Problem

We can use the Ford-Fulkerson algorithm in almost the same way, except we modify the way that the algorithm augments paths in order to account for node capacity. To do this, we let the bottleneck(P, f) = the minimum amount of flow that can pass through an edge or node along our path. An s-t cut in this node capacitated network is the same – it is a partition of the graph into subgraphs A and B such that $s \in A$ and $t \in B$. The difference is that the capacity of this cut is constrained by the capacity of the nodes in each partition.