Free All The Things

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Free All The Things

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- well known: free monads
- maybe known: free applicatives
- free monoids
- free <you name it>

This Talk

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- explain the technique behind "Free X"
- apply the technique to different examples
- takeaway: it's easier than you thought
- Source Code: https://github.com/markus1189/ free-all-the-things/tree/upnorth

Let's Get Started

Ready?



What's The Problem

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A free functor is left adjoint to a forgetful functor



What's The Problem

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A free functor is left adjoint to a forgetful functor

what's the problem?



What Is Free

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A free "thing" **FreeA** on a type(class) A is a A and a function

def inject(x: A): FreeA

such that for any other "thing" B and a function

val f: A => B

there exists a unique homomorphism g such that

What Is Free

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- still complicated!!!11
- a Free X is the minimal structure that satisfies the laws of the typeclass X (mine)
- rejoice: there is a pretty mechanical recipe
 - Define AST
 - Add Inject
 - Write Interpreter
 - Check laws

Why Free

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- nice API using typeclass
- use Free X as if it was X
- lift your DSL ops using inject (or lift)
- program reified into datastructure
- structure can be analyzed/optimized
- one program many interpretations

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Disclaimer Before We Start

- deep embeddings / initial encoding / data structure representation
- not: finally tagless, optimization, using Free Monads
- for some typeclasses, there are different sets of minimal operations (Monad: three!)
- fun thing: technique works for every one of them, homework

Freeing The Monad

The Monad Typeclass

we use the pure and flatMap version

```
trait Monad[F[ ]] {
1
       def pure [A](x:A):F[A]
2
3
       def flatMap[A, B](fa: F[A])(f: A \Rightarrow F[B]): F[B]
4
```

```
// Left identity
pure(a).flatMap(f) === f(a)

// Right identity
fa.flatMap(pure) === fa

// Associativity
fa.flatMap(f).flatMap(g) ===
fa.flatMap(a => f(a).flatMap(g))
```

Applying The Recipe

```
trait Monad[F[_]] {
    def pure[A](x: A): F[A]

def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

- remember our recipe
 - Define AST
 - Add Inject
 - Write Interpreter
 - · Check laws

Freeing The Monad

```
sealed abstract class Free [F], A
1
2
      final case class Pure[F[_], A](a: A)
3
          extends Free[F. A]
4
5
      final case class FlatMap[F[_], A, B](
6
          fa: Free[F, A],
7
          f: A \Rightarrow Free[F, B]
8
          extends Free[F, B]
10
      final case class Inject[F[_], A](fa: F[A])
11
          extends Free[F, A]
12
```

```
implicit def freeMonad[F[_], A]: Monad[Free[F, ?]] =
new Monad[Free[F, ?]] {
    def pure[A](x: A): Free[F, A] = Pure(x)

def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] =
FlatMap(fa, f)
}
```

Interpreter

```
def runFree[F[_], M[_]: Monad, A](nat: F ~> M)(
1
         free: Free[F, A]): M[A] = free match {
2
      case Pure(x) => Monad[M].pure(x)
3
      case Inject(fa) => nat(fa)
      case FlatMap(fa, f) =>
        Monad[M].flatMap(runFree(nat)(fa))(x =>
6
           runFree(nat)(f(x))
```

What about the laws?

```
// The associativity law
fa.flatMap(f).flatMap(g) ===
fa.flatMap(fa, a => f(a).flatMap(g))

val exp1 = FlatMap(FlatMap(fa, f), g)
val exp2 = FlatMap(fa, (a: Int) => FlatMap(f(a), g))
exp1 != exp2
```

What about the laws?





The Laws

- actually, we don't satisfy them
- programmer: after interpretation it's no longer visible
- mathematician: that's not the free monad!
- tradeoff: during construction vs during interpretation

Transforming Free Monads: Old Instance

```
implicit def freeMonad[F[_], A]: Monad[Free[F, ?]] =
new Monad[Free[F, ?]] {
def pure[A](x: A): Free[F, A] = Pure(x)

def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] =
FlatMap(fa, f)
}
```

Transforming Free Monads: Optimized flatMap

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

Transforming Free Monads: Optimized flatMap

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

- what do we learn?
- laws are not boring, they allow refactorings and optimization
- just by adhering to the law, we made our Free Monad more stack safe

```
FlatMap(FlatMap(FlatMap(fa, f), g), ...)
FlatMap(fa, x => FlatMap(f(x), y => FlatMap(q(y), z => ...)))
```

We Freed Monads

- DSL with monadic expressiveness
- context sensitive, branching, loops, fancy control flow
- familiarity with monadic style for DSL





And Once Again

- Define AST
- Add Inject
- Write Interpreter
- Check laws

The Functor Typeclass

```
trait Functor[F[_]] {
    def map[A, B](fa: F[A])(f: A => B): F[B]
}
```

The Functor Laws

```
1  // identity law
2  fa.map(identity) === fa
3
4  // composition
5  fa.map(f).map(q) === fa.map(f.compose(q))
```

```
sealed abstract class FreeFunctor[F[_], A]

case class Fmap[F[_], X, A](fa: F[X])(f: X => A)
    extends FreeFunctor[F, A]

case class Inject[F[_], A](fa: F[A])
    extends FreeFunctor[F, A]
```

- cool! that was easy
- can we get rid of something?



```
sealed abstract class FreeFunctor[F[], A]
2
     case class Fmap[F[_], X, A](fa: F[X])(f: X \Rightarrow A)
3
         extends FreeFunctor[F, A]
5
     case class Inject[F[_], A](fa: F[A])
6
         extends FreeFunctor[F, A]
7
```

Free Monad Free Functor Free Monoid Free Boolean Algebra Conclusio

Clean Code Police





Free Monad Free Functor Free Monoid Free Boolean Algebra Conclusion

Clean Code Police



• only one subclass?



```
sealed abstract class FreeFunctor [F], A
2
     case class Fmap[F[_], X, A](fa: F[X])(f: X \Rightarrow A)
3
         extends FreeFunctor[F, A]
4
5
     def inject[F[]], A](value: F[A]) =
6
       Fmap(value)(identity)
```

```
sealed abstract class Fmap[F[ ], A] {
       type X
2
       def fa: F[X]
       def f: X => A
6
     def inject[F[\ ], A](v: F[A]) = new Fmap[F, A] {
       type X = A
8
       def fa = v
       def f = identity
10
11
```

Freeing The Functor

```
sealed abstract class Coyoneda[F[ ], A] {
       type X
2
       def fa: F[X]
       def f: X => A
6
     def inject[F[\ ], A](v: F[A]) = new Coyoneda[F, A] {
       type X = A
8
       def fa = v
       def f = identity
10
11
```

```
implicit def coyoFun[F[_]]: Functor[Coyoneda[F, ?]] =
new Functor[Coyoneda[F, ?]] {
    def map[A, B](coyo: Coyoneda[F, A])(
        g: A => B): Coyoneda[F, B] =
    new Coyoneda[F, B] {
        type X = coyo.X
        def fa = coyo.fa
        def f = g.compose(coyo.f)
    }
}
```

Free Functor Interpreter

```
def runCoyo[F[_]: Functor, A](
coyo: Coyoneda[F, A]): F[A] =
Functor[F].map(coyo.fa)(coyo.f)
```

We Freed Functors

- DSL with hmm functorial expressiveness?
- map fusion! (functor law)
- interesting: combine with other structures
- boring interpreter, though
- still fun!



Freeing The Monoid

The Monoid Typeclass

```
trait Monoid[A] {
def empty: A
def combine(x: A, y: A): A
}
```

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Inject[A](x: A) extends FreeMonoid[A]

case class Combine[A](x: FreeMonoid[A],

y: FreeMonoid[A])

extends FreeMonoid[A]
```

The Laws

```
1  // left identity
2  empty |+| x === x
3
4  // right identity
5  x |+| empty === x
6
7  // associativity
8  1 |+| (2 |+| 3) === (1 |+| 2) |+| 3
```

The Laws and Free Monoid

- let's try to enforce those laws in our structure
- goal: correct by construction
- arbitrary decision: associate left vs right

Fixing Associativity

```
sealed trait NotCombine[+A]
1
2
      sealed abstract class FreeMonoid[+A]
3
4
      case object Empty
5
          extends FreeMonoid[Nothing]
6
          with NotCombine[Nothing]
7
8
      case class Inject[A](x: A)
          extends FreeMonoid[A]
10
          with NotCombine[A]
11
12
      case class Combine[A](x: NotCombine[A],
13
                             v: FreeMonoid[A])
14
          extends FreeMonoid[A]
15
```

The Problem With Neutral Elements

- get rid completely? not possible
- limit ourselves to a single element
- restrict Combine to have only real values on the left side
- goal: minimal canonical structure

```
case class Inject[A](x: A)

sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: Inject[A], y: FreeMonoid[A])

extends FreeMonoid[A]
```

Minimizing Structure — Remove Inject

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: A, y: FreeMonoid[A])
extends FreeMonoid[A]
```

The Monoid Instance

```
implicit def monoid[A]: Monoid[FreeMonoid[A]] =
new Monoid[FreeMonoid[A]] {
    override def empty = Empty
    override def combine(
        x: FreeMonoid[A],
        y: FreeMonoid[A]): FreeMonoid[A] = x match {
        case Empty => y
        case Combine(h, t) => Combine(h, combine(t, y))
    }
}
```

Minimizing Structure — Looks Familiar?

```
sealed abstract class FreeMonoid[+A]
2
     case object Empty extends FreeMonoid[Nothing]
3
4
     case class Combine[A](x: A, y: FreeMonoid[A])
5
         extends FreeMonoid[A]
6
```

Minimizing Structure — List

```
case object Nil extends List[Nothing]
case class Cons[A](head: A, tail: List[A])
extends List[A]
```

sealed abstract class List[+A]

We Freed Monoids

- DSL for combining things
- works beautifully with folds
- interpretation can be in parallel (associativity)



Free Monad Free Functor Free Monoid Free Boolean Algebra Conclusion

Now That We Can Free Anything



What should we free?



Credit Where It's Due

- Once upon a time: https://engineering.wingify.com/posts/Free-objects/
- use free boolean algebra to define DSL for event predicates
- credits to Chris Stucchio (@stucchio)

Let's Free A Boolean Algebra

- Define AST
- Add Inject
- Write Interpreter
- Check laws

```
trait BoolAlgebra[A] {
1
       def tru: A
2
       def fls: A
3
4
       def not(value: A): A
5
6
       def and (lhs: A, rhs: A): A
7
       def or(lhs: A, rhs: A): A
8
```

```
sealed abstract class FreeBool[+A]
1
2
     case object Tru extends FreeBool[Nothing]
3
     case object Fls extends FreeBool[Nothing]
5
     case class Not[A](value: FreeBool[A])
6
          extends FreeBool[A]
     case class And[A](lhs: FreeBool[A], rhs: FreeBool[A])
8
          extends FreeBool[A]
9
     case class Or[A](lhs: FreeBool[A], rhs: FreeBool[A])
10
          extends FreeBool[A]
11
     case class Inject[A](value: A) extends FreeBool[A]
12
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                      => B.tru
4
         case Fls => B.fls
5
         case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru => B.tru
         case Fls => B.fls
5
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                      => B.tru
4
         case Fls => B.fls
5
         case Inject(v) => f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
7
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                      => B.tru
4
                       => B.fls
         case Fls
5
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                      => B.tru
4
                       => B.fls
         case Fls
5
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
         case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
         case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

Using Free Bool

- that was simple (though boilerplate-y)
- what can we do with our new discovered structure
- reminder: boolean operators
 - true, false
 - and, or
 - xor, implies, nand, nor, nxor

sealed trait Search

```
case class Term(t: String) extends Search
2
     case class After(date: Date) extends Search
3
     case class InText(t: String) extends Search
     case class InUrl(url: String) extends Search
6
    // and the usual smart ctors
```

after sneaking in syntactic sugar behind the scenes:

```
val search = term("FP") &
after("20180101") &
 !(term("Java") | inText("spring")) &
inUrl("functionalconf")
```

The Site Type

```
case class Site(terms: List[String],
url: String,
indexedAt: Date,
text: String)
```

all predicates are run against a collection of these

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
         case Term(t)
                                  => site.terms.contains(t)
4
                                  => site.indexedAt > d
         case After(d)
5
          case InText(t: String) => site.text.contains(t)
6
          case InUrl(w)
                                 => site.url.contains(w)
8
        runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                 => site.terms.contains(t)
          case Term(t)
4
                                 => site.indexedAt > d
         case After(d)
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                                 => site.url.contains(w)
8
        runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
          case Term(t)
                                 => site.terms.contains(t)
                                 => site.indexedAt > d
         case After(d)
         case InText(t: String) => site.text.contains(t)
6
          case InUrl(w)
                                 => site.url.contains(w)
8
        runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
          case Term(t)
                                 => site.terms.contains(t)
4
                                 => site.indexedAt > d
          case After(d)
          case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                               => site.url.contains(w)
8
        runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

Demo Time



But Wait There's More

- short circuiting and other optimization
- what if you don't have all the information?
 - partially evaluate predicates
 - if evaluates successfully, done
 - else, send it on
- core language vs extension
 - Chris also demonstrates extension
 - translate a rich language to base instructions
 - with all the advantages

Optimizing Boolean Algebras

```
def optimize[A](fa: FreeBool[A]): FreeBool[A] =
1
       fa match {
2
                        => Tru
         case Tru
3
         case Fls => Fls
4
         case Inject(v) => Inject(v)
5
         case Not(Not(v)) => v
6
         case Not(v) => Not(v)
7
         case Or(Tru, ) => Tru
8
         case Or(_, Tru) => Tru
9
         case Or(x, y) => Or(x, y)
10
         case And(Fls, ) => Fls
11
         case And( , Fls) => Fls
12
         case And(x, y) \Rightarrow And(x, y)
13
14
```

Partial Evaluation

- idea: you might have only partial information
- evaluate as much as possible
- optimal: we can already reduce without needing more information
- otherwise: send it on (JSON, Protobuf, ...)

Partial Evaluation

```
def partialEvaluator[A, B](p: FreeBool[A])(
1
         f: A => Option[B])(implicit B: BoolAlgebra[B])
2
       : Either[FreeBool[A], B] = p match {
3
       case Tru => Right(B.tru)
4
       case Inject(v) => f(v).toRight(p)
5
       case Or(lhs, rhs) =>
6
         val(l, r) = (partialEvaluator(lhs)(f),
                        partialEvaluator(rhs)(f))
8
         // check if fully evaluated
9
         ???
10
11
       case => ???
12
```

Partial Evaluation

```
// fulltext not available
     case class SiteMetadata(terms: List[String],
                              url: String,
3
                              indexedAt: Date)
4
     def partially(meta: SiteMetadata)(
         p: Search): Option[Boolean] = p match {
6
                               => Some(meta.terms.contains(t))
       case Term(t)
7
       case After(d)
                               => Some(meta.indexedAt > d)
8
                               => Some(meta.url.contains(w))
       case InUrl(w)
       case InText(t: String) => None
10
11
```

Demo Time



We Freed Boolean Algebras

- good example of underused free structure
- partial evaluation
- serialize the AST (JSON, Protobuf, Avro, ...)
- exercise: minimize AST representation

Resources

- Free Boolean Algebra by Chris Stucchio https://engineering.wingify.com/posts/Free-objects/
- Source Code: https://github.com/markus1189/ free-all-the-things/tree/upnorth

Go And Free All The Things!

Markus Hauck (@markus1189)

Introduction

Free Monad

Free Functor

Free Monoid

Free Boolean Algebra

Conclusion

Freeing The Applicative

Freeing The Applicative

- free monads are great, but also limited
- we can't analyze the programs
- how about a smaller abstraction?

Recall

- Define AST
- Add Inject
- Write Interpreter
- Check laws

The Applicative Typeclass

```
trait Applicative[F[_]] {
   def pure[A](x: A): F[A]

def ap[A, B](fab: F[A => B], fa: F[A]): F[B]
}
```

AST for FreeApplicative

```
sealed abstract class FreeAp[F[ ]. A]
1
2
      final case class Pure[F[ ], A](a: A)
3
          extends FreeAp[F, A]
4
5
      final case class Ap[F[ ], A, B](
6
          fab: FreeAp[F, A \Rightarrow B],
7
          fa: FreeAp[F, A])
8
          extends FreeAp[F, B]
9
10
      final case class Inject[F[\ ], A](fa: F[A])
11
          extends FreeAp[F, A]
12
```

1

Laws

```
1 // identity
   Ap(Pure(identity), v) === v
3
  // composition
   Ap(Ap(Ap(Pure(_.compose), u), v), w) ===
     AD(u, AD(v, w))
  // homomorphism
   Ap(Pure(f), Pure(x)) === Pure(f(x))
10
  // interchange
11
   Ap(u, Pure(y)) === Ap(Pure((y)), u)
12
```

Don't Forget The Laws

Running FreeApplicatives

We Freed Applicatives

- DSL with applicative expressiveness
- context insensitive
- pure computation over effectful arguments
- more freedom during interpretation

Going Deeper

- try to encode one of the normal forms for boolean algebras
- try to remove Inject cases from Monad and Applicative
- free Magmas
- define free X using alternative minimal set of ops of the typeclass