

The Sudoku Puzzle

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Introduction

Sudoku is a placement puzzle, also known as Number Place in the U.S.A.

The game consists most frequently of a 9 x 9 grid, divided in 9 subgrids with dimension 3 x 3 called “regions”.

The purpose is to enter a digit from 1 to 9 (or other symbols e.g. letters, icons) in each cell of the grid so that each row, column and region contains only one instance of each digit.

We implemented a **SAT solver** (instead of combining backtracking and methods for constraint propagation as other Sudoku solver) to figure out a correct solution for the Sudoku.

Basically the Sudoku is translated into a propositional formula that can be satisfied only if the Sudoku has a solution.

Once the propositional formula is generated, the SAT solver tries to find a satisfying assignment that will become the solution for the original Sudoku.

Reduces Sudoku problem to a SAT clause

Digits are modelled by a datatype with nine elements (1, ..., 9).

We can say that the grid cells (x_1, \dots, x_9) are valid if they contain at least and at

most 1 digit each.

Definition 1

$$valid(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) \equiv \bigwedge_{d=1}^9 \bigvee_{i=1}^9 x_i = d$$

Labeling the 81 cells, we can check if the cells generate a correct solution for the Sudoku puzzle.

Sudoku definition

$$sudoku(\{x_{ij}\}_{i,j \in \{1, \dots, 9\}}) \equiv \bigwedge_{i=1}^9 valid(x_{i1}, x_{i2}, x_{i3}, x_{i4}, x_{i5}, x_{i6}, x_{i7}, x_{i8}, x_{i9})$$

$$\bigwedge_{j=1}^9 valid(x_{1j}, x_{2j}, x_{3j}, x_{4j}, x_{5j}, x_{6j}, x_{7j}, x_{8j}, x_{9j})$$

$$\bigwedge_{i,j \in \{1,4,7\}} valid(x_{ij}, x_{i(j+1)}, x_{i(j+2)}, x_{(i+1)j}, x_{(i+1)(j+1)}, x_{(i+1)(j+2)},$$

$$x_{(i+2)j}, x_{(i+2)(j+1)}, x_{(i+2)(j+2)})$$

SAT Solver

We're now introducing 9 boolean variables for each cell in the 9x9 grid ($9^3 = 729$ variables in total) in order to encode a Sudoku.

Each boolean value holds the truth value of the equation $x_{i,j} = d$.

A clause

$$\bigvee_{d=1}^9 p_{ij}^d$$

assures that a cell contains one of the nine accepted digits, whilst 36 clauses

$$\bigwedge_{1 \leq d < d' \leq 9} \neg p_{ij}^d \vee \neg p_{ij}^{d'}$$

assure that a cell doesn't hold two different digits.

Since the number of digits is equal to the number of cells in every row, column or region, then the nine grid cells (x_1, \dots, x_9) hold distinct values.

Lemma 1

$$\begin{aligned} \text{valid}(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) &\iff \bigwedge_{1 \leq i < j \leq 9} x_i \neq x_j \\ &\iff \bigwedge_{1 \leq i < j \leq 9} \bigwedge_{d=1}^9 x_i \neq d \vee x_j \neq d \end{aligned}$$

The given formula, when converted into **SAT**, is translated into 9 clauses * 36 inequations = 324 clauses, each of length 2. This allows more unit propagation at the boolean level than what was previously stated in **Def. 1**, which gives us the possibility of cross hatching digits (a technique used in Sudoku to reduce the search space).

Summarising, up to now, our **SAT** is composed of:

1. 81 definedness clauses of length 9
2. 81 · 36 uniqueness clauses of length 2
3. 27 · 324 validity clauses of length 2

For a total of **11745** clauses.

Since we know a priori the value of some cells (user input), we can just add these condition to our clause list.

Finally, our encoding produces a propositional formula already in **CNF**, so the conversion into **DIMACS CNF** (the input format for most sat solvers) is trivial.