

Lecture 11

Johnson - Lindenstraus Lemma

clarification

Suppose X is a zero mean sub-G(k^2) random variable

\Rightarrow

$$X \in \text{sub-G}(k^2) \Rightarrow \text{for } \lambda < \frac{1}{k}$$

$$\mathbb{E}[\exp(\lambda^2 X)] \leq \exp(k^2 \lambda^2)$$

$$\exists c \quad \mathbb{E}[e^{X^2/c}] \stackrel{?}{\leq} 2$$

$$\frac{1}{c} \stackrel{\text{orange}}{\leftarrow} \frac{k^2}{k} \leq e^{\frac{k^2}{c}} \stackrel{\text{orange}}{\leq} 2$$

$$\frac{k^2}{c} \stackrel{\text{orange}}{\uparrow} \leq \ln 2$$

set $c = \frac{k^2}{\ln 2}$ so both condition hold.

□

Norm of a vector of Gaussians

Let $\vec{v} = (v_1, \dots, v_d)$ be a vector in \mathbb{R}^d

Suppose $v_i \sim N(0, 1)$ are drawn from standard normal distribution.

what can we say about $\|v\|_2^2$?

we have shown $v_i^2 \in \text{Sub } E(2^2, 4)$

Let $X = \sum_{i=1}^d v_i^2$. X is sum of

d independent $\text{Sub } F \Rightarrow$

$X \in \text{Sub } E(2^2 \cdot d, 4)$

$$\Pr [\|v\|_2^2 - d \mid \geq d\varepsilon]$$

$$\leq \Pr \left[\left| \frac{1}{d} \sum v_i^2 - 1 \right| \geq \varepsilon \right]$$

$$\leq 2 \exp \left(-d \cdot \min \left(\frac{\varepsilon^2}{2}, \frac{\varepsilon}{2} \right) \right)$$

$$\leq \begin{cases} 2 e^{-\frac{d\varepsilon^2}{2}} & 0 \leq \varepsilon \leq 1 \\ 2 c^{-\frac{d\varepsilon}{2}} & \varepsilon > 1 \end{cases}$$

using CLT:

$$E[\sum v_i^2] = d$$

$$\text{Var}[v_i^2] = E[v_i^4] - \underbrace{E[v_i^2]^2}_{=d}$$

$$\leq 3 - 1 = 2$$

$$\Rightarrow \text{Var}(\sum v_i^2) = 2d$$

CLT: $\frac{\sum v_i^2 - d}{\sqrt{2d}} \rightarrow N(0,1)$

$$\Pr[|\sum v_i^2 - d| \geq \varepsilon d] \underset{d \rightarrow \infty}{\approx}$$

$$\Pr[|Z| \geq \sqrt{\frac{d}{2}} \varepsilon] \approx e^{-\frac{d\varepsilon^2}{2}}$$

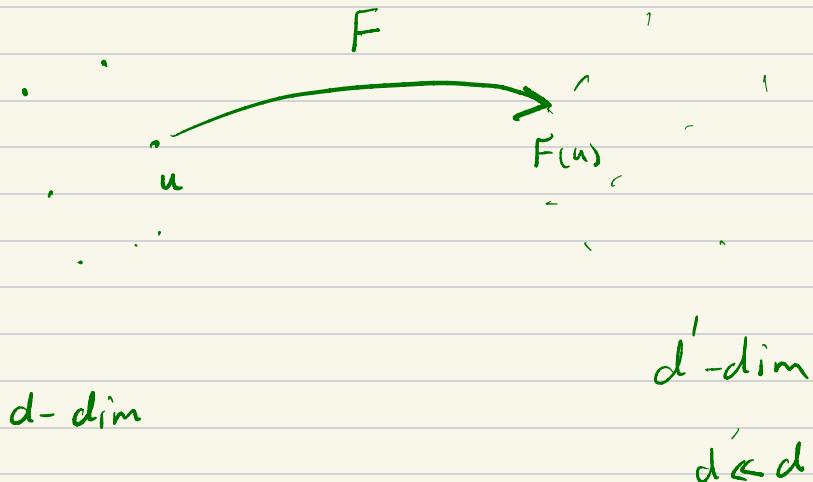
$\overbrace{\quad}^{\infty ?}$

$\infty ?$

Is our previous
bound loose?

Dimensionality reduction

Suppose we have n points in a d -dim space. Our hope is to embed the points to a lower dim (say d') such that the Euclidian distances between pair of points is preserved.



Example: k-means clustering

we have n points. our goal is to partition the points into k clusters such that

sum of distances to the mean of the cluster is minimized. This is equivalent

to ask for a partition $S = \{S_1, \dots, S_k\}$

that minimizes the following

$$\arg \min_S \sum_{S_i \in S} \sum_{x, y \in S_i} \|x - y\|_2^2$$

Generally, this problem is NP-hard. However, approximation algorithms exist with time complexity $\propto O(d)$. If we have an embedding to reduce the dimension, we can solve this problem faster.

Johnson - Lindenstrauss

Lemma

$\rightarrow u_1, \dots, u_n$

Given n points in \mathbb{R}^d and an integer

$d' \geq \frac{8 \log n/\delta}{\epsilon^2}$, there exists a randomized linear map $F: \mathbb{R}^d \rightarrow \mathbb{R}^{d'}$

such that with probability $1-\delta$:

$$(1-\epsilon) \|u_i - u_j\|_2^2 \leq \|F(u_i) - F(u_j)\|_2^2 \leq (1+\epsilon) \|u_i - u_j\|_2^2$$

The map:

Suppose we have an $d' \times d$ matrix

$M \in \mathbb{R}^{d' \times d}$ such that every

entry of M is drawn from a normal distribution $N(0, \frac{1}{d'})$

$$\forall i,j \quad M_{ij} \sim N(0, \frac{1}{d'})$$

$$\text{Let } F(u) = Mu$$

$$d' \begin{bmatrix} F(u) \end{bmatrix} = d \begin{bmatrix} u \end{bmatrix}$$

Diagram illustrating the dimensions of the matrices:

- The matrix $F(u)$ has dimension $d' \times d$.
- The matrix M has dimension $d \times d$.
- The vector u has dimension $d \times 1$.

Arrows indicate the dimensions: a pink arrow labeled d' points from the left side of the first matrix to its top; a pink arrow labeled d points from the right side of the second matrix to its top; a black arrow labeled d points from the bottom of the third matrix to its right side.

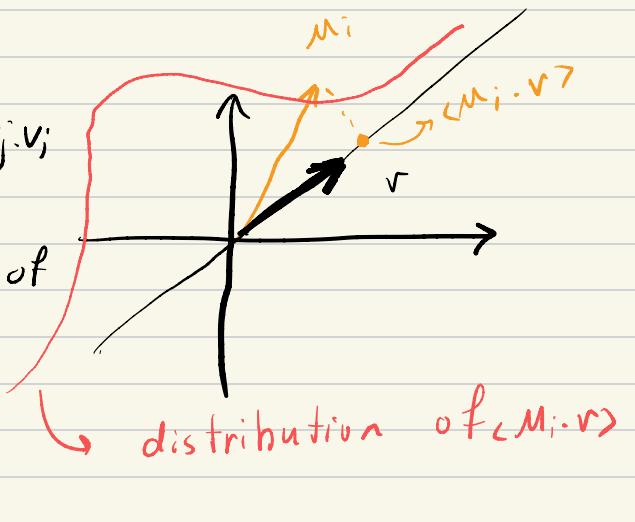
Proof of lemma

step 1 : projection of a random vector on to a fix direction

Suppose we have a fixed vector v of size $\|v\|_2 = 1$. Let $M_i \in \mathbb{R}^d$ such that $M_{ij} \sim N(0, \frac{1}{d})$. We claim that $\langle M_i \cdot v \rangle$ is a Gaussian random variable

$$Z = \langle M_i \cdot v \rangle = \sum_{j=1}^d M_{ij} v_j$$

Z is a linear combination of Gaussians.



$$\Rightarrow Z \sim N(0, ?)$$

$$\text{Var}(Z) = \text{Var} \left\{ \sum_{j=1}^d M_{ij} \cdot v_j \right\}$$

independence \rightarrow
of M_{ij} 's

$$= \sum_{j=1}^d v_j^2 \cdot \underbrace{\text{Var}[M_{ij}]}_{= d'}$$

$$= \frac{\|v\|_2^2}{d'} = \frac{1}{d'}$$

$$\Rightarrow \langle M_i \cdot v \rangle \sim N(0, \frac{1}{d'})$$

step 2 if $\|v\|_2^2 = 1$, then $\|F(v)\|_2^2$
is a sub-exponential r.v.

$$\|F(v)\|_2^2 = \|M \cdot v\|_2^2 = \sum_{i=1}^{d'} (\langle M_i \cdot v \rangle)^2$$

by step 1, $\langle M_i \cdot v \rangle \sim N(0, \frac{1}{d'})$

In distribution: $\langle M_i \cdot v \rangle \rightarrow \frac{1}{\sqrt{d'}} Z_i$

where $Z_i \sim N(0, 1)$

$$\Rightarrow \text{In distribution } \|F(v)\|_2^2 = \sum_{i=1}^{d'} \frac{Z_i^2}{d'}$$

Recall, earlier we have shown:

$$\Pr [| \| F(\mathbf{v}) \|_2^2 - 1 | \geq \varepsilon]$$

$$= \Pr [\left| \sum_{i=1}^{d'} \frac{z_i^2}{d'} - 1 \right| \geq \varepsilon]$$

$$\leq 2 \exp \left(- \frac{d' \varepsilon^2}{2} \right)$$

for $\varepsilon \in (0, 1]$

step 3 Suppose $\|u\|_2 \neq 0$

$\|F(u)\|_2^2$ has a similar tail

bound even when $\|u\|_2 \neq 1$

Let $v = \frac{u}{\|u\|_2}$. Clearly $\|v\|_2^2 = 1$

$$\Pr \left[\left| \frac{\|F(u)\|_2^2}{\|u\|_2^2} - 1 \right| > \varepsilon \right]$$

$$= \Pr \left[\left| \frac{\|F(u) \cdot \|u\|_2 \cdot v\|_2^2}{\|u\|_2^2} - 1 \right| > \varepsilon \right]$$

$$= \Pr \left[\left| \|F(v)\|_2^2 - 1 \right| > \varepsilon \right]$$

$$\leq 2 \exp \left(- \frac{\delta \varepsilon^2}{2} \right) \quad \text{for } \varepsilon \in (0, 1]$$

Step 4. proof of JL lemma

Our goal is to show $\|F(u_i) - F(u_j)\|_2^2$ is within $(1 \pm \varepsilon)$ factor of $\|u_i - u_j\|_2^2$.

More precisely, we show that

$$\Pr \left[\exists i, j : \frac{\|F(u_i) - F(u_j)\|_2^2}{\|u_i - u_j\|_2^2} \notin [1 - \varepsilon, 1 + \varepsilon] \right] \leq \delta$$

proof

For every $i, j \in [n]$, if $u_i = u_j$

the embedding does not change

the Euclidean distance.

Assume $u_i \neq u_j$, and let

$$u_{ij} = u_i - u_j$$

$$\text{Clearly } F(u_{ij}) = M(u_i - u_j) = F(u_i) - F(u_j)$$

$$\Pr \left\{ \frac{\|F(u_i) - F(u_j)\|_2^2}{\|u_i - u_j\|_2^2} \notin [1-\varepsilon, 1+\varepsilon] \right\}$$

$$= \Pr \left[\left| \frac{\|F(u_{ij})\|_2^2}{\|u_{ij}\|_2^2} - 1 \right| > \varepsilon \right]$$

$$\leq 2 \exp \left(- d \frac{\varepsilon^2}{2} \right)$$

↗
by part 3

Using union bound :

$$\Pr \left[\exists i, j : \frac{\|F(u_i) - F(u_j)\|_2^2}{\|u_i - u_j\|_2^2} \notin [1-\varepsilon, 1+\varepsilon] \right]$$

$$\leq \binom{n}{2} 2 \exp \left(- d \frac{\varepsilon^2}{2} \right) \underset{\uparrow}{\leq} \delta$$

by setting $d' = \Theta \left(\frac{\log n / \delta}{\varepsilon^2} \right)$