#### **DataBase Storage**

#### **Type Representation**

使用FLOAT,REAL/DOUBLE与使用NUMERIC,DECIMAL,VARCHAR的区别

FLOAT是根据IEEE-754 Standard用二进制来对浮点数进行编码,上过ICS都知道其实这种方式表示浮点数是有误差的。因此在一些不能容忍误差的场合如金融统计等精度要求极高的场景,不能使用float/double,而应该使用numeric/decimal。当然使用float/double速度会快很多

#### **Data Layout / Alignment**

• 如何表示NULL DATA TYPES IN DB?

choice 1: Special Values

如用INT\_MIN来表示NULL

choice 2: Null Column Bitmap Header

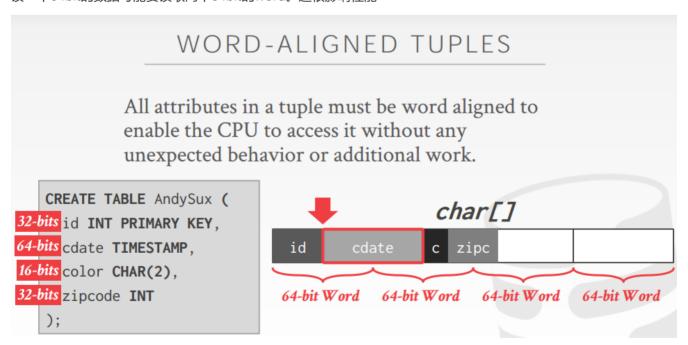
在tuple的header存一个bitmap,用单个bit来表示哪些属性是null

choice 3: Per Attribute Null Flag

用一个Flag来表示Null,但是会比较浪费存储空间,并且会有word alignment问题

• Unalignment data reading issue

下图所示:如果要读取cdate这个TS,由于其不是对齐存放的,跨越了第一个64bit word和第二个64bit word,导致读一个64bit的数据可能要读取两个64bit的word。这很影响性能



Possible Solution:

填充一些0字符,使得强制对齐。

### WORD-ALIGNMENT: PADDING

Add empty bits after attributes to ensure that tuple is word aligned.

```
CREATE TABLE AndySux (

32-bits id INT PRIMARY KEY,

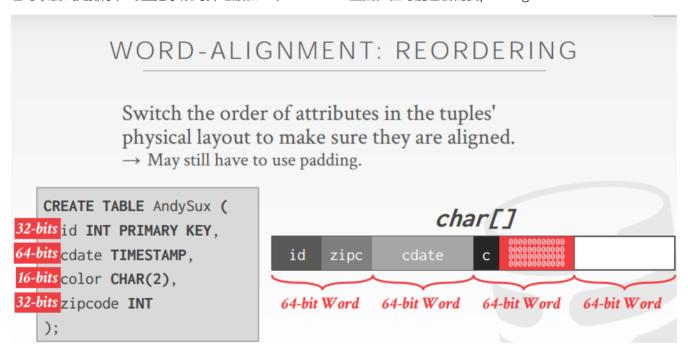
64-bits cdate TIMESTAMP,

16-bits color CHAR(2),

32-bits zipcode INT

);
```

也可以尝试先排序,尽量把字段对齐地放在一个64bit word里面。但可能也会需要padding



测试结果:可以看到如果解决好字段存储的对齐问题,性能提高非常明显

## CMU-DB ALIGNMENT EXPERIMENT

Processor: 1 socket, 4 cores w/ 2×HT Workload: Insert Microbenchmark

# Avg. Throughput

No Alignment 0.523 MB/sec
Padding 11.7 MB/sec
Padding + Sorting 814.8 MB/sec

#### **Storage Models**

N-ARY STORAGE MODEL(NSM)

#### 行式存储

一个tuple中所有的属性都连续存储在一起。适合OLTP因为OLTP通常涉及少量的tuple及tuple的所有attributes.

# N-ARY STORAGE MODEL (NSM)

The DBMS stores all of the attributes for a single tuple contiguously.

Ideal for OLTP workloads where txns tend to operate only on an individual entity and insertheavy workloads.

Use the tuple-at-a-time iterator model.

但是不适合OLAP场景,因为OLAP通常是需要某几个列,而不是全部列,举个极端例子如果你只需要海量数据某一个列的数据来做分析,那么你需要把海量数据所有的tuple的所有attribute都加载出来。

### N-ARY STORAGE MODEL (NSM)

### Advantages

- → Fast inserts, updates, and deletes.
- → Good for queries that need the entire tuple.
- → Can use index-oriented physical storage.

### Disadvantages

- → Not good for scanning large portions of the table and/or a subset of the attributes.
- DECOMPOSITION STORAGE MODEL

列式存储,海量数据的某一列属性连续存储在一起,减少了把数据从索引加载到内存的waste work,因为你只需要读取你所用到的数据

OLAP场景是查询多, 但对于增删改会比较慢

The DBMS stores a single attribute for all tuples contiguously in a block of data.

Ideal for OLAP workloads where read-only queries perform large scans over a subset of the table's attributes.

#### 具体实现细节:

• Tuple Identification

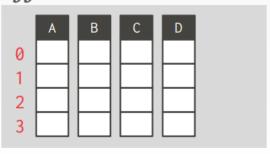
## Choice #1: Fixed-length Offsets

 $\rightarrow$  Each value is the same length for an attribute.

## Choice #2: Embedded Tuple Ids

 $\rightarrow$  Each value is stored with its tuple id in a column.

## Offsets



### Embedded Ids

	A		В		С		D
0		0		0		0	
1		1		1		1	
2		2		2		2	
3		3		3		3	

## DSM: DATA ORGANIZATION

## Choice #1: Insertion Order

→ Tuples are inserted into any free slot that is available in existing blocks.

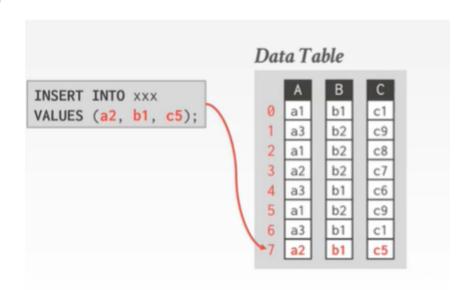
### Choice #2: Sorted Order

→ Tuples are inserted based into a slot according to some ordering scheme.

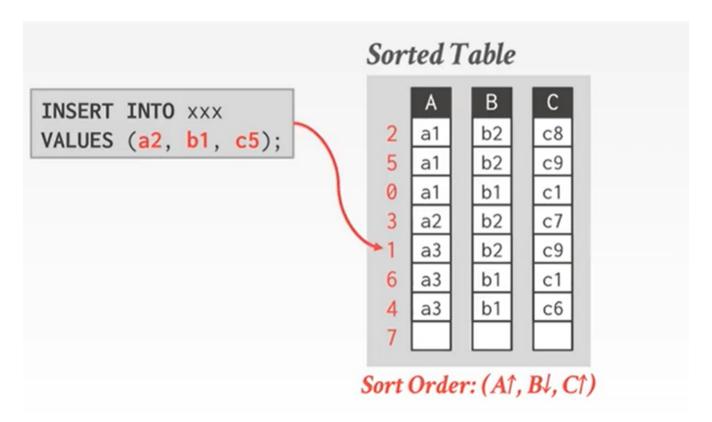
## Choice #3: Partitioned

→ Assign tuples to blocks according to their attribute values and some partitioning scheme (e.g., hashing, range).

#### **Insert Order**

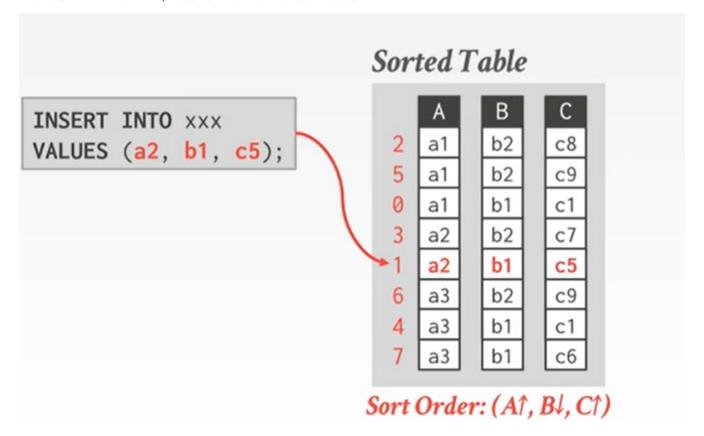


#### **Sorted Order**



插入的时候需要保证Sorted Table的全局有序

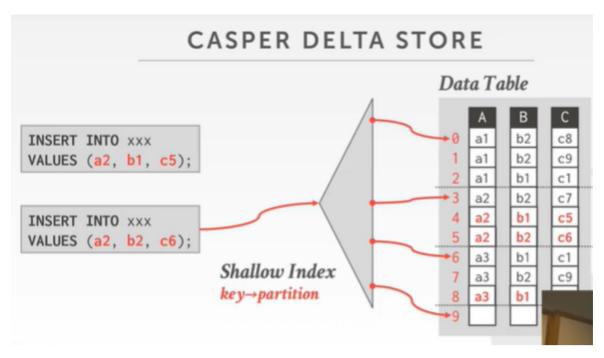
这个时候可能需要把该tuple插入位置及之后的数据全部后移



#### **Partitioned**

一个二级索引分别指向不同paritition(可以按照Key来分),插入数据的时候会定位到具体的partition.

那么寻找数据的时候只需要在这个数据的partition内查找,使得OLAP的查询效率得到提高。



但事实上我们可以结合两种storage model,因为可以把数据库的数据进行冷热划分:

hot data: 刚刚insert到DB的数据,很有可能最近还会被更新,适合增删改的OLTP,即行式存储。

cold data: 已经insert到DB很久的数据,被更新的几率已经不大,适合查的OLAP,即列式存储。

阿里云的AnalyticDB就是采用这种行列混存的形式存储冷热数据。

## HYBRID STORAGE MODEL

## Choice #1: Separate Execution Engines

→ Use separate execution engines that are optimized for either NSM or DSM databases.

### Choice #2: Single, Flexible Architecture

→ Use single execution engine that can efficiently operate on both NSM and DSM databases.

方案一: 对行列存储分别使用不同的执行引擎

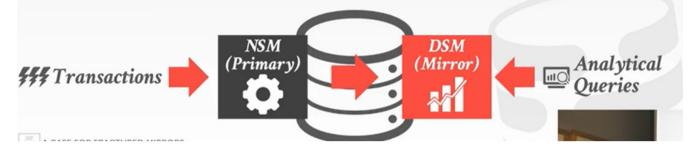
所有由事务插入到数据库的数据主要都是行式存储。

仅当分析性的Query来到要查询数据库数据的时候,要把行式存储的数据复制成列式存储的格式

### FRACTURED MIRRORS

Store a second copy of the database in a DSM layout that is automatically updated.

→ All updates are first entered in NSM then eventually copied into DSM mirror.



方案二: 只用一套执行引擎

关于delta store可以参考这篇https://aboutsqlserver.com/2014/05/06/clustered-columnstore-indexes-exploring-delta-store-and-delete-bitmap/

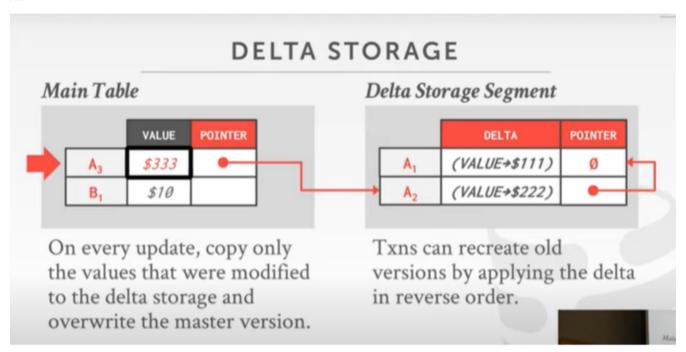
# **DELTA STORE**

Stage updates to the database in an NSM table.

A background thread migrates updates from delta store and applies them to DSM data.



回忆一下第三节课讲的Delta Storage: main table只存储最新版本的数据,然后通过指针把历史版本数据连接在一起。



#### **System Catalogs**

数据库的Catalog指的是什么?一个数据库系统会包括多个Catalog,每个Catalog下面又有多个Schema。每个Schema下面包含了DB的实例如table,view.可以理解为命名空间层面上用于进行租户隔离的机制。

#### **Schema Changes**

NSM:行存 DSM:列存

# SCHEMA CHANGES

## ADD COLUMN:

- → **NSM**: Copy tuples into new region in memory.
- → **DSM**: Just create the new column segment

### DROP COLUMN:

- → NSM #1: Copy tuples into new region of memory.
- → NSM #2: Mark column as "deprecated", clean up later.
- → **DSM**: Just drop the column and free memory.

## **CHANGE COLUMN:**

→ Check whether the conversion can happen. Depends on default values.

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