Merge sort

- Use a divide-and-conquer approach:
 - 1. If list is of length 0 or 1, already sorted
 - If list has more than one element, split into two lists, and sort each
 - 3. Merge results
 - 1. To merge, just look at first element of each, move smaller to end of the result
 - 2. When one list empty, just copy rest of other list

Example of merging

Left in list 1	Left in list 2	Compare	Result
[1,5,12,18,19,20]	[2,3,4,17]	1, 2	[]
[5,12,18,19,20]	[2,3,4,17]	5, 2	[1]
[5,12,18,19,20]	[3,4,17]	5, 3	[1,2]
[5,12,18,19,20]	[4,17]	5, 4	[1,2,3]
[5,12,18,19,20]	[17]	5, 17	[1,2,3,4]
[12,18,19,20]	[17]	12, 17	[1,2,3,4,5]
[18,19,20]	[17]	18, 17	[1,2,3,4,5,12]
[18,19,20]	[]	18, [[1,2,3,4,5,12,17]
	[]	[1,2,3,4	1,5,12,17,18,19,20]

Complexity of merge

- Comparison and copying are constant
- Number of comparisons O(len(L))
- Number of copyings O(len(L1) + len(L2))
- So merging is linear in length of the lists

```
def merge(left, right, compare):
    result = []
    i, j = 0, 0
    while i < len(left) and j < len(right):</pre>
        if compare(left[i], right[j]):
             result.append(left[i])
             i += 1
        else:
             result.append(right[j])
             j += 1
    while (i < len(left)):</pre>
         result.append(left[i])
         i += 1
    while (j < len(right)):</pre>
         result.append(right[j])
         j += 1
    return result
```

Putting it together

```
import operator

def mergeSort(L, compare = operator.lt):
    if len(L) < 2:
        return L[:]
    else:
        middle = int(len(L)/2)
        left = mergeSort(L[:middle], compare)
        right = mergeSort(L[middle:], compare)
        return merge(left, right, compare)</pre>
```

Complexity of merge sort

- Merge is O(len(L))
- Mergesort is O(len(L)) * number of calls to merge
 - O(len(L)) * number of calls to mergesort
 - O(len(L) * log(len(L)))
- Log linear O(n log n), where n is len(L)
- Does come with cost in space, as makes new copy of list