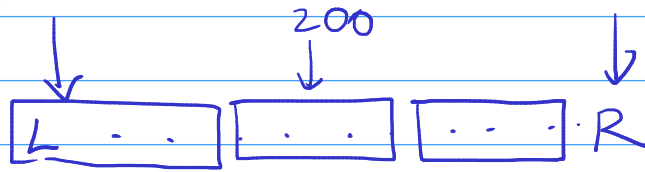


$$N = 30000$$

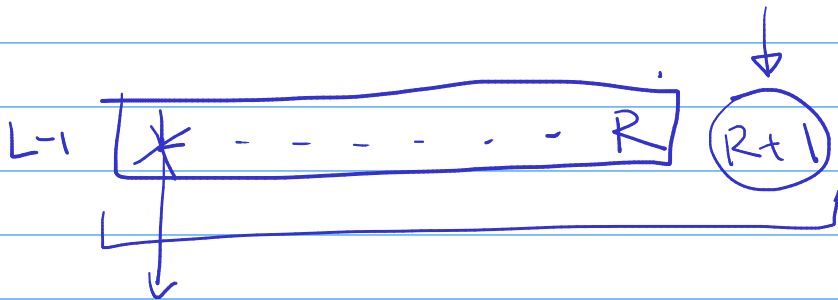
$$\sqrt{N} = \underline{200}$$

$$\underline{N\sqrt{N}},$$

$$\frac{N}{200} \quad 200 \quad q = 2 \times 10^5 \times 200 = \underline{4 \times 10^7} \quad \checkmark$$



MO's algorithm



Fractional Knapsack
0-1 Knapsack ✓

$$T = \underline{8 \text{ sec}}$$

$$R_1 = (\underline{100}, 7 \text{ sec}) = 14.28$$

$$R_2 = (\underline{50}, 4 \text{ sec}) = 12.5$$

$$R_3 = (\underline{55}, 4 \text{ sec}) = 13.75$$

$$f(i, k) = a_i - (k-1)^2 \cdot b_i$$

$$b_i = 1$$

$$\underline{30 \times 10^0}$$

$$k = \sqrt{a} = \underline{30}$$

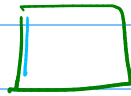
$$= \boxed{3000} \times 25000$$

$$= \underline{75 \times 10^6} \leftarrow \text{No. of itera}$$

$$= \underline{7.5 \times 10^7}$$

0-1 Knapsack

(W, N) ✓



$$\downarrow$$

$$\underline{dp[20][t]}$$

$$\downarrow$$

$$dp[19]$$

Ride i

Ride not taken

$$\underline{dp[i-1][t]}$$

$$\underline{dp[i][t]}$$

Ride Taken

$$\underline{dp[i-1][t-t_i]}$$

vector<int> dp[3005];

→ for(int i=1; i<=3004; i++) {

→ // dp[i] = vector<int> (25001, 0);

⇒ // dp[i-2].clear(); // 25000

for(int y=1; y<=25000; y++) {

{ dp[i][y] = - - - -

solve(int x, int y) → solve(n, y-1)
→ solve(x-1, y);

cout << solve(1000, 1000) << endl;

→ if dp[x][y] != -1 :

return dp[x][y];

solve(2,2)

solve(1,2)

solve(2,1)

for (x=0; x<1000; x++)

for (y=0; y<1000; y++)

solve(x, y);

cout << =

solve(0,0)

solve(1,1)

solve(0,1)

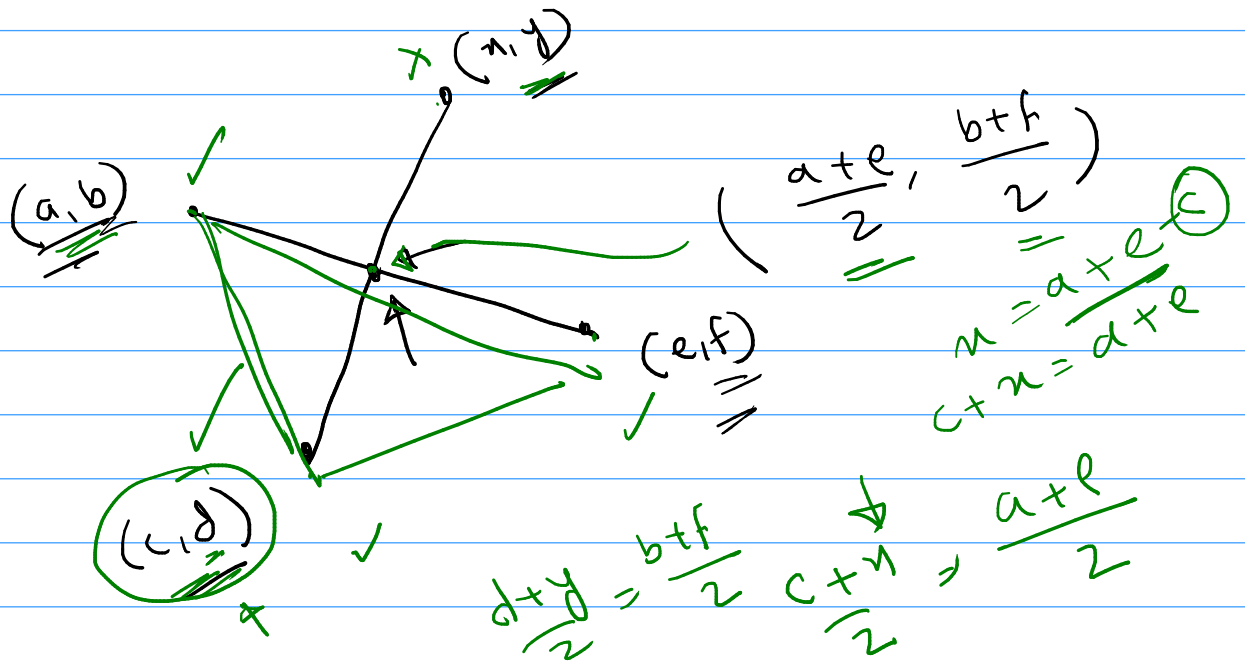
↓
x

solve(1,0)

↓
x

$\text{solve}(n, y) \rightarrow \text{solve}(n+1, y) \checkmark$
 $\text{solve}(n, y) \rightarrow \text{solve}(n, \underline{\underline{y-1}}) \checkmark$

Bitmask DP ✓



All submask of a mask \rightarrow $\textcircled{3^N}$
 \uparrow

$\underline{\underline{N}} \rightarrow \underline{\underline{2^N}}$
 \uparrow

$$s = (-s) \&$$

$$\underline{x} = \underline{c}_x + \underline{t}(\underline{d}_x - \underline{c}_x) \quad \leftarrow \text{parametric}$$

$$t=1, \quad x = c_x + d_x - c_x = d_x$$

$$t=0, \quad x = c_x$$

$$c(c_x, c_y) \quad t=0$$

$$D(d_x, d_y) \quad t=1$$

$$x_1 = c_x + t(d_x - c_x)$$

$$y_1 = c_y + t(d_y - c_y)$$

$$x_2 = a_x + t(b_x - a_x)$$

$$y_2 = a_y + t(b_y - a_y)$$

$$\underline{d}^2 = (\underline{x}_1 - \underline{x}_2)^2 + (\underline{y}_1 - \underline{y}_2)^2$$

$$\underline{d}^2(t) = \left(\underbrace{(c_x - a_x)}_p + t \underbrace{(d_x - c_x - b_x + a_x)}_q \right)^2$$

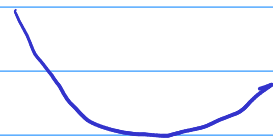
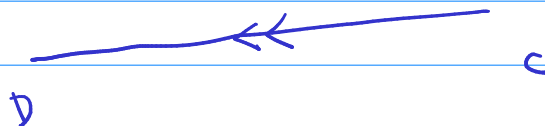
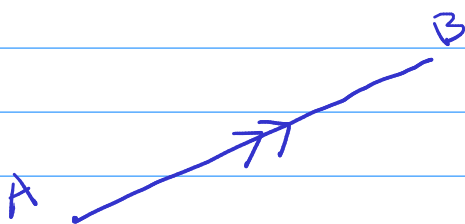
$$= \underline{(p + qt)^2} + \underline{(m + nt)^2}$$

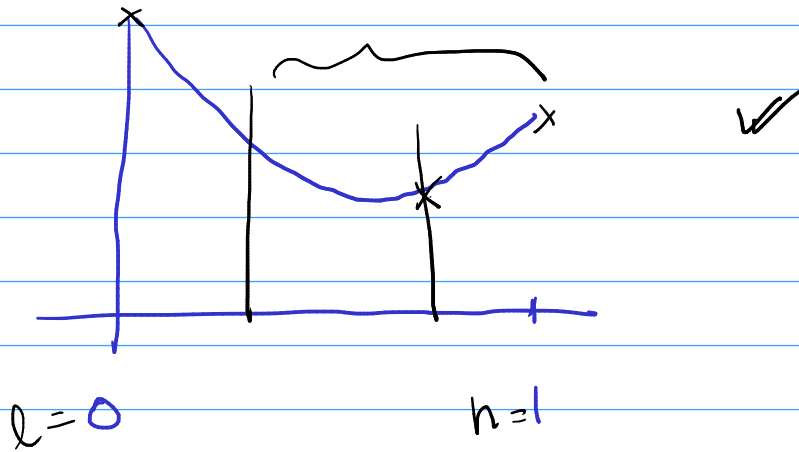
$$= \underline{p^2 + m^2} + t(2pq + 2mn) + t^2(q^2 + n^2)$$

$$\frac{d}{dt} \left[\underline{d^2} \right] = 0 + 2pq + 2mn + 2t(q^2 + n^2)$$

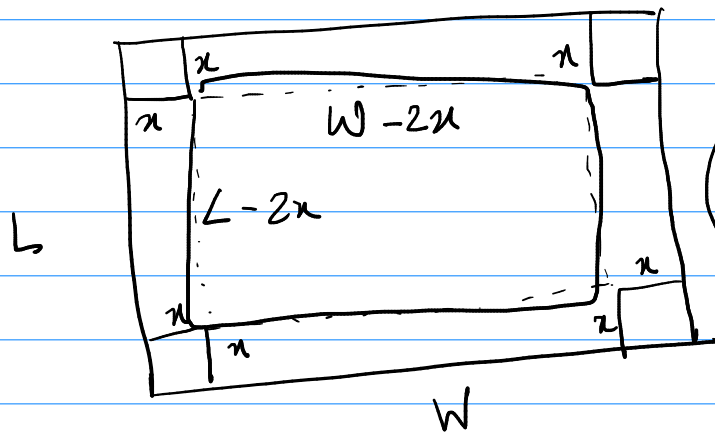
$$= 0$$

$$t = \frac{-pq - mn}{q^2 + n^2}$$





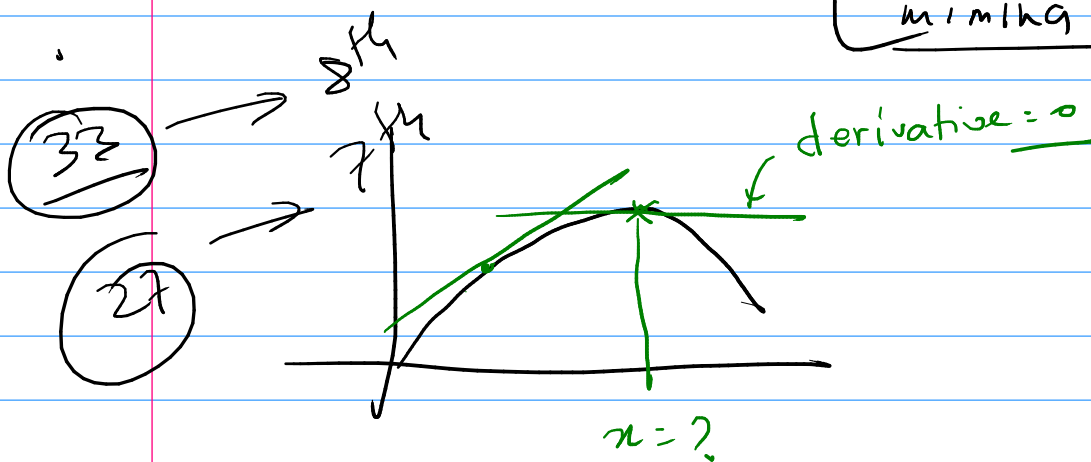
$$V = (W - 2n)(L - 2n)(n)$$

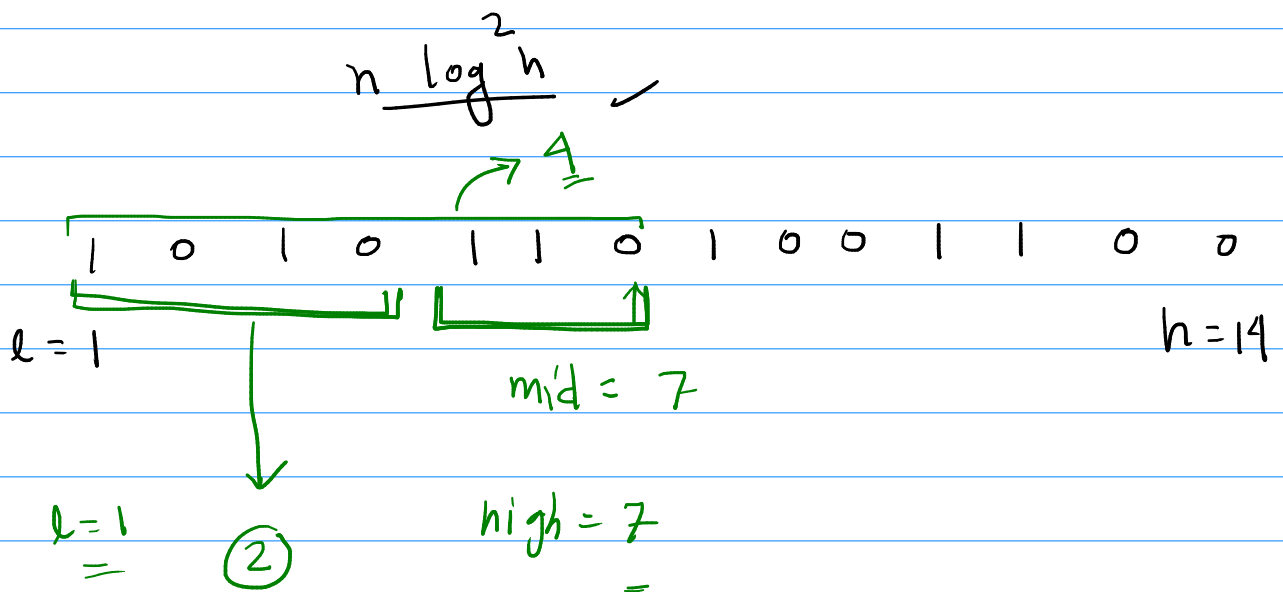
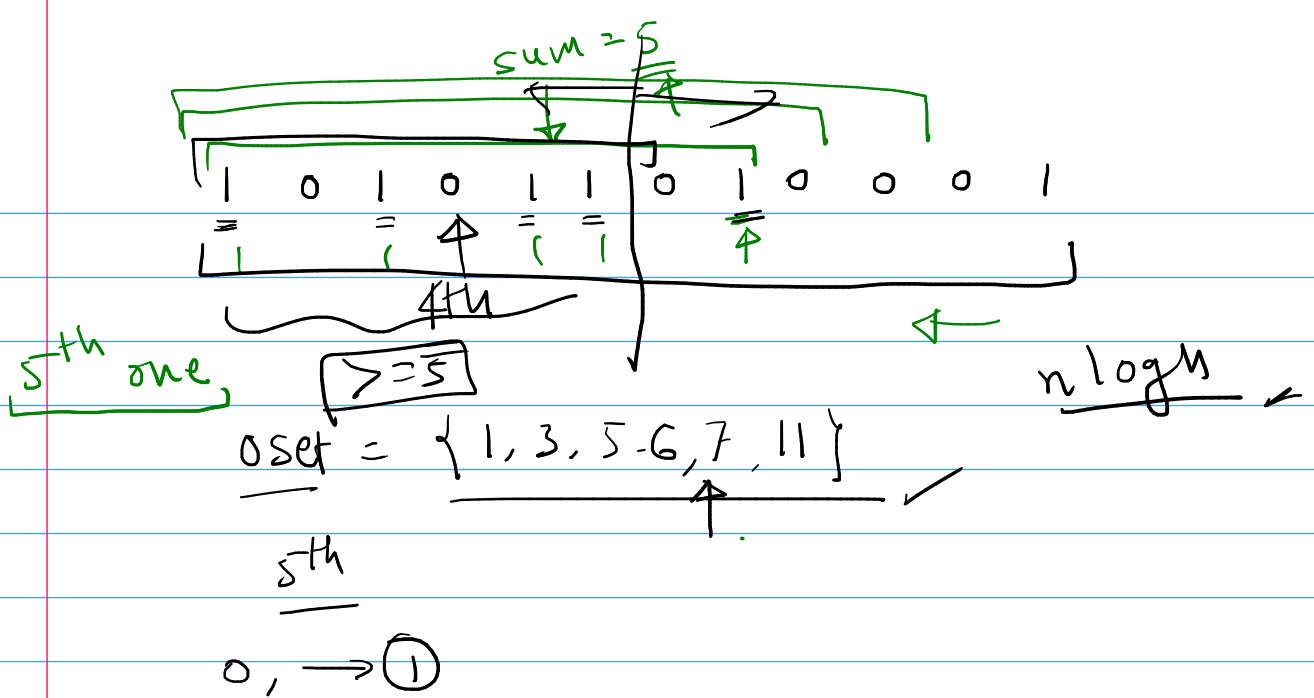


$$\frac{dV}{dn} = 0$$

$$\frac{\partial V}{\partial n} = 0$$

maxima
minima





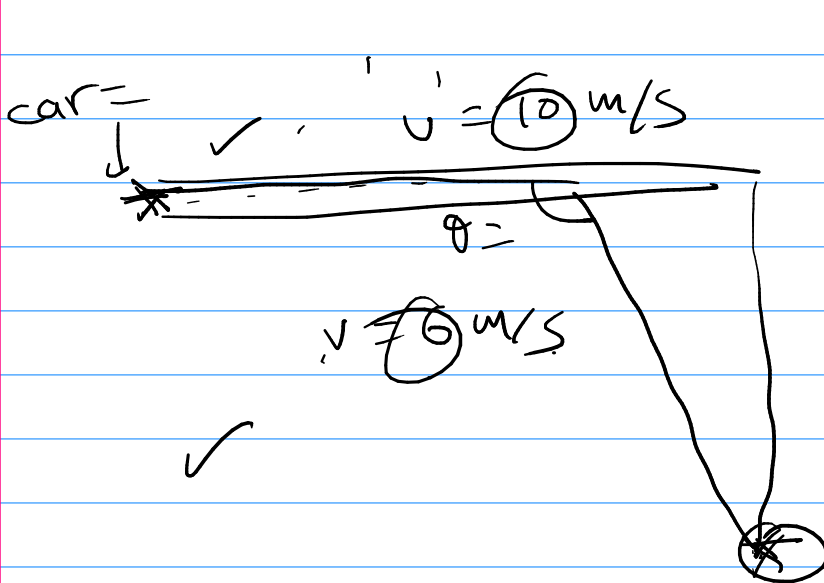
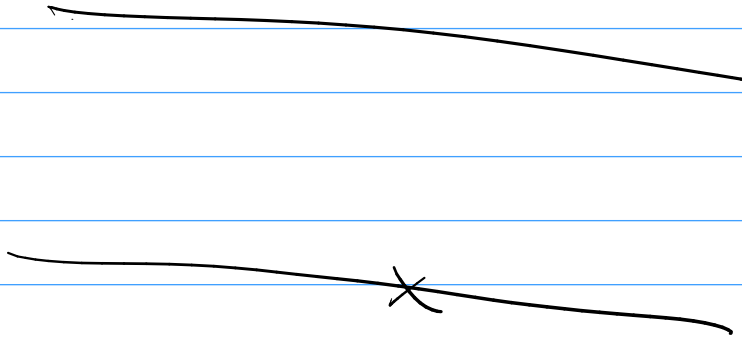
$K = 3$

Parse Tree

Sparse tree

RMQ $\rightarrow \underline{O(1)}$

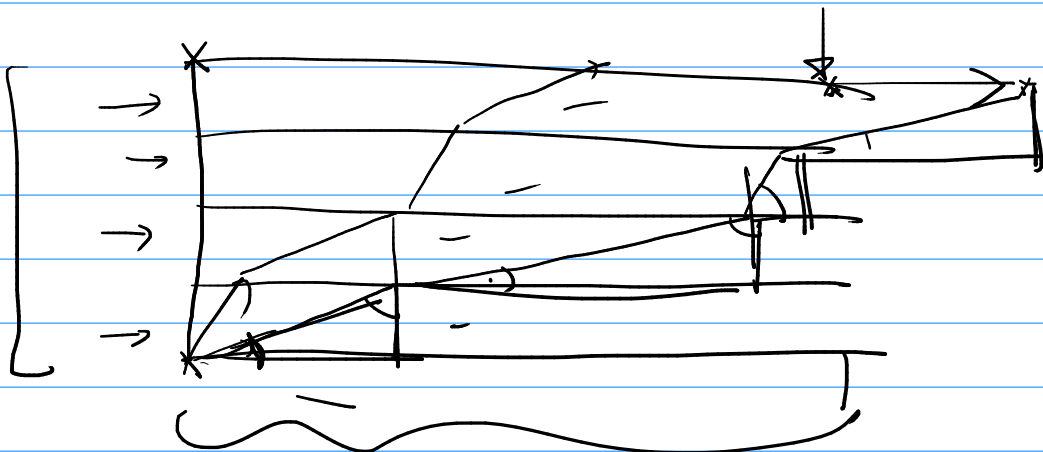
1 Rodou ←

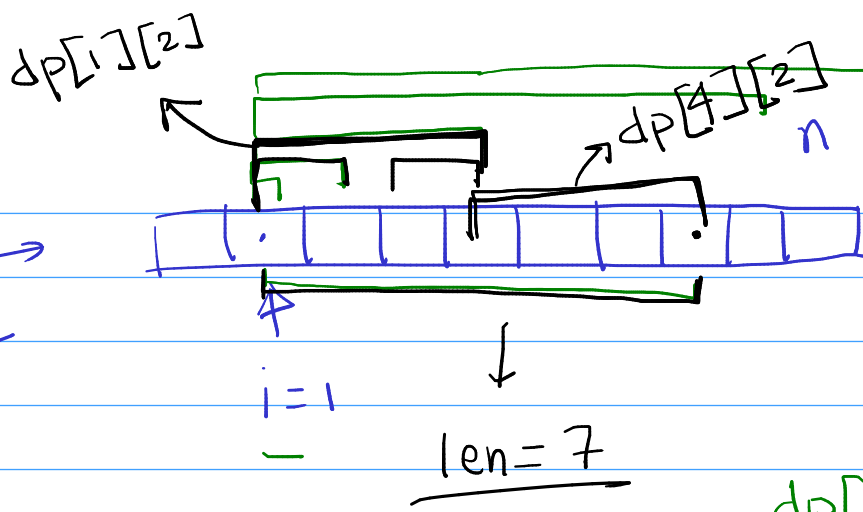


Kinematic

$\eta =$

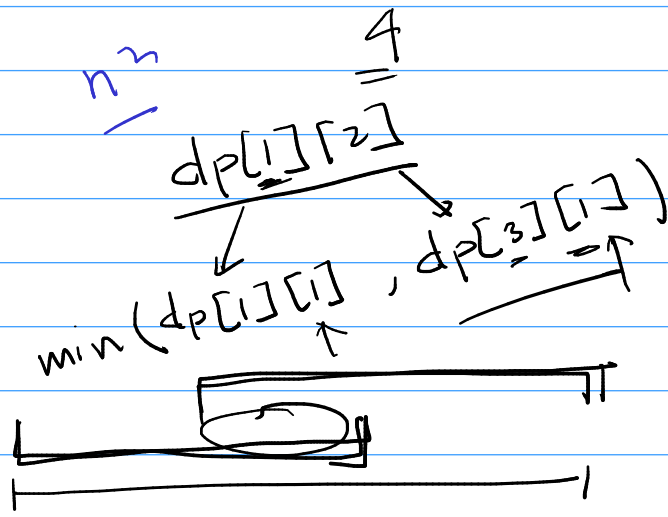
phonics





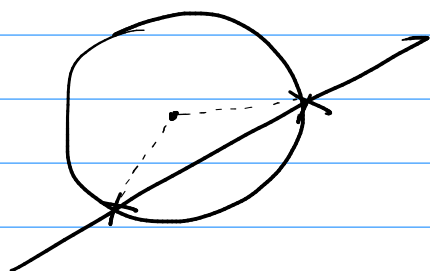
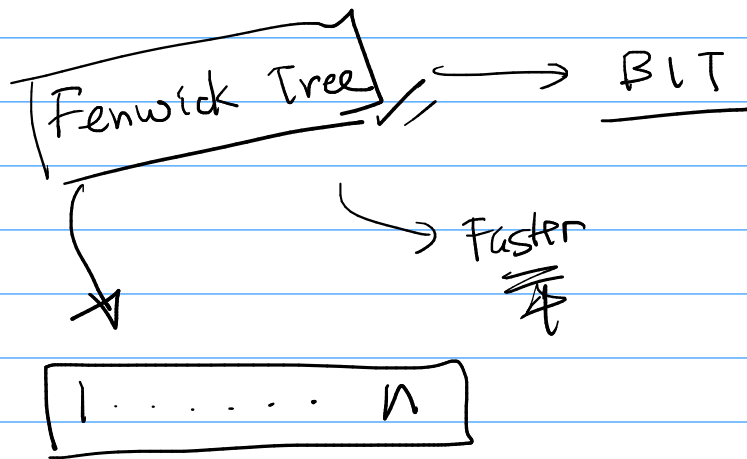
Seq Tree Mem comp
 $O(n)$
 \uparrow
 $\log n$

*
 Mem \rightarrow
 $n \log n$
 $O(1)$

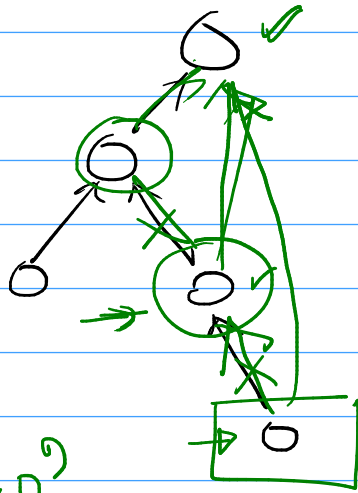


$dp[i][0]$ \rightarrow pos = 1
 leng = 2^0
 $dp[i][1]$ \rightarrow pos = 1
 leng = 2^1
 $dp[i][2]$ \rightarrow pos = 1
 leng = 2^2

min, max
gcd



parametric equation of
line



path compression

DSU

$\rightarrow 10^9$

$\frac{x \cdot n}{\downarrow}$
 $\frac{x < 4}{\downarrow}$

h^2

$\rightarrow \frac{a, b, c}{\text{given}}$

$$\boxed{ax + by = c}$$

$\boxed{x, y} \rightarrow \text{integer solution}$