

COMP9033 **DATA ANALYTICS** 

11/12

**BATCH DATA ANALYSIS** 

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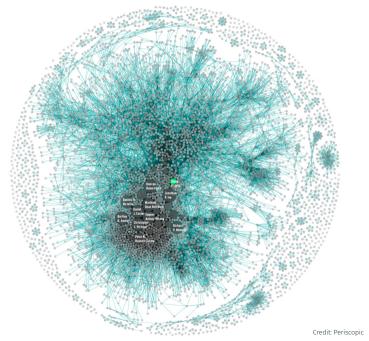
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### 0.1 / LAST WEEK

## 1. Sets and probability:

- · Set notation.
- · Set cardinality.
- · Probability.
- · Conditional probability.
- · Examples.

## 2. Association rule mining:

- · What it is.
- Terminology.
- · Support, confidence and lift.
- · Brute force mining.
- Efficient rule mining.

### 0.2 / THIS WEEK

- 1. Big data:
  - · The three Vs definition.
  - · Why it's relevant now.
- 2. High volume data:
  - · Why it's a problem.
  - · Cluster computing.
  - · Distributed file systems.

- 3. MapReduce:
  - · Cluster architecture.
  - · How it works.
  - · Examples.
- 4. Batch data frameworks:
  - · Mahout.
  - · TensorFlow.
  - · Spark.



## 1.1 / BIG DATA

- Over the past few weeks, we have covered a variety of topics in the area of data analysis, but we have yet to consider big data analysis. Why?
- · Big data is specialised:
  - Data analysis is a set of *general* tools for solving *general* problems.
  - Big data analysis is a set of specialised tools for solving specific problems.
  - Big data analysis problems are a subset of general data analysis problems we need to learn to walk before we can run!
- · Big data is hard:
  - Big data is a complex topic and comes with an additional layer of problems that must be solved.
  - Many "small data" solutions don't translate (e.g. due to performance).
  - · Generally, we have to adjust, adapt or re-architect our approach.
  - In some cases, we must resort to heuristic (i.e. suboptimal) approaches.

## 1.2 / WHAT IS BIG DATA?

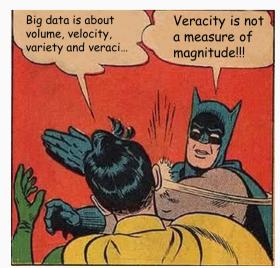
- Big data problems are encountered when analysing data at large scales.
- It is broadly (though not universally) agreed that there are three varieties:
  - 1. Volume.
  - 2. Velocity.
  - 3. Variety.
- *High volume* describes situations where the quantity of data to be analysed is so large that conventional storage and/or processing techniques cannot be used, *e.g.* due to hard disk, memory or CPU limitations.
- There is no specific threshold for the amount of data that constitutes a high volume problem, *e.g.* what we consider high volume today may not be in ten years due to Moore's Law.

## 1.3 / WHAT IS BIG DATA?

- High velocity describes situations where the rate at which data is produced is so fast that conventional processing systems cannot analyse the data in a reasonable amount of time.
- Increasingly, data is produced at faster rates and from more sources (e.g. stock markets, mobile devices, IoT), but we still require results quickly.
- High variety describes situations where there are a vast number of forms of data to be analysed: structured (e.g. SQL), semi-structured (e.g. CSV, JSON, XML) and, increasingly, unstructured (e.g. tweets, images, audio).
- In order to cope, we must build more robust systems, capable of extracting information from a plethora of data sources.

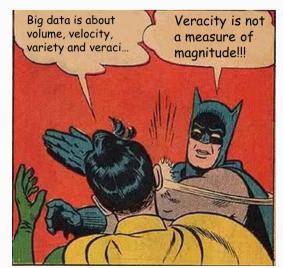
# 1.4 / WHAT IS BIG DATA?

- Volume, velocity and variety are known as the three Vs, and are the most commonly accepted definitions of big data.
- Other properties have also been proposed, but are controversial and not widely accepted.
- Generally, this is because the properties apply to both small data and big data problems, i.e. they do not become problems due to data scale alone.



## 1.5 / WHAT IS BIG DATA?

- One example is veracity, a property that defines the uncertainty or ambiguity of data.
- However, imprecise data is a problem at all scales of data magnitude, e.g.
  - Typos during manual data entry.
  - · Faulty sensors.
  - · Intentional misreporting.



## 1.6 / WHY IS BIG DATA A PROBLEM NOW?

- The rate at which data is being produced is accelerating:
  - Data producing devices (e.g. mobile, wearable, IoT) are growing in number.
  - · Monitoring tools and storage are becoming increasingly more affordable.
  - Data analysis tools are becoming more mature, incentivising the capture of more metrics and key performance indicators.
- It's estimated that:
  - 2.5 exabytes of data are produced every day (1 exabyte = 1 billion gigabytes!).
  - · About 75% of this data is unstructured.
  - The current growth in data *already* exceeds Moore's Law, *i.e.* conventional technology will become less effective at dealing with big data in the future.
  - The Internet of Things (IoT) will bring somewhere between thirty and two hundred billion devices online in the coming decade.
- Big data is a problem now and will be a problem in future!



### 2.1 / HIGH VOLUME DATA

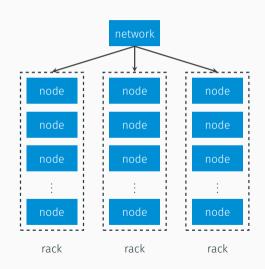
- High volume data problems occur when the amount of data becomes so large that conventional storage and processing systems begin to fail:
  - Data becomes too large to hold in memory.
  - Data becomes too large to store on disk.
  - The CPU time required to compute even trivial operations becomes so large that it is not feasible to execute tasks in a reasonable amount of time.
- There are two solutions to this problem:
  - 1. Build more powerful computers, which are capable of handling the additional demand.
  - 2. Use cluster computing to distribute the workload over many standard/commodity computers.

## 2.2 / HIGH VOLUME DATA

- Until very recently, building bigger computers was a popular solution:
  - · Relatively few problems required massive amounts of resources.
  - The high cost of purchasing a supercomputer was generally offset by the benefits of solving the problem.
- However, the rate at which data is currently being produced exceeds Moore's Law, which means that
  - · An increasing number of problems require massive compute power to solve.
  - The cost of sufficiently powerful supercomputers is increasing.
  - · The useful lifespan of supercomputers is decreasing.
- Clusters of compute nodes built from commodity hardware are generally less expensive than supercomputers and can be expanded in an ad hoc manner.
- However, conventional solutions (e.g. data storage, algorithms) must be adapted before they can operate on a clustered architecture.

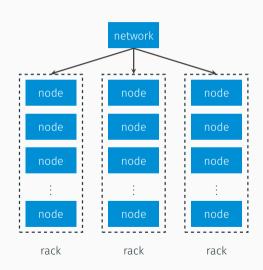
#### 2.3 / CLUSTER COMPUTING

- A cluster computing environment usually consists of three distinct components:
  - 1. Nodes.
  - 2. Racks.
  - 3. Networks.
- Nodes are simply commodity computers, with standard compute, memory and storage capabilities.



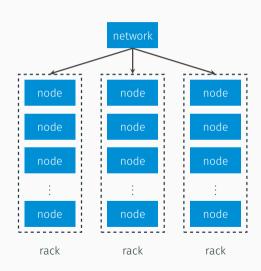
#### 2.4 / CLUSTER COMPUTING

- Racks are groups of nodes connected together by a local network.
- Groups of racks are then connected together by a wider area network, which may in turn connect to one or more other wider area networks.
- Typically, the network bandwidth across racks is smaller than the network bandwidth inside racks.



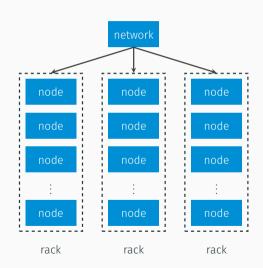
### 2.5 / CLUSTER COMPUTING

- Failure can be catastrophic in a clustered environment:
  - If a node fails, then its data is lost and any computation it was involved with may have to be aborted.
  - If a rack fails, then effectively all its nodes fail.
  - If a network fails, then effectively all its racks fail.



#### 2.6 / CLUSTER COMPUTING

- In practice, nodes, racks and networks can (and do) fail, so we must design with failure in mind.
- Generally, the problem is tackled in a two stage manner:
  - 1. Distributed file systems are used to mitigate against data loss.
  - 2. Distributed compute frameworks are used to mitigate against computation disruption.



#### 2.7 / DISTRIBUTED FILE SYSTEMS

- To mitigate against data loss due to node, rack or network failure, data is typically stored on a *distributed file system* (DFS).
- Many varieties of distributed file system have been proposed and/or implemented, e.g.
  - · Amazon S3.
  - · Ceph.
  - GlusterFS.
  - Google File System (GFS).

- Hadoop distributed file system (HDFS).
- Microsoft DFS.

• In general, different implementations have different design goal priorities, e.g. fault tolerance, data parallelism, read/write throughput, scalability.

#### 2.8 / THE HADOOP DISTRIBUTED FILE SYSTEM

- The *Hadoop distributed file system* (HDFS) is often used in batch data analysis applications.
- This is because the design of HDFS makes a number of assumptions about the kind of data being stored in the cluster, *i.e.* 
  - 1. Files are so large that it is not feasible to store them on a single node.
  - 2. Files are written once and subsequently read many times.
  - 3. When files are read, they are read in whole.
- These assumptions map well to many high volume data problems, i.e.
  - There is a massive quantity of data to be analysed, usually more than can be stored on a single computer.
  - Data is typically accumulated over time, and so is often stored in an append-only fashion.
  - Usually, the entire data history is analysed, *i.e.* data is processed in batch.



### 3.1 / MAPREDUCE

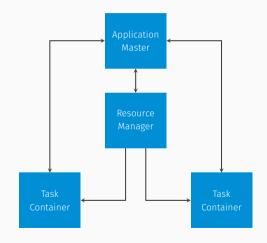
- MapReduce is a distributed computation paradigm for processing large, infrequently updated quantities of data on clustered computers:
  - · Large computations are broken down into small, isolated tasks.
  - These tasks are assigned to available worker nodes.
  - The outputs of the individual tasks are merged to form the final result.
- Google are credited with inventing MapReduce, and have been granted a patent on it, but several implementations are available.
- The most well known freely available implementation is Hadoop MapReduce, which operates on HDFS.
- While Hadoop MapReduce is written in Java, MapReduce as a concept is not language specific.

### 3.2 / MAPREDUCE

- MapReduce is used to solve a variety of high volume data problems in many different industries<sup>1</sup>, e.g.
  - · Amazon Elastic MapReduce.
  - · Comcast (network management).
  - · EBay (search optimisation).
  - · Facebook, LinkedIn, Spotify (recommendations/analytics).
  - · Google PageRank.
  - Hulu, Etsy, Rackspace (server log analysis).
  - · Salesforce (customer behaviour prediction).
  - · Telefonica (user modelling).
  - · Twitter (processing tweets and logs).
  - · Yahoo (research).

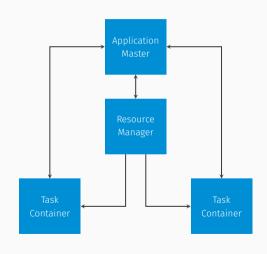
#### 3.3 / HADOOP MAPREDUCE ARCHITECTURE

- In Hadoop 2, MapReduce jobs are coordinated by a Resource Manager:
  - Schedules jobs according to available CPU, memory and disk space on worker nodes.
  - Knows where nodes are (i.e. rack aware).
- Typically, the Resource Manager is colocated with the cluster's HDFS NameNode on the cluster master.



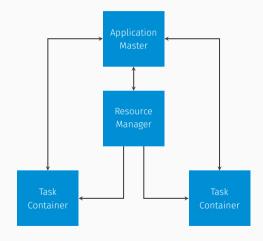
#### 3.4 / HADOOP MAPREDUCE ARCHITECTURE

- When a MapReduce job is run, the Resource Manager spawns an Application Master, which manages the state of the MapReduce job throughout its lifespan.
- The Application Master is responsible for negotiating compute resources from the Resource Manager that are "local" to the data required for the job.
- Usually the Application Master is located on a worker node, similar to the HDFS DataNode.



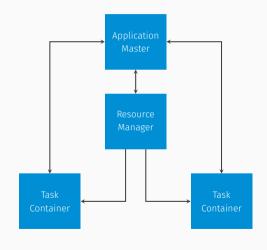
#### 3.5 / HADOOP MAPREDUCE ARCHITECTURE

- When the Application Master requests access to compute resources from the Resource Manager, the Resource Manager allocates a number of task containers, i.e. bounded compute resources on one or more worker nodes.
- The Application Master then uses the task containers to execute the map (and shuffle) and reduce tasks.



### 3.6 / HADOOP MAPREDUCE ARCHITECTURE

- If a task fails, the Application Master will attempt to rerun it on another node.
- If the Application Master fails, the Resource Manager will restart the MapReduce job.
- Depending on the cluster configuration, it may be possible to recover the outputs of previously completed tasks, avoiding the re-execution of all tasks.
- If the Resource Manager fails, then the MapReduce cluster will fail (although high availability is configurable).



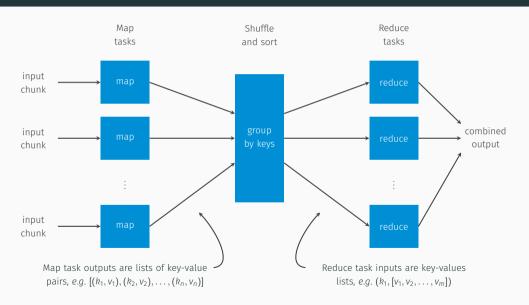
## 3.7 / HOW MAPREDUCE WORKS

- MapReduce works by decomposing computations into two distinct phases:
  - A mapping phase, where chunks of the input data are processed independently of one another (*i.e.* in parallel).
  - A reduction phase, where the results of the mapping phase are combined to formulate the final result (can also be run in parallel).
- The operations performed in the mapping and reduction phases are known as *map* and *reduce* tasks, respectively.
- Both map and reduce tasks *must* be stateless functions, so that if a task fails on one compute node, it can readily be reassigned to another without aborting the entire computation.
- Decomposing computations into map and reduce tasks can be a complicated procedure — MapReduce algorithms often don't resemble their "small" data equivalents.

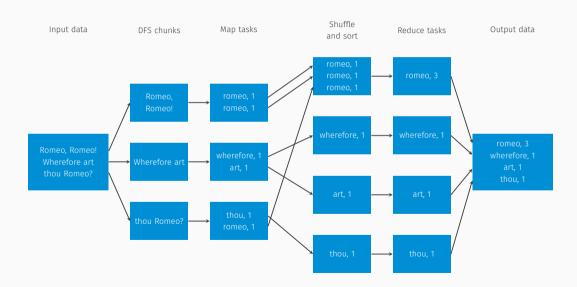
#### 3.8 / HOW MAPREDUCE WORKS

- A typical MapReduce job works as follows:
  - 1. A batch data processing job is decomposed into an operation consisting of map and reduce tasks *only*.
  - 2. The locations of the data chunks to be processed are identified (usually the data is stored on a DFS).
  - 3. Map tasks are assigned for each data chunk, and are then scheduled to be executed by a number of compute nodes "close" to those chunks (to reduce network I/O).
  - 4. The compute nodes execute the map tasks, each one processing its assigned chunk of data and converting it into a list of key-value pairs, according to some user-defined operation.
  - 5. The key-value pairs resulting from all of the map task executions are grouped by their keys.
  - 6. One or more reduce tasks are assigned to process the grouped key-values lists.
  - 7. The reduce tasks are assigned to a number of compute nodes, each one processing all of the values associated with a given key, according to some user-defined operation.
  - 8. The outputs of all of the reduce tasks are combined to produce the final result.

# 3.9 / HOW MAPREDUCE WORKS



# 3.10 / EXAMPLE: WORD COUNTING



## 3.11 / CLUSTER SIZE

- For maximum parallelism, we could build a cluster with enough nodes, so that all of the reduce tasks run simultaneously.
- · However, this is rarely done in practice.
- The number of unique keys produced by a set of mapping operations (and, therefore, the number of reduce tasks required) is typically very large, and so it would be costly to build a cluster with enough capacity to run them all in parallel.
- Also, the time required to complete a reduce operation can depend on its inputs, and so some reduce tasks will finish faster than others — it would be an inefficient use of resources to have a large number of inactive nodes.

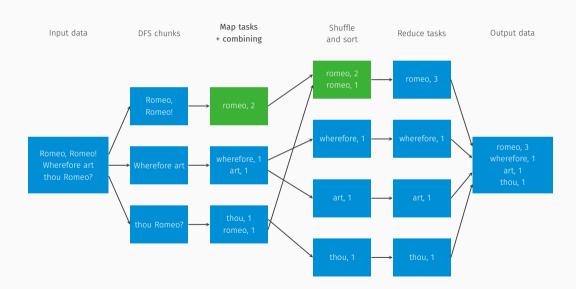
### 3.12 / CLUSTER PERFORMANCE

- Generally, map and reduce tasks are designed so that they do not consume large amounts of resources:
  - Usually, nodes are commodity hardware and so are not designed to execute complex tasks quickly.
  - Instead, data and task parallelism take care of computational complexity.
  - If individual tasks to take a long time to complete, node failures may significantly increase the overall compute time.
- In practice, the biggest bottleneck in the system is the limited bandwidth between compute nodes:
  - If map tasks produce large numbers of key-value pairs, we are forced to transfer large amounts of data across the DFS in order for the reduce tasks to produce the result.
  - The time required to transfer the data can delay the computation significantly.

## 3.13 / COMBINING

- One technique that mitigate against this bottleneck is *combining*:
  - Move some reduction logic to the mapping phase to reduce the amount of data produced by each map task.
  - If each map task produces less data, then we don't have to transfer as many key-value pairs across the DFS.
- · However, we can't combine map and reduce tasks in every situation:
  - Combination only works where the reduce operation is both associative and commutative, i.e. values can be processed in any order and the result will always be the same.
  - Associativity and commutativity are not general properties of computations, and so we can only use combination to decrease execution time in a limited subset of problems.

## 3.14 / EXAMPLE: WORD COUNTING



Batch data frameworks

# 4.1 / BATCH DATA FRAMEWORKS

- In the first release of Hadoop, the MapReduce layer operated directly on HDFS (see opposite).
- This architecture imposed a number of constraints on the kinds of applications that could be developed:
  - As Hadoop MadReduce was a Java application, all child applications had to be written in Java.
  - The Hadoop MapReduce API was complex, which increased both code complexity and application development time.

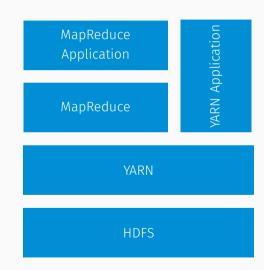
MapReduce Application

MapReduce

HDFS

### 4.2 / BATCH DATA FRAMEWORKS

- This architecture was significantly redesigned in Hadoop 2.
- The Yet Another Resource Negotiator
   (YARN) resource manager was added as
   an abstraction layer between HDFS and
   HDFS applications.
- YARN enables applications to be built directly on top of HDFS, while also allowing applications to be built on top of MapReduce.



## 4.3 / BATCH DATA FRAMEWORKS

- In recent years, there has been a trend towards using frameworks which further abstract these layers, making it simpler and faster to develop new applications, *e.g.* 
  - · Apache Flink.
  - Apache Giraph.
  - · Apache Mahout.
  - · Apache PredictionIO.

- · Apache Spark.
- Apache SystemML.
- · H20.
- · TensorFlow.

## 4.4 / APACHE MAHOUT

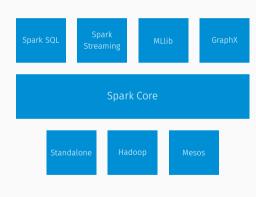
- Apache Mahout is an open source library of scalable, distributed machine learning algorithms.
- · It supports a number of different techniques, e.g.
  - · Classification: Naive Bayes.
  - · Clustering: K means.
  - Collaborative filtering: user-based and item-based.
  - · Dimensionality reduction: PCA.
- Initially, Mahout implemented its algorithms directly on top of MapReduce.
- However, in 2015, there was a shift in focus towards providing scalable algorithms in general, not just for MapReduce.
- Mahout has since deprecated most of its MapReduce algorithms in favour of Flink/H20/Spark implementations.

# 4.5 / TENSORFLOW

- TensorFlow is an open source library of machine learning algorithms, originally developed by Google.
- TensorFlow supports many machine learning algorithms:
  - Regression: linear regression, logistic regression, neural networks.
  - · Classification: neural networks.
  - · Clustering: K means.
- Like Spark, TensorFlow supports a number of different languages, including C++, Go, Java and Python.
- It also supports distributed computation using Hadoop.

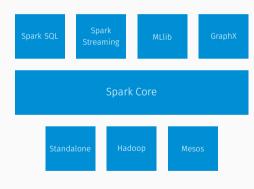
## 4.6 / APACHE SPARK

- Apache Spark is an open source software framework for processing both batch and streaming data in a compute cluster.
- Spark can be run in a number of different cluster configurations:
  - 1. Standalone, using Spark's built-in cluster manager.
  - 2. Hadoop, via YARN.
  - 3. Apache Mesos.



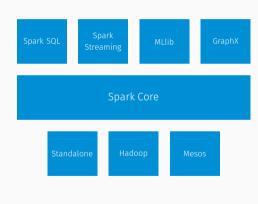
#### 4.7 / APACHE SPARK

- Spark's functionality is provided by four libraries:
  - 1. Spark SQL: Query data using SQL in Spark applications.
  - 2. Spark Streaming: Data stream processing support.
  - 3. MLlib: Distributed machine learning algorithms for batch and streamed data.
  - 4. GraphX: Graph data analysis.



## 4.8 / APACHE SPARK

- The Spark API is available in several languages:
  - · Java.
  - · Python.
  - R.
  - · Scala.
- Spark applications can be developed with code from one or more of these languages, enabling rapid development of experimental and enterprise software, side by side.



## 4.9 / APACHE SPARK

- Because the Spark API is available in many languages, we aren't limited to writing MapReduce jobs in Java (like in Hadoop).
- For instance, the code to the right implements a MapReduce job for counting words in a file on a HDFS in Python.
- For context, the equivalent code in Java is 60 lines without line breaks and comments!

```
import pyspark

# Create a Spark context object
sc = pyspark.SparkContext(master='local[*]')

# Load a text file from HDFS
text_file = sc.textFile('hdfs://...')

# Execute a MapReduce job
text_file.flatMap(lambda line: line.split()). \
    map(lambda word: (word, 1)). \
    reduceByKey(lambda x, y: x + y)
```

# 4.10 / APACHE SPARK

- Spark offers significant performance benefits over using MapReduce directly:
  - Where possible, Spark holds compute task results in-memory to speed up the next stage of calculation.
  - In contrast, MapReduce always reads from HDFS before running a task and writes to HDFS at the completion of a task.
- If no memory is available to store the results of a particular task, Spark defaults to writing them to disk, and so in a worst case scenario it is still roughly comparable to MapReduce.
- In 2014, Spark won the Daytona GraySort Benchmark Contest, beating MapReduce's previous winning score:
  - Hadoop sorted 102.5 TB of data in 72 minutes, using 2100 nodes.
  - Spark sorted 100 TB of data in 23 minutes, using just 206 nodes.



### X.1 / SUMMARY

- Big data is a big problem, but only in some cases:
  - We still need small data analytics for lots of problems.
  - · Big data is much harder, need to learn the basics then adapt.
- Using DFS + distributed compute cluster is a good solution for high volume data processing:
  - MapReduce is useful, but there can be a steep learning curve.
  - However, newer frameworks are making it easier to build complex distributed applications.
  - · Need to pick the right tool for the job!
- · Lab work:
  - · Create a standalone Spark cluster.
  - · Run a MapReduce job in Spark to analyse some data.
- · Next week: streaming data analysis!

#### X.2 / REFERENCES

- 1. Ullman et al. *Mining of Massive Data Sets.* Cambridge University Press, 2014. (stanford.io/1qtgAYh)
- 2. Landset et al. A survey of open source tools for machine learning with big data in the Hadoop ecosystem. Journal of Big Data, 2(1), p.24. (bit.ly/2oHv3RG)