Binary search

- Can we do better than (len(L)) for search?
- If know nothing about values of elements in list, then no.
- Worst case, would have to look at every element

What if list is ordered?

Suppose elements are sorted in ascending order

```
def search(L, e):
for i in range(len(L)):
    if L[i] == e:
        return True
    if L[i] > e:
        return False
return False
```

 Improves average complexity, but worst case still need to look at every element

Use binary search

- 1. Pick an index, i, that divides list in half
- 2. Ask if L[i] == e
- 3. If not, ask if L[i] larger or smaller than e
- 4. Depending on answer, search left or right half of \perp for \in

A new version of a divide-and-conquer algorithm

- Break into smaller version of problem (smaller list), plus some simple operations
- Answer to smaller version is answer to original problem

Binary search

```
def search(L, e):
def bSearch(L, e, low, high):
    if high == low:
        return L[low] == e
    mid = low + int((high - low)/2)
    if L[mid] == e:
        return True
    if L[mid] > e:
        return bSearch(L, e, low, mid - 1)
    else:
        return bSearch(L, e, mid + 1, high)
if len(L) == 0:
    return False
else:
    return bSearch(L, e, 0, len(L) - 1)
```

Analyzing binary search

- Does the recursion halt?
 - Decrementing function
 - 1. Maps values to which formal parameters are bound to non-negative integer
 - 2. When value is <= 0, recursion terminates
 - 3. For each recursive call, value of function is strictly less than value on entry to instance of function
 - Here function is high low
 - At least 0 first time called (1)
 - When exactly 0, no recursive call, returns (2)
 - Otherwise, halt or recursively call with value halved (3)
 - So terminates

Analyzing binary search

- What is complexity?
 - How many recursive calls? (work within each call is constant)
 - How many times can we divide high low in half before reaches 0?
 - $-\log_2(\text{high} \text{low})$
 - Thus search complexity is O (log(len(L)))