

Monoidal Structures, Operads and Stable Categories II

Last time: An ω -operad consists of an ω -category \mathcal{O}^\otimes and a functor $p: \mathcal{O}^\otimes \rightarrow \text{Fin}_*$, st.

(i) Every inert map $\alpha: p(\underline{X}) \rightarrow (n)$ in Fin_* has a p -coCart. lift $\bar{\alpha}: \underline{X} \rightarrow \underline{Y}$

(ii) The lifts of g_i exhibit $\mathcal{O}^\otimes \xrightarrow[\cong]{(g_i)} (\mathcal{O}^\otimes)^{\times n} =: \mathcal{O}^{\times n}$

(iii) $\text{Map}^x([X_1, \dots, X_n], [Y_1, \dots, Y_m]) \cong \prod_{1 \leq j \leq m} \text{Map}^{g_i^{\otimes x}}([X_1, \dots, X_n], Y_j)$

Ex/ • Every symmetric colored operad

• "Same thing" as Kan-enriched operads

• Symm. monoidal ω -cats, since
as $e_{(n)} = (e_{(1)})^{\times n}$ via g_i .

• \mathbb{E}_∞ commutative operad $\text{Id}: \text{Fin}_* \rightarrow \text{Fin}_*$ with one color a and

$$\text{Mul}_{\mathbb{E}_\infty}(\overbrace{a, \dots, a}^n; a) := \text{Map}_{\mathbb{E}_\infty}^e([a, \dots, a], [a]) = \{t\}$$

with $e: (n) \rightarrow (1)$, $1, \dots, n \mapsto 1$

• \mathbb{E}_n associative operad with one color a and

$$\text{Mul}_{\mathbb{E}_n}(a, \dots, a; a) = \{\text{total orders on } \{1, \dots, n\}\} = \Sigma_n$$

• LM^\otimes with two colors a, l and

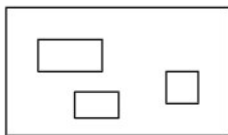
$$\text{Mul}_{\text{LM}}(a, \dots, a; a) = \Sigma_n = \text{Mul}(a, \dots, a, l; l) \text{ and rest empty}$$

• $\text{LM}_\text{pt}^\otimes$ with $\text{Mul}(a, \dots, a; l) = \Sigma_n$ as well no pointing for $n=0$.

• \mathbb{E}_k with a single color a , and

$$\text{Mul}(a, \dots, a; a) = \prod_{\infty} \text{RectEmb}(\square^k \times \{1, \dots, n\}, \square^k)$$

where RectEmb is the topol. space of rectilinear embeddings



$\text{RectEmb}(\square^k \times \{1, \dots, n\}, \square^k)$ is a closed subset of the real

vs $\mathbb{R}^{n \cdot 2^k} = \mathbb{R}\langle a_e^{(i)}, b_e^{(i)} \rangle$, which we equip with the std. topology

$$(x_1, \dots, x_n) \in \square^k \times \{i\} \mapsto (a_i x_1 + b_i^{(1)}, \dots, a_i x_k + b_i^{(k)})$$

Composition via iterative embeddings:

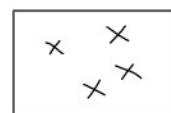
Unit given by the empty embedding

$$\emptyset \rightarrow \square^k$$

Alternatively, $\text{RectEmb}(\square^k \times \{1, \dots, n\}, \square^k) \simeq$



$$\simeq \text{Emb}^{\text{tr}}(\square^k \times \{1, \dots, n\}, \square^k) \simeq \text{Conf}_n(\mathbb{R}^k)$$



Rem: • $\mathbb{E}_n^\otimes = (\overbrace{\hookrightarrow \hookrightarrow \hookrightarrow}^n)$ is the associative operad

• $\mathbb{E}_n^\otimes \rightarrow \mathbb{E}_{n+1}^\otimes$ via \hookrightarrow 

Rem: $\mathbb{E}_n^\otimes = (\overbrace{\leftarrow \leftarrow \leftarrow}^{\text{map of operads no later}})$ is the associative operad

• $\mathbb{E}_n^\otimes \rightarrow \mathbb{E}_{n+1}^\otimes$ via $\leftarrow \leftarrow \leftarrow \rightarrow \rightarrow \rightarrow$

• $\mathbb{E}_\infty^\otimes \cong \text{colim}_n \mathbb{E}_n^\otimes$

Can define \mathbb{E}_M for M any topol. mfd., st. $\mathbb{E}_{\mathbb{R}^n} = \mathbb{E}_n$ no Factorization algebras!

Idea: \mathbb{E}_n^\otimes describes the algebraic structure on the n -th homotopy group.

e.g. in \mathbb{E}_2^\otimes one can commute , leading to braiding.

Algebras and Monoidal Structures

$\mathbb{E}_X / \mathbb{E}_1^\otimes \hookrightarrow \text{LM}^\otimes$
 $\mathbb{E}_1^\otimes \hookrightarrow \text{LM}_{pt}^\otimes$ via α

Def A map of ω -operads $f: \mathcal{O}^\otimes \rightarrow \mathcal{P}^\otimes$ is a functor

$\mathcal{O}^\otimes \xrightarrow{f} \mathcal{P}^\otimes$ that sends p -coCart. lifts of inert morph. to $!$.

$\mathcal{P} \xrightarrow{q} \text{Fin}_*$ Obtain $\text{Alg}_{\mathcal{O}}(\mathcal{P}) \subseteq \text{Fun}_{/\text{Fin}_*}(\mathcal{O}^\otimes, \mathcal{P}^\otimes) = \text{Fun}(\mathcal{O}^\otimes, \mathcal{P}^\otimes) \times \{p\}_{\text{Fun}(\mathcal{O}^\otimes, \text{Fin}_*)}$

Def Given symm monoidal ω -categories $\mathcal{C}^\otimes, \mathcal{D}^\otimes \rightarrow \text{Fin}_*$, a

\rightarrow lax monoidal functor is a map of operads $\mathcal{C}^\otimes \rightarrow \mathcal{D}^\otimes$

\rightarrow monoidal functor is a functor $\mathcal{C}^\otimes \rightarrow \mathcal{D}^\otimes$ sending

all p -coCart. morphisms to q -coCart. morph.

$\longleftrightarrow \text{Fin}_* \begin{array}{c} \text{C} \\ \Downarrow \\ \text{D} \end{array} \xrightarrow{\text{Cata}} \text{Cata}$
 Groth.constr.

(Rem: Factorization algebras are something "in between".)

Def Given an ω -operad \mathcal{O}^\otimes and an ω -category \mathcal{C} with products, an \mathcal{O} -monoid

in \mathcal{C} is a functor $M: \mathcal{O}^\otimes \rightarrow \mathcal{C}$ such that $\forall \underline{X} = [X_1, \dots, X_n] \in \mathcal{O}^\otimes$, the

(\otimes); exhibit $M(\underline{X}) \cong \prod_{i=1}^n M(X_i)$

Thm $\text{Mon}_{\mathcal{O}}(\mathcal{C}) \cong \text{Alg}_{\mathcal{O}}(\mathcal{C})$ where \mathcal{C} is equipped with its Cartesian monoidal str. \mathcal{C}^\times

In particular, for $\mathcal{C} = \text{Cat}_\infty$, we define \mathcal{O} -monoidal ω -categories

$\text{Alg}_{\mathcal{O}}(\text{Cat}_\infty) \cong \text{Mon}_{\mathcal{O}}(\text{Cat}_\infty) \cong \left\{ \begin{array}{l} \text{operibrations } \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes \text{ exhibiting} \\ \mathcal{C}_\circ \cong \prod_{i=1}^n \mathcal{C}_{\circ_i} \end{array} \right\}$

Ex \rightarrow Symm. mon ω -categories $\mathcal{C}^\otimes \rightarrow \text{Fin}_*$ are commutative algebras in Cat_∞^\times .

\rightarrow Monoidal ω -categories should be $\text{Alg}_{\mathbb{E}_1^\otimes}(\text{Cat}_\infty) \cong \left\{ \begin{array}{l} \mathcal{C}^\otimes \\ \downarrow_{\mathbb{E}_1^\otimes} \end{array} \text{ with } ! \right\}$

This agrees with our definition since $\Delta^{\text{op}} \rightarrow \mathbb{E}_1^\otimes$ is an "approximation"

(alternatively, can build a theory of non-symmetric ω -operads as

functors $\mathcal{O}^\otimes \rightarrow \Delta^{\text{op}}$ satisfying similar axioms, see [Gepner-Huang])

\rightarrow An LM^\otimes -monoid $\text{LM}^\otimes \rightarrow \text{Cat}_\infty$ consists of:

• A monoidal ω -category $L(\mathbb{E}_1^\otimes: \mathbb{E}_1^\otimes \rightarrow \text{Cat}_\infty$

→ An LM^{\otimes} -monoid $LM^{\otimes} \rightarrow \text{Cat}$ consists of:

- A monoidal ∞ -category $L(\mathbb{E}_1^{\otimes}: \mathbb{E}_1^{\otimes} \rightarrow \text{Cat})$
- An object $L(\ell) \in \text{Cat}$
- A left-tensoring of $L(\ell)$ over $L(\mathbb{E}_1^{\otimes})$

→ A LM_{pt}^{\otimes} -monoidal ∞ -cat. is a LM^{\otimes} -monoidal one with a specified pointing $L \in L(\ell)$.

Thm (Dunn additivity) For \mathcal{C} s.m., equip $\text{Alg}_{\mathbb{E}_k}(\mathcal{C})$ with the pointwise (ie. absolute) symm. mon. structure. Then, $\text{Alg}_{\mathbb{E}_{n+k}}(\mathcal{C}) = \text{Alg}_{\mathbb{E}_n}(\text{Alg}_{\mathbb{E}_k}(\mathcal{C}))$

Ex/ What are $\text{Alg}_{\mathbb{E}_2}(\text{Cat}_1)$? → $\text{Alg}(\text{monoidal categories})$, ie. an 1-category \mathcal{C} equipped with monoidal structures \otimes, \boxtimes such that

- $1_{\boxtimes}: (1, \otimes) \rightarrow (e, \otimes)$ is monoidal, ie. $1_{\boxtimes} = 1_{\otimes}$
- $\boxtimes: (e, \otimes) \times (e, \otimes) \rightarrow (e, \otimes)$ is monoidal, so (Eckmann-Hilton)

$$X \boxtimes Y \cong (X \otimes 1) \boxtimes (1 \otimes Y) \xrightarrow{\text{ptwise}} (X \boxtimes 1) \otimes (1 \boxtimes Y) \cong X \otimes Y$$

But: $(1 \otimes X) \boxtimes (Y \otimes 1) \cong (1 \boxtimes Y) \otimes (X \boxtimes 1) \cong Y \otimes X \Rightarrow \text{Braiding!}$

Conversely, a natural trafo $\beta_{X,Y}: X \otimes Y \xrightarrow{\cong} Y \otimes X$ induces

$$(W \otimes X) \otimes (Y \otimes Z) \cong (W \otimes Y) \otimes (X \otimes Z)$$

exhibiting $\boxtimes: (e \times e, \otimes_{\text{ptwise}}) \rightarrow (e, \otimes)$ as monoidal iff hexagon identities.

Rem: Recall $\text{Mul}_{\mathbb{E}_2}(X_1, \dots, X_n; Y) = \text{Conf}_n(\mathbb{R}^2)$, we have chosen one of these operations to decompose $\mathbb{E}_2 \cong \mathbb{E}_1 \otimes \mathbb{E}_1$ via Dunn additivity. Hence, the space of choices for $X_1 \otimes \dots \otimes X_n$ is acted on by the **Artin braid group** $\Pi_n \text{Conf}_n(\mathbb{R}^2)$

For $n=0$, empty or contractible \Rightarrow posets

(Def) $\text{Cat}_{(n,m)} \subseteq \text{Cat}_{(\infty,n)}$ on those \mathcal{C} where $\text{Map}_{\mathcal{C}}(X,Y)$ is $(n-1)$ -truncated $\forall X,Y \in \mathcal{C}$

	sets	1-cats	(2,1)-cats	...
\mathbb{E}_0	pointed	pointed	pointed	...
\mathbb{E}_1	monoid	monoidal	monoidal	...
\mathbb{E}_2	commut. monoid	braided	braided	
\mathbb{E}_3	"	symmetric	symplectic	
\mathbb{E}_4	"	"	symmetric	

Aside: Can also introduce braided or

\mathbb{E}_k -operads as certain functors

$$\mathcal{O}^{\otimes} \rightarrow (\Delta^{\text{op}})^{\times k} \quad \text{See}$$

[Haugseng, "The higher Morita category of \mathbb{E}_n -algebras"]

also works for (n,m) -categories

"Baez-Dolan stabilization" starting at

$$\text{Alg}_{\mathbb{E}_{n+2}}(\text{Cat}_{(n,m)})$$

Stable ∞ -categories

→ This is a property!

Recall: A category \mathcal{C} is called abelian if

- It admits a zero object, products & coproducts

- automatic
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 - The canonical map $X \sqcup Y \xrightarrow{\begin{pmatrix} \text{id} & 0 \\ 0 & \text{id} \end{pmatrix}} X \times Y$ is always an iso
 - It is enriched over abelian groups \hookrightarrow induces the addition
 - It admits kernels and cokernels
- additive
- $$\begin{array}{ccccccc} \ker(f) & \longrightarrow & X & \xrightarrow{f} & Y & \longrightarrow & \text{coker}(f) \\ \downarrow & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & \text{dom}(f) & \xrightarrow{\cong} & \text{im}(f) & \longrightarrow & 0 \end{array}$$

Def An ∞ -category \mathcal{C} is called stable : \Leftrightarrow

- It admits a zero object 0
- It admits fibers and cofibers (analogy of kernels & cokernels)
- A sequence $X' \xrightarrow{f} X \xrightarrow{g} X''$ exhibits X' as fiber of g iff it exhibits X'' as the cofiber of f

Def For $X \in \mathcal{C}$, let $\begin{array}{ccc} \Sigma X & \longrightarrow & 0 \\ \downarrow & \lrcorner & \downarrow \\ 0 & \longrightarrow & X \end{array}$ and $\begin{array}{ccc} X & \longrightarrow & 0 \\ \downarrow & \lrcorner & \downarrow \\ 0 & \longrightarrow & \Sigma X \end{array}$

then $\Sigma : \mathcal{C} \xrightarrow{\cong} \mathcal{C} : \tau$ are inverse equivalences if \mathcal{C} stable.

\rightarrow Explains the name. Also write $X[1] = \Sigma X$.

Reason: $\begin{array}{ccccc} \Sigma X & \longrightarrow & 0 & \longrightarrow & X \\ X & \longrightarrow & 0 & \longrightarrow & \Sigma X \end{array}$ (fiber sequences)

Note: $\pi_0 \text{Map}(X, Y) = \pi_0 \text{Map}(X, \Omega^2 \Sigma^2 Y) = \pi_0 \Omega^2 \text{Map}(X, \Sigma^2 Y) = \pi_2 \text{Map}(X, Y)$
 \Rightarrow We can add & subtract morphisms, abelian group

In fact, get a spectrum / infinite loop space.

Note: $\text{fib}(X \xrightarrow{0} \Sigma Y) \longrightarrow X$ shows $X \times Y = \text{fib}(X \xrightarrow{0} \Sigma Y) \in \mathcal{C}$
 $\Sigma \Sigma Y = Y \longrightarrow 0$ similarly $X \sqcup Y = \text{cofib}(X[-1] \xrightarrow{0} Y)$
 $\rightarrow X \sqcup Y \cong X \times Y$ since \mathcal{C} additive (Also enriched ✓ has $X, \sqcup, 0$ ✓)

Proposition: Any stable ∞ -category \mathcal{C} has biproducts, ie. $X \sqcup Y \xrightarrow{\cong} X \times Y$.
 In fact, \mathcal{C} is additive.

Fact: $X \rightarrow X'$ in \mathcal{C} stable is a pullback iff it is a pushout square.
 $\begin{array}{ccc} X & \longrightarrow & X' \\ \downarrow & & \downarrow \\ Y & \longrightarrow & Y' \end{array}$ Thus, we can apply the pasting lemma in both directions.

Rem: The mirroring $\Delta^* \times \Delta^* \rightarrow \Delta^* \times \Delta^*$ swapping the components acts on $X \xrightarrow{f} Y$ classified by $\begin{array}{ccc} X[1] & \longrightarrow & 0 \\ \downarrow & & \downarrow \\ 0 & \longrightarrow & Y \end{array}$ by sending it to $X \xrightarrow{-f} Y$, as it inverts the loops in $\text{Map}(X, Y) = \Omega \text{Map}(X[1], Y)$.

Theorem If \mathcal{C} is a stable ∞ -category, then \mathcal{C} is triangulated, with shift functor $X[1] := \Sigma X$ and dist. triangles $\begin{array}{ccccc} X & \xrightarrow{f} & Y & \longrightarrow & 0 \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow \\ 0 & \longrightarrow & Z & \longrightarrow & Y[1] \end{array}$

dist. triangles

$$\begin{array}{ccccc} X & \xrightarrow{f} & Y & \longrightarrow & 0 \\ \downarrow 1 & & \downarrow g & \downarrow 1 & \\ 0 & \longrightarrow & Z & \xrightarrow{h} & X[1] \end{array}$$

Proof: • $X \xrightarrow{f} Y$ can always be completed to dist. tr. $X \xrightarrow{f} Y \rightarrow \text{cofib}(f)$

- If $f = \text{id}_X$, then $\text{cofib}(f) = 0$
- Isomorphic to dist. tr. \Rightarrow still one

Shift

$$\begin{array}{ccccc} X & \longrightarrow & Y & \longrightarrow & 0 \\ \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & Z & \longrightarrow & X[1] \\ & & \downarrow & & \downarrow f[1] \\ & & 0 & \longrightarrow & Y[1] \end{array}$$

— since mirrored \rightarrow "turns loop around"

- $\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow & & \downarrow \\ X' & \xrightarrow{f'} & Y' \end{array}$ induce $\begin{array}{ccc} \text{cofib}(f) & & \\ \downarrow & & \\ \text{cofib}(f') & & \end{array}$ fitting into dist. tr. \Rightarrow functorial cone

Octahedron axiom

$$\begin{array}{ccccccc} X & \xrightarrow{f} & Y & \xrightarrow{g} & Z & \longrightarrow & 0 \\ \downarrow & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & Y/X & \xrightarrow{g'} & Z/X & \longrightarrow & X[1] \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ & & 0 & \longrightarrow & Y/X[1] & \longrightarrow & Y/X[1] \end{array}$$

(Could go on to find "higher octahedron".)

□

Not every triangulated category arises this way, but "all the interesting ones", e.g.

- $\mathcal{C}h(\mathcal{A}) = \{ \text{chain cplx, chain maps, chain homotopies, ...} \}$ is stable
- Any (pre)triangulated dg-category has an associated (stable) ∞ -category
- For \mathcal{A} Grothendieck abelian (for simplicity),

$D(\mathcal{A}) = \{ \text{injective chain cplx, chain maps, chain homot., ...} \}$ is stable

& has all limits and colimits (presentable stable)

- $D(\text{Sh}(X; \mathcal{A})) = \text{Sh}_{\infty}^{\text{inj}}(X; D(\mathcal{A}))$ is stable \rightarrow Factorization algebras...
- $\text{Sp} = \lim (S_* \xrightarrow{\Omega} S_* \xleftarrow{\Omega} \dots)$ is stable, in fact universal among them

Rem: The cat. of abelian groups Ab is universal among additive categories as every add. cat. is Ab -enriched, and conversely an Ab -enriched category is "Cauchy-complete" if it stems from an idempotent complete add. cat.

Similarly: A Sp -enriched ∞ -category is Cauchy-complete iff it stems from the natural Sp -enrichment of an idemp. compl. stable ∞ -category.