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Many Hardware Description Languages (HDL)

- Verilog, System Verilog
- Scala-based: Chisel, *SpinalHDL*
- Python-based: pyMTL, myHDL, pyRTL, migen
- Haskell-based: CλaSH^{*}
- PSHDL^{*}
- Bluespec

Some problems with current HDLs

- DSL artifacts
- HW artifacts
- Not fully synthesizable
- Unable to synthesize objects

HDLs tend to have DSL artifacts

```
// Chisel has == for SCALA, === for Chisel
io.v := y === UInt(0)

// pyRTL has special assignments
a <:= 3 // "assign, generated code"
a = 3   // "assign, in Python"
```

- DSL have faster development but mix two languages in one:

Chisel, CλaSH, myHDL, pyMTL, pyRTL

HDLs tend to have HW artifacts

```
a = 3  
a = a + 1  
assert(a==4); // may fail
```

- In several HDLs, the previous assertion may fail:

Chisel, SpinalHDL, CλaSH, PSHDL

Verilog (non-blocking)

HDLs can be not fully synthesize

```
a = 3  
#3 // Not synthesizable  
a = 4
```

- Some HDLs are not fully synthesizable which adds complexity:

myHDL, Verilog, System Verilog

HDLs can not synthesize objects well

```
# no methods in input/outputs  
a = input.get_value
```

- HDLs with synthesizable objects:

none

Informal poll for lack of adoption of HDLs

- Steep learning curve, language artifacts

Informal poll for lack of adoption of HDLs

- Steep learning curve, language artifacts
- Slow compilation and simulation

Informal poll for lack of adoption of HDLs

- Steep learning curve, language artifacts
- Slow compilation and simulation
- Verilog vs HDL (Most tools handle Verilog not X-HDL)
 - Difficult frequency/power/area feedback
 - Need to understand/debug generated verilog

Pyrope, a modern HDL with a live flow

- Steep learning curve, language artifacts
 - Modern and concise programming language, avoiding hardware specific artifacts
 - Static checks as long as they not produce false positives
 - Synthesis and simulation must be equal and deterministic
- Slow compilation and simulation
 - Live (under 30 secs) simulation, reload, and synthesis feedback goal
- Verilog vs HDL (Most tools handle Verilog not X-HDL)
 - Allows Pyrope 2 Verilog, edit Verilog, Verilog 2 Pyrope, edit Pyrope...

Things that Pyrope can not do

- Generic programming language, Pyrope is synthesizable
- No recursion, neither function nor variable instantiation recursion
- Loops/iterators unbound at compile time
- Any IO, syscall... besides debugging outputs
- rd/wr global variables
- Access objects with pointers. HDLs use hierarchy for references

A Counter with a pipeline stage

Quick Dive to Pyrope

Pyrope

```
# code/counter.prp file
if $enable {
  @total := @total + 1
}
```

A Counter with a pipeline stage

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Pyrope

```
# code/counter.prp file
if $enable {
  @total := @total + 1
}
```

Verilog

```
module s1 (input  clk,
           input  reset,
           input  enable,
           output [3:0] total);
  reg [3:0] total_flop;
  reg [3:0] total_next;

  assign total = total_flop;

  always_comb begin
    total_next = total;
    if (enable)
      total_next = total + 1'b1;
  end

  always @(posedge clk) begin
    if (reset) begin
      total_flop <= 4'b0;
    end else begin
      total_flop <= total_next;
    end
  end
end
endmodule
```

A Counter with a pipeline stage

Quick Dive to Pyrope

Pyrope

```
# code/counter.prp file
if $enable {
    @total := @total + 1
}
```

Pyrope unit test

```
# code/counter_test.prp file
b as counter __stage:true # pipeline type
b.total as __bits:4
b.enable = 0
I b.total == 0          # assertion
yield                  # advance clock
I b.total == 0
b.enable = 1
I b.total == 0
yield                  # advance clock
I b.total == 1
```

Verilog

```
module s1 (input  clk,
           input  reset,
           input  enable,
           output [3:0] total);

reg [3:0] total_flop;
reg [3:0] total_next;

assign total = total_flop;

always_comb begin
    total_next = total;
    if (enable)
        total_next = total + 1'b1;
end

always @(posedge clk) begin
    if (reset) begin
        total_flop <= 4'b0;
    end else begin
        total_flop <= total_next;
    end
end
endmodule
```

A Counter ~~with a pipeline stage~~

Quick Dive to Pyrope

Pyrope

```
# code/counter.prp file
if $enable {
    @total := @total + 1
}
```

Pyrope unit test

```
# code/counter_test.prp file
b as counter          #__stage:true
b.total as __bits:4
b.enable = 0
I b.total == 0
yield
I b.total == 0
b.enable = 1
I b.total == 1        # before it was 0
yield
I b.total == 2        # before it was 1
```

Verilog

```
module s2 (input  clk,
           input  reset,
           input  enable,
           output [3:0] total);

reg [3:0] total_flop;
reg [3:0] total_next;

    assign total = total_next;

always_comb begin
    total_next = total;
    if (enable)
        total_next = total + 1'b1;
end

always @(posedge clk) begin
    if (reset) begin
        total_flop <= 4'b0;
    end else begin
        total_flop <= total_next;
    end
end
endmodule
```


Testbenches

Quick Dive to Pyrope

- Pyrope language testbenches are synthesizable
- Complex tests can interface with C++

Pyrope

```
# code/test1.prp file
mytest = ::{
  puts "Hello World"
  I 1 == 0+1
  yield
  c = 1
  f.a = 2
  f.b = 3

  a = methodx c f
  I a.res == 6 and a.or == 0b11
}
```

C++

```
// mysrc/test2.cpp file
#include "prp_test1.hpp"

void prp_methodx(prp_test1_a &a
                 ,const prp_test1_b b
                 ,const prp_test1_f f) {
  a.res = b + f.a + f.b;
  a.or  = b | f.a | f.b;
}
```

```
$find . -type f
./code/test1.prp
./mysrc/test2.cpp
$prp --run test1.mytest ./mysrc/test2.cpp
```

A Ripple Carry Adder

Quick Dive to Pyrope

```
# libs/adder/code/rca.prp file
fa = :(a b cin):{
    tmp    = $a ^ $b
    %sum    = tmp ^ $cin
    %cout   = (tmp & $cin) | ($a & $b)
}

carry = $cin
for i:(0..a.__bits) {
    tmp = fa a[[i]] b[[i]] carry
    %sum[[i]] = tmp.sum
    carry     = tmp.cout
}
%cout = carry

test2 = ::{
    c = rca a:32 b:4 cin:0
    puts "sum is {0:b} {0}" c.sum
}
```

A Compact Ripple Carry Adder

Quick Dive to Pyrope

```
# libs/adder/code/rca2.prp file
c = $cin
for i:(0..$a.__bits) {
    %sum[[i]] = $a[[i]] ^ $b[[i]] ^ c
    c = ($a[[i]] & $b[[i]] | ($a[[i]] & c) | ($b[[i]] & c)
}

test = ::{
    for a:(1..100) b:(0..33) c:(0 1) {
        d = rca2 a:a b:b cin:c
        I d.sum == (a+b+c)
    }
}
```

```
$find . -type f
./libs/adder/code/rca2.prp
$prp --run libs.adder.rca2.test
```

A Carry Lookahead Adder

Quick Dive to Pyrope

```
# libs/adder/code/cla.prp file
%sum = rca.(a:$a b:$b cin:0).sum

g = $a & $b # Generate
p = $a ^ $b # Propagate

# 4 bit: c = g[[3]] | g[[2]] & p[[3]] | g[[1]] & p[[3]] & p[[2]] | ...
c = $cin & &.(p) # &.(p) is and reduction fcall with p as argument
for i:(0..a.__bits) {
  _tmp = g[[i]]
  for j:(i..(a.__bits-1)) {
    _tmp = _tmp & p[[j]]
  }
  c = c | _tmp
}
%cout = c

test = ::{
  for a:(1..40) b:(1..100) {
    c1 = cla a:a b:b cin:0
    c2 = rca a:a b:b cin:0
    I c1.cout == c2.cout
  }
}
```

Specializing the adders

Quick Dive to Pyrope

```
# libs/adder/code/scla.prp file
cla = :(a b) when a.__bits==8:{           # specialize cla when bits == 8
  s1 = cla a[[0..3]] b[[0..3]] cin:0      # cla for 4 bits
  s2 = cla a[[4..7]] b[[4..7]] cin:s1.cout # pass fast s1.cout as cin
  $sum = (s2.sum s1.sum)[[]]              # bit concatenation
}

cla = :(a b) when a.__bits==12:{          # specialize cla when bits == 12
  s1 = cla a[[0...6]] b[[0...6]] cin:0    # ... vs .. ranges like in Ruby
  s2 = cla a[[6..11]] b[[6..11]] cin:s1.cout
  $sum = (s2.sum s1.sum)[[]]
}

cla = :(a b):{                            # default CLA (not CLA, just RCA)
  $sum = rca.(a b cin:0).sum
}

test = ::{
  s = cla 3 5
  I s.sum == 8
}
```

```
$prp --run libs.adder.scla.test
```

Customizing the counter

Quick Dive to Pyrope

```
# code/counter.prp file
..+.. as __root.libs.adder.scla.cla      # Overload + operator
if $enable {
  @total := @total + 1
  I 3 ..+.. 4 == 7 == 3 + 4              # + is an alias for ..+..
}
```

2 Pipeline stage adder

Quick Dive to Pyrope

```
# code/add4.prp file
..+.. as __root.libs.adder.scla.cla
s1 as __root.libs.adder.rca

(%sum sum1 sum2) as __stage:true

sum1 = $a + $b
sum2 = $c + $c
%sum = s1 a:sum1.sum b:sum2.sum cin:0

test = ::{
  b as add4 a:1 b:2 c:3 d:4
  I b.sum == 0
  yield
  I b.sum == 0
  yield
  I b.sum == 10
}
```

Pyrope vs ...



vs Verilog

Verilog

```
module vsverilog (input clk,
                  input reset,
                  input [2:0] a,
                  input [2:0] b,
                  output reg [3:0] c);

  always @(posedge clk) begin
    if (reset) begin
      c <= 3'b0;
    end else begin
      c <= a + b;
    end
  end
endmodule
```

Pyrope

```
# code/vsverilog.prp file
($a $b) as __bits:2
%c = $a + $b
```

- No inputs/outputs
- Infer bit sizes
- Automatic reset to zero
- No reg/wire
- No blocking/non-blocking

vs Chisel

Chisel

```
import Chisel._
class GCD extends Module {
  val io = new Bundle {
    val a = UInt(INPUT, 16)
    val b = UInt(INPUT, 16)
    val e = Bool(INPUT)
    val z = UInt(OUTPUT, 16)
    val v = Bool(OUTPUT)
  }
  val x = Reg(UInt())
  val y = Reg(UInt())
  when (x > y) { x := x - y }
  unless (x > y) { y := y - x }
  when (io.e) { x := io.a; y := io.b }
  io.z := x
  io.v := y === UInt(0)
}
object Example {
  def main(args: Array[String]): Unit = {
    chiselMain(args, () => Module(new GCD()))
  }
}
```

Pyrope

```
if $a? and $b? {
  (@x @y) = ($a $b)
}else{
  if @x > @y { @x = @x - @y }
  else { @y = @y - @x }
  if @y == 0 { %z = @x }
}
```

```
test = ::{
  gcd as vschisel
  (a b z) as __bits:16
  z = gcd a:a.__rand b:b.__rand
  waitfor z
  puts "gcd for " a " and " b " is " z
}
```

- Global type inference
- No scala vs chisel syntax

vs Bluespec

BSV

```
module mkTb (Empty);
  Reg#(int) cycle <- mkReg (0);

  rule count_cycles;
    cycle <= cycle + 1
    if (cycle > 7) $finish(0);
  endrule

  int x = 10;
  rule r;
    int a = x;
    a = a * a;
    a = a - 5;
    if (pack(cycle)[0] == 0) a = a + 1;
    else a = a + 2;

    if (pack(cycle)[1:0] == 3) a = a + 3;

    for (int k=20;k<24;k=k+1)
      a = a + k;
    $display ("%0d: rule r, a=%0d",cycle,a);
  endrule
endmodule: mkTb
```

Pyrope

```
# code/vsbsv.prp file
@cycles = @cycles + 1
x = 10
a = x
a = a * a
a = a - 5
if cycle[[0]] == 0 { a = a + 1 }
else { a = a + 2 }

if cycle[[0..1]] == 3 { a = a + 3 }

for k:(20..24) {
  a = a + k
}
puts "{}: rule, a={}" cycle a
```

- More compact syntax
- More traditional language, no rules

vs migen (Python HDL)

migen

```
from migen.fhdl.std import *
from migen.fhdl import verilog
class Blinker(Module):
    def __init__(self, led, maxperiod):
        counter = Signal(max=maxperiod+1)
        period = Signal(max=maxperiod+1)
        self.comb += period.eq(maxperiod)
        self.sync += If(counter == 0,
            led.eq(~led),
            counter.eq(period)
        ).Else(
            counter.eq(counter - 1)
        )
led = Signal()
my_blinker = Blinker(led, 3000000)
print(verilog.convert(my_blinker, ios={led}))
```

Pyrope

```
# code/vsmigen.prp file
if @counter {
    @counter -= 1 # @counter-- does not work
}else{
    @counter = $maxperiod
    @led = ~@led # Not %, @ is always valid
}

test = ::{
    b = vsmigen maxperiod:300000
    puts "led is {}" b.led
    yield 300000
    puts "led is {}" b.led
}
```

- Avoid weird DSL syntax

vs pyRTL (Python HDL)

pyRTL

```
def fibonacci(n, req, bitwidth):
    a = pyrtl.Register(bitwidth, 'a')
    b = pyrtl.Register(bitwidth, 'b')
    i = pyrtl.Register(bitwidth, 'i')
    local_n = pyrtl.Register(bitwidth, 'local_n')
    done = pyrtl.WireVector(bitwidth=1, name='done')

    with pyrtl.conditional_assignment:
        with req:
            local_n.next |= n
            i.next |= 0
            a.next |= 0
            b.next |= 1
        with pyrtl.otherwise:
            i.next |= i + 1
            a.next |= b
            b.next |= a + b
    done <= i == local_n
    return a, done
```

Pyrope

```
# code/vs.pyrtl.prp file
(@a @b @i) as __bits:$bitwidth
if $n? { # new request
    (@a @b @i) = (0 0 n)
}else{
    (@a @b @i) = (@b,@a+@b, @i-1)
}
if @i == 0 { %result = @a }
```

```
test = ::{
    seq = (0 1 1 2 3 5 8 13 21 34)
    for n:(0..9) {
        b = vs.pyrtl bitwidth:6 n:n
        waitfor b.result # multiple clocks
        I b.result == seq[n]
    }
}
```

- same issues as chisel

vs PSHDL

PSHDL

```
module de.tuhh.ict.Timing {
    out uint a=1,b=2,c=3,d=4;
    a=b;
    b=c;
    c=d;
    d=5;
    // a == b == c == d == 5
}

module de.tuhh.ict.Timing {
    out register uint a=1,b=2,c=3,d=4;
    a=b;
    b=c;
    c=d;
    d=5;
    // a==2, b==3, c==4, d==5
}
```

Pyrope

```
# code/vspshdl.prp file
# % is the output vector
% = (a:1 b:2 c:3 d:4)
%a = %b
%b = %c
%c = %d
%d = 5
I % == (a:2 b:3 c:4 d:5)
```

- Avoid hardware driven syntax

vs CλaSH

CλaSH

```
upCounter :: Signal Bool -> Signal (Unsigned 8)
upCounter enable = s
  where
    s = register 0 (mux enable (s + 1) s)
```

- Easier to guess hw mapping
- More familiar syntax

Pyrope

```
# code/vsclash.prp file
@upCounter as __bits:8
if $enable {
  @upCounter += 1
}
```

vs Liberty (LXE)

Liberty

```
using corelib;
instance gen:source;
instance hole:sink;
gen.create_data = <<<
  *data = LSE_time_get_cycle(LSE_time_now);
  return LSE_signal_something | LSE_signal_enabled;
>>>;
gen.out ->[int] hole.in;
collector out.resolved on "gen" {
  header = <<<
#include <stdio.h>
  >>>;
  record = <<<
    if (LSE_signal_data_known(status) &&
        !LSE_signal_data_known(prevstatus)) {
      if(LSE_signal_data_present(status)) {
        printf(": %d\n", *datap);
      } else {
        printf(": No data\n");
      }
    }
  }
  >>>;
};
```

Pyrope

```
# code/vsliberty.prp file
gen = ::{
  @data = @data + 1
}
sink = ::{
  if $data? {
    puts ": {}" $data
  }else{
    puts ": No data"
  }
}
s = sink __stage:true
g = gen __stage:true
s.data = g.data
```

- Clean syntax
- No extra verbosity
- Similar handshake idea

vs Dart

dart

```
class Person {
  Person.fromJson(Map data) {
    print('in Person');
  }
}

class Employee extends Person {
  Employee.fromJson(Map data)
    : super.fromJson(data) {
    print('in Employee');
  }
}

main() {
  var emp = new Employee.fromJson({});
}

// Cascade operations
a..field1 = 1
..field2 = 2
```

Pyrope

```
# code/vsdart.prp file
person.fromJson = ::{
  puts "in Person"
}

employee = person
employee.fromJson = ::{
  super $ # Notice, no fromJson
  puts "in Employee"
}

emp = employee.fromJson
// No cascade operations
a.field1 = 1
a.field2 = 2
```

- Prototype inheritance
- No memory (new/delete)

vs Reason

Reason

```
type animal = Dog | Cat | Bird;
let result = switch (isBig, animal) {
| (true, Dog) => 1
| (true, Cat) => 2
| (true, Bird) => 3
| (false, Dog | Cat) => 4
| (false, Bird) => 5
};
```

Pyrope

```
# code/vsreason1.prp file
unique if isBig and animal is Dog {
    result = 1
}elif isBig and animal is Car {
    result = 2
}elif isBig and animal is Bird {
    result = 3
}elif !isBig
    and animal is Bird or animal is Cat {
    result = 4
}elif !isBig and animal is Bird {
    result = 5
}
```

- Pyrope mimics SystemVerilog unique keyword, no case or switch

vs Reason

Reason

```
let increment x => x + 1;
let double    x => x + x;

let eleven = increment (double 5);

let add = fun x y => x + y;
let addFive = add 5;
let eleven = addFive 6;
let twelve = addFive 7;
```

Pyrope

```
# code/vsreason.prp file
increment = :(x):{$x + 1 }
double    = :(x):{$x + $x}

eleven = increment.(double.(5))

add = :(x y):{$x + $y}
addFive = \add # add reference, no call
addFive = ::{ super x:5 y:$y }
eleven  = \addFive y:6
twelve  = \addFive y:7
```

- Pyrope has primitive currying

vs Python

Python

```
class objectTest():
    def __init__(self,a):
        self.value = a
    def get_value(self):
        return self.value

a = objectTest(1)
b = objectTest(1)

assert a != b
assert a.get_value() != b.get_value
assert a.get_value() == b.get_value()
assert a.get_value != b.get_value

total = [x*x for x in range(10) if x % 2]
```

Pyrope

```
# code/vspython.prp file
objectTest.get_value = ::{
    return __parent.myvalue
}
objectTest.set_value = :(a){
    __parent.myvalue = $a
    return __parent
}

a = objecttest.set_value 1
b = objecttest.set_value 1

I a == b == 1
I a.get_value.() == b.get_value
I a.get_value.() == b.get_value.()
I a.get_value == b.get_value
I a.__obj == b.__obj and a.__obj != 1.__obj

total = 0..10 |> filter ::{$ & 1} |> map ::{$*$}
I total == (1 9 25 49 81)
```

vs Javascript

Javascript Tricky

```
const a = {
  num: 0,
  valueOf: function() {
    return this.num += 1
  }
};
const equality = (a==1 && a==2 && a==3);
console.log(equality); // true

for(let pair of myMap) {
  var [key, value] = pair;
  console.log(key + " = " + value);
}
```

Pyrope

```
# code/vsjs1.prp file
a = 0
a.__read = ::{
  __parent += 1
}
equality = a == 1 and a == 2 and a == 3
I equality

for a:myMap ::{
  puts "{} = {}" a.__index a
}
```

- Some similarities in functionality

vs MATLAB

MATLAB

```
x = 1:10
y = 10:-2:0
A = [1 2; 3 4] # matrix 2x2

sum = 0;
for i=1:length(x)
    sum = sum + abs(x(i));
end

x3=(1:3).*2;
A = [1 0 3];
B = [2 3 7];
C = A.*B
% C = 2 0 21
C = A * B
% C = [[2 0 6] [3 0 9] [7 0 21]]
```

Pyrope

```
x = 1..10
y = 10..0 ..by.. 2 # (10 8 6 4 2 0)
A = ((1 2) (3 4))

sum = 0
for i:(1..x.__length) {
    sum = sum + abs.(x.(i))
}

x3=(1..3) ** 2 # compile error
I (2 4 6) == (1..3) * 2
A = (1 0 3)
B = (2 3 7)
C = A ** B # OK, matching sizes
I C == (2 0 21)
D = A * B
I C == ((2 0 6) (3 0 9) (7 0 21))
```

- Share tuple vs element operators
- Different applications/goals/...

vs Coffeescript

Coffeescript

```
square = (x) -> x * x
eat     = (x) -> alert square x

eat x for x in [1, 2, 3] when x isnt 2

r361 = square 3 + square 4
r25  = square(3) + square 4
# r361 == 361 and r25 == 25

# Minimum number of parenthesis
y = pow 10, floor log10 x
# Equivalent to
y = pow(10, floor(log10(x)))
```

Pyrope

```
square = :(x):{$ * $}
eat     = :(x):{puts square.($)}

for food:(1 2 3) {
  if food !=2 { eat food }
}

r=square.(3 + square.(4))# 361
r=square.(3) + square.(4)# 25

# Minimum number of parenthesis
y = pow.(10 floor.(log10.(x)))
# Simpler syntax with pipes
y = log2 x |> floor |> pow 10
```

- No iterators after statement
- Different rules about arguments

Pyrope Syntax



Basic Control Flow

ifs

```
# code/controlflow1.prp

if cond1 {
  I cond1
}elseif cond2 {
  I !cond1 and cond2
}

unique if cond3 {
  I cond3 and !cond4
}elseif cond4 {
  I !cond3 and cond4
}else{
  I !cond3 and !cond4
}

unique if cond5 {
  I cond5 and !cond6
}elseif cond6 {
  I !cond5 and cond6
}

I cond5 or cond6 # Unique implies full too
```

iterators

```
# code/controlflow2.prp
total = 0
for a:(1..3) { total += a }
I total == (1+2+3)

total = 0 # compact double nested loop
for a:(1..3) b:(1 2) { total += a }
I total == (1+2+3 + 1+2+3)

# Powerful library. Simple reduce example
reduce = ::{
  t = $0
  for a:$(1..) {
    t = $.__block t a
  }
  return t
}

a = (1 2 3) |> reduce ::{ $0+$1 }
I a == (1+2+3)
```

Element vs Tuple operator

Basic ops

```
# code/elementvstuple1.prp
# operators read left and right side

a = (1 2 3) + 4      # element op
I a == (5 6 7)

a = (1 2 3) ++ (4 5) # tuple concat
I a == (1 2 3 4 5)

a = (1 2 3) ..+.. 4  # element op
I a == (5 6 7)

# ..XX.. means operator XX
# .. are optional for non-alphanumeric operators
```

custom operators

```
# code/elementvstuple2.prp
..dox.. = :(a,b){ #.. .. is optional
  t = ()
  for a:$0 b:$1 ::{
    t += a+b
  }

  return t
}
I (1 3) ..dox.. (2 1) == (3 2 5 4)

sub1 = ::{
  t = ()
  b = $1[0] # first element in rhs
  for a:$0 {t += a+b}
  return t
}
# .. required in call to be operator
I (3 2) ..sub1.. 1 == (2 1)
I (3 2) ..sub1.. (2 3) == (1 0)
```

Operator precedence

- Unary operators (!,~,@,?,%...) bind stronger than binary operators (+,++,-,*,...)
- **Only** six levels of operator precedence (16 levels in c++)
- Always left-to-right evaluation

Priority	Category	Main operators in category
1	Unary	not ! ~ @ ? % \$
2	Mult/Div	*, /
3	bitwise ops	^, &
4	other bin	+, ++, --, <<, >>, >>>, <<<
5	comparators	<, <=, ==, !=, >=, >
6	logical	and, or

```
# code/precedence1.prp
# Typically expected results
I (true or !false) == (true or (!false))
I (3*5+5) == ((3*5) + 5)

a = true or false==false
b = true or (false==false)
I a == b

c = fcall 1 2
I c == fcall.(1 2)

#bar = true or false and true # compile error
#x = 3 ++ 4 -- 3                # compile error

c = a == 3 == b                # OK
I c == (a==3 and 3==b)
```

Operator precedence

explicit newline

```
# code/precedence2.prp
bar = x == 3
    or x == 3 and !(x!=3)
    or false

bar = false or
    true          # compile error, ops after newline

I (true or false==false) == (true or (false==false))

d = 1
    ,3
d = 1,
    3          # compile error, ',', after newline

bar = 3
    * 1 + 4
    * 3 - 1
I bar == 3 * (1+4) * (3-1)
```

explicit;

```
# code/precedence3.prp
x = 3 ++ 4 -- 3      # compile error, precedence?
x = 3 ; ++ 4 -- 3    # OK, 3 ++ (4 -- 3)

b = !a or d          # OK, ! higher precedence
b = !a or c == d     # OK, == breaks expression
I b == !a or (c == d)

bar = true or false and true # compile error
bar = true ; or true and false ; or true
I bar == true or (true and false) or true

I (1,3) == (1 3)
d = 1 3
I d == 1 3
I d == ;( ; 1 ;, 2;+1)  # Ugly but legal syntax
f = 1 +;3              # Ugly illegal syntax
```

Single line syntax

```
# code/singleline.prp
if true { x = 3 }          # OK
if true {
x = 3 }                    # OK
#if true
#{ x = 3 }                 # parse error, no newline

if true ::{ puts __parent.x} # new scope in block
if true { puts x }         # OK too

if true ::{ a = 3 ; puts a }

# parse error, no space between :: {
#if true :: {puts false}

c = 0
d = 0
if false ::{ c = 1 ; d = 2 }
I d == 0 and c == 0        # :: is a new scope

for a:(1..3) {puts a}
I a == 3                   # compile error

# ; is same as a newline
```

Code blocks

```
# code/codeblock.prp file
each as ::{
  I $__block is def
  for a:$ { $__block a }
}

each.(1 2 3)    ::{ puts $ }
(1 2 3) |> each ::{ puts $ }

map as ::{
  t = ()
  fun = $__block
  for a:$ {
    t += fun a
  }
  return t
}

a = ::{ 2+1 } # OK implicit return

# parse error, only last can be implicit return
#a = ::{ 1+1 ; 2+1 }

s = (1 2 3) |> map ::{ $+1 } |> map ::{ $*$ }
I s == (4 9 16)
```

Code blocks

```
# code/codeblock.prp file
each as ::{
  I $.__block is def
  for a:$ { $.__block a }
}

each.(1 2 3)    ::{ puts $ }
(1 2 3) |> each ::{ puts $ }

map as ::{
  t = ()
  fun = $.__block
  for a:$ {
    t += fun a
  }
  return t
}

a = ::{ 2+1 } # OK implicit return

# parse error, only last can be implicit return
#a = ::{ 1+1 ; 2+1 }

s = (1 2 3) |> map ::{ $+1 } |> map ::{ $*$ }
I s == (4 9 16)
```

```
# code/reduce.prp file
reduce = ::{
  if $.__size <= 1 { return $ }

  redop = \$.__block # code block reference
  tmp = $

  while true {
    tmp2 = ()
    for i:(0..tmp.__size by 2) {
      tmp2 += redop tmp[i] tmp[i+1]
    }
    if tmp2.__size <= 1 { return tmp2 }
    tmp = tmp2
    if tmp2.__size[[0]] { # odd number
      tmp = tmp2[..-2] # all but last two
      tmp += redop tmp2[-2..] # reduce last two
    }
  }
  I false
}

a = (1 2 3) |> reduce ::{ $0 + $1 }
I a == 6
```

Variable scope

Method constructs

```
# code/scope1.prp
a = 1
b = ::{
    d = 3    # b local scope
    %out = a # compile error
}
x = b
I a == 1
c = ::{
    a = 2    # local variable
    d = 4
    %out = a
}
I d==4      # compile error, d not defined
I c.out == 2
```


Variable scope

Method constructs

```
# code/scope1.prp
a = 1
b = ::{
  d = 3      # b local scope
  %out = a # compile error
}
x = b
I a == 1
c = ::{
  a = 2      # local variable
  d = 4
  %out = a
}
I d==4      # compile error, d not defined
I c.out == 2
```

Control flow constructs

```
# code/scope2.prp
a = 1
if a == 1 {
  a = 2
  b = 3
  _f = 4
}
I a == 2 and b == 3
I _f == 4 # compile error, undefined

total = 0 # needed, because read in loop
for _i:(1..3) { total += _i }

I total == 1+2+3
I _i == 3 # compile error, undefined

@val = 3
# root is always relative to the current file
I __root.scope2.val == 3
```

Scope outside code regions

Accessing outside scope

```
# code/scope3.prp
a = 1
if a == 1 ::{
  a = 2          # compile error
  __parent.a = 3 # OK
  __root.scope3.a = 5 # OK, same
  f = 3          # local scope
}
I f == 3          # compile error, undefined
I a == 5

b = ::{
  % = __parent.a
}

c.a = 5
c = b.()
b.a = 3
I c == 5
d = b.()
I d == 3
```

Code regions in ifs/fors

```
# code/scope4.prp
t = 0
for a:(1..3) { t += a;; }
I t == 1+2+3

t = 0
for a:(1..3) ::{t = a} # local scope
I t == 0

if t==0 ::{__parent.x = 3}
I x == 3
```

Implicit vs Explicit arguments

```
# code/impvsexp1.prp file
a = (1,2+3,3)           # tuple
a = f.(1,2,f2.(3))      # function call
b = (1
    ,2-3)               # 2 lines
a = f.(1
    ,2-4*fcall.(3-1))  # 2 lines function call
```

- Commas can be avoided if the elements are single line and have no expressions.

```
# code/impvsexp2.prp file
a = (1 2 3)
a = f.(1 2 3)
b = (1+23*fcall.(2+4))
```

- In function calls, when commas can be avoided, parenthesis are optional after a newline, an assignment, or a pipe operator.

```
# code/impvsexp3.prp file
a = (1 2 3)           # required, tuple
f 1 2 3               # after newline

a = 3 |> f 2 3 |> f 1  # after pipe
if f.(2 1 3) {        # must be explicit
    I true
}
I (1 2 3) == (1 2 3)  # must be explicit
```

Function call arguments

```
# code/fcalls.prp file
square = :(x):{$ * $}
#r=square 3 + square 4      # parse error, complex argument
#r=square(3 + square.(4))   # parse error, space required for arguments
#r=square (3 + square (4))  # parse error, missing explicit argument
r=square square 4          # compile error, square has 1 argument, 2 passed
r=square (3 + (square 4))   # compile error, two args, but first reqs argument
r=square (3 + square.(4))   # OK, 361 = (3+4^2)^2 ; ^ is exp, not power
r=square.(3 + square.(4))   # OK, 361
r=square.(3) + square.(4)   # OK, 25
pass = ::{
  if $__size == 1 { return 7 }
  if $__size == 2 { return 9 }
  11
}
puts 3 square 4 5           # compile error, missing required square arg
puts 3 square.(4) 5         # OK, prints "3 16 5"
puts 3 pass 4 5             # OK, prints "3 11 5"
puts 3 pass.(4) 5           # OK, prints "3 7 5"
```

Tuples

Basic tuples

```
# code/tuples1.prp
a = (b:1 c:2) # ordered, named
I a.b == 1 and a.c == 2
I a.0 == 1 and a.2 == 2

b = (3, false) # ordered, unnamed
I b.0 == 3 b[1] == false

c1 as (__bits:1, __bits:3) # final ordered unnamed
c as c1
c as (b:, c:) # final ordered named
c = (true 2)
c = (false 33) # compile error
c.bar = 3 # compile error

d as (a:3 5) # final, ordered, unnamed
I d.a == 3 and d[1] == 5

g = (1 2 3)
I (1 3) ..in.. g
g += (2 5)
```

Complex tuples

```
e.0 = 3 # unnamed, ordered
I e.0 == 3 and e == 3

s as __set:true
s = (1 2 3 3)
I s == (1 2 3)
s = s ++ 4 # add to tuple
s = s ++ (1 4 5)
I s == (1 2 3 4 5)

x = __size:32
x[(1 3)] = (3 1)
x[(1 2)] = (1 1)
I x[(1 3)] == (3 1)
I x[(0b1 0b11)][0] = 3
I x[0b1_11][1] = 1
```

Memories

Clear SRAMs

```
# code/mem1.prp
@a as __bits:3 __size:1024 __rdports:1
@b as @a __fwd:false # without cycle forwarding
@cycle as __bits:8

I @a[0] == @cycle

prev_val = @cycle
@cycle += 1
@a[0] = @cycle
@b[0] = @cycle

I @a[0] == @cycle
I @a[0].__flop == prev_val
I @b[0] == @b[0].__flop == prev_val

%out = @a[0] + @b.0
```

- Memory forward unless __flop used
- Reset to zero by default
- Enforces the rd/wr ports if indicated
- Moves logic to get addresses at posedge

Memories

Enforce SRAM constraints

```
# code/mem2.prp
# Enforce #rd and wr ports in SRAM
@a as __bits:8 __size:1024 __rdports:1 __wrports:1
@cycle as __bits:8

@cycle += 13

# ADDR must be stable at posedge. Push logic
@a[@cycle] = @cycle-1

%out = @a[~@cycle]
```

Becomes

```
# code/mem3.prp
# Enforce #rd and wr ports in SRAM
@a as __bits:8 __size:1024 __wrports:1
@cycle      as __bits:8

@cycle += 13

@_cycle_m1 = @cycle + 13 - 1
@_cycle_p13 = @cycle + 13
@_cycle_neg = ~(@cycle + 13)

@a[@_cycle_p13] == @_cycle_m1

%out = @a[@_cycle_neg]
```

Memories

Enforce SRAM constraints

```
# code/mem4.prp
# Enforce #rd and wr ports in SRAM
@a as __bits:8 __size:1024 __rdports:1 __wrports:1
@a as __posedge:false # posedge by default
@cycle as __bits:8

@cycle += 13

# SRAM can use pos/neg edge
@a[@cycle] = @cycle-1

%out = @a[~@cycle]
```


Ranges

Basic

```
# code/ranges1.prp
I (1 2 3) == 1...4 == 1..3

I (0..7 ..by.. 2) == (0 2 4 6)
I 0..15 ..by.. (2 3) == (0 2 5 7 10 12 15)

I (1..2) ..union.. 3 == (1..3)
I (1..10) ..intersect.. (2..20) == (2..10)

# Ranges can be open
I (3..) ..intersect.. (1..5) == (3..5)
I (..) ..intersect.. (1..2) == (1..2)
I (..4) ..union.. (2..3) == (..4)
I (2..) == (2..-1)
I (..3) == (-1..3)

# Ranges can be converted to values
# I (1..3)[[]] # compile error
```

Complex

```
# code/ranges2.prp
numbers = (1...10)
start = numbers[0..2]
middle = numbers[3...-2]
end = numbers[-2..]
copy = numbers[..]

I start == (1 2 3)
I middle == (4 5 6 7)
I end == (8 9)
I copy == (1 2 3 4 5 6 7 8 9)

val = 0b00_01_10_11

I val[[0..1]] == 0b11
I val[[..-2]] == 0b01_10_11
I val[[-2..]] == 0b00
I val[[-1]] == 0b1 # MSB

I (1..3) * 2 = (2 4 6)
I (1..3) + 2 == (3..5)
I (1 2 4) ++ 3 == (1..4)
```

Random number generation

rnd and rnd_bias interface. Seed controller by environment variable.

```
# code/rndtest.prp
a = __rnd 1..3          # rnd between 1 2 3
b as __bits:12
a = b.__rnd             # rnd from 0 to 4095
b.__rnd_bias = (1 0)    # weight 1 for value 0
b.__rnd_bias += (2 3)   # weight 2 for value 3
b.__rnd_bias += (2 4)   # weight 2 for value 4
b.__rnd_bias += (5 9)   # 0 10%, 3 20%, 4 20%, and 9 50% chance

c as __bits:8
c.__rnd_bias = (1 0)    # weight 1 for value 0
c.__rnd_bias += (2 255) # weight 2 for value 255
c.__rnd_bias += (7 1..254) # weight 7 for the rest
puts c.__rnd           # 10% chance 0, 20% chance 255, 70% other
```

```
$export PRP_RND_SEED=33
$prp --run rndtest
```

Resets

```
# code/reset1.prp
@a as __bits:3 __init:13

@b as __bits:3 __reset:false # disable reset

@mem0 as __bits:4 __init:3 __size:16

@mem1 as __bits:4 __reset:false __size:16

@mem2 as __bits:2 __size:32

# custom reset
@mem2.__init = ::{
  # Called during reset or after clear (!!)
```

```
  @_reset_pos as __bits:log2(__parent.__size) __reset:false
  __parent[@_reset_pos] = @_reset_pos
  @_reset_pos += 1
}
```

Constants

```
# code/constants.prp

a = 3                # implicit __bits:2 __sign:false
a = 3u               # explicit __sign:false, implicit __bits:2
a = 3u4bits          # explicit __sign:false, __bits:4

b = 0xFF_f__fs32bits # explicit __bits:32 __sign:true

c = 0b_111_010_111u32bits
c = 0b_111_010_111u2bits # compile error

c = 0xFF[[0..2]]      # explicit drop bits
```

Compile time assertions and checks

```
# code/assertions.prp

a = 3
C b = a # b and a must be known at Compile time

C if a==3 { # compile time if condition
  I true # runtime assertion
  %d = $0+a # no constant
}

C I a == b # Compile time assertion
I %d != a # runtime assertion
```

Bit precision

Explicit vs implicit

```
# code/precision1.prp
a = 3          # implicit, __range:3u2bits
a = a + 1      # OK

b = 3u2bits    # explicit, __bits:2 __range:3u2bits
b = b - 1      # OK, __range:2u2bits
b = b + 2      # compile error, __bits explicit 2
I b == 2
b := b + 2     # OK (drop bits)
I b == 0       # 4u2bits -> 0b100[[0..1]] == 0

# implicit unless all values explicit
c = 3 - 1u1bits # implicit, __bits:2 __range:2u2bits

@d as __range:(0 1 7) # allowed values
@d = 1            # OK
@d += 1
@d += 1          # compile error

I 0b11_1100 == (a 0b1100)[[]] # bit concatenation
```

Conditions

```
# code/precision2.prp
a as __range:(1..6)
a = 5
c = 5
if xx {
    a = a + 1 # OK
    c = c + 1
}else{
    a = a - 4 # OK
    c = c - 4
}
a = a + 1 # compile error, may be out range
I c.__range == (1,6) # all possible values
c = c + 2
I c.__range == (3,8) and c.__bits == 4
c = c ^ (c>>1) # Not predictable
I c.__range == (0..15) and c.__bits == 4
c = 300 # OK because c was explicit

d = 50u2bits # compile error
e = 3u2bits
e := 50      # OK, drop upper bits
e = e - 1
```

Fluid

Fluid syntax

```
# $i? = false # do not consume
# $i? = true  # consume
# $i! = true  # trigger retry to input
# $i! = false # do not retry input, consume if valid
# $i?        # is valid set?
# $i!        # is retry set?
# $i!!       # is clear set?
# $i!! = true # clear flop
# $i!! = false # do not clear flop

# %o? = false # do not generate output
# %o!        # is retry set?
# %o! = true  # compile error
# %o! = false # compile error
# %o?        # is valid set?
# %o!        # is retry set?
# %o!!       # is clear set?
# %o!! = true # clear flop
# %o!! = false # do not clear flop

# yield      # stop and start from here cycle
# waitfor    # block execution until input is ready
```

Dealing with valids

```
# code/fluid1.prp file
a as $c          # alias, no restart
try ::{
    I __parent.a == $a
    if __parent.a == 3 { %sum = __parent.a }
}
try ::{
    if %sum2! {
        %sum3 = $a # sum2 busy, try sum3
    }else{
        %sum2 = $a
    }
}
try ::{
    if $a? {
        $d? = false # do not consume b
        $e!! = true # clear input e
    }
}
```

Fluid Impacts

```
# code/fluid2.prp file
if a? and a.counter>0 { # Option 1
    @total += a.counter
}
try ::{ # Option 2 (same behavior)
    if a.counter>0 {
        @total += a.counter
    }
}
if a?.counter>0 { # Option 3 (same)
    @total += a.counter
}
@total += a?.counter # Option 4 (same)
```

```
# code/fluid3.prp file
puts "prints every cycle"
try ::{
    puts "odd cycles"
    yield # Yield applies to scope ::{}
    puts "even cycles"
}
puts "prints every cycle"
```

```
# code/fluid4.prp file
everyother = ::{
    if @conta {
        yield
    }
    @conta = ~@conta
    return 1
}

@total_all += 1
@total_yield += everyother.()
I @total_all == @total_yield
try ::{
    @total2_all += 1
}
try ::{
    @total2_yield += everyother.()
    I @total2_all == 2 then @total2_yield == 1
}
```


Fluid Restarts

```
# code/fluid5.prp file
%o1 = $in1?.0      # pass input, no restart

%o2 = $in2?.field  # pass field, no restart

%o3 = $in3?.a + 30 # use field, no restart

try ::{
  %o3 = $in3.a + 30 # same as last
}
```

```
# code/fluid6.prp file
```

Fluid Instantiation

Non Fluid Examples

```
# code/fluid7.prp file
sadd = ::{ %sum = $a + $b }
sinc = ::{ % = $ + 1 }
combinational = ::{
    % = ssum.(a:sinc.($a), b:sinc.($b))
}

one_stage_flop_out = ::{ # The output is flopped
    % = ssum.(a:sinc.($a), b:sinc.($b))
    % as __stage:true
}

one_stage_comb_out = ::{ # Not flopped output
    a1 as sinc
    a2 as ssum __stage:true
    % = a2.(a:a1.($a), b:a1.($b))
}

two_stage_comb_out = ::{ # Not flopped output
    a1 as sinc __stage:true
    a2 as ssum __stage:true
    % = a2.(a:a1.($a), b:a1.($b))
}
```

Fluid Examples

```
# code/fluid8.prp file

combinational = ::{
    % = ssum.(a:sinc.($a), b:sinc.($b))
}
incsum = combinational.(a:$a,b:$b)
incsum as __fluid:true # instance is fluid

one_stage_fluid = ::{ # Same as incsum
    % = ssum.(a:sinc.($a), b:sinc.($b))
    % as __fluid:true
}

mixed_weird_fluid = ::{
    %out1 = a2.(a:a1.($a), b:a1.($b))
    %out2 = a2.(a:$a b:$b)
    %out2 as __fluid:true
}

allfluid = mixed_weird_fluid
allfuild as __fluid:true # both out1 and out2
```

Connecting Stages

Traditional

```
sadd = ::{ %sum = $a + $b }
sinc = ::{ % = $ + 1 }

opt1_2stages = ::{
  s1_a = sinc.$a
  s1_b = sinc.$b
  s1_a as __stage:true
  s1_b as __stage:true
  % = sadd.(a:s1_a b:s1_b)
  % as __stage:true
}

opt2_2stages = ::{
  s1_a = sinc.$a
  s1_b = sinc.$b
  % = sadd.(a:s1_a b:s1_b)

  (s1_a s1_b %) as __stage:true
}
```

More Compact

```
opt3_2stages = ::{
  s1.a = sinc.$a
  s1.b = sinc.$b
  % = sadd.(a:s1.a b:s1.b)

  (s1 %) as __stage:true
}

opt4_2stages = ::{
  s1 = (a:sinc.$a, b:sinc.$b) __stage:true
  % = sadd.(s1) __stage:true
}
```

Assignments, as vs =

Base syntax

```
# code/assign1.prp
a = b          # potential restart in fluid
a as c         # alias a as c, no real read for fluid

b = __bits:3   # explicit bits
b = 3          # OK
b = __bits:10  # OK to redefine
b = 100        # OK

d as __bits:3  # explicit bits
d = __bits:4   # compile error, fixed with as
```

Fluid and assignments, as vs =

both out1 and out2 happens or nothing happens

```
# code/assign2.prp
_tmp1 = $a # read that can trigger restart
_tmp2 = $b
try ::{
    %out1 = _tmp1 + 1 # guarantee no restart (reread)
}
try ::{
    %out2 = _tmp2 + 1
}
```

```
# code/assign3.prp
try ::{
    %out1 = $a + 1
    %out2 = $b + 1
}
```

out1 and out2 can happen independently

```
# code/assign4.prp
try ::{
    %out1 = $a + 1
}
try ::{
    %out2 = $b + 1
}
```

```
# code/assign5.prp
_tmp1 as $a # alias, no restart trigger
_tmp2 = \ $b # pass reference, no restart
try ::{
    %out1 = _tmp1 + 1 # can trigger restart
}
try ::{
    %out2 = _tmp2 + 1 # can trigger restart
}
```

Objects

prototype inheritance

```
# code/objects1.prp
obj1.oneortwo = ::{return 1}
obj2.oneortwo = ::{return 2}
obj1.oneortwo2 = 1
obj2.oneortwo2 = 2

if $input[[0..1]] == 0 {
  tmp = obj1
  I tmp.oneortwo == 1
  I tmp.oneortwo2 == 1
}elseif $input[[0..1]] == 1 {
  tmp = obj2
  I tmp.oneortwo == 2
  I tmp.oneortwo2 == 2
}else{
  # Calling undefined method is __init value
  # NEVER runtime error
  I tmp.oneortwo == 0
  I tmp.oneortwo2 == 0
}

I tmp.oneortwo ..in.. (1 2 0)
```

overload

```
# code/objects1.prp
parent.dox = ::{return 1+$0}

child = parent # inherit
I child.__obj == parent.__obj
child.dox = ::{
  _tmp = super $
  __parent.val = 3 # new field in child
  return tmp + 3
}
I child.__obj != parent.__obj

I child.val == 0
t = child.dox 4
I t == (1+4+3)
I child.val == 3

grandson = child
grandson.dox = :when $0>20:{100}
t = grandson.dox 4
I t == (1+4+3)
t = grandson.dox 30
I t == 100
```

Objects 2

dealing with objects

```
# code/objects1.prp
obj1.foo as __bits:3
obj2.bar as __bits:2
I obj1 isnt obj2

obj1c = obj1
obj1.foo = 1
obj1c.foo = 3
I obj1 is obj1c

obj3 as obj1 or obj2 # Union type
if 3.__rnd == 2 {
  obj3 = obj1
  obj3.foo = 1
}else{
  obj3 = obj2
  obj3.bar = 2
}

if obj3 is obj1 {
  I obj3.foo == 1
}
```

Matching

binary matching

```
# code/objects1.prp
a = 0x73
I a == 0b111011
I a == 0b?11?11

c as __bits:4
I c.popcount <= 1 # Only 1 bit can be set

unique if c == 0b???1 { # ? only in binaries
    onehot = 1
}elif c == 0b??1? {
    onehot = 2
}elif c == 0b?1?? {
    onehot = 3
}elif c == 0b1??? {
    onehot = 4
}else{
    onehot = 0
}
```


Debugging

Debug statements have no impact

```
# code/debug1.prp
a = 3
I a == 3    # runtime check
C I a == 3   # compile time check

N a != 3    # N (never) is I !(xxx)
C if true { # condition known at compile time
  c = 3
}
C c = 3+4    # C is known at compile time
```

Strings have no compute impact, how at compile time

```
# code/debug2.prp
if c == 3 { # Error unless c is know at compile
  b = "potato"
}else{
  b = "carrot"
}
tup[b] = true
for a:tup ::{
  puts "index:{} value:{}" a.__index a
  I tup[a.__index] == a
}
```