





SIMD Acceleration for Index Structures: A Survey

Beyond Databases Architectures and Structures

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Agenda

Motivation

SIMD Style Processing

Surveyed Main-Memory Index Structures

Elf

Seg-Tree/Trie

FAST

VAST

ART

Qualitative Comparison

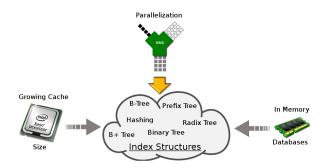
Conclusion







Motivation



- What are state-of-the-art index structures?
- Which optimizations do they have in common?



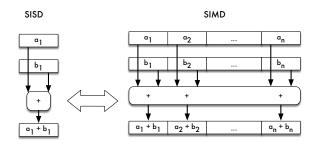


Agenda Motivation SIMD Style Processing Surveyed Main-Memory Index Structures

Qualitative Comparison



Single Instruction Multiple Data



• __m128i _mm_add_epi32 (__m128i a, __m128i b) Adds 4 signed 32-bit integers in a to 4 signed 32-bit integers in b.







Surveyed Main-Memory Index Structures

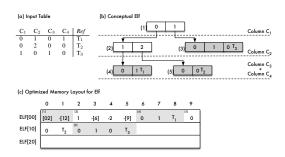
- Elf [Broneske et al. ICDE'17]
- Seg-Tree/Trie [Zeuch et al. EDBT'14]
- FAST: Fast Architecture Sensitive Tree [Kim et al. SIGMOD'10]
- VAST: Vector-Advanced and Compressed Structure Tree [Yamamuro et al. EDBT'12]
- ART: Adaptive Radix Tree [Leis et al. ICDE'13]







Elf



- Multi-dimensional index structure for column-wise storage
- Prefix-redundancy elimination on distinct column values
- Linearisation for optimized memory layout

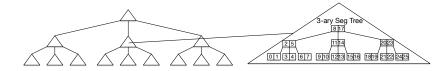
[Broneske et al. ICDE'17]







Seg-Tree/Trie



- Each node is a k-ary search tree
- Each node is linearised to use k-ary search
- $k = \frac{|SIMD|}{|Key|}$ partitions, k-1 separator keys are compared in parallel

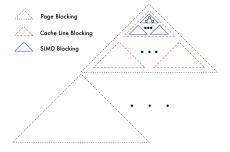
[Zeuch et al. EDBT'14]







Fast Architecture Sensitive Tree



- Based on binary tree
- Hierarchical blocking: page, cache line and SIMD blocks
- Efficient register, cache line and page usage

[Kim et al. SIGMOD'10]







Vector-Advanced and Compressed Structure Tree

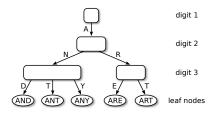
- Extension of FAST
- Decrease branch misses avoiding conditional branches
- Uses key compression
 - Lossy compression for inner nodes
 - Lossfree compression leaf nodes
- Decompression and error correction of lossy compression has less impact compared to the performance increase with SIMD

[Yamamuro et al. EDBT'12]





Adaptive Radix Tree



- Uses different node types with different number of keys and children
- Due to overfill or underfill of nodes, the node type is changed
- Reduced space consumption due to lazy expansion and path compression

[Leis et al. ICDE'13]







Qualitative Comparison

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	×	×	×	×	-	high







Qualitative Comparison

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	Х	×	Х	×	-	high
Minimized key size	o	-	×	×	-	high







Qualitative Comparison

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	Х	×	Х	×	-	high
Minimized key size	o	-	×	X	-	high
Adapted node sizes / types	-	×	-	X	-	low

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







Qualitative Comparison

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Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







Qualitative Comparison

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	×	×	х	×	-	high
Minimized key size	О	-	×	X	-	high
Adapted node sizes / types	-	×	-	X	-	low
Decreased branch misses	-	×	-	×	-	medium
Exploit cache lines using blocking and alignment	-	×	-	×	x	medium

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







Qualitative Comparison

Criterion	Seg-Tree/ Trie	FAST	ART	VAST	Elf	Impact
Horizontal vectorization	х	×	х	×	-	high
Minimized key size	О	-	×	×	-	high
Adapted node sizes / types	-	х	-	×	-	low
Decreased branch misses	-	X	-	×	-	medium
Exploit cache lines using blocking and alignment	-	×	-	×	×	medium
Use of Compression	0	-	×	×	×	medium

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







Qualitative Comparison

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Adapted node sizes / types	-	×	-	×	-	low
Decreased branch misses	-	×	-	×	-	medium
Exploit cache lines using blocking and alignment	-	×	-	×	×	medium
Usage of Compression	О	-	×	×	×	medium
Adapt search algorithm for linearized nodes	×	-	-	-	×	low

Legend: x = implements the issue; o = partially implements the issue; - = does not implement the issue







Conclusion

How to adapt index structures to modern database systems:

- Compare as many keys as possible in parallel with SIMD
 - Direct performance increase by a factor of X (where X keys fit a SIMD register)
- Efficient usage of cache line
- Decrease branch misses
- Use compression or/and adapted search algorithms







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Thank you for your attention!





