

Problem 1

*Return a non-consecutive subsequence from an array composing the maximum sum.
If all negative values or null, return an empty array.*

a) Solution

See external file [MaxSet.py](#).

b) Time Complexity: $O(n)$

At its upper bound, function `max_independent_set()` never exceeds $O(n)$.

It creates memo array **cache** of input size n , the size of **nums**.

The procedure loops $n - 2$ times, calling the recurrence relation at $O(1)$ complexity:

```
1 for indice in range(2, len(nums)):  
2   cache[indice] = max(cache[indice - 1], nums[indice] + cache[indice - 2])
```

After solving its subproblems, the code backtracks an entire n length.

At its worst, **indice** never decrements by 2, iterating down to indice 0.

```
1 indice = len(cache) - 1  
2 while indice > -1:  
3   if nums[indice] > 0 and cache[indice] != cache[indice - 1]:  
4     subseq.append(nums[indice])  
5     indice -= 2  
6   else:  
7     indice -= 1
```

Problem 2

Return the power set of set n of distinct numbers.

a) Solution

See external file [PowerSet.py](#).

b) Time Complexity: $O(2^n)$

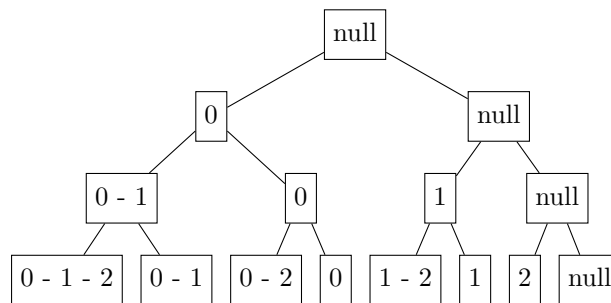
Function `powerset_helper()` is $O(2^n)$. Upper-bound, it generates 2^n subsets at each subproblem.

```

1 if pointer > len(input) - 1:
2     result.append(choices\_made[:])
3     return
4 choices\_made.append(input[pointer])
5 powerset\_helper(pointer + 1, choices\_made, input, result)
6 choices\_made.pop()
7 powerset\_helper(pointer + 1, choices\_made, input, result)
  
```

The code above starts with a null subset, recurses to the largest subset, and memoizes it. It breaks down the main array and recurses to a new subset, memoizing it and progressing. There are two choices at each recursion level: add the current pointer item or don't.

The tree shows subsets at each subtree and the present indices.



Root

Depth 1: Add the first indice or don't.

Depth 2: Add the second indice or don't.

Depth 3: Add the third indice or don't.

The code generates 2^n subsets for input size n .