



Symbols

General color codina

Best to Worst: Links and TBDs:

Symbols

Recommendations differ depending on organization

☼ Parameter set

Encryption / Ciphertext

Signing / Signature

Not yet standardized by NIST

Implementation Code

6 Implementation size

Key / Key Generation

▲ Decryption

Verification

CPU Cycles

Implementation complexity

Security categories of parameter sets











NIST Security Categories V, IV, III, II, I

Higher is more secure.

Implementation complexity and size









Low/Medium/High implementation complexity

Low/Medium/High implementation size

Lower is better.

Rating scales for parameter sizes and performance

Best to Worst: n < 2, n < 3, 4, $n \in \{5, 6\}$, $n \in \{7, 8\}$, n > 9



 $\mathcal{O}(5^n)$ CPU kilo cycles for key generation



 $\mathcal{O}(5^n)$ CPU kilo cycles for signing



 $\mathcal{O}(5^n)$ CPU kilo cycles for signature verification



 $\mathcal{O}(5^n)$ CPU kilo cycles for encryption / key encapsulation



 $\mathcal{O}(5^n)$ CPU kilo cycles for decryption / key decapsulation



 $\mathcal{O}(2^n)$ KB of signature size



 $\mathcal{O}(2^n)$ KB of ciphertext size



 $\mathcal{O}(2^{(n-5)})$ KB of signature algorithm public key size

 $\mathcal{O}(2^n)$ KB of encryption algorithm public key size

How To Interpret The Cheat Sheets

The goal of this series of cheat sheets is to make it as easy as possible to figure out which algorithm to pick for a given use case. Algorithm ID cards break down algorithm parameter sets, their important values and performance characteristics. The cheat sheets are intended to help users primarily in technical roles, such as engineers, architects or software developers working with post-quantum cryptography.

The focus is to avoid giving specific numbers measured in bits, bytes or cycles as this makes makes comparing numbers across algorithms difficult. Instead, this complexity is simplified by only providing a colorcoded number indicating the order of magnitude of each metric.

This approach prioritizes easy interpretation and comparability of metrics and in general quick informational gain over absolute precision of data – remember this is a cheat sheet, not a standard! This document is not intended to replace the study of algorithm specifications. It just aims to point you in the right direction quickly.

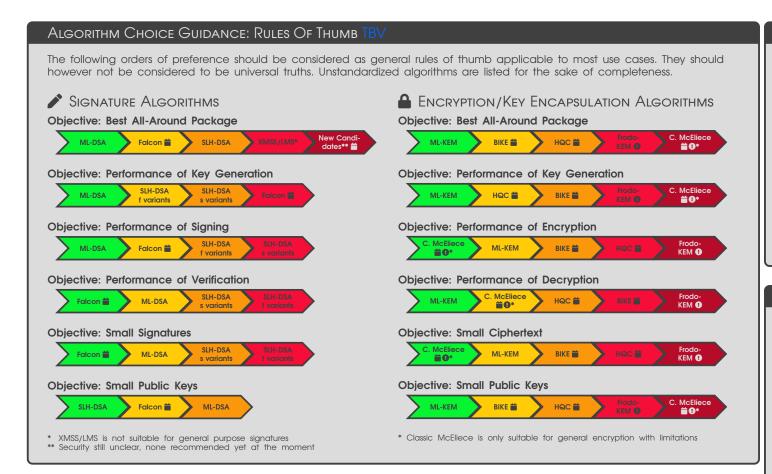
The approach of focusing on orders of magnitude walks a fine line between treating too many things as "equal" and not simplifying things enough to be easy to read and compare. "In the same order of magnitude" usually refers to "equal up to a factor of at most 10", which is a very coarse way of comparing numbers. Treating metrics that differ by a factor of e.g. 9.9 as "equal" because 9.9 < 10 paints a distorted picture. In cryptography, factors of 5 or even 2 can make a significant difference in security or performance, both in theory and in practice. In order to still tease out the differences in metrics without throwing too many things together that actually differ significantly, this cheat sheet applies different scaling and "orders of magnitude" (i.e., not regarding base 10) for different metrics.

It turns out that for metrics measured in (kilo) CPU cycle counts, i.e. algorithm performance, "up to a factor of 5" is a scale that is granular enough to work out the differences between algorithms while maintaining easy comparability. Those cycle counts heavily depend on the CPU used during measurement, hence the numbers need to be taken with a grain of salt, even if given exactly and not in terms of orders of magnitude.

For signature and ciphertext sizes as well as key sizes, measuring numbers in kilobytes "up to a factor of 2" is well suited to work out the differences between algorithms while allowing for quick comparison. Specifically for signature public key sizes only, we offset the corresponding color coding by 5 orders of magnitude. This is because SLH-DSA has extremely small pubic keys compared to all other signature algorithms, which would extend the scale into negative numbers (e.g., for SLH-DSA-SHA2-128s, the public key has $32=2^5$ bytes, which corresponds to an order of magnitude of -5 when measuring in orders of magnitude of base 2 and in kilobytes). This phenomenon of algorithm metrics spanning a very large range of orders of (base 2) magnitudes does not occur to this extent for encryption algorithms, making an offset unnecessary.

All values thus have a lower bound of 0. We do not limit the upper end of scales, but don't distinguish values areater than 10 anymore in terms of color codina. Please refer to the definitions on the left for symbol explanations, color coding and interpretation of numeric values.





Pure PQC vs PQ/T Hybrid

This topic depends on too many factors (e.g. cost of migration, security considerations, risk profile, GRC requirements) to give aeneral advice. Those aspects will differ greatly between different organizations. The main reasons to adopt PQ/T hybrids are to still have traditional algorithms in place in case a new algorithm turns out to be insecure, and that PQ/T might help to avoid a big bang migration as systems could simply ignore the PQC component if they do not support it yet. Consider recommendations from different government agencies: BSI and ANSSI recommend PQ/T hybrid strategies, whereas NIST is more reserved towards PQ/T. If using PQ/T hybrids, preferably use ECC (e.g. secp256r1, brainpoolP256r1, Curve25519) instead of RSA for the traditional component to keep key sizes as small as possible.

SECURITY CATEGORY CHOICES

- As a baseline, first consider using parameters.
- Use v or v for more security if possible (i.e., if a decrease in performance is not a concern and if no constraints apply).
- Use or iik if and only if iik or higher is not an option due to constraints (e.g. performance, memory, etc.).
- Comparing post-quantum security categories of algorithms to their traditional security level in bits is a complex subject and like comparing apples and oranges to some degree.

Pure vs. Pre-Hashing

- First, consider using pure (i.e., without pre-hashing) as this is the general recommendation.
- Pre-Hashing may be considered if one or more of the following apply:
 - The message M is too large to be sent to the cryptographic module (CM) for hashing without significantly impacting performance. This may be the case e.g. in CMS related use cases such as S/MIME or code signing, or in cases of very narrow communication channels to the CM (e.g. between APDUs exchanged between smartcard and smartcard reader).
 - The hash needs to be signed with different algorithms and would be computed repeatedly without pre-hashing.
 - The specific hash function is not supported in a CM.



SIGNATURE ALGORITHMS

ID Cards



ML-DSA (Module-Lattice-Based Digital Signature Algorithm)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
HASHING
NAMING

CRYSTALS-DILITHIUM \cite{O} FIPS 204 Signature Lattice Yes Standardized NIST, BSI, ANSSI Pure, Pre-Hashing by $k \times l$ matrix A (e.g. $6 \times 5 \rightarrow$ ML-DSA-65)

| IMPLEMENTATION | P | | 8 |
|----------------|-------------|------------|------------|
| Complexity | ? © | ? © | ? (S) |
| Size | ? \$ | ? S | ? ₽ |

| PARAMETER SET | OID | SECURITY CATEGORY | PER | FORMAI | NCE | SIGNATURE SIZE | PUBLIC KEY SIZE | Suitable Pre-Hashing |
|-------------------------------------|---|----------------------|----------------|----------------------------|----------------|-------------------|--------------------|---|
| ML-DSA-44 ML-DSA-65 ML-DSA-87 | 2.16.840.1.101.3.4.3.17 2.16.840.1.101.3.4.3.18 2.16.840.1.101.3.4.3.19 | | 30 30 40 | 3 3 3 0 4 0 | 20 30 30 | 100 | 5 5 6 | SHA-256, SHA3-256 SHA-384, SHA3-384 SHA-512, SHA3-512 |

/ Pros / Use If:

- You need a general-purpose signature algorithm with decent specs in all categories
- Usually the best option for a signature algorithm in most protocols (e.g. TLS, SSH, S/MIME, etc.).
- Good choice for X.509 certificates, including CA certificates.

/ Cons / Don't Use If:

You don't want a lattice-based algorithm



FALCON (FAST-FOURIER LATTICE-BASED COMPACT SIGNATURES OVER NTRU)



| Algorithm | Overall | Usability | Score |
|------------------|---------|-----------|--------------|

Falcon

Lattice

Yes

TBD

Signature

Pending

Project Page

PREVIOUS NAME SPECIFICATION TYPE **FAMILY** SUITABLE FOR GENERAL USE STANDARDIZATION STATUS RECOMMENDED BY HASHING

| Naming | TBD | | |
|----------------|----------------------|------------|------------|
| IMPLEMENTATION | P | ř | 8 |
| Complexity | H _{\$\psi}} | H | ? ₽ |
| Size | ? \$ | ? S | ? § |

| PARAMETER SET | OID | OID SECURITY PERFORMANCE SIGNATURE CATEGORY P P SIZE | | | | | | PUBLIC KEY SIZE | Suitable Pre-Hashing |
|---------------------------|------------|--|----|----------------|----------|----|----|--------------------|-------------------------|
| Falcon-512 Falcon-1024 | TBD TBD | U _® | 92 | 40 40 40 | 20 30 | 00 | 40 | TBD TBD | |

/ Pros / Use If:

- Generally good for use in most protocols (e.g. TLS, SSH, S/MIME, etc.). Exceptions may apply in cases of repeated generation of ephemeral keys (poor performance).
- Falcon has smaller signature sizes than ML-DSA:



• Falcon has smaller public key sizes than ML-DSA:

4, vs. 5, on Level I, 5, vs. 6, on Level V.

- You need a medium security category between 📭 and 🗸
- The algorithm is not yet standardized
- The algorithm requires expensive floating point arithmetic
- Key generation and signing are slower than for ML-DSA
- Less suitable in applications involving the creation of lots of ephemeral keys/certificates (key generation is slow).



SLH-DSA (Stateless Hash-Based Digital Signature Standard)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
HASHING
NAMING

SPHINCS+

FIPS 205
Signature
Hash (stateless)
Yes
Standardized
NIST, BSI, ANSSI
Pure, Pre-Hashing
based on various characteristics
(*s=small signatures, *f=fast)

| TATION | P | ř | 3 |
|--------|------------|------------|------------|
| TY | ? © | ? © | ? ₽ |

| | PARAMETER | OID | SECURITY PERFORMANCE | | NCE | SIGNATURE | PUBLIC KEY | SUITABLE | |
|---|--------------------|-------------------------|----------------------|--------------|----------|-----------|------------|-------------|-------------------|
| | SET | | CATEGORY | <i>P</i> / 0 | | Size | Size | Pre-Hashing | |
| - | | | | | | | | | |
| | SLH-DSA-SHA2-128s | 2.16.840.1.101.3.4.3.20 | | | 98 | 50 | 2 | 0 | SHA-256, SHA3-256 |
| | SLH-DSA-SHA2-128F | 2.16.840.1.101.3.4.3.21 | | 58 | 78 | 50 | 40 | 0 | SHA-256, SHA3-256 |
| | SLH-DSA-SHA2-192s | 2.16.840.1.101.3.4.3.22 | III | R | 93 | 50 | 40 | 0 | SHA-384, SHA3-384 |
| | SLH-DSA-SHA2-192F | 2.16.840.1.101.3.4.3.23 | III | 50 | R | 60 | 5 | 12 | SHA-384, SHA3-384 |
| | SLH-DSA-SHA2-256s | 2.16.840.1.101.3.4.3.24 | V | 7 8 | 93 | 50 | 40 | 0 | SHA-512, SHA3-512 |
| | SLH-DSA-SHA2-256F | 2.16.840.1.101.3.4.3.25 | Vø | 60 | | 60 | 5 | 100 | SHA-512, SHA3-512 |
| | SLH-DSA-SHAKE-128s | 2.16.840.1.101.3.4.3.26 | | R | 98 | 50 | 2 | 00 | SHA-256, SHA3-256 |
| | SLH-DSA-SHAKE-128F | 2.16.840.1.101.3.4.3.27 | | 50 | 78 | 50 | 40 | 00 | SHA-256, SHA3-256 |
| | SLH-DSA-SHAKE-192s | 2.16.840.1.101.3.4.3.28 | | R | 97 | 50 | 40 | 00 | SHA-384, SHA3-384 |
| ^ | SLH-DSA-SHAKE-192F | 2.16.840.1.101.3.4.3.29 | III | 50 | 78 | 60 | 5 | 10 | SHA-384, SHA3-384 |
| 5 | SLH-DSA-SHAKE-256s | 2.16.840.1.101.3.4.3.30 | V | R | 93 | 50 | 40 | 00 | SHA-512, SHA3-512 |
| | SLH-DSA-SHAKE-256F | 2.16.840.1.101.3.4.3.31 | V | 60 | B | 60 | 5 | 10 | SHA-512, SHA3-512 |

/ Pros / Use If:

IMPLEMENT

COMPLEXI

- Alternative to ML-DSA and Falcon that is not based on lattices
- Very small public keys
- Generally good for use in most protocols, even if ML-DSA is the preferred option.
- Use as a more conservative choice in use cases with less frequent algorithm execution, i.e. where latency or message size are not a big concern. For example, S/MIME, document or code signing, creation of CAs.

- Poor key generation and signing performance compared to other algorithms
- High complexity of the algorithm and the implementation
- Possible interoperability issues due to the many variants that may not all be supported everywhere



XMSS / XMSS-MT (eXtended Merkle Signature Scheme / eXtended Merkle Signature Scheme Multi Tree)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
HASHING
NAMING

| XMSS/XMSSMT |
|--|
| 𝚱 SP 800-208, 𝚱 RFC 8391 |
| Signature |
| Merkle Trees (stateful hash trees) |
| No |
| Standardized |
| NIST, BSI, ANSSI |
| TBD |
| XMSS-[Hashfamily]_[h]_[n] |
| XMSSMT-[Hashfamily]_[h]/[d]_[n] |
| where h is the tree height, |
| d is the number of layers, and |
| n is the message length in bits |
| |

| IMPLEMENTATION | P | • | 0 |
|----------------|-------------|------------|--------------|
| Complexity | ? <u>\$</u> | ? ₽ | ? ₽ |
| Size | ? \$ | ? S | ? () |

| PARAMETER SET | NUMERIC IDENTIFIER | SECURITY CATEGORY | PER | FORMAN | () | SIGNATURE SIZE | MAXIMUM SIGNATURES | NUMBER OF LAYERS |
|-----------------------|-----------------------|----------------------|--------------|------------|----------------|-------------------|-----------------------|------------------|
| XMSS-SHA2_10_256 | 0x00000001 | V | ?? | ? 8 | ?0 | ? | 2 ¹⁰ | 1 |
| XMSS-SHA2_16_256 | 0x00000002 | V | ₹ | ? <u>~</u> | 70 | ? | 216 | 1 |
| XMSS-SHA2_20_256 | 0x00000003 | Vo | ?8 | 76 | ?0 | ? | 2 ²⁰ | 1 |
| XMSSMT-SHA2_20/2_256 | 0x00000001 | V | ? <u>?</u> ? | ?? | ?0 | ? | 2 ²⁰ | 2 |
| XMSSMT-SHA2_20/4_256 | 0x00000002 | V | ? | ? | ₹ <u>7</u> 0 | ? | 2 ²⁰ | 4 |
| XMSSMT-SHA2_40/2_256 | 0x00000003 | V | ? <u>?</u> ? | 78 | 70 | ? | 2 ⁴⁰ | 2 |
| XMSSMT-SHA2_40/4_256 | 0x00000004 | V | ?? | ? | 70 | ? | 2 ⁴⁰ | 4 |
| XMSSMT-SHA2_40/8_256 | 0x00000005 | V | ? <u>?</u> ? | ? | ? 0 | ? | 2 ⁴⁰ | 8 |
| XMSSMT-SHA2_60/3_256 | 0x00000006 | V | ? | ? 8 | 70 | ? | 2 ⁶⁰ | 3 |
| XMSSMT-SHA2_60/6_256 | 0x00000007 | V | ? <u>?</u> ? | ? | ?0 | ? | 260 | 6 |
| XMSSMT-SHA2_60/12_256 | 0x00000008 | V | ?8 | ? 6 | ?0 | ? | 2 ⁶⁰ | 12 |

Note

₱ SP 800-208 defines further parameter sets not listed in ₱ RFC 8391 using other hash functions (SHA256/192, SHAKE256/256, SHAKE256/192). Furthermore, ₱ RFC 8391 lists optional parameter sets that are not approved in ₱ SP 800-208. All of those variants are omitted here as they are not likely to be widely used, in particular not after ML-DSA and SLH-DSA have been standardized in the meantime.

/ Pros / Use If:

- You can predict the maximum number of signatures that are going to be required
- Firmware signing use cases
- You want a signature scheme where the security only relies on the security of the hash function used without assuming the hardness of another mathematical problem.
- Cf. & SP 800-208, Section 1.1 for additional explanations

- You require an algorithm for general use
- You cannot predict the maximum number of signatures that are going to be required, or the number of required signatures exceeds the maximum number of signatures enabled through the approved parameter sets
- Your application does not allow for the careful state management and tracking of signatures performed that is required with this algorithm



LMS/HSS (Leighton-Micali Signature / Hierarchical Signature System)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
HASHING
NAMING

LMS, LMS/HSS

SP 800-208, RFC 8554
Signature
Merkle Trees (stateful hash trees)
No
Standardized
NIST, BSI, ANSSI
TBD
LMS_SHA256_M32_[h]
where h is the tree height

| IMPLEMENTATION | P | ř | • |
|----------------|------------|------------|-----|
| COMPLEXITY | ? © | ? <u>©</u> | 7°C |
| Size | ? | ? | ? |

| PARAMETER SET | NUMERIC IDENTIFIER | SECURITY CATEGORY | PER | FORMAN | (CE | SIGNATURE SIZE | MAXIMUM SIGNATURES |
|--|--|----------------------|----------|----------------|----------------|-------------------|---|
| LMS_SHA256_M32_H5 LMS_SHA256_M32_H10 | 0x00000005 0x00000006 | ? . | ?0 ?0 | ? 6 | ?0 ?0 | ? ₀ | 2 ⁵ 2 ¹⁰ |
| LMS_SHA256_M32_H15 LMS_SHA256_M32_H20 LMS_SHA256_M32_H25 | 0x00000007 0x00000008 0x00000009 | ? | ?0 | ?0 ?0 ?0 | ?0 ?0 ?0 | ? ₀ | 2 ¹⁵ 2 ²⁰ 2 ²⁵ |

NOTE:

♦ SP 800-208 defines further parameter sets not listed in ♠ RFC 8554 using other hash functions (SHA256/192, SHAKE256/256, SHAKE256/192). Those variants are omitted here as they are not likely to be widely used, in particular not after ML-DSA and SLH-DSA have been standardized in the meantime. Similar to XMSS multi-tree variants, a comparable structure exists with Hierarchical Signature System (HSS) variants of LMS that involves a number of LMS instances.

/ Pros / Use If:

- You can predict the maximum number of signatures that are going to be required
- Firmware signing use cases
- You want a signature scheme where the security only relies on the security of the hash function used without assuming the hardness of another mathematical problem.
- Cf. & SP 800-208, Section 1.1 for additional explanations

- You require an algorithm for general use
- You cannot predict the maximum number of signatures that are going to be required, or the number of required signatures exceeds the maximum number of signatures enabled through the approved parameter sets
- Your application does not allow for the careful state management and tracking of signatures performed that is required with this algorithm



ENCRYPTION / KEM ALGORITHMS

ID Cards



ML-KEM (Module-Lattice-Based Key-Encapsulation Mechanism Standard)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

CRYSTALS-KYBER

Encryption/KEM

Standardized

9 FIPS 203

Lattice

| Parameter Set | OID | SECURITY CATEGORY | Per P | FORMAN | CE CE | CIPHERTEXT SIZE | PUBLIC KEY SIZE |
|------------------|------------------------|----------------------|----------|--------|------------------|--------------------|--------------------|
| ML-KEM-512 | 2.16.840.1.101.3.4.4.7 | | 20 | 20 | 2 <mark>0</mark> | <u>о</u> | • |
| ML-KEM-768 | 2.16.840.1.101.3.4.4.2 | III | 28 | 20 | 3 0 | o _O | 0 |
| ML-KEM-1024 | 2.16.840.1.101.3.4.4.3 | V | 30 | 30 | 3 O | <u>o</u> | 00 |

/ Pros / Use If:

IMPLEMENTATION

COMPLEXITY

SIZE

- Currently the only post-quantum encryption/key encapsulation algorithm standardized by NIST
- Need a general-purpose encryption / key-encapsulation algorithm with decent specs in all categories
- Usually the best choice in most protocols involving an encryption/KEM component (e.g. TLS, S/MIME, etc.)

/ Cons / Don't Use If:

• You don't want a lattice-based algorithm



BIKE (BIT FLIPPING KEY ENCAPSULATION)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

BIKE
Project Page
Encryption/KEM
Codes
Yes
4th round candidate
TBD
after security categories I, III and V

| IMPLEMENTATION | P | | • |
|----------------|-------------|------------|------------|
| Complexity | ? <u>\$</u> | ? © | ₹ |
| Size | ?\$ \$ | ? \$ | ? S |

| PARAMETER SET | OID | SECURITY CATEGORY | PER | FORMAN | ICE | CIPHERTEXT SIZE | PUBLIC KEY SIZE |
|------------------|-----|----------------------|------------------|------------|------------|--------------------|--------------------|
| BIKE-L1 | TBD | | 3 <mark>2</mark> | 20 | 4 <u>C</u> | <u></u> | • |
| BIKE-L3 | TBD | III _® | 40 | 3 0 | 5 <u>O</u> | 10 | 100 |
| BIKE-L5 | TBD | V® | n/g | n/g | n/g | 2 ₀ | 20 |

NOTE:

No data available for BIKE-L5 for cycle counts. The algorithm overall usability score is computed over BIKE-L1 and BIKE-L3 only.

/ Pros / Use If:

• Possible alternative to ML-KEM not based on lattices.

- Not yet standardized. Note that NIST intends to standardize at most one of the algorithms HQC and BIKE.
- Metrics not as good as the ones of ML-KEM.



HQC (Hamming-Quasi-Cyclic)



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

HQC

Project Page
Encryption/KEM
Codes
Yes
4th round candidate
TBD
128, 192 and 256 bits of security
in reference to security categories

| IMPLEMENTATION | P | | 0 |
|----------------|-------------|------------|------------|
| Complexity | ? \$ | ? ₽ | ? ₽ |
| Size | ? S | ? S | ? S |

I, III and V

/ Pros / Use If:

• Possible alternative to ML-KEM not based on lattices.

- Not yet standardized. Note that NIST intends to standardize at most one of the algorithms HQC and BIKE.
- Metrics not as good as the ones of ML-KEM.



FRODOKEM



Algorithm Overall Usability Score

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

FrodoKEM

Project Page
Encryption/KEM
Lattice
Yes
Disregarded by NIST
BSI, ANSSI
FrodoKEM-[n]-[AES | SHAKE]
where n is a matrix dimension

| IMPLEMENTATION | P | | 8 |
|----------------|------------|------------|-------------|
| Complexity | ? % | ? ₽ | ? % |
| Size | ? 6 | ? S | ? \$ |

| Parameter Set | OID | SECURITY CATEGORY | PER | FORMAN | CE CE | CIPHERTEXT SIZE | PUBLIC KEY SIZE |
|-------------------------------------|------------|----------------------|----------|----------|----------|----------------------------------|--------------------|
| Frodokem-640-Aes | TBD TBD | <u>G</u> | 40 | 40 50 | 40 50 | 30 | 30 |
| FRODOKEM-976-AES FRODOKEM-976-SHAKE | TBD TBD | | 40 60 | 40 | 40 60 | ³ 0 | 30 |
| FRODOKEM-1344-AES | TBD | V _® | 50 | 60 50 | 50 | ³ 0 ⁴ 0 | ³ 0 |
| FRODOKEM-1344-SHAKE | TBD | V® | 60 | 60 | 60 | 40 | 40 |

NOTE

There are also ephemeral versions of FrodoKEM (eFrodoKEM-640, etc.), but BSI only recommends FrodoKEM-976 and FrodoKEM-1344 versions for long-term protection. The algorithm score also only includes the non-ephemeral versions listed above.

/ Pros / Use If:

• More conservative than ML-KEM since it is based on unstructured grids

/ Cons / Don't Use If:

• Not standardized by NIST because ML-KEM is more efficient.



CLASSIC MCELIECE



PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY

NAMING

Classic McEliece $\rat{9}$ Project Page Encryption/kEM Codes No / With Limitations 4th round candidate (NIST) BSI mceliece[n][deg(F(y))][[f]] where n, F(y) are parameters non-f versions: $(\mu, \nu) = (0, 0)$ f versions: $(\mu, \nu) = (32, 64)$

| IMPLEMENTATION | P | | • |
|----------------|------------|------------|------------|
| Complexity | ? <u>©</u> | ? © | ? ₽ |
| Size | ? S | ? S | ? § |

| PARAMETER SET | OID | SECURITY CATEGORY | PER | FORMAN | CE NCE | CIPHERTEXT SIZE | PUBLIC KEY SIZE |
|------------------|-----|----------------------|----------|------------|------------|--------------------|--------------------|
| MCELIECE348864 | TBD | | 60 | 20 | 30 | <u>о</u> | 8, |
| MCELIECE348864F | TBD | | 68 | n/8 | n/B | 00 | 8 |
| MCELIECE460896 | TBD | III | 8 | 20 | 30 | 0 0 | 9, |
| MCELIECE460896F | TBD | III | 8 | n/S | n/S | 0 0 | 9, |
| MCELIECE6688128 | TBD | V | R | 3 <u>0</u> | 3 O | o _O | 10, |
| MCELIECE6688128F | TBD | Vo | 8 | n/S | n/e | o _O | 10, |
| MCELIECE6960119 | TBD | V | R | 30 | 3 <u>O</u> | o _O | 10, |
| MCELIECE6960119F | TBD | V | R | n/e | n/e | 0 | 10, |
| MCELIECE8192128 | TBD | Vo | | 30 | 30 | ○ | 10, |
| MCELIECE8192128F | TBD | V | B | n/g | n/e | 00 | 10, |

Note:

No data available for the f versions for encryption/decryption cycle counts. The algorithm overall usability score is computed over non-f versions only. Non-f and f versions are interoperable, with f versions offering faster key generation and non-f versions having simpler key generation.

/ Pros / Use If:

- Possible alternative to ML-KEM not based on lattices.
- Considered safe and recommended by BSI even though not standardized by NIST.
- Use cases that allow for pre-sharing of keys where the large public key size is not a concern.

- Not yet standardized by NIST.
- Not suitable for general use due to the very large public keys and poor performance of key generation.
- Due to aforementioned point, likely not suitable for embedded devices.



NTRU (Number Theory Research Unit / Number Theorists 'R' Us)



Algorithm Overall Usability Score

NTRU

PREVIOUS NAME
SPECIFICATION
TYPE
FAMILY
SUITABLE FOR GENERAL USE
STANDARDIZATION STATUS
RECOMMENDED BY
NAMING

Project Page Encryption/KEM Lattice Yes Not Standardized n/a author initials and a prime number

| IMPLEMENTATION | P | | 8 |
|----------------|------------|------------|------------|
| Complexity | ? % | ? © | ? © |
| SIZE | ? ⑤ | ? S | ? ⑤ |

| PARAMETER SET | OID | SECURITY CATEGORY | Per P | FORMAN | CE NCE | CIPHERTEXT SIZE | PUBLIC KEY SIZE |
|------------------|-----|----------------------|----------|------------------|------------|--------------------|--------------------|
| NTRUHPS2048509 | N/A | | 50 | <mark>4</mark> 0 | 4 <u>0</u> | <u></u> | • |
| NTRUHRSS701 | N/A | III _® | eg eg | 40 | <u>5</u> 0 | 00 | 00 |
| NTRUHPS2048677 | N/A | III _® | eg . | 40 | 50 | ○ | 00 |
| NTRUHPS4096821 | N/A | Vo | 60 | 40 | 5 <u>0</u> | 00 | 00 |

/ Pros / Use If:

• The intellectual property rights of ML-KEM have been unclear for a while, with NTRU being a possible backup alternative without such issues. This topic is now considered sufficiently clear and resolved by NIST. In case legal concerns about ML-KEM still apply in some exceptional cases, NTRU might be an option. However, ML-KEM is usually the better option in terms of technical aspects.

- Not standardized
- Performance not as good as ML-KEM, especially for key generation. Similarly, BIKE and HQC offer better metrics.







Overview: Hash Algorithms and Extendable Output Functions (XOF) TBV

| FAMILY | VARIANT | OUTPUT SIZE [BITS] | ROUNDS | Construction | TRADITIONA | L PRE-QUANTUM SE | CURITY [BITS] | Post- | QUANTUM SECURITY | [Вітѕ] | PERFORMANCE |
|--------|-------------|--------------------|--------|-----------------------------|-------------------------------------|-------------------------------------|-------------------------------------|-------------------------------------|-------------------------------------|-------------------------------------|-------------|
| | | | | | Collision | Pre-Image | 2nd Pre-Image | Collision | Pre-Image | 2nd Pre-Image | |
| | SHA-224 | 224 | 64 | Merkle-Damgård construction | 112 | 224 | 224 | 74 | 112 | 112 | 7.62 |
| | SHA-256 | 256 | 64 | Merkle-Damgård construction | 128 | 256 | 256 | 85 | 128 | 128 | 7.63 |
| SHA-2 | SHA-384 | 384 | 80 | Merkle-Damgård construction | 192 | 384 | 384 | 128 | 192 | 192 | 5.12 |
| | SHA-512 | 512 | 80 | Merkle-Damgård construction | 256 | 512 | 512 | 170 | 256 | 256 | 5.06 |
| | SHA-512/224 | 224 | 80 | Merkle-Damgård construction | 112 | 224 | 224 | TBD | TBD | TBD | ≈ SHA-384 |
| | SHA-512/256 | 256 | 80 | Merkle-Damgård construction | 128 | 256 | 256 | TBD | TBD | TBD | ≈ SHA-384 |
| | SHA3-224 | 224 | 24 | Keccak sponge construction | 112 | 224 | 224 | 74 | 112 | 112 | 8.12 |
| | SHA3-256 | 256 | 24 | Keccak sponge construction | 128 | 256 | 256 | 85 | 128 | 128 | 8.59 |
| SHA-3 | SHA3-384 | 384 | 24 | KECCAK SPONGE CONSTRUCTION | 192 | 384 | 384 | 128 | 192 | 192 | 11.06 |
| 3HA-3 | SHA3-512 | 512 | 24 | Keccak sponge construction | 256 | 512 | 512 | 170 | 256 | 256 | 15.88 |
| | SHAKE128 | d (ARBITRARY) | 24 | Keccak sponge construction | $\min\left(\frac{d}{2};128\right)$ | $\min\left(\frac{d}{2};128\right)$ | $\min\left(\frac{d}{2};128\right)$ | $\min\left(\frac{d}{3};128\right)$ | $\min\left(\frac{d}{2};128\right)$ | $\min\left(\frac{d}{2};128\right)$ | 7.08 |
| | SHAKE256 | d (ARBITRARY) | 24 | Keccak sponge construction | $\min\left(\frac{d}{2}; 256\right)$ | $\min\left(\frac{d}{2}; 256\right)$ | $\min\left(\frac{d}{2}; 256\right)$ | $\min\left(\frac{d}{3}; 256\right)$ | $\min\left(\frac{d}{2}; 256\right)$ | $\min\left(\frac{d}{2}; 256\right)$ | 8.59 |

NOTES & TAKEAWAYS

- ullet SHAKE128 and SHAKE256 are exendable output functions and allow for an arbitrary output size d.
- SHA-512/224 and SHA-512/256 denotes the truncation of the SHA-512 hash to 224 resp. 256 bits.
- Security against collision attacks is specified using values according to Brassard et al. (Source)
- Performance is measured in median cycles/byte on a Skylake CPU (Source).
- Hash functions other than SHA-2 and SHA-3 variants (e.g. SHA-1, RipeMD-160, MD5, etc.) should generally be avoided as they are less secure. They are therefore not part of the above overview.







| Traditional vs. I | Post-Quantum | Security Level | . Comparison TB | |
|-------------------|--------------|----------------|-----------------|--|
|-------------------|--------------|----------------|-----------------|--|

| TRADITIONAL PRE | CLASSICAL ATTACK COMPLEXITY | ATTACK IN TERMS OF BIT SIZE N | ALGORITHM EXAMPLES | NIST SECURITY CATEGORIES DEFINED THROUGH ATTACK | Classical Attack | QUANTUM ATTACK COMPLEXITY | POST-QUANTUM SE SECURITY [BITS] | SIGNATURE PARAMETER SETS | Encryption/KEM Parameter Sets |
|------------------------|---|---|-----------------------|---|------------------|---|------------------------------------|-------------------------------------|--|
| ${n}$ | $\mathcal{O}(2^n)$ | $n	ext{-BIT}$ KEY SPACE EXHAUSTIVE KEY SEARCH | AFS | | GROVER | $O\left(2^{\frac{n}{2}}\right)$ | <u>n</u> 2 | FARAMETER SEIS | PARAMETER SEIS |
| n | $\mathcal{O}\left(2^{n}\right)$ | n-BIT HASH PRE-IMAGE ATTACK | SHA-2, SHA-3 | | GROVER | $\mathcal{O}\left(2^{\frac{n}{2}}\right)$ | $\frac{\overline{2}}{n}$ | | |
| | ` í | | SHA-2, SHA-3 | | | $\mathcal{O}\left(2^{\frac{2}{2}}\right)$ $\mathcal{O}\left(2^{\frac{n}{2}}\right)$ | 2 | | |
| n | $\mathcal{O}\left(2^{n}\right)$ | n-bit hash 2nd-pre-image attack | | | GROVER | \ / | $\frac{n}{2}$ | | |
| $\frac{\frac{n}{2}}{}$ | $\mathcal{O}\left(2^{\frac{n}{2}}\right)$ | n-BIT HASH COLLISION ATTACK | SHA-2, SHA-3 | | Brassard et al. | $\mathcal{O}\left(2^{\frac{n}{3}}\right)$ | $\frac{n}{3}$ | | |
| 128 | $\mathcal{O}\left(2^{128}\right)$ | 128-bit key space exhaustive key search | AES-128 | Category 1 | Grover | $\mathcal{O}\left(2^{64}\right)$ | 64 | Falcon-512, SLH-DSA 128 variants | ML-KEM-512, BIKE-L1, HQC-128, FrodoKEM-640 |
| | | 128-bit hash pre-image attack | N/A | | Grover | $\mathcal{O}\left(2^{64}\right)$ | 64 | | |
| | | 128-bit hash 2nd-pre-image attack | N/A | | GROVER | $\mathcal{O}\left(2^{64}\right)$ | 64 | | |
| | | | | | | | | | |
| 192 | $\mathcal{O}\left(2^{192}\right)$ | 192-bit key space exhaustive key search | AES-192 | Category 3 | Grover | $\mathcal{O}\left(2^{96}\right)$ | 96 | ML-DSA-67, SLH-DSA 192 variants | ML-KEM-768, BIKE-L3, HQC-192, FRODOKEM-976 |
| | | 192-bit hash pre-image attack | N/A | | Grover | $\mathcal{O}\left(2^{96}\right)$ | 96 | | |
| | | 192-bit hash 2nd-pre-image attack | N/A | | GROVER | $\mathcal{O}\left(2^{96}\right)$ | 96 | | |
| | | 384-bit hash collision attack | SHA-384, SHA3-384 | Category 4 | BRASSARD ET AL. | $\mathcal{O}\left(2^{128}\right)$ | 128 | | |
| 256 | $\mathcal{O}\left(2^{256}\right)$ | 256-bit key space exhaustive key search | AES-256 | Category 5 | Grover | $\mathcal{O}\left(2^{128}\right)$ | 128 | ML-DSA-85, SLH-DSA-256, variants | ML-KEM-1024, BIKE-L5, HQC-256, |
| | | 256-bit hash pre-image attack | N/A | | Grover | $\mathcal{O}\left(2^{128}\right)$ | 128 | | |
| | | 256-bit hash 2nd-pre-image attack | N/A | | Grover | $\mathcal{O}\left(2^{128}\right)$ | 128 | TBVXMS | FrodoKEM-1344 |
| | | 512-bit hash collision attack | SHA-512, SHA3-512 | | Brassard et al. | $\mathcal{O}\left(2^{170}\right)$ | 170 | | |
| 384 | $\mathcal{O}\left(2^{384}\right)$ | 384-bit key space exhaustive key search | N/A | | Grover | $\mathcal{O}\left(2^{192}\right)$ | 192 | | |
| | | 384-bit hash pre-image attack | SHA-384, SHA3-384 | | Grover | $\mathcal{O}\left(2^{192}\right)$ | 192 | | |
| | | 384-bit hash 2nd-pre-image attack | SHA-384, SHA3-384 | | Grover | $O\left(2^{192}\right)$ | 192 | | |
| | | 768-bit hash collision attack | N/A | | Brassard et al. | $\mathcal{O}\left(2^{256}\right)$ | 256 | | |
| 512 | $\mathcal{O}\left(2^{512}\right)$ | 512-bit key space exhaustive key search | N/A | | Grover | $\mathcal{O}\left(2^{256}\right)$ | 256 | | |
| | | 512-bit hash pre-image attack | SHA-512, SHA3-512 | | Grover | $\mathcal{O}\left(2^{256}\right)$ | 256 | | |
| | | 512-bit hash 2nd-pre-image attack | SHA-512, SHA3-512 | | Grover | $\mathcal{O}\left(2^{256}\right)$ | 256 | | |
| | | 1024-bit hash collision attack | N/A | | Brassard et al. | $\mathcal{O}\left(2^{341}\right)$ | 341 | | |

NOTES & TAKEAWAYS

- The above is still a simplified view. It does not dive into the specific assumptions made in the new algorithms' security proofs and the possible attacks on those algorithms (e.g., is a collision attack on a hash function enough, or would a pre-image be required to break the signature scheme using that hash function). It is also not including considering practical considerations and predictions about quantum computers, such as runtime or quantum circuit depth. Cf. 8 NIST PQC Evaluation Criteria, 4.A.5 for further details.
- A pre-quantum security level can have different post-quantum security levels equivalents. It all depends on the specific attack in question, resp. the security assumption a new post-quantum algorithm is based on or must be secure against. For hashing, the often cited "SHA-256 offers 128 bits of security" is not wrong, but only part of the picture: It offers 256 bits of security against pre-image and 2nd-preimage-attacks. The n bits of security against 2nd-preimage attacks is still a slightly simplified and harmonized view to avoid different formulas for different hash functions. Cf. 9 NIST: Hash Functions for details.
- The fact that the exact level of post-quantum security still depends on so many assumptions about CRQCs, comparing pre-quantum and post-quantum security levels remains an apples and oranges comparison to some degree. Such comparisons should be taken with a grain of salt, especially when done only on a rather superficial level.
- The lines marked with colors , and are the ones used by NIST to define categories I, II, III, IV and V. The lines marked in lighter colors , and have the same post-quantum security as the defining lines of categories I, III and V, respectively. They therefore fulfill the "comparable to or greater than" clause in NIST's category definitions. Note that categories IV and V are differentiated despite having 128 bits of post-quantum security again highlighting the fact that a precise categorization is difficult and requires a detailed analysis of many aspects.



ALGORITHM OVERALL USABILITY SCORE

We try to measure an algorithm's overall usability for a general use case without any special characteristics by calculating a single number between 0 (worst) and 10 (best) for the algorithm. This calculation is taking into account all parameter sets, performance metrics, public key size, signature/ciphertext size, the number of security categories provided, whether or not it is suitable for general use, and the complexity and size of its implementation. We define an algorithm's overall score as

$$\mathsf{score}_{\textit{algorithm}} = \mathsf{max}\left\{\mathsf{0}, \ \mathsf{avg}\left\{\mathsf{score}_{\textit{parameterSet}} \mid \textit{parameterSet} \ \text{is a parameter set of} \ \textit{algorithm}\right\} - \frac{1}{8} \cdot \left(\mathsf{5} - \mathsf{categories}_{\textit{algorithm}}\right) - \frac{1}{4} \cdot \mathsf{impl}_{\textit{algorithm}} - \mathsf{generality}_{\textit{algorithm}}\right\} \right\}$$

where score parameter. Set is a score for an individual algorithm parameter set. Depending on the type of algorithm, score parameter. Set is defined as

$$score_{signature-parameterSet} = 10 - avg \left\{ \begin{array}{cccc} \mathbf{n} & \mathbf{n} & \mathbf{n} & \mathbf{n} & \mathbf{n} \\ \mathbf{n} & \mathbf{n} & \mathbf{n} & \mathbf{n} \end{array} \right\}$$







 $score_{encryption-parameterSet} = 10 - avg \left\{ \begin{array}{ccc} n_{e}^{2}, & n_$







Furthermore, $1 \le$ categories $algorithm \le 5$ denotes the number of different NIST security categories offered by algorithm. By assigning numeric values of 0 = 100, 1 = 100



















Finally, we define

$$generality_{algorithm} = \begin{cases} 0 & \text{if } algorithm \text{ is a general purpose algorithm} \\ 2 & \text{else} \end{cases}$$

to take into account if the algorithm is suitable for general use.

TBD: In the algorithm scores given in ID cards, we currently use impl_{algorithm} = 0 in the respective computations because the necessary values for implementation complexity and size are still TBD. This will be corrected later.

Example: ML-DSA Overall Usability Score

We calculate









Similarly, we obtain $score_{ML-DSA-65} = 7.0$ and $score_{ML-DSA-65} = 6.2$. Furthermore, categories $m_{L-DSA} = 3$ since ML-DSA offers the three security categories $m_{L-DSA-65} = 7.0$ and $m_{L-DSA-65} = 0$ since ML-DSA is a general purpose signature algorithm. This results in an overall usability score of 6.55:







ML-DSA Overall Usability Score

$$score_{ML-DSA} = \max \left\{ 0, \text{ avg } \{ score_{ML-DSA-44}, \text{ score}_{ML-DSA-65}, \text{ score}_{ML-DSA-87} \} - \frac{1}{8} \cdot (5 - \text{categories}_{ML-DSA}) - \frac{1}{4} \cdot \text{impl}_{ML-DSA} - \text{generality}_{ML-DSA} \right\}$$

$$= \max \left\{ 0, \text{ avg } \{ 7.2, 7.0, 6.2 \} - \frac{1}{8} \cdot (5-3) - 0 - 0 \right\}$$

$$= \max \left\{ 0, 6.8 - 0.25 \right\}$$

$$= \max \left\{ 0, 6.55 \right\}$$

$$= 6.55$$

TBD: We use $impl_{MI-DSA} = 0$ because the necessary values to compute it are still TBD, cf. ML-DSA ID card. This will be corrected later.