

## 1 Question 1

As the graph  $G$  is composed of two connected components, we just have to count the number of edges in each connected component:

- The first component is a complete graph with 100 vertices, therefore there are  $\sum_{k=0}^{100-1} k = \frac{99*100}{2} = 4950$  edges in that component.
- In the second component, there are two partition set, with 50 nodes each which are linked to all the nodes of the other partition set. Therefore there are  $50 * 50 = 2500$  edges in that component. (It can be seen as a linear layer of neural network without bias)

So, in total there are **7450** edges in the graph.

Similarly, we count the number of triangles in each component:

- A triangle is just a triplet, therefore, counting the number of triangles is equivalent to counting the number of triplets in a set of 100 nodes: this is exactly  $\binom{100}{3} = \frac{100*99*98}{6} = 161,700$  triangles.
- In the complete bipartite component, since there are no edges between the nodes from the same partition set, there are no triangles.

So, in total there are **161,700** triangles in  $G$ .

## 2 Question 2

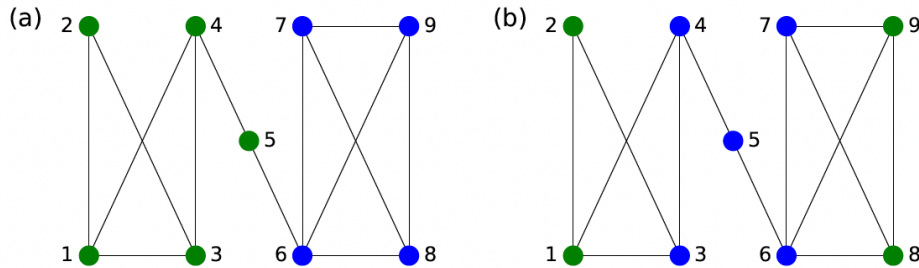


Figure 1: Two graphs where nodes have been assigned to 2 clusters. Cluster membership is indicated by node color.

In both graph (a) and (b),  $m = 13$ ,  $n_c = 2$ ,

- For the graph (a),  $l_{green} = 6$ ,  $l_{blue} = 6$ ,  $d_{green} = 13$  and  $d_{blue} = 13$ . Then,

$$Q = \left( \frac{6}{13} - \left( \frac{13}{26} \right)^2 \right) + \left( \frac{6}{13} - \left( \frac{13}{26} \right)^2 \right) = \frac{11}{26} \approx 0.42$$

- For the graph (b),  $l_{green} = 2$ ,  $l_{blue} = 4$ ,  $d_{green} = 11$  and  $d_{blue} = 15$ . Then,

$$Q = \left( \frac{2}{13} - \left( \frac{11}{26} \right)^2 \right) + \left( \frac{4}{13} - \left( \frac{15}{26} \right)^2 \right) = -\frac{17}{338} \approx -0.05$$

This results seems to be coherent with the natural idea that the clustering from graph (a) is better than the one from graph (b).

### 3 Question 3

If we note  $k$  the shortest path kernel, then there exists a function  $\phi$  such that:  $k(G_1, G_2) = \langle \phi(G_1), \phi(G_2) \rangle$ . The function  $\phi$  counts the frequency of shortest path distance in the graph i.e.  $\phi(G)[i]$  counts the number of shortest paths of distance  $i + 1$  in the graph.

Therefore,  $\phi(P_4) = [3, 2, 1]$  and  $\phi(C_4) = [4, 4, 0]$ .

Finally,  $k(C_4, C_4) = 32$ ,  $k(C_4, P_4) = 20$  and  $k(P_4, P_4) = 14$ .

### 4 Question 4

As the elements of the vectors  $f_G, f_{G'}$  are all non negative, if the kernel value :  $k(G, G') = f_G^T f_{G'} = 0$ , then it means that there are no types of graphlets of size 3 that both graphs  $G$  and  $G'$  share.

If we take  $G_1 = C_3$ , a complete graph with 3 nodes and  $G_2 = P_3$ , a path graph with 3 nodes. Then,  $f_G = [1, 0, 0, 0]$  and  $f'_G = [0, 1, 0, 0]$ , therefore  $k(G, G') = f_G^T f_{G'} = 0$ , as both graphs don't share a common 3-node subgraph.