CP431 - Assignment 1: MPI Programming for Gaps Between Prime Numbers

Wilfrid Laurier University

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Torin Borton-McCallum - 190824620 Grant Westerholm - 190462000 Brandon Parker - 191593730 Rhea Sharma - 200576620 Taha Amir - 190728860 William Mabia - 190133240 Riley Adams - 190416070

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1 Introduction

This document presents an MPI (Message Passing Interface) program developed in C to determine the largest gap between a pair of consecutive prime numbers up to the ranges of one billion and one trillion. The objective is to explore parallel programming techniques for prime number computations on the SHARCnet cluster. The program utilizes Number Theoretic Functions from the GMP (GNU Multiple Precision Arithmetic Library) for accurate prime number calculations, specifically mpz_probab_prime_p and mpz_nextprime.

The MPI program is designed to execute on 2, 3, 4, 5, 6, 7, and 8 processors, and the performance is benchmarked to analyze its scalability. Results include the identification of the largest prime gap and the two prime numbers within the specified ranges that realize this gap. The documentation addresses code correctness, documentation quality, legibility, and overall presentation.

2 Algorithm Overview

2.1 Design Approach

1. Range Calculation: The global prime number range is initialized, and each MPI process is assigned a specific subrange for computation. The calculation is distributed to maximize parallelization, with each process focusing on its assigned segment.

2. Prime Gap Calculation:

The algorithm iterates through each subrange, identifying consecutive prime numbers and calculating the gaps between them. The largest gap, along with the primes before and after it, is tracked for each process.

3. MPI Communication:

To consolidate results, MPI_Gather is employed to collect information about the largest gaps and associated primes from each process. Process 0 then determines the overall largest gap and its corresponding prime numbers.

4. Result Presentation:

The final results, including the largest gap, the primes before and after the gap, the total number of processes involved, and computation times, are then formatted and displayed. GMP (GNU Multiple Precision) library functions facilitate the handling of large integers for precise calculations.

2.2 Code Implementation

```
1 // include statements
2 #include <stdio.h>
3 #include <stdlib.h>
4 #include <string.h>
5 #include <mpi.h>
6 #include <gmp.h>
8 // Using max fixed size instead of dynamic allocation because I kept running into
      buffer problems
9 #define MAX_STRING_LENGTH 1024
10
int main(int argc, char** argv) {
      // Global variables to be used by each process
      int world_rank, world_size;
14
16
      MPI_Init(&argc, &argv);
      MPI_Comm_rank(MPI_COMM_WORLD, &world_rank);
      MPI_Comm_size(MPI_COMM_WORLD, &world_size);
18
      // Range for each process
20
21
      mpz_t range_start, range_end, global_range;
```

```
// Initialize and set the global range to 1 billion
23
24
      mpz_init_set_ui(global_range, 1000000000);
25
26
      mpz_t current_prime, next_prime, gap, max_gap, last_prime_in_range;
27
      mpz_t prime_before_gap, prime_after_gap;
28
      // Initialize all variables
29
      mpz_init(current_prime);
30
      mpz_init(next_prime);
31
      mpz_init(gap);
32
      mpz_init(max_gap);
33
34
      mpz_init(last_prime_in_range);
      mpz_init(prime_before_gap);
35
36
      mpz_init(prime_after_gap);
37
      mpz_set_ui(max_gap, 0);
38
      39
40
41
      // Calculate the range per process
      unsigned long long int range_per_process = mpz_get_ui(global_range) / world_size;
42
43
      // Set the start and end range for the current process
44
      mpz_init_set_ui(range_start, world_rank * range_per_process + 1);
45
      mpz_init_set_ui(range_end, (world_rank + 1) * range_per_process);
46
47
      // Adjusted for last process - let's it cover entire range
if (world_rank == world_size - 1) {
48
49
          mpz_set(range_end, global_range);
50
51
52
       // Synchronize all processes
53
      MPI_Barrier(MPI_COMM_WORLD);
54
55
56
      // Get the starting time
      double start_time = MPI_Wtime();
57
      // Get the next prime starting from range_start
59
      mpz_nextprime(current_prime, range_start);
60
61
      while (mpz_cmp(current_prime, range_end) <= 0) {</pre>
62
63
           // Get next prime and calculate gap
64
          mpz_nextprime(next_prime, current_prime);
65
66
           // Calculate the gap between the current and next prime
67
          mpz_sub(gap, next_prime, current_prime);
68
69
           // store if largest gap
          if (mpz_cmp(gap, max_gap) > 0) {
71
72
               mpz_set(max_gap, gap);
73
               mpz_set(prime_before_gap, current_prime);
               mpz_set(prime_after_gap, next_prime);
74
          }
75
76
77
           // Move to the next prime
          mpz_set(current_prime, next_prime);
78
      }
79
80
      // Set the last prime in the range
81
      mpz_set(last_prime_in_range, current_prime);
82
83
      // Synchronize all processes
84
      MPI_Barrier(MPI_COMM_WORLD);
85
86
      // Get the ending time
      double end_time = MPI_Wtime();
88
89
90
       // Calculate the elapsed time
      double elapsed_time = end_time - start_time;
91
92
      // Get the resolution of MPI_Wtime
93
```

```
double tick = MPI_Wtick();
94
95
       96
97
       // Prepare strings for communication
98
       char max_gap_str[MAX_STRING_LENGTH], prime_before_str[MAX_STRING_LENGTH],
99
       prime_after_str[MAX_STRING_LENGTH];
100
       // Initialize strings
101
       memset(max_gap_str, 0, MAX_STRING_LENGTH);
       memset(prime_before_str, 0, MAX_STRING_LENGTH);
103
104
       memset(prime_after_str, 0, MAX_STRING_LENGTH);
       mpz_get_str(max_gap_str, 10, max_gap);
106
       // Convert prime_before_gap and prime_after_gap to string
       mpz_get_str(prime_before_str, 10, prime_before_gap);
108
109
       mpz_get_str(prime_after_str, 10, prime_after_gap);
111
       // Allocate memory for string receive in process 0
       char (*all_max_gap_str)[MAX_STRING_LENGTH] = NULL;
       char (*all_prime_before_str)[MAX_STRING_LENGTH] = NULL;
       char (*all_prime_after_str)[MAX_STRING_LENGTH] = NULL;
114
       if (world_rank == 0) {
116
117
           // all_max_gap_str = malloc(world_size * MAX_STRING_LENGTH * sizeof(char));
118
           // all_prime_before_str = malloc(world_size * MAX_STRING_LENGTH * sizeof(char)
119
       ):
           // all_prime_after_str = malloc(world_size * MAX_STRING_LENGTH * sizeof(char))
120
           // This lets it compile with mpi c++ above only compiles with mpic
           all_max_gap_str = (char (*)[MAX_STRING_LENGTH])malloc(world_size *
123
       MAX_STRING_LENGTH * sizeof(char));
           all_prime_before_str = (char (*)[MAX_STRING_LENGTH])malloc(world_size *
124
       MAX_STRING_LENGTH * sizeof(char));
           all_prime_after_str = (char (*)[MAX_STRING_LENGTH])malloc(world_size *
       MAX_STRING_LENGTH * sizeof(char));
       }
127
128
       // Gather fixed size strings
       MPI_Gather(max_gap_str, MAX_STRING_LENGTH, MPI_CHAR, all_max_gap_str,
129
       MAX_STRING_LENGTH, MPI_CHAR, O, MPI_COMM_WORLD);
       MPI_Gather(prime_before_str, MAX_STRING_LENGTH, MPI_CHAR, all_prime_before_str,
130
       MAX_STRING_LENGTH, MPI_CHAR, O, MPI_COMM_WORLD);
       MPI_Gather(prime_after_str, MAX_STRING_LENGTH, MPI_CHAR, all_prime_after_str,
       MAX_STRING_LENGTH, MPI_CHAR, O, MPI_COMM_WORLD);
       if (world_rank == 0) {
133
           mpz_t global_max_gap, global_prime_before_gap, global_prime_after_gap;
134
135
           // Initialize global values
136
           mpz_init(global_max_gap);
137
           mpz_init(global_prime_before_gap);
138
           mpz_init(global_prime_after_gap);
139
140
           mpz_set_ui(global_max_gap, 0);
141
           // Process received strings
142
           for (int i = 0; i < world_size; ++i) {</pre>
143
               mpz_t temp_max_gap;
144
145
               mpz_init(temp_max_gap);
               mpz_set_str(temp_max_gap, all_max_gap_str[i], 10);
146
147
               // gmp_printf("temp_max_gap: %Zd\n", temp_max_gap);
148
149
               // gmp_printf("global_max_gap: %Zd\n", global_max_gap);
150
               if (mpz_cmp(temp_max_gap, global_max_gap) > 0) {
151
                   mpz_set(global_max_gap, temp_max_gap);
                   mpz_set_str(global_prime_before_gap, all_prime_before_str[i], 10);
153
                   mpz_set_str(global_prime_after_gap, all_prime_after_str[i], 10);
155
```

```
156
157
                // Clear temporary maximum gap
158
                mpz_clear(temp_max_gap);
           }
159
            // gmp_printf("Largest gap: %Zd, between primes %Zd and %Zd\n", max_gap,
161
       prime_before_gap, prime_after_gap);
            gmp_printf("Largest gap: %Zd, between primes %Zd and %Zd\n", global_max_gap,
162
       global_prime_before_gap, global_prime_after_gap);
           printf("Total processes: %d\n", world_size);
163
            \label{lem:printf("Total computation time: %e seconds n", elapsed_time);}
164
            printf("Average computation time per process: %e seconds\n", elapsed_time /
       world_size);
           printf("Resolution of MPI_Wtime: %e seconds\n", tick);
167
            // Cleanup
169
           mpz_clear(global_max_gap);
           mpz_clear(global_prime_before_gap);
           mpz_clear(global_prime_after_gap);
171
172
           free(all_max_gap_str);
           free(all_prime_before_str);
           free(all_prime_after_str);
174
176
       // Cleanup
177
       mpz_clear(range_start);
178
179
       mpz_clear(range_end);
       mpz_clear(global_range);
180
       mpz_clear(current_prime);
181
       mpz_clear(next_prime);
182
       mpz_clear(gap);
183
184
       mpz_clear(max_gap);
       mpz_clear(last_prime_in_range);
185
186
       mpz_clear(prime_before_gap);
       mpz_clear(prime_after_gap);
187
188
       MPI_Finalize();
189
190
       return 0;
191 }
```

3 Algorithm Results

3.1 One Billion

Running the algorithm listed above, the largest gap we produced between two contiguous prime numbers is 282. The two prime numbers that materialize this gap are 436273009 and 436273291. This is consistent with the result produced from a relatively equivalent serially programmed algorithm to accomplish the same task. The serial version took approximately 7 times longer to run to than the parallel implementation with 8 processes.

3.2 One Trillion

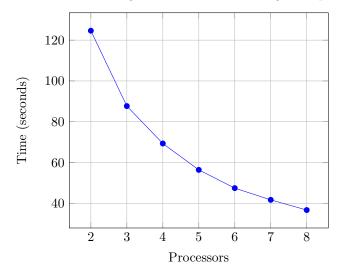
From one billion to one trillion, we expanded the algorithm to look for a larger gap between two contiguous prime numbers. The largest gap found increased from 282 to 540. The two prime numbers that materialize this largest gap are 738832927927 and 738832928467. Trying to run a serial implementation on one trillion is not feasible since the parallel solution's runtime was around 5 hours.

4 Benchmarking

The benchmarking results presented here evaluate the performance of a parallelized algorithm for computing prime gaps. The primary metric considered is the largest observed gap, which measures 282 between primes 436273009 and 436273291. The computations were conducted using MPI, a widely used message-passing interface for parallel computing.

4.1 Scalability Analysis - One Billion

Parallel Execution Time Scaling with Processors for Largest Gap Between One Billion



4.1.1 8 Processes

In the initial run with 8 processes, the total computation time was 124.63 seconds, yielding an average computation time per process of 15.58 seconds. This serves as a baseline for evaluating the scalability of the algorithm.

4.1.2 Increasing Processes (12, 16, 20)

As the number of processes increased to 12, 16, and 20, there was a notable improvement in performance. The total computation time decreased to 87.67 seconds, 69.32 seconds, and 56.39 seconds, respectively, showcasing the scalability of the parallelized algorithm.

4.1.3 Optimal Scaling (24 Processes)

Further scaling the computation to 24 processes increased efficiency, resulting in a total computation time of 47.44 seconds and an average computation time per process of 1.98 seconds. This suggests a near-linear scaling, indicative of effective parallelization.

4.1.4 Continued Improvement (28, 32 Processes)

Continuing the trend, 28 processes led to a total computation time of 41.68 seconds, and with 32 processes, the total computation time reduced to 36.69 seconds. The average computation time per process dropped to 1.49 seconds and 1.15 seconds, respectively, emphasizing the continued improvement with an increased number of processes.

4.2 MPI Wtime Resolution

It is important to note that the resolution of MPI Wtime, a function measuring time in MPI programs, remained consistent at 1.000000e-09 seconds across all experiments. This consistent resolution highlights the accuracy of the timing measurements.

4.3 One Trillion

This computation involved 10 processes, resulting in a total computation time of 16,651.70 seconds, translating to an average computation time per process of 1,665.17 seconds. The resolution of MPI Wtime for this experiment was set at 1.000000e-06 seconds. Comparing these results with the previous benchmarking conducted within the billion range, where the largest gap was 282, there is a noticeable

increase in both the size of the gap and the computation time. The larger search space in the trillion range naturally leads to more extensive computations, reflected in the increased gap size and the higher overall computation time. This highlights the scalability of the algorithm to handle significantly larger ranges.

5 Conclusion

In conclusion, the code implemented an MPI program in C designed to find the largest gap between consecutive prime numbers within the ranges of one billion and one trillion. The use of the GMP library for precise calculations such as the range and prime gap calculations, allows us to explore parallel programming by utilizing two to eight processors and assessing scalability through benchmarking. The algorithm efficiently distributes computations, identifies prime numbers and tracks gaps, consolidating results with MPI_Gather. The results demonstrate efficient parallelization, with a largest gap of 282 in the billion range and 540 in the trillion range. The scalability in terms of this algorithm is evident, with reduced computation time as the number of processes increases. The consistent resolution of MPI Wtime emphasizes accurate timing measurements and highlights the algorithm's reliability across varying computational scales. Overall, this research showcases the algorithm's scalability, efficiently handling larger search spaces and substantiating its effectiveness in parallelized prime number computations.

6 Appendices

6.1 Run Times for Parallel Implementation

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 8

Total computation time: 1.246293e+02 seconds

Average computation time per process: 1.557866e+01 seconds

Resolution of MPI_Wtime: 1.000000e-09 seconds

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 12

Total computation time: 8.767491e+01 seconds

Average computation time per process: 7.306242e+00 seconds

Resolution of MPI_Wtime: 1.000000e-09 seconds

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 16

Total computation time: 6.932332e+01 seconds

Average computation time per process: 4.332707e+00 seconds

Resolution of MPI_W time: 1.000000e-09 seconds

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 20

Total computation time: 5.639275e+01 seconds

Average computation time per process: 2.819637e+00 seconds

Resolution of MPI_Wtime: 1.000000e-09 seconds

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 24

Total computation time: 4.744021e+01 seconds

Average computation time per process: 1.976675e+00 seconds

Resolution of MPI_Wtime: 1.000000e-09 seconds

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 28

Total computation time: 4.167845e+01 seconds

Average computation time per process: 1.488516e+00 seconds

Resolution of MP_Wtime: 1.000000e-09 seconds

Largest gap: 282, between primes 436273009 and 436273291

Total processes: 32

Total computation time: 3.668536e+01 seconds

Average computation time per process: 1.146418e+00 seconds

Resolution of MPI_W time: 1.000000e-09 seconds

6.2 Run Time for Serial Implementation

Time taken: 938.300000 seconds

lprime: 436273009 rprime: 43627329821 gap: 282

6.3 Run Time for 1 trillion on Personal Computer

Largest gap: 540, between primes 738832927927 and 738832928467

Total processes: 10

Total computation time: 1.665170e+04 seconds

Average computation time per process: 1.665170e+03 seconds

Resolution of MPI_Wtime: 1.000000e-06 seconds