§1 TAKE_RISC INTRODUCTION 1

Important: Before reading TAKE_RISC, please read or at least skim the program for GB_GATES.

1. Introduction. This demonstration program uses graphs constructed by the *risc* procedure in the GB_GATES module to produce an interactive program called take_risc, which multiplies and divides small numbers the slow way—by simulating the behavior of a logical circuit, one gate at a time.

The program assumes that UNIX conventions are being used. Some code in sections listed under 'UNIX dependencies' in the index might need to change if this program is ported to other operating systems.

To run the program under UNIX, say 'take_risc \(\text{trace} \)', where \(\text{trace} \) is nonempty if and only if you want the machine computations to be printed out.

The program will prompt you for two numbers, and it will use the simulated RISC machine to compute their product and quotient. Then it will ask for two more numbers, and so on.

2. Here is the general layout of this program, as seen by the C compiler:

```
#include "gb_graph.h"
                             /* the standard GraphBase data structures */
#include "gb_gates.h"
                             /* routines for gate graphs */
  ⟨ Preprocessor definitions ⟩
  (Global variables 3)
  main(argc, argv)
                    /* the number of command-line arguments */
      int argc;
                      /* an array of strings containing those arguments */
                                  /* we'll show registers 0-7 if tracing */
    trace = (argc > 1 ? 8 : 0);
    if ((g = risc(8_L)) \equiv \Lambda) {
      printf("Sorry, _I_couldn't_generate_the_graph_(trouble_code_%ld)!\n", panic_code);
      return (-1);
    printf("Welcome_to_the_world_of_microRISC.\n");
    while (1) {
       ⟨ Prompt for two numbers; break if unsuccessful 4⟩;
       (Use the RISC machine to compute the product, p 7);
      printf("The_product_of_%ld_oad_%ld_is_%ld%s.\n", m, n, p, o? "_o(overflow_occurred)": "");
       \langle Use the RISC machine to compute the quotient and remainder, q and r \otimes\rangle;
      printf("The\_quotient\_is\_%ld,\_and\_the\_remainder\_is\_%ld.\n",q,r);
                  /* normal exit */
    return 0;
  }
3. \langle \text{Global variables } 3 \rangle \equiv
  Graph *g;
                 /* graph that defines a simple RISC machine */
                   /* overflow, product, quotient, remainder */
  long o, p, q, r;
                  /* number of registers to trace */
  long trace;
                 /* numbers to be multiplied and divided */
  long m, n;
  char buffer[100]; /* input buffer */
See also section 6.
```

This code is used in section 2.

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```
4. \#define prompt(s)
          { printf(s); fflush(stdout); /* make sure the user sees the prompt */
            if (fgets(buffer, 99, stdin) \equiv \Lambda) break; }
\langle Prompt \text{ for two numbers; } \mathbf{break} \text{ if unsuccessful } 4 \rangle \equiv
  prompt("\nGimme_a_number:_\");
step 0:
  if (sscanf(buffer, "%ld", \&m) \neq 1) break;
step 1:
  if (m < 0) {
     prompt(\texttt{"Excuse\_me,\_I\_meant\_a\_positive\_number:\_"});
     if (sscanf(buffer, "%ld", \&m) \neq 1) break;
     if (m \le 0) break;
  while (m > {}^{\#}7fff) {
     prompt("That_number's_too_big; _please_try_again: _");
     if (sscanf(buffer, "%ld", \&m) \neq 1) goto step\theta; /* step\theta will break out */
     if (m \le 0) goto step1;
  \langle Now do the same thing for n instead of m 5\rangle;
This code is used in section 2.
5. (Now do the same thing for n instead of m \ 5) \equiv
  prompt("OK, _now_gimme_another: _");
  if (sscanf(buffer, "%ld", &n) \neq 1) break;
step 2:
  if (n \leq 0) {
     prompt("Excuse\_me, \_I\_meant\_a\_positive\_number: \_");
     if (sscanf(buffer, "%ld", &n) \neq 1) break;
     if (n \le 0) break;
  while (n > #7fff) {
     prompt(\texttt{"That} \verb"unumber's \verb"utoo \verb"big; \verb"uplease \verb"try \verb"again: \verb"");
     if (sscanf(buffer, "%ld", \&n) \neq 1) goto step\theta; /* step\theta will break out */
     if (n < 0) goto step2;
This code is used in section 4.
```

 $\S6$ Take_risc a risc program 3

6. A RISC program. Here is the little program we will run on the little computer. It consists mainly of a subroutine called tri, which computes the value of the ternary operation $x \lfloor y/z \rfloor$, assuming that $y \geq 0$ and z > 0; the inputs x, y, z appear in registers 1, 2, 3, respectively, and the exit address is assumed to be in register 7. As special cases we can compute the product xy (letting z = 1) or the quotient $\lfloor y/z \rfloor$ (letting z = 1). When the subroutine returns, it leaves the result in register 4; it also leaves the value $(y \mod z) - z$ in register 2. Overflow will be set if and only if the true result was not between -2^{15} and $2^{15} - 1$, inclusive.

It would not be difficult to modify the code to make it work with unsigned 16-bit numbers, or to make it deliver results with 32 or 48 or perhaps even 64 bits of precision.

```
/* location 'div' in the program below */
#define div 7
#define mult 10
                      /* location 'mult' in the program below */
\#define memry\_size 34
                             /* the number of instructions in the program below */
\langle \text{Global variables } 3 \rangle + \equiv
  unsigned long memry[memry\_size] = \{ /* a "read-only memory" used by run\_risc */
  #2ff0.
             /* start: r2 = m (contents of next word) */
  #1111,
             /* (we will put the value of m here, in memry[1]) */
  #1a30,
                  r1 = n (contents of next word) */
             /* (we will put the value of n here, in memry[3]) */
  #3333.
                   jumpto (contents of next word), r? = return address */
  #7f70,
             /* (we will put either mult or div here, in memry[5]) */
  #5555,
  #0f8f,
             /* halt without changing any status bits */
  #3a21,
             /* \ div: \ r3 = r1 \ */
  #1a01,
             /*
                   r1 = 1 */
  #0a12,
             /*
                   goto tri (literally, r\theta += 2) */
  #3a01,
             /* mult: r3 = 1 */
  #4000.
             /* tri: r_4 = 0 */
             /*
  <sup>#</sup>5000,
                 r5 = 0 */
                   r6 = 0 */
  #6000,
             /*
             /*
                   r2 - = r3 */
  #2a63,
  #0f95,
             /*
                   goto l2 */
  #3063,
             /* l1: r3 \ll = 1 */
  #1061,
             /*
                   r1 \ll = 1 */
                   if (overflow) r6 = 1 */
  #6ac1.
             /*
                   r5 ++ */
  #5fd1.
  #2a63,
             /* l2: r2 -= r3 */
  #039b.
             /*
                   if (\geq 0) goto l1 */
  #0843.
             /*
                   goto 14 */
  #3463,
             /* l3: r3 \gg = 1 */
  #1561,
             /* r1 \gg = 1 */
             /* l4: r2 += r3 */
  #2863,
  #0c94,
                  if (< 0) goto l5 */
  #4861,
             /*
                  r_4 + = r_1 */
             /*
  #6ac1,
                   if (overflow) r\theta = 1 */
  #2a63,
             /*
                   r2 - = r3 */
             /* l5: r5 -- */
  #5a41.
             /*
                   if (\geq 0) goto l3 */
  #0398,
  <sup>#</sup>6666,
                   if (r6) force overflow (literally r6 \gg = 4) */
  #0fa7};
                   return (literally, r\theta = r7, preserving overflow) */
```

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```
7. \( \) Use the RISC machine to compute the product, \( p \) 7 \\ \) \( = \) \( memry[1] = m; \) \( memry[3] = n; \) \( memry[5] = mult; \) \( run_risc(g, memry, memry_size, trace); \) \( p = (\long) \) \( risc_state[4]; \) \( o = (\long) \) \( risc_state[16] \& 1; \) \/* the overflow bit \( */ \)
8. \( \) Use the RISC machine to compute the quotient and remainder, \( q \) and \( r \) 8 \\ \( = \) \( memry[5] = \) \( div; \) \( run_risc(g, memry, memry_size, trace); \) \( q = (\long) \) \( risc_state[4]; \) \( r = ((\long)(risc_state[2] + n)) \\ \& \( #7fff; \)
This code is used in section 2.
```

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9. Index. Finally, here's a list that shows where the identifiers of this program are defined and used.

 $argc: \underline{2}.$ $argv: \underline{2}.$ buffer: $\underline{3}$, 4, 5. $div: \underline{6}, 8.$ fflush: 4.fgets: 4.g: $\underline{3}$. Graph: 3. $l1: \underline{6}.$ $l2: \underline{6}.$ *l3*: 6. $l4: \underline{6}.$ $l5: \underline{6}.$ $m: \underline{3}.$ main: $\underline{2}$. memry: $\underline{6}$, 7, 8. $memry_size$: $\underline{6}$, 7, 8. $mult: \underline{6}, 7.$ n: $\underline{3}$. o: $\underline{3}$. p: $\underline{3}$. $panic_code$: 2. printf: 2, 4.prompt: $\underline{4}$, 5. q: $\underline{3}$. r: $\underline{3}$. risc: 1, 2. $risc_state$: 7, 8. run_risc: 6, 7, 8. sscanf: 4, 5. $start: \underline{6}.$ stdin: 4. stdout: 4. $step\theta: \underline{4}, 5.$ $step 1: \underline{4}.$ $step 2: \underline{5}.$ trace: $2, \underline{3}, 7, 8.$ $tri: \underline{6}.$

UNIX dependencies: 2.

6 NAMES OF THE SECTIONS TAKE_RISC

TAKE_RISC

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Preliminary work on the Stanford Graph Base project was supported in part by National Science Foundation grant CCR-86-10181.