Symbols

General color codina

Best to Worst: TBD:

Symbols

- Recommendations differ depending on organization
- ☼ Parameter set
- Encryption / Ciphertext
- Signing / Signature
- Not yet standardized by NIST
- Implementation Code
- **6** Implementation size

Key Generation

- ▲ Decryption
- Verification
- E CPU Cycles
- Implementation complexity

Security categories of parameter sets











NIST Security Categories V, IV, III, II, I

Higher is more secure.

Implementation complexity and size









Low/Medium/High implementation complexity

Low/Medium/High implementation size

Lower is better.

Rating scales for parameter sizes and performance

Best to Worst: n < 2, n < 3, 4, $n \in \{5, 6\}$, $n \in \{7, 8\}$, n > 9



 $\mathcal{O}(5^n)$ CPU kilo cycles for key generation



 $\mathcal{O}(5^n)$ CPU kilo cycles for signing

 $\mathcal{O}(5^n)$ CPU kilo cycles for signature verification

 $\mathcal{O}(5^n)$ CPU kilo cycles for encryption / key encapsulation

 $\mathcal{O}(5^n)$ CPU kilo cycles for decryption / key decapsulation



 $\mathcal{O}(2^n)$ KB of signature size

 $\mathcal{O}(2^n)$ KB of ciphertext size

 $\mathcal{O}(2^{(n-5)})$ KB of signature algorithm public key size

 $\mathcal{O}(2^n)$ KB of encryption algorithm public key size

How To Interpret This Cheat Sheet

The goal of this cheat sheet is to make it as easy as possible to figure out which algorithm to pick for a given use case. Algorithm ID cards break down algorithm parameter sets, their important values and performance characteristics. The cheat sheet is intended to help users primarily in technical roles, such as engineers, architects or software developers working with post-guantum cryptography.

The focus is to avoid giving specific numbers measured in bits, bytes or cycles as this makes makes comparing numbers across algorithms difficult. Instead, this complexity is simplified by only providing a colorcoded number indicating the order of magnitude of each metric.

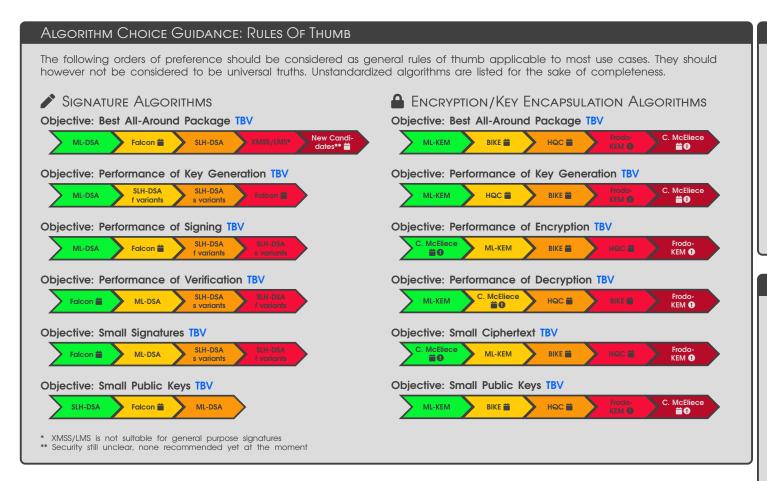
This approach prioritizes easy interpretation and comparability of metrics and in general quick informational gain over absolute precision of data – remember this is a cheat sheet, not a standard! This document is not intended to replace the study of algorithm specifications. It just aims to point you in the right direction quickly.

The approach of focusing on orders of magnitude walks a fine line between treating too many things as "equal" and not simplifying things enough to be easy to read and compare. "In the same order of magnitude" usually refers to "equal up to a factor of at most 10", which is a very coarse way of comparing numbers. Treating metrics that differ by a factor of e.g. 9.9 as "equal" because 9.9 < 10 paints a distorted picture. In cryptography, factors of 5 or even 2 can make a significant difference in security or performance, both in theory and in practice. In order to still tease out the differences in metrics without throwing too many things together that actually differ significantly, this cheat sheet applies different scaling and "orders of magnitude" (i.e., not regarding base 10) for different metrics.

It turns out that for metrics measured in (kilo) CPU cycle counts, i.e. algorithm performance, "up to a factor of 5" is a scale that is granular enough to work out the differences between algorithms while maintaining easy comparability. Those cycle counts heavily depend on the CPU used during measurement, hence the numbers need to be taken with a grain of salt, even if given exactly and not in terms of orders of magnitude.

For signature and ciphertext sizes as well as key sizes, measuring numbers in kilobytes "up to a factor of 2" is well suited to work out the differences between algorithms while allowing for quick comparison. Specifically for signature public key sizes only, we offset the corresponding color coding by 5 orders of magnitude. This is because SLH-DSA has extremely small pubic keys compared to all other signature algorithms, which would extend the scale into negative numbers (e.g., for SLH-DSA-SHA2-128s, the public key has $32=2^5$ bytes, which corresponds to an order of magnitude of -5 when measuring in orders of magnitude of base 2 and in kilobytes). This phenomenon of algorithm metrics spanning a very large range of orders of (base 2) magnitudes does not occur to this extent for encryption algorithms, making an offset unnecessary.

All values thus have a lower bound of 0. We do not limit the upper end of scales, but don't distinguish values areater than 10 anymore in terms of color codina. Please refer to the definitions on the left for symbol explanations, color coding and interpretation of numeric values.



PURE PQC VS PQ/T HYBRID

© Mario Schiener 🛅 😱

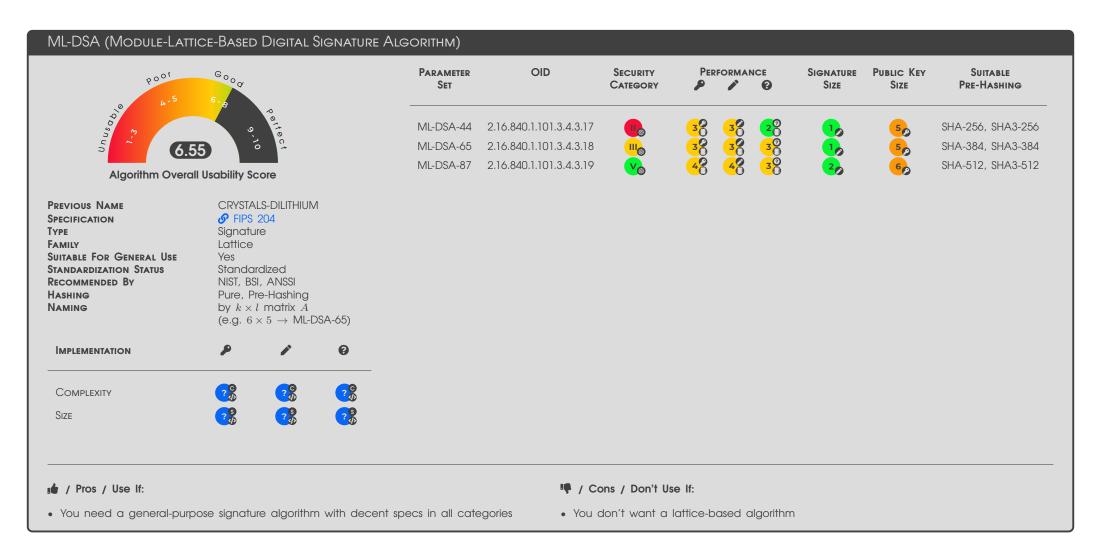
This topic depends on too many factors (e.g. cost of migration, security considerations, risk profile, GRC requirements) to give general advice. Those aspects will differ greatly between different organizations. The main reasons to adopt PQ/T hybrids are to still have traditional algorithms in place in case a new algorithm turns out to be insecure, and that PQ/T might help to avoid a big bang migration as systems could simply ignore the PQC component if they do not support it yet. Consider recommendations from different government agencies: BSI and ANSSI recommend PQ/T hybrid strategies, whereas NIST is more reserved towards PQ/T. If using PQ/T hybrids, preferably use ECC (e.g. secp256r1, brainpoolP256r1, Curve25519) instead of RSA for the traditional component to keep key sizes as small as possible.

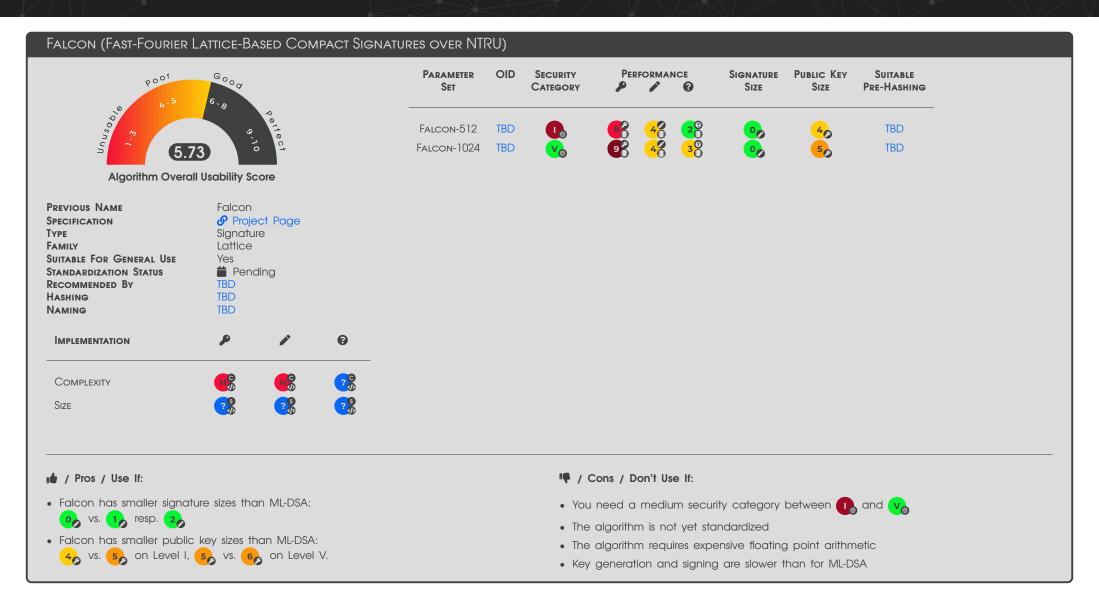
SECURITY CATEGORY CHOICES

- First, consider using $\overline{\mathbf{u}}$ as a baseline.
- Use IV or V for more security if possible (i.e., if a decrease in performance is not a concern and if no constraints apply).
- Use or is and only if or higher is not an option due to constraints (e.g. performance, memory, etc.).
- Comparing post-quantum security categories of algorithms to their traditional security level in bits is a complex subject and like comparing apples and oranges to some degree.

Pure vs. Pre-Hashing

- First, consider using pure (i.e., without pre-hashing) as this is the general recommendation.
- Pre-Hashing may be considered if one or more of the following applies:
 - The message *M* is too large to be sent to cryptographic module (CM) for hashing without significantly impacting performance. This may be the case e.g. in CMS related use cases such as S/MIME or code signing, or in cases of very narrow communication channels to the CM (e.g. between APDUs exchanged between smartcard and smartcard reader).
 - The hash needs to be signed with different algorithms and would be computed repeatedly without pre-hashing.
 - The specific hash function is not supported in a CM.





SLH-DSA (Stateless Hash-Based Digital Signature Standard) Goog POOT **PARAMETER** OID SECURITY PERFORMANCE SIGNATURE PUBLIC KEY SUITABLE SET CATEGORY SIZE SIZE PRE-HASHING 2.16.840.1.101.3.4.3.20 SHA-256, SHA3-256 SLH-DSA-SHA2-128s SLH-DSA-SHA2-128F 2.16.840.1.101.3.4.3.21 SHA-256, SHA3-256 2.16.840.1.101.3.4.3.22 SLH-DSA-SHA2-192s SHA-384, SHA3-384 Algorithm Overall Usability Score SLH-DSA-SHA2-192F 2.16.840.1.101.3.4.3.23 SHA-384, SHA3-384 SLH-DSA-SHA2-256s 2.16.840.1.101.3.4.3.24 SHA-512, SHA3-512 PREVIOUS NAME SPHINCS+ **SPECIFICATION 9** FIPS 205 SLH-DSA-SHA2-256F 2.16.840.1.101.3.4.3.25 SHA-512, SHA3-512 TYPE Signature SLH-DSA-SHAKE-128s 2.16.840.1.101.3.4.3.26 SHA-256, SHA3-256 Hash (stateless) FAMILY SUITABLE FOR GENERAL USE Yes SLH-DSA-SHAKE-128F 2.16.840.1.101.3.4.3.27 SHA-256, SHA3-256 STANDARDIZATION STATUS Standardized SLH-DSA-SHAKE-192s 2.16.840.1.101.3.4.3.28 SHA-384, SHA3-384 RECOMMENDED BY NIST, BSI, ANSSI Pure, Pre-Hashina HASHING SLH-DSA-SHAKE-192F 2.16.840.1.101.3.4.3.29 SHA-384, SHA3-384 NAMING based on various characteristics SLH-DSA-SHAKE-256s 2.16.840.1.101.3.4.3.30 SHA-512, SHA3-512 (*s=small signatures, *f=fast) SLH-DSA-SHAKE-256F 2.16.840.1.101.3.4.3.31 SHA-512, SHA3-512 **IMPLEMENTATION** 0 COMPLEXITY SIZE / Pros / Use If: / Cons / Don't Use If: Alternative to ML-DSA and Falcon that is not based on lattices Poor key generation and signing performance compared to other algorithms Very small public keys • High complexity of the algorithm and the implementation • Possible interoperability issues due to the many variants that may not all be supported everywhere

XMSS / XMSS-MT (eXtended Merkle Signature Scheme / extended Merkle Signature Scheme Multi Tree)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
HASHING
NAMING

XMSS/XMSSMT
𝚱 SP 800-208, 𝚱 RFC 8391
Signature
Merkle Trees (stateful hash trees)
No
Standardized
NIST, BSI, ANSSI
TBD
XMSS-[Hashfamily]_[h]_[n]
XMSSMT-[Hashfamily]_[h]/[d]_[n]
where h is the tree height,
d is the number of layers, and
n is the message length in bits

	n is ine	message i	engin in bi
N	P	•	•
	? ©	? ©	? ©
	? S	? ∯	? ®

Parameter Set	NUMERIC IDENTIFIER	SECURITY CATEGORY	PER	FORMAN	(CE	SIGNATURE SIZE	MAXIMUM SIGNATURES	NUMBER OF LAYERS
XMSS-SHA2_10_256	0x00000001	V®	??	· **	?0	?	210	1
XMSS-SHA2_16_256	0x00000002	V	?	? <u>?</u>	? <u>@</u>	?	216	1
XMSS-SHA2_20_256	0x00000003	Vo	?8	? 8	?0	?	2 ²⁰	1
XMSSMT-SHA2_20/2_256	0x00000001	Vo	??	?8	?0	?	2 ²⁰	2
XMSSMT-SHA2_20/4_256	0x00000002	V	₹	?	70	?	2 ²⁰	4
XMSSMT-SHA2_40/2_256	0x00000003	V	₹ <u>₹</u>	78	70	?	2 ⁴⁰	2
XMSSMT-SHA2_40/4_256	0x00000004	V	<u>₹</u>	?	? <u>@</u>	?	2 ⁴⁰	4
XMSSMT-SHA2_40/8_256	0x00000005	V	? <u>?</u>	? 8	₹ <u>0</u>	?	2 ⁴⁰	8
XMSSMT-SHA2_60/3_256	0x00000006	V	? <u>?</u>	? <u>?</u>	70	?	2 ⁶⁰	3
XMSSMT-SHA2_60/6_256	0x00000007	V	₹ <u>₹</u>	₹	₹ <u>0</u>	?	260	6
XMSSMT-SHA2_60/12_256	0x00000008	Vo	?8	₹	?0	?	2 ⁶⁰	12

Note

₱ SP 800-208 defines further parameter sets not listed in ₱ RFC 8391 using other hash functions (SHA256/192, SHAKE256/256, SHAKE256/192). Furthermore, ₱ RFC 8391 lists optional parameter sets that are not approved in ₱ SP 800-208. All of those variants are omitted here as they are not likely to be widely used, in particular not after ML-DSA and SLH-DSA have been standardized in the meantime.

/ Pros / Use If:

IMPLEMENTATION

COMPLEXITY

Size

- You can predict the maximum number of signatures that are going to be required
- Firmware signing use cases
- You want a signature scheme where the security only relies on the security of the hash function used without assuming the hardness of another mathematical problem.
- Cf. § SP 800-208, Section 1.1 for additional explanations

- You require an algorithm for general use
- You cannot predict the maximum number of signatures that are going to be required, or the number of required signatures exceeds the maximum number of signatures enabled through the approved parameter sets
- Your application does not allow for the careful state management and tracking of signatures performed that is required with this algorithm

LMS/HSS (Leighton-Micali Signature / Hierarchical Signature System)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
HASHING
NAMING

LMS, LMS/HSS

SP 800-208, RFC 8554
Signature
Merkle Trees (stateful hash trees)
No
Standardized
NIST, BSI, ANSSI
TBD
LMS_SHA256_M32_[h]
where h is the tree height

IMPLEMENTATION	P	•	•
COMPLEXITY	? ©	? ©	? €
Size	?#	₹	?

PARAMETER SET	NUMERIC IDENTIFIER	SECURITY CATEGORY	PER	FORMAN	(CE	SIGNATURE SIZE	MAXIMUM SIGNATURES
LMS_SHA256_M32_H5 LMS_SHA256_M32_H10	0x00000005 0x00000006	?	?0	? 6	?0 ?0	?0	2 ⁵ 2 ¹⁰
LMS_SHA256_M32_H15 LMS_SHA256_M32_H20 LMS_SHA256_M32_H25	0x00000007 0x00000008 0x00000009	? ? •	?0	?0	?0 ?0 ?0	? ₀	2 ¹⁵ 2 ²⁰ 2 ²⁵

NOTE:

♦ SP 800-208 defines further parameter sets not listed in ♠ RFC 8554 using other hash functions (SHA256/192, SHAKE256/256, SHAKE256/192). Those variants are omitted here as they are not likely to be widely used, in particular not after ML-DSA and SLH-DSA have been standardized in the meantime. Similar to XMSS multi-tree variants, a comparable structure exists with Hierarchical Signature System (HSS) variants of LMS that involves a number of LMS instances.

/ Pros / Use If:

- You can predict the maximum number of signatures that are going to be required
- Firmware signing use cases
- You want a signature scheme where the security only relies on the security of the hash function used without assuming the hardness of another mathematical problem.
- Cf. SP 800-208, Section 1.1 for additional explanations

- You require an algorithm for general use
- You cannot predict the maximum number of signatures that are going to be required, or the number of required signatures exceeds the maximum number of signatures enabled through the approved parameter sets
- Your application does not allow for the careful state management and tracking of signatures performed that is required with this algorithm

♠ ENCRYPTION / KEY ENCAPSULATION ALGORITHM ID CARDS



BIKE (BIT FLIPPING KEY ENCAPSULATION)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

BIKE

Project Page
Encryption/KEM
Codes
Yes
4th round candidate
TBD
after security categories I, III and V

IMPLEMENTATION	•		8
Complexity	? ©	? ©	?g
Size	? S	? S	? S

PARAMETER SET	OID	SECURITY CATEGORY	PER	FORMAN	ICE	CIPHERTEXT SIZE	PUBLIC KEY SIZE
BIKE-L1	TBD		30	20	4 <mark>0</mark>	<u></u>	
BIKE-L3	TBD	III _®	48	3 0	5 <u>Ö</u>	10	1
BIKE-L5	TBD	V®	n/g	n/g	n/g	2 ₀	20

NOTE

No data available for BIKE-L5 for cycle counts. The algorithm overall usability score is computed over BIKE-L1 and BIKE-L3 only.

/ Pros / Use If:

• Possible alternative to ML-KEM not based on lattices.

- Not yet standardized. Note that NIST intends to standardize at most one of the algorithms HQC and BIKE.
- Metrics not as good as the ones of ML-KEM.

HQC (HAMMING-QUASI-CYCLIC) PARAMETER SECURITY **PERFORMANCE** CIPHERTEXT PUBLIC KEY SET CATEGORY SIZE SIZE HQC-128 HQC-192 **TBD** HQC-256 **TBD** Algorithm Overall Usability Score PREVIOUS NAME HQC Project Page SPECIFICATION TYPE Encryption/KEM **FAMILY** Codes SUITABLE FOR GENERAL USE 4th round candidate STANDARDIZATION STATUS RECOMMENDED BY **TBD** NAMING 128, 192 and 256 bits of security in reference to security categories I, III and V **IMPLEMENTATION** 0 COMPLEXITY / Pros / Use If: / Cons / Don't Use If: Possible alternative to ML-KEM not based on lattices. • Not yet standardized. Note that NIST intends to standardize at most one of the algorithms HQC and BIKE. • Metrics not as good as the ones of ML-KEM.

FrodokeM Good **PARAMETER** SECURITY PERFORMANCE CIPHERTEXT PUBLIC KEY POOT SET SIZE SIZE CATEGORY Unusab/o FRODOKEM-640-AES **TBD** FRODOKEM-640-SHAKE TBD FRODOKEM-976-AES TBD Algorithm Overall Usability Score FRODOKEM-976-SHAKE TBD FRODOKEM-1344-AES TBD PREVIOUS NAME **FrodoKEM** SPECIFICATION Project Page FRODOKEM-1344-SHAKE Encryption/KEM TYPE Lattice **FAMILY** SUITABLE FOR GENERAL USE Yes Disregarded by NIST STANDARDIZATION STATUS Note: RECOMMENDED BY BSI, ANSSI There are also ephemeral versions of FrodoKEM (eFrodoKEM-640, etc.), but BSI only recommends FrodoKEM-976 Frodokem-[n]-[AES | SHAKE] NAMING and FrodoKEM-1344 versions for long-term protection. The algorithm score also only includes the non-ephemeral where n is a matrix dimension versions listed above. 0 **IMPLEMENTATION** COMPLEXITY / Pros / Use If: / Cons / Don't Use If:

• More conservative than ML-KEM since it is based on unstructured grids

• Not standardized by NIST because ML-KEM is more efficient.

3.31

Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

CLASSIC McELIECE

Classic McEliece $\ref{Project Page}$ Project Page Encryption/KEM Codes No / With Limitations 4th round candidate (NIST) BSI mceliece[n][deg(F(y))][[f]] where n, F(y) are parameters non-f versions: $(\mu, \nu) = (0, 0)$ f versions: $(\mu, \nu) = (32, 64)$

IMPLEMENTATION	P	ř	0
Complexity	? ©	? ©	? ©
Size	? 6	? 6	? (5)

PARAMETER SET	OID	SECURITY CATEGORY	PER	FORMAN	CE	CIPHERTEXT SIZE	PUBLIC KEY SIZE
MCELIECE348864	TBD		68	20	3 O	<u></u>	8,
MCELIECE348864F	TBD		60	n/eS	n/e	0 0	6
мсецесе460896	TBD	III	R	20	3 0	<u></u>	9,
MCELIECE460896F	TBD	III	R	n/eS	n/e	<u></u>	9,
MCELIECE6688128	TBD	V		3 <u>0</u>	3 O	0 0	10,
MCELIECE6688128F	TBD	Vo	R	n/e	n/e	<u></u>	10,
MCELIECE6960119	TBD	V	R	3 <u>0</u>	3 <u>0</u>	0	10,
MCELIECE6960119F	TBD	V	R	n/a	n/e	<u></u>	10,
MCELIECE8192128	TBD	Vo		30	3 0	<u></u>	10,
MCELIECE8192128F	TBD	V®	8	n/g	n/g	∞	10,

Note:

No data available for the f versions for encryption/decryption cycle counts. The algorithm overall usability score is computed over non-f versions only. Non-f and f versions are interoperable, with f versions offering faster key generation and non-f versions having simpler key generation.

/ Pros / Use If:

- Possible alternative to ML-KEM not based on lattices.
- Considered safe and recommended by BSI even though not standardized by NIST.
- Use cases that allow for pre-sharing of keys where the large public key size is not a concern.

- Not yet standardized by NIST.
- Not suitable for general use due to the very large public keys and poor performance of key generation.
- Due to aforementioned point, likely not suitable for embedded devices.

DITIONAL VS	s. Post-Quan	itum Security Level Compa	arison <mark>TBV</mark>						
TRADITIONAL PRE	E-QUANTUM SECURITY CLASSICAL ATTACK COMPLEXITY	ATTACK IN TERMS OF BIT SIZE N	ALGORITHM EXAMPLES	NIST SECURITY CATEGORIES DEFINED THROUGH ATTACK	Classical Attack	Quantum Attack Complexity	POST-QUANTUM SE SECURITY [BITS]	ecurity Signature Parameter Sets	Encryption/KEM Parameter Sets
n	$\mathcal{O}\left(2^{n}\right)$	$n ext{-BIT}$ KEY SPACE EXHAUSTIVE KEY SEARCH			Grover	$\mathcal{O}\left(2^{\frac{n}{2}}\right)$	$\frac{n}{2}$		
n	$\mathcal{O}\left(2^{n}\right)$	$n ext{-BIT}$ hash pre-image attack			Grover	$\mathcal{O}\left(2^{\frac{n}{2}}\right)$	$\frac{n}{2}$		
n	$\mathcal{O}\left(2^{n}\right)$	n-bit hash 2nd-pre-image attack			Grover	$\mathcal{O}\left(2^{\frac{n}{2}}\right)$	$\frac{n}{2}$		
$\frac{n}{2}$	$\mathcal{O}\left(2^{\frac{n}{2}}\right)$	$n ext{-BIT}$ hash collision attack			Brassard et al.	$\mathcal{O}\left(2^{\frac{n}{3}}\right)$	$\frac{n}{3}$		
		128-bit key space exhaustive key search	AES-128	Category 1	Grover	$\mathcal{O}\left(2^{64}\right)$	64		ML-KEM-512.
	(122)	128-bit hash pre-image attack	N/A		Grover	$\mathcal{O}\left(2^{64}\right)$	64	FALCON-512,	BIKE-L1, HQC-128,
128	$\mathcal{O}\left(2^{128}\right)$	128-bit hash 2nd-pre-image attack	N/A		GROVER	$\mathcal{O}\left(2^{64}\right)$	64	SLH-DSA 128 variants	FrodoKEM-640
		256-bit hash collision attack	SHA-256, SHA3-256	CATEGORY 2	Brassard et al.	$\mathcal{O}\left(2^{85}\right)$	85	MLA-DSA-44	
		192-bit key space exhaustive key search	AES-192	Category 3	Grover	$\mathcal{O}\left(2^{96}\right)$	96		ML-KEM-768.
	. (102)	192-bit hash pre-image attack	N/A		Grover	$\mathcal{O}\left(2^{96}\right)$	96	ML-DSA-67, SLH-DSA 192 variants	BIKE-L3, HQC-192,
192	$\mathcal{O}\left(2^{192}\right)$	192-bit hash 2nd-pre-image attack	N/A		GROVER	$\mathcal{O}\left(2^{96}\right)$	96	SLH-DSA 192 VARIANTS	FrodoKEM-976
		384-bit hash collision attack	SHA-384, SHA3-384	Category 4	Brassard et al.	$\mathcal{O}\left(2^{128}\right)$	128		
		256-bit key space exhaustive key search	AES-256	Category 5	Grover	$\mathcal{O}\left(2^{128}\right)$	128	ML-DSA-85,	ML-KEM-1024.
251	- (-256)	256-bit hash pre-image attack	N/A		Grover	$\mathcal{O}\left(2^{128}\right)$	128	SLH-DSA 256 VARIANTS	BIKE-L5, HQC-256,
256	$\mathcal{O}\left(2^{256}\right)$	256-bit hash 2nd-pre-image attack	N/A		Grover	$\mathcal{O}\left(2^{128}\right)$	128	XMSS**	FrodoKEM-1344
		512-bit hash collision attack	SHA-512, SHA3-512		Brassard et al.	$\mathcal{O}\left(2^{170}\right)$	170		
		384-bit key space exhaustive key search	N/A		Grover	$\mathcal{O}\left(2^{192}\right)$	192		
	(-384)	384-bit hash pre-image attack	SHA-384, SHA3-384		Grover	$\mathcal{O}\left(2^{192}\right)$	192		
384	$\mathcal{O}\left(2^{384}\right)$	384-bit hash 2nd-pre-image attack	SHA-384, SHA3-384		Grover	$\mathcal{O}\left(2^{192}\right)$	192		
		768-bit hash collision attack	N/A		Brassard et al.	$\mathcal{O}\left(2^{256}\right)$	256		
		512-bit key space exhaustive key search	N/A		Grover	$\mathcal{O}\left(2^{256}\right)$	256		
510	- (-512)	512-bit hash pre-image attack	SHA-512, SHA3-512		Grover	$\mathcal{O}\left(2^{256}\right)$	256		
512	$\mathcal{O}\left(2^{512}\right)$	512-bit hash 2nd-pre-image attack	SHA-512, SHA3-512		Grover	$\mathcal{O}\left(2^{256}\right)$	256		
		1024-bit hash collision attack	N/A		Brassard et al.	$\mathcal{O}\left(2^{341}\right)$	341		

TBD ADDITIONAL COMMENTS AND EXPLANATIONS TO READERS FOR HOW THIS IS TO BE INTERPRETED; EXPLAIN LIGHTER COLORS, XMSS, LMS?

ALGORITHM OVERALL USABILITY SCORE

We try to measure an algorithm's overall usability for a general use case without any special characteristics by calculating a single number between 0 (worst) and 10 (best) for the algorithm. This calculation is taking into account all parameter sets, performance metrics, public key size, signature/ciphertext size, the number of security categories provided, whether or not it is suitable for general use, and the complexity and size of its implementation. We define an algorithm's overall score as

$$\mathsf{score}_{\textit{algorithm}} = \mathsf{max}\left\{\mathsf{0}, \ \mathsf{avg}\left\{\mathsf{score}_{\textit{parameterSet}} \mid \textit{parameterSet} \ \text{is a parameter set of} \ \textit{algorithm}\right\} - \frac{1}{8} \cdot \left(\mathsf{5} - \mathsf{categories}_{\textit{algorithm}}\right) - \frac{1}{4} \cdot \mathsf{impl}_{\textit{algorithm}} - \mathsf{generality}_{\textit{algorithm}}\right\} \right\}$$

where score parameter. Set is a score for an individual algorithm parameter set. Depending on the type of algorithm, score parameter. Set is defined as

$$score_{signature-parameterSet} = 10 - avg \left\{ \begin{array}{ccc} \mathbf{n_0^2}, & \mathbf{n_0^2}, & \mathbf{n_0^2}, & \mathbf{n_0^2}, \end{array} \right\}$$







 $impl_{algorithm} = avg \left\{ \begin{array}{c} L_{o}^{o}, M_{o}^{o}, L_{o}^{o}, M_{o}^{o}, \end{array} \right\}$

score_{encryption-parameterSet} =
$$10 - \text{avg} \left\{ \begin{array}{c} \mathbf{n_0^2}, \quad \mathbf{n_0^2}, \quad \mathbf{n_0^2}, \quad \mathbf{n_0^2}, \end{array} \right.$$







Furthermore, $1 \le$ categories $algorithm \le 5$ denotes the number of different NIST security categories offered by algorithm. By assigning numeric values of 0 = 100, 1 = 100













Finally, we define

$$\text{generality}_{\textit{algorithm}} = \begin{cases} 0 & \text{if } \textit{algorithm} \text{ is a general purpose algorithm} \\ 2 & \text{else} \end{cases}$$

to take into account if the algorithm is suitable for general use.

TBD: In the algorithm scores given in ID cards, we currently use impl_{algorithm} = 0 in the respective computations because the necessary values for implementation complexity and size are still TBD. This will be corrected later.

Example: ML-DSA Overall Usability Score

We calculate









Similarly, we obtain $score_{ML-DSA-65} = 7.0$ and $score_{ML-DSA-67} = 6.2$. Furthermore, categories $m_{L-DSA} = 3$ since ML-DSA offers the three security categories $m_{L-DSA-65} = 7.0$ and $m_{L-DSA-65} = 0$ since ML-DSA is a general purpose





signature algorithm. This results in an overall usability score of 6.55:



$$score_{ML-DSA} = \max \left\{ 0, \text{ avg } \{ score_{ML-DSA-44}, \text{ score}_{ML-DSA-65}, \text{ score}_{ML-DSA-87} \} - \frac{1}{8} \cdot (5 - \text{categories}_{ML-DSA}) - \frac{1}{4} \cdot \text{impl}_{ML-DSA} - \text{generality}_{ML-DSA} \right\}$$

$$= \max \left\{ 0, \text{ avg } \{ 7.2, 7.0, 6.2 \} - \frac{1}{8} \cdot (5-3) - 0 - 0 \right\}$$

$$= \max \left\{ 0, 6.8 - 0.25 \right\}$$

$$= \max \left\{ 0, 6.55 \right\}$$

$$= 6.55$$

TBD: We use $impl_{MI-DSA} = 0$ because the necessary values to compute it are still TBD, cf. ML-DSA ID card. This will be corrected later.