-: Data Base Management System: Database - Collection of related Data (RDBMS) Unstructured — webpages DBMS - SQL server → Oracle 91,11,12cetc 4-TK CTC → MySQL -university $\rightarrow DB_2$ # File Vs System DBWZ - searching is fast - efficient memory utilization -> Concurrency (protocols exist) Security (Role based) Access control) > Reduce Redundancy # 2 Tier and 3 Tier Architecture client C1 Application C1 [C2] C3 client louger dient, dunt 2 déent machine Application server Database Server -) dient server architecture F Bank [offline]

-> tails when no. of users are more

Database: It is a collection of related data and storage area where we store it for retrievening information.

2Tier > Scalability Problem (No of users are large)

-> Good Maintanance

- Low Security (clients directly interact with database)

3 Tiur

- query will be processed at bussiness layer

-> less load at database

-> Mose Seture

→ Hard to maintain

ef web applications

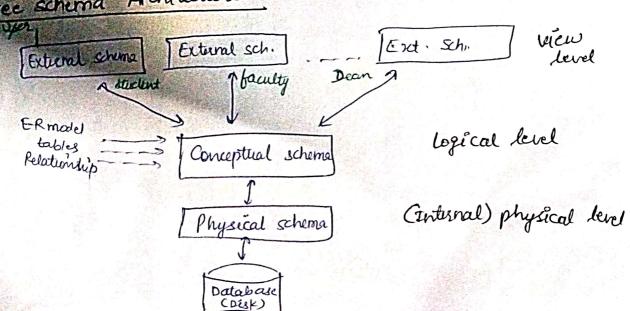
logical representation of database # schema:

W Barically the foordend

-> In RDBMS, table is a schema

> In soi, Data Definition language is used to implement the scheme

Three schema Architecture :-



(Wew)

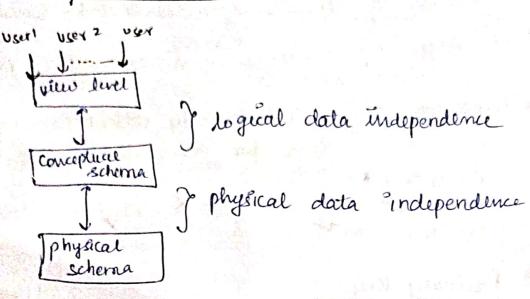
(what data is Data Base designer (logical)

Data Base Administrator (physical)

@ In Harddisk, data is stored in files

71 Data Independence Abstraction:

(No change in)
application



- The change at conceptual level (adding or removing columns etc) Will not affect the view level. This is called logical data in dependence.
- The change at physical level (ep. shifting data from hard disk 1 to 2 hard disk 2) will not affect the conceptual schema. This is called physical data independence.

y) storage structure

- 2) Pote structure change
- 3) Index

Key.
Attribute iniquely identification.

of suderit table

Aadhar Card, Roll no, Registration no., Licence no, voterid, email ID, phone no.

Collection of all the Keys is called Candidate Key Set thousing the sol of personal Key.

-> Cardidate Keys include the attributes which can never have same values for more than one entry.

PAN P Regestration no

Age X father NameX

Primary Key [Unique + NOT NULL]

ez Registration Number

-sonly one primary Key

Foreign Key:

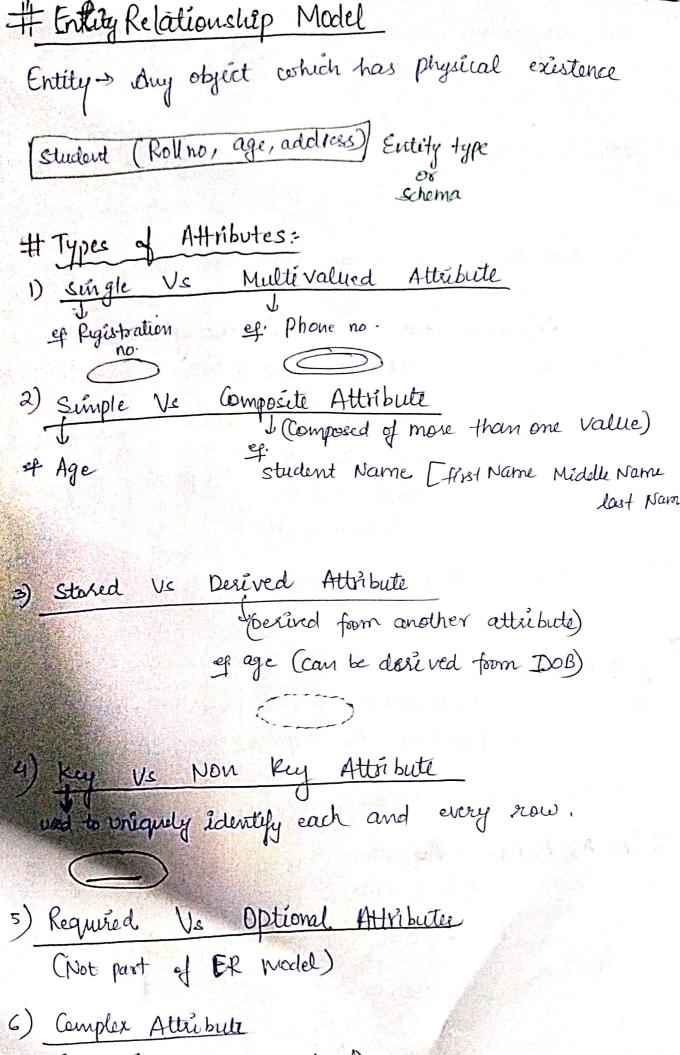
primary ky of same table or another table (relation)

-maintains referential Integrity

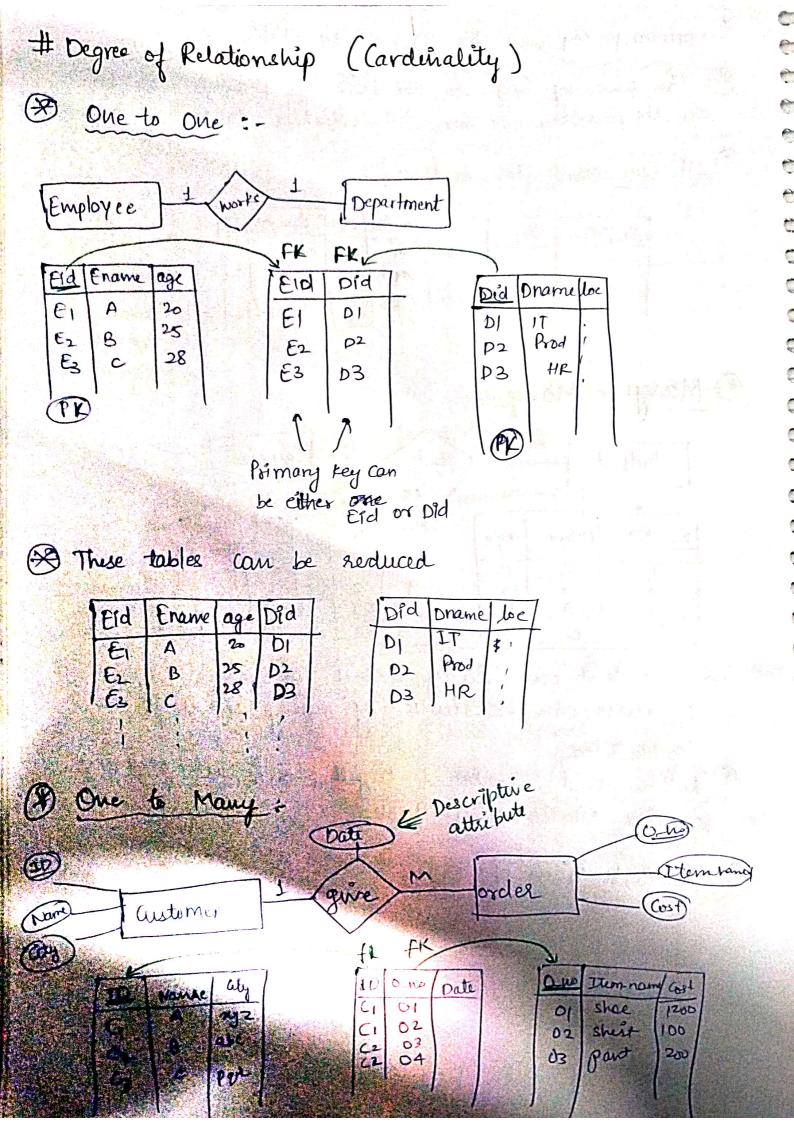
es. DX X	Pasy Referenced table		
Roll No	Name	Address	
1	A	Delhi	
3	B	Mumbai	
	0	Hapur	

	Referen	cing table	FK
•	THE RESERVE THE PERSON NAMED IN COLUMN TWO IS NOT THE OWNER, THE PERSON NAMED IN THE OWNER, THE PERSO	Course Name	Roll No.
	9	DBMS	1
	C2	Networks	2

(1) Jourse name varchas(1)
Creale table Course (Course id varchas (10), Course _name varchas(10), Rollno int references Student (Rollno));
사람이 다른 사람들이 되었다. 아이들은 사람들은 이번 사람들은 이번 사람들이 되었다. 그는 사람들이 되었다면 보다 되었다면 보다 되었다. 그는 사람들이 되었다면 보다 되었다면 보다 되었다면 보다 되었다.
-> Attertable Course ADD constraint fk foreign key (Rollno)
references Student (Rollno);
- One table can have more than one foreign key.
Aller table Course ADD Constitution of the preferences Student (Rollno); The table can have more than one foreign key. Referenced table Preferenced table
2) Delete - May cause violation 2) Delete - will not Cause I
2) Delete - May cause violation (Solution -> On delete Cas cade) -> (On delete set Null) -> (On delete no action) 2) Delete - will not Cause any violation 3) updation - May cause violation
3) Updation - May Cause Violation
Super key is a combination of all possible attributes which can uniquely identify two tuples in a table
-> Super key set of any condidate key is super key,
O: R (A) A: A A4 - An) then thow many super keys are possible? If A is candidate key =? 2 ⁿ⁻¹
If A1, A2 are Condidate Kuys = 2 2 + 2 - 2 - 2 - 2 - 2
$A_1 A_2 A_3$ $A_1 \left(\begin{array}{c} A_1 \\ A_2 \end{array} \right) = 2^{m-1}$ $A_1 \left(\begin{array}{c} A_1 \\ A_2 \end{array} \right)$



(Composite + Muttivalued)



-> primary key will be O_no in this example The primary key to the side of many to acts as the primary key for Irelationship table. (1) Tables can be reduced: (x) Many to Many Course Study student RollNo age Name 16 17 16 -> 1 student can study many courses

1 course can be studied by many students

(Rollno, cid) Combined to gether acts as prinary Key for relationship table.

(No Reduction of tables

IF Normalisation: It is a technique to remove or Reduce redundancy from a table. cid chame Fid Fram subu Sname Sname Age DBMS RAM John RAM 1 50 Ravi TAVA Bob Varun 25 Mitin DBMS John) RAM 30000 20 A meit John Row level duplicary DBMS -> Can be hardled with primary key Column level duplicary 1) Insert Anamoly 2) Deletion Anamoly 3) Updation anamoly # first Normal form: Table should not contain ceny multivaleced Attribute Coursez Course Name Rollis Wante Course Rollno Course CFP CICHT Sai 1 Sac Harry Java Harsh Java NULL! onkar C/DBMS onkaz DBMS C Name Con Pollino drawback Sai Sai

Retino Name Court

Sai C

Sai CH

2 Harsh Java

3 onkar C

3 toukar Dems

Roll no	Name
)	sai
2	flassh
3	ontor

Roll no	Course
$\begin{bmatrix} 1 \\ 4 \end{bmatrix}$	C
2	C+1 Fava
3	C
3	DBMS