Moscow IPT

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Contest (1)

CMakeLists.txt

```
cmake_minimum_required(VERSION 3.24)
project(main)

set(CMAKE_CXX_STANDARD 20)

file(GLOB sources *.cpp)

add_compile_definitions(LOCAL)

if(CMAKE_BUILD_TYPE STREQUAL "Debug")
    set(CMAKE_CXX_FLAGS "${CMAKE_CXX_FLAGS} -D_GLIBCXX_DEBUG -
        fsanitize=address,undefined")
elseif(CMAKE_BUILD_TYPE STREQUAL "Release")
    add_compile_options(-O3)
endif()

foreach(src ${sources})
    get_filename_component(name ${src} NAME_WE)
    add_executable(${name} ${src})
endforeach()
```

Numerical (2)

2.1 Newton's method

To compute $B = \frac{1}{A}$ modulo x^m: define $B_1 = \{inv(A[0])\}$ and $B_{2n} = B_n(2 - A \cdot B_n)$.

To compute $B = \sqrt{A}$ modulo x^m: define $B_1 = \{\sqrt{A[0]}\}$ and $B_{2n} = \frac{B_n}{2} + \frac{A}{2B_n}$.

To compute $B = \log(1 + xA)$ modulo x^m: $B = \int \frac{(1+xA)'}{1+xA}$.

If T is EGF for some objects then $C = -\log(1 - T) = \sum_{k=1}^{+\infty} \frac{T^k}{k}$ is EGF for cycles of them.

To compute $B = e^{xA}$ modulo x^m: define $B_1 = \{1\}$ and $B_{2n} = B_n(1 + A - \log B_n)$.

If T is EGF for some objects then $F = e^T = \sum_{k=0}^{\infty} \frac{T^k}{k!}$ is EGF for their unordered combinations.

In general case you have equation $P(B, x) = 0$. (f.e. $\log B - A(x) = 0$).

The transition is $B_{2n-\alpha} = B_n - \frac{P(B,x)}{P'_B(B,x)}$. (if $P'_B(B, x)$ is divisible by x^α).

2.2 Additional modulos for FFT

Modulo	Other form	Roots
998244353	$(119 << 23) + 1$	$3, 62$
167772161	$(5 << 25) + 1$	$3, 62$
469762049	$(7 << 26) + 1$	$3, 62$
1004535809	$(479 << 21) + 1$	$3, 62$
1012924417	$(483 << 21) + 1$	62

FFT
Description: Applies the discrete Fourier transform to a sequence of numbers modulo MOD.
Time: $\mathcal{O}(n \log n)$

```cpp
int rev[N], root[N];

void init(int n) {
    static int last_init = -1;
    if (n == last_init) return;
    last_init = n;
    for (int i = 1; i < n; ++i) {
        rev[i] = (rev[i >> 1] >> 1) | (i & 1) * (n >> 1);
    }
    const int root_n = binpow(ROOT, (MOD - 1) / n);
    for (int i = 0, cur = 1; i < n / 2; ++i) {
        root[i + n / 2] = cur;
        cur = mul(cur, root_n);
    }
    for (int i = n / 2 - 1; i >= 0; --i) {
        root[i] = root[i << 1];
    }
}

void dft(int* f, int n, bool inverse = false) {
    init(n);
    for (int i = 0; i < n; ++i) {
        if (i < rev[i]) swap(f[i], f[rev[i]]);
    }
    for (int k = 1; k < n; k <<= 1)
        for (int i = 0; i < n; i += (k << 1))
```

```
            for (int j = 0; j < k; ++j) {
                int z = mul(f[i + j + k], root[j + k]);
                f[i + j + k] = add(f[i + j], MOD - z);
                f[i + j] = add(f[i + j], z);
            }
    if (inverse) {
        reverse(f + 1, f + n);
        const int inv_n = inv(n);
        for (int i = 0; i < n; ++i) f[i] = mul(f[i], inv_n);
    }
}
```

Middle product
Description: Calculates middle-product of two arrays using Tellegen's principle.
Time: $\mathcal{O}(n \log n)$

```
vector<int> mulT(const vector<int>& a, const vector<int>& b) {
    int n = sz(a), m = sz(b), k = 1;
    while (k < n) k <<= 1;
    fill_n(fft1, k, 0), fill_n(fft2, k, 0);
    copy(all(a), fft1), copy(all(b), fft2);
    dft(fft1, k, true), dft(fft2, k);
    for (int i = 0; i < k; ++i) fft1[i] = mul(fft1[i], fft2[i]);
    dft(fft1, k);
    return {fft1, fft1 + n - m + 1};
}
```

Polynomial division
Description: Finds first n coefficients of $1/a$ modulo MOD.
Time: $\mathcal{O}(n \log n)$

```
vector<int> inv(vector<int> a, int n) {
    assert(a[0] != 0);
    vector res{inv(a[0])};
    for (int k = 1; k < n; k <<= 1) {
        auto tmp = mul(res, res);
        if ((int)tmp.size() > n) {
            tmp.resize(n);
        }
        tmp = mul(tmp, {a.begin(), a.begin() + min(2 * k, sz(a))});
        tmp.resize(2 * k);
        for (int i = 0; i < k; ++i) {
            tmp[i] = sub(add(res[i], res[i]), tmp[i]);
            tmp[k + i] = sub(0, tmp[k + i]);
        }
        res.swap(tmp);
    }
```

Berlekamp-Massey
Description: Returns the polynomial of a recurrent sequence of order n from the first $2n$ terms.
Usage: `berlekamp_massey({0, 1, 1, 3, 5, 11})` // {1, -1, -2}
Time: $\mathcal{O}(n^2)$

```
vector<int> berlekamp_massey(vector<int> s) {
    int n = sz(s), L = 0, m = 0;
    vector<int> c(n), b(n), t;
    c[0] = b[0] = 1;
    int eval = 1;
    for (int i = 0; i < n; ++i) {
        m++;
        int delta = 0;
        for (int j = 0; j <= L; ++j) {
            delta = add(delta, mul(c[j], s[i - j]));
        }
        if (delta == 0) continue;
        t = c;
        int coef = mul(delta, inv(eval));
        for (int j = m; j < n; ++j) {
            c[j] = sub(c[j], mul(coef, b[j - m]));
        }
        if (2 * L > i) continue;
        L = i + 1 - L, m = 0, b = t, eval = delta;
    }
    c.resize(L + 1);
    return c;
}
```

2.3 Linear recurrence

Let A be generating function for our recurrence, C be its characteristic polynomial and $k = |C|$.

Let $D = C \cdot A$. Then $D \bmod x^k = D$

$$A = \frac{(A \bmod x^k)C \bmod x^k}{C} = \frac{A_0 C \bmod x^k}{C}$$

$$[x^n]\frac{P(x)}{Q(x)} = [x^n]\frac{P(x)Q(-x)}{Q(x)Q(-x)}$$

2.4   Xor-And-Or Convolutions

To compute

$$c_i = \sum_{j,k:j|k=i} a_j \cdot b_k$$

we transform (a) and (b) into their subset sums, take the elementwise product, and then apply the inverse transform.

For

$$c_i = \sum_{j,k:j\&k=i} a_j \cdot b_k$$

we transform (a) and (b) into their superset sums and do the same.

To compute

$$c_i = \sum_{j,k:j\oplus k=i} a_j \cdot b_k$$

we applay Walsh-Hadamard transform, take the elementwise product, and then apply the inverse transform.

Walsh-Hadamard
Description: Applies Walsh-Hadamard transform
Time: $\mathcal{O}\left(2^n \cdot n\right)$

```cpp
vector<int> f(1 << n);
for (int k = 0; k < n; ++k) {
    for (int mask = 0; mask < (1 << n); ++mask) {
        if (~mask >> k & 1) {
            int u = f[mask], v = f[mask ^ (1 << k)];
            f[mask] = u + v;
            f[mask ^ (1 << k)] = u - v;
        }
    }
}
```

Flows (3)

Dinitz
Description: Finds maximum flow using Dinitz algorithm.
Time: $\mathcal{O}\left(n^2 m\right)$

```cpp
struct Edge {
    int to, cap, flow;
};
```

```cpp
vector<Edge> E;
vector<int> gr[N];

int n;
int d[N], ptr[N];

bool bfs(int v0 = 0, int cc = 1) {
    fill(d, d + n, -1);
    d[v0] = 0;
    vector<int> q{v0};
    for (int st = 0; st < sz(q); ++st) {
        int v = q[st];
        for (int id : gr[v]) {
            auto [to, cp, fl] = E[id];
            if (d[to] != -1 || cp - fl < cc) continue;
            d[to] = d[v] + 1;
            q.emplace_back(to);
        }
    }
    return d[n - 1] != -1;
}

int dfs(int v, int flow, int cc = 1) {
    if (v == n - 1 || !flow) return flow;
    for (; ptr[v] < sz(gr[v]); ++ptr[v]) {
        auto [to, cp, fl] = E[gr[v][ptr[v]]];
        if (d[to] != d[v] + 1 || cp - fl < cc) continue;
        int pushed = dfs(to, min(flow, cp - fl), cc);
        if (pushed) {
            int id = gr[v][ptr[v]];
            E[id].flow += pushed;
            E[id ^ 1].flow -= pushed;
            return pushed;
        }
    }
    return 0;
}

ll dinitz() {
    ll flow = 0;
    for (int c = INF; c > 0; c >>= 1) {
        while (bfs(0, c)) {
            fill(ptr, ptr + n, 0);
            while (int pushed = dfs(0, INF, c))
                flow += pushed;
        }
    }
```

```
    }
    return flow;
}
```

MCMF
Description: Finds Minimal Cost Maximal Flow.

```
struct Edge {
    ll to, f, c, w;
};

vector<Edge> E;
vector<int> gr[N];

void add_edge(int u, int v, ll c, ll w) {
    gr[u].push_back(sz(E));
    E.emplace_back(v, 0, c, w);
    gr[v].push_back(sz(E));
    E.emplace_back(u, 0, 0, -w);
}

pair<ll, ll> mcmf(int n) {
    vector<ll> dist(n);
    vector<ll> pr(n);
    vector<ll> phi(n);
    auto dijkstra = [&] {
        fill(all(dist), INF);
        dist[0] = 0;
        priority_queue<pair<ll, int>, vector<pair<ll, int>>,
            greater<>> pq;
        pq.emplace(0, 0);
        while (!pq.empty()) {
            auto [d, v] = pq.top();
            pq.pop();
            if (d != dist[v]) continue;
            for (int idx : gr[v]) {
                if (E[idx].c == E[idx].f) continue;
                int to = E[idx].to;
                ll w = E[idx].w + phi[v] - phi[to];
                if (dist[to] > d + w) {
                    dist[to] = d + w;
                    pr[to] = idx;
                    pq.emplace(d + w, to);
                }
            }
        }
    }
```

```
    };

    ll total_cost = 0, total_flow = 0;
    while (true) {
        dijkstra();
        if (dist[n - 1] == INF) break;
        ll min_cap = INF;
        int cur = n - 1;
        while (cur != 0) {
            min_cap = min(min_cap, E[pr[cur]].c - E[pr[cur]].f);
            cur = E[pr[cur] ^ 1].to;
        }
        cur = n - 1;
        while (cur != 0) {
            E[pr[cur]].f += min_cap;
            E[pr[cur] ^ 1].f -= min_cap;
            total_cost += min_cap * E[pr[cur]].w;
            cur = E[pr[cur] ^ 1].to;
        }
        total_flow += min_cap;
        for (int i = 0; i < n; ++i) {
            phi[i] += dist[i];
        }
    }

    return {total_flow, total_cost};
}
```

Number Theory (4)

Extended GCD
Description: Finds two integers x and y, such that $ax + by = gcd(a, b)$.

```
ll exgcd(ll a, ll b, ll &x, ll &y) {
    if (!b) return x = 1, y = 0, a;
    ll d = exgcd(b, a % b, y, x);
    return y -= a/b * x, d;
}
```

CRT
Description: Chinese Remainder Theorem. crt(a, m, b, n) computes x s.t. $x \equiv a \mod m, x \equiv b \mod n$.
Time: $\mathcal{O}(\log n)$

```
ll crt(ll a, ll m, ll b, ll n) {
```

```
    if (n > m) swap(a, b), swap(m, n);
    ll x, y, g = exgcd(m, n, x, y);
    assert((a - b) % g == 0); // else no solution
    x = (b - a) % n * x % n / g * m + a;
    return x < 0 ? x + m * n / g : x;
}
```

Miller-Rabin

Description: Deterministic Miller-Rabin primality test. Guaranteed to work for numbers up to $7 \cdot 10^{18}$
Time: 7 times the complexity of $a^b \bmod c$.

```
bool is_prime(ll n) {
    if (n < 2 || n % 6 % 4 != 1) return (n | 1) == 3;
    ll A[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022};
    ll s = __builtin_ctzll(n - 1), d = n >> s;
    for (ll a : A) {
        ll p = binpow(a % n, d, n), i = s;
        while (p != 1 && p != n - 1 && a % n && i--)
            p = mul(p, p, n);
        if (p != n - 1 && i != s) return 0;
    }
    return 1;
}
```

Pollard-Rho

Description: Pollard-rho randomized factorization algorithm. Returns prime factors of a number, in arbitrary order.
Time: $\mathcal{O}\left(n^{1/4}\right)$, less for numbers with small factors.

```
ll pollard(ll n) {
    auto f = [n](ll x) { return mul(x, x, n) + 1; };
    ll x = 0, y = 0, t = 30, prd = 2, i = 1, q;
    while (t++ % 40 || __gcd(prd, n) == 1) {
        if (x == y) x = ++i, y = f(x);
        if ((q = mul(prd, max(x,y) - min(x,y), n))) prd = q;
        x = f(x), y = f(f(y));
    }
    return __gcd(prd, n);
}

vector<ll> factor(ll n) {
    if (n == 1) return {};
    if (is_prime(n)) return {n};
    ll x = pollard(n);
    auto l = factor(x), r = factor(n / x);
    l.insert(l.end(), all(r));
```

```
    return l;
}
```

4.1 Möbius function

$$\mu(n) = \begin{cases} 0, n \text{ is not square free} \\ 1, n \text{ has even number of prime factors} \\ -1, n \text{ has odd number of prime factors} \end{cases}$$

$$\sum_{d|n} \mu(d) = \text{int}(n = 1)$$

Möbius inversion:

$$g(n) = \sum_{d|n} f(d) \iff f(n) = \sum_{d|n} \mu(d)g(n/d) = \sum_{d|n} \mu(n/d)g(d)$$

4.2 Sums of powers

$$\sum_{s=1}^{n} s^k = \sum_{s=1}^{k} s! S(k,s) C_{n+1}^{s+1}$$

$\sum_{s=1}^{n} s^1 = 1\frac{(n+1)n}{2!}$

$\sum_{s=1}^{n} s^2 = 1\frac{(n+1)n}{2!} + 2\frac{(n+1)n(n-1)}{3!} = \frac{n(n+1)(2n+1)}{6}$

$\sum_{s=1}^{n} s^3 = 1\frac{(n+1)n}{2!} + 6\frac{(n+1)n(n-1)}{3!} + 6\frac{(n+1)n(n-1)(n-2)}{4!} = \frac{(n+1)^2 n^2}{4}$

4.3 Floor sum

$$f(a, b, c, n) = \sum_{x=0}^{n} \left\lfloor \frac{ax+b}{c} \right\rfloor$$

Let $m = \left\lfloor \frac{an+b}{c} \right\rfloor$. When $a < c$ and $b < c$, we have:

$$f(a, b, c, n) = mn - f(c, c - b - 1, a, m - 1)$$

Otherwise:

$$f(a, b, c, n) = \frac{n(n+1)}{2} \left\lfloor \frac{a}{c} \right\rfloor + (n+1) \left\lfloor \frac{b}{c} \right\rfloor + f(a \bmod c, b \bmod c, c, n)$$

4.4   Square of floor sum

$$g(a, b, c, n) = \sum_{x=0}^{n} x \left\lfloor \frac{ax + b}{c} \right\rfloor$$

$$h(a, b, c, n) = \sum_{x=0}^{n} \left\lfloor \frac{ax + b}{c} \right\rfloor^2$$

Let $m = \left\lfloor \frac{an+b}{c} \right\rfloor$. When $a < c$ and $b < c$, we have:

$$g(a, b, c, n) = \frac{mn(n + 1)}{2} - \frac{f(c, c - b - 1, a, m - 1)}{2} - \frac{h(c, c - b - 1, a, m - 1)}{2}$$

$$h(a, b, c, n) = mn(m + 1) - 2g(c, c - b - 1, a, m - 1) - 2f(c, c - b - 1, a, m - 1)$$
$$-f(a, b, c, n)$$

Otherwise:

$$g(a, b, c, n) = \frac{n(n + 1)(2n + 1)}{6} \left\lfloor \frac{a}{c} \right\rfloor + \frac{n(n + 1)}{2} \left\lfloor \frac{b}{c} \right\rfloor + g(a \bmod c, b \bmod c, c, n)$$

$$h(a, b, c, n) = \frac{n(n + 1)(2n + 1)}{6} \left\lfloor \frac{a}{c} \right\rfloor^2 + (n + 1) \left\lfloor \frac{b}{c} \right\rfloor^2 + n(n + 1) \left\lfloor \frac{a}{c} \right\rfloor \left\lfloor \frac{b}{c} \right\rfloor +$$

$$+h(a \bmod c, b \bmod c, c, n) + 2 \left\lfloor \frac{b}{c} \right\rfloor f(a \bmod c, b \bmod c, c, n) +$$

$$+2 \left\lfloor \frac{a}{c} \right\rfloor g(a \bmod c, b \bmod c, c, n)$$

Floor Sum implementation
Description: Calculates a tuple containing $f(a, b, c, n)$, $g(a, b, c, n)$ and $h(a, b, c, n)$.
Time: $\mathcal{O}(\log C)$

```cpp
tuple<ll, ll, ll> floor_sum(ll a, ll b, ll c, ll n) {
    if (a == 0) {
        return {
                mul(n + 1, b / c),
                mul(b / c, n, n + 1, INV_2),
                mul(n + 1, b / c, b / c)
        };
    }
    if (a >= c || b >= c) {
        auto [f, g, h] = get(a % c, b % c, c, n);
```

```cpp
        return {
                (f + mul(n, n + 1, INV_2, a / c) + mul(n + 1, b / c
                    )) % MOD,
                (g + mul(a / c, n, n + 1, 2 * n + 1, INV_6) + mul(b
                    / c, n, n + 1, INV_2)) % MOD,
                (h + 2 * mul(b / c, f) + 2 * mul(a / c, g) + mul(a
                    / c, a / c, n, n + 1, 2 * n + 1, INV_6) + mul(b
                    / c, b / c, n + 1) + mul(a / c, b / c, n, n +
                    1)) % MOD
        };
    }
    ll m = (a * n + b) / c;
    auto [f, g, h] = get(c, c - b - 1, a, m - 1);
    return {
            (mul(m, n) + MOD - f) % MOD,
            mul((mul(m, n, n + 1) + MOD * 2ll - f - h) % MOD, INV_2
                ),
            (mul(n, m, m) + MOD * 3ll - f - 2 * g) % MOD
    };
}
```

Combinatorics (5)

5.1   Derangements
Number of n-permutations that none of the elements appears in their original position.

$$D_n = (n - 1)(D_{n-1} + D_{n-2}) = nD_{n-1} + (-1)^n \approx \frac{n!}{e}$$

5.2   Burnside's lemma
Given a group G of symmetries and a set Ω, the number of elements of Ω up to symmetry equal

$$\frac{1}{|G|} \sum_{g \in G} N(g)$$

where $N(g)$ is the number of elements fixed by g ($g(x) = x$).

5.3 Stirling numbers (first kind)

Unsigned Stirling numbers of the first kind $c_{n,k}$ is the number of permutations of n elements with k cycles as well as the coefficient on x^k in the expansion $x(x+1)(x+2)\ldots(x+(n-1))$.

Signed Stirling numbers of the first kind $s_{n,k}$ is the coefficient on x^k in the expansion $x(x-1)(x-2)\ldots(x-(n-1))$.

$$c_{0,0} = s_{0,0} = 1 \quad c_{k,0} = s_{k,0} = 0 \quad c_{n,k} = s_{n,k} = 0 \text{ for } k > n$$

$$c_{n,k} = c_{n-1,k-1} + (n-1)c_{n-1,k}$$

$$s_{n,k} = s_{n-1,k-1} - (n-1)s_{n-1,k}$$

$$s_{n,k} = (-1)^{n+k}c_{n,k}$$

$$\text{EGF: } \sum_{n=0}^{\infty}\sum_{k=0}^{n} s_{n,k}\frac{x^n}{n!}y^k = (1+x)^y$$

$$\text{EGF: } \sum_{n=k}^{\infty} s_{n,k}\frac{x^n}{n!} = \frac{(\log(1+x))^k}{k!}$$

5.4 Stirling numbers (second kind)

The Stirling numbers of the second kind $S(n,k)$, count the number of ways to partition a set of n labelled objects into k nonempty unlabelled subsets.

$$S(n,n) = 1 \text{ for } n \geq 0 \quad S(0,n) = S(n,0) = 0 \text{ for } n > 0$$

$$S(n+1,k) = k \cdot S(n,k) + S(n,k-1)$$

$$S(n,k) = \sum_{t=0}^{k} \frac{(-1)^{k-t}t^n}{(k-t)!t!}$$

$$\text{EGF: } \sum_{n=0}^{\infty}\sum_{k=0}^{n} S(n,k)\frac{x^n}{n!}y^k = e^{y(e^x-1)}$$

$$\text{EGF: } \sum_{n=k}^{\infty} S(n,k)\frac{x^n}{n!} = \frac{(e^x-1)^k}{k!}$$

5.5 Bell numbers

Bell number B_n is the number of partitions of n labeled elements.

$$B_0 = B_1 = 1$$

$$B_n = \sum_{k=0}^{n-1} C_{n-1}^k B_k = \sum_{k=0}^{n} S(n,k)$$

$$\text{EGF: } \sum_{n=1}^{\infty} \frac{B_n}{n!}x^n = e^{e^x-1}$$

5.6 Narayana numbers

Narayana number $N(n,k)$ is the number of correct bracket sequences with length $2n$ and exactly k distinct nestings. Also the number of unlabeled ordered rooted trees with $n+1$ vertices and k leaves.

$$N(n,k) = \frac{1}{n}C_n^k C_n^{k-1}$$

5.7 Labeled unrooted trees

Every tree on n vertices has unique sequence of $n-2$ integers from $\{1\ldots n\}$ assosiated with the tree.

Vertex with degree d appears in sequence $d-1$ times.

On n vertices: n^{n-2}.

With degrees $d_1, d_2, \ldots, d_n$: $\frac{(n-2)!}{(d_1-1)!\ldots(d_n-1)!}$.

Strings (6)

KMP
Description: Calculates prefix function and Z-function of the given string.
Time: $\mathcal{O}(n)$

```cpp
vector<int> pi(const string& s) {
    vector<int> p(sz(s));
    for (int i = 1; i < sz(s); ++i) {
        int g = p[i - 1];
        while (g && s[i] != s[g]) g = p[g - 1];
        p[i] = g + (s[i] == s[g]);
    }
}
```

```
    return p;
}


vector<int> zf(const string& s) {
    vector<int> z(sz(s));
    int l = -1, r = -1;
    for (int i = 1; i < sz(s); ++i) {
        z[i] = i >= r ? 0 : min(r - i, z[i - l]);
        while (i + z[i] < sz(s) && s[i + z[i]] == s[z[i]]) ++z[i];
        if (i + z[i] > r) l = i, r = i + z[i];
    }
    return z;
}
```

Manacher
Description: Manacher's algorithm
Time: $\mathcal{O}(n)$

```
vector<int> manacher(const std::string& s) {
    int n = s.size();
    vector<int> res(2 * n + 1);
    int longest = 0, root_longest = 0;
    for (int i = 1; i < 2 * n + 1; ++i) {
        int pal;
        if (i > longest) {
            pal = 1;
        } else {
            pal = min(longest - i, res[2 * root_longest - i]);
        }
        while (pal < i && i + pal < 2 * n && s[(i - pal - 2) / 2]
            == s[(i + pal) / 2]) {
            pal += 2;
        }
        res[i] = pal;
        if (i + pal > longest) {
            longest = i + pal;
            root_longest = i;
        }
    }
    return res;
}
```

Aho-Corasick
Description: Builds Aho-Corasick
Time: $\mathcal{O}(nC)$

```
const int C = 26;
```

```
struct node {
    int nx[C], first = -1, suff = -1;//, zsuff = -1;
    vector<int> idx;
    node() {
        fill(nx, nx + C, -1);
    }
};
vector<node> t(1);
void add_word(string& s, int id) {
    int v = 0;
    for (char ch : s) {
        int x = ch - 'a';
        if (t[v].nx[x] == -1) {
            t[v].nx[x] = sz(t);
            t.emplace_back();
        }
        v = t[v].nx[x];
    }
    t[v].idx.emplace_back(id);
}
void build_aho() {
    vector<pair<int, int>> q;
    for (int x = 0; x < C; ++x) {
        if (t[0].nx[x] == -1) {
            t[0].nx[x] = 0;
        } else {
            q.emplace_back(0, x);
        }
    }
    for (int st = 0; st < sz(q); ++st) {
        auto [par, x] = q[st];
        int a = t[par].nx[x];
        if (t[par].suff == -1) {
            t[a].suff = 0;
        } else {
            t[a].suff = t[t[par].suff].nx[x];
            // t[a].zsuff = t[t[a].suff].idx.empty() ? t[t[a].suff
            //    ].zsuff : t[a].suff;
        }
        for (int y = 0; y < C; ++y) {
            if (t[a].nx[y] == -1) {
                t[a].nx[y] = t[t[a].suff].nx[y];
            } else {
                q.emplace_back(a, y);
            }
        }
    }
```

```
        }
    }
```

Suffix array
Description: Calculates suffix array of the given string.
Time: $\mathcal{O}(n \log n)$

```cpp
vector<int> suffix_array(const auto& s, bool sort_suffixes=true) {
    int n = s.size() + sort_suffixes;
    auto [mn, mx] = minmax_element(s.begin(), s.begin() + n);
    vector<int> c(n), cur(n), sa(n), cnt(max<int>(*mx - *mn + 1, n)
        + 1);
    iota(sa.begin(), sa.end(), 0);
    for (int i = 0; i < n; ++i) {
        c[i] = s[i] - *mn + 1;
    }
    ranges::sort(sa, [&](int i, int j) { return c[i] < c[j]; });
    for (int k = 1; k < n; k <<= 1) {
        ranges::fill(cnt, 0);
        for (int x : c) cnt[x]++;
        for (int i = 1; i < cnt.size(); ++i) cnt[i] += cnt[i - 1];
        for (int i : sa) {
            i -= k;
            if (i < 0) i += n;
            int c2 = c[i] - 1;
            cur[cnt[c2]++] = i;
        }
        swap(cur, sa);
        int x = -1, y = -1, p = 0;
        for (int i : sa) {
            int nxt = i + k;
            if (nxt >= n) nxt -= n;
            if (c[i] != x || c[nxt] != y) {
                x = c[i], y = c[nxt], p++;
            }
            cur[i] = p;
        }
        swap(cur, c);
    }
    return sa;
}
```

LCP
Description: Calculates LCP array of the given string and its suffix array.
Time: $\mathcal{O}(n)$

```cpp
vector<int> find_lcp(const string& s, const vector<int>& sa) {
```

```cpp
    int n = sa.size();
    vector<int> pos(n), lcp(n);
    for (int i = 0; i < n; ++i) pos[sa[i]] = i;
    int l = 0;
    for (int i = 0; i < n; ++i) {
        if (pos[i] == n - 1) {
            l = 0;
        } else {
            while (s[(i + l) % n] == s[(sa[pos[i] + 1] + l) % n])
                ++l;
            lcp[pos[i]] = l;
            l = max(0, l - 1);
        }
    }
    return lcp;
}
```

Linear suffix array
Description: Calculates suffix array of the given string.
Time: Fastest

```cpp
#define For(i, j, k) for ( int i = (j) ; i <= (k) ; i++ )
#define Fol(i, j, k) for ( int i = (j) ; i >= (k) ; i-- )

namespace SA {
    template <class IT> inline void induced_sort(int n, int m, IT s
        , int o,
                                                    int *val, int *
                                                      ty, int *cnt
                                                      , int *sa) {
        int *c = cnt + m + 1; fill(sa, sa + n + 1, 0), copy(cnt, c,
            c);
        Fol(i, o, 0) sa[--c[s[val[i]]]] = val[i];
        copy(cnt, cnt + m, c + 1);
        For(i, 0, n) if ( sa[i] && !ty[sa[i] - 1] ) sa[c[s[sa[i] -
            1]]++] = sa[i] - 1;
        copy(cnt, c, c);
        Fol(i, n, 0) if ( sa[i] &&  ty[sa[i] - 1] ) sa[--c[s[sa[i]
            - 1]]] = sa[i] - 1;
    }
    template <class IT> inline bool lms_equal(int x, int y, IT s,
        int *ty) {
        if ( s[x] != s[y] ) return false;
        while ( s[++x] == s[++y] && ty[x] == ty[y] ) if ( ty[x] ==
            2 ) return true;
        return false;
```

```cpp
    }
    template <class IT> inline void sa_is(int n, int m, IT s, int *
        ty, int *lms, int *cnt, int *sa) {
        int o = -1, c = -1, *t = sa;
        ty[n] = 1; Fol(i, n - 1, 0) ty[i] = s[i] == s[i + 1] ? ty[i
            + 1] : s[i] < s[i + 1];
        fill(cnt, cnt + m + 1, 0); For(i, 0, n) cnt[s[i]]++;
            partial_sum(cnt, cnt + m + 1, cnt);
        For(i, 1, n) if ( !ty[i - 1] && ty[i] ) ty[i] = 2, lms[++o]
            = i;
        induced_sort(n, m, s, o, lms, ty, cnt, sa);
        For(i, 0, n) if ( ty[sa[i]] == 2 ) *t++ = sa[i];
        For(i, 0, o) c += c <= 0 || !lms_equal(sa[i], sa[i - 1], s,
            ty), t[sa[i] >> 1] = c;
        For(i, 0, o) t[i] = t[lms[i] >> 1];
        if ( c < o ) sa_is(o, c, t, ty + n + 1, lms + o + 1, cnt +
            m + 1, sa);
        else For(i, 0, o) sa[t[i]] = i;
        For(i, 0, o) lms[o + 1 + i] = lms[sa[i]];
        induced_sort(n, m, s, o, lms + o + 1, ty, cnt, sa);
    }
    vector<int> suffix_sort(const auto& str) {
        using T = std::decay_t<decltype(str[0])>;
        int n = str.size();
        auto s = str.begin();
        vector<T> t(n + 1);
        auto o = minmax_element(s, s + n); int d = *o.first - 1, m
            = *o.second - d;
        vector<int> ty(2 * (n + m + 2)), lms(n + 1), cnt(2 * (n + m
            + 2)), sa(n + 1);
        For(i, 0, n - 1) t[i] = s[i] - d; t[n] = 0;
        sa_is(n, m, t.data(), ty.data(), lms.data(), cnt.data(), sa
            .data()); For(i, 1, n) sa[i]++;
        for (auto& x : sa) --x;
        ++sa[0];
        return sa;
    }
}
```

Minimal rotation

Description: Rotates the given string until it is lexicographically minimal, returns shift.
Time: $\mathcal{O}(n)$

```cpp
int min_rotation(string& s, int len) {
    s += s;
    int i = 0, ans = 0;
```

```cpp
    while (i < len) {
        ans = i;
        int j = i + 1, k = i;
        while (j < len * 2 && s[k] <= s[j]) {
            if (s[k] < s[j]) {
                k = i;
            } else {
                k += 1;
            }
            j += 1;
        }
        while (i <= k) {
            i += j - k;
        }
    }
    s = s.substr(ans, len);
    return ans;
}
```

Palindromic tree

Description: Maintains palindromic tree with support of adding new symbols to the end.
Time: $\mathcal{O}(n)$

```cpp
template<int C> struct EERTREE {
    int top = 1, last = 0, sz = 2;
    vector<int> str, suff, len;
    vector<array<int, C>> nx;

    EERTREE(int n) : str(n + 1, -1), suff(n + 2)
                    , len(n + 2), nx(n + 2, array<int, C>()) {
        len[1] = -1, suff[0] = 1;
    }
    int get_link(int v) {
        while (str[top - len[v] - 2] != str[top - 1]) v = suff[v];
        return v;
    }
    void add_letter(int c) {
        str[top++] = c;
        last = get_link(last);
        if (!nx[last][c]) {
            len[sz] = len[last] + 2;
            suff[sz] = nx[get_link(suff[last])][c];
            nx[last][c] = sz++;
        }
        last = nx[last][c];
    }
```

```
};
```

Suffix Automaton
Description: Builds suffix automaton
Time: $\mathcal{O}(nC)$

```
const int C = 26;

struct Node {
    inline static int timer = 0;

    array<Node*, C> nxt{};
    Node* par = nullptr;
    Node* suf = nullptr;

    // for multi-string automaton
    int last_used = -1;

    // more fields if you want
    // bool is_terminal = false;
};

Node* root = new Node();

Node* add(Node* a, int ch) {
    Node* b = new Node();
    b->suf = root;
    b->par = a;
    while (a != nullptr) {
        if (a->nxt[ch] == nullptr) {
            a->nxt[ch] = b;
            a = a->suf;
            continue;
        }
        Node* c = a->nxt[ch];
        if (c->par == a) {
            b->suf = c;
            break;
        }
        Node* d = new Node();
        d->suf = c->suf;
        c->suf = d;
        b->suf = d;
        d->par = a;
        d->nxt = c->nxt;
        while (a != nullptr && a->nxt[ch] == c) {
```

```
            a->nxt[ch] = d;
            a = a->suf;
        }
        break;
    }
    return b->par->nxt[ch];
}

void visit_subautomaton(Node* v, auto&& cb = [](){}) {
    if (v->last_used == Node::timer) {
        return;
    }
    v->last_used = Node::timer;
    cb(v);
    visit_subautomaton(v->par, cb);
    visit_subautomaton(v->suf, cb);
}
```

Graphs (7)

SCC
Description: Finds SCC for oriented graph in reverse topsort order.
Time: $\mathcal{O}(n+m)$

```
vector<vector<int>> g;
vector<int> tin, up, color, comp, stk;
int timer = 0, comp_col = 0;

void dfs(int v) {
    color[v] = 1;
    tin[v] = timer++;
    up[v] = tin[v];
    stk.push_back(v);
    for (int to : g[v]) {
        if (color[to] == 0) {
            dfs(to);
            up[v] = min(up[v], up[to]);
        } else if (color[to] == 1) {
            up[v] = min(up[v], tin[to]);
        } else if (comp[to] == -1) {
            up[v] = min(up[v], up[to]);
        }
    }
    if (up[v] >= tin[v]) {
        while (true) {
```

```
            int u = stk.back();
            stk.pop_back();
            comp[u] = comp_col;
            if (u == v) break;
        }
        ++comp_col;
    }
    color[v] = 2;
}
```

Directed MST

Description: Finds a minimum spanning tree/arborescence of a directed graph, given a root
node. If no MST exists, returns -1.
Time: $\mathcal{O}(E \log V)$

```
struct RollbackUF {
    vi e; vector<pii> st;
    RollbackUF(int n) : e(n, -1) {}
    int size(int x) { return -e[find(x)]; }
    int find(int x) { return e[x] < 0 ? x : find(e[x]); }
    int time() { return sz(st); }
    void rollback(int t) {
        for (int i = time(); i --> t;)
            e[st[i].first] = st[i].second;
        st.resize(t);
    }
    bool join(int a, int b) {
        a = find(a), b = find(b);
        if (a == b) return false;
        if (e[a] > e[b]) swap(a, b);
        st.push_back({a, e[a]});
        st.push_back({b, e[b]});
        e[a] += e[b]; e[b] = a;
        return true;
    }
};

struct Edge { int a, b; ll w; };
struct Node {
    Edge key;
    Node *l, *r;
    ll delta;
    void prop() {
        key.w += delta;
        if (l) l->delta += delta;
        if (r) r->delta += delta;
```

```
        delta = 0;
    }
    Edge top() { prop(); return key; }
};
Node *merge(Node *a, Node *b) {
    if (!a || !b) return a ?: b;
    a->prop(), b->prop();
    if (a->key.w > b->key.w) swap(a, b);
    swap(a->l, (a->r = merge(b, a->r)));
    return a;
}
void pop(Node*& a) { a->prop(); a = merge(a->l, a->r); }

pair<ll, vi> dmst(int n, int r, vector<Edge>& g) {
    RollbackUF uf(n);
    vector<Node*> heap(n);
    for (Edge e : g) heap[e.b] = merge(heap[e.b], new Node{e});
    ll res = 0;
    vi seen(n, -1), path(n), par(n);
    seen[r] = r;
    vector<Edge> Q(n), in(n, {-1,-1}), comp;
    deque<tuple<int, int, vector<Edge>>> cycs;
    rep(s,0,n) {
        int u = s, qi = 0, w;
        while (seen[u] < 0) {
            if (!heap[u]) return {-1,{}};
            Edge e = heap[u]->top();
            heap[u]->delta -= e.w, pop(heap[u]);
            Q[qi] = e, path[qi++] = u, seen[u] = s;
            res += e.w, u = uf.find(e.a);
            if (seen[u] == s) {
                Node* cyc = 0;
                int end = qi, time = uf.time();
                do cyc = merge(cyc, heap[w = path[--qi]]);
                while (uf.join(u, w));
                u = uf.find(u), heap[u] = cyc, seen[u] = -1;
                cycs.push_front({u, time, {&Q[qi], &Q[end]}});
            }
        }
        rep(i,0,qi) in[uf.find(Q[i].b)] = Q[i];
    }

    for (auto& [u,t,comp] : cycs) { // restore sol (optional)
        uf.rollback(t);
        Edge inEdge = in[u];
        for (auto& e : comp) in[uf.find(e.b)] = e;
```

```
        in[uf.find(inEdge.b)] = inEdge;
    }
    rep(i,0,n) par[i] = in[i].a;
    return {res, par};
}
```

Link-Cut
Description: Represents a forest of unrooted trees. You can add and remove edges (as long as the result is still a forest), and check whether two nodes are in the same tree.
Time: All operations take amortized $\mathcal{O}(\log N)$.

```
struct Node { // Splay tree. Root's pp contains tree's parent.
    Node *p = 0, *pp = 0, *c[2];
    bool flip = 0;
    Node() { c[0] = c[1] = 0; fix(); }
    void fix() {
        if (c[0]) c[0]->p = this;
        if (c[1]) c[1]->p = this;
        // (+ update sum of subtree elements etc. if wanted)
    }
    void pushFlip() {
        if (!flip) return;
        flip = 0; swap(c[0], c[1]);
        if (c[0]) c[0]->flip ^= 1;
        if (c[1]) c[1]->flip ^= 1;
    }
    int up() { return p ? p->c[1] == this : -1; }
    void rot(int i, int b) {
        int h = i ^ b;
        Node *x = c[i], *y = b == 2 ? x : x->c[h], *z = b ? y : x;
        if ((y->p = p)) p->c[up()] = y;
        c[i] = z->c[i ^ 1];
        if (b < 2) {
            x->c[h] = y->c[h ^ 1];
            y->c[h ^ 1] = x;
        }
        z->c[i ^ 1] = this;
        fix(); x->fix(); y->fix();
        if (p) p->fix();
        swap(pp, y->pp);
    }
    void splay() {
        for (pushFlip(); p; ) {
            if (p->p) p->p->pushFlip();
            p->pushFlip(); pushFlip();
            int c1 = up(), c2 = p->up();
```

```
            if (c2 == -1) p->rot(c1, 2);
            else p->p->rot(c2, c1 != c2);
        }
    }
    Node* first() {
        pushFlip();
        return c[0] ? c[0]->first() : (splay(), this);
    }
};

struct LinkCut {
    vector<Node> node;
    LinkCut(int N) : node(N) {}

    void link(int u, int v) { // add an edge (u, v)
        assert(!connected(u, v));
        makeRoot(&node[u]);
        node[u].pp = &node[v];
    }
    void cut(int u, int v) { // remove an edge (u, v)
        Node *x = &node[u], *top = &node[v];
        makeRoot(top); x->splay();
        assert(top == (x->pp ?: x->c[0]));
        if (x->pp) x->pp = 0;
        else {
            x->c[0] = top->p = 0;
            x->fix();
        }
    }
    bool connected(int u, int v) { // are u, v in the same tree?
        Node* nu = access(&node[u])->first();
        return nu == access(&node[v])->first();
    }
    void makeRoot(Node* u) {
        access(u);
        u->splay();
        if(u->c[0]) {
            u->c[0]->p = 0;
            u->c[0]->flip ^= 1;
            u->c[0]->pp = u;
            u->c[0] = 0;
            u->fix();
        }
    }
    Node* access(Node* u) {
        u->splay();
```

```
        while (Node* pp = u->pp) {
            pp->splay(); u->pp = 0;
            if (pp->c[1]) {
                pp->c[1]->p = 0; pp->c[1]->pp = pp; }
            pp->c[1] = u; pp->fix(); u = pp;
        }
        return u;
    }
};
```

Maximum Clique

Description: Quickly finds a maximum clique of a graph (given as symmetric bitset matrix; self-edges not allowed). Can be used to find a maximum independent set by finding a clique of the complement graph.
Time: Runs in about 1s for n=155 and worst case random graphs (p=.90). Runs faster for sparse graphs.

```
typedef vector<bitset<200>> vb;
struct Maxclique {
    double limit=0.025, pk=0;
    struct Vertex { int i, d=0; };
    typedef vector<Vertex> vv;
    vb e;
    vv V;
    vector<vi> C;
    vi qmax, q, S, old;
    void init(vv& r) {
        for (auto& v : r) v.d = 0;
        for (auto& v : r) for (auto j : r) v.d += e[v.i][j.i];
        sort(all(r), [](auto a, auto b) { return a.d > b.d; });
        int mxD = r[0].d;
        rep(i,0,sz(r)) r[i].d = min(i, mxD) + 1;
    }
    void expand(vv& R, int lev = 1) {
        S[lev] += S[lev - 1] - old[lev];
        old[lev] = S[lev - 1];
        while (sz(R)) {
            if (sz(q) + R.back().d <= sz(qmax)) return;
            q.push_back(R.back().i);
            vv T;
            for(auto v:R) if (e[R.back().i][v.i]) T.push_back({v.i}
                );
            if (sz(T)) {
                if (S[lev]++ / ++pk < limit) init(T);
                int j = 0, mxk = 1, mnk = max(sz(qmax) - sz(q) + 1,
                    1);
```

```
                C[1].clear(), C[2].clear();
                for (auto v : T) {
                    int k = 1;
                    auto f = [&](int i) { return e[v.i][i]; };
                    while (any_of(all(C[k]), f)) k++;
                    if (k > mxk) mxk = k, C[mxk + 1].clear();
                    if (k < mnk) T[j++].i = v.i;
                    C[k].push_back(v.i);
                }
                if (j > 0) T[j - 1].d = 0;
                rep(k,mnk,mxk + 1) for (int i : C[k])
                    T[j].i = i, T[j++].d = k;
                expand(T, lev + 1);
            } else if (sz(q) > sz(qmax)) qmax = q;
            q.pop_back(), R.pop_back();
        }
    }
    vi maxClique() { init(V), expand(V); return qmax; }
    Maxclique(vb conn) : e(conn), C(sz(e)+1), S(sz(C)), old(S) {
        rep(i,0,sz(e)) V.push_back({i});
    }
};
```

Weighted Matroid Intersection

Description: Finds maximum weighed matroid intersection. m.may_be_added(set as a vector of items) - items that may be added to independent set for it to still be independent.
Time: $\mathcal{O}\left(|answer| \cdot n \cdot (m.may\_be\_added + n)\right)$

```
template <class M1, class M2>
vector<int> matroid_intersection(M1&& m1, M2&& m2, int n, auto&&
    cost = [](int) {return 1;}) {
    const ll INF18 = 4'000'000'000'000'000'000;
    vector<int> result;
    vector<ll> potential(n, 0);
    while (true) {
        vector<int> start_zone = m1.may_be_added(result);
        vector<int> end_zone = m2.may_be_added(result);
        vector<vector<int>> g(n);
        for (int i = 0; i < result.size(); ++i) {
            int val = result[i];
            swap(result[i], result.back());
            result.pop_back();
            vector<int> forward_edges = m1.may_be_added(result);
            vector<int> backward_edges = m2.may_be_added(result);
            for (auto x : forward_edges) {
                if (x == val) continue;
```

```cpp
            g[val].push_back(x);
        }
        for (auto x : backward_edges) {
            if (x == val) continue;
            g[x].push_back(val);
        }
        result.push_back(val);
        swap(result[i], result.back());
    }
    set<pair<pair<ll, int>, int>> d;
    vector<pair<ll, int>> dist(n, {INF18, -1});
    vector<int> parent(n, -1);
    vector<char> is_used(n);
    for (auto v : start_zone) {
        dist[v] = {-cost(v) - potential[v], 0};
    }
    while (true) {
        int v = -1;
        for (int i = 0; i < n; ++i) {
            if (is_used[i] || dist[i] == pair{INF18, -1}) {
                continue;
            }
            if (v == -1 || dist[i] < dist[v]) {
                v = i;
            }
        }
        if (v == -1) break;
        is_used[v] = 1;
        for (auto u : g[v]) {
            ll w = potential[v] - potential[u] + (dist[v].
                second & 1 ? cost(v) - cost(u) : 0);
            if (w < 0) {
                exit(1);
            }
            pair<ll, int> u_dist = dist[v];
            ++u_dist.second;
            u_dist.first += w;
            if (dist[u] > u_dist) {
                dist[u] = u_dist;
                parent[u] = v;
            }
        }
    }
    pair<ll, int> best_dist = {1, -1};
    int best_end = -1;
    for (auto v : end_zone) {
```

```cpp
            pair<ll, int> dist_now = dist[v];
            if (dist_now == pair{INF18, -1}) {
                continue;
            }
            dist_now.first += potential[v];
            if (dist_now < best_dist) {
                best_dist = dist_now;
                best_end = v;
            }
        }
        if (best_end == -1) {
            break;
        }
        vector<int> xr;
        while (best_end != -1) {
            xr.push_back(best_end);
            best_end = parent[best_end];
        }
        ranges::sort(xr);
        ranges::sort(result);
        vector<int> new_result;
        new_result.reserve(xr.size() + result.size());
        ranges::merge(xr, result, back_inserter(new_result));
        int sz = 0;
        for (int i = 0; i < new_result.size(); ++i) {
            if (i + 1 != new_result.size() && new_result[i] ==
                new_result[i + 1]) {
                ++i;
                continue;
            }
            new_result[sz++] = new_result[i];
        }
        new_result.resize(sz);
        result = std::move(new_result);
        for (int i = 0; i < n; ++i) {
            potential[i] += dist[i].first;
        }
    }
    return result;
}
```

Dominator Tree
Description: Finds dominators for all vertices in a graph for a given starting vertex.
Usage: Look at example(). Modify rmax if N > 1e6.
Time: $\mathcal{O}\left(M \log N\right)$

```cpp
class Dominator {
public:
    Dominator(vector<vector<int>> h, int s) {
        n = sz(h);
        newId.assign(n, -1);
        g = h;
        timer = 0;
        parent.resize(n);
        dfs(s);
        n = timer;
        vector<vector<int>> tmp(n);
        for (int i = 0; i < sz(newId); ++i) {
            if (newId[i] == -1) {
                continue;
            }
            for (int to : g[i]) {
                tmp[newId[i]].push_back(newId[to]);
            }
        }
        g.swap(tmp);
        gr.resize(n);
        for (int i = 0; i < n; ++i) {
            for (int to : g[i]) {
                gr[to].push_back(i);
            }
        }
        sdom.resize(n);
        label.resize(n);
        dsu.resize(n);
        for (int i = 0; i < n; ++i) {
            label[i] = i;
            dsu[i] = i;
        }
        for (int w = n - 1; w >= 0; --w) {
            sdom[w] = w;
            for (int v : gr[w]) {
                if (v < w) {
                    sdom[w] = min(sdom[w], v);
                } else {
                    pair<int, int> p = getDSU(v);
                    sdom[w] = min(sdom[w], sdom[p.second]);
                }
            }
            if (w) {
                dsu[w] = parent[w];
            }
```

```cpp
        }
        calcIdom();
    }

    void get(vector<int>& res1, vector<int>& res2) const {
        res1 = newId;
        res2 = idom;
    }
private:
    vector<int> newId;
    vector<vector<int>> g, gr;
    vector<int> sdom;
    vector<int> idom;
    int timer;
    const int rmax = 20;
    vector<vector<int>> shifts;
    vector<int> parent;
    vector<int> label;
    vector<int> dsu;
    vector<int> depth;
    int n;

    void dfs(int v) {
        newId[v] = timer++;
        for (int to : g[v]) {
            if (newId[to] == -1) {
                dfs(to);
                parent[newId[to]] = newId[v];
            }
        }
    }

    pair<int, int> getDSU(int v) {
        if (dsu[v] == v) {
            return {-1, -1};
        }
        pair<int, int> p = getDSU(dsu[v]);
        if (p.first == -1) {
            return {v, label[v]};
        }
        int u = p.second;
        if (sdom[u] < sdom[label[v]]) {
            label[v] = u;
        }
        dsu[v] = p.first;
        return {p.first, label[v]};
```

```
    }

    int lca(int u, int v) {
        if (depth[u] > depth[v]) {
            swap(u, v);
        }
        for (int r = rmax - 1; r >= 0; --r) {
            if (depth[v] - (1 << r) >= depth[u]) {
                v = shifts[r][v];
            }
        }
        assert(depth[u] == depth[v]);
        if (u == v) {
            return u;
        }
        for (int r = rmax - 1; r >= 0; --r) {
            if (shifts[r][u] != shifts[r][v]) {
                u = shifts[r][u];
                v = shifts[r][v];
            }
        }
        assert(u != v && shifts[0][u] == shifts[0][v]);
        return shifts[0][u];
    }

    void calcIdom() {
        idom.resize(n);
        depth.resize(n);
        shifts = vector<vector<int>>(rmax, vector<int>(n));
        for (int i = 0; i < n; ++i) {
            if (i == 0) {
                idom[i] = i;
                depth[i] = 0;
            } else {
                idom[i] = lca(parent[i], sdom[i]);
                depth[i] = depth[idom[i]] + 1;
            }
            shifts[0][i] = idom[i];
            for (int r = 0; r + 1 < rmax; ++r) {
                int j = shifts[r][i];
                shifts[r + 1][i] = shifts[r][j];
            }
        }
    }
};
```

```
void example() {
    vector<vector<int>> g(n);
    // ...
    Dominator D(g, starting_vertex);
    vector<int> newId, idom;
    D.get(newId, idom);
    for (int i = 0; i < sz(edges); ++i) {
        int u = edges[i].first, v = edges[i].second;
        u = newId[u], v = newId[v];
        // ...
    }
}
```

Miscellaneous (8)

8.1   Nesterov's gradient descent

$x^{k+1} = x^k + \tau(x^k - x^{k-1}) - \gamma_k \nabla f(x^k + \tau(x^k - x^{k-1}))$

$\tau \simeq 0.85 \ldots 0.95$

Integrate
Description: Function integration over an interval using Simpson's rule. The error is proportional to h^4.

```
double integrate(double a, double b, auto&& f, int n = 1000) {
    double h = (b - a) / 2 / n, rs = f(a) + f(b);
    for (int i = 1; i < n * 2; ++i) {
        rs += f(a + i * h) * (i & 1 ? 4 : 2);
    }
    return rs * h / 3;
}
```

Fractional binary search
Description: Finds the smallest fraction $p/q \in [0, 1]$ s.t. $f(p/q)$ is true and $p, q \le N$.
Time: $\mathcal{O}(\log N)$

```
struct frac { ll p, q; };

frac fracBS(auto&& f, ll N) {
    bool dir = true, A = true, B = true;
    frac lo{0, 1}, hi{1, 1}; // Set hi to 1/0 to search (0, N]
    if (f(lo)) return lo;
    assert(f(hi));
    while (A || B) {
        ll adv = 0, step = 1;
```

```
        for (int si = 0; step; (step *= 2) >>= si) {
            adv += step;
            frac mid{lo.p * adv + hi.p, lo.q * adv + hi.q};
            if (abs(mid.p) > N || mid.q > N || dir == !f(mid)) {
                adv -= step; si = 2;
            }
        }
        hi.p += lo.p * adv;
        hi.q += lo.q * adv;
        dir = !dir;
        swap(lo, hi);
        A = B; B = !!adv;
    }
    return dir ? hi : lo;
}
```