

Computer Architecture - Lecture 5

- Intel 8080/8086
 - Jump (branch) instructions are either unconditional or conditional.
 - Unconditional jump:
 Jump <address>
 - The address can be near or far. Near address is a relative address (measured relative the current value of ip). Far address can be a complete address (segment:offset).
 - Technically, the jump instruction modifies the program counter(ip).
 - In assembly we write:
 Jump <lable>
 - The label is interpreted as an address.
 - Example:
 jmp L1
 mov ax, 7
 L1: mov ax, 10
 - Conditional branch (jump):
 - For conditional jump, a register called 'flags' is added to register files. The flags register is accessed bitwise. Examples of flags bits:
 - Carry flag (1 if the last operation produces a carry)
 - Example:
 - Mov AL,190
 - Add AL, 67 ; AL should be 257 but since max value is 255, a carry occurs and AL=2
 - The carry flag is set (equals 1).
 - The sign flag is set when the sign bit is 1 (negative number).
 Example:
 Mov AX,76
 Sub AX-80 ; negative result (-4) and sign flag is set(1).
 - The overflow flag: if the content of the last instruction destination overflows (becomes greater than the maximum value)
 - The zero flag: if the content of the destination is zero.
 - Conditional jump instructions are based on flag values,
 - JC <address/label>; jump if the carry flag is set.
 - JNC <address, label>; jump if the carry flag is clear.
 - JS <address/label> ; jump if the sign flag is set.
 - JNS <address/label> ; jump if the sign flag is clear.
 - Also, higher level forms are supported by assemblers:
 - JL <address/label> ; jump if less (destination <source> in last subtraction) (signed)

- JG <address/label> ; jump if greater (signed).
- JNL <address/label> ; jump if not less.
- JGE <address/label> ; jump if greater or equal.
- JA <address/label> ; jump if above (unsigned).
- JB <address/label> ; jump if below (unsigned).
- Example:
 - Mov AL, 190
 - Add AL, 67
 - JC L1; will jump to L1 since the carry flag is set.
- Example:
 - Mov AL, 10
 - Sub AL, 15 (AL changes to -5)
 - JC L1; will jump since AL < 15.
- CMP and text instruction:
 - The CMP instruction compares the destination with the source (by subtraction destination – source without storing the result, just change the flags)
- Example:
 - Mov AL, 10
 - Cmp AL, 15 (AL remains 10)
 - JL L1 ; jumps since AL < 15.
- Example:
 - Mov AL, 130; (binary 10000010) As a signed number (2's complement) = - (01111110)
 - Cmp AL, 132; (binary 10000100) As a signed number (2's complement) = - (01111100)
 - JA and JB will behave different than JG and JL.
 - JB will jump but JL will not jump.
- Example:
 - Mov AL, -3 - (00000011) → 11111101
 - Cmp AL, 5 (00000101)
 - JA will jump but JG will not jump)
- Loop instruction:
 - Loop <address/label>
 - This is a composite instruction equivalent to:
 - Dec CX
 - Cmp CX, 0
 - Jnz <address/label>

- To implement a loop of 10 iterations:
 - Mov cx, 10
 - L1: xxx (several lines of code)
 - Loop L1
- Stack instructions:
 - Stack is a memory area accessed in a LIFO order (last in first out).
 - The top of stack is pointed to SS (stack segment) : SP (stack pointer).
 - Basic stack operations
 - Push <src>; stores the value of the src to the top of stack.
 - Example:
 - Mov AX, 50
 - Push AX; stores 50 on the top of the stack
 - Mov AX, 90
 - Push AX; stores 90 on the top of the stack (above 50)
 - Pop <dst>; removes the topmost stack value and store it in destination.
 - Example:
 - Pop BX; removes 90 from the stack and store it in BX.
 - Pop CX; removes 50 from the stack and store it in CX.
 - Technically, push is a composite instruction.
 - Push <src> is equivalent to:
 - Sub SP, <size of the source>
 - Mov ss:[sp], src
 - Example: Let the SP = 500 then if we write:
 - Mov AX, 50
 - Push AX
 - Then, SP will be 498. Memory at the address ss:[498] contains 50.
 - Mov AX, 50
 - Push AX
 - Then, SP will be 496. Memory at the address ss:[496] contains 90.
- The pop instruction (POP <dst>) is a composite instruction equivalent to:
 - Mov dst, ss:[sp]
 - Add ss:[sp], size of <dst>
- Example:
 - Pop BX
 - BX = ss:[496] = 90
 - SP will be 498
 - Pop CX
 - CX = ss:[498] = 50
 - Sp will be 500

- In programming languages, stack is used for temporary storage [local variables and formal parameters of functions. It also supports the function calls.
- Call and return instruction:
 - Call <address/label>, calls a procedure (jump and wait to return)
 - Ret; returns to address after the last call.
- Example:
 - ...
 - Call L1; Pushes L2 into the stack and jumps to L1.
 - L2: xxx (several lines of code)
 - L1: xxx (several lines of code)
 - Ret; pops the top of stack into instruction pointer (ip) [i.e., jump to L2]
- X86 32-bit processors:
 - Processors have the same names, but a prefix e added:
 - Registers: EAX, EBX, ECX, EDX, ESI, EDI, ESP, EBP, EIP, EFlags, etc
 - Each is 32-bit registers.
 - AX is the lower 16-bits of EAX, BX is the lower 16-bits of EBX, etc.
 - AL is the lower 8 bits of EAX, AH is the higher 8 bits of (AX). i.e., the second lower 8 bits of EAX.
- X86 64-bit processors:
 - Registers have the same names with a prefix r.
 - Examples: RAX, RBX, RCX, RDX, etc.
 - Also, some registers added like: R8, R9, ... R15.
 - Each register is 64 bits.
 - EAX is the lower 32-bits of RAX, etc.