I've done the math enough to know the dangers of our second guessing Doom to crumble unless we grow and strengthen our communication.

 $\sim \text{Tool}$

1 The Model

There are two players, a sender and a receiver, and a state of nature $q \in [0,1]$ that is known to the sender but not the receiver. Let g denote the density of q. The sender sends an intended message $m(q) \in [0,1]$. The receiver receives a noisy version of the intended message, which we call the received message, $\widetilde{m} = m(q) + e$. e is distributed according to a continuous density f with support $[-\overline{e}, \overline{e}]$, where for some integer $N \geq 2$, $\overline{e} = \frac{1}{2N}$. We specify f as follows. Let $x : [-\overline{e}, \overline{e}] \to [0, 1]$ be given by

$$x(e) = \frac{1}{2} \left[1 + \frac{e}{\overline{e}} \right] \tag{1}$$

Consider, as an example, the Beta PDF:

$$f(e) = x(e)^{a-1} (1 - x(e))^{b-1}$$
(2)

for constants a > 0 and b > 0. The receiver then takes an action $a(\widetilde{m}) \in [0, 1]$. Both sender and receiver have utility over the state, q, and the receiver's action, a, of

$$U(q,a) = -\frac{1}{2}(q-a)^2 \tag{3}$$

Prior to the start of the game, the sender can specify an intended message function m(q) that she will use. The receiver chooses an action based upon the function m(q) and the received message \widetilde{m} , and we denote this function $a(\widetilde{m})$. We work backward and start with the optimal action, given m(q) and \widetilde{m} .

Fix a message function m. Let $Q(\widetilde{m})$ denote the set of all states $q \in [0,1]$ such that for some noise $e \in [-\bar{e}, \bar{e}], \ \widetilde{m} = m(q) + e$. In set builder notation,

$$Q(\widetilde{m}) = \{ q \in [0, 1] \mid m(q) - \overline{e} \le \widetilde{m} \le m(q) + \overline{e} \}. \tag{4}$$

Finally, define

$$q_{+}(\widetilde{m}) \equiv \sup Q(\widetilde{m}) \tag{5}$$

$$q_{-}(\widetilde{m}) \equiv \inf Q(\widetilde{m}) \tag{6}$$

 $q_{+}(\widetilde{m})$ and $q_{-}(\widetilde{m})$ are the highest and lowest states that could possibly be associated with the received message \widetilde{m} . Define

$$I_{(\alpha,\beta,\gamma)}(\tilde{m}) \equiv \int_{\alpha}^{\beta} q^{\gamma} f(\tilde{m} - m(q)) dq$$
 (7)

Suppose that the receiver receives the message $\widetilde{m} \in [-\overline{e}, 1 + \overline{e}]$. $q | \widetilde{m}$ has support $[q_{-}(\widetilde{m}), q_{+}(\widetilde{m})]$ and density

$$g(q|\widetilde{m}) = \frac{f(\widetilde{m} - m(q))g(q)}{\int_0^1 f(\widetilde{m} - m(t))g(t)dt}$$
(8)

The receiver's problem is to

$$\max_{a(\widetilde{m})} \int_{q_{-}(\widetilde{m})}^{q_{+}(\widetilde{m})} U(q, a) g(q | \widetilde{m}) dq. \tag{9}$$

The receiver's optimal action is simply the expected value of the state, q, given the received message \widetilde{m} :

$$a(\widetilde{m}) = \int_{q_{-}(\widetilde{m})}^{q_{+}(\widetilde{m})} qg(q|\widetilde{m}) dq. \tag{10}$$

It will be helpful to refer to the cost of a message function. Let the cost functional C be given by

$$C[m] \equiv \int_{0}^{1} \int_{-\bar{e}}^{\bar{e}} (q - a(m(q) + e))^{2} f(e) de dq, \tag{11}$$

where a is the receiver's optimal action from Equation (10). C[m] is the expected loss for a given message function m. The integrand is the loss for a given state q and action $a(\tilde{m})$. The interior integral integrates over the possible exogenous errors, to generate the expected loss given the state. The exterior integral integrates over possible states. Therefore, the sender's problem is to choose a message function m that minimizes C[m]:

$$\min_{m \in M} C[m] \tag{12}$$

where M is the space of weakly increasing piece-wise continuous functions on [0,1]. The change of variables $\widetilde{m}=m(q)+e$ and an application of Fubini's Theorem yield

$$C[m] = \int_0^1 \int_{m(q)-\bar{e}}^{m(q)+\bar{e}} (q - a(\widetilde{m}))^2 f(\widetilde{m} - m(q)) g(q) d\widetilde{m} dq$$
 (13)

$$= \int_{-\bar{e}}^{1+\bar{e}} \int_{q_{-}(\widetilde{m})}^{q_{+}(\widetilde{m})} (q - a(\widetilde{m}))^{2} f(\widetilde{m} - m(q)) g(q) dq d\widetilde{m}.$$
 (14)

We consider the costs of identity and discrete message functions.

2 Identity Message

We consider the identity message function, m(q) = q, only for the case in which q is uniformly distributed.

$$a(\widetilde{m}) = \begin{cases} \overline{a}(\widetilde{m}) & \text{if } 1 - \overline{e} < \widetilde{m} \le 1 + \overline{e} \\ a(\widetilde{m}) & \text{if } \overline{e} < \widetilde{m} \le 1 - \overline{e} \\ \underline{a}(\widetilde{m}) & \text{if } - \overline{e} \le \widetilde{m} \le \overline{e} \end{cases}$$
(15)

where

$$\overline{a}(\widetilde{m}) = \frac{I_{(\widetilde{m} - \overline{e}, 1, 1)}(\widetilde{m})}{I_{(\widetilde{m} - \overline{e}, 1, 0)}(\widetilde{m})}$$
(16)

$$a(\widetilde{m}) = \frac{I_{(\widetilde{m} - \bar{e}, \widetilde{m} + \bar{e}, 1)}(\widetilde{m})}{I_{(\widetilde{m} - \bar{e}, \widetilde{m} + \bar{e}, 0)}(\widetilde{m})}$$
(17)

$$\underline{a}(\widetilde{m}) = \frac{I_{(0,\widetilde{m}+\bar{e},1)}(\widetilde{m})}{I_{(0,\widetilde{m}+\bar{e},0)}(\widetilde{m})}$$
(18)

Note that the normalizing constant cancels out when computing the conditional expectation. The cost of the identity message function is given by

$$C[m_{\mathcal{I}}] = \int_{-\bar{e}}^{1+\bar{e}} \int_{q_{-}(\widetilde{m})}^{q_{+}(\widetilde{m})} (q - a(\widetilde{m}))^{2} f(\widetilde{m} - q) dq d\widetilde{m}.$$
 (19)

where $q_{+}(\widetilde{m}) = \min\{\widetilde{m} + \overline{e}, 1\}$ and $q_{-}(\widetilde{m}) = \max\{\widetilde{m} - \overline{e}, 0\}$. Define

$$\overline{z} = \int_{1-\bar{e}}^{1+\bar{e}} \int_{\widetilde{m}-\bar{e}}^{1} (q - \overline{a}(\widetilde{m}))^2 f(\widetilde{m} - q) dq d\widetilde{m}$$
 (20)

$$z = \int_{\bar{e}}^{1-\bar{e}} \int_{\widetilde{m}-\bar{e}}^{\widetilde{m}+\bar{e}} (q - a(\widetilde{m}))^2 f(\widetilde{m} - q) dq d\widetilde{m}$$
 (21)

$$\underline{z} = \int_{-\bar{e}}^{\bar{e}} \int_{0}^{\tilde{m}+\bar{e}} (q - \underline{a}(\tilde{m}))^{2} f(\tilde{m} - q) dq d\tilde{m}$$
 (22)

so that $C[m_{\mathcal{I}}] = \overline{z} + z + \underline{z}$.

3 Discrete Message

Fix an integer $M \geq 1$ and define $K = M \times N$. Consider the partition

$$0 = x_0 < x_1 < \dots < x_K < x_{K+1} = 1 \tag{23}$$

of [0,1]. For each $i \in \{0,\ldots,K\}$, define $X_i = [x_i,x_{i+1})$. A discrete message with K+1 messages is given by

$$m_{\mathcal{D}}(q) = \frac{1}{K} \sum_{i=0}^{K} \chi_{X_i}(q).$$
 (24)

Each $\widetilde{m} \in [-\bar{e}, 1+\bar{e}]$ maps to M potential bins. The action function is

$$a(\widetilde{m}) = \tag{25}$$

Let $j_+, j_- : [-\bar{e}, 1 + \bar{e}] \to \{0, \dots, K\}$ be given by

$$j_{+}(\widetilde{m}) = \min\{|\widetilde{m}K| + M, K\}$$
(26)

$$j_{-}(\widetilde{m}) = \max\{0, \lfloor \widetilde{m}K \rfloor - M\}$$
(27)

Note that

$$q_{+}(\widetilde{m}) = x_{j_{+}(\widetilde{m})} \tag{28}$$

$$q_{-}(\widetilde{m}) = x_{j_{-}(\widetilde{m})} \tag{29}$$

The cost is given by

$$C[m_{\mathcal{D}}] = \int_{-\bar{e}}^{1+\bar{e}} \int_{x_{j-(\widetilde{m})}}^{x_{j+(\widetilde{m})}} (q - a(\widetilde{m}))^{2} f(\widetilde{m} - m_{\mathcal{D}}(q)) g(q) dq d\widetilde{m}.$$
 (30)

$$= \sum_{j=0}^{j=K} \int_{x_j}^{x_j} (q - a(\widetilde{m}))^2 f(\widetilde{m} - m_{\mathcal{D}}(q)) g(q) dq d\widetilde{m}.$$
 (31)

(32)

It is sufficient to choose the partition in Equation 23. The first-order conditions are

$$x$$
 (33)