

#### **Profitable Bandits**

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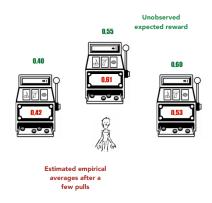
## The classical Multi-Armed Bandit problem

At each time t = 1, ..., T

- ▶ Pull an arm  $a_t \in \{1, ..., K\}$
- Receive reward  $X_{a_t,t} \sim \nu_{a_t}$

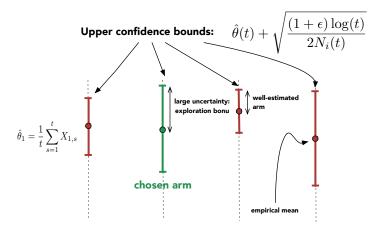
Goal: maximize  $\mathbb{E}[\sum_{t=1}^{T} X_{a_t,t}]$ 

Dilemma: exploration vs exploitation



# A successful approach: UCB algorithm (Auer et al., 2002)

- ▶ Initialization: pull each arm once
- ► Then:



## The Profitable Bandit problem

At each time t = 1, ..., T

- ▶ Choose arms  $A_t \subset \{1, ..., K\}$
- ▶ Observe rewards  $X_{a,c,t} \sim \nu_a$  for all  $a \in A_t$ ,  $c \in \{1, ..., C_a(t)\}$

#### Objective

maximize 
$$S_T := \mathbb{E}[\sum_{t=1}^T \sum_{a \in A_t} \sum_{c=1}^{C_a(t)} (X_{a,c,t} - \tau_a)]$$

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Hence, optimal choice:  $A^* = \{a \in \{1, \dots, K\}, \Delta_a > 0\}$  with  $\Delta_a = \mu_a - \tau_a$  and  $\mu_a = \mathbb{E}[X_{a,1,1}]$ .

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Equivalently, minimize the expected regret

$$\begin{split} R_T &= \sum_{a \in A^*} \Delta_a \, \tilde{C}_a(T) - S_T \\ &= \sum_{a \in A^*} \Delta_a \, \left( \, \tilde{C}_a(T) - \mathbb{E}[N_a(T)] \right) + \sum_{a \notin A^*} |\Delta_a| \mathbb{E}[N_a(T)], \end{split}$$

where 
$$\tilde{C}_a(T) = \mathbb{E}[\sum_{t=1}^T C_a(t)], N_a(T) = \sum_{t=1}^T C_a(t)\mathbb{I}\{a \in A_t\}.$$

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#### Lower bound

#### Theorem

If the  $\nu_a$ 's belong to an one-dimensional exponential family, for all uniformly efficient strategies, for all non-profitable arms a such that  $\mu_a < \tau_a$ ,

$$\liminf_{T \to \infty} \frac{\mathbb{E}[N_a(T)]}{\log T} \ge \frac{1}{d(\mu_a, \tau_a)},$$

with d the KL-divergence of the family parametrized by the mean:  $d(\mu_a, \mu_{a'}) = KL(\nu_a, \nu_{a'})$ .

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#### Consequence:

$$R_T \gtrsim \sum_{a \notin A^*} \frac{|\Delta_a|}{d(\mu_a, \tau_a)} \log T.$$

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# **Index policy**

An index policy is fully characterized by the choice of index  $u_a(t)$ .

# Algorithm 1 Generic index policy

**Require:** time horizon T, thresholds  $(\tau_a)_{a \in \{1,...,K\}}$ 

- 1: Pull all arms:  $A_1 = \{1, ..., K\}$
- 2: **for** t = 1 **to** T 1 **do**
- 3: Compute  $u_a(t)$  for all arms  $a \in \{1, ..., K\}$
- 4: Choose  $A_{t+1} \leftarrow \{a \in \{1, \dots, K\}, u_a(t) \ge \tau_a\}$
- 5: end for

# Three index policies

▶ kl-UCB-4P

$$u_a(t) = \sup \left\{ q > \hat{\mu}_a(t) : N_a(t)d(\hat{\mu}_a(t),q) \leq \log t + c \log \log t \right\}$$

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Bayes-UCB-4P

$$u_a(t) = Q(1 - 1/(t(\log t)^c); \lambda_a^{t-1}),$$

with  $\lambda_a^{t-1}$  the posterior distribution on  $\mu_a$  after round t-1.

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► Thompson-Sampling-4P

$$u_a(t) = \mu(\theta_{a,t}),$$

where  $\theta_{a,t} \sim \pi_a^{t-1}$  with  $\pi_a^{t-1}$  the posterior distribution on  $\theta_a$  after round t-1.

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## **Upper bound**

#### **Theorem**

For the three policies defined above (kl-UCB-4P, Bayes-UCB-4P, TS-4P),

$$R_T \leq \sum_{a \notin A^*} \frac{c_a^+}{c_a^-} \frac{|\Delta_a|}{d(\mu_a, \tau_a)} \log T + o(\log \log T),$$

where for all  $t \ge 1$ :  $c_a^- \le C_a(t) \le c_a^+$ .

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where for all  $t \ge 1$ :  $c_a^- \le C_a(t) \le c_a^+$ .

Conclusion: the three algorithms are asymptotically optimal when  $C_a(1) = \cdots = C_a(T)$  for all  $a \notin A^*$ .

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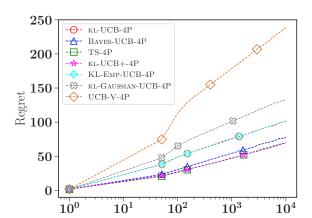
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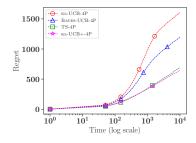
#### Bernoulli distributions

- $T = 10^4$
- $ightharpoonup K = 5 \text{ arms, } \nu_a = \text{Bernoulli}(\mu_a)$
- $\blacktriangleright$  ( $\mu_a$ ,  $\tau_a$ ): (0.1, 0.2), (0.3, 0.2), (0.5, 0.4), (0.5, 0.6), (0.7, 0.8)
- $ightharpoonup C_a(t) 1 \sim \mathsf{Poisson}(a+1)$



# **Poisson and Exponential distributions**

#### Poisson



#### Exponential

