

# **PROLOG ACADEMY**



## **DATA STRUCTURE**

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□ Book followed - Data structures by Seymour Lipschutz (Schaum Series)

LET'S START!

# Circular Linked List

1. Circular linked list is a linked list where all nodes are connected to form a circle. There is no NULL at the end. A circular linked list can be a singly circular linked list or doubly circular linked list.

## Advantages of Circular Linked Lists:

- 1) Any node can be a starting point. We can traverse the whole list by starting from any point. We just need to stop when the first visited node is visited again.
- 2) Useful for implementation of queue.
- 3) It is convenient for the operating system to use a circular list so that when it reaches the end of the list it can cycle around to the front of the list.

# Traversing in circular Linked list

```
void printList()
{
    temp = start;
    do
    {
        printf("%d ", temp->data);
        temp = temp->next;
    }
    while (temp != start)
}
```

# Try These!

- | Sorted insertion in LL
- ▯ Find the mid point of linked list in  $O(n)$
- ▯ Tortoise and hare algorithm
- ▯ Representation of polynomial using linked list.

# Header Linked Lists

- Header linked list is a linked list which always contains a special node called the Header Node, at the beginning of the list.
- It has two types:
  - a) Grounded Header List  
Last Node Contains the NULL Pointer.
  - b) Circular Header List  
Last Node Points Back to the Header Node.

# Grounded HLL

1.A grounded header list is a header list where the last node contains the null pointer.

2.The term “**grounded**” comes from the fact that many texts use the electrical ground symbol to indicate the null pointer.

# Circular HLL

A circular header linked list is a header list where the last node points back to the header node.

- \* START always points to the header node.

Start ->link = NULL means grounded header list is empty.

Start->link = START indicates that a circular header list is empty.

# Traversing a Circular Header List

1. Set  $PTR := LINK[START]$ . [Initializes pointer PTR.]
2. Repeat Steps 3 and 4 while  $PTR \neq START$ .
3. Apply PROCESS to  $INFO[PTR]$ .
4. Set  $PTR := LINK[PTR]$ . [PTR now points to the next node.]  
[End of Step 2 loop.]
5. Exit.



# Doubly Linked list

\* A two-way list is a linear collection of data elements, called nodes, where each node N is divided into three parts:

- 1.Information field

- 2.Forward Link which points to the next node

- 3.Backward Link which points to the previous node

The starting address or the address of first node is stored in START / FIRST pointer .

Another pointer can be used to traverse list from end. This pointer is called END.

# Insertion in the doubly LL

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Bring a node from Avail and call it new

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At beginning

```
new->right=start  
start->left=new  
new->left=NULL  
new->info=Data  
start=new
```

At End

```
temp=start  
while(temp->right!=NULL)  
{ temp=temp->right  
}  
temp->right=new  
new->left=temp  
new->right=NULL  
new->info=Data
```

- At the middle  
new->left=temp  
new->right=temp->right  
temp->right=new  
new->right->left=new  
new->info=Data