

# Automata in Software Verification and Testing

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March 5, 2024

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  - ▶ Shape analysis (majority of the work).
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  - ▶ Have to deal with infinite state spaces.
  - ▶ Application: Proving correctness of critical software systems (e.g., operating system kernel).

# Automata Theory in Shape Analysis

- Data structures can be viewed as graphs.
- Graphs can be accepted by automata.
- An automaton can represent a set of graphs.
- **Goal of analysis:** Derive an automaton for each program location representing all possible shapes of data structures at the location,
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## Advantages

- Existing **efficient algorithms** for handling finite automata.
- **Genericity** — can represent various kinds of data structures.



# Counterexample Validation and Abstraction Refinement for Shape Analysis based on Forest Automata

Holík, L., [Hruška, M.](#), Lengál, O., Rogalewicz, A., Vojnar, T.: Counterexample Validation and Interpolation-Based Refinement for Forest Automata. In *Proc. of VMCAI'17*, LNCS 10145, Springer, 2017.

# Forest Automata Encoding of Heap

- A **Forest Automaton (FA)**\* is a tuple of **tree automata (TA)**.

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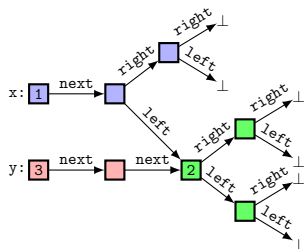
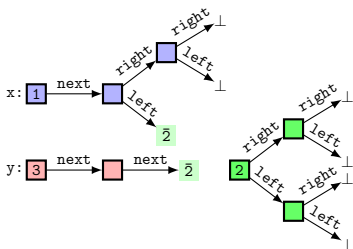
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- An FA  $F = (TA_1, \dots, TA_n)$  represents **tree decompositions** (tuples of trees  $t_1, \dots, t_n$ ) of heap graphs such that  $\forall 1 \leq i \leq n : t_i \in L(TA_i)$ .

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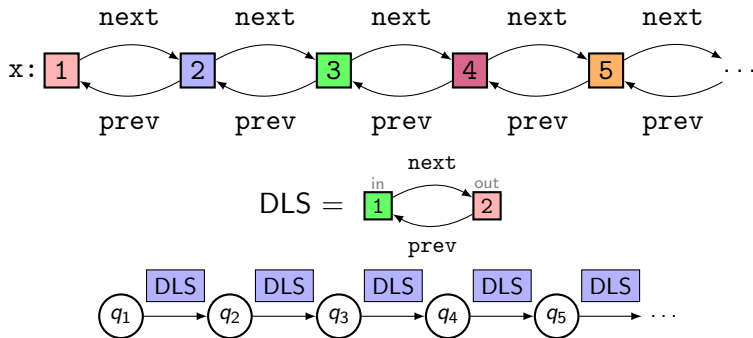
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- Encoded heap graphs obtained by connecting leaves with the referenced roots.



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# Boxes

- **Boxes** are used to extend the expressive power of FA.
- A box is an FA that can be used as a symbol of another FA.
- A box represents repeating subgraphs of a heap.
- FA having boxes in the alphabet are called **hierarchical**.



# An Overview of Verification Method

- Based on [Abstract Regular Tree Model Checking \(ARTMC\)](#)<sup>†</sup> and [Counterexample-guided Abstraction Refinement \(CEGAR\)](#)<sup>‡</sup>.
  - ▶ Sets of heap configurations are represented by automata.
  - ▶ Employs [abstraction](#) over automata to overapproximate the set of reachable configurations (allowing termination on  $\infty$  state spaces).

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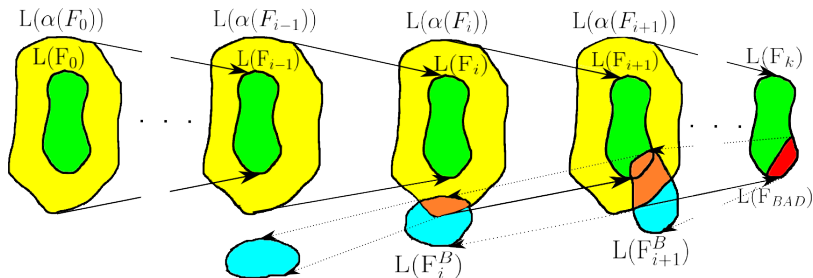
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# Abstraction

- Overapproximates reachable configurations and accelerates analysis
- Collapses states in the same equivalence class of a relation  $\sim$
- Height Abstraction
  - ▶ Equivalence  $\sim_H$  is defined as

$$q_1 \sim_H q_2 \stackrel{DEF}{\equiv} L^n(q_1) = L^n(q_2)$$

where  $L^n(q)$  is the language of prefixes of trees accepted from  $L(q)$  with height up to  $n$

- ▶ Refinement: By increasing the height
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- Predicate Language Abstraction

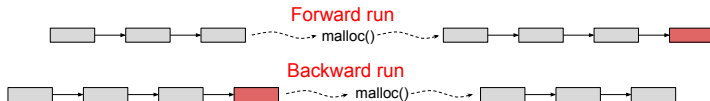
- ▶ Given a set of predicate languages  $P = \{p_1, \dots, p_n\}$ , equivalence  $\sim_P$  is defined as

$$q_1 \sim_P q_2 \stackrel{DEF}{\equiv} \forall p \in P : L(q_1) \cap p \neq \emptyset \Leftrightarrow L(q_2) \cap p \neq \emptyset$$

- ▶ Refinement: Interpolating languages from counterexamples
  - ▶ Informed refinement

# Backward Run and Abstraction for Forest Automata

- Counterexample validation via backward run.
- Ingredients for backward run:
  - ▶ Reversion of abstract transformations.

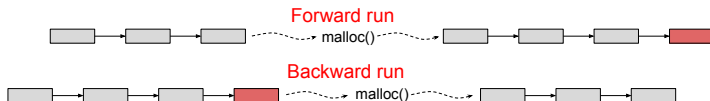


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- ▶ Reversion of abstraction  $\rightarrow$  intersection of FA.

- Consider FA  $F_1 = (TA_1^1, \dots, TA_n^1)$  and  $F_2 = (TA_1^2, \dots, TA_n^2)$ , intersection is done component-wise using TA intersection, i.e.,  $F_1 \cap F_2 = (TA_1^1 \cap TA_1^2, \dots, TA_n^1 \cap TA_n^2)$ .

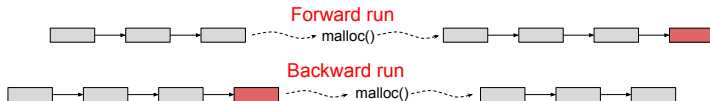
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- ▶ Precise reversion of folding and unfolding of boxes  $\Rightarrow$  compatible form.

- Ingredients for predicate language abstraction:

- ▶ Predicate languages represented by TA.
- ▶ New predicate TAs obtained by intersecting FAs from FW and BW run.

# Experimental Evaluation

- Shape analysis based on Forest automata implemented in the **Forester** tool

| Program            | Status | LoC | Time [s] | Refnm | Preds | Program            | Status | LoC | Time [s] | Refnm | Preds |
|--------------------|--------|-----|----------|-------|-------|--------------------|--------|-----|----------|-------|-------|
| SLL (delete)       | safe   | 33  | 0.02     | 0     | 0     | DLL (rev)          | safe   | 39  | 0.70     | 0     | 0     |
| SLL (bubblesort)   | safe   | 42  | 0.02     | 0     | 0     | CDLL               | safe   | 32  | 0.02     | 0     | 0     |
| SLL (insertsort)   | safe   | 36  | 0.04     | 0     | 0     | DLL (insertsort)   | safe   | 42  | 0.56     | 0     | 0     |
| SLLOfCSLL          | safe   | 47  | 0.02     | 0     | 0     | DLLOfCDLL          | safe   | 54  | 1.76     | 0     | 0     |
| <b>SLL01</b>       | safe   | 70  | 1.20     | 1     | 1     | <b>DLL01</b>       | safe   | 73  | 0.65     | 2     | 2     |
| <b>CircularSLL</b> | safe   | 49  | 3.57     | 3     | 3     | <b>CircularDLL</b> | safe   | 52  | 37.22    | 18    | 24    |
| <b>OptPtrSLL</b>   | safe   | 59  | 1.90     | 3     | 3     | <b>OptPtrDLL</b>   | safe   | 62  | 1.87     | 5     | 5     |
| <b>QueueSLL</b>    | safe   | 71  | 11.32    | 10    | 10    | <b>QueueDLL</b>    | safe   | 74  | 44.68    | 14    | 14    |
| <b>GBSLL</b>       | safe   | 64  | 0.84     | 3     | 3     | <b>GBDLL</b>       | safe   | 71  | 1.89     | 4     | 4     |
| <b>GBSLLSent</b>   | safe   | 68  | 0.85     | 3     | 3     | <b>GBDLLSent</b>   | safe   | 75  | 2.19     | 4     | 4     |
| <b>RGSLL</b>       | safe   | 72  | 14.41    | 22    | 38    | <b>RGDLL</b>       | safe   | 76  | 78.76    | 26    | 26    |
| <b>WBSLL</b>       | safe   | 62  | 0.84     | 5     | 5     | <b>WBDLL</b>       | safe   | 71  | 1.37     | 7     | 7     |
| <b>SortedSLL</b>   | safe   | 76  | 227.12   | 15    | 15    | <b>SortedDLL</b>   | safe   | 82  | 36.67    | 11    | 11    |
| <b>EndSLL</b>      | safe   | 45  | 0.07     | 2     | 2     | <b>EndDLL</b>      | safe   | 49  | 0.10     | 3     | 3     |
| <b>TreeRB</b>      | error  | 130 | 0.08     | 0     | 0     | <b>TreeWB</b>      | error  | 125 | 0.05     | 0     | 0     |
| <b>TreeCnstr</b>   | safe   | 52  | 0.31     | 0     | 0     | <b>TreeCnstr</b>   | error  | 52  | 0.03     | 0     | 0     |
| <b>TreeOfCSLL</b>  | safe   | 109 | 0.57     | 0     | 0     | <b>TreeOfCSLL</b>  | error  | 109 | 0.56     | 1     | 3     |
| <b>TreeStack</b>   | safe   | 58  | 0.20     | 0     | 0     | <b>TreeStack</b>   | error  | 58  | 0.01     | 0     | 0     |
| <b>TreeDSW</b>     | safe   | 72  | 1.87     | 0     | 0     | <b>TreeDSW</b>     | error  | 72  | 0.02     | 0     | 0     |
| <b>TreeRootPtr</b> | safe   | 62  | 1.43     | 0     | 0     | <b>TreeRootPtr</b> | error  | 62  | 0.17     | 2     | 6     |
| <b>SkipList</b>    | safe   | 84  | 3.36     | 0     | 0     | <b>SkipList</b>    | error  | 84  | 0.08     | 1     | 1     |

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- **2017**
  - ▶ Generating correctness witness (forest automaton representing shape invariants of program).
  - ▶ Counterexample analysis and abstraction refinement for hierarchical forest automata.
- Forester has never won any medal, but was able to verify difficult data structures (skip-lists, or various trees) which were not solved by any other tool.

# Towards Efficient Shape Analysis with Tree Automata

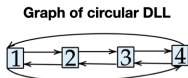
Holík, L., [Hruška, M.](#) NETYS'21.

# Towards Efficient Shape Analysis with Tree Automata (TA)

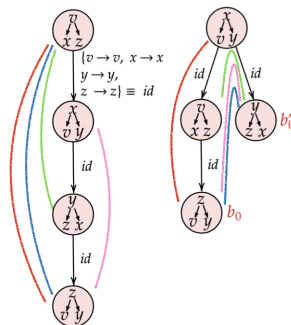
- We proposed new automata able to represent **graphs with bounded tree-width**.
- They are based on tree automata but we work with single TA instead of tuple of TA (like in FA).
- We sketched **efficient algorithm for entailment** for these automata.

# Tree Decomposition with Variables

- How to represent a graph with bounded tree-width by a tree automaton? By tree decompositions with **variables**.



Tree decompositions:



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- We can create TA representing all tree decompositions that is of a **polynomial size assuming a fixed tree-width** of the original automaton.
- Since automata inclusion is EXPTIME-complete, we have entailment which is **singly exponential** (assuming a fixed maximum tree width).

## Shape Analysis based on SMT Solving in 2LS Framework

Malík, V., [Hruška, M.](#), Schrammel P., Vojnar T. FMCAD'18.

# Shape Analysis based on SMT Solving in 2LS Framework

- Shape analysis using **SMT solving** to compute the points-to relation between pointers and (abstract) memory addresses.
- Domain designed for **representation of linked-lists**.
- Therefore more straightforward than the automata-based approaches, but lacks their generality.

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- Implemented within **the 2LS framework** for program analysis.
  - ▶ Combination with other domains in the framework, e.g., the numerical domain.

# Shape Analysis based on SMT Solving in 2LS Framework

- Verification procedure takes:
  - ▶ A **first order formula over combination of SMT theories** that represents the program in SSA form (i.e., transition relation).
  - ▶ A **set of invariants based on predefined templates** (proposed for various domains).
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- And verifies that there is no reachable invariants violate the properties of interest.
- The template for shape analysis is:  $\mathcal{T}^S \equiv \bigwedge_{p_i^{lb} \in P_{tr}^{lb}} \mathcal{T}_{p_i^{lb}}^S(d_{p_i^{lb}})$ ,
  - ▶ where  $\mathcal{T}_{p_i^{lb}}^S(d_{p_i^{lb}}) \equiv (\bigvee_{a \in d_{p_i^{lb}}} p_i^{lb} = a)$ .
  - ▶ Basically, it describes the points-to relation between a pointer ( $p_i^{lb}$ ) and addresses  $a$  that the pointer may point to.

# Experimental Evaluation

**Table:** Comparison of 2LS with other tools on examples that need reasoning about unbounded data structures and their stored data.

|                 | 2LS    | CPA-Seq | PredatorHP | Forester | Symbiotic | UAutomizer |
|-----------------|--------|---------|------------|----------|-----------|------------|
| Calendar        | 2.88   | timeout | false      | unknown  | timeout   | timeout    |
| Cart            | 23.70  | timeout | false      | unknown  | timeout   | timeout    |
| Hash Function   | 3.65   | 8.51    | unknown    | unknown  | unknown   | timeout    |
| MinMax          | 5.14   | timeout | false      | unknown  | timeout   | timeout    |
| Packet Filter   | 431.00 | timeout | timeout    | unknown  | unknown   | timeout    |
| Process Queue   | 6.62   | 7.68    | timeout    | unknown  | timeout   | timeout    |
| Quick Sort      | 18.20  | 3.50    | timeout    | unknown  | unknown   | 5.75       |
| Running Example | 1.24   | timeout | timeout    | unknown  | timeout   | unknown    |
| SM1             | 0.53   | timeout | 0.31       | false    | timeout   | timeout    |
| SM2             | 0.55   | 5.41    | false      | false    | timeout   | 14.50      |



## Generating Scenarios for Digital Twins of Distributed Manufacturing Execution Systems.

Fiedor, T., [Hruška, M.](#), Smrčka, A. EUROCAST'22.

- Manufacturing Execution System (MES)
  - ▶ Software managing production in factory.
  - ▶ This includes communication with information system, machines, etc.
  - ▶ Difficult to test:
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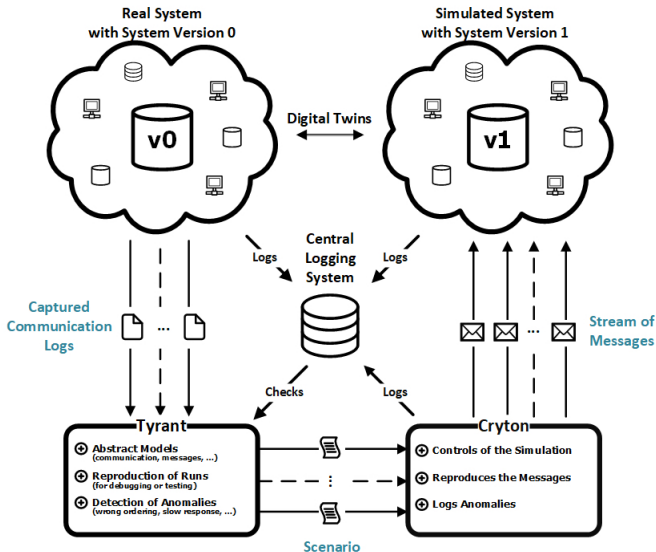
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- Project solved with the Unis company ([MES Pharis](#)) and Masaryk University ([Cryton](#) for orchestrating digital twin)

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  - ▶ Direct generation of scenario for digital twin.
  - ▶ Can be transformed to finite automaton — application of abstraction to overapproximate language.
  - ▶ New strings in overapproximated language  $\Rightarrow$  new series of events  $\Rightarrow$  new scenario  $\Rightarrow$  **new test case**.



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  - ▶ New strings in overapproximated language  $\Rightarrow$  new series of events  $\Rightarrow$  new scenario  $\Rightarrow$  **new test case**.
- Implemented in the **Tyrant** tool.
  - ▶ Generated valid scenarios for orchestration of digital twin with deployed Pharis.
  - ▶ More engineering work needed to be deployed in production.

# Conclusion

- Work presented in the thesis:
  - ▶ Counterexample analysis and abstraction refinement for Forest automata.
  - ▶ Forester participation in SV-COMP, editions 2015-2018.
  - ▶ A chapter on forest automata in book on software verification with refined presentation of the approach.
  - ▶ Automata over graphs with bounded tree-width.
  - ▶ Shape analysis based on SMT solving.
  - ▶ Automated testing of distributed manufacturing execution systems.

# Conclusion

- Work presented in the thesis:
  - ▶ Counterexample analysis and abstraction refinement for Forest automata.
  - ▶ Forester participation in SV-COMP, editions 2015-2018.
  - ▶ A chapter on forest automata in book on software verification with refined presentation of the approach.
  - ▶ Automata over graphs with bounded tree-width.
  - ▶ Shape analysis based on SMT solving.
  - ▶ Automated testing of distributed manufacturing execution systems.
- Work not presented (or only mildly mentioned) in the thesis:
  - ▶ Connection of Predator and Symbiotic participating in [SV-COMP'20](#).
  - ▶ The design and implementation of efficient automata library called MATA ([TACAS'24](#)).

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# The Questions of Reviewers

- Is a user-non-friendliness of library for finite automata a problem if it used as backend of verification tool?
  - ▶ Context: In conclusion of my thesis I proposed that one of the future directions is to design a new efficient, simple, and user-friendly automata library.
  - ▶ Depends on definition of user-friendliness.
  - ▶ In general, if you work in small team with limited amount of resources you want library that has a clear interface and predictable behaviour, which is a kind of user friendliness.

# The Questions of Reviewers

- When generating abstract messages, could the abstraction lead to generation of messages that really do not make sense in practice? Is it still worth to examine the behaviour of the system after transferring such messages?
  - ▶ Context: In work on testing MES we proposed methods for generating new test cases using abstraction over messages and communication models.
  - ▶ Yes, it can lead to messages or sequences of messages that does not make sense but the system under testing should not fail when these messages.
  - ▶ From practical point of view, it is necessary to tune the abstractions to avoid generating such things.