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## MATH 135: SET THEORY SOLUTIONS TO FINAL EXAM

**1a.** State precisely the Power Set Axiom in the language  $\mathcal{L}(\in)$  having only the binary relation symbol  $\in$ .

$$(\forall x)(\exists y)(\forall t)[t \in y \leftrightarrow (\forall u)(u \in t \to u \in x)]$$

**1b.** Give a formal definition of the expression X is a transitive set in the language of set theory  $\mathcal{L}(\in)$  having only the binary relation symbol  $\in$ .

$$(\forall t)(\forall u)[(u \in t \& t \in X) \to u \in X]$$

1c. Give a formal definition of the expression  $\alpha$  is an ordinal (or as we wrote this in class  $\alpha \in \mathbb{ON}$ ) in the language  $\mathcal{L}(\in)$  having only the binary relation symbol  $\in$ .

 $\alpha$  is transitive &  $(\forall x)(\forall y)[(x \in \alpha \& y \in \alpha) \to (x \in y \lor x = y \lor y \in x)] \& \& (\forall x)[x \in \alpha \to x \text{ is transitive}]$ 

**1d.** Give a formal definition of the expression  $Y = \operatorname{ran}(R)$  in the language of set theory  $\mathcal{L}(\in)$  having only the binary relation symbol  $\in$ .

We define first:

$$t = \{x,y\} : \Longleftrightarrow (\forall u)[u \in t \leftrightarrow (u = x \lor u = y)]$$

We then define

$$t = \langle x, y \rangle : \longleftarrow t = \{ \{x\}, \{x, y\} \}$$

Finally,

$$Y = \operatorname{ran}(R) : \iff (\forall t)[t \in Y \leftrightarrow (\exists x)(\langle x, t \rangle \in R)]$$

1e. State precisely the Replacement Axiom Scheme in the language of set theory  $\mathcal{L}(\in)$  having only the binary relation symbol  $\in$ .

Let  $\varphi = \varphi(x, y, t_1, \dots, t_n)$  be a formula of set theory having free variables amongst  $x, y, t_1, \dots, t_n$  and for which the variables A and B do not appear.

This instance of Replacement states:

**1f.** State precisely the Empty Set Axiom in the language of set theory  $\mathcal{L}(\in)$  having only the binary relation symbol  $\in$ .

$$(\exists x)(\forall y)(\neg y \in x)$$

**2.** Let X be any set and  $R \subseteq X \times X$  any relation on X. By the recursion theorem on  $\omega$  there is a unique function  $f: \omega \to \mathcal{P}(X \times X)$  satisfying  $f(0) = I_X \cup R \cup R^{-1}$  and  $f(n^+) = f(n) \circ f(n)$  for all  $n \in \omega$ . Let  $E := \bigcup \operatorname{ran}(f)$ . Show that for any equivalence relation E' on X with  $R \subseteq E'$ , we have  $E \subseteq E'$ .

Proof. We argue by induction that for every  $n \in \omega$  we have  $f(n) \subseteq R'$ . Consider the case of n=0. By hypothesis,  $R \subseteq E'$ . Since E' is an equivalence relation on X and is thus reflexive on X,  $I_X \subseteq E'$ . Finally, since E' is symmetric and  $R \subseteq E'$ , we have  $R^{-1} \subseteq E'$ . Thus,  $f(0) = R \cup I_X \cup R^{-1} \subseteq E'$ . Suppose now that we know  $f(n) \subseteq E'$ . Suppose that  $t \in f(n^+)$ . Since  $f(n^+) = f(n) \circ f(n)$ , there are x, y and z for which  $t = \langle x, z \rangle$ ,  $\langle x, y \rangle \in f(n)$  and  $\langle y, z \rangle \in f(n)$ . Since  $f(n) \subseteq E'$ , we have  $\langle x, y \rangle \in E'$  and  $\langle y, z \rangle \in E'$ . As E' is a transitive relation,  $t = \langle x, z \rangle \in E'$ . Thus,  $f(n^+)$ . From we conclude that  $E = \bigcup \operatorname{ran}(f) \subseteq E'$ , as required.

**3. Show** that for every set X there is some set K so that for every  $x \in X$  one has  $x \prec K$ .

*Proof.* Let  $K := \mathcal{P}(\bigcup X)$ . Consider any  $x \in X$ . Then  $x \subseteq \bigcup X$ , so that  $x \preceq \bigcup X$  and by Cantor's theorem,  $\bigcup X \prec \mathcal{P}(\bigcup X) = K$ . Thus,  $x \prec K$ .

**4. Show** that if X is a nonempty finite set, then X has at least one maximal element (with respect to  $\subseteq$ ).

Proof. We argue by induction on  $\operatorname{card}(X)$  with the case of  $\operatorname{card}(X) = 0$  being trivial. Suppose now that  $\operatorname{card}(X) = n^+$  are we already know the result for n. Fix a bijection  $f: n^+ \to X$  and set X' := f[[n]] and x := f(n). If n = 0, then x is the maximal element. Otherwise,  $X' \neq \emptyset$  so by induction there is some  $y \in X'$  maximal in X'. Observe that because f is a bijection,  $x \neq y$ . So there are two cases to consider:  $y \subset x$  or  $y \not\subset x$ . In the former case, x is maximal in X for if there were some  $z \in X$  with  $x \subset z$ , then necessarily  $z \in X'$  and  $y \subset z$  contradicting maximality of y in X'. In the latter case, y is maximal in X for we know that  $y \not\subset x$  and any other z would have to come form X' and we know that y is maximal in X'.

**5. Show** that there is an **infinite** set N and an **onto** function  $f: \omega \to N$  for which  $f(0) = \emptyset$  and for every  $n \in \omega$  one has  $f(n^+) = \{f(n)\}.$ 

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**6.** Prove that for sets A, B and C, one has  ${}^{C}(A \times B) \approx {}^{C}A \times {}^{C}B$ . (Yes, we talked about this in class. I want to see your detailed proof.)

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7. Recall that for a set X, the product is the set

$$\prod X := \{ f \in \mathcal{P}(X \times \bigcup X) \ : \ (\forall x)[x \in X \to f(x) \in x \} \ .$$

**Show** that if X is a set of nonempty disjoint sets, then  $\bigcup X \prec \prod X$ . (Note: There are two parts to this. You need to establish that  $\bigcup X \preceq \prod X$  and that  $\bigcup X \not\approx \prod X$ .)

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**8. Show** that for every set X there is an inductive set I with  $X \subseteq I$ .

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