Foundations of Distributed Systems

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Motivation

- Handle a large number of clients and requests with a single server
- The "c10k problem" in 1999:
 - http://www.kegel.com/c10k.html
- Examples:
 - financial, games, ...
 - notifications in mobile apps
 - machine-to-machine (M2M)

Case study

- Simple chat server:
 - Forward all messages to all clients
- Consider:
 - Large number of clients
 - Slow connections



First threaded solution

- For each connection:
 - Handler thread
- When reading, write to all other connections
- Use buffering:
 - To minimize system calls
 - To cope with slow readers

Sockets in java.net

```
ServerSocket ss=new ServerSocket(12345);
while(true) {
    Socket s=ss.accept();
    InputStream is=s.getInputStream();
    OutputStream os=s.getOutputStream();
    // i/o
    s.close();
}
```

Buffers in java.net

```
ServerSocket ss=new ServerSocket(12345);
while(true) {
   Socket s=ss.accept();
   InputStream is=new BufferedInputStream(s.getInputStream());
   OutputStream os=new BufferedOutputStream(s.getOutputStream());
   // i/o
   os.flush();
   s.close();
                            Needed to
                          actually write!
```

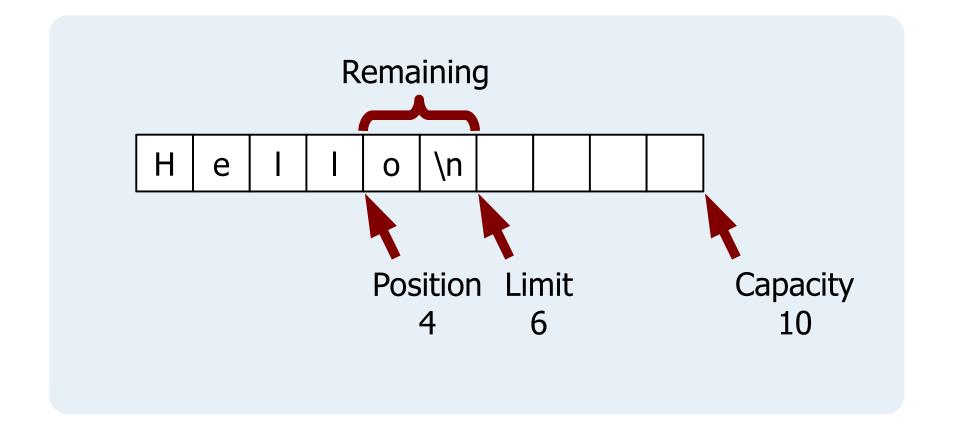
Memory

- Memory used: n connections x messages in transit (~ n²)
 - Caused by data copying in stacked abstractions
 - Serialization!
 - Overhead in allocation and garbage collection
- Solution: Store transient data in reusable shared buffers
 - Pointers/indexes into statically allocated data

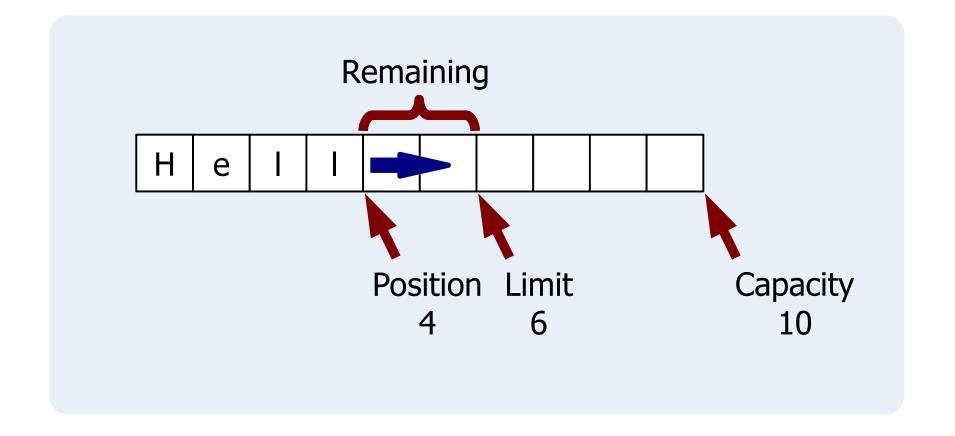
Sockets in java.nio

```
ServerSocketChannel ss=ServerSocketChannel.open();
ss.bind(new InetSocketAddress(12345));
while(true) {
    SocketChannel s=ss.accept();
    // i/o
    s.close();
}
```

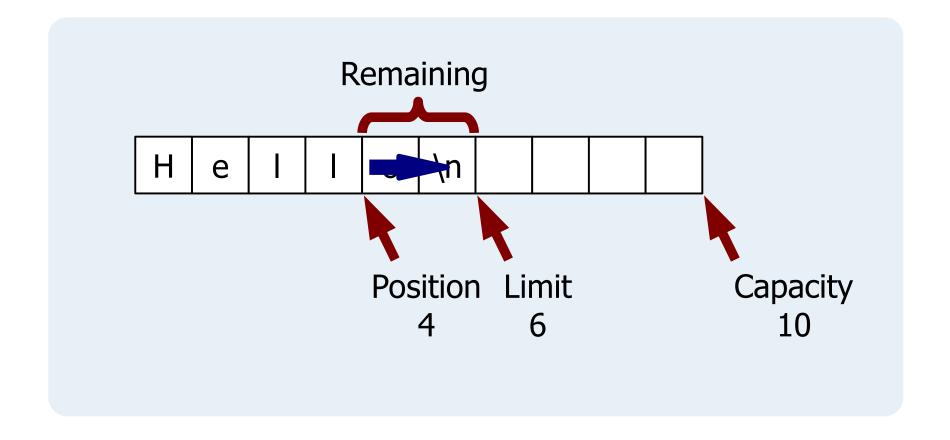
Buffer = Array + Indexes:



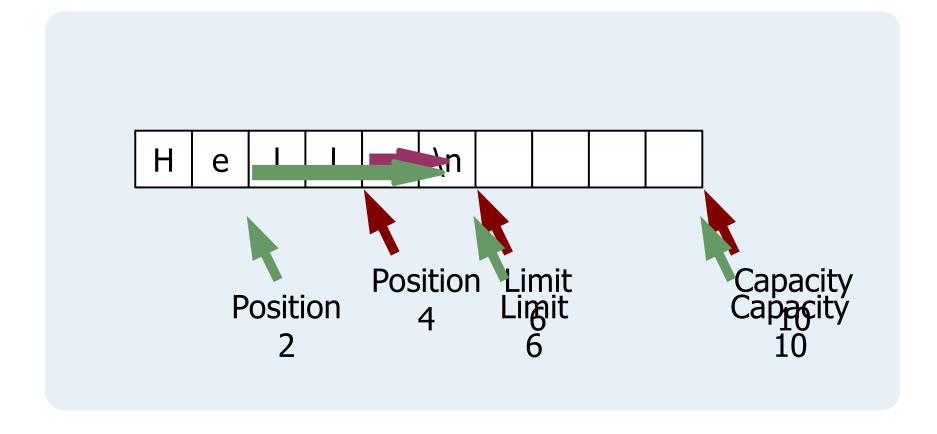
Put/read: advances position, sets content

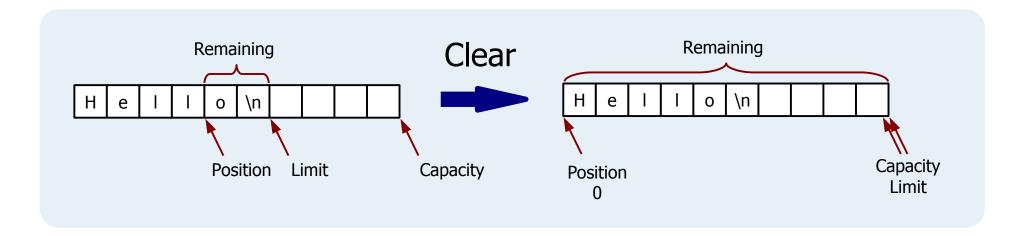


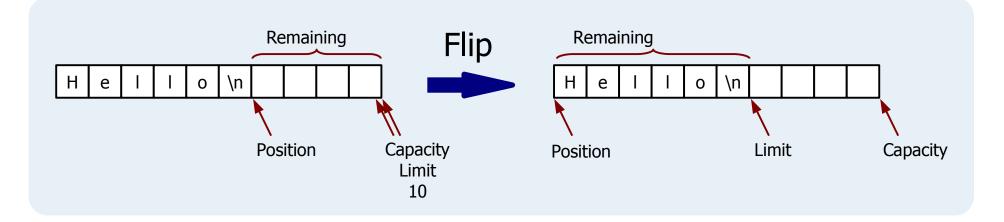
Get/write: advances position, gets content



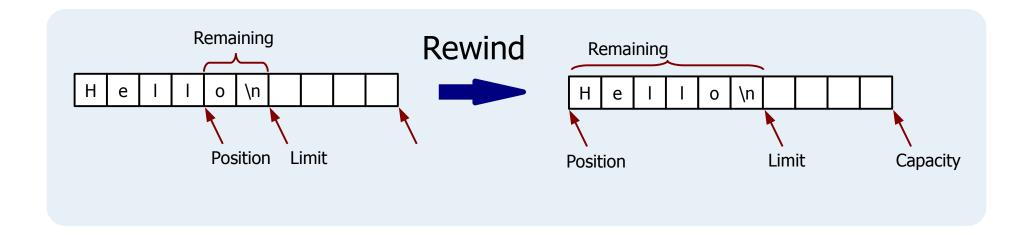
 Duplicate and wrap: multiple pointers into the same array

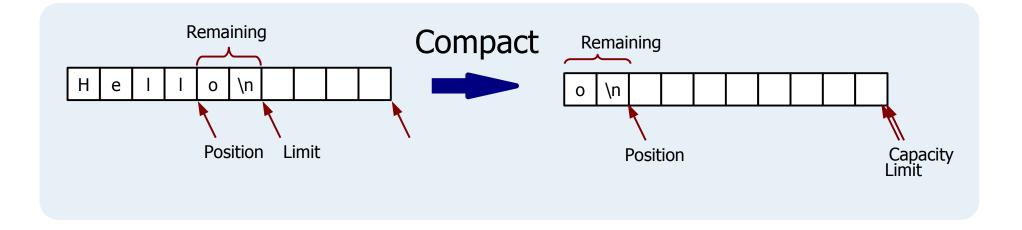






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Shared buffers

- Memory used: messages in transit (~ n)
- Ideally, never allocate or dispose of memory in normal operation:
 - No overhead, but...
 - Needs reference counting to know when to reuse

Flushing buffers

```
ServerSocket ss=new ServerSocket(12345);
while(true) {
   Socket s=ss.accept();
   InputStream is=new BufferedInputStream(s.getInputStream());
   OutputStream os=new BufferedOutputStream(s.getOutputStream());
   // i/o
   os.flush();
   s.close();
                              What if
                          write blocks?
```

Second threaded solution

- For each connection:
 - Reader thread
 - Writer thread
 - Pending queue
- When reading, insert in queues and notify writer threads
- When writing, remove from queue and notify reader threads

Threads

Problems:

- Memory overhead (stacks)
- Context switches and lock contention
- "Thundering herd", hidden queue, and fairness

Blocking sockets

```
try {
    ByteBuffer buf=ByteBuffer.allocate(100);

s.read(buf);
buf.flip();

r.write(buf);
} catch(IOException e) {
    report(e);
}
```

Asynchronous sockets

```
ByteBuffer buf=ByteBuffer.allocate(100);
s.read(buf, null, new CompletionHandler() {
    public void completed(Integer result, Object a) {
        buf.flip();

        r.write(buf, ...);
    }
    public void failure(Throwable t, Object a) {
        report(t);
    }
);
```

Thread pools

- For non-blocking, short-lived events:
 - One pool thread for hardware thread
- While all threads are blocked, the application stops handling events

```
AsynchronousChannelGroup g =
    AsynchronousChannelGroup.withFixedThreadPool(...);

AsynchronousSocketChannel s =
    AsynchronousSocketChannel.open(g);

... /* callbacks use g.shutdown() to exit */
g.awaitTermination(Long.MAX_VALUE, TimeUnit.SECONDS);
```