Free All The Things

Markus Hauck @markus1189



Free All The Things

- · well known: free monads
- maybe known: free applicatives
- free monoids
- free <you name it>

This Talk

- explain the technique behind "Free X"
- apply the technique to different examples
- takeaway: you can safely do this at home

Let's Get Started

Ready?



What's The Problem

A free functor is left adjoint to a forgetful functor



What's The Problem

A free functor is left adjoint to a forgetful functor

what's the problem?



What Is Free

A free "thing" **FreeA** on a type(class) A is a A and a function

def inject(x: A): FreeA

such that for any other "thing" B and a function

val $f: A \Rightarrow B$

there exists a unique homomorphism g such that

What Is Free

- still complicated!!!11
- a Free X is the minimal structure that satisfies the laws of the typeclass X (mine)
- rejoice: there is a pretty mechanical recipe
 - Define AST
 - Add Inject
 - Write Interpreter
 - Check laws

Why Free

- but why should I care at all?
- nice API using typeclass
- use Free X as if it was X
- lift your DSL ops using inject (or lift)
- program reified into datastructure
- structure can be analyzed/optimized
- one program many interpretations

Disclaimer Before We Start

- deep embeddings / initial encoding / data structure representation
- not: finally tagless, optimization, using Free Monads
- finally tagless possible, but transformation/inspection of AST is more complex
- for some typeclasses, there are different sets of minimal operations (Monad: three!)
- talk might include some hand waving, but for your own good
- fun thing: technique works for every one of them, homework

Freeing The Monad

The Monad Typeclass

we use the pure and flatMap version

```
trait Monad[F[_]] {
    def pure[A](x: A): F[A]

def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

Give Me The Laws

```
1 // Left identity
   pure(a).flatMap(f) === f(a)
  // Right identity
   fa.flatMap(pure) === fa
6
  // Associativity
   fa.flatMap(f).flatMap(g) ===
     fa.flatMap(a \Rightarrow f(a).flatMap(g))
```

```
trait Monad[F[ ]] {
1
       def pure [A](x:A):F[A]
3
       def flatMap[A, B](fa: F[A])(f: A \Rightarrow F[B]): F[B]
```

- remember our recipe
 - Define AST
 - Add Inject
 - Write Interpreter
 - Check laws

Freeing The Monad

```
sealed abstract class Free [F], A
1
2
      final case class Pure[F[_], A](a: A)
3
          extends Free[F. A]
4
5
      final case class FlatMap[F[_], A, B](
6
          fa: Free[F, A],
7
          f: A \Rightarrow Free[F, B]
8
          extends Free[F, B]
10
      final case class Inject[F[_], A](fa: F[A])
11
          extends Free[F, A]
12
```

Freeing The Monad

• the Monad typeclass instance just wraps things up

Interpreter

going from our FreeMonad back into regular Monad-Land

```
def runFree[F[_], M[_]: Monad, A](nat: F ~> M)(
    free: Free[F, A]): M[A] = free match {
    case Pure(x) => Monad[M].pure(x)

    case Inject(fa) => nat(fa)

    case FlatMap(fa, f) =>
    Monad[M].flatMap(runFree(nat)(fa))(x =>
    runFree(nat)(f(x)))

}
```

What about the laws?

```
// The associativity law
   fa.flatMap(f).flatMap(g) ===
     fa.flatMap(fa, a \Rightarrow f(a).flatMap(q))
  val exp1 = FlatMap(FlatMap(fa, f), q)
   val exp2 = FlatMap(fa, (a: Int) => FlatMap(f(a), q))
3
   exp1 != exp2
```

What about the laws?





The Laws

- actually, we don't satisfy them
- programmer: after interpretation it's no longer visible
- mathematician: that's not the free monad!
- tradeoff: during construction vs during interpretation
- let's see what we can do here

Transforming Free Monads: Old Instance

```
implicit def freeMonad[F[_], A]: Monad[Free[F, ?]] =
new Monad[Free[F, ?]] {
    def pure[A](x: A): Free[F, A] = Pure(x)

def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] =
FlatMap(fa, f)
}
```

Transforming Free Monads: Optimized flatMap

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

Transforming Free Monads: Optimized flatMap

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

Transforming Free Monads: Done

- Quiz: what about fa.flatMap(pure) === fa
- what do we learn?
- laws are not boring, they allow refactorings and optimization
- just by adhering to the law, we made our Free Monad more stack safe

```
FlatMap(FlatMap(fa, f), q), ...)
FlatMap(fa, x \Rightarrow FlatMap(f(x), y \Rightarrow FlatMap(g(y), z \Rightarrow ...)))
```

We Freed Monads

- DSL with monadic expressiveness
- context sensitive, branching, loops, fancy control flow
- familiarity with monadic style for DSL

And Once Again

- Define AST
- Add Inject
- Write Interpreter
- Check laws

The Functor Typeclass

```
trait Functor[F[_]] {
    def map[A, B](fa: F[A])(f: A => B): F[B]
}
```

The Functor Laws

```
1  // identity law
2  fa.map(identity) === fa
3
4  // composition
5  fa.map(f).map(q) === fa.map(f.compose(q))
```

```
sealed abstract class FreeFunctor[F[_], A]

case class Fmap[F[_], X, A](fa: F[X])(f: X => A)
    extends FreeFunctor[F, A]

case class Inject[F[_], A](fa: F[A])
    extends FreeFunctor[F, A]
```

- cool! that was easy
- next step is to try to minimize the structure
- can we get rid of something?

Clean Code Police



troduction Free Monad Free Functor Free Monoid Free Boolean Algebra Conclusion

Clean Code Police



- only one subclass?
- what do we currently use the trait for?

Freeing The Functor

```
sealed abstract class Fmap[F[_], A] {
1
       type X
2
       def fa: F[X]
       def f: X => A
6
     def inject[F[\ ], A](v: F[A]) = new Fmap[F, A] {
       type X = A
       def fa = v
9
       def f = identity
10
11
```

wait a moment, does that look familiar?

```
sealed abstract class Coyoneda[F[ ], A] {
       type X
2
       def fa: F[X]
       def f: X => A
6
     def inject[F[\_], A](v: F[A]) = new Coyoneda[F, A] {
       type X = A
8
       def fa = v
       def f = identity
10
11
```

Free Functor Instance

```
implicit def coyoFun[F[_]]: Functor[Coyoneda[F, ?]] =
new Functor[Coyoneda[F, ?]] {
    def map[A, B](coyo: Coyoneda[F, A])(
        g: A => B): Coyoneda[F, B] =
    new Coyoneda[F, B] {
        type X = coyo.X
        def fa = coyo.fa
        def f = g.compose(coyo.f)
    }
}
```

Free Functor Interpreter

```
def runCoyo[F[_]: Functor, A](
coyo: Coyoneda[F, A]): F[A] =
Functor[F].map(coyo.fa)(coyo.f)
```

We Freed Functors

- DSL with hmm functorial expressiveness?
- map fusion! (functor law)
- interesting: combine with other structures
- boring interpreter, though
- still fun!

Freeing The Monoid



The Monoid Typeclass

```
trait Monoid[A] {
    def empty: A
    def combine(x: A, y: A): A
}
```

let's start blindly following our mechanical approach

The Free Monoid — First Try

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Inject[A](x: A) extends FreeMonoid[A]

case class Combine[A](x: FreeMonoid[A],

y: FreeMonoid[A])

extends FreeMonoid[A]
```

The Laws

```
1  // left identity
2  empty |+| x === x
3
4  // right identity
5  x |+| empty === x
6
7  // associativity
8  1 |+| (2 |+| 3) === (1 |+| 2) |+| 3
```

The Laws and Free Monoid

- let's try to enforce those laws in our structure
- goal: correct by construction
- (almost) arbitrary decision: associate left vs right
- you will see why we chose to associate to the right soon!

Fixing Associativity

- so this is okay:
- combine(Inject(1), Combine(Inject(2), Inject(3)))
 - but this should not be possible:
- combine(Combine(Inject(1), Inject(2)), Inject(3))

Fixing Associativity

```
sealed trait NotCombine[+A]
2
      sealed abstract class FreeMonoid[+A]
3
4
      case object Empty
          extends FreeMonoid[Nothing]
6
          with NotCombine[Nothing]
7
8
      case class Inject[A](x: A)
          extends FreeMonoid[A]
10
          with NotCombine[A]
11
12
      case class Combine[A](x: NotCombine[A],
13
                             v: FreeMonoid[A])
14
          extends FreeMonoid[A]
15
```

The Problem With Neutral Elements

- still problem: Combine(Empty, Empty)
- get rid completely? not possible
- limit ourselves to a single element
- restrict Combine to have only Inject values on the left side
- goal: minimal canonical structure

Minimizing Structure — Extract Inject

```
case class Inject[A](x: A)

sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: Inject[A], y: FreeMonoid[A])

extends FreeMonoid[A]
```

but now Inject looks silly...

Minimizing Structure — Remove Inject

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: A, y: FreeMonoid[A])
extends FreeMonoid[A]
```

The Monoid Instance

note: that might start to look familiar

Minimizing Structure — Looks Familiar?

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: A, y: FreeMonoid[A])
extends FreeMonoid[A]
```

Minimizing Structure — List

```
sealed abstract class List[+A]

case object Nil extends List[Nothing]

case class Cons[A](head: A, tail: List[A])
extends List[A]
```

We Freed Monoids

- DSL for combining things
- works beautifully with folds
- interpretation can be in parallel (associativity)

troduction Free Monad Free Functor Free Monoid Free Boolean Algebra Conclusion

Now That We Can Free Anything



What should we free?

Credit Where It's Due

- Once upon a time: https://engineering.wingify.com/posts/Free-objects/
- use free boolean algebra to define DSL for event predicates
- credits to Chris Stucchio (@stucchio)

Let's Free A Boolean Algebra

- Define AST
- Add Inject
- Write Interpreter
- Check laws

Boolean Algebras

```
trait BoolAlgebra[A] {
  def tru: A
  def fls: A

def not(value: A): A

def and(lhs: A, rhs: A): A
  def or(lhs: A, rhs: A): A
}
```

```
sealed abstract class FreeBool[+A]
1
2
     case object Tru extends FreeBool[Nothing]
3
     case object Fls extends FreeBool[Nothing]
5
     case class Not[A](value: FreeBool[A])
6
          extends FreeBool[A]
     case class And[A](lhs: FreeBool[A], rhs: FreeBool[A])
8
          extends FreeBool[A]
9
     case class Or[A](lhs: FreeBool[A], rhs: FreeBool[A])
10
          extends FreeBool[A]
11
     case class Inject[A](value: A) extends FreeBool[A]
12
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru => B.tru
         case Fls => B.fls
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru => B.tru
         case Fls => B.fls
5
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru => B.tru
         case Fls => B.fls
         case Inject(v) => f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
7
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru => B.tru
         case Fls => B.fls
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
7
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                     => B.tru
         case Fls => B.fls
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
         case Or(lhs, rhs) =>
8
           B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
         case And(lhs, rhs) =>
10
           B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

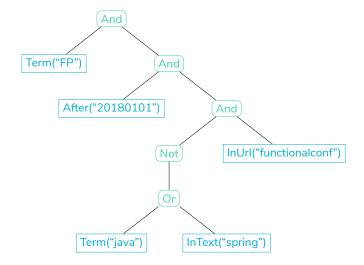
Using Free Bool

- that was simple (though boilerplate-y)
- skip trying to encode laws (because I don't know how)
- but as long as target is lawful, morally okay
- what can we do with our new discovered structure
- reminder: boolean operators
 - true, false
 - and, or
 - xor, implies, nand, nor, nxor

```
sealed trait Search
case class Term(t: String) extends Search
case class After(date: Date) extends Search
case class InText(t: String) extends Search
case class InUrl(url: String) extends Search
// and the usual smart ctors
```

after sneaking in syntactic sugar behind the scenes:

```
val search = term("FP") &
after("20180101") &
 !(term("Java") | inText("spring")) &
inUrl("functionalconf")
```



The Site Type

```
case class Site(terms: List[String],
url: String,
indexedAt: Date,
text: String)
```

- all predicates are run against a collection of these
- we know how to tear down FreeBool, but not yet how to interpret Search values

```
def evalSearch(pred: FreeBool[Search])(
1
         site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
         case Term(t)
                                 => site.terms.contains(t)
                                => site.indexedAt > d
         case After(d)
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                              => site.url.contains(w)
8
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

```
def evalSearch(pred: FreeBool[Search])(
1
         site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                 => site.terms.contains(t)
         case Term(t)
                                => site.indexedAt > d
         case After(d)
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                              => site.url.contains(w)
8
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

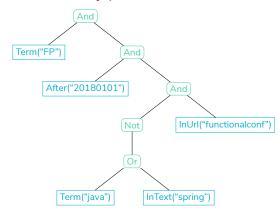
Free Bool Example: Search

```
def evalSearch(pred: FreeBool[Search])(
1
         site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
         case Term(t)
                                => site.terms.contains(t)
                                => site.indexedAt > d
         case After(d)
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                              => site.url.contains(w)
8
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

Free Bool Example: Search

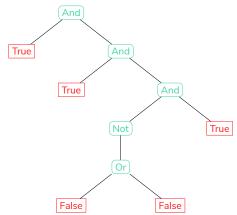
```
def evalSearch(pred: FreeBool[Search])(
1
         site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
         case Term(t)
                         => site.terms.contains(t)
                               => site.indexedAt > d
         case After(d)
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                             => site.url.contains(w)
8
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

```
Site(terms = List("FP", "bangalore"),
url = "functionalconf.com",
indexedAt = "20181212",
text = "lots of body")
```



Free Bool Example: Search

```
Site(terms = List("FP", "bangalore"),
url = "functionalconf.com",
indexedAt = "20181212",
text = "lots of body")
```



But Wait There's More

- short circuiting and other optimization
- what if you don't have all the information?
 - partially evaluate predicates
 - if evaluates successfully, done
 - else, send it on
- core language vs extension
 - Chris also demonstrates extension
 - translate a rich language to base instructions
 - with all the advantages

Optimizing Boolean Algebras

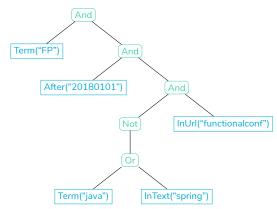
```
def optimize[A](fa: FreeBool[A]): FreeBool[A] =
1
       fa match {
2
                    => Tru
         case Tru
3
         case Fls => Fls
4
         case Inject(v) => Inject(v)
5
         case Not(Not(v)) => optimize(v)
6
         case Not(v) => Not(optimize(v))
         case Or(Tru, ) => Tru
8
         case Or( , Tru) => Tru
9
         case Or(x, y) => Or(optimize(x), optimize(y))
10
         case And(Fls, ) => Fls
11
         case And( , Fls) => Fls
12
         case And(x, y) => And(optimize(x), optimize(y))
13
14
```

- problem: you might have only partial information
- e.g. full text of indexed website slow to load
- evaluate as much as possible
- optimal: we can already reduce without needing more information
- otherwise: send it on (JSON, Protobuf, ...)

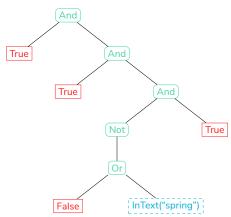
```
def partialEvaluator[A, B](p: FreeBool[A])(
1
         f: A => Option[B])(implicit B: BoolAlgebra[B])
2
       : Either[FreeBool[A], B] = p match {
3
       case Tru => Right(B.tru)
4
       case Inject(v) => f(v).toRight(p)
       case Or(lhs, rhs) =>
6
         val(l, r) = (partialEvaluator(lhs)(f),
                        partialEvaluator(rhs)(f))
8
         // check if fully evaluated
9
         ???
10
11
       case => ???
12
```

```
case class SiteMetadata(terms: List[String],
                              url: String,
2
                              indexedAt: Date)
3
4
     def partially(meta: SiteMetadata)(
         p: Search): Option[Boolean] = p match {
6
                               => Some(meta.terms.contains(t))
       case Term(t)
                               => Some(meta.indexedAt > d)
       case After(d)
8
                               => Some(meta.url.contains(w))
       case InUrl(w)
       case InText(t: String) => None
10
11
```

```
SiteMetadata(terms = List("FP", "bangalore"),
url = "functionalconf.com",
indexedAt = "20180929")
```



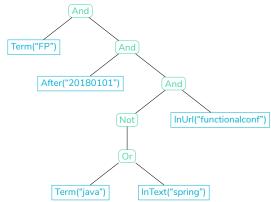
```
SiteMetadata(terms = List("FP", "bangalore"),
url = "functionalconf.com",
indexedAt = "20180929")
```



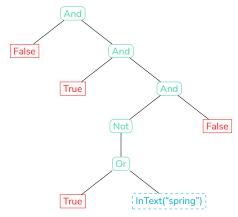
```
SiteMetadata(terms = List("FP", "bangalore"),
url = "functionalconf.com",
indexedAt = "20180929")
```

```
Not InText("spring")
```

```
SiteMetadata(terms = List("Java", "spring"),
url = "spring.io",
indexedAt = "20180929")
```



```
SiteMetadata(terms = List("Java", "spring"),
url = "spring.io",
indexedAt = "20180929")
```



```
SiteMetadata(terms = List("Java", "spring"),
url = "spring.io",
indexedAt = "20180929")
```



We Freed Boolean Algebras

- good example of underused free structure
- partial evaluation
- whole-program optimizations
- serialize the AST (JSON, Protobuf, ...)

Resources

- Free Boolean Algebra by Chris Stucchio https://engineering.wingify.com/posts/Free-objects/
- Source Code: https://github.com/markus1189/ free-all-the-things/tree/functionalconf

Go And Free All The Things!

Markus Hauck (@markus1189)

Introduction

Free Monad

Free Functor

Free Monoid

Free Boolean Algebra

Conclusion

Freeing The Applicative

Freeing The Applicative

- free monads are great, but also limited
- we can't analyze the programs
- how about a smaller abstraction?

Recall

- Define AST
- Add Inject
- Write Interpreter
- Check laws

The Applicative Typeclass

```
trait Applicative[F[_]] {
   def pure[A](x: A): F[A]

def ap[A, B](fab: F[A => B], fa: F[A]): F[B]
}
```

AST for FreeApplicative

```
sealed abstract class FreeAp[F[ ]. A]
1
2
      final case class Pure[F[ ], A](a: A)
3
          extends FreeAp[F, A]
4
5
      final case class Ap[F[ ], A, B](
6
          fab: FreeAp[F, A \Rightarrow B],
7
          fa: FreeAp[F, A])
8
          extends FreeAp[F, B]
9
10
      final case class Inject[F[\ ], A](fa: F[A])
11
          extends FreeAp[F, A]
12
```

1

Laws

```
1 // identity
   Ap(Pure(identity), v) === v
3
  // composition
   Ap(Ap(Ap(Pure(_.compose), u), v), w) ===
     AD(u, AD(v, w))
  // homomorphism
   Ap(Pure(f), Pure(x)) === Pure(f(x))
10
  // interchange
11
   Ap(u, Pure(y)) === Ap(Pure((y)), u)
12
```

Don't Forget The Laws

Running FreeApplicatives

We Freed Applicatives

- DSL with applicative expressiveness
- context insensitive
- pure computation over effectful arguments
- more freedom during interpretation

Going Deeper

- try to encode one of the normal forms for boolean algebras
- try to remove Inject cases from Monad and Applicative
- free Magmas
- define free X using alternative minimal set of ops of the typeclass