 120 K= 120 X1=-0.05533645 X2=-0.44743656 X3=-0.71097339 X4=-1.10309806
N=3450 J=   1 K=   1 X1=+1.00000000 X2=-1.00000000 X3=-1.00000000 X4=-1.00000000
N=2100 J=   1 K=   2 X1=+6.00000000 X2=+6.00000000 X3=-0.71097339 X4=-1.10309806
N= 320 J=   1 K=   2 X1=+0.49040732 X2=+0.49040732 X3=+0.49039250 X4=+0.49039250
N=8990 J=   1 K=   2 X1=+1.00000000 X2=+1.00000000 X3=+0.99993750 X4=+0.99993750
N=6160 J=   1 K=   2 X1=+3.00000000 X2=+2.00000000 X3=+3.00000000 X4=-1.10309806
N=   0 J=   2 K=   3 X1=+1.00000000 X2=-1.00000000 X3=-1.00000000 X4=-1.00000000
N= 930 J=   2 K=   3 X1=+0.83466552 X2=+0.83466552 X3=+0.83466552 X4=+0.83466552
```

```
THAT TOOK 23.8 SECONDS, FOR 42.0 KWIPS
FINISHED ENDS 0
```

- If an OUT 8 stream is numbered in the 50-57 range it is taken to be written in Ferranti 5-hole paper tape code. In that case **mtp** attempts to emulate an original teleprinter transcription of the data. For example, this KDF9 code:

ABCDEFGHIJKLMNPOQRSTUVWXYZ}|\Ø

is transcribed thus in Figures Shift:

12*4 () 78±=-vlf , 0>» 3 | 56/×9+ .ncr

and thus in Letters Shift:

ABCDEFGHIJKLMNPOQRSTUVWXYZ. ?f

Note that the Ø is suppressed.

- If a magnetic tape contains only legible text in paper tape code, e.g. a Usercode object program as written out by the Kids Grove Algol compiler, append a T to the tape deck name, e.g. **mtp** MT3T, to get its transcription into plain Latin-1.
- Invoked with just the deck name, e.g. **mtp** MT0, it gives a slice-by-slice analysis the tape file. Each character in a data block is displayed using its Case Normal Latin-1 transliteration. Each slice is displayed with its record number in the direct access file, the KDF9 data is enclosed in « » quotes, an asterisk flags a block written with a 'last block' marker, and complete blocks, possibly consisting of several slices, are separated by blank lines. Here is the analysis of the above OUT8 tape:

TAPE LABEL TSN '-00-1478', IDENTIFIER '+KOUTPUT/0023001'

```

2      * «0001ØØP"P000001©»
3      * «0001ØA&"©ØØØØP© »
4      «0003ØØ&"WHETSTONEUPU   23/09/69®»
5      «0001ØØ&"   STR30®»
6      «0001ØØ&"©ØØØØP© »
7      «0023ØØ&"N?ØØØØØØ   ØØØØØ J?ØØØØØ   ØØØØØ K?ØØØØØ   ØØØØØ X1?ØØØØ+1.00000000000000 »
8      «X2?ØØØØ-1.00000000000000 X3?ØØØØ-1.00000000000000 X4?ØØØØ-1.00000000000000000000®»
9      «0023ØØ&"N?ØØØØØØ 1200000 J?ØØØØØØ 1400000 K?ØØØØØØ 1200000 X1?ØØØØ-0.06834220000000 »
10     «X2?ØØØØ-0.46263766000000 X3?ØØØØ-0.72971839000000 X4?ØØØØ-1.1239790700000000000000®»
11     «0023ØØ&"N?ØØØØØØ 1400000 J?ØØØØØØ 1200000 K?ØØØØØØ 1200000 X1?ØØØØ-0.05533645000000 »
12     «X2?ØØØØ-0.44743656000000 X3?ØØØØ-0.71097339000000 X4?ØØØØ-1.1030980600000000000000®»
13     «0023ØØ&"N?ØØØØØØ3450000 J?ØØØØØØ   1ØØØØ K?ØØØØØØ   1ØØØØ X1?ØØØØ+1.00000000000000 »
14     «X2?ØØØØ-1.00000000000000 X3?ØØØØ-1.00000000000000 X4?ØØØØ-1.00000000000000000000®»
15     «0023ØØ&"N?ØØØØØØ2100000 J?ØØØØØØ   1ØØØØ K?ØØØØØØ   2ØØØØ X1?ØØØØ+6.00000000000000 »
16     «X2?ØØØØ+6.00000000000000 X3?ØØØØ-0.71097339000000 X4?ØØØØ-1.1030980600000000000000®»
17     «0023ØØ&"N?ØØØØØØ 3200000 J?ØØØØØØ   1ØØØØ K?ØØØØØØ   2ØØØØ X1?ØØØØ+0.49040732000000 »
18     «X2?ØØØØ+0.49040732000000 X3?ØØØØ+0.49039250000000 X4?ØØØØ+0.4903925000000000000000®»
19     «0023ØØ&"N?ØØØØØØ8990000 J?ØØØØØØ   1ØØØØ K?ØØØØØØ   2ØØØØ X1?ØØØØ+1.00000000000000 »
20     «X2?ØØØØ+1.00000000000000 X3?ØØØØ+0.99993750000000 X4?ØØØØ+0.9999375000000000000000®»
21     «0023ØØ&"N?ØØØØØØ6160000 J?ØØØØØØ   1ØØØØ K?ØØØØØØ   2ØØØØ X1?ØØØØ+3.00000000000000 »
22     «X2?ØØØØ+2.00000000000000 X3?ØØØØ+3.00000000000000 X4?ØØØØ-1.1030980600000000000000®»
23     «0023ØØ&"N?ØØØØØØ   ØØØØØ J?ØØØØØØ   2ØØØØ K?ØØØØØØ   3ØØØØ X1?ØØØØ+1.00000000000000 »
24     «X2?ØØØØ-1.00000000000000 X3?ØØØØ-1.00000000000000 X4?ØØØØ-1.00000000000000000000®»
25     «0023ØØ&"N?ØØØØØØ 9300000 J?ØØØØØØ   2ØØØØ K?ØØØØØØ   3ØØØØ X1?ØØØØ+0.83466552000000 »
26     «X2?ØØØØ+0.83466552000000 X3?ØØØØ+0.83466552000000 X4?ØØØØ+0.8346655200000000000000®»
27     «0001ØØ&"ØØØØØØØ®»
28     «0007ØØ&"THAT TOOK ØØØØØØ 23.8 ØØSECONDS~ FOR ØØØ 42.0 ØØKWIPSØØ®»
29     «0002ØØ&"FINISHED ENDS Ø®»
30     * «0001ØZ&"©ØØØØP© »
31     * «}ØØØØØØØ»
32     WIPE of 1906 erased characters
EOF
```

The final 'EOF' is not from the tape; it is a confirmation by **mtp** that all of the data on the tape has been shown.

APPENDIX 11: ANCILLARY PROGRAMS FOR USE WITH **ee9**

This section gives specifications for the utility programs—some shell command files, some in KDF9 Usercode, and others in Ada 2012—that are provided with distributions of **ee9** to facilitate its use.

THE **a2b** UTILITY

a2b reads data from standard input in a stated code and writes it to standard output in another form. Conversions are available between raw bytes and paper tape code, between paper tape code and Latin-1, and from paper tape code to a legible display. The command takes the form:

```
a2b { -reg | -r2p | -L2p | -L2A | -p2c [ F | D | M | P ] | -p2L | -p2t | -p2o | -p2r | -F2L }
```

Exactly one flag parameter must be given; it is actually case-insensitive. The effect is as follows:

-r2p	Take the input to be ‘raw’ bytes, perhaps representing KDF9 syllables; read groups of 6 to form 48-bit words; output each such word as 8 KDF9 characters in paper tape code .
-L2p	Take the input to be 8-bit characters in Latin-1 code; transcribe them individually to KDF9 characters in paper tape code .
-L2A, -reg	Take the input to be 8-bit characters in Latin-1 code; transcribe them to characters in 7-bit ASCII code that are equivalent for ka13 , namely: % for ÷, X for ×, ~ for ∅ and # for ±. If the flag is -reg , in addition delete any CR characters.
-p2L	Take the input to be KDF9 characters in paper tape code ; transcribe them individually to 8-bit characters in Latin-1 code; ignore NULs (runout); report parity errors.
-p2c	Take the input to be a KDF9 object program in paper tape code and output a valid call tape for it, using the name extracted from words 2 and 3 of the program’s A block; report parity errors. If one of the parameters D, M, or P is given a call type of that kind is created. The parameter F (for file) makes a call tape for the medium specified in the A block of the program file, and this is the default.
-p2t	Like -p2L , but effect case shifting and ignore redundant or non-printing characters, so as to give a clean output text ; report parity errors.
-p2o	Take the input to be KDF9 characters in paper tape code ; read groups of 8 to form 48-bit words; output each such word in octal half/word, Q-store, syllable and character formats; report parity errors and incomplete words.
-p2r	Take the input to be KDF9 characters in paper tape code ; read groups of 8 to form 48-bit words; output each such word as 6 ‘raw’ bytes; report parity errors and incomplete words.
-F2L	Take the input to be characters in Ferranti paper tape code ; transcribe them to 8-bit characters in Latin-1 code.

THE **RLT** UTILITY

RLT is a Usercode program that reads a set of tape labels from TR1 and writes them in succession to the magnetic tape files MT0-MT5 and ST0, having erased a gap in which to write the new label on each tape. The data given in TR1 consists of one or more lines of text. The first 8 characters of the line give the ‘Tape Serial Number’ (TSN). The next 8 or 16 characters give the mnemonic ‘identifier’ of the tape. If the tape is intended to be a scratch work tape, its identifier must consist of 8 blank spaces (all zeros in KDF9 character code). A 16-character identifier must start with the character ‘+’. There must be an EM character (‘|’ in **ee9**) after the last character of the identifier. A line starting with an EM character terminates the data.

Here is an example, showing the sort of data that is used to label the tapes by **ee9**’s regression test suite:

```
-00-1234WHETLIST |
00000000          |
-00-0552EFPBEAAG |
77777777MS-DUMP.  |
-00-2339+KOUTPUT/0023004 |
0-00-929PRINTEND  |
0-00-777AMPEXTM4  |
|
```

For ease of running, the command file **rlt** is provided to set up the environment and the **ee9** mode parameters for **RLT**. **rlt** truncates the former contents of all of the tape files and then runs **RLT** in privileged program mode, so that it can access the physical tape buffers without having had them allocated by Director (which would require the tapes to be already labelled.)

rlt takes the new labels from the file **RLT_data.txt** by default. For a different set of labels, supply the name of an alternative file as a parameter.

ports [-b | -k | -f *name*]

ports ('stop' backwards) converts from one KDF9 Algol format to another. With a first parameter of -f and a *name* parameter, it converts from either (or both) of the **whet** and **kalgol** stropped formats (see Appendix 12) to the Flexowriter format needed by **walgol**. With a first parameter of -b or -k it converts in the opposite sense: to 'bangs' format, with ! stropping, for -b; or to 'kwote' format, with ' ' stropping, for -k. When going from Flexowriter format to a stropped format it transfers verbatim any Flexowriter symbols that are accepted by the Kidsgrove compiler. See the table below. **ports** reads from its standard input and writes the result to its standard output.

bangs format	kwote format	Flexowriter format
!library	'library'	<u>c o m m e n t _ l i b r a r y</u>
!ALGOL	'ALGOL'	<u>A L G O L</u>
!and	'and'	<u>a n d</u>
!array	'array'	<u>a r r a y</u>
!begin	'begin'	<u>b e g i n</u>
!boolean	'boolean'	<u>b o o l e a n</u>
!Boolean	'Boolean'	<u>B o o l e a n</u>
!comment	'comment'	<u>c o m m e n t</u>
!do	'do'	<u>d o</u>
!else	'else'	<u>e l s e</u>
!end	'end'	<u>e n d</u>
!eqv	'eqv'	<u>e q v</u>
!EXIT	'EXIT'	<u>E X I T</u>
!false	'false'	<u>f a l s e</u>
!for	'for'	<u>f o r</u>
!goto	'goto'	<u>g o t o</u>
!to	'to'	<u>t o</u>
!go	'go'	<u>g o</u>
!if	'if'	<u>i f</u>
!imp	'imp'	<u>i m p</u>
!integer	'integer'	<u>i n t e g e r</u>
!KDF9	'KDF9'	<u>K D F 9</u>
!label	'label'	<u>l a b e l</u>
!not	'not'	<u>n o t</u>
!or	'or'	<u>o r</u>
!own	'own'	<u>o w n</u>
!procedure	'procedure'	<u>p r o c e d u r e</u>
!real	'real'	<u>r e a l</u>
!step	'step'	<u>s t e p</u>
!string	'string'	<u>s t r i n g</u>
!switch	'switch'	<u>s w i t c h</u>
!then	'then'	<u>t h e n</u>
!true	'true'	<u>t r u e</u>
!until	'until'	<u>u n t i l</u>
!value	'value'	<u>v a l u e</u>
!while	'while'	<u>w h i l e</u>
!ge or >= or _>	'ge' or >= or _>	<u>> back to >=</u>
!le or <= or _<	'le' or <= or _<	<u>< back to <=</u>
{ or _[{ or _[<u>_ back to _[</u>
} or _]	} or _]	<u>_ back to _]</u>
!ne or # or ±	'ne' or # or ±	<u>± back to ±</u>
!up or ^	'up' or ^	<u>^ back to ^</u>
!div or % or ÷	'div' or % or ÷	<u>÷ back to ÷</u>
\$	\$	<u>* back to \$</u>
* or ×	* or ×	<u>× back to ×</u>
~ or �	~ or �	<u>� back to �</u>

In this table '*x back to y*' indicates that *x* is the result of transcribing the stropped-format symbol to the Flexowriter format and that if *x* is found in Flexowriter-format input, it is transcribed as *y* to the stropped output. An underscore *followed by a blank* is accepted as an alternative to \$ in a stropped text.

bangs *prog*

kwote *prog*

bangs converts its standard input from Flexowriter format to the new Kidsgrove format in its standard output, using the ‘bang’ (!) stropping convention for basic symbols (e.g. `_e_n_d` becomes `!end`). **kwote** does the same thing, but using the quoting convention (e.g. `_e_n_d` becomes `'end'`). The result may need some amendment to pass compilation by **kalgol**. These shell commands merely invoke **ports** with appropriate parameters.

glance -o|-r

glance -o (a kind of ‘peep’) performs a range of peephole optimizations—like those of Healey and Wichmann—on the Usercode object programs generated by the Kidsgrove compiler. **glance -r** merely reformats the input, but makes no changes of substance. **glance** reads from its standard input and writes the result to its standard output.

kidopt

kidopt takes in a list of symbolic Kidsgrove Algol compilation options and converts them to the bit patterns used to transfer options to the compiler. See Appendix 12.

st_tl [*store* [*time*]]

st_tl amends, or creates if absent, the ST (store size) and TL (time limit) fields of a Usercode object program such as is generated by **kalgol**. If the *time* parameter is omitted, a value of 3000 (seconds) is used. If both are omitted, a store size of 7680 (words) is used. **st_tl** reads from its standard input and writes the result to its standard output.

The **PLT** utility (not yet available)

PLT is a Usercode program that generates POST-format program library tapes that are used by the Time Sharing Director to hold programs that can be loaded into store at the command of an M-type call tape. If given no data **PLT** claims a **ZERO** (scratch) work tape, relabels it +PROGRAM/LIBRARY, and writes the sentinel blocks that structure it as an empty library. If given a KDF9 object program in TR1, it appends that program to the library. The command file **plt** is provided to set up the environment and the **ee9** input mode parameters correctly for **PLT**. See the EE KDF9 System Routine Library Manual, §22.1, ‘Library Tape and Program Structure’, 1965.

ee9_regr_tests

ee9_regr_tests runs a comprehensive regression test of **ee9**. It does not include any tests of boot mode, because running a Director forces user interaction. This is considered a *regression* test because it compares its logged output with what is taken to be the correct output, and differences (if any) are reported.

cr_regr_tests

cr_regr_tests runs a comprehensive regression test of card reader and card punch emulation.

dr_regr_tests

dr_tests runs a set of regression tests of drum store emulation.

fd_regr_tests

fd_tests run a set of regression tests of fixed disc emulation.

kids_tests

kids_tests runs a selected set of tests of the Kidsgrove Algol compiler and runtime system.

liv_tests

liv_tests runs a comprehensive set of tests of basic CPU error-detection functionality.

mt_regr_tests

mt_regr_tests runs a comprehensive regression test of magnetic tape emulation.

nouv_tests

nouv_tests runs a comprehensive set of tests of basic CPU nesting store functionality.

si_tests

si_tests runs a set of tests for Standard Interface buffer emulation.

walgol_tests

walgol_tests runs a selected set of tests of the Whetstone Algol compiler and runtime system.

wich_tests

wich_tests runs a selected set of optimised tests of the Kidsgrove Algol compiler and runtime system.

to_9_from_1934

to_9_from_1934 reads the file `Data/wabbit_data_1900`, which must contain commands for the ICT 1934 plotter, converts them to KDF9 plotter commands, and writes them to the file `Data/wabbit_data_kdf9` in KDF9 8-track paper tape code. See ICT 1900 Technical publication 4003/C, ‘System Manual Volume X: Visual Input/Output Equipment’, 1966.

APPENDIX 12: WORKING WITH THE WHETSTONE AND KIDSGROVE ALGOL 60 COMPILERS

The Algol 60 Report defines a ‘reference language’ that ignores the practicalities of implementation on the graphically impoverished computers of the 1960s. Indeed it uses symbols that are not trivially available even today, such as \uparrow , \times , \div , \supset , and \equiv . It further specifies that an implementation may work in a ‘hardware language’ that represents the reference language using a computer’s native character set. The Whetstone (‘WAlgol’) and Kidsgrove (‘KAlgol’) hardware languages are based on the Friden Flexowriter character set listed in Appendix 2A. Symbols not present there are synthesized, primarily by means of the non-escaping underline character, so that **end**, for example, is prepared for input to a compiler as end and \wedge as and.

ee9 uses the Latin-1 break character as its transliteration of the Flexowriter underline, so that end appears thus: e_n_d.

Another scheme, used on computers with even less flexible input than the KDF9, is the ‘stropped’ format: each word symbol is decorated with flag characters and *ad hoc* arrangements are made for other symbols. Many Algol systems represent **end** as 'end', with the word enclosed in apostrophes. The Eldon 2 Teletype-based interactive system flags reserved words with an initial exclamation mark, thus: !end, and this is the form I find easiest to use.

The input routines for the resurrected Kidsgrove system allow great flexibility and convenience, with well-considered alternative lexical forms. The following table lists the transliterations used with **ee9** for both the Whetstone and Kidsgrove compilers. The original Whetstone representations are accepted by the Kidsgrove compiler, but not vice versa; the Whetstone representations are the only ones accepted by the 1964-vintage Whetstone compiler.

You may find it easiest to write a program in stropped format and convert it to original format, even if you only want to use the Whetstone compiler. See the **ports** commands specified in Appendix 11; it provides easy ways of converting programs from one format to the other. The **whet** command takes a program in Kidsgrove format, automatically invokes **ports** to convert it, and then uses **walgol** to compile and run the reformatted program.

Algol 60	Flexowriter	walgol	kalgol and whet
\uparrow	\uparrow	\wedge	\wedge !up 'up'
\times	\times	\times	\times * (for space in string see ‘_’)
\div	\div	\div	\div % !div 'div'
\supset	\supset	\supset	\supset ~
\neq	\neq	\pm	\pm # != !ne 'ne'
\geq	\geq	$_>$	$_>$ >= !ge 'ge'
\leq	\leq	$_<$	$_<$ <= !le 'le'
	\rightarrow		
end etc	<u>end</u> etc	<u>e_n_d</u> etc	!end 'end' etc
‘string’	[string]	<u>_[string_]</u>	{string} <u>_[string_]</u>
‘_’	[*l]	<u>_[*_]</u>	{_ } {\$} (note blank after _)
\neg	<u>not</u>	<u>_n_o_t</u>	!not 'not'
\wedge	<u>and</u>	<u>_a_n_d</u>	!and 'and'
\vee	<u>or</u>	<u>_o_r</u>	!or 'or'
\supset	<u>imp</u>	<u>_i_m_p</u>	!imp 'imp'
\equiv	<u>eqv</u>	<u>_e_q_v</u>	!eqv 'eqv'

WHETSTONE ALGOL (WALGOL)

A selection of Algol source programs can be found in the directory **Testing/Whetstone**. I provide a shell command, **walgol**, in the **Testing** directory, that makes it easy to compile and run Algol programs using WAlgol. It expects programs to be in text files named with the ‘.a60’ suffix. The suffix is optional in a name given as parameter to the **walgol** command; if supplied, it must be correct; if omitted, it is automatically restored; so that, e.g.:

```
walgol FLUID
```

compiles and runs **Whetstone/FLUID.a60**. There is actually a lot more flexibility in the location of WAlgol source programs; see Appendix 1. A source program and its stream 20 data may both be kept in the same file. Alternatively, only the program need be held in that file, with its data being provided by the interactive input feature of **ee9**; see §3.3.

Additional *mode* and *miscellany* parameters are optional and as given in §1 for the **nine** command. More comprehensive options, as described in §5, may be specified in the **settings_1.txt** file. If you want to specify different options for the program run, as opposed to its compilation, put them in the second options file, **settings_2.txt**.

If either TP0 or LP0 is found to be non-empty at the end of a run, its content is displayed on the terminal.

The **walgol** command sets up the FW0 file for the Whetstone system so that you do not need to type in any responses to its Flexowriter prompts. If the given replies are unsuitable, change the data in **FW0_for_Whetstone**.

A compilation listing is produced only if the response to the compiler’s **OUT**; prompt is an OUT 8 stream number: either 10. | or 30. |, rather than N. |. Stream 10 is directed to TP0; stream 30 goes to LP0. You may find that the listing disappoints: the source code is not included. Instead the compiler outputs a table relating source line numbers and object code addresses, which may later be quoted in runtime error messages. The line numbers are somewhat deceptive, as they seem to ignore blank lines and lines before the first **begin**.

When the WAlgol object program starts, it prompts **STREAM**; for the number of the output stream to be used; again, reply 10. | or 30. |. This should be the same as the stream number nominated for output in the source program. All of this is automated with the provided FW0_for_whetstone file, which suppresses the listing and sends output to stream 30.

You are likely to encounter both compilation errors and execution errors. They are explained in the documents *Whetstone Compilation Errors* and *Whetstone Runtime Errors*. Runtime errors may be accompanied by a ‘retrospective trace’ which lists the procedures called but not yet exited, the value of any failing operand, and a source code line number.

TO AVOID HAVING TO TYPE PROGRAMS IN THE CLUMSY FLEXOWRITER FORMAT, USE THE **whet** COMMAND INSTEAD. SEE §1.

KIDSGROVE ALGOL (KALGOL)

The KAlgol system has been resurrected by David Holdsworth, David Huxtable (one of its original authors) and Brian Wichmann. It is something of a chimaera. Several of its original components are lost and have been replaced with reverse engineered substitutes. The object program is generated in Usercode, which needs to be assembled. For these reasons running the Kidsgrove system is complex, and a little flaky. Nevertheless it is well worth the effort to see what a pioneering and very clever compiler was able to make of an intractable language like Algol 60 six decades ago.

A shell command, **kalgol**, in the **Testing** directory, hides this complexity and makes it easy to compile KAlgol programs. It has one parameter, the name of the program, taken from the directory **Testing/Kidsgrove**. Execution *mode* and *miscellany* parameters for the run of the compiler may be given in the environment variables KIDSMODE and KIDSMISC respectively. It is particularly useful to set KIDSMISC=w to suppress the compiler’s voluminous Flexowriter output, and this is the default. Store size and time limits may be specified by means of the ST and TL environment variables. (The defaults are 30000kW and 99999sec; note that these attributes affect only runs under a Director). For example:

```
ST=8192 TL=3600 kalgol FLUID
```

compiles Kidsgrove/FLUID.a60, allowing the object program 8kW of core and an hour of CPU time.

A shell command, **kids**, in the **Testing** directory, combines a call of **kalgol** to compile a program and a call of **nine** to run it. It has one mandatory parameter: the name of the program. Additional parameters are taken to be *data*, *mode* and *miscellany* options for the run of the object program by **nine**. Mode and miscellany parameters for the run of the compiler may be given as for **kalgol**. For example:

```
export KIDSMODE=t; export KIDSMISC=rsw
kids FLUID FLUID_data f
```

After a successful compilation with **kids** or **kalgol** the Usercode object program is left in the **Assembly** directory. The KDF9 executable is left in the **Binary** directory and can be run with the **nine** command, e.g.:

```
nine FLUID FLUID_data f
```

There is actually a lot more flexibility in the location of KAlgol source programs, object programs, their data files, and other components of the Kidsgrove compilation system provided with **ee9**; see Appendix 1.

Error messages are explained in the documents *Kidsgrove Compilation Errors* and *Kidsgrove Runtime Errors*.

KAlgol procedure bodies can be written in a version of Usercode that allows access to formal parameters. This is how the standard functions and the I/O libraries are implemented. See the document *Kidsgrove Code Procedures* and, for a simple example, the **PLOT** program in Kidsgrove.

The standard functions of Algol 60, and the KDF9 Algol I/O procedures, are incorporated into a program by means of the **library** (!library, 'library') directive, which is placed immediately after the first **begin** of a program and lists the (source code) libraries to be included. The document *Kidsgrove IO Libraries* names the libraries that were originally provided. The table below indicates the libraries (A0 through A15) that are available at the moment, and the procedures they contain. This is subject to occasional enhancements.

A0	the standard functions of Algol 60: sqrt , ln , exp , sin , cos , arctan , entier , abs , sign
A1	open and close ; all input must come from stream 20, and all output must go to stream 30
A4	read and read boolean (but read boolean is unlikely to work with A1 as yet)
A5	write , output , format , write boolean
A6	the union of A1, A4, A5, and A15
A12	newline , space , tab , gap
A13	character I/O and Algol Basic Symbol I/O
A14	copy text
A15	write text

It is an error to import duplicates, so that A6 should not be specified in combination with any of A1, A4, A5, and A15.

Kidsgrove compilation options, affecting the object code, e.g. by enabling optimisation, may be given on the command line, thus:

```
kalgol FLUID with OPTIMISER
```

or by setting an environment variable, thus:

```
export KIDSOPTS="OPTIMISER"
kalgol FLUID
```

(The “with ...” syntax is a nod to the **TRANSLATE** command of the KDF9’s POST software development system.)

To communicate compilation options to **kalgol** when using the **kids** command, which has an ample sufficiency of optional parameters, use KIDSOPTS, thus:

```
export KIDSOPTS="OPTIMISER"
kids FLUID FLUID_data
```

The presently effective compilation options are as follows:

NONE	#0	apply default options
OPTIMISER	#40	enable global optimization
option	code	meaning

- **NONE** specifies the compiler's default behaviour: an unoptimised, overflow-checking, assignment-tracing compilation with none of the other compile-time or run-time diagnostic outputs.

- **OPTIMISER** specifies the enabling of the compiler's global optimisations. The resulting program will probably run quite a bit faster. The optimiser may produce incorrect object code in difficult cases, but it is instructive to compare the Usercode produced with and without optimisation. Simple array-handling loops, in particular, benefit enormously. However, there is a restriction that the number of iterations of an optimised **for .. step .. until** loop must be less than 65536, because the iteration count is stored in the 16-bit C field of a Q store register. Optimised loops evaluate the **step** and **until** expressions only once. This makes a difference if the body of the loop has a potential side effect on either of them, or on the controlled variable: the side effect may be inoperative in optimised code, or may cause a malfunction.

Options are given to the compiler by means of the P, 'poke', flag in a settings file. The **kidopt** program takes in the list of symbolic options and outputs the corresponding P flag setting. **kalgol** copies **settings_for_kalgol.txt** into **settings_1.txt** and then appends the output from **kidopt**. This enables the user to set some persistent options for **ee9**'s runs of the compiler and allows them to be augmented with Kidsgrove-specific compilation options, per-run, by means of **kidopt**'s output. For more information about how **kalgol** works, see §7 of *ee9 Implementation Overview*.

(N.B. This correspondence between option bits and object code is that which is implemented in the version of the compiler that has been restored to working order and is now used with **ee9**. It is not consistent with English Electric documentation.)

An alternative to **kids** is **wich**, which post-processes the Usercode object program from an optimised compilation with **glance**, performing peephole transformations like those suggested by Healey and Wichmann in 1971.

The command **wich f** leaves the Usercode object program in a file called *f-OPT.k3*. Overflow checks and calls to diagnostic routines are removed from the program, and many longer, slower sequences are replaced by shorter, faster coding. The resulting program typically runs between 0% and 300% faster. However, the main reason for using **wich** is that it enables the user to see how a compiler could better exploit the KDF9 architecture. (For further steps in that direction, see the KDF9 Pascal cross-compiler, PASKAL.)

SEE ALSO

The Algol 60 language, and its usage on the KDF9 computer, are described in the following English Electric publications:

Report On The Algorithmic Language ALGOL 60

Ed. P. Naur; 1960. Available in: *Egdon System Reference Manual, Chapter 7, 'EGDON ALGOL'*.

Publication 103 550566; English Electric–Leo–Marconi Computers Limited; 1968.

KDF9 ALGOL Programming

Publication 1002 mm (R) 1000565, English Electric–Leo–Marconi Computers Limited; 1969.

The following papers describing the Kidsgrove and Whetstone systems were written by the original designers and implementers and constitute an invaluable historical resource:

KIDSGROVE

Implementation of ALGOL 60 for the English Electric KDF9

F. G. Duncan

Computer Journal, Vol.5, No.2, pp. 130–132; 1962.

Input and output for ALGOL 60 on KDF9

F. G. Duncan

Computer Journal, Vol.5, No.4, pp. 341–344; 1963.

A multi-pass translation scheme for ALGOL 60

E.N. Hawkins and D.H.R. Huxtable

Annual Review in Automatic Programming, Vol.3, pp. 163-205, Pergamon Press; 1963.

On writing an optimizing translator for ALGOL 60

D.H.R. Huxtable

Introduction to System Programming, APIC Studies in Data Processing, No.4, pp. 137-155, Academic Press; 1964.

Improving the Usercode Generated by the Kidsgrove Algol Compiler

R. Healey and B. A. Wichmann; National Physical Laboratory CCU TM3; 1971.

WHETSTONE

The Whetstone KDF9 ALGOL Translator

B. Randell

Introduction to System Programming, APIC Studies in Data Processing, No.4, pp. 122-136, Academic Press; 1964.

ALGOL 60 Implementation

B. Randell and L.J. Russell

Academic Press; 1964.

These documents are available online, at time of writing, at www.findlayw.plus.com/KDF9.

APPENDIX 13: USING USERCODE

The definitive guide to writing Usercode programs is the Manual, however some remarks about Usercode in the context of **ee9** may be helpful. Usercode programs can be conveniently compiled using the **ucc** command (see §1). **ucc** invokes **ka13**, a modern-day Usercode compiler written by David Holdsworth as part of the effort to resurrect the Whetstone Algol system. As a result of its history it does not implement every feature of the Usercode language as described in the Manual; only those features found to be necessary for its original purpose have been implemented. Here is a list of Usercode language features that are not supported by **ka13** and alternatives that can be used in each instance:

Usercode	ka13 alternative
device-specific I/O mnemonics (e.g. LPQ, TWQ, etc)	generic I/O mnemonics (e.g. POAQ)
floating point V store constants containing 0	scale the number explicitly
double-length V store constants	give the two words separately
V store printer constants using * for a space	use {S}

The input routines for **ka13** can use some of the same lexical forms as the resurrected Kidsgrove Algol compiler. The following table lists ASCII character transliterations that you may find convenient for use with **ka13**, particularly if your computer's editor does not readily support the Latin-1 forms:

Original KDF9 Usercode	Latin-1 for ee9	ASCII alternative
×	×	X
÷	÷	%
10	0	~
≠	±	#
→		
* (for a label starting a new word of code)	*	\$
string quotes ⌊ and ⌋	_ [and _]	{ and }

The ancillary program **a2b** has a facility to convert single characters in column 2 to those in column 3. (See Appendix 11.)

The **ee9** distribution contains a hierarchy of directories to locate its files systematically and conveniently for the casual user. The **Assembly** and **Binary** directories within **Testing** contain materials for KDF9 Usercode programs to work with. However, these locations are defaults only, and others can be substituted if that better meets the user's needs. To specify alternative places shell environment variables maybe used, as follows:

KDF9_USERCODE	Assembly
KDF9_BINARY	Binary
KDF9_DATA	Data
environment variable name	default path

The default path in each case is used if the environment variable is unset (either undefined or containing the empty string).

The defaults for unset variables conform to the distributed structure when the current working directory is **Testing**, where executable binaries and shell files are normally held. See the description of the **workplace** command in Appendix 1.

APPENDIX 14: PASKAL, THE KDF9 PASCAL CROSS-COMPILER

PASKAL runs natively on your computer and produces a KDF9 object program in the form of a Usercode text, with the aim of offering a better way to see the KDF9 architecture in action. If you are set up to compile and run ALGOL 60 or Usercode with **ee9** [see §1], you can easily compile and run Pascal programs thus:

```
pasc prog [ data | - ] [ mode | - ] [ miscellany | - ] [ tt [ tt ] ]
```

where the parameters are the same as those for the **kids** command [see Appendix 12], except that the *prog* parameter designates the source file `Pascal/prog.p`. For example:

```
pasc FLUID - t
```

compiles `Pascal/FLUID.p`, generating the Usercode program `Assembly/FLUID.k3` and compiling that into the KDF9 executable `Binary/FLUID`. It then runs the latter in tracing mode.

At the moment, the *data* parameter is ignored and a hyphen must be supplied if any later parameter is needed.

As for other KDF9 compilers, alternative source, Usercode and binary directories can be specified [see Appendix 1].

PASKAL generates very efficient code for many special cases, with the objective of reducing code size, a very important consideration with only 48K syllables of code available. This also shows the KDF9 architecture in a good light, which is the *raison d'être* of PASKAL. However, being a 1-pass compiler, it cannot perform global optimizations, but is helped in making non-local improvements by post-processing with **glance**, which has access to the whole of the Usercode object program.

A more complete description of the Pascal system is available in the *PASKAL: Users' Guide*, and a summary account of how the compiler works is given in the *PASKAL: Implementation Overview*.

SEE ALSO

PASKAL: Implementation Overview
BILL FINDLAY, 2022

PASKAL: Users' Guide
BILL FINDLAY, 2022

These documents are available online, at time of writing, at www.findlayw.plus.com/KDF9.