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# MIMESIS AS RANDOM GRAPH COLORING

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## ABSTRACT

Inspired by the thought-provoking masterpiece by René Girard, *Le Bouc Emissaire*, a simple and tractable model for mimetic behaviour occurred to me. When we change our beliefs, we do so not because of their intrinsic value. Our desire to switch from belief  $A$  to belief  $B$  is proportional to the number of adherents of belief  $B$  that we know.

## 1 Mimesis as a decentralised process

In this article, I propose that when we change our beliefs, we do so not because of their intrinsic value. Our desire to switch from belief  $A$  to belief  $B$  is proportional to the number of adherents of belief  $B$  that we know. Technically, I modelled the problem of two conflicting beliefs that propagate through a network with  $N$  nodes in a decentralised manner. These beliefs are in some sense competing for adherents.

Using vertex notation, two individuals  $v_i$  and  $v_j$  with identical beliefs are connected with probability  $q$ , and  $1 - q$  otherwise.  $v_i$  changes its belief with a probability proportional to the number of nodes connected to  $v_i$  that have opposing views.

Two key motivating questions are:

1. Under what circumstances does a belief get completely wiped out?
2. Under what circumstances does a belief completely dominate(i.e. wipe out) all other beliefs?

In the scenario where there are only two possible beliefs these two questions are equivalent and I show that on average it's sufficient that  $q > 1 - q$  and that initially, one belief has a greater number of adherents than the other.

Nodes carrying the first belief were assigned to the set of red vertices,  $R$ , and nodes carrying the second belief were assigned to the set of blue vertices,  $B$ . After further reflection, I chose  $+1$  and  $-1$  as labels. The reason being that a change of belief using this representation would be equivalent to multiplication by  $-1$ . As a result, the  $N$  vertices could be represented by an  $N$ -dimensional vector:

$$\vec{v} \in \{-1, 1\}^N \tag{1}$$

where  $N = |v_i \in R| + |v_j \in B|$ .

## 2 A random graph model

Using this representation, between each pair of vertices we may define a virtual weight matrix  $W$ :

$$w_{ij} = v_i \cdot v_j \quad (2)$$

where  $w_{ij} = +1$  implies identical beliefs and we have  $w_{ij} = -1$  otherwise.

Now, we note that  $W$  may be conveniently decomposed as follows:

$$W = W^+ + W^- \quad (3)$$

where  $W^-$  denotes potential connections between nodes of different color and  $W^+$  denotes potential connections between nodes of identical colors.

In order to simulate variations in connectivity we may assume that nodes of the same color are connected with probability  $\frac{1}{2} < q < 1$  and nodes of different color are connected with probability  $1 - q$ . Given  $W$  we may therefore construct the adjacency matrix  $A$  by sampling random matrices:

$$M_1, M_2 \sim \mathcal{U}([0, 1])^{N \times N} \quad (4)$$

$$M^+ = 1_{[0, q)} \circ M_1 \quad (5)$$

$$M^- = 1_{(1-q, 1]} \circ M_2 \quad (6)$$

where  $1_{[0, q)}$  denotes the characteristic function over the set  $[0, q)$  and then we compute the Hadamard products:

$$A^+ = M^+ \cdot W^+ \quad (7)$$

$$A^- = M^- \cdot W^- \quad (8)$$

so the adjacency matrix is given by  $A = A^+ + A^-$ .

## 3 Stochastic dynamics for computer simulation

Now, in order to simulate stochastic dynamics we simply use majority vote:

$$p(v_i^{n+1} = v_i^n) = \frac{\bar{N}_i}{N_i} \quad (9)$$

$$p(v_i^{n+1} = -1 \cdot v_i^n) = 1 - \frac{\bar{N}_i}{N_i} \quad (10)$$

$$\bar{N}_i = |A(i, -) > 0| - 1 \quad (11)$$

$$N_i = \bar{N}_i + |A(i, -) < 0| \quad (12)$$

where  $|A(i, -) > 0| - 1$  denotes the number of connections between  $v_i$  and nodes sharing the same belief without counting a connection to itself.

## 4 Analysis

If we denote the number of red vertices at instant  $n$  by  $\alpha_n$  and the number of blue vertices by  $\beta_n$  we may observe that the expected number of neighbors is given by:

$$\langle N(v_i \in R) \rangle = q \cdot (\alpha_n - 1) + (1 - q) \cdot \beta_n \quad (13)$$

$$\langle N(v_i \in B) \rangle = q \cdot (\beta_n - 1) + (1 - q) \cdot \alpha_n \quad (14)$$

Using the above equations we may define the expected value:

$$\langle \alpha_{n+1} \rangle = \alpha_n \left( \frac{q \cdot (\alpha_n - 1)}{q \cdot (\alpha_n - 1) + (1 - q) \cdot \beta_n} \right) + \beta_n \left( \frac{(1 - q) \cdot \alpha_n}{q \cdot (\beta_n - 1) + (1 - q) \cdot \alpha_n} \right) \quad (15)$$

and we may deduce that  $\langle \beta_{n+1} \rangle = N - \langle \alpha_{n+1} \rangle$ .

**4.1**  $\alpha_n > \beta_n$  **implies that**  $\lim_{n \rightarrow \infty} \langle \alpha_n \rangle = N$

Assuming that  $q > 1 - q$ , a simple calculation shows that:

$$\langle \alpha_{n+1} \rangle - \alpha_n \geq 0 \iff \alpha_n \geq \beta_n \quad (16)$$

and since:

$$\langle \alpha_{n+1} \rangle - \alpha_n = 0 \iff \alpha_n = \beta_n \quad (17)$$

we may deduce that:

$$\lim_{n \rightarrow \infty} \langle \alpha_n \rangle = N \quad (18)$$

### 4.2 Analysis of $\Delta\alpha$

Using the fact that  $\beta_n = N - \alpha_n$  we may derive the following continuous-space variant of  $\Delta\alpha_n = \langle \alpha_{n+1} \rangle - \alpha_n$ :

$$\Delta\alpha(\alpha, \gamma) = \frac{\alpha \cdot (N - \alpha)}{\gamma \cdot (\alpha - 1) + (N - \alpha)} - \frac{\alpha \cdot (N - \alpha)}{\gamma \cdot (N - \alpha - 1) + \alpha} \quad (19)$$

where  $\gamma = \frac{q}{1-q}$ .

## References

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- [3] P. Erdős and A. Rényi. On the evolution of random graphs. 1960.
- [4] G. Grimmett and C. McDiarmid. On colouring random graphs. 1975.