# FUNCTIONAL PROGRAMMING AT SIMCORP



## THE SPEAKER

#### FLORIAN BIERMANN

#### ITU alumn

- M.Sc. in Software Development (2014)
- Ph.D. in Computer Science (2018)
  - "Data Parallel Spreadsheet Programming"

#### Developer at SimCorp since August 2018

- OTC Core
- C# and OCaml, sometimes a bit of F# and Emacs Lisp
- Financial contracts, derivatives, framework, application integration







## **SIMCORP**

#### INVESTMENT MANAGEMENT SOLUTIONS PROVIDER

- Established in 1971
- EUR 382.6 million revenue
- 25 offices globally, main development in DK and UA
- 1900 employees (2020)
- 300+ SimCorp Dimension clients worldwide



Fully integrated front-to-back investment management solution, powered by an award-winning Investment Book of Record, offered globally

190+ clients



## FINANCIAL CONTRACTS

#### WHAT DO WE DO – AND WHY?

#### Bond (loan)

- Pay amount up front (nominal)
- Pay interest on nominal over time
- Pay back nominal in the end

#### Option (financial instrument)

- Right to exercise an underlying contract.
- E.g. take up loan on already agreed conditions.
- In a period or at specific points in time.

#### Over-the-Counter (OTC)

- Highly customizable financial instruments
- Not "centrally cleared", bilateral agreement.

Like an insurance: if it rains in May, you may buy crops at \$5 a pound.

Financial contracts have market value!

Cross-currency trades, interest swaps, Sell-buy-back,

- -



## SIMCORP TECHNOLOGIES

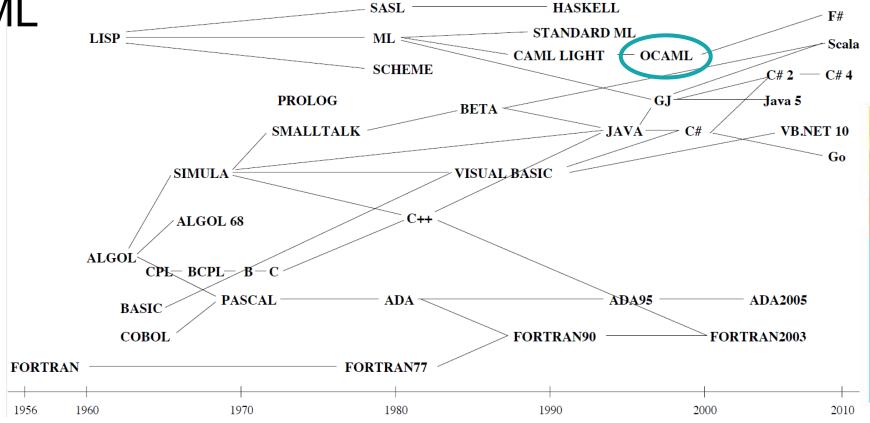
A SMALL SELECTION



## WHAT THIS TALK IS NOT ABOUT

- Dyalog APL
  - Most code at SimCorp is APL code
  - In use since the 1970's
  - Dynamically typed, interpreted, declarative, array-oriented
- C#
  - Most new development on .NET
  - Threading, services, GUI
- F#
  - Some components of SimCorp Dimension
  - SimCorp internal static APL type checker

## **OCAML**



Peter Sestoft

Programming Language Concepts
Second Edition

- Multi-paradigm: functional, imperative, object-oriented
- Static types and type inference
- Strict evaluation
- Interactive top-level (REPL)

## INDUCTIVE TYPES STRUCTURE PROGRAMS

```
let rec eval : expr -> int = function
type expr =
                                                  | Cst i -> i
  | Cst of int
  | Add of expr * expr
                                                  | Add (e1, e2) -> eval e1 + eval e2
  | Mul of expr * expr
                                                  | Mul (e1, e2) -> eval e1 * eval e2
  | Neg of expr
                                                  | Neg e -> -(eval e)
           # let e1 = Add (Cst 1, Neg (Cst 1));;
           val e1 : expr = Add (Cst 1, Neg (Cst 1))
           # let e2 = Add ("one", Cst 1);;
           Error: This expression has type string but an expression was expected of
           type expr
           # eval e1;;
           -: int = 0
```

## Composing contracts: an adventure in financial engineering

Functional pearl

Simon Peyton Jones
Microsoft Research, Cambridge
simonpj@microsoft.com

Jean-Marc Eber
LexiFi Technologies, Paris
jeanmarc.eber@lexifi.com

Julian Seward
University of Glasgow
v-sewardj@microsoft.com

#### 23rd August 2000

#### Abstract

Financial and insurance contracts do not sound like promising territory for functional programming and formal semantics, but in fact we have discovered that insights from pro-

At this point, any red-blooded functional programmer should start to foam at the mouth, yelling "build a combinator library". And indeed, that turns out to be not only possible, but tremendously beneficial.

The finance industry has an enormous vocabulary of jargon

## FINANCIAL CONTRACTS MODELLING

- Two parties agree on an amortized loan.
- Pay 30000 DKK up front.
- Pay back over three years with interest.
- Decide when payments are due.

```
let amortized_loan =
  let principal = cst 30000. in
  let coupon = cst 11000. in
  all [ give (flow 2019-01-01 DKK principal);
     flow 2020-01-01 DKK coupon;
     flow 2021-01-01 DKK coupon;
     flow 2022-01-01 DKK coupon ]
```

## A COMBINATOR LIBRARY FOR FINANCIAL CONTRACTS

#### OCAML DIALECT "LEXIFI OCAML"

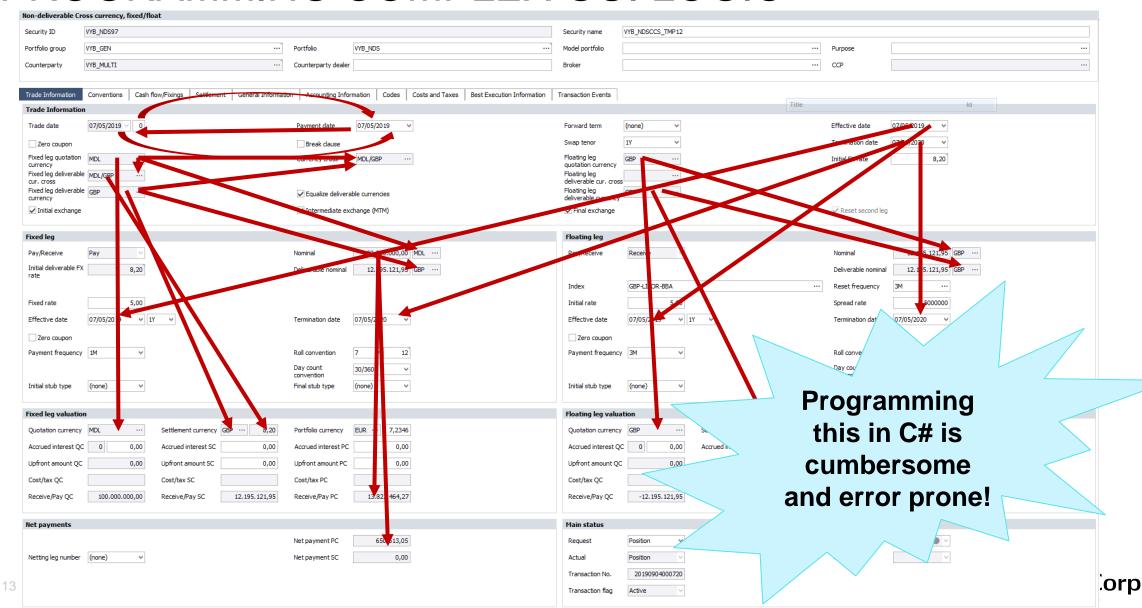
```
type binop = Add | Sub | Mul | Div | Max
                                                       let flow t cur obs =
type obs =
                                                         Acquire (t, Scale (obs, One cur))
   Underlying of string
    Const of float
                                                       let amortized loan =
    Binop of obs * binop * obs
                                                         let principal = Const 30000. in
                                                         let coupon = Const 11000. in
   Fixed of date * obs
                                                         all [ Give (flow 2019-01-01 DKK principal);
                           Fix an underlying market rate to
type contract =
                                                               flow 2020-01-01 DKK coupon;
                              its value at some date.
                                                               flow 2021-01-01 DKK coupon;
   One of currency
    Scale of obs * contract
                                                               flow 2022-01-01 DKK coupon ]
    Acquire of date * contract
    All of contract list
   Give of contract
    Either of contract * contract
                                                                     Evaluates to a value of type
   Anytime of date * date * contract * contract
                                  Exercise contract in between
```

dates, or exercise other contract after.

contract that we can valuate.



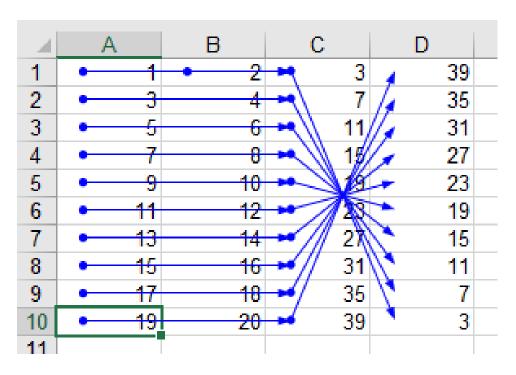
## PROGRAMMING COMPLEX GUI LOGIC



## LOOKS FAMILIAR?

#### SPREADSHEET-LIKE EVALUATION MODEL

- Fields depend on each other
  - When the user updates a field, all depending fields must be updated, too.
  - "Reactive programming"
  - "Self-adjusting computation"
- Free of side effects
  - From a programmer's point of view!
- Programming challenge for domain experts
  - What happens if the user changes the number of settlement days to 5?
- Solution: pure, type-safe, declarative programming



## Typelets — A Rule-Based Evaluation Model for Dynamic, Statically Typed User Interfaces

Martin Elsman<sup>1</sup> and Anders Schack-Nielsen<sup>2</sup>

- <sup>1</sup> University of Copenhagen, Universitetsparken 5, DK-2100 Copenhagen, Denmark mael@diku.dk
  - SimCorp, Weidekampsgade 16, DK-2300 Copenhagen, Denmark anders.schack-nielsen@simcorp.com

Abstract. We present the concept of typelets, a specification technique for dynamic graphical user interfaces (GUIs) based on types. The technique is implemented in a dialect of ML, called MLFi,<sup>3</sup> which supports dynamic types, for migrating type-level information into the object level, so-called type properties, allowing easy specification of, for instance, GUI control attributes, and type paths, which allows for type-safe access to type components at runtime. Through the use of Hindley-Milner style type-inference in MLFi, the features allow for type-level programming of user interfaces. The dynamic behavior of typelets are specified us-



## RULES DESCRIBE BUSINESS LOGIC

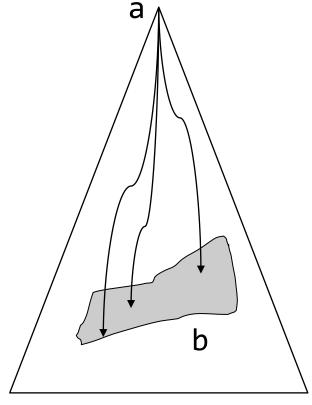
#### DECLARATIVELY COMPUTE FACTORIAL

```
type t = {
  number: int;
                                          Use record types to declare fields in the
  result: (int [@t readonly])
                                                     business logic.
let rule =
  let rec fact = function
                                               Define "normal" OCaml function.
      0 -> 1
      n -> n * fact (n - 1)
  in
  Rule.update
                                            Use it as projection in "update" rule.
    (Fields.value [%p number])
    (Fields.value [%p result])
    fact
                                            Box Factorial
let layout =
                                             Number
                                                                             Result
                                                                                                     40320
  let open Layout in
  box "Factorial"
      (lpick [%p number] % lpick [%p result])
```

## ACCESSING FIELDS THROUGH FIELD API

#### INTERNAL API FOR DEVELOPING BUSINESS LOGIC

```
module type FIELDS = sig
  type ('i, 'a) t (* 'i : type of the root *)
               (* 'a : type of elements pointed to *)
 val const : 'a ttype -> 'a -> ('i, 'a) t
  val value : ('i, 'a) path -> ('i, 'a) t
 val enabled : ('i, _) path -> ('i, bool) t
  val readonly : ('i, _) path -> ('i, bool) t
  val restrict : ('i, 'a) path -> ('i, 'a list) t
 val ( & ) : ('i, 'a) t -> ('i, 'b) t -> ('i, 'a * 'b) t
end
type vec2d = { x : float; y : float }
type vec3d = { x : vedc2; z : float }
```



```
let f = Fields.value [%p x.x] & Fields.value [%p x.y] & Fields.value [%p .z]
```

## What is the type of f?

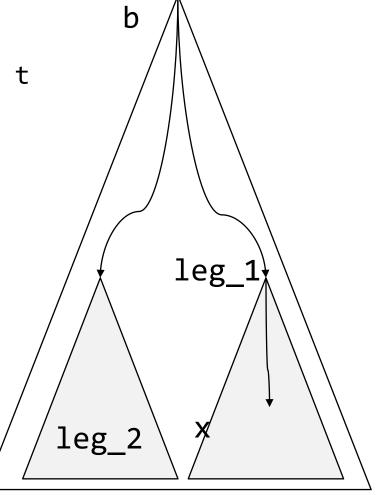
## PATHS COMPOSE

#### ALLOWS TO DIFFERENTIATE BETWEEN REPLICATED STRUCTUES

```
module type PATHS = sig
  val subpath : ('a, 'b) t -> ('b, 'c) t -> ('a, 'c) t
end
```

```
type a = { x : int; y : string }
type b = { leg_1 : a; leg_2 : a }
```

let x\_leg\_1 : (b, int) t = subpath [%p leg\_1] [%p x]



## CONSTRUCTING INSTRUMENTS THROUGH RULE API

#### INTERNAL API FOR DEVELOPING BUSINESS LOGIC

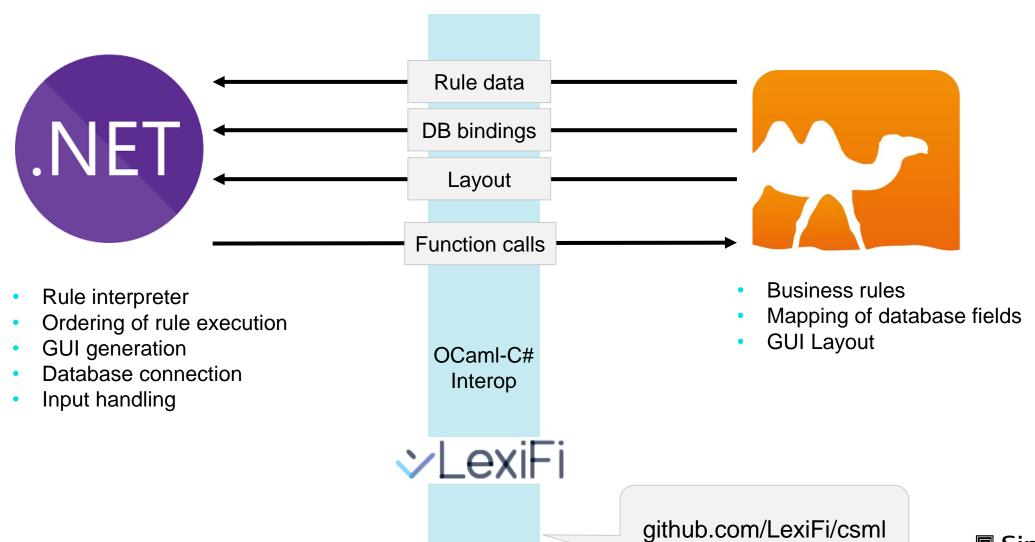
```
module type RULE = sig
  type 'i t
  type ('i, 'a) fields = ('i, 'a) Fields.t
  val update : ('i,'a)fields -> ('i,'b)fields -> ('a -> 'b) -> 'i t
  val validate : ('i,'a)fields -> ('a -> string option) -> 'i t
  val button : ('i,'a)fields -> ('i,'b)fields -> ('a -> 'b) -> ('i,unit)path -> 'i t
  val subpath : ('i,'a)path -> 'a t -> 'i t
  val all : 'i t list -> 'i t
               : ('i, 'a)fields -> ('i, 'b)fields -> ('a -> 'b) -> ('b -> 'a) -> 'i t
  val iso
  . . .
                                                                Check whether some
end
                                                               property holds for given
                                     Lift rule of type 'a t into
                                                                      fields.
                                    context of 'i if path from 'i
   Isomorphism between
                                           to 'a exists.
        two fields.
```

## What does function all: 'i t list -> 'i t do?



## RULE EXECUTION

#### INTERPETER IMPLEMENTED IN .NET - BUSINESS RULES IN OCAML



# STOP: DEMO TIME!



#### MONADIC RULE API

#### SPLITTING RULES INTO SMALLER STEPS

- OCaml run-time is not reentrant
- Data base access & calls to APL are long-running and blocking.
- Solution: split rule computation into small chunks compute step-wise.

```
module Rules : sig =
          val update m : ('i,'a)fields -> ('i,'b)fields -> ('a -> 'b m) -> 'i t
          val update : ('i,'a)fields -> ('i,'b)fields -> ('a -> 'b) -> 'i t
        end = struct
          let update src tgt f =
                                                              val return : 'a -> 'a m
            update_m src tgt (fun x -> return (f x))_
                                                              val abort : unit -> 'a m
        end
                               update m
                                 (value [%p portfolio_ik])
                                 (value [%p currency_ik])
(value [%p x])
                                 (function
(value [%p w])
                                      0 -> abort ()
(call_apl_1 >>= fun (_, y) ->
                                      por ik -> dbhandle >>= fun h ->
return y >>=
                                      match select1 h pCURIK (all [pPORIK =. por_ik]) with
call apl 2)
                                          [cur ik] -> return cur ik
```

-> abort ())

Monadic type 'a m!

update m

## EVALUATING RULES AFTER A USER CHANGE

#### TWO LOOPS TO RULE THE RULES

```
void EvalRules() {
  while (Rq.HasNext) {
    RuleInfo r = Rq.Pop();
    Value inp = Env.GetValue(r.Source);
    Value res = EvalRule(r.Run(inp));
    Env.SetValue(r.Target, res);
}}
    Takes first step
    SetValue() also
```

updates Rq.

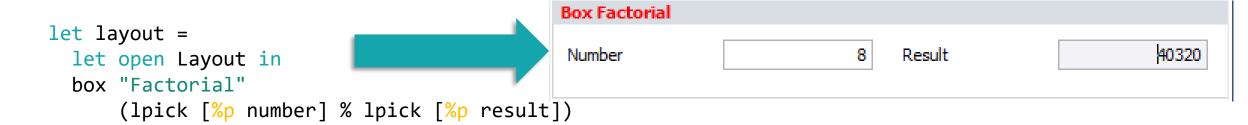
- Iterative algorithm on a queue of rules
- A rule is enqueued when one of its source fields is written to.
- Rules are evaluated step-wise to enable interleaving
  - Two clients:
  - One runs a C#-rule
  - The other runs an OCaml rule
  - Minimal mutual exclusion
- Rule programmer does not really care.

## SEPARATION OF CONCERNS

#### **GUI LAYOUT GENERATION VIA LAYOUT API**

```
type base = Number | Button | CheckBox | TextBox | Date | DropDown

type 'p t =
    | Pick of access_path * string option * base
    | Hseq of 'p t * 'p t
    | Vseq of 'p t * 'p t
    | Halign of halign * 'p t
    | Valign of valign * 'p t
    | Text of string
    | NamedButton of string
    | Box of string * 'p t
```



## HOW WELL DOES THIS WORK IN PRACTICE?

#### HIGH RE-USE, FAST TIME-TO-MARKET

- Plug-in system using OCaml Functors:
  - A system to generate modules from other modules.
  - Highly composable, type safe, no run-time overhead.
  - Rule composition possible thanks to purity!

- Many business rules are generic!
  - E.g. business calendar functionality.
  - Financial instruments differ only in few places.

Re-use factor for business rules is ~10

- 274'845 lines of OCaml code
- 428'771 business rules in production
- 42'132 unique business rules
- Over 100 financial instruments

## SUMMARY

#### FUNCTIONAL PROGRAMMING AT SIMCORP

Combinator library for modelling financial contracts

Every-day business for us, revolutionary for the finance sector.

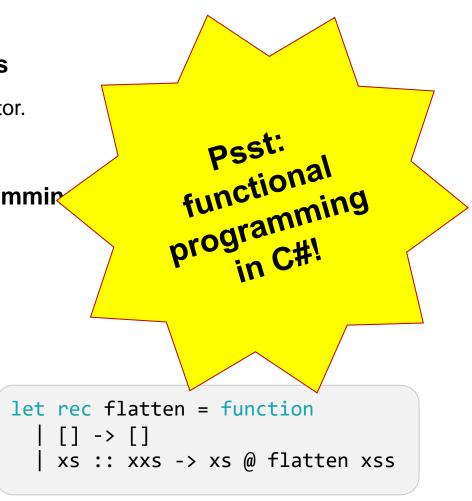
Domain experts model contracts.

Declarative business logic for type-safe GUI programming

- Similar to constructing a spreadsheet.
- Focus on what, not how.

#### Challenges ahead

- You gotta know your stuff:
  - Polymorphism, existential types, phantom types
  - Monads & API design
  - Compositionality & catamorphisms



# Thank you!

Florian.Biermann@SimCorp.com



#### LEGAL NOTICE

The contents of this publication are for general information and illustrative purposes only and are used at the reader's own risk. SimCorp uses all reasonable endeavors to ensure the accuracy of the information. However, SimCorp does not guarantee or warrant the accuracy, completeness, factual correctness, or reliability of any information in this publication and does not accept liability for errors, omissions, inaccuracies, or typographical errors. The views and opinions expressed in this publication are not necessarily those of SimCorp A/S. All rights reserved. Without limiting rights under copyright, no part of this document may be reproduced, stored in, or introduced into a retrieval system, or transmitted in any form, by any means (electronic, mechanical, photocopying, recording, or otherwise), or for any purpose without the express written permission of SimCorp A/S. SimCorp, the SimCorp logo, SimCorp Services are either registered trademarks or trademarks or trademarks of SimCorp A/S in Denmark and/or other countries. Refer to www.simcorp.com/trademarks for a full list of SimCorp A/S trademarks. Other trademarks referred to in this document are the property of their respective owners.