

# Automata in Software Verification and Testing

Martin Hruška

supervisors: Tomáš Vojnar, Lukáš Holík

Brno University of Technology, Czech Republic

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# Introduction

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  - ▶ Shape analysis (majority of the work).
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  - ▶ Have to deal with infinite state spaces.
  - ▶ Application: Proving correctness of critical software systems (e.g., operating system kernel).

# Automata Theory in Shape Analysis

- Data structures can be viewed as graphs.
- Graphs can be accepted by automata.
- An automaton can represent a set of graphs.
- **Goal of analysis:** Derive an automaton for each program location representing all possible shapes of data structures at the location.
- Applied in **Abstract Regular Tree Model Checking (ARTMC)**.

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## Advantages

- Existing **efficient algorithms** for handling finite automata.
- **Genericity** - can represent various kinds of data structures.



# Counterexample Validation and Interpolation-Based Refinement for Forest Automata

Holík, L., [Hruška, M.](#), Lengál, O., Rogalewicz, A., Vojnar, T. In *Proc. of VMCAI'17*.

# Forest Automata Encoding of Heap

- A **Forest Automaton (FA)**\* is a tuple of **tree automata (TA)**.

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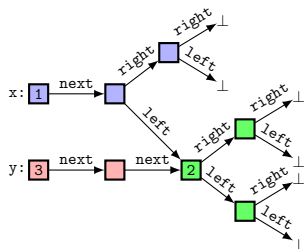
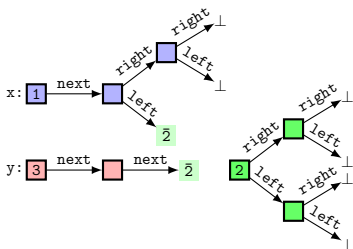
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- An FA  $F = (TA_1, \dots, TA_n)$  represents **tree decompositions** (tuples of trees  $t_1, \dots, t_n$ ) of heap graphs such that  $\forall 1 \leq i \leq n : t_i \in L(TA_i)$ .

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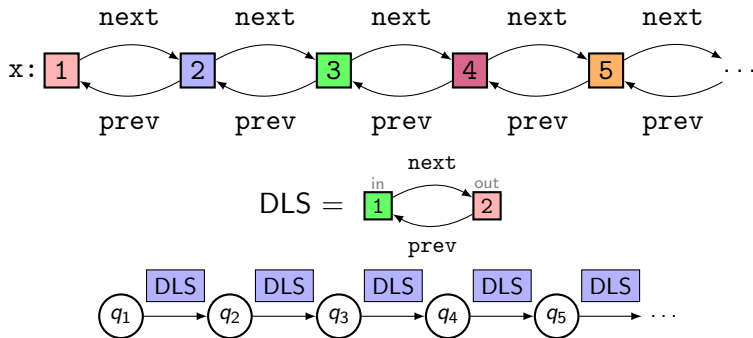
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- Encoded heap graphs obtained by connecting leaves with the referenced roots.



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# Boxes

- **Boxes** are used to extend the expressive power of FA.
- A box is an FA that can be used as a symbol of another FA.
- FA having boxes in the alphabet are called **hierarchical**.
- A box represents repeating subgraphs of a heap.



# An Overview of Verification Method

- Based on [Abstract Regular Tree Model Checking \(ARTMC\)](#)<sup>†</sup> and [Counterexample-guided Abstraction Refinement \(CEGAR\)](#)<sup>‡</sup>.
  - ▶ Sets of heap configurations are represented by automata.
  - ▶ Employs [abstraction](#) over automata to overapproximate the set of reachable configurations (allowing termination on  $\infty$  state spaces).

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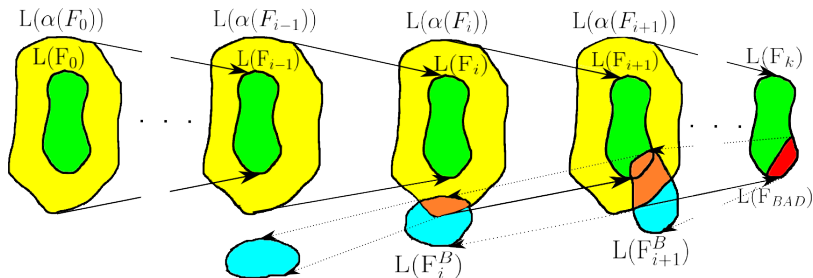
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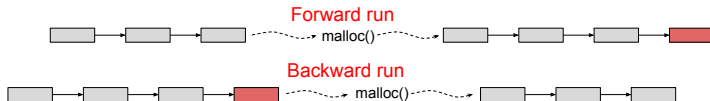
- Overapproximates reachable configurations and accelerates analysis.
- Collapses states in the same equivalence class of a relation  $\sim$ .
- Predicate Language Abstraction
  - ▶ Given a set of predicate languages  $P = \{p_1, \dots, p_n\}$ , equivalence  $\sim_P$  is defined as

$$q_1 \sim_P q_2 \stackrel{DEF}{\equiv} \forall p \in P : L(q_1) \cap p \neq \emptyset \Leftrightarrow L(q_2) \cap p \neq \emptyset$$

- ▶ **Refinement:** Interpolating languages from counterexamples.
- ▶ **Informed refinement.**

# Backward Run and Abstraction for Forest Automata

- Counterexample validation via backward run.
- Ingredients for backward run:
  - ▶ Reversion of abstract transformations.

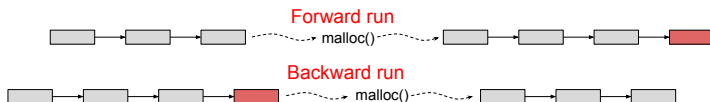


# Backward Run and Abstraction for Forest Automata

- Counterexample validation via backward run.

- **Ingredients for backward run:**

- ▶ Reversion of **abstract transformations**.



- ▶ Reversion of **abstraction**  $\rightarrow$  **intersection** of FA.

- Consider FA  $F_1 = (TA_1^1, \dots, TA_n^1)$  and  $F_2 = (TA_1^2, \dots, TA_n^2)$ , intersection is done **component-wise** using TA intersection, i.e.,  $F_1 \cap F_2 = (TA_1^1 \cap TA_1^2, \dots, TA_n^1 \cap TA_n^2)$ .

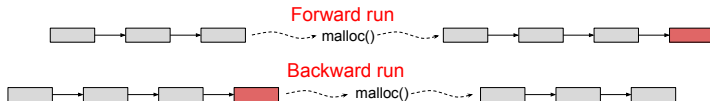
- ▶ Precise reversion of folding and unfolding of boxes  $\Rightarrow$  **compatible form**.

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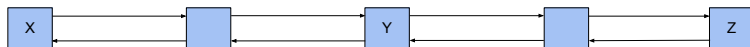
- ▶ Precise reversion of folding and unfolding of boxes  $\Rightarrow$  compatible form.

- Ingredients for predicate language abstraction:

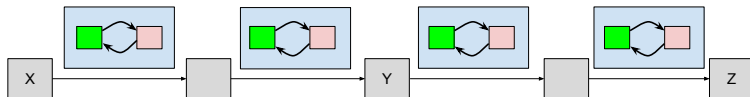
- ▶ Predicate languages represented by TA.
- ▶ New predicate TAs obtained by intersecting FAs from FW and BW run.

# Backward run with Hierarchical Forest Automata

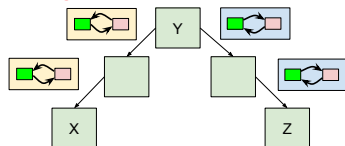
- To enable component-wise hierarchical FA intersection, so called **compatible form** of FA is needed:
  - ▶ They have same number of TA.
  - ▶ The same subgraphs are folded into the same boxes.
- Compatible form is preserved during backward run.



**Folding from X:**



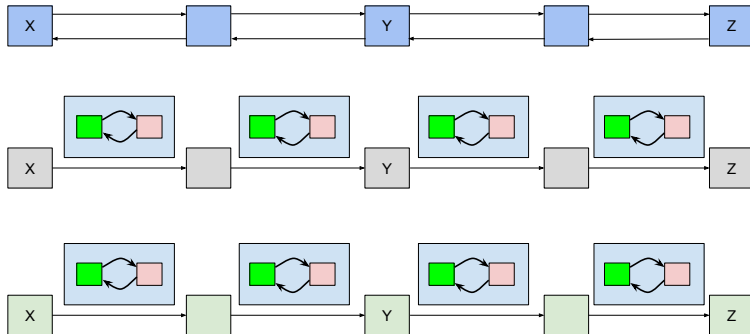
**Folding from Y:**



# Intersection with Hierarchical Forest Automata

## ■ Intersection of FA is done:

- ▶ Component-wise using TA intersection.
- ▶ When TA intersection reaches boxes in the transitions of both automata it calls the whole procedure for **FA intersection recursively on the boxes** and uses its results as a new box.



# Experimental Evaluation

- Shape analysis based on Forest automata implemented in the **Forester** tool

Program	Status	LoC	Time [s]	Refnm	Preds	Program	Status	LoC	Time [s]	Refnm	Preds
SLL (delete)	safe	33	0.02	0	0	DLL (rev)	safe	39	0.70	0	0
SLL (bubblesort)	safe	42	0.02	0	0	CDLL	safe	32	0.02	0	0
SLL (insertsort)	safe	36	0.04	0	0	DLL (insertsort)	safe	42	0.56	0	0
SLLOfCSLL	safe	47	0.02	0	0	DLLOfCDLL	safe	54	1.76	0	0
SLL01	safe	70	1.20	1	1	DLL01	safe	73	0.65	2	2
CircularSLL	safe	49	3.57	3	3	CircularDLL	safe	52	37.22	18	24
OptPtrSLL	safe	59	1.90	3	3	OptPtrDLL	safe	62	1.87	5	5
QueueSLL	safe	71	11.32	10	10	QueueDLL	safe	74	44.68	14	14
GBSLL	safe	64	0.84	3	3	GBDLL	safe	71	1.89	4	4
GBSLLSent	safe	68	0.85	3	3	GBDLLSent	safe	75	2.19	4	4
RGSLL	safe	72	14.41	22	38	RGDLL	safe	76	78.76	26	26
WBSLL	safe	62	0.84	5	5	WBDLL	safe	71	1.37	7	7
SortedSLL	safe	76	227.12	15	15	SortedDLL	safe	82	36.67	11	11
EndSLL	safe	45	0.07	2	2	EndDLL	safe	49	0.10	3	3
TreeRB	error	130	0.08	0	0	TreeWB	error	125	0.05	0	0
TreeCnstr	safe	52	0.31	0	0	TreeCnstr	error	52	0.03	0	0
TreeOfCSLL	safe	109	0.57	0	0	TreeOfCSLL	error	109	0.56	1	3
TreeStack	safe	58	0.20	0	0	TreeStack	error	58	0.01	0	0
TreeDSW	safe	72	1.87	0	0	TreeDSW	error	72	0.02	0	0
TreeRootPtr	safe	62	1.43	0	0	TreeRootPtr	error	62	0.17	2	6
SkipList	safe	84	3.36	0	0	SkipList	error	84	0.08	1	1



- Forester competed in SV-COMP in years 2015-2018 in the categories related to memory safety (*MemSafety*, *HeapReach*).

- Forester competed in SV-COMP in years 2015-2018 in the [categories related to memory safety](#) (*MemSafety*, *HeapReach*).
- Forester has never won any medal, but was able to verify difficult data structures (skip-lists or various trees) which were not solved by any other tool.

# Towards Efficient Shape Analysis with Tree Automata

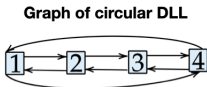
Holík, L., [Hruška, M.](#) NETYS'21.

# Towards Efficient Shape Analysis with Tree Automata (TA)

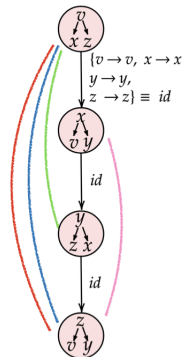
- We proposed new automata able to represent **graphs with a bounded tree-width**.
- They are based on tree automata but we work with a single TA instead of tuple of TA (like in FA).
- We sketched **an efficient algorithm for entailment** for these automata.

# Tree Decomposition with Variables

- How to represent a graph with bounded tree-width by a tree automaton? By tree decompositions with **variables**.



**Tree decompositions:**



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- We can create a TA representing all tree decompositions that is of **a polynomial size assuming a fixed tree-width** of the original automaton.
- Since automata TA inclusion is EXPTIME-complete, we have entailment which is **singly exponential** (assuming a fixed maximum tree width).



## Shape Analysis based on SMT Solving in 2LS Framework

Malík, V., [Hruška, M.](#), Schrammel P., Vojnar T. FMCAD'18.

# Shape Analysis based on SMT Solving in 2LS Framework

- Shape analysis using **SMT solving** to compute the points-to relation between pointers and (abstract) memory addresses.
- A domain designed for **representation of linked-lists**.
- Therefore more straightforward than the automata-based approaches, but lacks their generality.

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- Shape analysis using **SMT solving** to compute the points-to relation between pointers and (abstract) memory addresses.
- A domain designed for **representation of linked-lists**.
- Therefore more straightforward than the automata-based approaches, but lacks their generality.
- Implemented within **the 2LS framework** for program analysis.
  - ▶ Combination with other domains in the framework, e.g., the numerical domain.

# Shape Analysis based on SMT Solving in 2LS Framework

- The verification procedure takes:
  - ▶ A **first order formula over combination of SMT theories** that represents the program in SSA form (i.e., transition relation).
  - ▶ A **set of invariants based on predefined templates** (proposed for various domains).
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- And verifies that no reachable invariants violate the properties of interest.
- The template for shape analysis is:  $\mathcal{T}^S \equiv \bigwedge_{p_i^{lb} \in P_{tr}^{lb}} \mathcal{T}_{p_i^{lb}}^S(d_{p_i^{lb}})$ ,
  - ▶ where  $\mathcal{T}_{p_i^{lb}}^S(d_{p_i^{lb}}) \equiv (\bigvee_{a \in d_{p_i^{lb}}} p_i^{lb} = a)$ .
  - ▶ Basically, it describes the points-to relation between a pointer ( $p_i^{lb}$ ) and addresses  $a$  that the pointer may point to.

# Experimental Evaluation

**Table:** Comparison of 2LS with other tools on examples that combine reasoning about unbounded data structures and their stored data.

	2LS	CPA-Seq	PredatorHP	Forester	Symbiotic	UAutomizer
Calendar	2.88	timeout	false	unknown	timeout	timeout
Cart	23.70	timeout	false	unknown	timeout	timeout
Hash Function	3.65	8.51	unknown	unknown	unknown	timeout
MinMax	5.14	timeout	false	unknown	timeout	timeout
Packet Filter	431.00	timeout	timeout	unknown	unknown	timeout
Process Queue	6.62	7.68	timeout	unknown	timeout	timeout
Quick Sort	18.20	3.50	timeout	unknown	unknown	5.75
Running Example	1.24	timeout	timeout	unknown	timeout	unknown
SM1	0.53	timeout	0.31	false	timeout	timeout
SM2	0.55	5.41	false	false	timeout	14.50

## Generating Scenarios for Digital Twins of Distributed Manufacturing Execution Systems.

Fiedor, T., [Hruška, M.](#), Smrčka, A. EUROCAST'22.



## ■ Manufacturing Execution System (MES)

- ▶ A software managing production in factory.
- ▶ Communicates with information system, machines, and other parts of factory.
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- ▶ Generally, a digital copy of a cyber-physical system — simulation of reality on computer, often with a graphical interface.
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## ■ Manufacturing Execution System (MES)

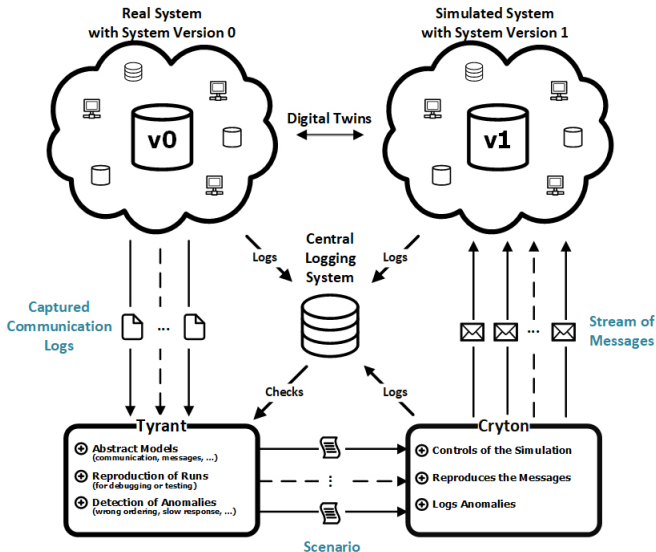
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- The project solved with the Unis company ([MES Pharis](#)) and Masaryk University ([Cryton](#) for orchestrating digital twin)

# Our Solution



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  - ▶ Direct generation of a scenario for digital twin.
  - ▶ Transformation to finite automaton - apply abstraction to overapproximate language.
  - ▶ New strings in overapproximated language  $\Rightarrow$  new series of events  $\Rightarrow$  new scenario  $\Rightarrow$  **new test case**.

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  - ▶ New strings in overapproximated language  $\Rightarrow$  new series of events  $\Rightarrow$  new scenario  $\Rightarrow$  **new test case**.
- Implemented in the **Tyrant** tool.
  - ▶ Generated valid scenarios for testing MES Pharis.
  - ▶ More engineering work needed to be deployed in production.

# Conclusion



- Work presented in the thesis:
  - ▶ Counterexample analysis and abstraction refinement for Forest automata.
  - ▶ Forester participation in SV-COMP, editions 2015-2018.
  - ▶ A chapter on forest automata in book on software verification with refined presentation of the approach.
  - ▶ Automata over graphs with bounded tree-width.
  - ▶ Shape analysis based on SMT solving.
  - ▶ Automated testing of distributed manufacturing execution systems.

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  - ▶ Automated testing of distributed manufacturing execution systems.
- Work not presented and only mildly mentioned in the thesis:
  - ▶ Connection of Predator and Symbiotic participating in [SV-COMP'20](#).
  - ▶ The design and implementation of efficient automata library called MATA ([TACAS'24](#)).

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# The Questions of Reviewers

- Is a user-non-friendliness of library for finite automata a problem if it used as backend of verification tool?
  - ▶ Context: In conclusion of my thesis I proposed that one of the future directions is to design a new efficient, simple, and user-friendly automata library.
  - ▶ Depends on definition of user-friendliness.
  - ▶ In general, if you work in small team with limited amount of resources you want library that has a clear interface and predictable behaviour, which is a kind of user friendliness.

# The Questions of Reviewers

- When generating abstract messages, could the abstraction lead to generation of messages that really do not make sense in practice? Is it still worth to examine the behaviour of the system after transferring such messages?
  - ▶ Context: In work on testing MES we proposed methods for generating new test cases using abstraction over messages and communication models.
  - ▶ Yes, it can lead to messages or sequences of messages that do not make sense but the system under testing should not fail when these messages.
  - ▶ From practical point of view, it is necessary to tune the abstractions to avoid generating such things (common scenario. interesting if CEGAR can be used)