

CS 3510

Goals:

Thinking abstractly

Thinking rigorously

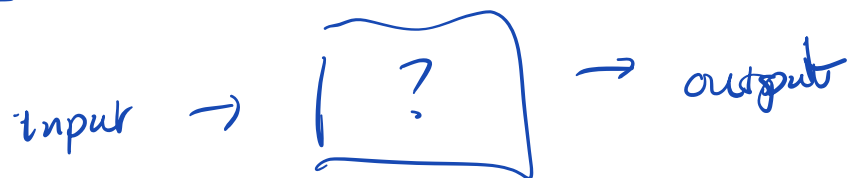
Design + analysis of algorithms

Binary search

DFS

Sorting

addition



correct?
How fast?

of operations: $\left\{ \begin{array}{l} \# \text{ (comparisons)} \\ \# \text{ bit / digit operations} \\ \# \text{ data structure operations} \end{array} \right.$

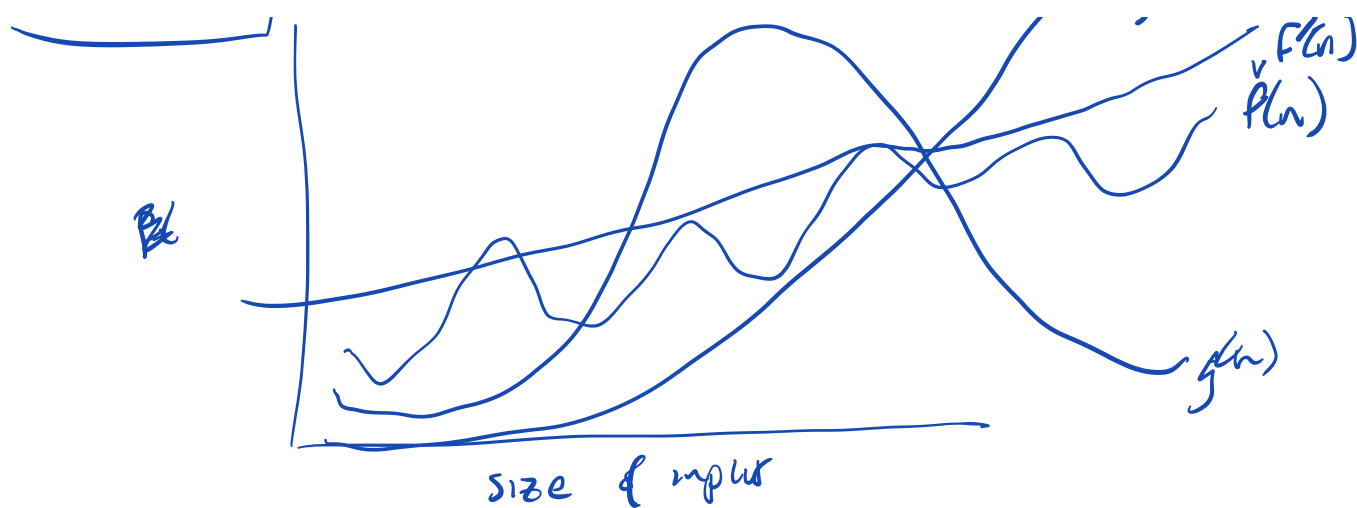
How fast
operations
comparisons
data str
operations

input

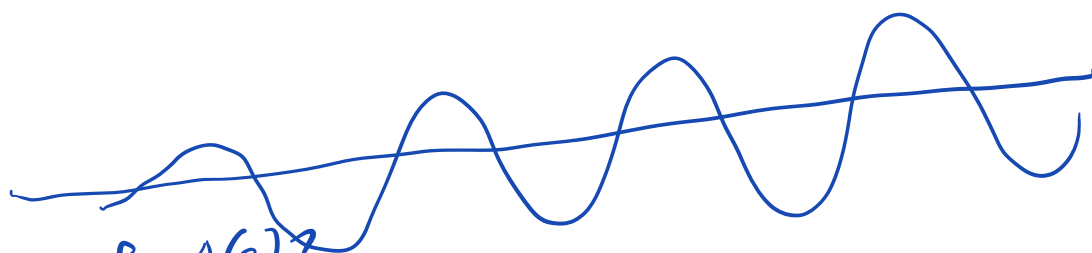
\rightarrow
 $f(n)$ steps

output

How fast
does it grow?
 $h(n)$



Def: $f(n)$ is $O(g(n))$ if
 $\exists c > 0$ s.t. $f(n) \leq c \cdot g(n) \quad \forall n \geq n_0$



Is $f = O(g)$?

Will $c \cdot g(n)$ overtake f at some point
 + from then on?

- I Divide & conquer
- II Dynamic programming
- III Graph algorithms
- IV NP-completeness
- V optional topics

I Divide + conquer.

A) Divide the problem into smaller subproblems (of the same type)

B) Solve subproblems recursively

C) Merge the solutions to solve the original problem.

Merge Sort

Given a_1, \dots, a_n ← distinct and we want to output them in order.

Operation: # comparisons
Is $a_i < a_j$?

Assume $n = 2^k$ for some $k \in \mathbb{Z}^+$

Ex:

sort

6 1 10 12 3 8 7 2

L R

Dividing

... and

7

1 6 7 8

mp \rightarrow Solve i

Merge;
 $< n$

Is $1 < 2$ yes

①

Is $6 < 2$ no

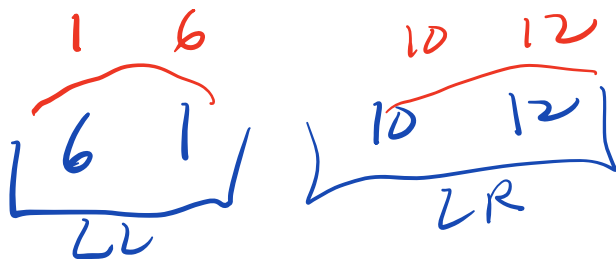
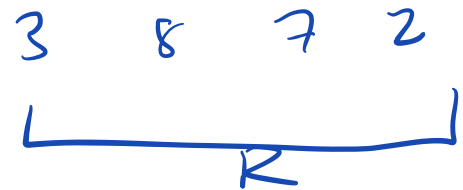
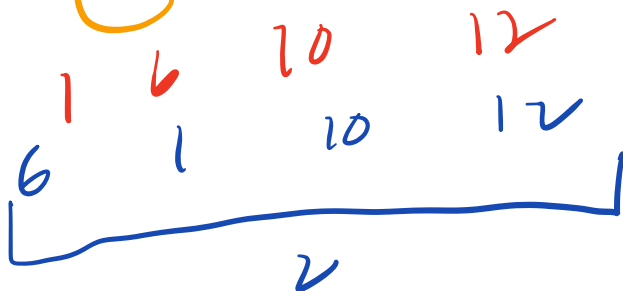
②

Is $6 < 3$ no

③

Is $6 < 7$ yes

④



6 11

10

12

n items

Time to sort n items is $T(n)$

$$T(n) < 2T\left(\frac{n}{2}\right) + n$$

$$T(n) < 2\left(2T\left(\frac{n}{4}\right) + \frac{n}{2}\right) + n$$

$$= 4T\left(\frac{n}{4}\right) + 2 \cdot \frac{n}{2} + n$$

$$= 4T\left(\frac{n}{4}\right) + 2n$$

$$T(n) = 4\left(2T\left(\frac{n}{8}\right) + \frac{n}{4}\right) + 2n$$

$$= 8T\left(\frac{n}{8}\right) + \frac{4n}{4} + 2n$$

$$= 8T\left(\frac{n}{8}\right) + 3n$$

$$T(n) \approx 16 T\left(\frac{n}{16}\right) + 4n$$

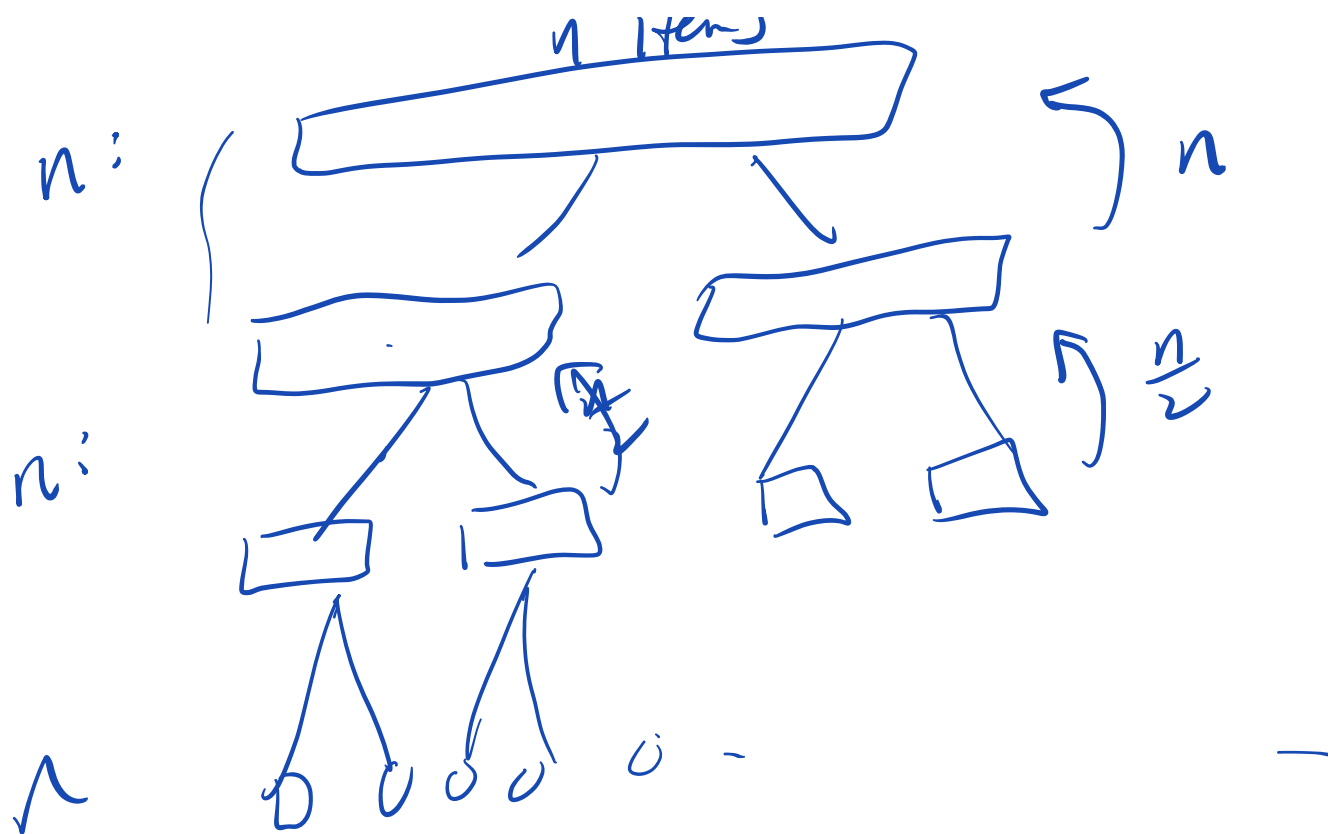
Guess

$$T(n) = 2^k T\left(\frac{n}{2^k}\right) + kn$$

$$\text{If } T(n) \leq 2^k T\left(\frac{n}{2^k}\right) + kn$$

$$\text{then } T(n) \leq 2^{k+1} T\left(\frac{n}{2^{k+1}}\right) + (k+1)n$$

$$T(1) = 0$$



→ / 0 0 0 0 0 0 0 0

Total # comp $\leq k \cdot n$

$$= \boxed{\log_2 n \cdot n}$$