

UAP_Nightjars CP Reference

University of Asia Pacific

Team: UAP_Nightjars (Maruf,Rhythm,Arafat)

Contents

| | | | | | |
|---|-----------|--|----|---|-----------|
| 1 Data Structure | 1 | 4.7 How Many Digits in $n!$ | 9 | 6.4 Largest Substring More than K | 15 |
| 1.1 BIT | 1 | 4.8 Derangement | 9 | 6.5 Longest Common Prefix Of Two Substrings | 16 |
| 1.2 BST | 2 | 4.9 Euler Phi | 10 | 6.6 Longest Common Substring | 16 |
| 1.3 Bit Binary Search | 2 | 4.10 Fibonacci Number Faster | 10 | 6.7 Pattern Matching | 16 |
| 1.4 Sparse Table | 2 | 4.11 GCD LCM | 10 | 6.8 Two Strings Are Equal Or Not | 17 |
| 1.5 DSU | 2 | 4.12 Josephus | 10 | 7 Mathematical Formulas | 17 |
| 1.6 Merge Sort Tree | 2 | 4.13 Large Number GCD | 10 | 7.1 Math Formula | 17 |
| 1.7 Monotonous Queue | 2 | 4.14 Lengenders Formula | 10 | 7.2 Catalan Numbers | 17 |
| 1.8 Ordered Set | 3 | 4.15 Modular for Subtraction & Multiplication . | 10 | 7.3 Geometry Formulas 1 | 18 |
| 1.9 Segment Tree All | 3 | 4.16 Optimized Sieve | 10 | 7.4 Geometry Formulas 2 | 18 |
| 1.10 Sparse Table | 4 | 4.17 Lower Prime Factorization | 11 | 7.5 Matrices and Determinants | 19 |
| 2 Dynamic Programming | 4 | 4.18 Prime or Not | 11 | 7.6 Vectors | 19 |
| 2.1 Longest Common Subsequence print using next state | 4 | 4.19 Number of Divisor A Number Have | 11 | 7.7 Straight Lines | 20 |
| 2.2 Longest Increasing Subsequence Print . . . | 4 | 4.20 Prime Factors | 11 | | |
| 3 Graphs | 5 | 4.21 Formulas Using Prime Factorization Math | 11 | | |
| 3.1 BFS | 5 | 4.22 Sum & Product of Divisor | 11 | | |
| 3.2 Bellman Ford | 5 | 4.23 Sieve | 11 | | |
| 3.3 Bipartite Garphs | 5 | 4.24 $n!$ Trailing Zero | 11 | | |
| 3.4 Cycle or not Print Cycle | 5 | 4.25 Number of Divisors 1 to n | 11 | | |
| 3.5 Cycle or not in Directed Graph | 6 | 4.26 NcR (when n and r to big)(Lucas Theorem) | 12 | | |
| 3.6 DFS | 6 | 4.27 how many numbers (0.. n) have k-th bit=1 . | 12 | | |
| 3.7 Dijkstra | 6 | 5 Mixed | | | |
| 3.8 Floyd Warshall | 7 | 5.1 Run Code Technique | | | |
| 3.9 Krushkals MST | 7 | 5.2 Converts Any Integer Number to Any Base | | | |
| 3.10 LCA | 7 | 5.3 i128 Bit | | | |
| 3.11 Prims MST | 8 | 5.4 Base Conberstion | | | |
| 3.12 Strongly Connected Components | 8 | 5.5 Bitset | | | |
| 3.13 Topological Sorting | 8 | 5.6 Bit Operation | | | |
| 3.14 Zero One BFS | 8 | 5.7 Merge Sort | | | |
| 4 Number Theory | 9 | 5.8 STL | | | |
| 4.1 2D Prefix Sum | 9 | 5.9 Sqrt Decomposition | | | |
| 4.2 Number of Integers Coprime with N (1 to n) | 9 | 5.10 Binary Search | | | |
| 4.3 Binary Exponentiation | 9 | 5.11 lower_bound & upper_bound | | | |
| 4.4 Binary Multiplication with Mod | 9 | 5.12 Ternary Search | | | |
| 4.5 Modular Arithmetic | 9 | 5.13 Custom Sorting | | | |
| 4.6 Binomial Coefficients nCr | 9 | 5.14 Permutation & Combination | | | |
| | | 5.15 python | 15 | | |
| 6 Strings | 15 | 6.1 All String Operation | 15 | | |
| | | 6.2 Single Hashing | 15 | | |
| | | 6.3 Single Hashing | 15 | | |

1. Data Structure

1.1. BIT

```
#include <bits/stdc++.h>
using namespace std;
const int N = 3e5 + 9;
struct BIT{ // 1-indexed
    int n;
    vector<int> t;
    BIT() {}
    BIT(int _n){
        n = _n;
        t.assign(n + 5, 0);
    }
    int qry(int i){
        int ans = 0;
        for (; i >= 1; i -= (i & -i))
            ans += t[i];
        return ans;
    }
    void upd(int i, int val){
        if (i <= 0)
            return;
        for (; i <= n; i += (i & -i))
            t[i] += val;
    }
    void upd(int l, int r, int val){
        upd(l, val);
        upd(r + 1, -val);
    }
    int qry(int l, int r)
```

```

    return qry(r) - qry(l - 1);
};

int32_t main(){
    int n, q;
    cin >> n >> q;
    BIT bit(n);
    bit.upd(x, 1); // demo
    bit.query(10); // ddemo, change this
    return 0;
}

```

1.2. BST

```

#include <bits/stdc++.h>
using namespace std;
const int N = 3e5 + 9;

// the code returns a BST which will create if we
// add the values one by one
// here nodes are indicated by values and every node
// must be distinct
set<int> se;
map<int, int> l, r; // l contains the left child of
// the node, r contains right child of the node
int main() {
    int n;
    cin >> n;
    int k;
    cin >> k; // root of the tree
    se.insert(k);
    for (int i = 1; i < n; i++) {
        int k;
        cin >> k;
        auto it = se.upper_bound(k);
        if (it != se.end() && l.find(*it) == l.end())
            l[*it] = k;
        else
            --it, r[*it] = k;
        se.insert(k);
    }
    for (int i = 1; i <= n; i++)
        cout << l[i] << ',' << r[i] << '\n';
    return 0;
}

```

1.3. Bit Binary Search

```

// --- Bit Binary Search in O(log(n)) ---
const int M = 20 const int N = 1 << M int
lower_bound(int val){int ans = 0, sum = 0;
for (int i = M - 1; i >= 0; i--){
    int x = ans + (1 << i);
    if (sum + bit[x] < val)ans = x, sum += bit[x];
    }return ans + 1; }

```

1.4. Sparse Table

```

// build time = nlog(n)
// query time = O(1)
// lg(n) = log^*(n)
struct SparseTable {
    vector<vector<int>> st;
    int n, K;
    SparseTable() {}
    SparseTable(vector<int> a) {
        n = a.size(); K = __lg(n) + 1;
        st = vector<vector<int>>(K, vector<int>(n));
        for (int i = 0; i < n; i++) st[0][i] = a[i];
        for (int i = 1; i < K; i++) {
            for (int j = 0; j + (1 << i) <= n; j++) {
                st[i][j] = min(st[i - 1][j], st[i - 1][j
                    + (1 << (i - 1))]);
            }
        }
        int get(int L, int R) {
            int i = __lg(R - L + 1);
            return min(st[i][L], st[i][R - (1 << i) + 1]);
        }
    }
};

```

1.5. DSU

```

int connected_components;
int maximum_size;
struct DSU {
    vector<int> parent, component_size;
    // O(N)
    DSU(int _n) {
        parent = vector<int>(_n);
        component_size = vector<int>(_n);
        for (int i = 0; i < _n; i++) {
            parent[i] = i;
            component_size[i] = 1;
        }
    }
    int find_parent(int a) {
        if (parent[a] == a) return a;
        else {
            int temp = find_parent(parent[a]);
            parent[a] = temp;
            return parent[a];
        }
    }
    bool unite(int a, int b) {
        a = find_parent(a);
        b = find_parent(b);
        if (a == b) return false;
        connected_components--;

```

```

        if (component_size[b] > component_size[a]) swap
            (a, b);
        component_size[a] += component_size[b];
        maximum_size = max(maximum_size,
            component_size[a]);
        parent[b] = a; // now the components of a and
        b have been merged
    }
    return true;
}
};


```

1.6. Merge Sort Tree

```

// Mergesort Tree - Time <O(nlogn), O(log^2n)> -
// Memory O(nlogn)
// Mergesort Tree is a segment tree that stores the
// sorted subarray
// on each node.

vi st[4 * N];
void build(int p, int l, int r) {
    if (l == r) {
        st[p].pb(s[l]);
        return;
    }
    build(2 * p, l, (l + r) / 2);
    build(2 * p + 1, (l + r) / 2 + 1, r);
    st[p].resize(r - l + 1);
    merge(st[2 * p].begin(), st[2 * p].end(),
        st[2 * p + 1].begin(), st[2 * p + 1].end(),
        st[p].begin());
}

int query(int p, int l, int r, int i, int j, int a,
        int b) {
    if (j < l || i > r)
        return 0;
    if (i <= l & j >= r)
        return upper_bound(st[p].begin(), st[p].end(),
            b) -
            lower_bound(st[p].begin(), st[p].end(), a
            );
    return query(2 * p, l, (l + r) / 2, i, j, a, b) +
        query(2 * p + 1, (l + r) / 2 + 1, r, i, j, a
            , b);
}

```

1.7. Monotonous Queue

```

#include <bits/stdc++.h>
using namespace std;
const int N = 3e5 + 9;
struct monotonous_queue { // max, stores strictly
    decreasing sequence of the current queue

```

```

int a[N + 10], b[N + 10], l = 0, r = -1;

void push(int val) {
    int cnt = 0;
    while (l <= r && a[r] <= val) {
        cnt += b[r] + 1;
        r--;
    }
    a[++r] = val;
    b[r] = cnt;
}
int top() {
    return a[l];
}

void pop() {
    if (l > r)
        return;
    if (b[l] > 0) {
        b[l]--;
        return;
    }
    l++;
}
});
```

1.8. Ordered Set

```

// Find by Order (find_by_order) ->O(log(n));
// Order of Key (order_of_key)->O(log(n));
// Deletion (erase)->O(log(n));

#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds; //additional line

// Define an indexed set
typedef tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> indexed_set;

//greater(x ar the koto guia boro value ache)
//greater_equal(for multiset)
// typedef tree<int, null_type, less_equal<int>,
rb_tree_tag, tree_order_statistics_node_update>
indexed_set;

//for multiset or same value count
// typedef tree<pair<int, int>, null_type, less<pair
<int, int>>, rb_tree_tag,
tree_order_statistics_node_update> indexed_set;

// s.find_by_order(1)->first
// s.order_of_key(make_pair(20, 2))

int main() {
    indexed_set s;
    // Insert elements
```

```

s.insert(10);

// Find the 2nd smallest element (0-based
// indexing)
cout << "2nd smallest: " << *s.find_by_order(1)
<< "\n";

// Find the number of elements less than 20
cout << "Number of elements < 20: " << s.
order_of_key(20) << "\n";

// Check if an element exists
if (s.find(15) != s.end()) {
    cout << "15 is present in the set.\n";
}

// Erase an element
s.erase(15);
return 0;
}
```

1.9. Segment Tree All

```

#include <bits/stdc++.h>
#define int long long
using namespace std;
const long long inf = 1e18;
struct SegmentTree {
    vector < int > stree, lazy;
    int sz;
    vector < int > v;
    SegmentTree (vector < int > &ara) {
        v = ara;
        sz = ara.size();
        stree = vector < int > (4 * sz, inf);
        lazy = vector < int > (4 * sz, 0);
        build(1, 0, sz - 1);
    }
    int merge(int v1, int v2) {
        return min(v1, v2);
    }
    int size() {
        return sz;
    }
    void push (int u) {
        int val = lazy[u];
        lazy[u] = 0;
        stree[u * 2] += val;
        stree[u * 2 + 1] += val;
        lazy[u * 2] += val;
        lazy[u * 2 + 1] += val;
    }
    // Definition : lazy propagation for range update
    // (SUM update)
    void push(int u, int tl, int tr) {
        if(lazy[u] == 0 || tl == tr) return;
        int mid = (tl + tr) / 2;
```

```

        stree[u * 2] += lazy[u] * (mid - tl + 1);
        lazy[u * 2] += lazy[u];
        stree[u * 2 + 1] += lazy[u] * (tr - mid);
        lazy[u * 2 + 1] += lazy[u];
        lazy[u] = 0;
    }

    void build (int u, int tl, int tr) {
        if (tl > tr) return;
        if (tl == tr) {
            if (tl < v.size()) stree[u] = v[tl];
            return;
        }
        int mid = (tl + tr) / 2;
        build(u * 2, tl, mid);
        build(u * 2 + 1, mid + 1, tr);
        stree[u] = merge(stree[u * 2], stree[u * 2 +
1]);
    }
    int query(int l, int r) {
        return query(l, 0, sz - 1, l, r);
    }
    void update(int l, int r, int val) {
        update(l, 0, sz - 1, l, r, val);
    }
    void update (int id, int val) {
        update(l, 0, sz - 1, id, val);
    }
    // Definition : return minimum of the range that
    // overlaps with l, r
    int query(int u, int tl, int tr, int l, int r) {
        if (tl > tr) return inf;
        if (l > r) return inf;
        if (tr < l || tl > r) return inf;
        if (tl >= l && tr <= r) return stree[u];
        int mid = (tl + tr) / 2;
        push(u);
        int lc = query(u * 2, tl, mid, l, r);
        int rc = query(u * 2 + 1, mid + 1, tr, l, r);
        return merge(lc, rc);
    }
    void update (int u, int tl, int tr, int id, int
val) {
        if (id < tl || id > tr) return;
        if (tl == tr) {
            stree[u] = val;
            return;
        }
        if (tl > tr) return;
        int mid = (tl + tr) / 2;
        push(u);
        update(u * 2, tl, mid, id, val);
        update(u * 2 + 1, mid + 1, tr, id, val);
        stree[u] = merge(stree[u * 2], stree[u * 2 +
1]);
    }
    void update (int u, int tl, int tr, int l, int r,
int val) {
        if (tl > tr) return;
        if (l > r) return;
```

```

if (tr < l || tl > r) return;
if (l <= tl && tr <= r) {
    stree[u] += val; //if range add -> stree[u]
    += (val * (tr - tl + 1));
    lazy[u] += val;
    return;
}
int mid = (tl + tr) / 2;
push(u);
update(u * 2, tl, mid, l, r, val);
update(u * 2 + 1, mid + 1, tr, l, r, val);
stree[u] = merge(stree[u * 2], stree[u * 2 +
    1]);
}
void printTree() {
    printTree(1, 0, sz - 1);
}
void printTree(int u, int tl, int tr) {
    cout << "Node " << u << " -> [" << tl << ", "
        << tr << "] = " << stree[u] << "\n";
    if (tl == tr) return; // leaf node
    int mid = (tl + tr) / 2;
    printTree(u * 2, tl, mid); // left child
    printTree(u * 2 + 1, mid + 1, tr); // right
        child
}
// using alive array
int Kth_one(int k) {
    return Kth_one(1, 0, sz - 1, k);
}
int Kth_one(int u, int tl, int tr, int k) {
    if(k > stree[u] || k <= 0) return -1; // not
        enough 1's
    if (tl == tr) {
        return tl;
    }
    // push(u);
    int mid = (tl + tr) / 2;
    if (stree[u * 2] >= k) {
        return Kth_one(u * 2, tl, mid, k);
    } else {
        return Kth_one(u * 2 + 1, mid + 1, tr, k -
            stree[u * 2]);
    }
}
signed main () {
    return 0;
}

```

1.10. Sparse Table

```

#include <bits/stdc++.h>
using namespace std;
const int N = 1e5 + 9;

```

```

int t[N][18], a[N];
void build(int n) {
    for (int i = 1; i <= n; ++i)
        t[i][0] = a[i];
    for (int k = 1; k < 18; ++k) {
        for (int i = 1; i + (1 << k) - 1 <= n; ++i) {
            t[i][k] = min(t[i][k - 1], t[i + (1 << (k -
                1))][k - 1]);
        }
    }
}
int query(int l, int r) {
    int k = 31 - __builtin_clz(r - l + 1);
    return min(t[l][k], t[r - (1 << k) + 1][k]);
}
int32_t main() {
    int n;
    cin >> n;
    for (int i = 1; i <= n; i++)
        cin >> a[i];
    build(n);
    int q;
    cin >> q;
    while (q--) {
        int l, r;
        cin >> l >> r;
        ++l;
        ++r;
        cout << query(l, r) << '\n';
    }
    return 0;
}

// Longest Common Subsequence print using next state
#include <bits/stdc++.h>
#define ll long long
using namespace std;
int n, m;
vector<int> a, b;
vector<vector<int>> dp;
vector<vector<pair<int, int>>> next_state;

int lcs(int i, int j) {
    if(dp[i][j] != -1) return dp[i][j];
    if(i == n || j == m) return 0;
    if(a[i] == b[j]) {
        next_state[i][j] = {i + 1, j + 1};
        dp[i][j] = 1 + lcs(i + 1, j + 1);
    }
}

```

```

else{
    int x = lcs(i+1, j);
    int y = lcs(i, j+1);
    if(x>y){
        next_state[i][j] = {i + 1, j};
    }
    else{
        next_state[i][j] = {i, j + 1};
    }
    dp[i][j] = max(dp[i][j], x);
    dp[i][j] = max(dp[i][j], y);
}
return dp[i][j];
}

void I_Am_Here() {
    cin >> n >> m;
    a = vector<int>(n);
    b = vector<int>(m);
    dp = vector<vector<int>>(n + 1, vector<int>(m +
        1, -1));
    next_state = vector<vector<pair<int, int>>>(n +
        1, vector<pair<int, int>>(m + 1, {0, 0}));
    for (int i = 0; i < n; i++) {
        cin >> a[i];
    }
    for (int i = 0; i < m; i++) {
        cin >> b[i];
    }
    cout << lcs(0, 0) << endl;

    int i=0 , j=0;
    while(i<n && j<m){
        int nxPosI=next_state[i][j].ff;
        int nxPosJ=next_state[i][j].ss;
        if(a[i]==b[j] && nxPosI==i+1 && nxPosJ==j+1){
            cout << a[i] << " ";
        }
        i = nxPosI;
        j = nxPosJ;
    }
    cout << endl;
}

int32_t main() {
    I_Am_Here();
}

```

2. Dynamic Programming

2.1. Longest Common Subsequence print using next state

```

#include <bits/stdc++.h>
using namespace std;

int lengthOfLIS(vector<int>& arr) {
    vector<int> lis;
    for(int x : arr) {

```

2.2. Longest Increasing Subsequence Print

```

auto it = lower_bound(lis.begin(), lis.end(),
                     x);
if(it == lis.end()) lis.push_back(x);
else *it = x;
}
return lis.size();
}

int main() {
vector<int> arr = {10, 22, 9, 33, 21, 50, 41,
                   60};
cout << lengthOfLIS(arr) << "\n";
return 0;
}

```

3. Graphs

3.1. BFS

```

const int N = 1e5 + 10; // Maximum num of nodes
int dist[N], par[N];
vector<int> adj[N];
queue<int> q;
void bfs(int s) {
    memset(dist, 63, sizeof(dist));
    dist[s] = 0;
    q.push(s);
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        for (auto v : adj[u])
            if (dist[v] > dist[u] + 1) {
                par[v] = u;
                dist[v] = dist[u] + 1;
                q.push(v);
            }
    }
}

```

3.2. Bellman Ford

```

const int N = 3e5 + 9;
struct st {
    int a, b, cost;
} e[N];
const int INF = 2e9;
int32_t main() {
    int n, m;
    cin >> n >> m;
    for (int i = 0; i < m; i++)
        cin >> e[i].a >> e[i].b >> e[i].cost;
    int s;
    cin >> s; // is there any negative cycle which is
               // reachable from s ?
}

```

```

vector<int> d(n, INF); // for finding any cycle(
                       // not necessarily from s) set d[i] = 0 for all
                       // i
d[s] = 0;
vector<int> p(n, -1);
int x;

for (int i = 0; i < n; ++i) {
    x = -1;
    for (int j = 0; j < m; ++j) {
        if (d[e[j].a] < INF) {
            if (d[e[j].b] > d[e[j].a] + e[j].cost) {
                d[e[j].b] = max(-INF, d[e[j].a] + e[j].
                                  cost); // for overflow
                p[e[j].b] = e[j].a;
                x = e[j].b;
            }
        }
    }
    if (x == -1)
        cout << "No negative cycle from " << s;
    else {
        int y = x; // x can be on any cycle or
                   // reachable from some cycle
        for (int i = 0; i < n; ++i)
            y = p[y];
        vector<int> path;
        for (int cur = y;; cur = p[cur]) {
            path.push_back(cur);
            if (cur == y && path.size() > 1)
                break;
        }
        reverse(path.begin(), path.end());
        cout << "Negative cycle: ";
        for (int i = 0; i < path.size(); ++i)
            cout << path[i] << ' ';
    }
    return 0;
}

```

3.3. Bipartite Graphs

```

const int N = 1e5 + 9;
vector<int> g[N];
bool vis[N];
int col[N];
bool ok;
void dfs(int u) {
    vis[u] = true;
    for (auto v : g[u]) {
        if (!vis[v]) {
            col[v] = col[u] ^ 1;
            dfs(v);
        } else {
            if (col[u] == col[v])
                ok = false;
        }
    }
}

```

```

int32_t main() {
    int n, m;
    cin >> n >> m;
    while (m--) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    ok = true;
    for (int i = 1; i <= n; i++)
        if (!vis[i])
            dfs(i);
    if (ok) cout << "YES\n";
    else cout << "NO\n";
}

```

3.4. Cycle or not Print Cycle

```

vector < vector < int > > adj;
vector < bool > vis;
vector < int > parent;
bool isCycle = 0;
int stop;
int start;
void dfs(int s,int p){
    vis[s]=1;
    parent[s]=p;
    for(auto x : adj[s]){
        if(vis[x]){
            if(parent[s]==x){
                continue; //agar tar thake aysi
            }
            else{
                isCycle = 1;
                start = x;
                stop = s;
                return;
            }
        }
        else{
            dfs(x,s);
        }
    }
}
void I_Am_Here() {
    int n;
    cin>> n;
    int m;
    cin>>m;
    adj = vector < vector < int > >(n + 1, vector <
                                      int > ());
    vis = vector < bool > (n + 1, false);
    parent = vector < int > (n + 1, 0);

    for(int i=0 ; i<m ; i++){
        int u,v;
        cin>>u>>v;
        adj[u].push_back(v);
        adj[v].push_back(u);
    }
}

```

```

adj[u].push_back(v);
adj[v].push_back(u);
}

for(int i=1 ; i<=n ; i++) {
    if(!vis[i]){
        dfs(i,0);
        if(isCycle) break;
    }
}
if(!isCycle){
    cout<<"IMPOSSIBLE"<<endl;
    return;
}
int t = start;
vector<int>ans;
ans.push_back(start);
while(start != stop){
    start = parent[start];
    ans.push_back(start);
}
ans.push_back(t);
cout<<ans.size()<<endl;
for(auto x : ans) cout<<x<< " ";cout<<endl;
}

```

3.5. Cycle or not in Directed Graph

```

vector<vector<int>>adj;
vector<int>vis;
vector<int>parent;
int n,m;
bool isCycle = 0;
int start,End;
void dfs(int s){
    vis[s]=1;
    for(auto child : adj[s]){
        if(!vis[child]){
            parent[child]=s;
            dfs(child);
            if(isCycle) return;
        }
        else if (vis[child]==1){
            isCycle=1;
            start = child;
            End = s;
            return;
        }
    }
    vis[s]=2;
}
void I_Am_Here() {
    cin >> n >> m;
    adj = vector<vector<int>>(n+1, vector<int>());
    vis = vector<int>(n+1,0);
    parent = vector<int>(n+1,0);

    for(int i=0,x,y ; i<m ; i++) {

```

```

        cin >> x >> y;
        adj[x].push_back(y);
    }

    for(int i=1 ; i<=n ; i++) {
        if(!vis[i]){
            dfs(i);
        }
        if(isCycle){
            break;
        }
    }
    if(!isCycle){
        cout<<"IMPOSSIBLE"<<endl;
        return ;
    }
    vector<int>ans;
    ans.push_back(start);
    ans.push_back(End);
    while(start != End){
        End = parent[End];
        ans.push_back(End);
    }
    reverse(full(ans));
    cout<<ans.size()<<endl;
    for(auto i:ans){
        cout<<i<< ' ';
    }
    cout<<endl;
}

```

3.6. DFS

```

const int N = 1e5 + 9;
vector<int> g[N];
bool vis[N];
void dfs(int u) {
    vis[u] = true;
    for (auto v : g[u])
        if (!vis[v])
            dfs(v);
}
int32_t main() {
    int n, m;
    cin >> n >> m;
    while (m--) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    dfs(u);
}

```

3.7. Dijkstra

```

const int N = 3e5 + 9, mod = 998244353;
```

```

int n, m;
vector<pair<int, int>> g[N], r[N];
vector<long long> dijkstra(int s, int t, vector<int> &cnt) {
    const long long inf = 1e18;
    priority_queue<pair<long long, int>, vector<pair<long long, int>>, greater<pair<long long, int>> q;
    vector<long long> d(n + 1, inf);
    vector<bool> vis(n + 1, 0);
    q.push({0, s});
    d[s] = 0;
    cnt.resize(n + 1, 0); // number of shortest paths
    cnt[s] = 1;

    while (!q.empty()) {
        auto x = q.top();
        q.pop();
        int u = x.second;
        if (vis[u]) continue;
        vis[u] = 1;

        for (auto y : g[u]) {
            int v = y.first;
            long long w = y.second;
            if (d[u] + w < d[v]) {
                d[v] = d[u] + w;
                q.push({d[v], v});
                cnt[v] = cnt[u];
            } else if (d[u] + w == d[v]) {
                cnt[v] = (cnt[v] + cnt[u]) % mod;
            }
        }
    }
    return d;
}

int u[N], v[N], w[N];
int32_t main() {
    int s, t;
    cin >> n >> m >> s >> t;
    for (int i = 1; i <= m; i++) {
        cin >> u[i] >> v[i] >> w[i];
        g[u[i]].push_back({v[i], w[i]});
        r[v[i]].push_back({u[i], w[i]});
    }
    vector<int> cnt1, cnt2;
    auto d1 = dijkstra(s, t, cnt1);
    auto d2 = dijkstra(t, s, cnt2);
    long long ans = d1[t];

    for (int i = 1; i <= m; i++) {
        int x = u[i], y = v[i];
        long long nw = d1[x] + w[i] + d2[y];
        if (nw == ans && 1LL * cnt1[x] * cnt2[y] % mod == cnt1[t])
            cout << "YES\n";
        else if (nw - ans + 1 < w[i]) cout << "CAN " << nw - ans + 1 << '\n';
    }
    else cout << "NO\n";
}

```

```

    }
    return 0;
}

```

3.8. Floyd Warshall

```

const int N = 105;
int d[N][N];
int main() {
    int n = 10;
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= n; j++) {
            if (i != j)
                d[i][j] = 1e9;
        }
    }

    for (int k = 1; k <= n; ++k) {
        for (int i = 1; i <= n; ++i) {
            for (int j = 1; j <= n; ++j) {
                d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
            }
        }
    }

    return 0;
}

```

3.9. Krushkals MST

```

const int N = 3e5 + 9, mod = 1e9;
struct dsu {
    vector<int> par, rnk, size;
    int c;

    dsu(int n) : par(n + 1), rnk(n + 1, 0), size(n + 1, 1), c(n) {
        for (int i = 1; i <= n; ++i)
            par[i] = i;
    }

    int find(int i) {
        return (par[i] == i ? i : (par[i] = find(par[i])));
    }

    bool same(int i, int j) {
        return find(i) == find(j);
    }

    int get_size(int i) {
        return size[find(i)];
    }

    int count() {
        return c; // connected components
    }
}

```

```

int merge(int i, int j) {
    if ((i = find(i)) == (j = find(j)))
        return -1;
    else
        --c;
    if (rnk[i] > rnk[j])
        swap(i, j);
    par[i] = j;
    size[j] += size[i];
    if (rnk[i] == rnk[j])
        rnk[j]++;
    return j;
}

int32_t main() {
    int n, m;
    cin >> n >> m;
    vector<array<int, 3>> ed;
    for (int i = 1; i <= m; i++) {
        int u, v, w;
        cin >> u >> v >> w;
        ed.push_back({w, u, v});
    }

    sort(ed.begin(), ed.end());
    long long ans = 0;
    dsu d(n);

    for (auto e : ed) {
        int u = e[1], v = e[2], w = e[0];
        if (d.same(u, v))
            continue;
        ans += w;
        d.merge(u, v);
    }
    cout << ans << '\n';
    return 0;
}

```

3.10. LCA

```

const int N = 3e5 + 9, LG = 18;
vector<int> g[N];
int par[N][LG + 1], dep[N], sz[N];
void dfs(int u, int p = 0) {
    par[u][0] = p;
    dep[u] = dep[p] + 1;
    sz[u] = 1;

    for (int i = 1; i <= LG; i++)
        par[u][i] = par[par[u][i - 1]][i - 1];

    for (auto v : g[u])
        if (v != p) {
            dfs(v, u);
            sz[u] += sz[v];
        }
}

```

```

}

int lca(int u, int v) {
    if (dep[u] < dep[v])
        swap(u, v);
    for (int k = LG; k >= 0; k--)
        if (dep[par[u][k]] >= dep[v])
            u = par[u][k];
    if (u == v)
        return u;
    for (int k = LG; k >= 0; k--)
        if (par[u][k] != par[v][k])
            u = par[u][k], v = par[v][k];
    return par[u][0];
}

int kth(int u, int k) {
    assert(k >= 0);
    for (int i = 0; i <= LG; i++)
        if (k & (1 << i))
            u = par[u][i];
    return u;
}

int dist(int u, int v) {
    int l = lca(u, v);
    return dep[u] + dep[v] - (dep[l] << 1);
}

// kth node from u to v, 0th node is u
int go(int u, int v, int k) {
    int l = lca(u, v);
    int d = dep[u] + dep[v] - (dep[l] << 1);
    assert(k <= d);
    if (dep[l] + k <= dep[u])
        return kth(u, k);
    k -= dep[u] - dep[l];
    return kth(v, dep[v] - dep[l] - k);
}

int32_t main() {
    int n;
    cin >> n;
    for (int i = 1; i < n; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    dfs(1);
    int q;
    cin >> q;
    while (q--) {
        int u, v;
        cin >> u >> v;
        cout << dist(u, v) << '\n';
    }
    return 0;
}

```

3.11. Prims MST

```

const int N = 2020;
int g[N][N], w[N], to[N], selected[N];

long long Prims(int n, vector<pair<int, int>> &edges)
{
    long long ans = 0;
    for (int i = 1; i <= n; i++)
        w[i] = 1e9, selected[i] = 0, to[i] = -1;
    w[1] = 0;

    for (int i = 1; i <= n; i++) {
        int u = -1;
        for (int j = 1; j <= n; j++)
            if (!selected[j] && (u == -1 || w[j] < w[u]))
                u = j;

        if (w[u] == 1e9)
            return -1; // NO MST

        selected[u] = 1;
        ans += w[u];

        if (to[u] != -1)
            edges.emplace_back(u, to[u]); // order of
                                         // the edges may be changed

        for (int v = 1; v <= n; v++)
            if (g[u][v] < w[v])
                w[v] = g[u][v], to[v] = u;
    }

    return ans;
}

string s[N];
int main()
{
    int n, m;
    cin >> n >> m;

    for (int i = 1; i <= n; i++)
        for (int j = 1; j <= n; j++)
            g[i][j] = 1e9;

    for (int i = 1; i <= n; i++)
        cin >> s[i];

    for (int i = 1; i <= n; i++) {
        for (int j = i + 1; j <= n; j++) {
            int w = 0;
            for (int k = 0; k < m; k++)
                w = max(w, (int)abs(s[i][k] - s[j][k]));
            g[i][j] = min(g[i][j], w);
            g[j][i] = min(g[j][i], w);
        }
    }

    vector<pair<int, int>> ed;
}

```

```

long long ans = Prims(n, ed);
int res = 0;
for (auto e : ed) res = max(res, g[e.first][e.second]);
cout << res << '\n';
return 0;
}

// https://www.codechef.com/ICL2016/problems/ICL16A/

```

3.12. Strongly Connected Components

```

const int N = 3e5 + 9;
// given a directed graph return the minimum number
// of edges to be added so that the whole graph
// become an SCC

bool vis[N];
vector<int> g[N], r[N], G[N], vec; // G is the
                                         // condensed graph
void dfs1(int u) {
    vis[u] = 1;
    for (auto v : g[u])
        if (!vis[v])
            dfs1(v);
    vec.push_back(u);
}

vector<int> comp;
void dfs2(int u) {
    comp.push_back(u);
    vis[u] = 1;
    for (auto v : r[u])
        if (!vis[v])
            dfs2(v);
}

int idx[N], in[N], out[N];

int main() {
    int n, m;
    cin >> n >> m;
    for (int i = 1; i <= m; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        r[v].push_back(u);
    }

    for (int i = 1; i <= n; i++)
        if (!vis[i])
            dfs1(i);

    reverse(vec.begin(), vec.end());
    memset(vis, 0, sizeof vis);
    int scc = 0;

```

```

for (auto u : vec) {
    if (!vis[u]) {
        comp.clear();
        dfs2(u);
        scc++;
        for (auto x : comp)
            idx[x] = scc;
    }
}

for (int u = 1; u <= n; u++) {
    for (auto v : g[u]) {
        if (idx[u] != idx[v]) {
            in[idx[v]]++;
            out[idx[u]]++;
            G[idx[u]].push_back(idx[v]);
        }
    }
}

int needed_in = 0, needed_out = 0;
for (int i = 1; i <= scc; i++) {
    if (!in[i]) needed_in++;
    if (!out[i]) needed_out++;
}
int ans = max(needed_in, needed_out);
if (scc == 1) ans = 0;
cout << ans << '\n';
return 0;
}

```

3.13. Topological Sorting

```

vector<int> adj[100005];
vector<int> ans;
vector<int> vis;
void dfs(int u) {
    vis[u] = 1;
    for (int v : adj[u])
        if (!vis[v]) dfs(v);
    ans.push_back(u);
}

vector<int> topSort(int n) {
    vis.assign(n + 1, 0);
    ans.clear();
    for (int i = 1; i <= n; i++) {
        if (!vis[i]) dfs(i);
    }
    reverse(ans.begin(), ans.end());
    return ans;
}
// https://cses.fi/problemset/task/1679

```

3.14. Zero One BFS

```
// 0-1 BFS - O(V+E)
```

```

const int N = 1e5 + 5;
int dist[N];
vector<pii> adj[N];
deque<pii> dq;

void zero_one_bfs(int x) {
    cl(dist, 63);
    dist[x] = 0;
    dq.push_back({x, 0});

    while (!dq.empty()) {
        int u = dq.front().st;
        int ud = dq.front().nd;
        dq.pop_front();

        if (dist[u] < ud)
            continue;

        for (auto x : adj[u]) {
            int v = x.st;
            int w = x.nd;
            if (dist[u] + w < dist[v]) {
                dist[v] = dist[u] + w;
                if (w)
                    dq.push_back({v, dist[v]});
                else
                    dq.push_front({v, dist[v]});
            }
        }
    }
}

```

4. Number Theory

4.1. 2D Prefix Sum

```

for (int i = 1; i <= n; i++) {
    for (int j = 1; j <= n; j++) {
        prefixSum[i][j] = prefixSum[i-1][j] + prefixSum[
            i][j-1] - prefixSum[i-1][j-1] + grid[i-1][j-
            1];
    }
}
cin>>x1>>y1>>x2>>y2;
cout<<prefixSum[x2][y2] + prefixSum[x1-1][y1-1] -
    prefixSum[x1-1][y2] - prefixSum[x2][y1-1]<<endl;

```

4.2. Number of Integers Coprime with N(1 to n)

```

void phi_1_to_n(int n) {
    phi = vector<int>(n+1);
    for (int i = 0; i <= n; i++)
        phi[i] = i;

    for (int i = 2; i <= n; i++) {
        if (phi[i] == i) {
            for (int j = i; j <= n; j += i)
                phi[j] -= phi[j] / i;
        }
    }
}

```

```
}
```

4.3. Binary Exponentiation

```

long long binpow(long long a, long long b, long long
    m) {
    a %= m;
    long long res = 1;
    while (b > 0) {
        if (b & 1) res = res * a % m;
        a = a * a % m;
        b >>= 1;
    }
    return res;
}

```

4.4. Binary Multiplication with Mod

```

long long binmul(long long a, long long b, long long
    m) {
    long long res = 0;
    a %= m;
    while (b > 0) {
        if (b & 1) res = (res + a) % m;
        a = (a + a) % m;
        b >>= 1;
    }
    return res;
}

```

4.5. Modular Arithmetic

```

// Greatest Common Divisor & Lowest Common Multiple
ll gcd(ll a, ll b) {
    return b ? gcd(b, a % b) : a;
}

ll lcm(ll a, ll b) {
    return a / gcd(a, b) * b;
}

// Multiply caring overflow
ll mulmod(ll a, ll b, ll m = MOD) {
    ll r = 0;
    for (a %= m; b; b >>= 1, a = (a * 2) % m)
        if (b & 1)
            r = (r + a) % m;
    return r;
}

// Another option for mulmod is using long double
ull mulmod(ull a, ull b, ull m = MOD) {
    ull q = (ld)a * (ld)b / (ld)m;
    ull r = a * b - q * m;
    return (r + m) % m;
}

```

```

// Fast exponential
ll fexp(ll a, ll b, ll m = MOD) {
    ll r = 1;
    for (a %= m; b; b >>= 1, a = (a * a) % m)
        if (b & 1)
            r = (r * a) % m;
    return r;
}

```

4.6. Binomial Coefficients nCr

```

const int N = 2e5 + 9, mod = 998244353;
long long fact[N], finv[N], inv[N];
void precal_factorial(int n) {
    inv[0] = inv[1] = finv[0] = finv[1] = fact[0] =
        fact[1] = 1;
    for (int i = 2; i <= n; i++) {
        inv[i] = inv[mod % i] * (mod - mod / i) % mod;
        finv[i] = (finv[i - 1] * inv[i]) % mod;
        fact[i] = (fact[i - 1] * i) % mod;
    }
}

long long ncr(long long n, long long r) {
    if (r > n) return 0;
    long long a = (finv[r] * finv[n - r]) % mod;
    return (a * fact[n]) % mod;
}

```

4.7. How Many Digits in n!

```

#include <bits/stdc++.h>
using namespace std;
int findDigits(int n){ if (n < 0) return 0;if (n <=
    1) return 1; double digits = 0;
    for (int i = 2; i <= n; i++) digits += log10(i);
    return floor(digits) + 1;
}

```

4.8. Derangement

```

const int N = 1e6 + 9, mod = 1e9 + 7;
int d[N];
int32_t main() {
    d[0] = 1;
    d[1] = 0;
    for (int i = 1; i < N; i++)
        d[i] = 1LL * (i - 1) * (d[i - 1] + d[i - 2]) %
            mod;

    int n;
    cin >> n;
}

```

```

cout << d[n] << '\n';
return 0; }

/* There are n children at a Christmas party, and
each of them has brought a gift.
The idea is that everybody will get a gift brought
by someone else.
In how many ways can the gifts be distributed? */

```

4.9. Euler Phi

```

// Euler phi (totient)
int ind = 0, pf = primes[0], ans = n;
while (1ll * pf * pf <= n) {
    if (n % pf == 0)
        ans -= ans / pf;
    while (n % pf == 0)
        n /= pf;
    pf = primes[++ind];
if (n != 1)ans -= ans / n;

// Euler Totient Function 1 to n in O(nloglog(n))
const int N = 1e5 + 9;
int phi[N];

void totient() {
    for (int i = 1; i < N; i++)
        phi[i] = i;
    for (int i = 2; i < N; i++) {
        if (phi[i] == i)
            for (int j = i; j < N; j += i)
                phi[j] -= phi[j] / i;
    }
}

```

4.10. Fibonacci Number Faster

```

int fib(long long n, int mod) {
    assert(n >= 0);
    if (n <= 1) return n;
    int a = 0, b = 1;
    long long i = 1ll << (63 - __builtin_clzll(n) - 1);

    for (; i; i >= 1) {
        int na = (a * (long long)a + b * (long long)b) %
mod;
        int nb = (2ll * a + b) * b % mod;
        a = na;
        b = nb;

        if (n & i) {
            int c = a + b;
            if (c >= mod) c -= mod;
            a = b;
            b = c;
        }
    }
}

```

```

    }
    return b;
}


```

4.11. GCD LCM

```

int gcd(int a, int b){if (a == 0) return b; return
gcd(b % a, a);}
long long lcm(long long a, long long b){ return (a /
__gcd(a, b)) * b;}

```

4.12. Josephus

```

/* Josephus Problem - It returns the position to be,
in order to not die. O(n) */
/* With k=2, for instance, the game begins with 2
being killed and then n+2, n+4, */
ll josephus(ll n, ll k) {
    if (n == 1)
        return 1;
    else
        return (josephus(n - 1, k) + k - 1) % n + 1;
}

/* Another Way to compute the last position to be
killed - O(d * log n) */
ll josephus(ll n, ll d) {
    ll K = 1;
    while (K <= (d - 1) * n)
        K = (d * K + d - 2) / (d - 1);
    return d * n + 1 - K;
}

```

4.13. Large Number GCD

```

ll gcd(ll a, ll b) {
    if (!a)
        return b;
    return gcd(b % a, a);
}
ll reduceB(ll a, char b[]) {
    ll mod = 0;
    for (int i = 0; i < strlen(b); i++)
        mod = (mod * 10 + b[i] - '0') % a;
    return mod; // return modulo
}
ll gcdLarge(ll a, char b[]) {
    ll num = reduceB(a, b);
    return gcd(a, num);
}
int main() {
    ll a = 1221;
}

```

```

char b[] = "
123456789101112131415161718192012223242526272829
";
if (a == 0)
    cout << b << endl;
else
    cout << gcdLarge(a, b) << endl;
return 0;
}

```

4.14. Legendre's Formula

```

// n and a prime number p, find the largest x such
that px divides n !(factorial) in O(logn).
#include <bits/stdc++.h>
using namespace std;
int legendre(long long n, long long p){ int ans = 0;
    while (n){ans += n / p; n /= p;}
    return ans;}

```

4.15. Modular for Subtraction & Multiplication

```

int vagsesh = (a % m - b % m + m) % m; // For
Subtraction
long long res = 1; // For Multiplication
for (int i = 1; i <= n; i++)res = (res * a) % m;
cout << res << endl;

```

4.16. Optimized Sieve

```

// (Fast Sieve, Using bit set, Works till 10^8 in
less than 1s, Memory Complexity: O (n/64))

const int N = 1e8 + 9;
bitset<N> f;
int32_t main() {
    int n = N - 9;
    vector<int> primes;
    f[1] = true;

    for (int i = 2; i * i <= n; i++)
        if (!f[i])
            for (int j = i * i; j <= n; j += i)
                f[j] = true;

    for (int i = 2; i <= n; i++)
        if (!f[i])
            primes.push_back(i);

    cout << primes.size() << '\n';
    return 0;
}

```

4.17. Lower Prime Factorization

```
// Linear Sieve for prime factorization
//time complexity: O(n)
//space complexity: O(n)
const int N = 10000000; // up to 1e7
vector<int> lp; // lowest prime factor
vector<int> primes; // list of primes
void linear_sieve(int n) {
    lp = vector<int>(n + 1, 0);
    for (int i = 2; i <= n; i++) {
        if (lp[i] == 0) { // i is prime
            lp[i] = i;
            primes.push_back(i);
        }
        for (int p : primes) {
            if (p > lp[i] || 1LL * i * p > n) break;
            lp[i * p] = p;
        }
    }
}
```

4.18. Prime or Not

```
bool prime(int n) {
    if(n < 2) return false;
    if(n == 2) return true;
    for(int i = 2; i * i <= n; i++) {
        if(n % i == 0)
            return false;
    }
    return true;
}
```

4.19. Number of Divisor A Number Have

```
long long numDivisors[1000005] = {};
void calculateDivisors(long long numDivisors[])
{
    //calculate number of divisor
    for (int i = 1; i < 1000005; i++) {
        for (int j = i; j < 1000005; j += i) {
            numDivisors[j]++;
        }
    }
}
```

4.20. Prime Factors

```
void calcFp(int n) {
    fp = vector < int >(n + 1, 0);
    for (int i = 2; i <= n; i++) {
        if (fp[i] == 0) { // i is prime
            for (int j = i; j <= n; j += i) {
                if (fp[j] == 0) fp[j] = i;
            }
        }
    }
}
```

```
if (fp[j] == 0) fp[j] = i;
}

int main () {
    calcFp((int)1e6);
    while (1) {
        int x;
        cin >> x;
        if (x == 0) break;
        vector < int > primes, pw;
        while (x > 1) {
            int p = fp[x];
            int cnt = 0;
            while (x % p == 0) {
                x /= p;
                cnt++;
            }
            primes.push_back(p);
            pw.push_back(cnt);
        }
        for (int i = 0; i < primes.size(); i++) {
            cout << primes[i] << "^" << pw[i] << endl;
        }
    }
    return 0;
}
```

4.21. Formulas Using Prime Factorization Math

```
// if a number prime factor =  $p_1^{x_1} \cdot p_2^{x_2} \cdot p_3^{x_3} \cdots p_n^{x_n}$ 
// then number of divisors =  $(x_1+1) \cdot (x_2+1) \cdot (x_3+1) \cdots (x_n+1)$ 
// total divisor sum =  $(p_1^{(x_1+1)-1/p_1-1}) \cdot (p_2^{(x_2+1)-1/p_2-1}) \cdot (p_3^{(x_3+1)-1/p_3-1}) \cdots (p_n^{(x_n+1)-1/p_n-1})$ 
// -----sum = sum-N
// product of divisor =  $n^{(\text{number of divisors}/2)}$ 
// b/a = b*a^(-1);
// a^(-1) == a^(MOD-2)%MOD;
// product of divisor
int div_prod = expo(div_prod, k[i] + 1) *
    expo(expo(p[i], (k[i] * (k[i] + 1) / 2)), div_cnt2) % MOD;
int div_cnt2 = div_cnt2 * (k[i] + 1) % (MOD - 1);
```

4.22. Sum & Product of Divisor

```
for (int i = 0; i < n; i++) cin >> p[i] >> k[i];
ll div_cnt = 1, div_sum = 1, div_prod = 1, div_cnt2
    = 1;
for (int i = 0; i < n; i++) {
    div_cnt = div_cnt * (k[i] + 1) % MOD;
```

```
div_sum = div_sum * (expo(p[i], k[i] + 1) - 1) %
    MOD * expo(p[i] - 1, MOD - 2) % MOD;
div_prod = expo(div_prod, k[i] + 1) * expo(
    expo(p[i], (k[i] * (k[i] + 1) / 2)), div_cnt2) %
    MOD;
div_cnt2 = div_cnt2 * (k[i] + 1) % (MOD - 1);
}
cout << div_cnt << ' ' << div_sum << ' ' <<
    div_prod;
```

4.23. Sieve

```
// primes which are less than n in O(nloglog(n))
const int N = 1e7 + 9;
bool f[N];

int32_t main() {
    int n = N - 9;
    vector<int> primes;
    f[1] = true;
    for (int i = 2; i <= n; i++) {
        if (!f[i]) {
            primes.push_back(i);
            for (int j = i + i; j <= n; j += i)
                f[j] = true;
        }
    }
    cout << primes.size() << '\n';
    return 0;
}
```

4.24. n! Trailing Zero

```
int findTrailingZeros(int n) {
    if (n < 0) // Negative Number Edge Case
        return -1;
    int count = 0;
    for (int i=5;n/i >= 1; i *= 5) count += n / i;
    return count;
}
```

4.25. Number of Divisors 1 to n

```
int d[104];
int32_t main(){
int n = 100;
    for (int i = 1; i <= n; i++) for (int j = i; j <=
    n; j += i)d[j]++; // d[j] += i // for sum of
    divisors
    for (int i = 1; i <= n; i++)cout << d[i] << ;
    return 0;
}
```

4.26. NcR (when n and r to big)(Lucas Theorem)

$$\binom{n}{r} \bmod p = \left(\frac{n!}{r!(n-r)!} \right) \bmod p$$

n and r too big step 1:

- write n in base p
- write r in base p
- (Let the digits be $n_1, n_2, n_3\dots$ and $k_1, k_2, k_3\dots$)

$$\binom{n}{k} \bmod p = \binom{n_0}{k_0} \cdot \binom{n_1}{k_1} \cdots \bmod p$$

If $k_i > n_i$, then the answer is 0.

```
-----code section-----
#include <bits/stdc++.h>
using namespace std;

int p; // prime modulus
vector<int> fact; // factorials modulo p

// Fast exponentiation: a^b % mod
int modPow(int a, int b, int mod) {
    int result = 1;
    a %= mod;
    while (b > 0) {
        if (b & 1) result = (int)((1LL * result * a) % mod);
        a = (int)((1LL * a * a) % mod);
        b >>= 1;
    }
    return result;
}

// Precompute factorials modulo p
void precomputeFactorials(int prime) {
    p = prime;
    fact = vector<int>(p+1, 1);
    for (int i = 1; i < p; ++i)
        fact[i] = ((1LL * fact[i - 1] * i) % p);
}

// Compute nCr % p using modular inverse on the fly
int nCr(int n, int r) {
    if (r < 0 || r > n) return 0;
    int numerator = fact[n];
    int denominator = ((1LL * fact[r] * fact[n - r]) % p);
    return ((1LL * numerator * modPow(denominator, p - 2, p)) % p);
}

// Lucas Theorem
```

```
int nCrModpLucas(int n, int r) {
    int res = 1;
    while (n > 0 || r > 0) {
        int ni = n % p;
        int ri = r % p;
        if (ri > ni) return 0; // if any digit ri > ni
        , result is 0
        res = (int)((1LL * res * nCr(ni, ri)) % p);
        n /= p;
        r /= p;
    }
    return res;
}
int main() {
    int n = 100, r = 25, prime = 5;
    precomputeFactorials(prime);
    cout << "Value of nCr % p is " << nCrModpLucas(n, r) << endl;
    return 0;
}
```

4.27. how many numbers (0..n) have k-th bit=1

$$\text{count}(k) = \left(\frac{n+1}{2^{k+1}}\right) \cdot 2^k + \max(0, ((n+1) \bmod 2^{k+1}) - 2^k)$$

5. Mixed

5.1. Run Code Technique

```
//Pi value
const long double PI = 3.14159265358979323846;
const double PI = acos(-1.0);

//C++ Shell for run
g++ file_name.cpp -o a
.\a

//Python shell for run
python -u "g:\foldername\filename.py"

//first io
ios::sync_with_stdio(false);
cin.tie(nullptr);

//input, output in file
#ifndef ONLINE_JUDGE
freopen("input.txt", "r", stdin);
freopen("output.txt", "w", stdout);
#endif
```

5.2. Converts Any Integer Number to Any Base

```
string convertToBase(int number, int base) {
    if (base < 2 || base > 36) return ""; // invalid base
    if (number == 0) return "0"; // edge case

    string digits = "0123456789
    ABCDEFGHIJKLMNOPQRSTUVWXYZ";
    string result = "";
    bool isNegative = false;

    if (number < 0) {
        isNegative = true;
        number = -number;
    }

    while (number > 0) {
        int remainder = number % base;
        result += digits[remainder];
        number /= base;
    }

    if (isNegative) result += '-';
    reverse(result.begin(), result.end());
    return result;
```

5.3. i128 Bit

```
#include <bits/stdc++.h>
#define i128 __int128_t
using namespace std;
int32_t main() {
    i128 x = 1;
    int y = (int)(x+1);
    cout << y << endl;
    return 0;
}
// i128 can be stored not printed
// range -2^127 to 2^127-1
```

5.4. Base Conversion

```
string base_convert(int n, int b) {
    string s = "";
    while (n > 0) {
        s = to_string(n % b) + s;
        n /= b;
    }
    return s;
}
int convert_to_decimal(string s, int base) {
    int n = 0, power = 1;
    for (int i = (int)s.size() - 1; i >= 0; i--) {
        n += power * (s[i] - '0');
```

```

    power *= base;
}
return n;

```

5.5. Bitset

```

// Declaration with size N
bitset<100005> bs; // all bits 0 initially

// Set all bits to 1
bs.set();

// Set i-th bit
bs.set(i);

// Reset i-th bit (make 0)
bs.reset(i);

// Reset all bits
bs.reset();

// Flip i-th bit
bs.flip(i);

// Flip entire bitset
bs.flip();

// Check if i-th bit is set
bool f = bs.test(i);

// Count number of 1s
int cnt = bs.count();

// Check if all bits are 1
bool ok1 = bs.all();

// Check if any bit is 1
bool ok2 = bs.any();

// Check if none bit is 1
bool ok3 = bs.none();

// Convert bitset to unsigned long long
unsigned long long x = bs.to_ullong();

// Convert bitset to string
string s = bs.to_string();

```

5.6. Bit Operation

```

const int inf = numeric_limits<int>::max() - 5;
#define int long long

struct ST {

```

```

int right_shift(int a, int b) { return (a >> b); }
int left_shift(int a, int b) { return (a << b); }
int bitwise_not(int a) { return (~a); }
int ON_BIT(int a) { return __builtin_popcount(a); }
int leading_zero(int a) { return __builtin_clz(a); }
int tailing_zero(int a) { return __builtin_ctz(a); }
int LSB(int a) { return (a & 1); }
bool kth_bit(int a, int k) { return (a & (1 << k)); }
int msb(int a) { return a ? 32 - __builtin_clz(a) : 0; }

string base_convert(int n, int b) {
    string s = "";
    while (n > 0) {
        s = to_string(n % b) + s;
        n /= b;
    }
    return s;
}

int convert_to_decimal(string s, int base) {
    int n = 0, power = 1;
    for (int i = (int)s.size() - 1; i >= 0; i--) {
        n += power * (s[i] - '0');
        power *= base;
    }
    return n;
} bit;

// other
// Count number of set bits (1s)
int pc = __builtin_popcount(x); // for int
int pcLL = __builtin_popcountll(x); // for long long

// Count trailing zeros
int tz = __builtin_ctz(x);
int tzLL = __builtin_ctzll(x);

// Count leading zeros
int lz = __builtin_clz(x);
int lzLL = __builtin_clzll(x);

// Check power of two
bool isP2 = (x > 0) && ((x & (x - 1)) == 0);

// Merge-sort with inversion count - O(nlog n)
int n, inv;
vector<int> v, ans;

void mergesort(int l, int r, vector<int> &v) {

```

```

    if (l == r) return;
    int mid = (l + r) / 2;
    mergesort(l, mid, v);
    mergesort(mid + 1, r, v);

    int i = l, j = mid + 1, k = l;
    while (i <= mid || j <= r) {
        if (i <= mid && (j > r || v[i] <= v[j]))
            ans[k++] = v[i++];
        else
            ans[k++] = v[j++], inv += j - k;
    }

    for (int i = l; i <= r; i++)
        v[i] = ans[i];
}
// in main: ans.resize(v.size());

```

5.8. STL

```

auto [a, b] = p; pair<int, int> p; pair<int, int> p3;

deque<int> dq; dq.push_front(1); dq.push_back(3);
    pop_front(); pop_back(); //all operation O(1)

stack<int> st; st.push(2); st.pop(); st.top(); st.size();

queue<int> q; q.push(5); q.pop(); q.front(); q.back();
    size(); empty();
//all operation O(1)

-----priority queue-----
// Max priority first (largest value removed first)
priority_queue<long long> pq_max;

// Min priority first (smallest value removed first)
priority_queue<long long, vector<long long>, greater<long long> pq_min;

// Clear
pq_max = priority_queue<long long>(); //same ot other operation
pq.push(x); pq.top(); pq.pop();
//push and pop log(n) and other o(1)

-----map-----
map<int, int> mp1; map<int, pair<int, int>> mp2; mp2[0].first; mp2[0].second;
int index = lower_bound(v.begin(), v.end(), val) - v.begin();
int index = upper_bound(v.begin(), v.end(), val) - v.begin();
auto it = s.lower_bound(6); *it; it--;
auto it = s.upper_bound(6); *it; it--;
//all log(n)

```

5.7. Merge Sort

```

// Merge-sort with inversion count - O(nlog n)
int n, inv;
vector<int> v, ans;

void mergesort(int l, int r, vector<int> &v) {

```

```
-----unordered_map-----
unordered_map<int,int> ump;
ump[key] = value; // O(1) average, O(n) worst
ump.insert({key,value}); // O(1) average, O(n) worst
ump.find(key); // O(1) average, O(n) worst
ump.count(key); // O(1) average, O(n) worst
ump.erase(key); // O(1) average, O(n) worst
ump.size();ump.empty();ump.clear();//>o(1)

-----multiset-----
multiset < int > m; insert(); erase(); begin(); end()
(); size(); count(); empty(); ms.count(x),ms.
lower_bound(x),ms.upper_bound(x); //all operation
TC = log(n)
//all log(n) but begin,size,empty->o(1)
```

5.9. Sqrt Decomposition

```
// Square Root Decomposition (Mos Algorithm) - O(n^(3/2))
const int N = 1e5 + 1, SQ = 500;
int n, m, v[N];

void add(int p) { /* add value to aggregated data
structure */ }
void rem(int p) { /* remove value from aggregated
data structure */ }

struct query {
    int i, l, r, ans;
} qs[N];

bool c1(query a, query b) {
    if (a.l / SQ != b.l / SQ) return a.l < b.l;
    return a.l / SQ & 1 ? a.r > b.r : a.r < b.r;
}
bool c2(query a, query b) { return a.i < b.i; }
/* inside main */
int l = 0, r = -1;
sort(qs, qs + m, c1);

for (int i = 0; i < m; ++i) {
    query &q = qs[i];
    while (r < q.r) add(v[++r]);
    while (r > q.r) rem(v[r--]);
    while (l < q.l) rem(v[l++]);
    while (l > q.l) add(v[--l]);
    q.ans = /* calculate answer */;
}

sort(qs, qs + m, c2); // sort to original order
```

5.10. Binary Search

\it give maximum valid answer

```
while (high-low>1){
    mid = (high+low)/2;
    //mid jokhon ans hote pare tokhon mid+1/mid-1
    kora jabe nah
    if (canPlaceCow(mid,a)){
        low = mid;
    } else{
        high = mid;
    }
}\>Finds minimum value
while (low <= high) {
    long long mid = (low + high) / 2;
    if (check(mid)) {
        ans = mid; // mid is valid
        high = mid - 1; // try smaller
    } else {
        low = mid + 1;
    }
}
```

5.11. lower_bound & upper_bound

.....lower bound.....
it = lower_bound(arr.begin(), arr.end(), x);
Iterator to the first element x

.....upper bound.....
it = upper_bound(arr.begin(), arr.end(), x);
Iterator to the first element > x

it give 0-based indexing

5.12. Ternary Search

```
// Ternary Search - O(log(n))
// Max version, for minimum version just change
signals
ll ternary_search(ll l, ll r) {
    while (r - l > 3) {
        ll m1 = (l + r) / 2;
        ll m2 = (l + r) / 2 + 1;
        ll f1 = f(m1), f2 = f(m2); // if(f1 > f2) l =
        m1;
        if (f1 < f2)
            l = m1;
        else
            r = m2;
    }
    ll ans = 0;
    for (int i = l; i <= r; i++) {
        ll tmp = f(i); // ans = min(ans, tmp);
        ans = max(ans, tmp);
    }
    return ans;
}
```

```
}
// Faster version - 300 iterations up to 1e-6
precision
double ternary_search(double l, double r, int No =
300) {
    // for(int i = 0; i < No; i++){
    while (r - l > EPS) {
        double m1 = l + (r - l) / 3;
        double m2 = r - (r - l) / 3;
        // if (f(m1) > f(m2))
        if (f(m1) < f(m2))
            l = m1;
        else
            r = m2;
    }
    return f(l);
}
```

5.13. Custom Sorting

```
//single value
bool descending(int a, int b) {
    return a > b;
}
int main() {
    vector<int> nums = {5, 3, 1, 4, 2};
    sort(nums.begin(), nums.end(), descending); // custom comparator
}

//three value
bool customSort(const pair<int, pair<int,int>> &a,
const pair<int, pair<int,int>> &b) {
    if (a.first != b.first)
        return a.first < b.first;// 1st ascending
    if (a.second.first != b.second.first)
        return a.second.first > b.second.first; // 2nd
descending
    return a.second.second < b.second.second; // 3rd
ascending
}

int main() {
    vector<pair<int, pair<int,int>>> v = {
        {1, {5, 10}},};
    sort(v.begin(), v.end(), customSort);
}
```

5.14. Permutation & Combination

$$\text{Permutation: } {}^n P_r = \frac{n!}{(n-r)!}$$

$$\text{Combination: } {}^n C_r = \frac{n!}{r!(n-r)!}$$

5.15. python

```
# reopen
import sys
sys.stdout = open(out, w )
sys.stdin = open(in , r )
//Dummy example
R = lambda: map(int, input().split())
n, k = R()
v, t = [], [0]*n
for p, c, i in sorted(zip(R(), R(), range(n))):
    t[i] = sum(v)+c
    v += [c]
    v = sorted(v)[::-1]
    if len(v) > k:
        v.pop()
print(''.join(map(str, t)))
```

6. Strings

6.1. All String Operation

```
string s = "abcdef";
//Substring
string sub = s.substr(0, 5+1);

//Find substring or character
int pos = s.find("cd"); // -1 if not found
int posChar = s.find('e');

// Compare strings
bool eq = (s1 == s2);
bool lt = (s1 < s2); // lexicographical comparison

// Convert int <-> string
string numStr = to_string(45);
int x = stoi("123");
long long y = stoll("123456789"); // string to long
long

// Convert float <-> string
double d = 3.14;
string ds = to_string(d);
double dd = stod("2.718");

// Uppercase / Lowercase
string t = "AbCd";
transform(t.begin(), t.end(), t.begin(), ::toupper);
transform(t.begin(), t.end(), t.begin(), ::tolower);

//Reverse string
reverse(rev.begin(), rev.end());

//Remove spaces
string str = "a b c d";
```

```
str.erase(remove(str.begin(), str.end(), ' '), str.end());

// Subsequence check
string small = "ace";
int i = 0;
for (char c : s) if (i < small.size() && small[i] == c) i++;
bool isSubseq = (i == small.size());
```

6.2. Single Hashing

```
//single string hashing
struct StringHash {
    const long long MOD = 1000000007; // large prime
    const long long P = 31; // base (for lowercase letters)
    vector<long long> prefix, power;

    StringHash(const string &s) {
        int n = s.size();
        prefix = vector<long long>(n+1, 0);
        power = vector<long long>(n+1, 1);

        for (int i = 1; i <= n; i++) {
            power[i] = (power[i-1] * P) % MOD;
            prefix[i] = (prefix[i-1] * P + (s[i-1] - 'a' + 1)) % MOD;
        }
    }

    // get hash of substring [l, r] (0-indexed inclusive)
    long long getHash(int l, int r) {
        long long res = (prefix[r+1] - (prefix[l] * power[r-l+1])) % MOD + MOD) % MOD;
        return res;
    }
};
```

6.3. Single Hashing

```
// double string hashing to further reduce collision
struct StringHash {
    // two large primes to reduce collision
    const long long MOD1 = 1000000007;
    const long long MOD2 = 1000000009;
    const long long P1 = 31;
    const long long P2 = 37;

    vector<long long> prefix1, prefix2;
    vector<long long> power1, power2;

    StringHash(const string &s) {
```

```
int n = s.size();
prefix1.resize(n+1, 0);
prefix2.resize(n+1, 0);
power1.resize(n+1, 1);
power2.resize(n+1, 1);

for (int i = 1; i <= n; i++) {
    int val = s[i-1]; // works for all ASCII characters
    power1[i] = (power1[i-1] * P1) % MOD1;
    power2[i] = (power2[i-1] * P2) % MOD2;

    prefix1[i] = (prefix1[i-1] * P1 + val) % MOD1;
    prefix2[i] = (prefix2[i-1] * P2 + val) % MOD2;
}

// get double hash of substring [l, r] inclusive
pair<long long, long long> getHash(int l, int r) {
    long long h1 = (prefix1[r+1] - (prefix1[l] * power1[r-l+1])) % MOD1 + MOD1;
    long long h2 = (prefix2[r+1] - (prefix2[l] * power2[r-l+1])) % MOD2 + MOD2;
    return {h1, h2};
}

// quick compare of substrings from two StringHash objects
bool isEqual(int l1, int r1, int l2, int r2,
StringHash &other) {
    return getHash(l1, r1) == other.getHash(l2, r2);
}
};
```

6.4. Largest Substring More than K

```
int n;
int max_oc(int len) {
    map<pair<int, int>, int> mp;
    for (int i = 0; i + len - 1 < n; i++)
        mp[get_hash(i, i + len - 1)]++;

    int ans = 0;
    for (auto [x, y] : mp)
        ans = max(ans, y);

    return ans;
}

int32_t main() {
    prec();
    string s;
    cin >> s;
    build(s);
```

```

int k;
cin >> k;
n = s.size();
int l = 1, r = s.size(), ans = -1;
while (l <= r) {
    int mid = (l + r) >> 1;
    if (max_oc(mid) >= k)
        ans = mid, l = mid + 1;
    else
        r = mid - 1;
}
cout << ans << '\n';
return 0;

```

6.5. Longest Common Prefix Of Two Substrings

```

int lcp(int i, int j, int x, int y) { // O(log n)
    int l = 1, r = min(j - i + 1, y - x + 1), ans = 0;
    while (l <= r) {
        int mid = l + r >> 1;
        if (get_hash(i, i + mid - 1) == get_hash(x, x + mid - 1)) {
            ans = mid;
            l = mid + 1;
        } else {
            r = mid - 1;
        }
    }
    return ans;
}

/*
given a string s of size n and q queries of type i,
j, x, y.
find the LCP of Substring s[i...j] and s[x...y].
1 <= n, q <= 1e5
*/

```

6.6. Longest Common Substring

```

const int N = 1e5 + 9;
const int p1 = 137, mod1 = 127657753, p2 = 277, mod2 = 987654319;

int power(long long n, long long k, int mod) {
    int ans = 1 % mod; n %= mod;
    if (n < 0) n += mod;
    while (k) {
        if (k & 1) ans = (long long)ans * n % mod;
        n = (long long)n * n % mod;
        k >>= 1;
    }
    return ans;
}

int ip1, ip2;

```

```

pair<int, int> pw[N], ipw[N];

void prec(){ // O(n)
    pw[0] = {1, 1};
    for (int i = 1; i < N; i++) {
        pw[i].first = 1LL * pw[i - 1].first * p1 % mod1;
        pw[i].second = 1LL * pw[i - 1].second * p2 % mod2;
    }
    ip1 = power(p1, mod1 - 2, mod1);
    ip2 = power(p2, mod2 - 2, mod2);
    ipw[0] = {1, 1};
    for (int i = 1; i < N; i++) {
        ipw[i].first = 1LL * ipw[i - 1].first * ip1 % mod1;
        ipw[i].second = 1LL * ipw[i - 1].second * ip2 % mod2;
    }
}

pair<int, int> string_hash(string s){ // O(n)
    int n = s.size();
    pair<int, int> hs({0, 0});
    for (int i = 0; i < n; i++) {
        hs.first += 1LL * s[i] * pw[i].first % mod1;
        hs.first %= mod1;
        hs.second += 1LL * s[i] * pw[i].second % mod2;
        hs.second %= mod2;
    }
    return hs;
}

struct Hashing{
    pair<int, int> pref[N];
    void build(string s){ // O(n)
        int n = s.size();
        for (int i = 0; i < n; i++) {
            pref[i].first = 1LL * s[i] * pw[i].first % mod1;
            if (i) pref[i].first = (pref[i].first + pref[i - 1].first) % mod1;
            pref[i].second = 1LL * s[i] * pw[i].second % mod2;
            if (i) pref[i].second = (pref[i].second + pref[i - 1].second) % mod2;
        }
    }

    pair<int, int> get_hash(int i, int j){ // O(1)
        pair<int, int> hs({0, 0});
        hs.first = pref[j].first;
        if (i) hs.first = (hs.first - pref[i - 1].first + mod1) % mod1;
        hs.first = 1LL * hs.first * ipw[i].first % mod1;
        hs.second = pref[j].second;
        if (i) hs.second = (hs.second - pref[i - 1].second + mod2) % mod2;
        hs.second = 1LL * hs.second * ipw[i].second % mod2;
        return hs;
    }
}

} A, B;

int n;
string a, b;
string res;

bool ok(int k){ // is there a k length substring
    set<pair<int, int>> substring_hashes_in_a;
    for (int i = 0; i + k - 1 < n; i++)
        substring_hashes_in_a.insert(A.get_hash(i, i + k - 1));
    for (int i = 0; i + k - 1 < n; i++){
        auto substring_hash_in_b = B.get_hash(i, i + k - 1);
        if (substring_hashes_in_a.find(substring_hash_in_b) != substring_hashes_in_a.end()){
            res = b.substr(i, k);
            return true;
        }
    }
    return false;
}

int32_t main(){
    prec();
    cin >> n;
    cin >> a >> b;
    A.build(a);
    B.build(b);
    int l = 1, r = n, ans = 0;
    while (l <= r){
        int mid = (l + r) / 2;
        if (ok(mid)){ ans = mid; l = mid + 1; }
        else { r = mid - 1; }
    }
    ok(ans);
    cout << res << '\n'; // O(n log^2 n)
    return 0;
} // Find the Longest Common Substring of Two Strings

```

6.7. Pattern Matching

```

const int N = 1e6 + 9;
const int p1 = 137, mod1 = 127657753, p2 = 277, mod2 = 987654319;

int power(long long n, long long k, int mod) {
    int ans = 1 % mod;
    n %= mod;
    if (n < 0) n += mod;
    while (k) {
        if (k & 1) ans = (long long)ans * n % mod;
        n = (long long)n * n % mod;
        k >>= 1;
    }
    return ans;
}

```

```

}

int ip1, ip2;
pair<int, int> pw[N], ipw[N];

void prec(){
    pw[0] = {1, 1};
    for (int i = 1; i < N; i++){
        pw[i].first = 1LL * pw[i - 1].first * p1 %
                      mod1;
        pw[i].second = 1LL * pw[i - 1].second * p2 %
                      mod2;
    }
    ip1 = power(p1, mod1 - 2, mod1);
    ip2 = power(p2, mod2 - 2, mod2);
    ipw[0] = {1, 1};
    for (int i = 1; i < N; i++){
        ipw[i].first = 1LL * ipw[i - 1].first * ip1 %
                      mod1;
        ipw[i].second = 1LL * ipw[i - 1].second * ip2 %
                      mod2;
    }
}

pair<int, int> string_hash(string s){
    int n = s.size();
    pair<int, int> hs({0, 0});
    for (int i = 0; i < n; i++){
        hs.first += 1LL * s[i] * pw[i].first % mod1;
        hs.first %= mod1;
        hs.second += 1LL * s[i] * pw[i].second % mod2;
        hs.second %= mod2;
    }
    return hs;
}

pair<int, int> pref[N];

void build(string s){
    int n = s.size();
    for (int i = 0; i < n; i++){
        pref[i].first = 1LL * s[i] * pw[i].first % mod1;
        if (i) pref[i].first = (pref[i].first + pref[i - 1].first) % mod1;
        pref[i].second = 1LL * s[i] * pw[i].second % mod2;
        if (i) pref[i].second = (pref[i].second + pref[i - 1].second) % mod2;
    }
}

pair<int, int> get_hash(int i, int j){
    assert(i <= j);
    pair<int, int> hs({0, 0});
    hs.first = pref[j].first;
    if (i) hs.first = (hs.first - pref[i - 1].first +
                      mod1) % mod1;
    hs.first = 1LL * hs.first * ipw[i].first % mod1;
    hs.second = pref[j].second;
}

```

```

if (i) hs.second = (hs.second - pref[i - 1].second +
                     mod2) % mod2;
hs.second = 1LL * hs.second * ipw[i].second % mod2;
return hs;
}

int32_t main(){
    prec();
    string a, b;
    cin >> a >> b;
    build(a);
    int ans = 0, n = a.size(), m = b.size();
    auto hash_b = string_hash(b);
    for (int i = 0; i + m - 1 < n; i++)
        ans += get_hash(i, i + m - 1) == hash_b;
    cout << ans << '\n';
    return 0;
}

/*
Given a string a and pattern b, count the number of
positions where the pattern occurs in the string
.

Input:
saippuaakaappias
pp
Output:2*/

```

6.8. Two Strings Are Equal Or Not

7. Mathematical Formulas

7.1. Math Formula

```

summation a to b = ((b-a+1)*(a+b))/2

Orientation of 3 point
// a(x,y),b(x,y),c(x,y),d(x,y)
// slop ab = (yb-ya)/(xb-xa);
// slop bc = (yc-yb)/(xc-xb);
#include <bits/stdc++.h>
using namespace std;
int main() {
    // Orientation of 3 points(a,b,c)->
    // Counterclockwise (abc<0): ab slop uoper c point
    // Clockwise (abc>0): ab slop nicha c point
    // Collinear(abc=0): ab slop soja c point

    //how to calculate
    ABCslop = slopAB - slopBC
    = (By-Ay)/(Bx-Ax) - (By-Cy)/(Bx-Cx);
    = ((By-Ay))*(Cx-Bx) - (Cy-By)*(Bx-Ax);
}

```

$$\sum_{1 \leq i <= j \leq n} a_i a_j = \frac{(\sum_i a_i)^2 - \sum_i a_i^2}{2}.$$

The sum of $a_i a_j$ means we take all pairs of elements in the range $(i - j)$, multiply each pair, and then add them together

For the array [1, 2, 3, 4, 5]:

$$\sum_{i=1}^5 a_i^2 = 1^2 + 2^2 + 3^2 + 4^2 + 5^2,$$

$$\sum_{i=1}^5 a_i = 1 + 2 + 3 + 4 + 5.$$

7.2. Catalan Numbers

1, 1, 2, 5, 14, 42, 132, 429, 1430.....

Binomial coefficient: $C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{(n+1)! n!}$

Recursive Formula: $C_n = C_0 C_{n-1} + C_1 C_{n-2} + \dots + C_{n-1} C_0$

Another recurrence relation: $C_{n+1} = \frac{2(2n+1)}{n+2} C_n$

use : Counting Valid Parentheses Sequences, Number of Binary Search Trees, Ways to Triangulate a Convex Polygon,...etc

7.3. Geometry Formulas 1

1. Triangle Given: base b , height h , sides a, b, c

$$\text{semiperimeter } s = \frac{a+b+c}{2}, \text{Area} = \frac{1}{2}bh, \text{Perimeter} = a+b+c$$

$$\text{Heron Formula Area} = \sqrt{s(s-a)(s-b)(s-c)}$$

2. Square Given: side s

$$\text{Area} = s^2, \text{Perimeter} = 4s, \text{Diagonal} = s\sqrt{2}$$

3. Rectangle Given: length l , width w

$$\text{Area} = lw, \text{Perimeter} = 2(l+w), \text{Diagonal} = \sqrt{l^2 + w^2}$$

4. Trapezoid (Trapezium) Given: bases a, b , height h

$$\text{Area} = \frac{1}{2}(a+b)h, \text{Midsegment} = \frac{a+b}{2}, \text{Perimeter} = a+b+c+d$$

5. Parallelogram Given: base b , height h , sides a, b

$$\text{Area} = bh, \text{Perimeter} = 2(a+b)d_1 = \sqrt{a^2 + b^2 + 2ab \cos \theta}$$

Diagonal formulas:

$$d_1 = \sqrt{a^2 + b^2 + 2ab \cos \theta}, \quad d_2 = \sqrt{a^2 + b^2 - 2ab \cos \theta}$$

6. Circle

Given: radius r

$$\text{Area} = \pi r^2, \text{Circumference} = 2\pi r, \text{Diameter} = 2r$$

7. Sphere

Given: radius r

$$\text{Surface Area} = 4\pi r^2, \text{Volume} = \frac{4}{3}\pi r^3$$

8. Cylinder

Given: radius r , height h

$$\text{Lateral Surface Area} = 2\pi rh$$

$$\text{Total Surface Area} = 2\pi r(h+r), \text{Volume} = \pi r^2 h$$

9. Cone

Given: radius r , height h , slant height $l = \sqrt{r^2 + h^2}$

$$\text{Lateral Surface Area} = \pi rl$$

$$\text{Total Surface Area} = \pi r(r+l), \text{Volume} = \frac{1}{3}\pi r^2 h$$

10. Pyramid

Given: base area B , height h , slant height l

$$\text{Volume} = \frac{1}{3}Bh, \text{Lateral Surface Area (square base)} = 2al$$

$$\text{Total Surface Area (square base)} = a^2 + 2al$$

11. Rectangular Prism (Cuboid)

Given: length l , width w , height h

$$\text{Surface Area} = 2(lw + lh + wh), \text{Volume} = lwh$$

$$\text{Diagonal} = \sqrt{l^2 + w^2 + h^2}$$

12. Cube

Given: side s

$$\text{Surface Area} = 6s^2, \text{Volume} = s^3, \text{Diagonal} = s\sqrt{3}$$

7.4. Geometry Formulas 2

Distance, Dot, Cross, Orientation :

$$\text{Distance of two points: } d = \sqrt{(x_2 - x_1)^2 + (y_2 - y_1)^2}$$

$$\text{Dot product: } \vec{a} \cdot \vec{b} = a_x b_x + a_y b_y$$

$$\text{Cross product (2D): } \vec{a} \times \vec{b} = a_x b_y - a_y b_x$$

$$\text{Orientation: } \text{orient}(a, b, c) = (b - a) \times (c - a)$$

Line through

$$(x_1, y_1), (x_2, y_2) : Ax + By + C = 0$$

where

$$A = y_1 - y_2, \quad B = x_2 - x_1, \quad C = x_1 y_2 - x_2 y_1$$

Distance from point to line:

$$d = \frac{|Ax_0 + By_0 + C|}{\sqrt{A^2 + B^2}}$$

Polygon Formula

$$\text{Polygon Area} = \frac{1}{2} \left| \sum_{i=1}^n (x_i y_{i+1} - x_{i+1} y_i) \right|$$

Polygon Perimeter:

$$\sum_{i=1}^n d(p_i, p_{i+1})$$

Circle Geometry:

$$\text{Chord length} = 2\sqrt{r^2 - d^2}$$

$$\text{Sector area: } \frac{1}{2}r^2\theta, \quad \text{Arc length: } s = r\theta$$

$$\text{Triangle area using circumradius } A = \frac{abc}{4R}$$

Vectors

$$\text{Angle between vectors } \theta = \cos^{-1} \left(\frac{\vec{a} \cdot \vec{b}}{|\vec{a}| |\vec{b}|} \right)$$

$$\text{Projection: } \text{proj}_{\vec{b}}(\vec{a}) = \frac{\vec{a} \cdot \vec{b}}{|\vec{b}|^2} \vec{b}$$

Number Theory:

Modular inverse (for prime modulus): $a^{-1} \equiv a^{m-2} \pmod{m}$

$$\text{Combination: } \binom{n}{r} = \frac{n!}{r!(n-r)!}$$

$$\text{Catalan numbers: } C_n = \frac{1}{n+1} \binom{2n}{n}$$

7.5. Matrices and Determinants

- The number of rows or columns (Order) in a square matrix is called its order.
- A matrix whose all elements are zero is called a null matrix.
- Matrix multiplication AB is possible only when the number of columns in matrix A is equal to the number of rows in matrix B .
- If A is a non-singular matrix, then the inverse is given by:

$$A^{-1} = \frac{1}{|A|} \text{Adj}(A)$$

- Adjoint of A : $\text{Adj}(A) = (\text{Cofactor Matrix})^T$.
- Determinant minors and cofactors (for 3x3 D):

$$D = \begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix}$$

Minor of a_{11} : $m_{11} = a_{22}a_{33} - a_{23}a_{32}$ Cofactor of a_{ij} : $A_{ij} = (-1)^{i+j}m_{ij}$ Example: $A_{11} = a_{22}a_{33} - a_{23}a_{32}$, $A_{12} = -m_{12}$.

7. Important Properties of Determinants:

- If any two columns (or rows) are identical, its value is zero.

$$\begin{vmatrix} a & a & c \\ b & b & f \\ d & d & g \end{vmatrix} = 0$$

- If all elements in any one row (or column) of a determinant are multiplied by a scalar k , the value is k times the value of the original determinant.

$$\begin{vmatrix} ka & kb & kc \\ d & e & f \\ g & h & i \end{vmatrix} = k \begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix}$$

- The value of the determinant remains unchanged if a multiple of one row (or column) is added to another.

$$D = \begin{vmatrix} a_1 & b_1 & c_1 \\ a_2 & b_2 & c_2 \\ a_3 & b_3 & c_3 \end{vmatrix} \implies D = \begin{vmatrix} a_1 + mb_1 & b_1 & c_1 \\ a_2 + mb_2 & b_2 & c_2 \\ a_3 + mb_3 & b_3 & c_3 \end{vmatrix}$$

- If two adjacent rows (or columns) are interchanged, the sign of the determinant's value changes.

- Cramer's Rule for solving simultaneous linear equations: Given: $a_1x + b_1y + c_1z = d_1$, $a_2x + b_2y + c_2z = d_2$, $a_3x + b_3y + c_3z = d_3$.

$$D = \begin{vmatrix} a_1 & b_1 & c_1 \\ a_2 & b_2 & c_2 \\ a_3 & b_3 & c_3 \end{vmatrix} \neq 0$$

$$D_x = \begin{vmatrix} d_1 & b_1 & c_1 \\ d_2 & b_2 & c_2 \\ d_3 & b_3 & c_3 \end{vmatrix}, \quad D_y = \begin{vmatrix} a_1 & d_1 & c_1 \\ a_2 & d_2 & c_2 \\ a_3 & d_3 & c_3 \end{vmatrix}$$

$$D_z = \begin{vmatrix} a_1 & b_1 & d_1 \\ a_2 & b_2 & d_2 \\ a_3 & b_3 & d_3 \end{vmatrix}$$

Solution: $x = \frac{D_x}{D}$, $y = \frac{D_y}{D}$, and $z = \frac{D_z}{D}$.

7.6. Vectors

- If $\vec{a} = a_1\mathbf{i} + a_2\mathbf{j} + a_3\mathbf{k}$, the magnitude is:

$$|\vec{a}| = \sqrt{a_1^2 + a_2^2 + a_3^2}$$

- Position Vector: The position vector of point (x, y, z) is:

$$\vec{r} = x\mathbf{i} + y\mathbf{j} + z\mathbf{k}$$

- Vector Subtraction: If $\vec{OP} = x_1\mathbf{i} + y_1\mathbf{j} + z_1\mathbf{k}$ and $\vec{OQ} = x_2\mathbf{i} + y_2\mathbf{j} + z_2\mathbf{k}$, then:

$$\vec{PQ} = \vec{OQ} - \vec{OP} = (x_2 - x_1)\mathbf{i} + (y_2 - y_1)\mathbf{j} + (z_2 - z_1)\mathbf{k}$$

- Section Formula:

- Internal Division: If R divides AB internally in the ratio $m : n$, and \vec{a}, \vec{b} are position vectors of A and B , then:

$$\vec{r} = \frac{m\vec{b} + n\vec{a}}{m + n}$$

- Midpoint: If C is the midpoint of AB , the position vector of C is:

$$\vec{c} = \frac{\vec{a} + \vec{b}}{2}$$

- Unit vector parallel to \vec{a} : $\pm \frac{\vec{a}}{|\vec{a}|}$.

- Unit vector perpendicular to \vec{a} and \vec{b} : $\pm \frac{\vec{a} \times \vec{b}}{|\vec{a} \times \vec{b}|}$.

- Unit vector perpendicular to the plane containing \vec{a} and \vec{b} : $\pm \frac{\vec{a} \times \vec{b}}{|\vec{a} \times \vec{b}|}$.
- Vector addition: If $\vec{a} = a_1\mathbf{i} + a_2\mathbf{j} + a_3\mathbf{k}$ and $\vec{b} = b_1\mathbf{i} + b_2\mathbf{j} + b_3\mathbf{k}$, then:

$$\vec{a} + \vec{b} = (a_1 + b_1)\mathbf{i} + (a_2 + b_2)\mathbf{j} + (a_3 + b_3)\mathbf{k}$$

- Scalar Product (Dot Product):

$$\vec{a} \cdot \vec{b} = |\vec{a}||\vec{b}| \cos \theta = a_1b_1 + a_2b_2 + a_3b_3$$

- Unit Vector Dot Products:

$$\mathbf{i} \cdot \mathbf{i} = \mathbf{j} \cdot \mathbf{j} = \mathbf{k} \cdot \mathbf{k} = 1, \quad \mathbf{i} \cdot \mathbf{j} = \mathbf{j} \cdot \mathbf{k} = \mathbf{k} \cdot \mathbf{i} = 0$$

- Vector Product (Cross Product):

- Magnitude: $|\vec{a} \times \vec{b}| = |\vec{a}||\vec{b}| \sin \theta$.

- Cartesian form:

$$\vec{a} \times \vec{b} = \begin{vmatrix} \mathbf{i} & \mathbf{j} & \mathbf{k} \\ a_1 & a_2 & a_3 \\ b_1 & b_2 & b_3 \end{vmatrix}$$

- Property: $\vec{a} \times \vec{b} = -(\vec{b} \times \vec{a})$.

- Unit Vector Cross Products: $\mathbf{i} \times \mathbf{i} = 0$. Cyclic relations:

$$\mathbf{i} \times \mathbf{j} = \mathbf{k}, \quad \mathbf{j} \times \mathbf{k} = \mathbf{i}, \quad \mathbf{k} \times \mathbf{i} = \mathbf{j}$$

- Condition for Perpendicularity: If $\vec{a} \perp \vec{b}$, then $\vec{a} \cdot \vec{b} = 0$.

- Scalar projection of \vec{a} onto \vec{b} :

$$a \cos \theta = \frac{\vec{a} \cdot \vec{b}}{|\vec{b}|}$$

- Vector projection of \vec{a} onto \vec{b} (along \vec{b}):

$$\frac{(\vec{a} \cdot \vec{b})}{|\vec{b}|^2} \vec{b}$$

16. Scalar projection of \vec{b} onto \vec{a} :

$$\frac{\vec{a} \cdot \vec{b}}{|\vec{a}|}$$

17. Volume of the parallelepiped formed by $\vec{a}, \vec{b}, \vec{c}$:

$$\begin{vmatrix} a_1 & a_2 & a_3 \\ b_1 & b_2 & b_3 \\ c_1 & c_2 & c_3 \end{vmatrix} \text{ cubic units}$$

18. Condition for Coplanarity: If vectors $\vec{a}, \vec{b}, \vec{c}$ are coplanar, then $\vec{a} \cdot (\vec{b} \times \vec{c}) = 0$.

19. Area Formulas (Parallelogram and Triangle):

- i. Area of parallelogram formed by \vec{a} and \vec{b} : $|\vec{a} \times \vec{b}|$ sq units.

- ii. Area of triangle formed by \vec{a} and \vec{b} : $\frac{1}{2}|\vec{a} \times \vec{b}|$ sq units.

7.7. Straight Lines

1. Coordinates on the axes: A point on the x -axis is $(x, 0)$; a point on the y -axis is $(0, y)$.

2. Distance of point (x, y) : From the x -axis is $|y|$; from the y -axis is $|x|$.

3. Cartesian (x, y) and Polar (r, θ) Coordinate relations:

$$x = r \cos \theta, \quad y = r \sin \theta$$

$$r = \sqrt{x^2 + y^2}, \quad \theta = \tan^{-1}\left(\frac{y}{x}\right)$$

(Conditions: $0 \leq r < \infty$ and $0 \leq \theta < 2\pi$).

4. Location of Polar Coordinates (r, θ) from Cartesian (x, y) :

- i. For the **first quadrant** ($x > 0, y > 0$): $\theta = \tan^{-1}\left(\frac{y}{x}\right); 0 < \theta < 90^\circ$.

- ii. For the **second quadrant** ($-x, y$): $\theta = 180^\circ - \tan^{-1}\left(\frac{y}{|x|}\right); 90^\circ < \theta < 180^\circ$.

- iii. For the **third quadrant** ($-x, -y$): $\theta = 180^\circ + \tan^{-1}\left(\frac{|y|}{|x|}\right); 180^\circ < \theta < 270^\circ$.

- iv. For the **fourth quadrant** ($x, -y$): $\theta = 360^\circ - \tan^{-1}\left(\frac{|y|}{x}\right); 270^\circ < \theta < 360^\circ$.

- v. Alternatively, for Cartesian (x, y) , the angle θ is: $\theta = \tan^{-1}\left(\frac{y}{x}\right)$ (If $x > 0$) or $\pi + \tan^{-1}\left(\frac{y}{x}\right)$ (If $x < 0$)

5. **Distance Formula** between two points (x_1, y_1) and (x_2, y_2) :

$$d = \sqrt{(x_2 - x_1)^2 + (y_2 - y_1)^2}$$

6. Distance between two points (x_1, b) and (x_2, b) : $|x_2 - x_1|$ (since $y_1 = y_2 = b$).

7. Distance between two points (a, y_1) and (a, y_2) : $|y_2 - y_1|$ (since $x_1 = x_2 = a$).

8. Distance between two **Polar Points** (r_1, θ_1) and (r_2, θ_2) :

$$d = \sqrt{r_1^2 + r_2^2 - 2r_1r_2 \cos(\theta_1 - \theta_2)}$$

9. **Division Formula (Internal)**: Coordinates of point dividing line segment (x_1, y_1) and (x_2, y_2) **internally** in ratio $m_1 : m_2$:

$$\left(\frac{m_1x_2 + m_2x_1}{m_1 + m_2}, \frac{m_1y_2 + m_2y_1}{m_1 + m_2} \right)$$

10. **Division Formula (External)**: Coordinates of point dividing line segment (x_1, y_1) and (x_2, y_2) **externally** in ratio $m_1 : m_2$:

$$\left(\frac{m_1x_2 - m_2x_1}{m_1 - m_2}, \frac{m_1y_2 - m_2y_1}{m_1 - m_2} \right)$$

11. **Midpoint Formula**: Midpoint of line segment between (x_1, y_1) and (x_2, y_2) :

$$\left(\frac{x_1 + x_2}{2}, \frac{y_1 + y_2}{2} \right)$$

12. **Centroid Formula**: Centroid of a triangle with vertices $(x_1, y_1), (x_2, y_2), (x_3, y_3)$:

$$\left(\frac{x_1 + x_2 + x_3}{3}, \frac{y_1 + y_2 + y_3}{3} \right)$$

13. **Area of Triangle** $\triangle ABC$ with vertices $A(x_1, y_1), B(x_2, y_2), C(x_3, y_3)$:

$$\text{Area} = \frac{1}{2} \begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix} \text{ sq units}$$

14. Area of Triangle with origin $(0, 0)$ and vertices $(x_1, y_1), (x_2, y_2)$:

$$\text{Area} = \frac{1}{2} |x_1y_2 - x_2y_1| \text{ sq units}$$

15. **Collinearity Condition**: Three points $A(x_1, y_1), B(x_2, y_2), C(x_3, y_3)$ are collinear if the area of $\triangle ABC$ is zero:

$$\frac{1}{2} |x_1(y_2 - y_3) + x_2(y_3 - y_1) + x_3(y_1 - y_2)| = 0$$

(or using the determinant form: $\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix} = 0$)

16. **Area of Quadrilateral** $ABCD$ with vertices $(x_1, y_1), (x_2, y_2), (x_3, y_3), (x_4, y_4)$:

$$\text{Area} = \frac{1}{2} \begin{vmatrix} x_1 & x_2 & x_3 & x_4 & x_1 \\ y_1 & y_2 & y_3 & y_4 & y_1 \end{vmatrix}$$

17. **Slope (Gradient)** of a Straight Line:

- i. The **slope** is $m = \tan \theta$, where θ is the angle with the positive x -axis ($0^\circ \leq \theta < 180^\circ$).

- ii. Slope of the **x -axis** (or parallel to x -axis) is $m = 0$.

- iii. Slope of the **y -axis** (or parallel to y -axis) is **undefined**.

- iv. Slope of the line passing through (x_1, y_1) and (x_2, y_2) : $m = \frac{y_2 - y_1}{x_2 - x_1}$.

- v. Slope of the line $ax + by + c = 0$: $m = -\frac{a}{b}$.

18. Equation of the **x -axis**: $y = 0$.

19. Equation of the **y -axis**: $x = 0$.

20. Equation of a line **parallel to the x -axis** (or horizontal line): $y = b$.

21. Equation of a line **parallel to the y -axis** (or vertical line): $x = a$.

22. Equation of a line **passing through the origin**: $y = mx$.

23. **Different Forms of the Equation of a Straight Line**:

- i. **Point-Slope form** (passing through (x_1, y_1) with slope m): $y - y_1 = m(x - x_1)$.

- ii. **Slope-Intercept form** (with slope m and y -intercept c): $y = mx + c$.

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| <p>iii. Two-Point form (passing through (x_1, y_1) and (x_2, y_2)): $\frac{y-y_1}{y_2-y_1} = \frac{x-x_1}{x_2-x_1}$.</p> <p>iv. Intercept form (with x-intercept a and y-intercept b): $\frac{x}{a} + \frac{y}{b} = 1$.</p> | <p>v. General form: $ax + by + c = 0$.</p> <p>24. Area of the triangle formed by the line $\frac{x}{a} + \frac{y}{b} = 1$ and the coordinate axes: $\frac{1}{2} ab$ sq units.</p> <p>25. Perpendicular Distance from a point (x_1, y_1) to the line $ax + by + c = 0$:</p> $d = \frac{ ax_1 + by_1 + c }{\sqrt{a^2 + b^2}}$ |
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