

Title

subtitle

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24th March, 2018

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1 Introduction

2 Basic Definition

We first recall definitions of linear transformation and translation, and further introduce affine transformation, which we will use to prove stability of uniform mesh refinement and bisection in sections later. Then we introduce basic notions which we will use to present mesh refinement and show how they would work out.

2.1 Translation, Linear Transformation and Affine Transformation(pf for stability)

We define several classes of transformations that we use frequently. linear, translation, affine transformation

Definition. Let \mathbf{v} be a fixed vector, a **Translation** T_v on a figure applies as $T_v(\mathbf{p}) = \mathbf{p} + \mathbf{v}$, for a vector \mathbf{p} in the figure which we translate.

A *translation* moves every point of a figure or space by the same distance in the same direction. A translation T can be represented by an addition of a constant vector to every point.

Definition. Let V and W be vector spaces over the same field \mathbf{K} . We say a function $f : V \rightarrow W$ is **linear transformation** if the following satisfied:

$$\begin{aligned} f(\mathbf{u} + \mathbf{v}) &= f(\mathbf{u}) + f(\mathbf{v}) & \forall \mathbf{u}, \mathbf{v} \in V, \\ f(c\mathbf{u}) &= cf(\mathbf{u}), & \forall \mathbf{u} \in V, c \in \mathbf{K}. \end{aligned}$$

In other words, a linear transformation is a mapping between two vector spaces which preserves the operations of addition and scalar multiplication. Moreover, if V and W are finite dimensional, we can represent the linear transformation f by a matrix M . For example, if M is an $m \times n$ matrix, then f is a linear transformation from \mathbf{R}^n to \mathbf{R}^m .

Definition. An **affine transformation** from \mathbf{R}^n to \mathbf{R}^n is of the form

$$F(x) = Ax + v, \quad x \in \mathbf{R}^n,$$

where A is a linear transformation $\in \mathbf{R}^{n \times n}$, and v is a translation vector in \mathbf{R}^n

Affine transformation preserves points, lines and planes, but need not preserve point zero in a linear space in contrast to linear transformation. So we see that translation and linear transformation is affine, but the opposite is not true.

The inverse mapping $F^{-1} = x \mapsto A^{-1}(x - v)$ is also an affine transformation.

Clearly, we see that an affine transformation F is invertible if and only if A is invertible by the way we defined F^{-1} . Affine transformation helps carry results from one simplex to another simplex in our discussion, and more details are covered after introducing *Simplices* and *Triangulations* in the next section.

2.2 Simplices

simplices, subsimplices [DONE]

affine set, affine transformation [DONE] how to affine transformation act on simplices [UPDATED]

Definition. A **k -simplex** $T \in \mathbb{R}^n$ is a k -dimensional convex hull of $k + 1$ vertices $x_0, \dots, x_k \in \mathbb{R}^n$, which are affinely independent.

$$\begin{aligned} T &= [x_0, \dots, x_k] \\ &:= \left\{ x = \sum_{i=0}^k \lambda_i x_i \mid \sum_{i=0}^k \lambda_i = 1 \text{ and } 0 \leq \lambda_i \leq 1, 0 \leq i \leq k \right\} \\ &:= \left\{ \lambda_0 x_0 + \dots + \lambda_k x_k \mid \sum_{i=0}^k \lambda_i = 1 \text{ and } 0 \leq \lambda_i \leq 1, 0 \leq i \leq k \right\}. \end{aligned}$$

If $k = n$, we can call k -simplex without addressing the dimension. 2-simplices are also called *triangle*, and 3-simplices are called *tetrahedra*.

Definition. An l -simplex $S = [y_0, \dots, y_l]$ is called an **l -subsimplex** of k -simplex $T = [x_0, \dots, x_k]$, if indices $0 \leq i_0 \leq \dots \leq i_l \leq k$ with $y_i = x_{i_i}$, for $0 \leq l \leq k \leq n$.

Since there are $k+1$ vertices in k -simplex T , and $l+1$ vertices in l -subsimplex S , the number of l -subsimplex of k -simplex is $\binom{k+1}{l+1}$.

Consider simplices $T = [x_0, \dots, x_k]$ and $T' = [y_0, \dots, y_k]$. We say these two simplices T and T' are equal, i.e. $T = T'$, if $x_i = y_i$ for $0 \leq i \leq k$. Note that the vertex ordering of simplex is fixed, so if two simplices T and T' denote the same set but with different vertex ordering, they are not equal; instead, we say that T coincides with T' , i.e. $T \cong T'$.

Simplex under Affine Transformation

Instead of taking a single variable $x \in \mathbb{R}^n$ for affine transformation, we can take a subset $S \subset \mathbb{R}^n$, which contains $x \in \mathbb{R}^n$. Then the transformed set S' ,

$$S' = F(S) = \{F(x) \mid x \in S\}.$$

Similarly, if we regard a k -dimensional simplex $T = [x_0, \dots, x_k]$ as a subset $\in \mathbb{R}^n$, then the image of T under affine transformation, denoted by T' , is defined as below

$$\begin{aligned} T' &= F(T) \\ &= \{F(x_0), \dots, F(x_k)\}. \end{aligned}$$

We can see that T' is still a k -dimensional simplex. Furthermore, we can define $F(T) = AT + v$, where T is k -dimensional simplex $\in \mathbb{R}^n$. We might be curious about the relationship between simplices T and T' . Since vertices of a simplex are in a specific given order, so different vertex ordering leads to different simplices. Therefore, there exists an unique affine transformation such that $T = T'$. Another important property of simplices is congruence.

Definition. Two simplices T, T' are defined to be congruent if they can be obtained from each other by rotation, mirroring, scaling, and translation, i.e. if there exists a scaling factor $c \neq 0$, a translation vector $v \in \mathbb{R}^n$, and an orthogonal matrix $Q \in \mathbb{R}^{n \times n}$ such that

$$T' \cong v + cQT$$

When the two simplices T, T' have same vertices but with different vertex ordering, we can translate it as $T' = F(T)$. Based on how we define the affine transformation, it's not hard to see that $T \cong T'$. Then we say that T and T' are in a same congruent class.

2.3 Shape Regularity Measure

Shape measure offers an objective mathematical measure on the overall quality of an Finite Element mesh, and this is helpful to explore the simplex regularity and to improve the quality of shapes of the elements. Different definitions are used for shape measure to present the quality of simplex, and we simply introduce the *geometric shape measure* $\mu(T)$ of simplex T , the one we use in this paper.

Simplex Diameter and Volume

Let $T \in \mathbb{R}^n$ be a k -simplex where $k \leq n$, with vertices $x_0, \dots, x_k \in \mathbb{R}^n$. We let $\text{diam}(T)^k$ denote the diameter of T , and we define

$$\text{diam}(T)^k = \max \|x_i - x_j\|, \quad 0 \leq i \leq j \leq k.$$

In other words, $\text{diam}(T)^k$ is the longest distance between two vertices of T , which is equivalent to the length of the longest edge of T . If T is a single vertex, then $\text{diam}(T) = 0$.

Let $\text{vol}^m(T)$ denote k -dimensional volume of T , and we have

$$\text{vol}^k(T) = \frac{1}{k!} \cdot |\det(x_1 - x_0, x_2 - x_0, \dots, x_k - x_0)|$$

If $k = 0$, then T is a 0-dimensional simplex, i.e. a vertex. Then clearly, $\text{diam}(T) = 0$, and $\text{vol}^0(T) = 1$, the vol^0 of a single vertex is one.

Shape Measure

Simplex diameter and volume are important to introduce shape measures. Here we define the geometric shape measure $\mu(T)$ of simplex T by

$$\mu(T) = \frac{\text{diam}(T)^k}{\text{vol}^k(T)}, \quad \text{vol}^k(T) \neq 0$$

If $\text{vol}^k(T) = 0$, then we define $\mu(T) = \infty$.

To understand this definition, we can translate $\mu(T)$ as a measurement of how different the two variables, i.e. $\text{diam}(T)^k$ and $\text{vol}^k(T)$, are. For example, for a 3-dimensional simplex, i.e. a triangle, shape measure helps measure how narrow the triangle is. In other words, it measures how small the smallest angle of the triangle is.

Stability of a Simplex

The reason why we need shape measure is to help understanding whether a simplex T is non-degenerate, and to quantify how degenerate or non-degenerate. Let T be a k -dimensional simplex in \mathbb{R}^n . We say that a simplex T is degenerate if and only if $\mu(T) = \infty$, i.e. $\text{vol}_k(T) = 0$ approximating to 0.

Observing two pictures above, we actually want the interior angles of the simple T , i.e. triangles in this example, to be uniformly bounded from zero. Thus $\text{vol}_2(T)$ will never go to 0. While cutting a simplex into smaller pieces, we want to keep those pieces uniformly bounded and avoid degenerate simplices.

[ADDED LEMMA1: Congruent T1 T2 have same shape measure - ADDED]
[Q1. CHECK]

Lemma. *If T, T' are simplices that are congruent to each other, then $\mu(T) = \mu(T')$.*

Proof.

Since T is congruent to T' , by definition, we have $T' \cong v + cQT$, where c is scaling factor, v is a translation vector and Q is an orthogonal matrix. In fact, we will show that scaling, translation, rotation or mirroring does not influence the shape measure of a simplex.

To be specific, when scaling a simplex T by a non zero factor c to obtain T' , we have $\text{vol}^k(T') = \frac{1}{k!} \cdot |\det(cx_1 - cx_0, cx_2 - cx_0, \dots, cx_k - cx_0)| = \frac{c}{k!} \cdot |\det(x_1 - x_0, x_2 - x_0, \dots, x_k - x_0)| = c \cdot \text{vol}^k(T)$. Since it scales over all vertices, $\text{diam}(T')^k = c \cdot \text{diam}(T)^k$. Therefore, we see $\mu(T') = \frac{\text{diam}(T')^k}{\text{vol}^k(T')} = \frac{c \cdot \text{diam}(T)^k}{c \cdot \text{vol}^k(T)} = \frac{\text{diam}(T)^k}{\text{vol}^k(T)} = \mu(T)$. Moreover, translation over simplex T by a nonsingular vector v to obtain T' will not influence the shape measure as well. Under translation, the ordering of vertices may be changed. However, by definition of simplex diameter and volume, difference in vertex ordering will not affect the absolute value of the determinant nor the longest distance between two vertices. Thus, we obtain that $\mu(T) = \frac{\text{diam}(T)^k}{\text{vol}^k(T)} = \frac{\text{diam}(T')^k}{\text{vol}^k(T')} = \mu(T')$.

Similarly, rotations and mirroring represented by an orthogonal matrix Q will change the ordering of vertices, but again, simplex volume and diameter are independent of vertex ordering. Thus, we still obtain that $\mu(T) = \mu(T')$. Now we see the shape measure is independent of scaling, translation, rotation or mirroring. Thus a simplex T' which is obtained by these motions shares a same shape measure with T .

[ORIGINAL:] First, we show that volumes of k -simplices T and T' are the same, i.e. $\text{vol}^k(T) = \text{vol}^k(T')$. By definition of congruence, we can see one difference between equality and congruence of two simplices is that congruence is independent of vertex ordering. Recall the definition of simplex volume.

$\text{vol}^k(T) = \frac{1}{k!} \cdot |\det(x_1 - x_0, x_2 - x_0, \dots, x_k - x_0)|$. Clearly, permutation of vertex x_i , where $0 \leq i \leq k$, will not influence the absolute value of the determinant. Both dividing by $k!$, we have $\text{vol}^k(T) = \text{vol}^k(T')$.

Moreover, by definition of congruence above, we can obtain T' by applying rotation, mirroring, scaling, and translation over T . Thus, $\text{diam}(T')^k$, the longest distance between two vertices of T' , is equal to the longest distance between two vertices of T , $\text{diam}(T)^k$. Therefore, we have $\mu(T) = \frac{\text{diam}(T)^k}{\text{vol}^k(T)} = \frac{\text{diam}(T')^k}{\text{vol}^k(T')} = \mu(T')$. \square

2.4 Simplicial Complex(OR Triangulation)

Definition. A *simplicial complex* \mathcal{T} in \mathbb{R}^n is a finite set of simplices in \mathbb{R}^n that satisfies the following:

1. Any face of a simplex from \mathcal{T} is also in \mathcal{T} .
2. The intersection of any two simplices $T_1, T_2 \in \mathcal{T}$ is a face of both T_1 and T_2 .

In other words, the first condition asks \mathcal{T} to be closed under subsimplices, and the second condition asks that the intersection of any two simplices is either a common subsimplex or empty because the empty set is a face of every simplex. Examples in 2D are shown in Figure.1.

Any subset $T' \in \mathcal{T}$ that is itself a simplicial complex is called a *subcomplex* of \mathcal{T} . A *simplicial k -complex* \mathcal{T} is a simplicial complex where the largest dimension of any simplex in \mathcal{T} is k . So a simplicial 2-complex must not contain tetrahedra or higher dimension simplices. The 0-complex of \mathcal{T} is called a *vertex set* of \mathcal{T} . We can also think simplicial complex as a space with a triangulation, which

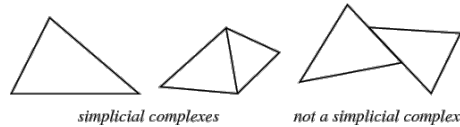


Figure 1: Simplicial Complex Example and Counterexample

is the division of a surface or a plane polygon into a set of 2-simplices. The constraints of triangulation will be discussed in ? section.

Simplicial Complex under Affine Transformation

Extending further from simplex under affine transformation, now we know that simplicial complex is just a finite set of simplices. Therefore, we can define the Transformed Simplicial Complex $F(\mathcal{T})$ as follows

$$F(\mathcal{T}) := \{F(T) \mid T \in \mathcal{T}\}$$

If \mathcal{T} is consistent, then $F(\mathcal{T})$ is also consistent by inheriting this property from \mathcal{T} .

Shape Measure of Simplicial Complex [DELETE??]

We define shape measure of a simplex $T \in \mathbb{R}^n$ as $\mu(T) = \frac{\text{diam}(T)^k}{\text{vol}^k(T)}$. Now consider a simplicial complex $\mathcal{T} \in \mathbb{R}^n$, we define the geometric shape measure $\mu(\mathcal{T})$ as follows,

$$\mu(\mathcal{T}) = \max_{T \in \mathcal{T}} \mu(T)$$

By definition, we see that the shape measure of a simplicial complex \mathcal{T} is the supreme of the set of shape measures of all simplex $T \in \mathcal{T}$. If the largest shape measure of a simplex in this simplicial complex is bounded, then none of simplices in \mathcal{T} is degenerate. In other words, if simplex $T_0 \in \mathcal{T}$ is non-degenerate, then simplicial complex \mathcal{T} non-degenerate. [Correct?? Pf needed??]

3 Refinement Strategy in General

Refinement is a procedure of mesh modification in which we can divide a regular domain into smaller pieces under boundary constraints. This process can be applied recursively to simplify some differential equations by generating smaller pieces of the domain. Let us first introduce triangulation to help understand refinement on a simplex. Generally speaking, we can think triangulation as a subdivision of a plane into triangles. Definition below is a more formal way to take when extending to higher dimension.

Definition. A triangulation of \mathbb{R}^n is subdivision into n -dimensional simplices such that intersection of any two simplices is either empty or sharing a common face, and any face of a simplex is in the triangulation.

Indeed, we say that this triangulation is consistent as it is not simply subdividing of a space. Moreover, the triangulation defined here can be treated equivalently as simplicial complex as it is a finite set of simplices satisfying

1. Any face of a simplex from a triangulation is also in the triangulation,
2. The intersection of any two simplices T_1, T_2 in a triangulation is a face of

both T_1 and T_2 or empty. (Denote triangulation same as simplicial complex \mathcal{T})

We can think a refinement of a simplex T as a triangulation \mathcal{T} which consists of smaller pieces of simplices of same type of the simplex T . Now consider a refinement of a simplicial complex. Let \mathcal{T} and \mathcal{T}' be two different simplicial complex covering a same domain Ω . This means that the domain $\Omega = \bigcup(T|T \in \mathcal{T}) = \bigcup(T'|T \in \mathcal{T}')$. We say that \mathcal{T}' is a refinement of \mathcal{T} if each simplex $T \in \mathcal{T}$ is in \mathcal{T}' or the triangulation of T is in \mathcal{T}' .

As mentioned before, we may recursively apply a refinement strategy to help simplify some problems. By recursively taking refinement process from \mathcal{T}_0 , we have a hierarchy triangulation $\mathcal{T}_k, k \in \mathbb{N}$, where \mathcal{T}_k is a refinement of \mathcal{T}_{k-1} .

Definition. Let \mathcal{T}_0 be the initial simplicial complex in \mathbb{R}^n where it starts from, then we define the hierarchy triangulation \mathcal{T}_k as follows

$$\mathcal{T}_k := \bigcup \{\text{refinement of simplex } T \mid T \in \mathcal{T}_{k-1}\}, \quad k \in \mathbb{N}$$

3.1 Stability of Refinement

Definition. We say a refinement strategy is **stable** if there exists a constant $C > 0$ such that $\mu(T) < C$ for all simplices T .

Theorem. If the number of congruence classes, obtained by applying the refinement of a non degenerate simplex T initially, is finite, then the refinement strategy is stable.[UPDATED]

Proof. Idea:

1. \mathcal{T}_0 is non-degenerate, then \mathcal{T}_i is non-degenerate.

Claim. Refinement strategy over initial simplicial complex \mathcal{T}_i produces only non-degenerate simplices T .

Proof. We prove this claim by induction. Clearly, the base case is true since it is given that all simplices T in \mathcal{T} are non-degenerate. For induction, suppose simplices in simplicial complex \mathcal{T}_k is non-degenerate, i.e. given $C > 0, \mu(T) < C, \forall T \in \mathcal{T}_k$. Apply the refinement strategy on \mathcal{T}_k , and then we obtain $\mathcal{T}_{k+1} = \bigcup \{\text{refinement of simplex } T \mid T \in \mathcal{T}_k\}, \quad k \in \mathbb{N}$. [connection???: if simplex $T_0 \in \mathcal{T}$ is non-degenerate, then simplicial complex \mathcal{T} non-degenerate.] \square

Claim. If the number of congruence classes is finite, then the number of shape measure is finite, and there exist a common bound $C > 0$ such that $C \geq \mu(T)$.

Proof. We proved that simplices in same congruence classes share the same shape measure. If we have finite number of congruence classes, clearly we have finite number of shape measure. When all simplices are non degenerate, we always have an upper bound for their shape measure $\mu(T)$. With the finite number of shape measure, we may set C as the maximum of all upper bounds of shape measures. And therefore $C \geq \mu(\mathcal{T})$ \square

Since T_0 is non-degenerate, then \mathcal{T}_i is non-degenerate. Moreover, we know there exists a common bound C for all shape measures since the number of congruence classes is finite. Therefore, we proved the stability.[UPDATED] \square

2. Finite number of congruence classes, then finite number of bounds?
3. Summary: each \mathcal{T} is non-degenerate + finite number of \mathcal{T} = all \mathcal{T} is non-degenerate. By def, the refinement strategy is stable.
[Q2. NEED CORRECTION]

3.2 Consistency of Refinement

4 Uniform Refinement

[Al of uniform refinement in 2-d, show it's stable and consistent]

4.1 2D -[UPDATED]

One popular refinement strategy is *red/green* refinement proposed by R. E. Bank et al. The red refinement here is regular refinement which divides a triangle into four congruent smaller triangles by connecting midpoints of its three edges. The green refinement is irregular refinement which connects the refined edge midpoint to its opposite corner.[PIC-NEED TO BE UPDATE]

Let $T = [x^0, x^1, x^2]$ be the triangle to be refined, and denote x^{ij} as the midpoint of the edge between x^i and x^j .

Algorithm Red refinement in 2D {
 $T_1 := [x^0, x^{01}, x^{02}]; \quad T_2 := [x^{01}, x^1, x^{12}];$
 $T_3 := [x^{02}, x^{12}, x^2]; \quad T_4 := [x^{01}, x^{12}, x^{02}];$

} Subdividing a triangle by connecting its midpoint, we obtain four congruent simplices T_1, T_2, T_3 and T_4 . The green refinement is applied on one refined edge on which we may confront with degenerate simplices. We will never touch and refine such simplices to avoid them from degenerating.

Clearly, the red refinement strategy is stable since it produces a finite number of triangles congruent to the original simplex. Meanwhile, we preserve consistency by bisecting triangles with one refined edge and never refine them any further. Therefore, we obtain stability and consistency through red and green refinement.

4.2 3D

5 Bisection

FIGHTING!!! ALMOST THERE WOOOOOOOT WOOOOOOOOOT $\backslash(\bullet>\omega<\bullet)/$

[Bibliography - ADD] [CORRECT TEXTIT for T, F T'] -Simplicial complex is always consistent by definition as touch edge is always a subsimplex...
x-finite number of consistency class-¿finite number of shape measure-¿shape
measure is bounded lemma: simplex in one consistency class has same shape
measure(pf)

Prev Q:

1.
Lemma 1 [UPDATED]
Shape measure of simplicial complex [UPDATED]
2.
Simplex T degenerate [UPDATED]
Shape measure of simplicial complex [UPDATED]
3.
thm pf: finite number of congruence classes leads to stable refinement strategy [UPDATED]
4.
uniform refinement 2D[UPDATED]

Q:

1. More for Uniform refinement 2d? code?
2. bisection: newest vertex? longest?
[Lemma 1: If T_0 and T_1 in a same congruence class, then the shape measure of these two simplices are the same] Iff?? [If T_0 is non degenerate, then simplicial complex is non degenerate] [Thm] A refinement strategy is stable if and only if the number of congruence classes, obtained by applying the refinement of a non degenerate simplex T initially, is finite. [DELETED] [Al of uniform refinement in 2-d, show it's stable and consistent]
note: distinguish refinement and uniform refinement? regular, global = uniform, local = bisection?

References

- [1] Weisstein, Eric W. "Simplicial Complex."
<http://mathworld.wolfram.com/SimplicialComplex.html>