

# Assignment 4

MATH 667 Quantum Information Theory

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## 1 Problem 1

**Problem 1.** Let  $\mathcal{E} : H_n \rightarrow H_m$  be a linear map from the space of  $n \times n$  Hermitian matrices to the space of  $m \times m$  Hermitian matrices. The dual map  $\mathcal{E}^* : H_m \rightarrow H_n$  is the linear map defined by the relation

$$\text{Tr}[\rho \mathcal{E}^*(\sigma)] = \text{Tr}[\mathcal{E}(\rho) \sigma] \quad \forall \rho \in H_n \text{ and } \sigma \in H_m.$$

- (a) Show that if  $\mathcal{E}$  is a CP map then  $\mathcal{E}^*$  is also a CP map.
- (b) Show that if  $\mathcal{E}$  is TP then  $\mathcal{E}^*$  is unital (i.e. show that  $\mathcal{E}^*(I_m) = I_n$ ).
- (c) Let  $\rho, \sigma \in H_{n,+}$  be two density matrices. Show that  $\sigma \prec \rho$  (i.e. the vector of eigenvalues of  $\sigma$  is majorized by that of  $\rho$ ) if and only if there exists a unital CPTP map  $\mathcal{E} : H_n \rightarrow H_n$  such that  $\sigma = \mathcal{E}(\rho)$ .

**Solution.** In the following, we use  $\text{id}_{H_\ell}$  to denote the identity mapping on  $H_\ell$  for some integer  $\ell$ . Recall that a linear map  $\Lambda : H_n \rightarrow H_m$  is completely positive if and only if the map  $\Lambda \otimes \text{id}_{H_\ell}$  is positivity preserving for some integer  $\ell \geq n$ .

- (a) *Proof.* Suppose that  $\mathcal{E}$  is a completely positive map, let  $\ell = \max\{m, n\}$  and let  $P \in H_m \otimes H_\ell$  be a positive operator. We need to show that  $(\mathcal{E}^* \otimes \text{id}_{H_\ell})(P) \in H_n \otimes H_\ell$  is a positive operator. Recall that an operator  $A \in H_n \otimes H_\ell$  is positive if and only if it holds that  $\text{Tr}[AQ] \geq 0$  for all positive operators  $Q$ . Let  $Q \in H_n \otimes H_\ell$  be a positive, then

$$\text{Tr}[Q(\mathcal{E}^* \otimes \text{id}_{H_\ell})(P)] = \text{Tr}[(\mathcal{E} \otimes \text{id}_{H_\ell})(Q)P] \geq 0$$

where  $(\mathcal{E} \otimes \text{id}_{H_\ell})(Q)$  is positive by the complete positivity of  $\mathcal{E}$ . Hence  $\mathcal{E}^*$  is completely positive.  $\square$

- (b) *Proof.* Suppose that  $\mathcal{E}$  is trace preserving. Note that the identity matrix  $I_n$  is the only operator  $A \in H_n$  with the property that  $\langle \psi | A | \psi \rangle = 1$  holds for all unit vectors  $|\psi\rangle \in \mathbb{C}^n$ . Let  $|\psi\rangle \in \mathbb{C}^n$  be an arbitrary unit vector. Then

$$\langle \psi | \mathcal{E}^*(I_m) | \psi \rangle = \text{Tr}[|\psi\rangle\langle\psi| \mathcal{E}^*(I_m)] = \text{Tr}[\mathcal{E}(|\psi\rangle\langle\psi|) I_m] = \text{Tr}[\mathcal{E}(|\psi\rangle\langle\psi|)] = \text{Tr}[|\psi\rangle\langle\psi|] = 1$$

since  $\mathcal{E}$  is trace-preserving. It follows that  $\mathcal{E}^*(I_m) = I_n$ , as desired.  $\square$

- (c) *Proof.* Let  $\rho, \sigma \in H_{n,+}$  be two density matrices and let  $\vec{\lambda}$  and  $\vec{\mu}$  be the vectors of the eigenvalues of  $\rho$  and  $\sigma$  respectively. There exist orthonormal bases  $\{|u_i\rangle\}$  and  $\{|v_i\rangle\}$  of  $\mathbb{C}^n$  such that

$$\rho = \sum_i \lambda_i |u_i\rangle\langle u_i| \quad \text{and} \quad \sigma = \sum_i \mu_i |v_i\rangle\langle v_i|.$$

Furthermore, define unitary operators  $U = \sum_i |u_i\rangle\langle i|$  and  $V = \sum_i |v_i\rangle\langle i|$  such that  $\rho = U \text{diag}(\vec{\lambda}) U^*$  and  $\sigma = V \text{diag}(\vec{\mu}) V^*$ .

We first suppose that  $\sigma \prec \rho$ , that is  $\vec{\mu} \prec \vec{\lambda}$ . By assumption, there exists a doubly stochastic matrix  $D$  such that  $\vec{\mu} = D\vec{\lambda}$ . Since  $D$  is doubly stochastic, it can be written as a convex combination of permutation matrices

$$D = \sum_{\pi} p_{\pi} P_{\pi},$$

where the sum is taken over all permutations  $\pi$  of  $n$  elements,  $P_{\pi}$  is the permutation operator corresponding to  $\pi$  defined by

$$P_{\pi} = \sum_{i=1}^n |\pi(i)\rangle \langle i|,$$

and each  $p_{\pi} \geq 0$  is a probability such that  $\sum_{\pi} p_{\pi} = 1$ . Note that  $P_{\pi}^* = P_{\pi}^{-1}$  and thus  $P_{\pi}^* P_{\pi} = P_{\pi} P_{\pi}^* = I_n$  holds for every permutation  $\pi$ . Therefore

$$\sum_{\pi} p_{\pi} P_{\pi}^* P_{\pi} = \sum_{\pi} p_{\pi} P_{\pi} P_{\pi}^* = \sum_{\pi} p_{\pi} I_n = I_n.$$

Consider the collection of operators  $\{K_{\pi}\}$  where  $K_{\pi} = \sqrt{p_{\pi}} V P_{\pi}^* U^*$  for each  $\pi$ . We see that

$$\sum_{\pi} K_{\pi}^* K_{\pi} = \sum_{\pi} p_{\pi} U P_{\pi} V^* V P_{\pi}^* U^* = U \left( \sum_{\pi} p_{\pi} P_{\pi} P_{\pi}^* \right) U^* = U I_n U^* = I_n$$

and thus the collection  $\{K_{\pi}\}$  is a Kraus representation for the channel  $\mathcal{E} : \mathcal{H}_n \rightarrow \mathcal{H}_n$  defined by

$$\mathcal{E}(A) = \sum_{\pi} K_{\pi} A K_{\pi}^* = \sum_{\pi} p_{\pi} V P_{\pi}^* U^* A U P_{\pi} V$$

for all  $A \in \mathcal{H}_n$ . Note that  $\mathcal{E}$  is unital. Indeed, the dual map  $\mathcal{E}^*$  can be given by  $\mathcal{E}^*(A) = \sum_{\pi} K_{\pi}^* A K_{\pi}$  for all operators  $A$  and

$$\sum_{\pi} K_{\pi} K_{\pi}^* = \sum_{\pi} p_{\pi} V P_{\pi}^* U^* U P_{\pi} V^* = V \left( \sum_{\pi} p_{\pi} P_{\pi} P_{\pi}^* \right) V^* = V I_n V^* = I_n.$$

Hence  $\mathcal{E}^*$  is also a CPTP map with Kraus operators  $\{K_{\pi}^*\}$ . Since  $\mathcal{E}^*$  is trace preserving and  $(\mathcal{E}^*)^* = \mathcal{E}$ , it follows from part (b) that  $\mathcal{E}$  is unital.

Finally, note that  $P_{\pi}^* \text{diag}(\vec{\lambda}) P_{\pi} = \text{diag}(P_{\pi} \vec{\lambda})$  holds for every  $\pi$ , since

$$P_{\pi}^* \text{diag}(\vec{\lambda}) P_{\pi} = \sum_{i,j,k} |i\rangle \langle \pi(i)| (\lambda_j |j\rangle \langle j|) |\pi(k)\rangle \langle k| = \sum_i \lambda_i |\pi(i)\rangle \langle \pi(i)| = \sum_i \lambda_{\pi^{-1}(i)} |i\rangle \langle i| = \text{diag}(P_{\pi} \vec{\lambda})$$

and  $P_{\pi} \vec{\lambda} = \sum_i \lambda_i |\pi(i)\rangle = \sum_i \lambda_{\pi^{-1}(i)} |i\rangle$ . Now

$$\begin{aligned} \mathcal{E}(\rho) &= \sum_{\pi} p_{\pi} V P_{\pi}^* U^* \rho U P_{\pi} V^* = \sum_{\pi} p_{\pi} V P_{\pi}^* \text{diag}(\vec{\lambda}) P_{\pi} V^* \\ &= \sum_{\pi} p_{\pi} V \text{diag}(P_{\pi} \vec{\lambda}) V^* \\ &= V \text{diag} \left( \sum_{\pi} p_{\pi} P_{\pi} \vec{\lambda} \right) V^* = V \text{diag}(\vec{\mu}) V^* = \sigma, \end{aligned}$$

and thus  $\mathcal{E}(\rho) = \sigma$  for the unital CPTP map  $\mathcal{E}$ .

Now suppose that there exists a unital CPTP map  $\mathcal{E}$  such that  $\sigma = \mathcal{E}(\rho)$ . Define an  $n \times n$  matrix  $D$  whose elements are given by

$$D_{ij} = \text{Tr}[|v_i\rangle \langle v_i| \mathcal{E}(|u_j\rangle \langle u_j|)].$$

Note that  $\sum_i |u_i\rangle\langle u_i| = \sum_i |v_i\rangle\langle v_i| = I_n$ , since both  $\{|u_j\rangle\}$  and  $\{|v_i\rangle\}$  are orthonormal bases. All of the columns and rows of  $D$  each sum to 1, since

$$\sum_i D_{ij} = \text{Tr}[I_n \mathcal{E}(|u_j\rangle\langle u_j|)] = \text{Tr}[\mathcal{E}(|u_j\rangle\langle u_j|)] = 1$$

holds for all  $j$  by the fact that  $\mathcal{E}$  is trace-preserving and

$$\sum_j D_{ij} = \text{Tr}[|v_i\rangle\langle v_i| \mathcal{E}(I_n)] = \text{Tr}[|v_i\rangle\langle v_i|] = 1$$

holds for all  $i$  by the fact that  $\mathcal{E}$  is unital. Furthermore, each  $D_{ij}$  is nonnegative, since  $\mathcal{E}(|u_j\rangle\langle u_j|)$  is a positive operator by positivity of  $\mathcal{E}$ , and thus

$$D_{ij} = \langle v_i | \mathcal{E}(|u_j\rangle\langle u_j|) | v_i \rangle \geq 0$$

for all  $i$  and  $j$ . It follows that  $D$  is a doubly stochastic matrix.

We now show that  $\vec{\mu} = D\vec{\lambda}$ . The  $i^{\text{th}}$  entry of  $D\vec{\lambda}$  is

$$\begin{aligned} (D\vec{\lambda})_i &= \sum_j \text{Tr}[|v_i\rangle\langle v_i| \mathcal{E}(|u_j\rangle\langle u_j|)] \lambda_j \\ &= \text{Tr}\left[|v_i\rangle\langle v_i| \mathcal{E}\left(\underbrace{\sum_j \lambda_j |u_j\rangle\langle u_j|}_{\rho}\right)\right] = \text{Tr}[|v_i\rangle\langle v_i| \sigma] = \langle v_i | \sigma | v_i \rangle = \mu_i, \end{aligned}$$

and thus  $\vec{\mu} = D\vec{\lambda}$  for a doubly stochastic matrix  $D$ . It follows that  $\vec{\mu} \prec \vec{\lambda}$  and thus  $\sigma \prec \rho$ , as desired.  $\square$

## 2 Problem 2

**Problem 2.** Find necessary and sufficient conditions for which the following equality holds:

$$S(\rho^{AB}) = |S(\rho^A) - S(\rho^B)|. \quad (2.1)$$

Give an example.

**Solution.** The necessary and sufficient conditions for equality come from the following observation together with Proposition 1. The problem can be broken down into two sub-statements:

- (i)  $S(\rho^{AB}) = S(\rho^B) - S(\rho^A)$  if and only if  $\rho^{AR} = \rho^A \otimes \rho^R$  holds for all purifications  $\rho^{ABR}$  of  $\rho^{AB}$ , and
- (ii)  $S(\rho^{AB}) = S(\rho^A) - S(\rho^B)$  if and only if  $\rho^{BR} = \rho^B \otimes \rho^R$  holds for all purifications  $\rho^{ABR}$  of  $\rho^{AB}$ .

The following proposition proves only part (i), but flipping the A and B yields part (ii). Note that  $S(\rho^{AB})$  must be nonnegative. Hence equality in (2.1) holds if and only if the condition in either (i) or (ii) holds.

**Proposition 1.** *Let  $\rho^{AB}$  be a bipartite state. Then the equality  $S(\rho^{AB}) = S(\rho^B) - S(\rho^A)$  holds if and only if it holds that  $\rho^{AR} = \rho^A \otimes \rho^R$  for all possible purifications  $\rho^{ABR}$  of  $\rho^{AB}$ .*

*Proof.* Let  $\rho^{ABR}$  be a purification of  $\rho^{AB}$  such that  $\rho^{ABR} = |\psi\rangle\langle\psi|^{ABR}$  for a pure state vector  $|\psi\rangle^{ABR}$  and

$$\rho^{AB} = \text{Tr}_R \rho^{ABR}.$$

Suppose that  $\rho^{AR} = \rho^A \otimes \rho^R$ . Recall that the von Neumann entropy is sub-additive,  $S(\rho^{AR}) \leq S(\rho^A) + S(\rho^R)$  with equality if and only if  $\rho^{AR} = \rho^A \otimes \rho^R$  (i.e., systems A and R are uncorrelated). It follows that

$$S(\rho^{AR}) = S(\rho^A) + S(\rho^R) \quad (2.2)$$

from the assumption. Since the state  $\rho^{ABR} = |\psi\rangle\langle\psi|^{ABR}$  is pure, it holds that

$$S(\rho^R) = S(\rho^{AB}) \quad \text{and} \quad S(\rho^B) = S(\rho^{AR}). \quad (2.3)$$

Putting together equations (2.2) and (2.3) yields the equality  $S(\rho^{AB}) = S(\rho^B) - S(\rho^A)$ .

For the converse, the exact same argument works in reverse. That is, if we suppose that  $S(\rho^{AB}) = S(\rho^B) - S(\rho^A)$ , we can make the same replacements as above to yield the equality  $S(\rho^{AR}) = S(\rho^A) + S(\rho^R)$  for any possible purification  $\rho^{ABR}$ . This equality occurs if and only if  $\rho^{AR} = \rho^A \otimes \rho^R$ , as desired.  $\square$

**Example.** The “trivial” examples of states  $\rho^{AB}$  that satisfy this equality in (2.1) are those for which either

- $\rho^{AB}$  is pure, in which case  $S(\rho^{AB}) = 0$  and  $S(\rho^A) = S(\rho^B)$ , or
- $\rho^{AB} = \rho^A \otimes \rho^B$  and at least one of  $\rho^A$  or  $\rho^B$  is pure, in which case  $S(\rho^{AB}) = S(\rho^A)$  if  $\rho^B$  is pure or  $S(\rho^{AB}) = S(\rho^B)$  if  $\rho^A$  is pure.

We can construct a non-trivial example of a state  $\rho^{AB}$  that satisfies the equality in (2.1) in the following. Consider the pure state vectors  $|u\rangle$  and  $|v\rangle$  in  $\mathbb{C}^2 \otimes \mathbb{C}^4$  defined by

$$|u\rangle^{AB} = \frac{1}{\sqrt{2}}(|00\rangle^{AB} + |11\rangle^{AB}) \quad \text{and} \quad |v\rangle^{AB} = \frac{1}{\sqrt{2}}(|02\rangle^{AB} + |13\rangle^{AB}).$$

Note that  $\langle u|v\rangle = 0$ . Define the following mixed state (which is certainly not pure)

$$\rho^{AB} = \frac{1}{2}|u\rangle\langle u| + \frac{1}{2}|v\rangle\langle v|$$

and note that  $S(\rho^{\text{AB}}) = -\frac{1}{2} \log \frac{1}{2} - \frac{1}{2} \log \frac{1}{2} = \log 2 = 1$ . The reduced density operators on systems A and B are

$$\rho^{\text{A}} = \frac{1}{2}(|0\rangle\langle 0| + |1\rangle\langle 1|) \quad \text{and} \quad \rho^{\text{B}} = \frac{1}{4}(|0\rangle\langle 0| + |1\rangle\langle 1| + |2\rangle\langle 2| + |3\rangle\langle 3|),$$

both of which are clearly not pure. The entropies of these reduced density operators are  $S(\rho^{\text{A}}) = \log 2 = 1$  and  $S(\rho^{\text{B}}) = \log 4 = 2$ . We see that

$$|S(\rho^{\text{A}}) - S(\rho^{\text{B}})| = |1 - 2| = 1 = S(\rho^{\text{AB}}),$$

as desired.

### 3 Problem 3

**Problem 3.** Prove the strong sub-additivity of the Shannon entropy: For three random variables  $X, Y, Z$ ,

$$H(X, Y, Z) + H(Y) \leq H(X, Y) + H(Y, Z)$$

where  $H$  is the Shannon entropy.

**Solution.** *Proof.* Define the mutual information of two variables  $X$  and  $Y$  as

$$I(X : Y) = H(X) + H(Y) - H(X, Y).$$

From the sub-additivity of the Shannon entropy,  $H(X, Y) \leq H(X) + H(Y)$ , it is clear that  $I(X : Y) \geq 0$ . Define the conditional mutual information as

$$I(X : Z | Y) = \sum_y p(y) I(X|_{Y=y} : Z|_{Y=y}).$$

Since both  $p(y)$  and  $I(X|_{Y=y} : Z|_{Y=y})$  are nonnegative for every  $y$ , it follows that  $I(X : Z | Y) \geq 0$ . Recall that  $p(x|Y=y) = \frac{p(x,y)}{p(y)}$ . By definition of the mutual information, we have

$$\begin{aligned} p(y) I(X|_{Y=y} : Z|_{Y=y}) &= p(y) H(X|_{Y=y}) + p(y) H(Z|_{Y=y}) - p(y) H(X, Z|_{Y=y}) \\ &= - \sum_x p(x, y) \log \frac{p(x, y)}{p(y)} - \sum_z p(z, y) \log \frac{p(z, y)}{p(y)} + \sum_{x, z} p(x, y, z) \log \frac{p(x, y, z)}{p(y)} \\ &= - \sum_x p(x, y) \log p(x, y) - H(Y) - \sum_z p(z, y) \log p(z, y) - H(Y) \\ &\quad + \sum_{x, z} p(x, y, z) \log p(x, y, z) + H(Y) \\ &= - \sum_x p(x, y) \log p(x, y) - \sum_z p(z, y) \log p(z, y) + \sum_{x, z} p(x, y, z) \log p(x, y, z) - H(Y). \end{aligned}$$

By taking the sum over all  $y$ 's in the definition of  $I(X : Z | Y)$ , we have

$$\begin{aligned} I(X : Z | Y) &= \sum_y p(y) I(X|_{Y=y} : Z|_{Y=y}) \\ &= H(X, Y) + H(Y, Z) - H(X, Y, Z) - H(Y). \end{aligned}$$

Since  $I(X : Z | Y) \geq 0$ , from this it follows that

$$H(X, Y) + H(Y, Z) - H(X, Y, Z) - H(Y) \geq 0$$

and thus  $H(X, Y, Z) + H(Y) \leq H(X, Y) + H(Y, Z)$  as desired.  $\square$

## 4 Problem 4

**Problem 4.** Consider an i.i.d. source characterized by a random variable  $X$  with alphabet  $x \in \{1, \dots, d\} = \mathcal{X}$  corresponding to probability  $p(x) > 0$ . Consider a sequence of size  $n$ , denoted as  $x^n = (x_1, \dots, x_n) \in \mathcal{X}^n$ . The empirical distribution of the sequence  $x^n$  is defined by

$$q_{x^n}(x) = \frac{1}{n} N(x|x^n)$$

where  $N(x|x^n)$  is the number of times the symbol  $x \in \{1, \dots, d\}$  appears in  $x^n$ . For  $\delta > 0$  denote

$$T(n, \delta) = \{x^n \in \mathcal{X}^n \mid |q_{x^n}(x) - p(x)| < \delta \ \forall x \in \mathcal{X}\}.$$

(a) Show that for any  $\epsilon, \delta > 0$  and sufficiently large  $n$

$$\Pr(T(n, \delta)) \geq 1 - \epsilon.$$

(b) Show that for any  $\epsilon, \delta > 0$  and sufficiently large  $n$

$$(1 - \epsilon)2^{n(H(X) - c\delta)} \leq |T(n, \delta)| \leq 2^{n(H(X) + c\delta)}$$

for some positive constant  $c$ .

(c) Show that if  $x^n \in T(n, \delta)$  then

$$2^{-n(H(X) + c\delta)} \leq p(x^n) \leq 2^{-n(H(X) - c\delta)} \quad (4.1)$$

for some positive constant  $c$ , and  $p(x^n) = p(x_1)p(x_2) \cdots p(x_n)$ .

**Solution.** The idea for this solution is due to Wilde<sup>1</sup> and Yeung<sup>2</sup>.

(a) Let  $\epsilon, \delta > 0$ . For each  $x \in \mathcal{X}$ , consider the i.i.d. indicator random variables  $I_1(x), \dots, I_n(x)$  obtained by sampling  $x_i$  and then setting

$$I_i(x) = \begin{cases} 1 & \text{if } x_i = x \\ 0 & \text{if } x_i \neq x \end{cases}.$$

We can write  $N(x|x^n)$  as

$$N(x|x^n) = \sum_{i=1}^n I_i(x).$$

Since  $\Pr(I_i(x) = 1) = p(x)$ , we have the expected values  $E[I_i(x)] = p(x)$  for all  $x \in \mathcal{X}$ . By the weak law of large numbers, for every  $a \in \mathcal{X}$ , there is a sufficiently large  $n_a$  such that

$$\Pr(\{|q_{x^n}(a) - p(a)| \geq \delta\}) = \Pr\left(\left|\frac{1}{n} \sum_{i=1}^n I_i(a) - p(a)\right| \geq \delta\right) < \frac{\epsilon}{d} \quad (4.2)$$

holds for all  $n > n_a$ , where  $d = |\mathcal{X}|$ . Let  $n_0 = \max\{n_a \mid a \in \mathcal{X}\}$ . Then for all  $n > n_0$ , we have

$$\begin{aligned} \Pr(\{|q_{x^n}(a) - p(a)| \geq \delta \text{ for some } a \in \mathcal{X}\}) &= \Pr\left(\left|\frac{1}{n} \sum_{i=1}^n I_i(a) - p(a)\right| \geq \delta \text{ for some } a \in \mathcal{X}\right) \\ &= \Pr\left(\bigcup_{a \in \mathcal{X}} \left\{\left|\frac{1}{n} \sum_{i=1}^n I_i(a) - p(a)\right| \geq \delta\right\}\right) \end{aligned}$$

<sup>1</sup>See Section 14.7.1 - 14.7.3 in *Classical to Quantum Shannon Theory* by Mark Wilde.

<sup>2</sup>See Chapter 6 in *Information Theory and Network Coding* by Raymond Yeung.

$$\begin{aligned} &\leq \sum_{a \in \mathcal{X}} \Pr \left( \left\{ \left| \frac{1}{n} \sum_{i=1}^n I_i(a) - p(a) \right| \geq \delta \right\} \right) \\ &< \sum_{a \in \mathcal{X}} \frac{\epsilon}{d} = \epsilon, \end{aligned}$$

where the first inequality follows from the union bound<sup>3</sup> and the last inequality follows from (4.2). The probability of the complement of the above event is thus

$$\Pr \left( \{ |q_{x^n}(a) - p(a)| < \delta \text{ for all } a \in \mathcal{X} \} \right) = 1 - \epsilon$$

for all  $n > n_0$ , as desired.

- (b) We will prove part (b) from part (c). Hence, assuming part (c) is true, we can use the upper bound in (4.1) (i.e.  $p(x^n) \leq 2^{-n(H(X)-c\delta)}$ ) to find that

$$\Pr(T(n, \delta)) = \sum_{x^n \in T(n, \delta)} p(x^n) \leq |T(n, \delta)| 2^{-n(H(X)-c\delta)}$$

for some positive constant  $c$ . For  $n$  sufficiently large, we have  $1 - \epsilon \leq \Pr(T(n, \delta))$ , and it follows from part (a) that

$$1 - \epsilon \leq \Pr(T(n, \delta)) \leq |T(n, \delta)| 2^{-n(H(X)-c\delta)}.$$

Thus  $(1 - \epsilon) 2^{n(H(X)-c\delta)} \leq |T(n, \delta)|$ , which yields the desired lower bound. For the upper bound, we use the lower bound in (4.1) (i.e.  $\leq 2^{-n(H(X)+c\delta)} \leq p(x^n)$ ). Now

$$|T(n, \delta)| 2^{-n(H(X)+c\delta)} \leq \sum_{x^n \in T(n, \delta)} p(x^n) = \Pr(T(n, \delta)) \leq 1,$$

and rearranging yields the desired upper bound  $|T(n, \delta)| \leq 2^{n(H(X)+c\delta)}$ . (Note that this upper bound holds for all  $n$ .)

- (c) If  $p(a) \neq 0$  for every  $a \in \mathcal{X}$ , for any sequence  $x^n$  we can write

$$p(x^n) = \prod_{a \in \mathcal{X}} p(a)^{N(a|x^n)}.$$

If  $p(a) = 0$  for any  $a \in \mathcal{X}$ , then we must consider the subset  $\mathcal{X}^+$  of  $\mathcal{X}$  containing all  $a$  for which  $p(a) > 0$ . For  $n$  large enough, it must hold that

$$p(x^n) = \prod_{a \in \mathcal{X}^+} p(a)^{N(a|x^n)}. \quad (4.3)$$

for any  $x^n \in T(n, \delta)$ . That is, no symbols  $a \in \mathcal{X}$  with probability  $p(a) = 0$  can appear in typical sequences  $x^n$  if  $n$  is sufficiently large. Taking the logarithm of (4.3) yields

$$\begin{aligned} \log p(x^n) &= \sum_{a \in \mathcal{X}^+} N(a|x^n) \log p(a) \\ &= \sum_{a \in \mathcal{X}^+} (N(a|x^n) + np(a) - np(a)) \log p(a) \\ &= n \sum_{a \in \mathcal{X}^+} p(a) \log p(a) + n \sum_{a \in \mathcal{X}^+} \left( \frac{1}{n} N(a|x^n) - p(a) \right) \log p(a) \end{aligned}$$

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<sup>3</sup>The union bound refers to the fact that  $\Pr(\bigcup_i A_i) \leq \sum_i \Pr(A_i)$  for any set of events  $A_i$ .



$$= -n \left( H(X) + \sum_{a \in \mathcal{X}^+} \left( \frac{1}{n} N(a|x^n) - p(a) \right) (-\log p(a)) \right).$$

Since  $x^n \in T(n, \delta)$ , we know

$$\left| \frac{1}{n} N(a|x^n) - p(a) \right| \leq \delta$$

which implies

$$\begin{aligned} \left| \sum_{a \in \mathcal{X}^+} \left( \frac{1}{n} N(a|x^n) - p(a) \right) (-\log p(a)) \right| &\leq \sum_{a \in \mathcal{X}^+} \left| \left( \frac{1}{n} N(a|x^n) - p(a) \right) \right| (-\log p(a)) \\ &\leq \delta \sum_{a \in \mathcal{X}^+} (-\log p(a)) \\ &= \delta c \end{aligned}$$

where we take  $c$  to be the positive constant

$$c = - \sum_{a \in \mathcal{X}^+} \log p(a).$$

From this we see that

$$-\delta c \leq \sum_{a \in \mathcal{X}^+} \left( \frac{1}{n} N(a|x^n) - p(a) \right) (-\log p(a)) \leq \delta c$$

and thus

$$-n(H(X) + \delta c) \leq -n \underbrace{\left( H(X) + \sum_{a \in \mathcal{X}^+} \left( \frac{1}{n} N(a|x^n) - p(a) \right) \log p(a) \right)}_{\log p(x^n)} \leq -n(H(X) - \delta c),$$

which yields

$$2^{-n(H(X) + \delta c)} \leq p(x^n) \leq 2^{-n(H(X) - \delta c)},$$

as desired.