

Hw 2 due today

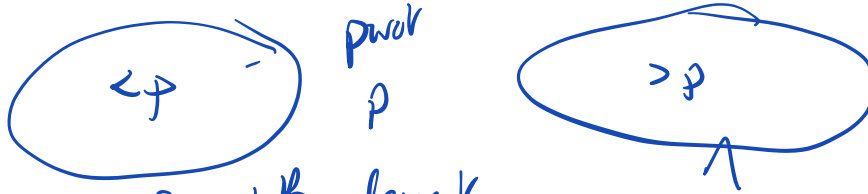
Hw 3 → Practice exam


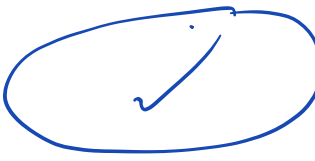
Quiz 1 next Thursday

Review Sessions Tu, W
1-3, 5-7


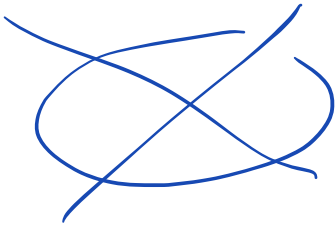
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1. Quicksort takes $O(n \log n)$ time if we have good pivots
 2. If we could find the median, we have a good (the best) pivot
 3. Median is not self-reducible, but Selection is: You can break selection into smaller problems of the same type, and median is a special case
 4. Selection is $O(n)$ if we get good pivots each time.

[Q] = [$\frac{1}{4}n$ | Good | $\frac{3}{4}n$]_n
sorted order


 Looking for k^{th} element

if $k > p$'s position
 \rightarrow  

if $k < p$'s position

$k = p$'s position, output p .

Want to eliminate $\geq \frac{1}{4}$ of remaining items at each step

because then $T(n) \leq T\left(\frac{3n}{4}\right) + O(n)$

$\Rightarrow T(n) = O(n)$

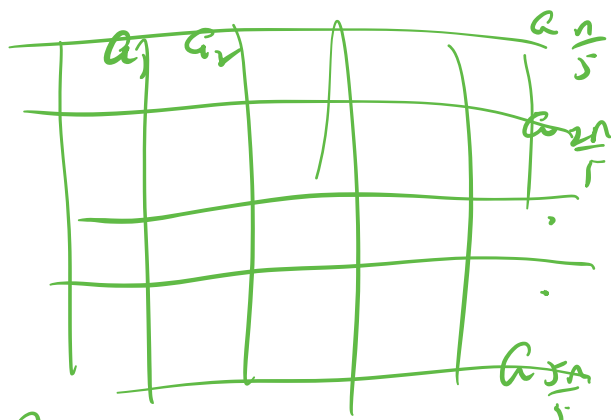
Finding a "good" pivot:

input $a_1 \dots a_n$

Selection takes $T(n)$ time

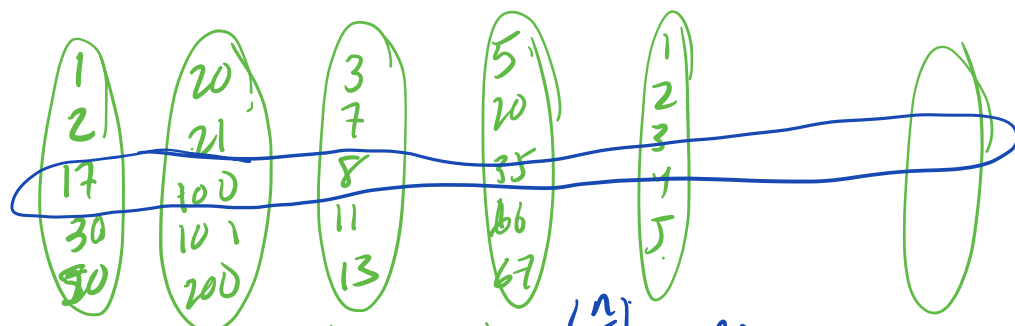
\therefore \dots in an array $O(n)$

1. Rearrange



$\frac{n}{5}$ columns

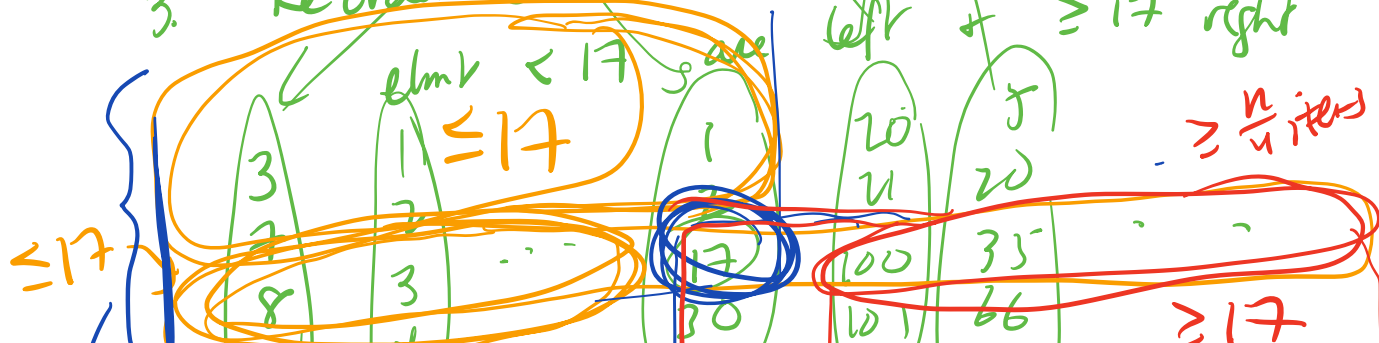
2. Sort each of the $\frac{n}{5}$ columns. $O(n)$ time

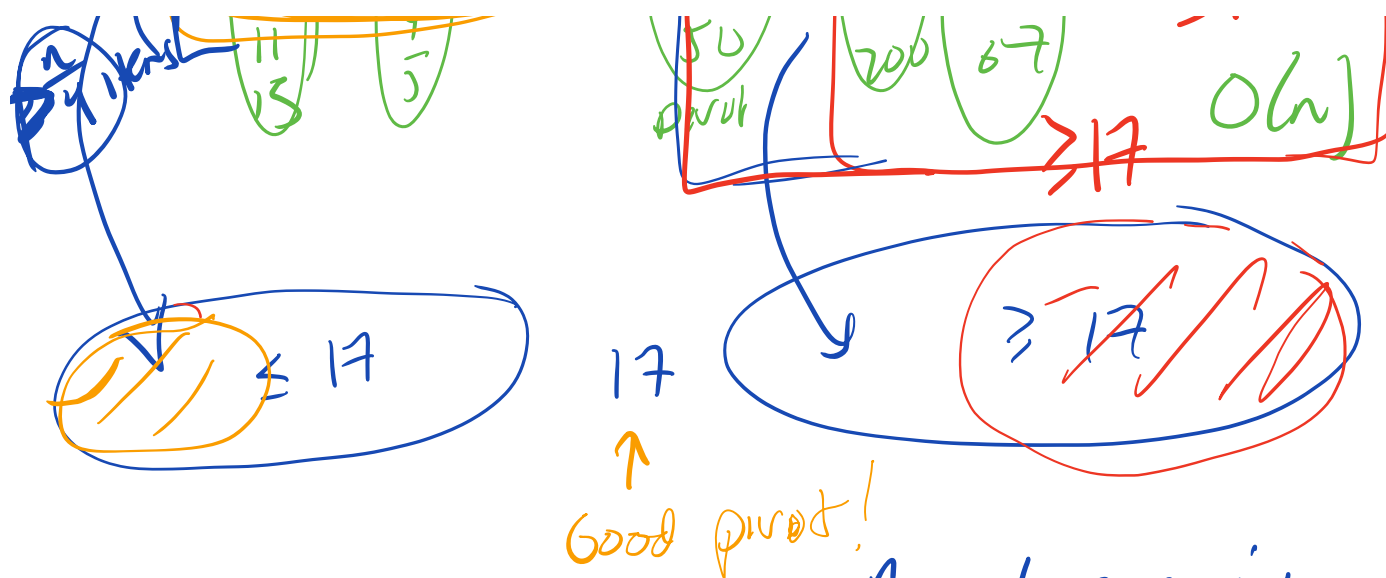


Find the median of the $\frac{n}{5}$ middle elements
selection $T(\frac{n}{5})$

Maybe 17 is media

3. Re order columns so that ones with middle left & ≥ 17 right





Now I eliminate $\geq \frac{n}{4}$ of remaining items

$$T(n) \leq T\left(\frac{n}{5}\right) + T\left(\frac{3n}{4}\right) + O(n)$$

$T(n) = O(n)$ ↑ recursive

Version of Master Theorem

$$\text{If } T(n) = T(c_1 n) + T(c_2 n) + \dots + T(c_k n) + O(n^d)$$

$$\text{And } c_1 + c_2 + \dots + c_k < 1$$

$$\text{then } T(n) = \begin{cases} O(n^d) & \text{if } d \geq 1 \\ O(\log n) & \text{if } d = 0 \end{cases}$$

Dynamic Programming

Fibonacci sequence

$$a_0 = 1$$

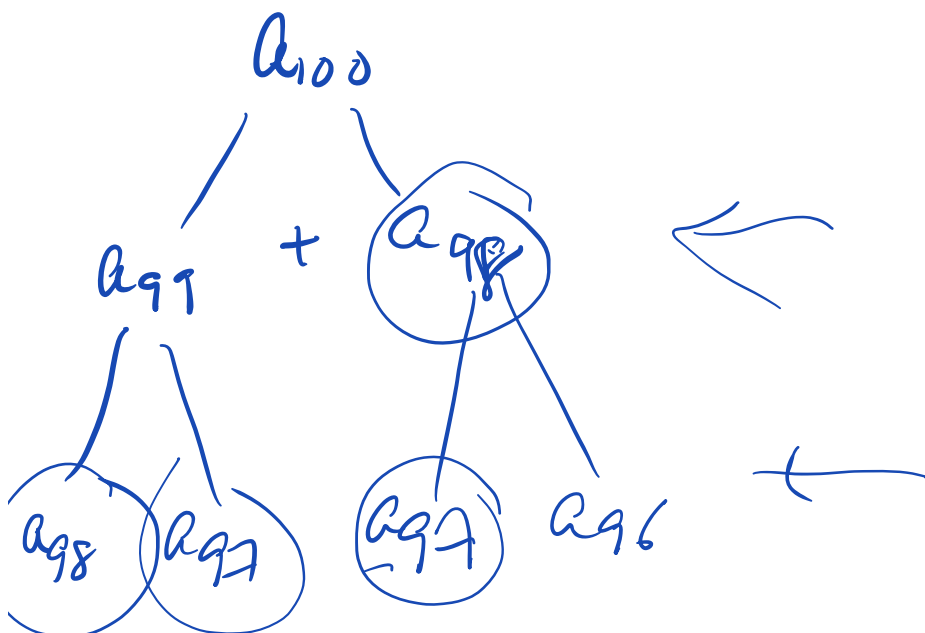
$$a_1 = 1$$

$$a_n = a_{n-1} + a_{n-2}$$

(1 1 2 3 5 8 13 21 ... 100)

What is a_{100} ?

$O(n)$
time



2^n leaves



← n levels
or double each time