





## SIMD Acceleration for Index Structures: A Survey

**Beyond Databases Architectures and Structures** 

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September 18, 2018







## **Agenda**

Motivation

SIMD Style Processing

**Surveyed Main-Memory Index Structures** 

Elf

Seg-Tree/Trie

**FAST** 

**VAST** 

ART

**Qualitative Comparison** 

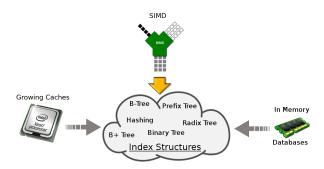
Conclusion







## **Motivation**



- What are state-of-the-art index structures?
- Which optimizations do they have in common?



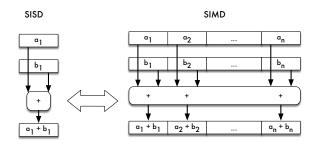


#### Agenda Motivation SIMD Style Processing Surveyed Main-Memory Index Structures

Qualitative Comparison



## Single Instruction Multiple Data



• \_\_m128i \_mm\_add\_epi32 (\_\_m128i a, \_\_m128i b) Adds 4 signed 32-bit integers in a to 4 signed 32-bit integers in b.







## **Surveyed Main-Memory Index Structures**

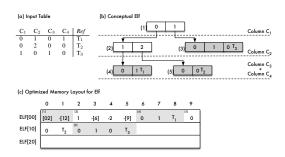
Index Structure	Based on	Reference
Elf	Prefix Tree	[Broneske et al. ICDE'17]
Seg-Tree/Trie	CSB-Tree	[Zeuch et al. EDBT'14]
FAST: Fast Architecture Sensitive Tree	Binary Tree	[Kim et al. SIGMOD'10]
VAST: Vector-Advanced and Compressed Structure Tree	Binary Tree	[Yamamuro et al. EDBT'12]
ART: Adaptive Radix Tree	Radix Tree	[Leis et al. ICDE'13]







### **Elf**



- Multi-dimensional index structure for column-wise storage
- Prefix-redundancy elimination on distinct column values
- Linearisation for optimized memory layout

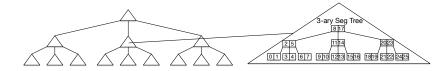
[Broneske et al. ICDE'17]







## Seg-Tree/Trie



- Each node is a k-ary search tree
- Each node is linearised to use k-ary search
- $k = \frac{|SIMD|}{|Key|}$  partitions, k-1 separator keys are compared in parallel

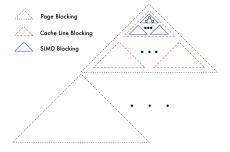
[Zeuch et al. EDBT'14]







## **Fast Architecture Sensitive Tree**



- Based on binary tree
- Hierarchical blocking: page, cache line and SIMD blocks
- Efficient register, cache line and page usage

[Kim et al. SIGMOD'10]







## Vector-Advanced and Compressed Structure Tree

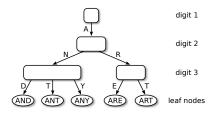
- Extension of FAST
- Decrease branch misses avoiding conditional branches
- Uses key compression
  - Lossy compression for inner nodes
  - Lossfree compression leaf nodes
- Decompression and error correction of lossy compression has less impact compared to the performance increase with SIMD

[Yamamuro et al. EDBT'12]





## **Adaptive Radix Tree**



- Uses different node types with different number of keys and children
- Due to overfill or underfill of nodes, the node type is changed
- Reduced space consumption due to lazy expansion and path compression

[Leis et al. ICDE'13]







## **Qualitative Comparison**

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	×	×	×	×	-	high







## **Qualitative Comparison**

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	Х	×	Х	×	-	high
Minimized key size	o	-	×	×	-	high







## **Qualitative Comparison**

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	Х	×	Х	×	-	high
Minimized key size	o	-	×	X	-	high
Adapted node sizes / types	-	x	-	X	-	low

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







## **Qualitative Comparison**

-	X	high
		nign
-	x	high
-	x	low
-	×	medium
ŀ		
	×	-

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







## **Qualitative Comparison**

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	×	×	Х	×	-	high
Minimized key size	o	-	×	x	-	high
Adapted node sizes / types	-	×	-	x	-	low
Decreased branch misses	-	×	-	x	-	medium
Exploit cache lines using blocking and alignment	-	×	-	×	×	medium

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







## **Qualitative Comparison**

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	х	×	х	×	-	high
Minimized key size	О	-	×	×	-	high
Adapted node sizes / types	-	X	-	×	-	low
Decreased branch misses	-	x	-	×	-	medium
Exploit cache lines using blocking and alignment	-	×	-	×	×	medium
Use of Compression	О	_	×	×	×	medium

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







## **Qualitative Comparison**

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	×	×	Х	×	-	high
Minimized key size	o	-	×	×	-	high
Adapted node sizes / types	-	×	-	×	-	low
Decreased branch misses	-	×	-	×	-	medium
Exploit cache lines using blocking and alignment	-	×	-	×	×	medium
Usage of Compression	o	-	×	×	×	medium
Adapt search algorithm for linearized nodes	×	-	-	-	×	low

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







## **Conclusion**

How to adapt index structures to modern database systems:

- Compare as many keys as possible in parallel with SIMD
  - Direct performance increase by a factor of X (where X keys fit a SIMD register)
- Efficient usage of cache line
- Decrease branch misses
- Use compression or/and adapted search algorithms







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## Thank you for your attention!





