

# Canonical Standard Forms

CLASS 10

# Boolean functions - Canonical form vs. standard forms

## Independent Switches Problem

Note:

$$\left. \begin{array}{l} A = 1 \\ \text{and} \\ B = 1 \end{array} \right\} \iff AB = 1$$

$$\left. \begin{array}{l} A = 1 \\ \text{or} \\ B = 1 \end{array} \right\} \iff A + B = 1$$

Let's write the expression of function f, from its truth table.

### Sum of minterms representation:

x	y	z	f
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

$$f = 1 \iff \left\{ \begin{array}{l} x = 0 \text{ and } y = 0 \text{ and } z = 1 \\ \text{or} \\ x = 0 \text{ and } y = 1 \text{ and } z = 0 \\ \text{or} \\ x = 1 \text{ and } y = 0 \text{ and } z = 0 \\ \text{or} \\ x = 1 \text{ and } y = 1 \text{ and } z = 1 \end{array} \right\} \iff \left\{ \begin{array}{l} x' = 1 \text{ and } y' = 1 \text{ and } z = 1 \\ \text{or} \\ x' = 1 \text{ and } y = 1 \text{ and } z' = 1 \\ \text{or} \\ x = 1 \text{ and } y' = 1 \text{ and } z' = 1 \\ \text{or} \\ x = 1 \text{ and } y = 1 \text{ and } z = 1 \end{array} \right\}$$

$$\iff \left\{ \begin{array}{l} x'y'z = 1 \\ \text{or} \\ x'yz' = 1 \\ \text{or} \\ xy'z' = 1 \\ \text{or} \\ xyz = 1 \end{array} \right\} \iff \begin{array}{c} \text{minterms:} \\ x'y'z + x'yz' + xy'z' + xyz = 1 \\ \updownarrow \\ \boxed{f = x'y'z + x'yz' + xy'z' + xyz} \end{array}$$

This is called the canonical sum of products form.

The canonical form is unique for a given function, and it can be used to compare functions/expressions. Most times, the canonical form may be reduced, simplified.

Any sum of products, not necessarily containing only minterms, is called standard.

Similarly, we have a dual form, called the canonical product of sums form; however in this course we will only deal with the sum of minterms representations, as the other is handled similarly.

**Example (Short Cut):** Obtain the canonical form directly from the truth table of an arbitrary function

	x	y	z	f
0	0	0	0	1
1	0	0	1	1
2	0	1	0	0
3	0	1	1	0
4	1	0	0	1
5	1	0	1	1
6	1	1	0	0
7	1	1	1	1

Canonical sum of products form:

$$f = \underline{x'y'z'} + \underline{x'y'z} + \underline{xy'z'} + \underline{xy'z} + xyz \quad \text{also a standard form}$$

$$f = \underline{x'y'} + \underline{xy'} + xyz$$

$$f = y' + xyz$$

$$[= (y' + y)(y' + xz) = y' + xz]$$

$$f = y' + xz$$

} other standard forms for f

Because the canonical sum of products form is unique we express this function above as:

$$f = \sum (0, 1, 4, 5, 7)$$

# Minimization of functions

## Time and space trade-offs

Variety of techniques to obtain gate simplification.

However: simplification depends on the metric we use:

- The number of literals it contains  
= amount of wiring needed to implement the function: # inputs: 3-4 usual,  
> 8,9 very rare.
- The number of gates  
= strong correlation with # components needed for implementation;  
simplest design to manufacture is the one with fewest gates, not literals.
- Number of cascaded levels of gates  
= reducing # logic levels would reduce the overall delay in the path input --> output;  
however an implementation with minimum delay rarely yields an implementation with fewest #  
literals or gates.

Traditional minimization techniques: reduce delay at expense of adding more gates.

Other methods: Trade-off between increased circuit delay and reduced gate count.

# Example

$$F = \underline{a'b'c} + \underline{a'bc} + \underline{ab'c} + abc'$$

$$F_1 = abc' + a'c + b'c$$

- 2-level implementation
- 7 literals (<12 as the original)

$$F_2 = tc' + t'c,$$

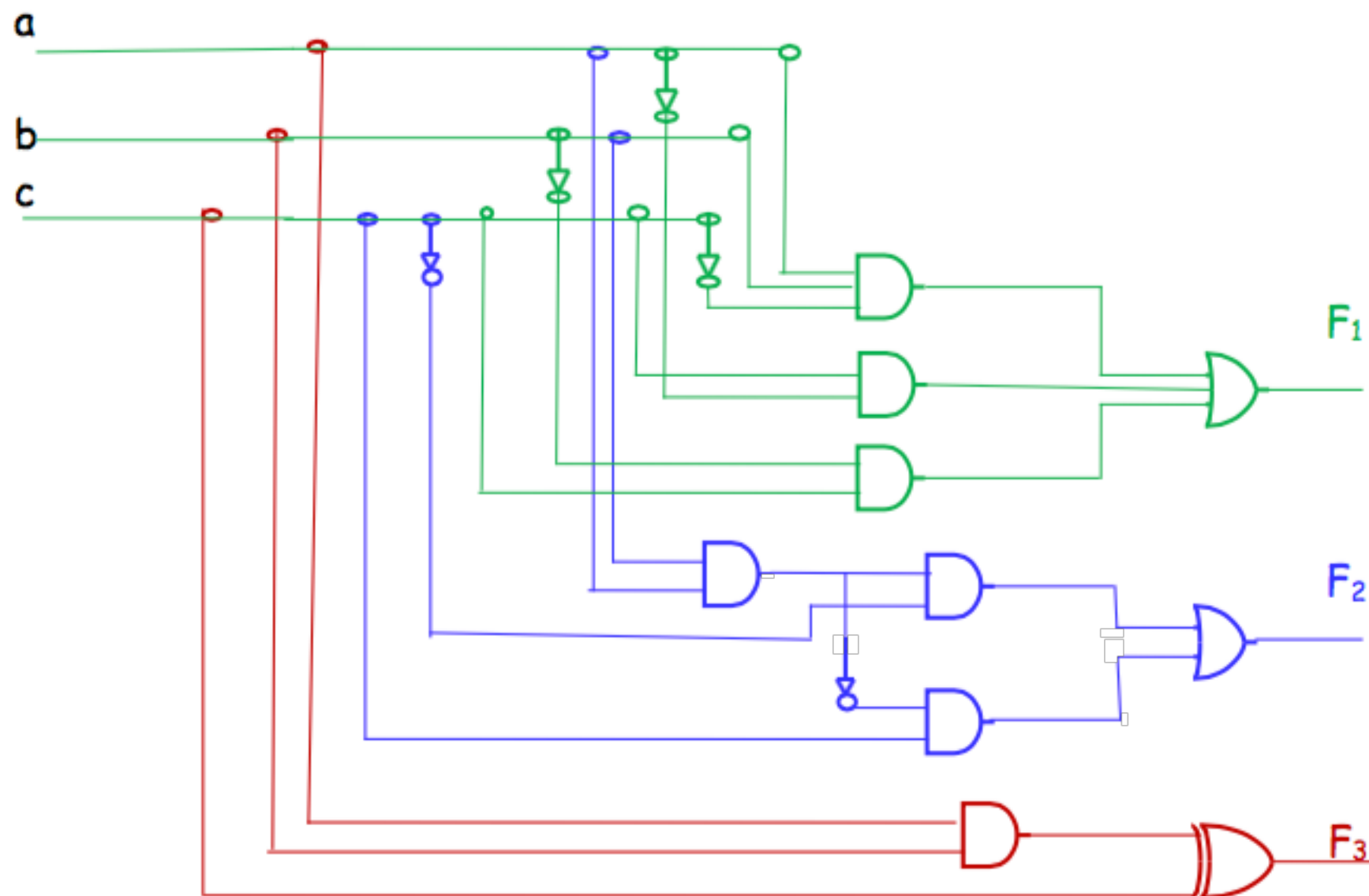
$$\text{set } t = ab$$

- 3-level implementation
- 4 literals
- longest path: 4 gates (> 3 in  $F_1$ )  
--> not as fast as  $F_1$ .
- total # gates in  $F_2 \leq$  in  $F_1$ .

$$F_3 = (ab) \oplus c$$

- XOR = complex gate: implement by combining NAND, NOR
- lowest gate count, but also worst signal delay, for XOR slow compared to simple AND, OR.

## Logical Diagrams (all three functions in one diagram):



## HW 12 - assigned

12-A Given the Boolean functions  $F_1$  and  $F_2$ ,

- (a) Show that the Boolean function  $E = F_1 + F_2$  contains the sum of the minterms of  $F_1$  and  $F_2$ .
- (b) Show that the Boolean function  $G = F_1 F_2$  contains only the minterms that are common to  $F_1$  and  $F_2$ .

12-B Give the truth table of the function:

$$F = xy + xy' + y'z$$