

# SIMD Acceleration for Index Structures

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## Abstract—

- **summary:**

**Give short an overview of SIMD and modern index structures**

**Explain what are the problems of the “old” index structures made for disk-based database systems**

**Explain which approaches were made to adapt index structures to modern systems and what they have in common and what are differences**

- **Why is this work important:**

**Give a state of current development of the index structures**

**Collect common approaches to adapt other index structures TODO: ReThink**

- **K-ary search trees, FAST, VAST and ART compared**

- **Contribution: What are important approaches used by different implementations to adapt index structures to modern systems**

**Index Terms**—SIMD, index

## I. INTRODUCTION

After decades of creating and improving index structures for disk-based database systems, nowadays even large databases fit into the main memory. Since index structures like the  $B^+$ -tree [TODO: ref] or the radix tree [TODO: ref] have an important part in database-systems to realise SCAN or range-based search operations, these index structures experienced many adaptations to fulfill the needs of modern database-systems. Instead of overcoming the bottleneck of IO-operations from disk to RAM, the target of modern index-structures is to improve the usage of CPU cache and processor architectures.

In this paper we compare different approaches to adapt index structures to use Single Instruction Multiple Data (SIMD) [1] operations and to better overcome the bottleneck from RAM to CPU cache. We consider the K-ary Search Tree (Seg-tree) [TODO: ref], Adapted Radix Tree (ART) [4], Fast Architecture Sensitive Tree (FAST) [6], and Vector-Advanced and Compressed Structure Tree (VAST) [5]. As the authors of VAST-Tree show, important causes for increased calculation time are cache misses and branch mispredictions. To overcome branch mispredictions and to decrease CPU cycles, SIMD [TODO: ref] is used in modern index structures for tree traversal. The authors of the k-ary search show how to use SIMD to compare multiple keys in one CPU cycle. To decrease cache misses, the authors of FAST and ART show how to adapt index structures to the cache line size.

All approaches use SIMD only for key comparison within tree traversal and try to decrease the key size to fit more keys

into one SIMD register. Therefore FAST and Seg-tree only provide implementations for search algorithms. We consider the approaches VAST and ART make to implement operations like update and insert and name ideas to use SIMD for them. Consequently, with this work we make the following contributions:

- We compare different adaptations of index structures to fulfill requirements of modern database systems
- We highlight the usage of SIMD and the cache line adaptations in all approaches
- We show openings for adaption for other approaches to use SIMD

We organized the rest of the paper as follows. In section 2 we give the preliminaries for SIMD in general and for the use in index structures. In section 3 we analyse the different approaches of adapted index structures and evaluate the comparison in section 4. In section 5 we name related work. In section 5 we present our conclusion and describe future work in section 7.

## II. PRELIMINARIES

A common approach of decreasing CPU cycles for algorithms is to adapt the algorithm to pipelining. While one instruction is executed, the next instruction is already fetched. In contrast to execute one operation on one data after another, the idea of SIMD is to execute a single instruction on multiple data. The authors of Rethinking SIMD Vectorization for In-Memory Databases [TODO: ref] show two general approaches to use SIMD in in-memory databases, horizontal and vertical vector processing. They name the comparison of one search key to multiple other keys horizontal vectorization, whereas processing a different input key per vector lane is named vertical vectorization. TODO: Exact difference?

Since FAST, Seg-Tree, ART and VAST only use horizontal vectorization, we focus on this approach. Modern CPUs have additional SIMD registers along with an additional instruction set adapted to process multiple data with a single instruction. For example, the authors of Seg-Tree [3] use 128-bit SIMD registers and adjusted SIMD operations to load data into a register and to compare the data of one SIMD register with another. A 128-bit SIMD register processes sixteen 8-bit or eight 16-bit data items with one instruction.

## III. ADAPTED TREE STRUCTURES

TODO: Compare all 4 or merge FAST and VAST together/ extend FAST with VAST??

- A. *K-ary search tree*
- B. *FAST*
- C. *ART*
- D. *VAST*

#### IV. EVALUATION

In common:

- SIMD instructions used to compare the search key with multiple keys of the index
- Segmenting tree to blocks for a better usage of cache lines, save the data of the nodes in an adapted way
- The keys should be as short as possible to compare more keys in one step and to decrease the passed data to the cache line
- Each approach improves the tree traversal

Differences:

- Node compression in VAST, Path compression in ART and K-ary seg trie
- FAST and K-ary trees readonly to improve traversal, ART and FAST adapt insert too
- FAST uses and K-ary trees will use GPU calculation instead of CPU

Why performance can not be compared in a useful way...

#### V. RELATED WORK

TODO:

- ART and VAST compared to FAST??
- Ideas and implementations of the adapted trees already in III...
- KD-Tree with SIMD

#### VI. CONCLUSION

#### VII. FUTURE WORK

Open questions, use SIMD for tree creation/updates instead of only for traversal

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