

Adversarial Classification on Social Networks

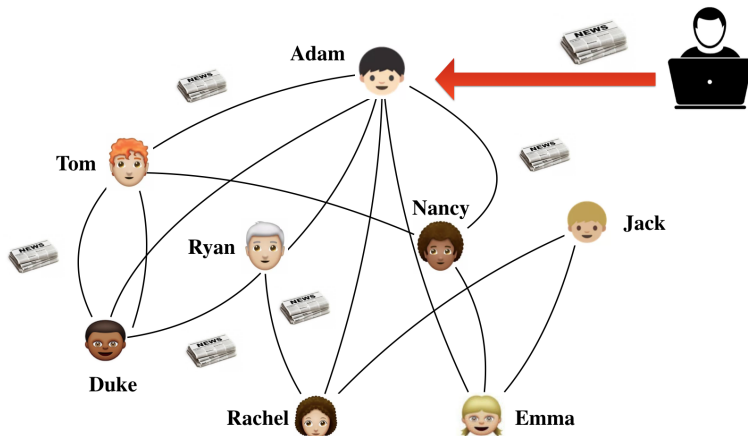
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Problem Setting



Motivation

- Over 50% adults in the U.S. regard social media as primary sources for news. [**holcomb2013news**].
- Over 37 million news stories in 2016 U.S. Presidential election later proved fake. [**allcott2017social**]
- Anti-social posts/discussions are negatively affecting users and damage online communities. [**cheng2015antisocial**]
- Social network spams and phishing can defraud users and spread malwares.

Traditional Defense

- Train a “global” detector from past data and deploy it everywhere.
- Ignore network structures, propagation of messages, and adversarial behavior.

Not Adequate

- Adversaries can tune content to avoid being detected.
- Traditional learning approaches ignore network structures.
 - The impact of detection errors.
 - Being able to detect malicious content at multiple nodes creates a degree of redundancy.

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- A node is affected if its shortest path to s is above T , which is externally supplied.
- The influence of a message initially affecting a node s is defined as $\sigma(s, x)$, which is the expected number of affected nodes over time window T .

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Innovations

- Learn and deploy *heterogeneous* detectors at different nodes.
- Explicitly considering both *propagation* of messages and *adversarial manipulation* during learning.

$$U_d = \alpha \sum_{x \in D^-} \sum_{i \in V} \sigma(i, \Theta, x) - (1 - \alpha) \sum_{x \in D^+} \sigma(s, \Theta, z(x)) \quad (1)$$

- D^- , D^+ are benign and malicious data, respectively.
- $\Theta = \{\theta_1, \theta_2, \dots, \theta_{|V|}\}$ being parameters of detectors at different nodes.
- The expected influence is now a function of the parameters of detectors (Θ), as well as manipulated messages ($z(x)$).
- $x \rightarrow z(x)$: adversarial manipulation.

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Attacker's actions

- Find a node $s \in V$ to start propagation (reminiscent of the famous influence maximization problem).
- Transform $x \rightarrow z(x)$ in order to avoid detection.

For any original malicious instance $x \in D^+$:

$$\begin{aligned} \max_{i,z} \quad & \sigma(i, \Theta, z) \\ \text{s.t.} \quad & \|z - x\|_p \leq \epsilon \\ & \mathbb{1}[\theta_j(z) = 1] = 0, \forall j \in V \end{aligned} \tag{2}$$

- ϵ : the attacker's budget.
- $\theta_j(z) = 1$: the manipulated message is detected at node j .

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Stackelberg Game

The interaction between the **defender** and the **attacker** is modeled as a Stackelberg game. which proceeds as follow:

- The **defender** first learns Θ (the parameters of detectors at different nodes).
- The **attacker** observes Θ and construct its optimal attack against the defender.

$$\begin{aligned} \max_{\Theta} \quad & \alpha \sum_{x \in D^-} \sum_i \sigma(i, \Theta, x) - (1 - \alpha) \sum_{x \in D^+} \sigma(s, \Theta, z(x)) \\ \text{s.t. :} \quad & \forall x \in D^+ : \quad (s, z(x)) \in \arg \max_{j, z} \sigma(j, \Theta, z) \\ & \forall x \in D^+ : \quad \|z(x) - x\|_p \leq \epsilon \\ & \forall x \in D^+ : \quad \mathbb{1}[\theta_k(x) = 1] = 0, \forall k \in V \end{aligned}$$

The equilibrium of this game: $(\Theta, s(\Theta), z(x; \Theta))$.

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Solution Approach

Assumption

The defender *knows* the node being attacked.

- This assumption enables us to collapse the bi-level optimization into a single-level optimization.
- Assume the defender knows the node s will be attacked, by leveraging *Implicit Function Theorem*, we can solve the single-level optimization, which results in the optimal defense strategy Θ_s^* .

Relax the assumption

We relax the assumption that the defender *knows* the node being attacked, and introduce a heuristic algorithm to solve for $(\Theta, s(\Theta), z(x; \Theta))$.

Heuristic algorithm:

- For each node $i \in V$ we solve for the Θ_i^* .
- $\Theta^* = \arg \max_{\Theta_i^*} U_d(\Theta_i^*)$

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- In our experiments, we consider a specific detection model: logistic regression (LR)
- $\Theta = \{\theta_1, \theta_2, \dots, \theta_{|V|}\}$: thresholds of detectors
- We compare our defense strategy against three others:
 - Baseline: simply learn a LR on training data and deploy it at all nodes
 - Re-training: iteratively augment the original training data with attacked instances, re-training the LR each time, until convergence
 - Personalized-single-threshold: this strategy is only allowed to tune a single node's threshold.

Experiments

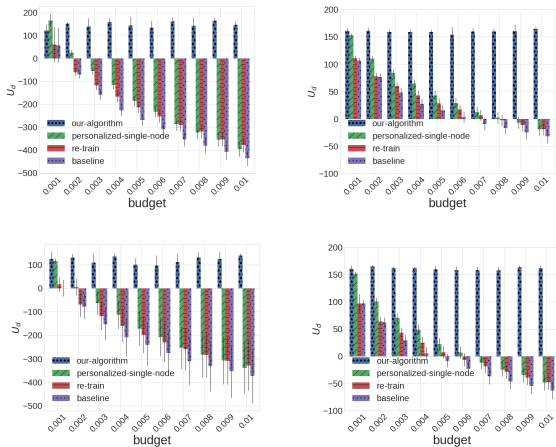


Figure: The performance of each defense strategy. Each bar is averaged over 10 random topologies. Left: BA. Right: Small-world

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- Instead of deploying a “global” detector, we learn and deploy a collection of *heterogeneous* detectors, which takes network structures, propagation of messages, and adversarial behavior into account.
- Formalize the overall problem as a Stackelberg game between a defender and an attacker.
- Utilize *Implicit Function Theorem* to design a novel approach for solving the resulted Stackelberg game.

Thank you!

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