



# Recursion for decision problems

# Knapsack problem

**Given** : a set  $S$  of  $n$  objects each with a weight  $w_i$  and value  $v_i$ ,  $0 \leq i \leq n - 1$  and a knapsack with weight capacity  $W$

**Question** : What is the maximum value that can be obtained by packing the knapsack with objects from  $S$  such that the sum of weights of chosen objects does not exceed  $W$  ?

- To answer this question you don't need to examine every possible subset of  $S$
- **Optimality principle** : Suppose in an optimal solution, after choosing objects from subset of first  $m$  objects, capacity  $C$  remains. Then the choices for the remaining  $n-m$  objects must be an optimal solution for the subproblem of knapsack capacity  $C$ .

# Knapsack examples

- Can we use a greedy approach ?
- Consider weights =  $\{6, 4, 2, 1\}$  and  
values =  $\{6, 4, 5, 3\}$  and  
capacity of knapsack  $W = 6$
- Fill objects by weights, smallest first.  
weights chosen =  $\{1, 2\}$ , total value =  $3 + 5 = 8$
- Fill objects by values, largest first.  
weights chosen =  $\{6\}$ , total value = 6
- Best solution :  
Weights chosen =  $\{4, 2\}$ , total value =  $4 + 5 = 9$

# Recursion subproblem for knapsack

- **Given:** A subset of remaining objects indexed from  $k$  (i.e.  $k, k+1, \dots, n-1$ ) and remaining capacity  $c$  of knapsack where  $0 \leq c \leq W$ ,
- **Needed :** What is the maximum value  $F_k(c)$  that can be obtained ?

Answer to original problem : Set  $k = 0$  and  $c = W$

**Base step:**

$F_n(c) = 0$  (as no more objects available)

**Recursion step ( $k=0, 1, 2, \dots, n-1$ ):**

$F_k(c) = F_{k+1}(c)$  if  $w_i > c$  (object cannot be chosen)

$F_k(c) = \max(v_i + F_{k+1}(c - w_i), F_{k+1}(c))$  otherwise

i.e. choose between use it or lose it decisions

Need  $F_0(W)$  to answer the original problem.

```
public static int knapsack(int [] weights, int [] values, int start, int
capacity) {
    if (start == weights.length) {
        return 0;
    }
    if (weights[start] > capacity) {
        return knapsack(weights, values, start+ 1, capacity);
    }
    return Math.max(knapsack(weights, values, start+ 1, capacity),
        values[start] + knapsack(weights, values, start+ 1, capacity-
weights[start]));
}
```

```
public static int knapsack(int [] weights, int [] values, int capacity) {
    return knapsack(weights, values, 0, capacity);
}
```

Some of recursive calls starting from  $K=0, C = 6$

Order of calls shown with alphabets, dotted lines show call returns, numbers show values returned.

Remaining knapsack capacity  $C \rightarrow$

