

Taking "algebraically" seriously
in the definition of
algebraically inductive type

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Abstract

Theorem. In a 1-topos, the following two categories are isomorphic, with an isomorphism that is the identity on objects:

1. Pullback-natural, associative, algebraically injective objects.
2. Algebras of the partial-map classifier (aka lifting) monad.

- Partial results towards the ∞ -topos situation.
- We work in HoTT/UF.

I. Algebraic injectives (MScs'2021)

Def. Algebraic injective structure on a type D consists of

1. An extension operation, for any types X and Y ,

$$(-) | (-) : (X \rightarrow D) \times (X \hookrightarrow Y) \rightarrow (Y \rightarrow D).$$

fibers are propositions.

2. For each map $f : X \rightarrow D$ and embedding $j : X \hookrightarrow Y$,

a choice of an identification $(f | j) \circ j = f$, as illustrated by

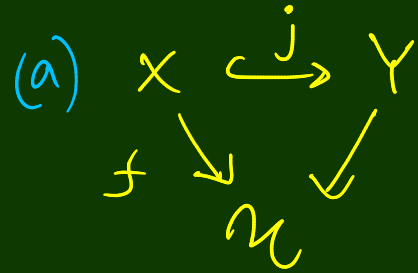
$$\begin{array}{ccc} X & \xhookrightarrow{j} & Y \\ f \searrow & & \swarrow f | j \\ & D & \end{array}$$

Some examples

MSCS'2021

(They need univalence.)

1. $D := \mathcal{U}$



$$(f|j)(y) := \prod_{(x, -) : \text{fiber } j \ y} f(x) .$$

(Right Kan extension.)

(b)

Use \sum instead.

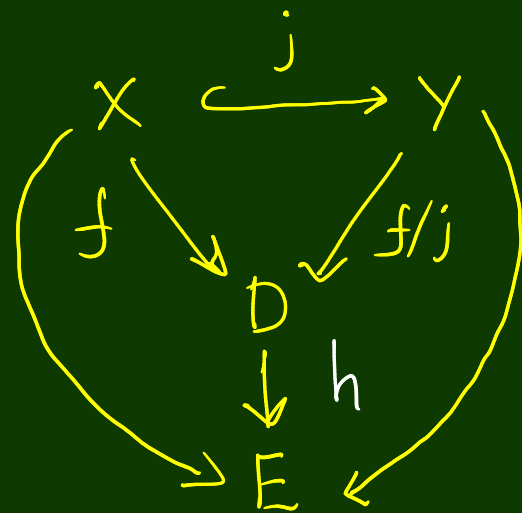
(Left Kan extension.)

2. The type Ω of propositions. Use \forall or \exists .

3. Universes of n -types.

4. Algebras of the lifting monad. We'll come back to this.

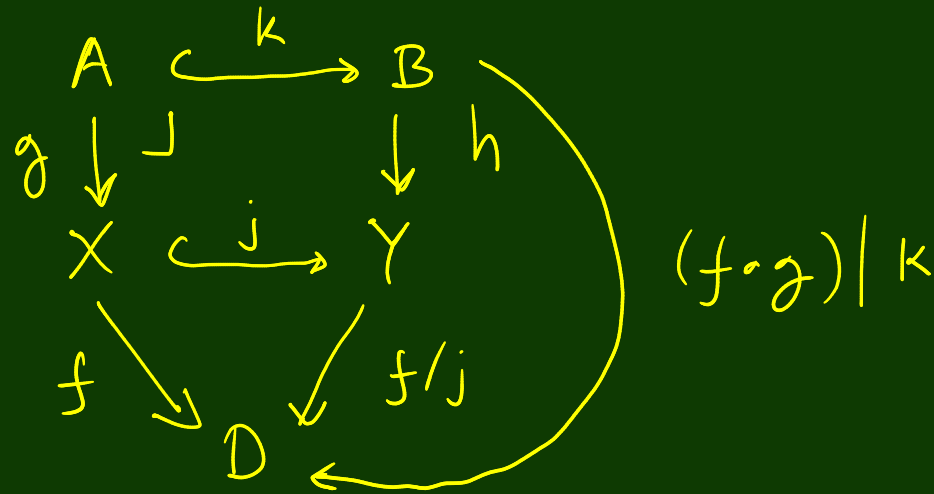
Homomorphisms of algebraic injectives



$$h \circ (f|j) = (h \circ f) | j$$

Pullback naturality

Previous examples are all pullback natural.



$$(f|j) \circ h = (f \circ g) | k$$

It is essential that the square is a pullback.

Consider the non-pullback square

$$\begin{array}{ccc} 0 & \xrightarrow{\quad} & 1 \\ \downarrow & & \downarrow \\ 1 & \xrightarrow{\quad} & 1 \end{array}$$

for a counter-example.

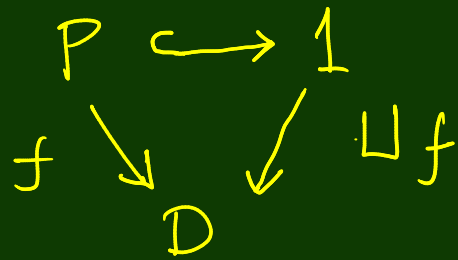
Associativity

$$\begin{array}{ccccc}
 X & \xrightarrow{j} & Y & \xrightarrow{k} & Z \\
 & \searrow & \downarrow f|j & \swarrow & \\
 f & & D & & f|(k \circ j) = (f|j)|k
 \end{array}$$

Examples (Mscs' 2024) $D := \mathcal{U}$ with extension given by Π or Σ .

II. Algebraic floppy structure

MSCS'2021



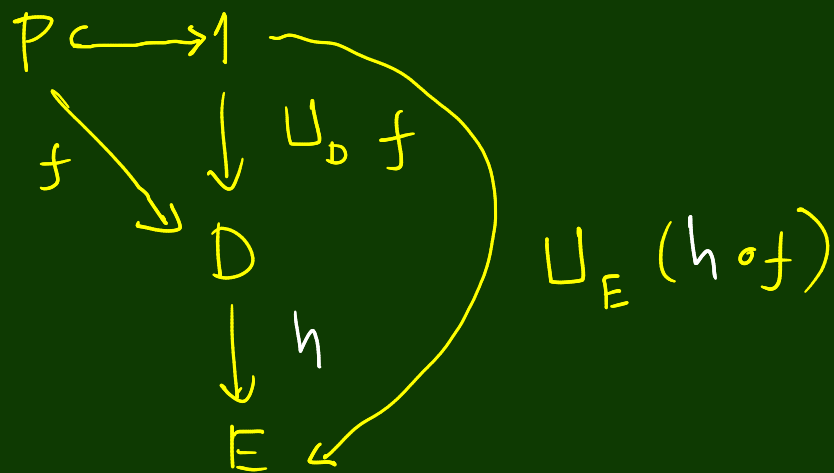
Every partial element
of D can be extended
to a total element

N.B. T.F.A.E.

1. The map $P \rightarrow 1$ is an embedding.
2. The type P is a proposition.

Trivial fact. Algebraic inductive structure is, in particular,
an algebraic floppy structure.

Homomorphisms



$$h \circ U_D f = U_E(h \circ f)$$

construction | Algebraic fibby \rightarrow algebraic injective | MSCS'2021

$$\begin{array}{ccc}
 X & \xrightarrow{j} & Y \\
 f \searrow & & \swarrow \\
 & D &
 \end{array}
 \quad (f|j)(\gamma) := \bigsqcup_{(x,-) \in \underbrace{\text{fiber } j}_Y} f x$$

Fact (new observation)

This algebraic injective structure is pullback natural.
(Not only fiber-natural)

$$\begin{array}{ccc}
 P_Y & \longrightarrow & 1 \\
 p_{r_1} \downarrow & & \downarrow \bigsqcup (f \circ p_{r_1}) \\
 X & \xrightarrow{f} & D
 \end{array}$$

Fiberwise extension.

$$\begin{array}{ccc}
 P_Y & \xrightarrow{c} & 1 \\
 p_{r_1} \downarrow & \lrcorner & \downarrow y \\
 X & \xrightarrow{j} & Y \\
 f \searrow & & \swarrow f|j \\
 & D &
 \end{array}
 \quad \bigsqcup (f \circ p_{r_1})$$

Get this by fibbiness.

III. The lifting monad

M.H.E & C. Kuopp CSL'2017 &
TypeTopology 2018

$$\mathcal{L}X := \sum_{P:\Omega} (P \xrightarrow{\varphi} X)$$

The type of partial elements of X .

$$\begin{array}{ccc} X & & \\ \downarrow & & \\ X & \xrightarrow{f} & Y \\ \eta_X \downarrow & & \downarrow \eta_Y \\ \mathcal{L}X & \xrightarrow{\mathcal{L}f} & \mathcal{L}Y \end{array}$$

$$(\underline{1}, \lambda \cdot x) \quad (P, \varphi) \mapsto (P, f \circ \varphi)$$

$$\text{is-def(ined)} : \mathcal{L}X \rightarrow \Omega := p_{c1}$$

$$\text{value} : (\ell : \mathcal{L}X) \rightarrow \text{is-def } \ell \rightarrow X := p_{c2}$$

$$\begin{array}{ccc} X & \xrightarrow{g} & \mathcal{L}Y \\ \mathcal{L}X & \xrightarrow{g^\#} & \mathcal{L}Y \end{array}$$

$$(P, \varphi) \mapsto$$

is a prop,
as required.

$$\left(\sum_{P:\Omega} \text{is-def}(g(\varphi P)) \right) \lambda(p, d) . \text{value}(g(\varphi P)) d$$

Monad algebras

1. Structure map

$$\sqcup : \mathcal{Z}A \rightarrow A$$

Extend a partial element to a total element!
So \mathcal{Z} -algebras give algebraic fluffy structure.
(MScS'2021)

2. Unit law

$$\sqcup (\lambda(-:1). a) = a \text{ for every } a:A.$$

Extension "property", as for fluffy types.
(extension data !)

3. Associativity law

$$\sqcup_{p:P} \sqcup_{q:Q_p} f(p,q) = \sqcup_{r: \sum_{p:P} Q_p} f r$$

$$P : \Omega$$

$$Q : P \rightarrow \Omega$$

$$f : \sum_{p:P} Q_p \rightarrow A$$

Previously not counted for
when discussing injectivity

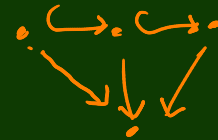
IV. Putting I-III together

So \mathcal{Z} -algebra structure = associative algebraic flabby structure.

Lemma Let \sqcup be the algebraic flabby structure induced by
a given algebraic injective structure \mid that is pullback natural.

Then \sqcup is associative iff \mid is associative

$$\sqcup \sqcup = \sqcup$$



Ongoing work. Replace "iff" by " \simeq ".
(For sets we have this.)

Lemma Let $|$ be the algebraic injective structure induced by
a given algebraic fibbing structure \sqcup .
Then $|$ is always pullback natural. (We've already discussed this.)

Lemma The round trip $\sqcup \mapsto | \mapsto \sqcup'$
is always the identity on both extension operators and extension data.

Lemma The round trip $| \mapsto \sqcup \mapsto |'$
is the identity on extension operators
iff $|$ is pullback natural.

But what about
extension data?
Ongoing.

Theorem. Let D be any type.

1. Then

$$\begin{aligned} & \text{pullback-natural, associative injective structure on } D \\ & \iff \text{associative algebraic flabby structure on } D \\ & = \\ & \mathcal{L}\text{-algebra structure on } D. \end{aligned}$$

2. If D is a set, then " \iff " in (1) becomes an equivalence " \simeq ".

- What is missing to always have 3 type equivalence?
 - check that the pullback-naturality data is unchanged by round trips.
 - check that the associativity data is unchanged by round trips.

(ongoing work, perhaps not difficult.)

- But there is still something else missing.

V. Is \mathcal{L} really a monad?

Nobody knows what a monad on types is in HoTT/UF.

People do know what monads on ∞ -toposes are, though.

But we don't know how to say that in the language of HoTT/UF.

The problem is how to specify coherence data for the monad laws.