

CL=CSCI 160

CLASS 5

JUSTIFY THE SUBTRACTION ALGORITHM USING THE r 'th COMPLEMENT NUMBER REPRESENTATION WE GAVE IN CLASS

In our base r representation, where we allocate n digits to the integer part and m digits to the fraction part, we define the r 's complement of a number N as follows:

$$N_{rc} = r^n - N$$

Show that in the above representation, given two non-negative numbers M and N as inputs, the algorithm below computes the value of $M - N$, in the same representation. Compute all steps of the algorithm below in the boxes, as indicated. **Justify each step in its box. Compare $M - N$ to 0 in each of the branches of step 2) below, and show the exact inequality covered by each branch in the places indicated.**

ALGORITHM to compute $M - N$:

- 1) Compute $M + N_{rc}$
- 2) (a) If there is an e.a.c. then ignore it. Stop.
 (b) If there is no e.a.c. then compute:

$$-(M + N_{rc})_{rc}$$

where $(M + N_{rc})$ is already computed at step 1) above. Stop.

$$N_{rc} = r^n - N$$

$$N = \underbrace{\hspace{1.5cm}}_{n} \cdot \underbrace{\hspace{1.5cm}}_{m}$$

integer part fraction part

$n, m = \text{number of locations}$

ALGORITHM to compute $M - N$:

- 1) Compute $M + N_{rc}$
- 2) (a) If there is an e.a.c. then ignore it. Stop.
 (b) If there is no e.a.c. then compute:

$$-(M + N_{rc})_{rc}$$

where $(M + N_{rc})$ is already computed at step 1) above. Stop.

Justification

$$1) \quad M + N_{rc} = M + r^n - N$$

2) e. a. c. has the magnitude r^n . On branch a) there is an e. a. c., which means we have:

$$M + r^n - N \geq r^n \leftrightarrow M - N \geq 0. \text{ Branch a) says: "if there is an e.a.c. ignore it".}$$

This means we have to subtract the value of the e.a.c., r^n , from the value found at 1); the answer is

$$M + r^n - N - r^n = M - N \geq 0$$

Branch b) is taken when there is no e.a.c., and we know that then $M - N < 0$. Following the algorithm

the answer is: $-(M + N_{rc})_{rc} = -(M + r^n - N)_{rc} = -(r^n - (M + r^n - N)) =$

$$= -(r^n - M - r^n + N) = -(N - M) = -|M - N| < 0 \quad \text{since } N - M > 0.$$

Note In the algorithm using r 's complement representation 0 appears in branch (a) as 0, and not in branch (b), as -0 , as it appears in the algorithm using $(r-1)$'s complement representation.

Addendum to Q1.2

Answer this **Addendum to Q1.2** before you answer the other Quiz 1.2 questions.

Suppose $M = N$.

After we apply the two subtraction algorithms we studied:

1. $(r - 1)$'s complement representation, and
2. r 's complement representation,

which sign (+, or -) should the exact result have in each of the above two algorithms?

Encircle the appropriate sign in spaces below:

1. $(r - 1)$'s complement

$$M - N = + \textcircled{-} 0$$

2. r 's complement

$$M - N = \textcircled{+} - 0$$

1.24 Formulate a weighted binary code for the decimal digits, using weights

a) *6, 3, 1, 1

b) 6, 4, 2, 1

Solution

	<u>6 3 1 1</u>	<u>6 4 2 1</u>
<u>0</u>	0 0 0 0	0 0 0 0
<u>1</u>	0 0 0 1 ^(*)	0 0 0 1
<u>2</u>	0 0 1 1	0 0 1 0
<u>3</u>	0 1 0 0	0 0 1 1
<u>4</u>	0 1 1 0 ^(*)	0 1 0 0
<u>5</u>	0 1 1 1	0 1 0 1
<u>6</u>	1 0 0 0	0 1 1 0 ^(*)
<u>7</u>	1 0 1 0 ^(*)	1 0 0 1 ^(*)
<u>8</u>	1 0 1 1	1 0 1 0
<u>9</u>	1 1 0 0	1 0 1 1

Table

Four Different Binary Codes for the Decimal Digits

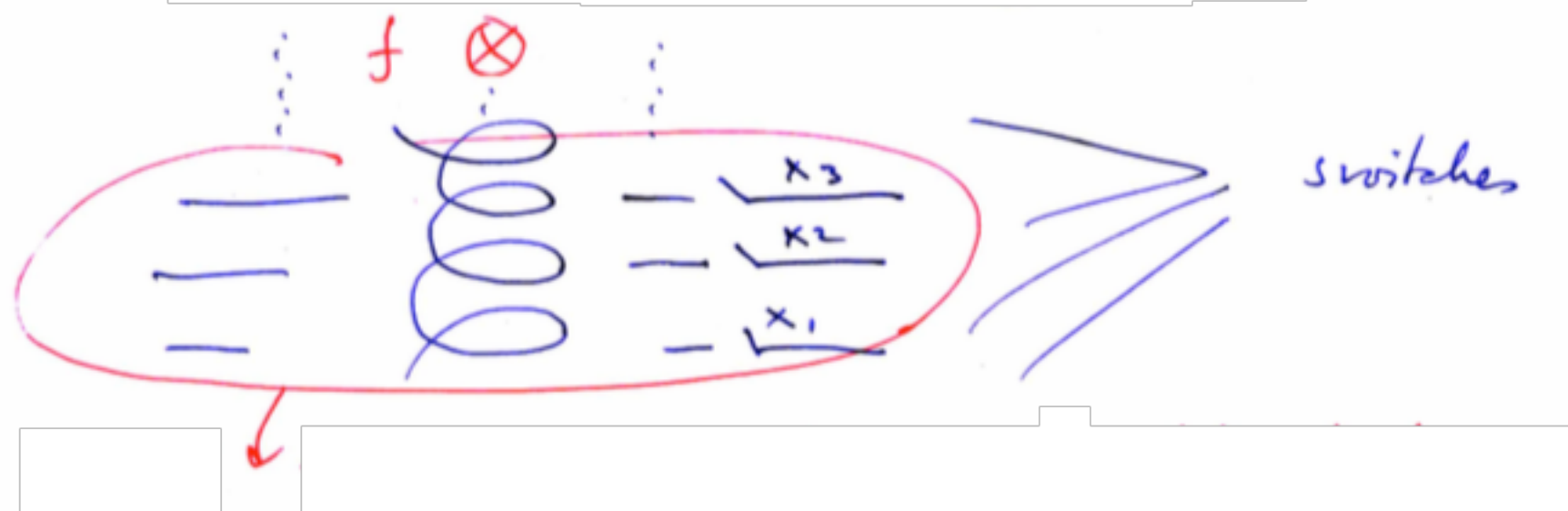
Decimal Digit	BCD 8421	2421	Excess-3	8, 4, -2, -1
0	0000	0000	0011	0000
1	0001	0001	0100	0111
2	0010	0010	0101	0110
3	0011	0011	0110	0101
4	0100	0100	0111	0100
5	0101	1011	1000	1011
6	0110	1100	1001	1010
7	0111	1101	1010	1001
8	1000	1110	1011	1000
9	1001	1111	1100	1111
Unused bit combinations	1010	0101	0000	0001
	1011	0110	0001	0010
	1100	0111	0010	0011
	1101	1000	1101	1100
	1110	1001	1110	1101
	1111	1010	1111	1110

(* Here we have a choice

Note None of these codes is **self-complementary**, i.e. ((r-1)s complements correspond to each other, that is decimal digits complementary in 9's compl. have binary representations that are 1's complements of each other.

Problem of the Independent Switches: example of a Boolean function

There is a staircase leading to n floors, all being lit by one lamp, which is independently operated by any of n switches, which are placed one at each of the n floors.



To simplify, consider $n = 3$; that is there are **3** floors and 3 switches.

Design a function f of the 3 switches. f represents this lamp. Find its truth table.

Say: **on = 1** and **off = 0**

x_1	x_2	x_3	f_1	f_2
0	0	0	0	1
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	0

Note: the value of f for any initial state of the switches is arbitrary, as we can start with any configuration; however once we choose one initial state, the values of f for every other configuration of the switches are determined by this initial state.

f_1 and f_2 are **complementary**

Hamming Distance (d_H) between two strings equals the number of locations in which the two strings differ.

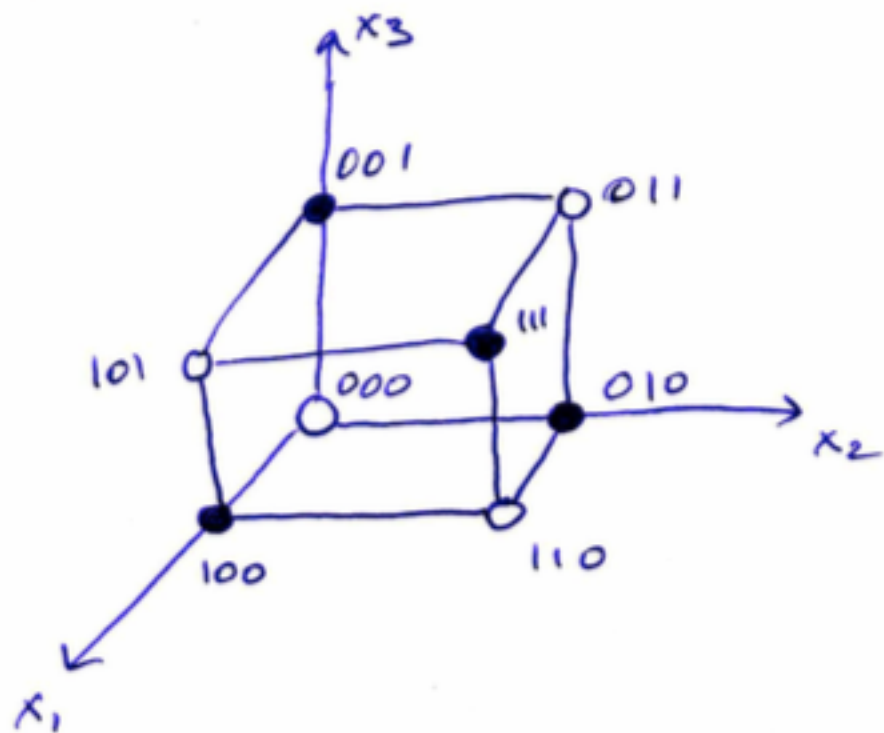
ex: $d_H(1 \underline{0} \underline{0}, 1 \underline{1} \underline{0}) = 1$; $d_H(\underline{0} \underline{0} \underline{1}, \underline{1} \underline{1} \underline{0}) = 3$

x_1	x_2	x_3	f_1	f_2
0	0	0	0	1
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	0

Hamming Distance (d_H) between two strings equals the number of locations in which the two strings differ.

Note: When the H distance between two $x_1 x_2 x_3$ values is odd then the corresponding values for f are different. Conversely, when the H distance is even then the corresponding values for f are the same.

Let's represent the problem of the 3 independent switches using a 3-D cube.



The vertices of the cube have as coordinates all possible inputs for the 3 switches, while the color of each vertex shows the state of the lamp (f).

● = on $\leftrightarrow f = 1$; ○ = off $\leftrightarrow f = 0$

HW: Solve the problem of the independent switches for $n = 4$, which means: 1) Truth table

2) Draw the 4-D cube

Boolean Algebra - Definition

A Boolean algebra is a 6-tuple:
that follows the next 6 axioms:

1) Closure for $+$, \cdot

if $x, y \in B$, then $x \cdot y, x + y \in B, \forall x, y \in B$

2) Identity elements:

$$x + 0 = 0 + x = x$$

$$x \cdot 1 = 1 \cdot x = x, \forall x \in B$$

3) $+$, \cdot Commutative:

$$x + y = y + x$$

$$x \cdot y = y \cdot x, \forall x, y \in B$$

4) $+$, \cdot Distributive: one relative to the other

$$(x \cdot y) + z = (x + z)(y + z)$$

$$(x + y) \cdot z = x \cdot z + y \cdot z \quad \forall x, y, z \in B$$

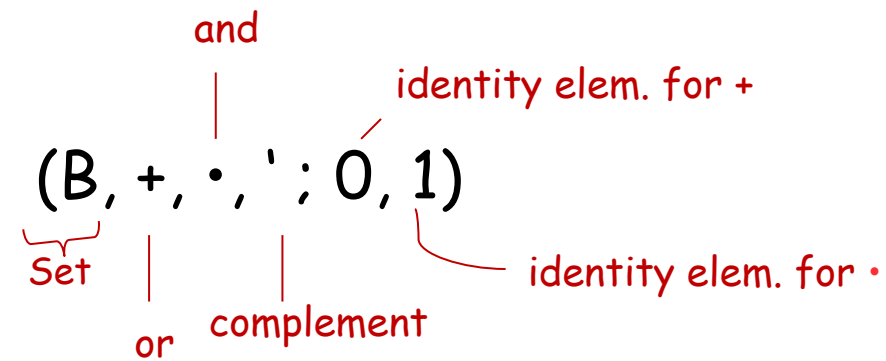
5) $\forall x \in B, \exists x' \text{ or } \overline{x} \in B$ such that:

$$x + x' = 1$$

$$x \cdot x' = 0$$

6) **There are at least 2 elements:**

$$0 \neq 1.$$



HW: Write all axioms and properties for the B. algebra of sets:

$S = \text{set}, S \neq \emptyset$

$(P(S), \cup, \cap, \complement; \emptyset, S).$

Set of all subsets of S

[Remember: if S has n elements then $P(S)$ has 2^n]

Axioms of the Boolean Algebra $(B, +, \cdot, ' ; 0, 1)$

1) $+$, \cdot are closed operations in B:

$$x + y \in B$$

$$x \cdot y \in B, \quad \forall x, y \in B$$

2) identity elements:

$$x + 0 = 0 + x = x$$

$$x \cdot 1 = 1 \cdot x = x, \quad \forall x \in B$$

3) $+$, \cdot are commutative:

$$x + y = y + x$$

$$x \cdot y = y \cdot x, \quad \forall x, y \in B$$

4) $+$, \cdot are distributive one relative to the other one:

$$(x + y) \cdot z = x \cdot z + y \cdot z$$

$$(x \cdot y) + z = (x + z)(y + z), \quad \forall x, y, z \in B$$

5) $\forall x \in B, \exists x' \in B$, the complement of x , such that:

$$x + x' = 1 \text{ and}$$

$$x \cdot x' = 0$$

6) There are at least two elements in B :

$$0 \neq 1$$

Properties of Boolean Algebras (B)

7) $+$, \cdot are idempotent:

$$x + x = x$$

$$x \cdot x = x, \quad \forall x \in B$$

8) $x + 1 = 1,$

$$x \cdot 0 = 0, \quad \forall x \in B$$

9) $x'' = x, \quad \forall x \in B$

10) DeMorgan's Law:

$$(x + y)' = x'y',$$

$$(x \cdot y)' = x' + y', \quad \forall x, y \in B$$

11) $+$, \cdot are associative:

$$x + (y + z) = (x + y) + z,$$

$$x \cdot (y \cdot z) = (x \cdot y) \cdot z, \quad \forall x, y, z \in B$$

12) Absorption for $+$, \cdot :

$$x + xy = x$$

$$x(x + y) = x, \quad \forall x, y \in B$$

Axioms + Properties for Boolean Algebra ($B, +, \cdot, ' ; 0, 1$)

1) $+, \cdot$ are closed in B

2) identity elements:

$$x + 0 = x = 1 \cdot x, \quad \forall x$$

3) $+, \cdot$ are commutative

4) $+, \cdot$ are distributive one relative to the other one:

$$(x + y) \cdot z = x \cdot z + y \cdot z$$

$$(x \cdot y) + z = (x + z)(y + z), \quad \forall x, y, z$$

5) $\forall x \in B, \exists x' \in B$, the complement of x , such that:

$$x + x' = 1 \text{ and } x \cdot x' = 0$$

6) $0 \neq 1$

7) $+, \cdot$ are idempotent:

$$x + x = x \text{ and } x \cdot x = x, \quad \forall x$$

8) $x + 1 = 1$ and $x \cdot 0 = 0, \quad \forall x$

9) $x'' = x, \quad \forall x$

10) DeMorgan's Law:

$$(x + y)' = x'y', \text{ and } (x \cdot y)' = x' + y', \quad \forall x, y$$

11) $+, \cdot$ are associative

12) Absorption for $+, \cdot$:

$$x + xy = x$$

$$x(x + y) = x, \quad \forall x, y$$

HW 3-D The solutions to the quadratic equation

$$x^2 - 11x + 22 = 0$$

are $x = 3$ and $x = 6$.

Determine the base of the numbers in the equation.

Solution

Use sol. $x = 3$: $3^2 - 11_r + 3 + 22_r = 0$ ^{base 10} (\Rightarrow)

$$(\Rightarrow) 9 - (r+1) \cdot 3 + 2r + 2 = 0 \quad (\Rightarrow) r = 8$$

[one could use other solution $r = 6$ and obtain the same].

HW 3-E

Convert the hexadecimal number 68BE to binary, and then convert it from binary to octal and then to base 32.

Solution

digit

$r=32$

15 - F

16 - G

17 - H

18 - I

19 - J

20 - K

21 - L

22 - M

23 - N

24 - O

25 - P

26 - Q

27 - R

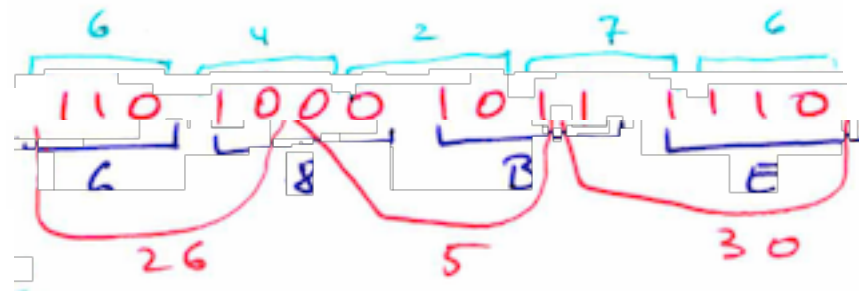
28 - S

29 - T

30 - U

31 - V

68BE_{hex} =



= 64276₈

= Q5U₃₂