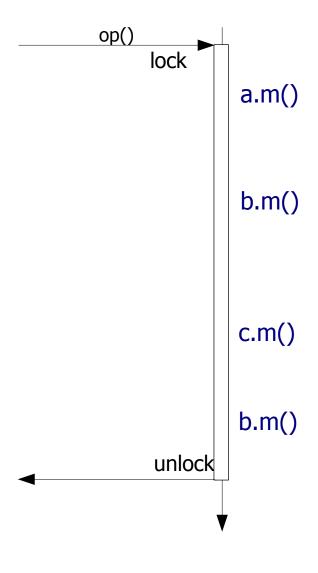
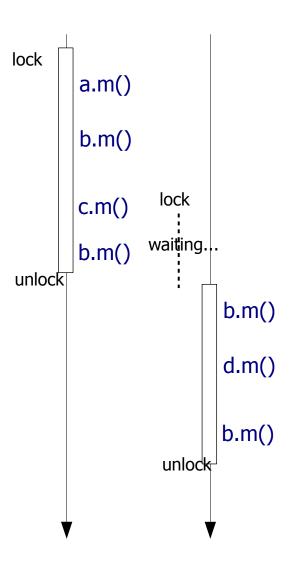
### Challenges

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- Atomic operations on multiple distributed objects
  - Concurrent client invocations
  - Client and server failure

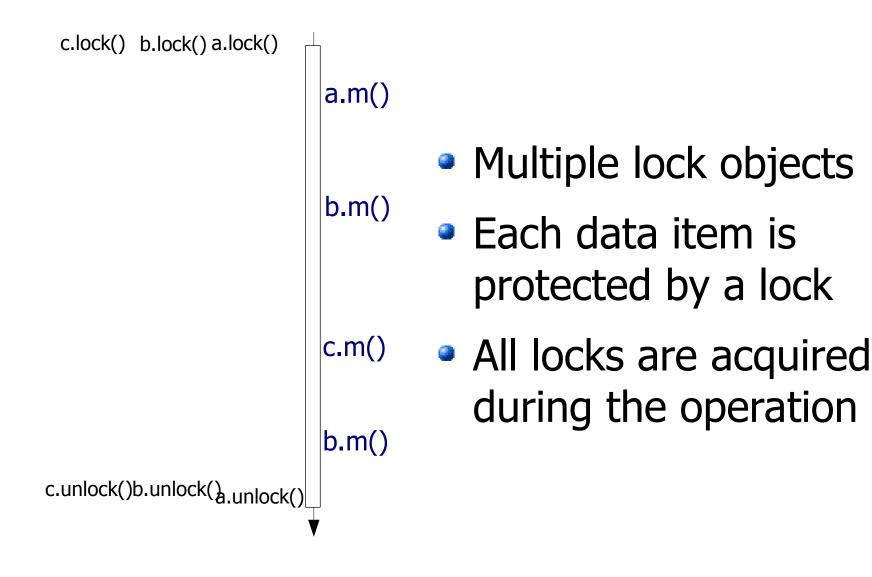


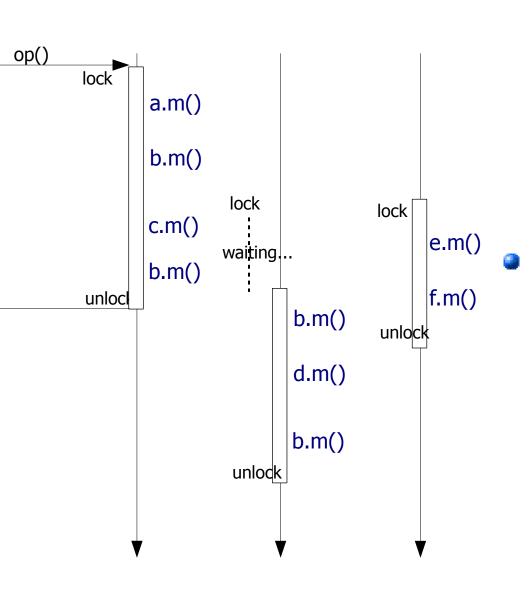
- Lock/unlock operations in a lock server
- A single global lock protects all data items
- Acquired for the duration of the composite operation



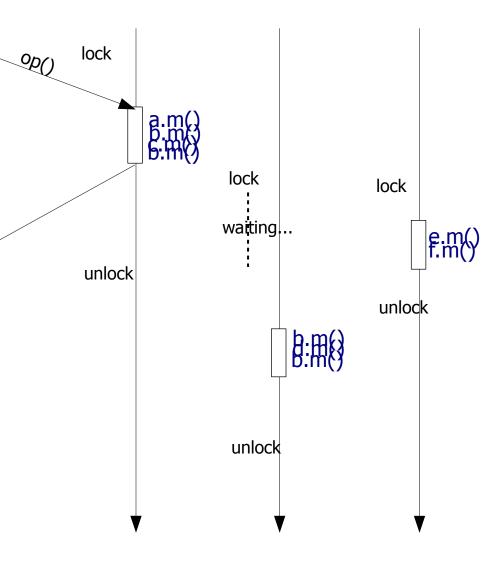
- Operations never overlap
- History is a sequence of operations

- Performance bottleneck:
  - At most one operation executing
- Remote invocations needed for lock / unlock
- Single point of failure:
  - Lock server fails
  - Lock client fails without unlock()





Operations on different items may overlap

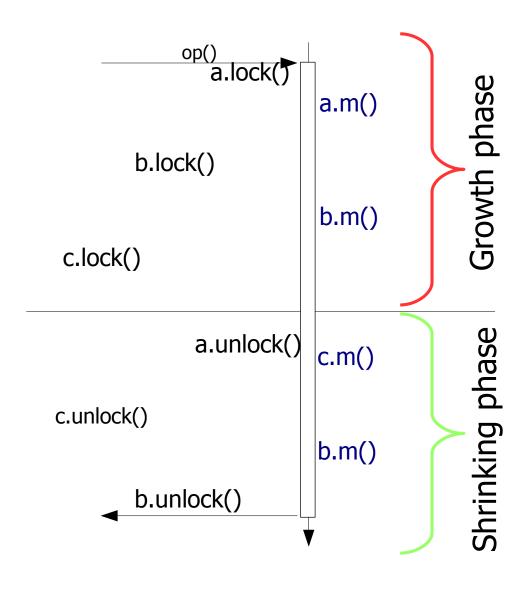


Indistinguishable from:

- Instantaneous execution with delay on round-trip
- Clients cannot deny that operations do not overlap

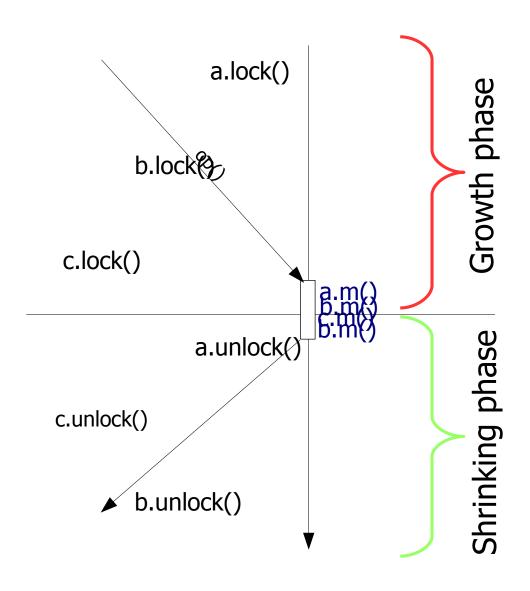
- Performance bottleneck
- Remote invocations needed for lock / unlock
- Single point of failure:
  - Lock server fails
  - Lock client fails without unlock()
- Must know all items accessed on start
- Deadlocks

# 2-Phase locking



- All lock requests precede all unlock requests
- Lock can be implicit in operation

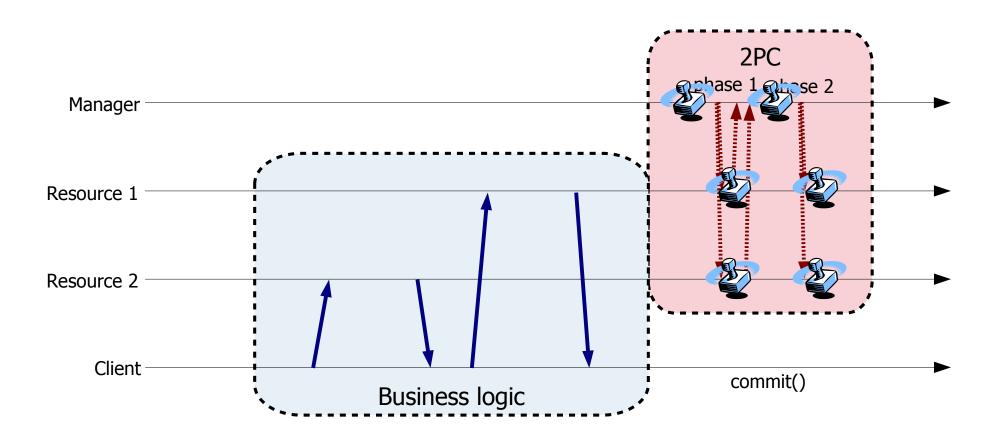
### 2-Phase locking



 An observer cannot deny that all happens in-between phases

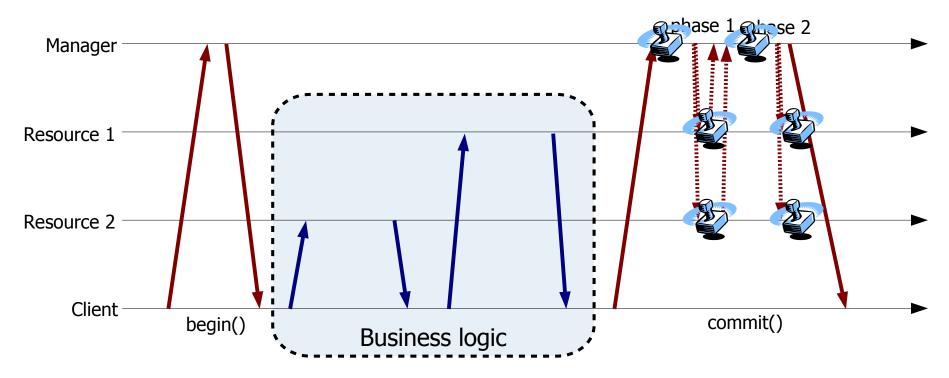
- Performance bottleneck
- Remote invocations needed for <del>lock</del> / unlock
  - Lock operation is implicit
- Single point of failure:
  - Lock server fails
  - Client fails without unlock()
- Must know all items accessed on start
- Deadlocks

# **Application vs Transactions**



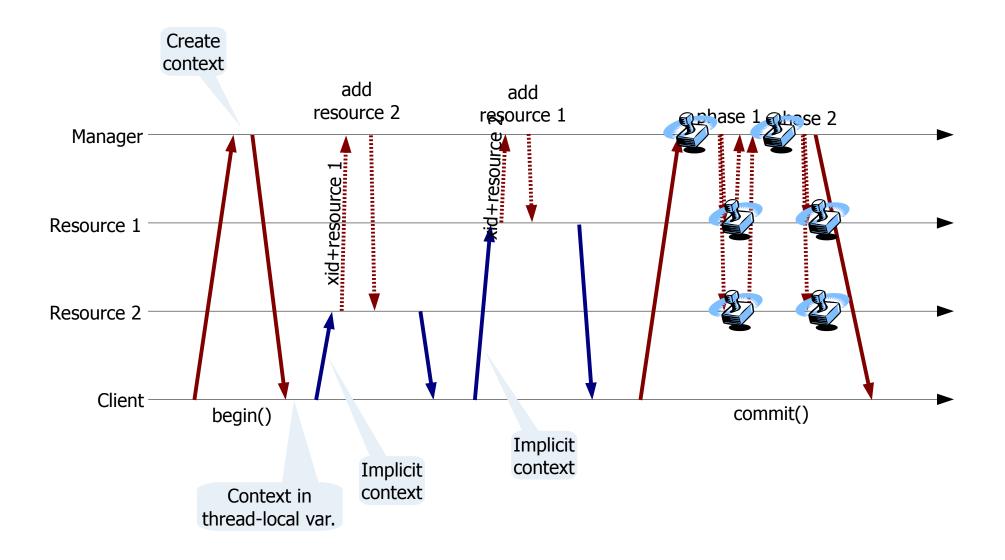
- How to connect them?
  - Don't change business logic

#### Transaction demarcation

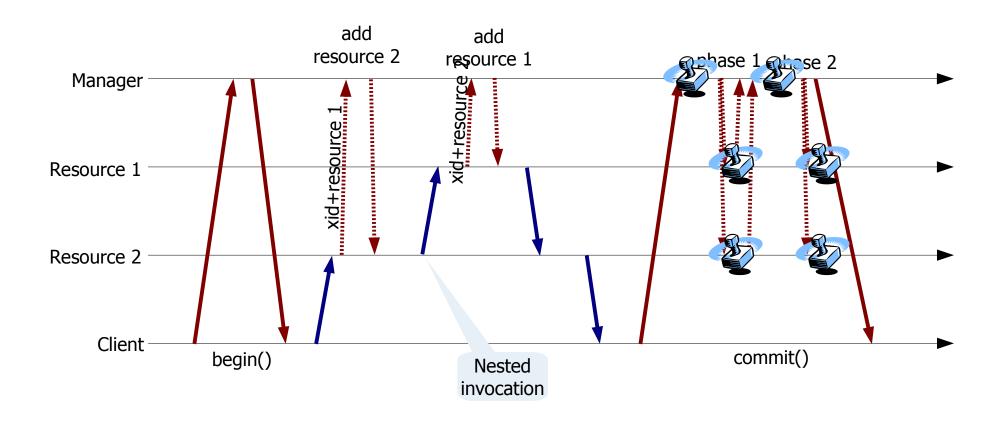


- Call transaction manager before/after
- How does manager discover participants?

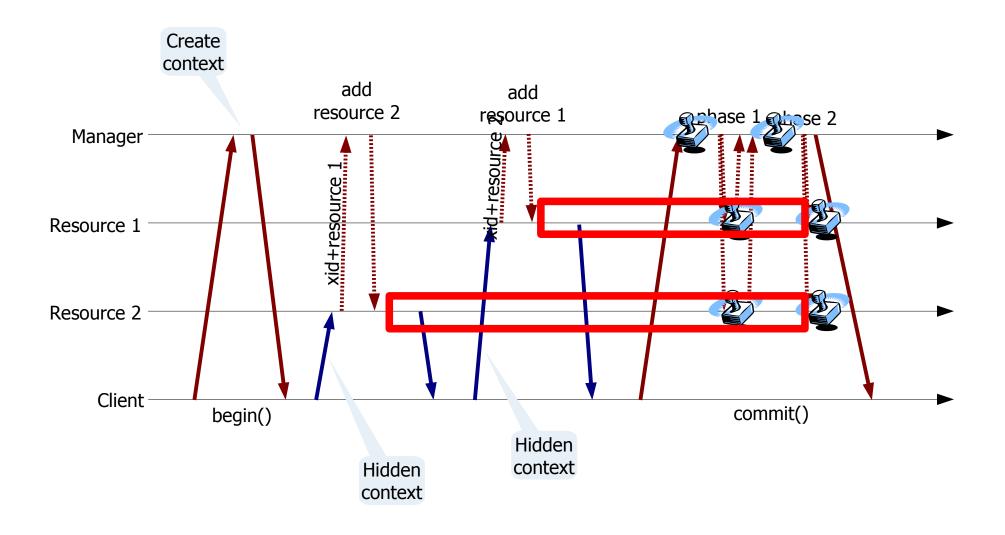
# Transactional RPC + 2PC



# Transactional RPC + 2PC



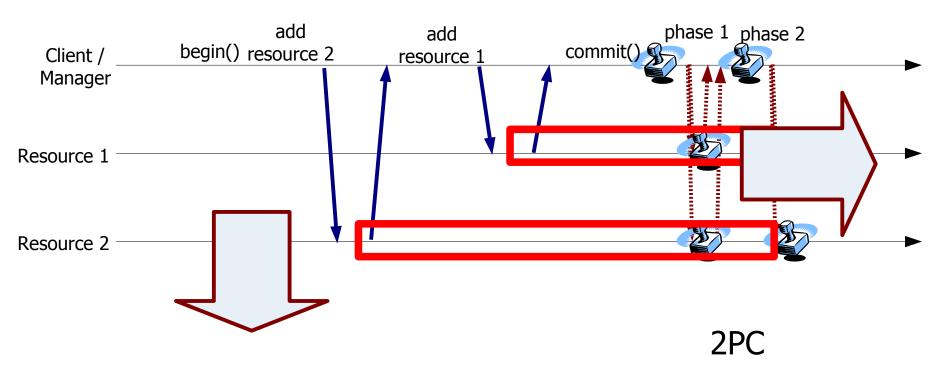
# Transactional RPC + 2PC + 2PL



- Performance bottleneck
- Remote invocations needed for lock / unlock
  - Lock operation is implicit
  - Unlock operation is implicit (txn. commit)
- Single point of failure:
  - Lock server fails
  - Lock client fails without unlock() (rollback)
- Must know all items accessed on start
- Deadlocks (rollback)

- Atomicity with faults: 2PC
- Atomicity with concurrent clients: 2PL
- With 2PL + 2PC:
  - Rollback on deadlock / client failure
  - Implicitly release locks on commit / rollback

### 2PC+2PL: Challenges

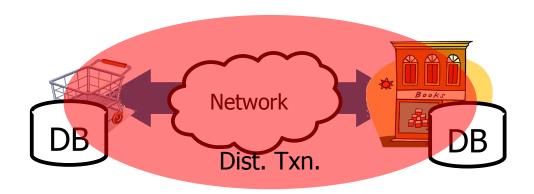


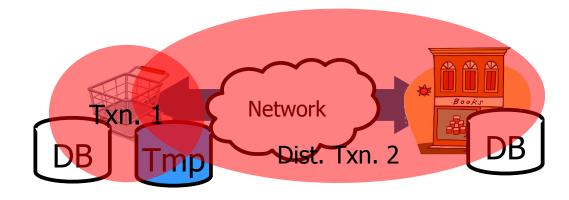
Scalability with:

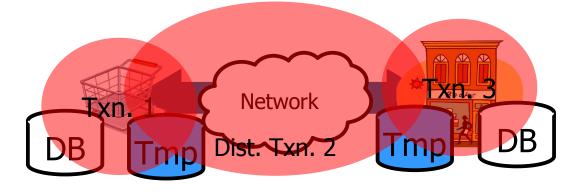
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- More resources?
- Slower 2PC (remote resources)?

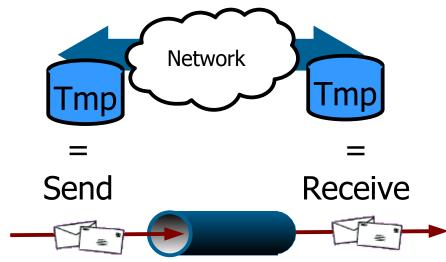
#### Store and Forward







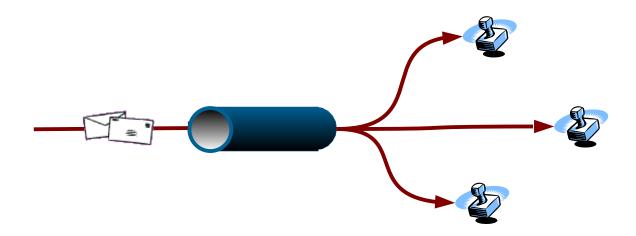
#### Transactional queues



- Insert and remove from temporary DBs abstracted as transactional send/receive
- Advantages:
  - Performance and fault decoupling
  - Asynchronous (pull) operations

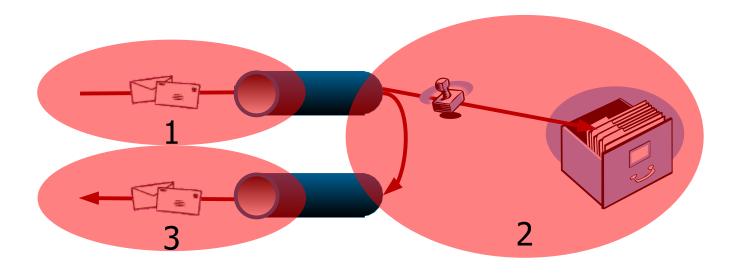
#### Server cluster

- Multiple servers polling the same queue:
  - Each message processed exactly once
  - No order enforced



# Request/Reply

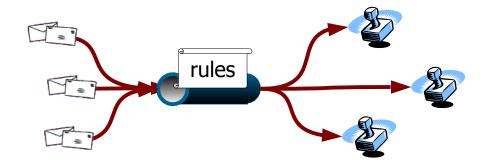
- Client/server:
  - Server polls for requests
  - Client polls for replies
- Multiple transaction contexts:



### Limitations of queues

- Consider adding another application...
- Addressing is explicit in applications:
  - Each sender must know exactly which receivers are interested in a message

### Message brokers



- Goal: Decouple message routing from sending and receiving
- Receivers subscribe to messages that fit some criteria
- Senders publish messages

# Topic-based routing

- Publisher tags messages with a topic id
- Subscriber requests a topic
- Options:
  - Flat topics
  - Hierarchical topics

### Content-based routing

- Publisher structures message content using a common structure
- Subscriber registers a content filter
- Options:
  - Relational data / SQL WHERE clauses
  - XML / XQuery

- Transactional messaging decouples:
  - Faults
  - Performance
  - Application logic
- Two options:
  - Tightly coupled applications: 2PL + 2PC + TRPC
  - Loosely coupled: 2PC + MQ