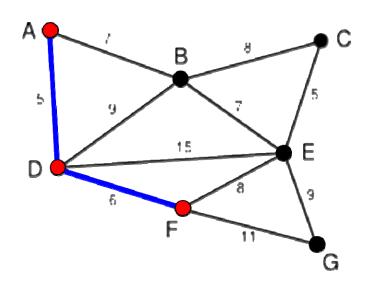
"본 강의 동영상 및 자료는 대한민국 저작권법을 준수합니다. 본 강의 동영상 및 자료는 상명대학교 재학생들의 수업목적으로 제작·배포되는 것이므로, 수업목적으로 내려받은 강의 동영상 및 자료는 수업목적 이외에 다른 용도로 사용할 수 없으며, 다른 장소 및 타인에게 복제, 전송하여 공유할 수 없습니다. 이를 위반해서 발생하는 모든 법적 책임은 행위 주체인 본인에게 있습니다."

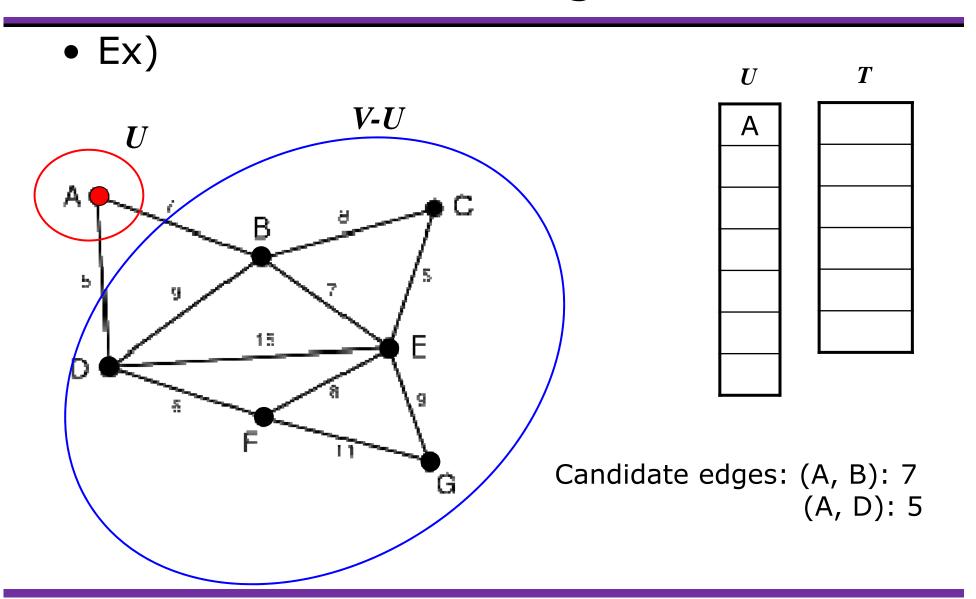
- Vertex-based algorithm
  - Basic strategy
    - Vertices on a Graph is classified into three categories:
      - U → Vertices in minimum-cost spanning tree (T)
      - Vertices incident to U
      - Other vertices

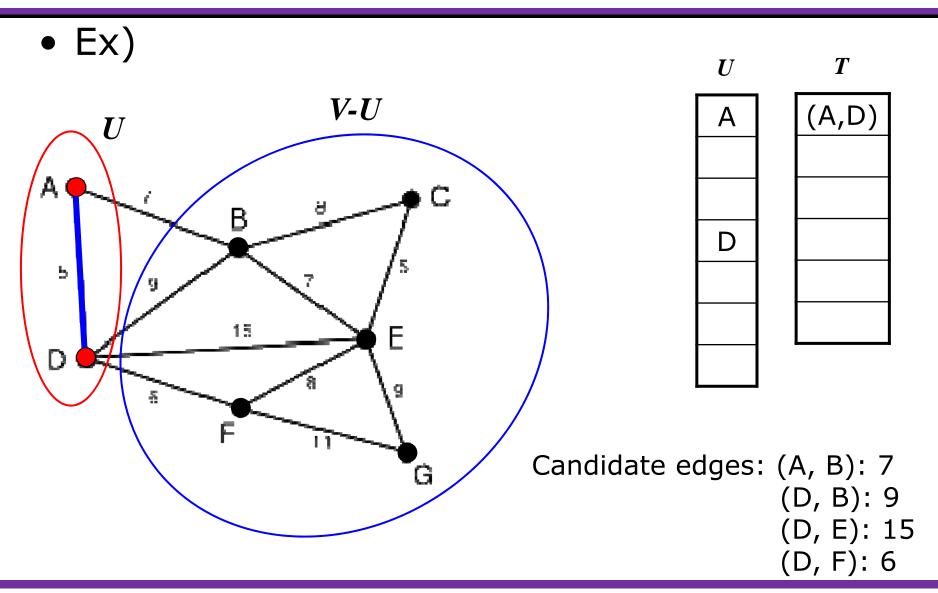


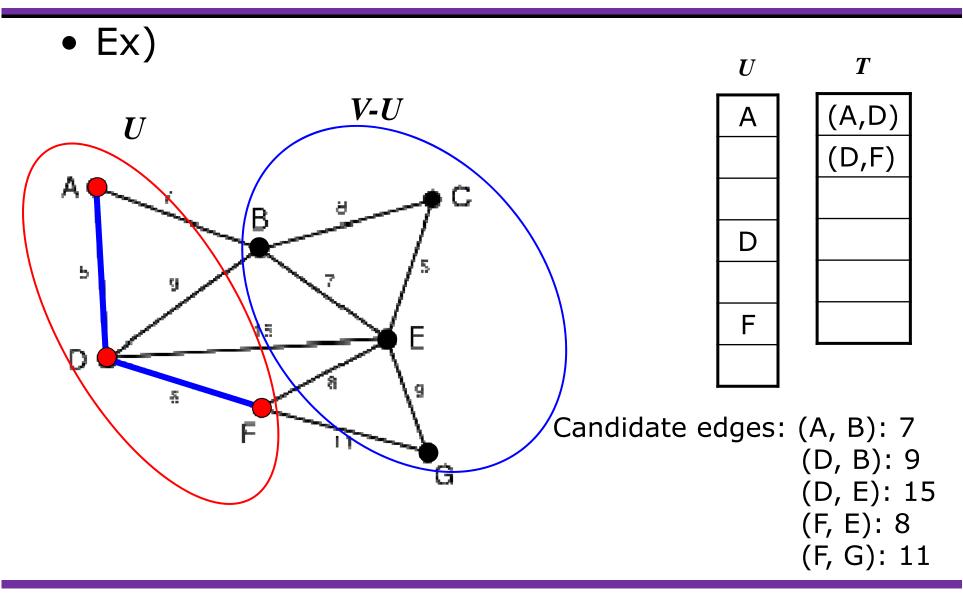
#### Algorithm

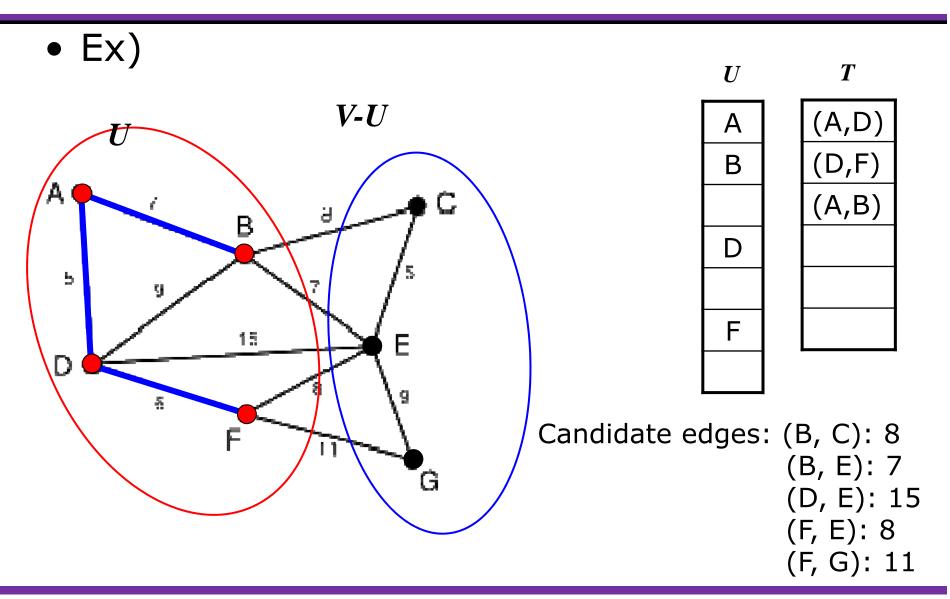
- Initially, set T as Φ.
- Find all the vertices incident to the vertices in T.
- Find the minimum-weight edge among the edges that connect a vertex belongs to T and a vertex that does not belong to T.
- Add the vertex on the edge to T.
- Repeat this process until all vertices belong to U.

```
Tree Prim( Vertex V, Edge E )
   Vertex *U;
    vertex u,v;
    T = \{ \};
    U = \{ A \};
    while (U != V) {
        (u,v) = lowest cost edge with u in U and v in (V - U);
        if ((u,v) is NULL)
            return NULL;
        T += (u, v);
        U += v;
    return T;
```

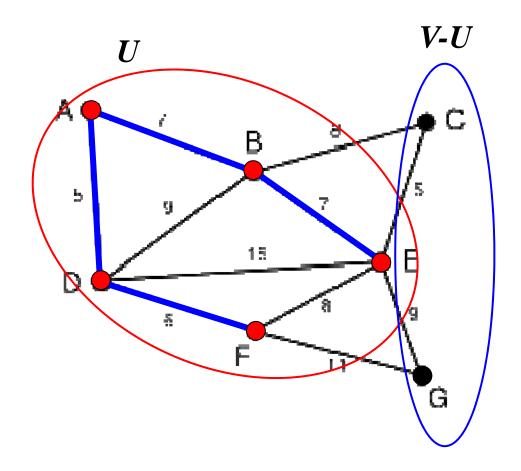


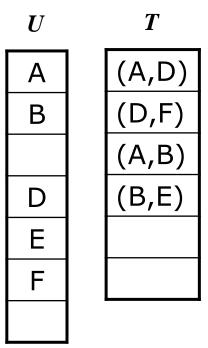








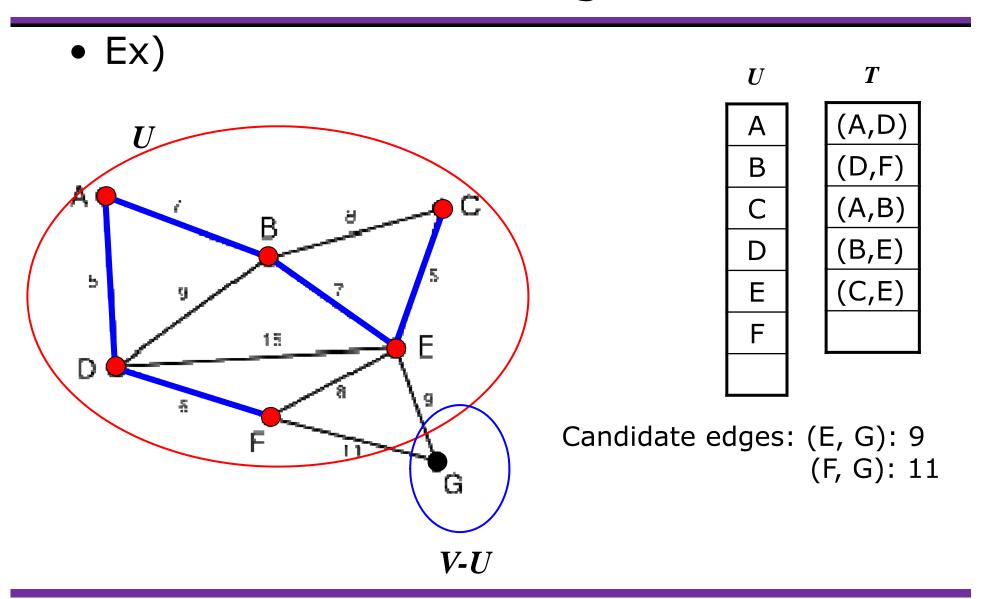


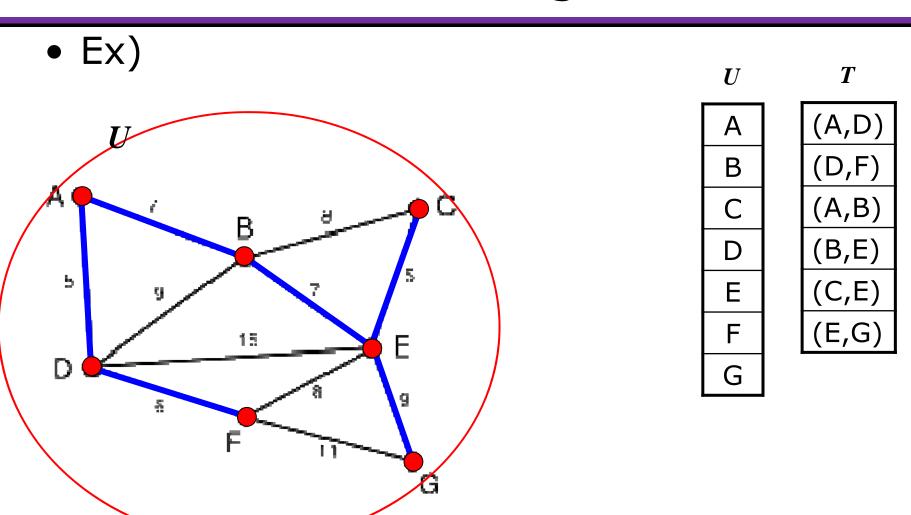


Candidate edges: (B, C): 8

(E, C): 5 (E, G): 9

(F, G): 11

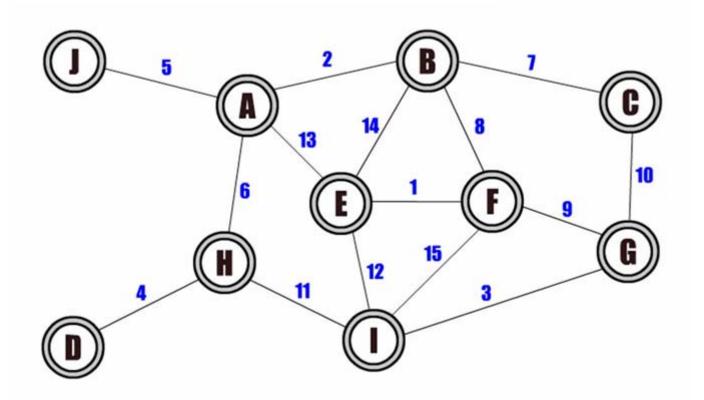


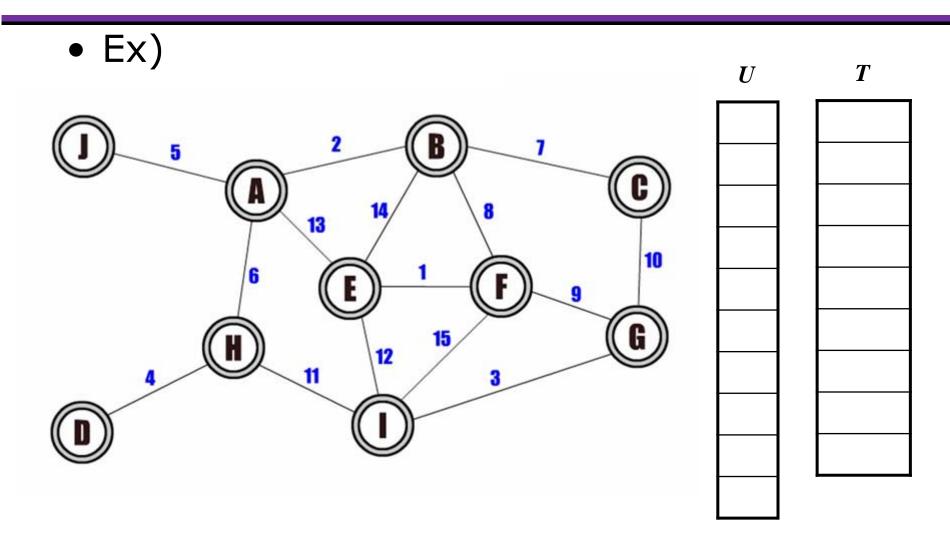


- Performance analysis
  - All edges in the graph are processed  $\rightarrow$  O(m)
  - All edges are managed in a heap  $\rightarrow$  O(m log m)

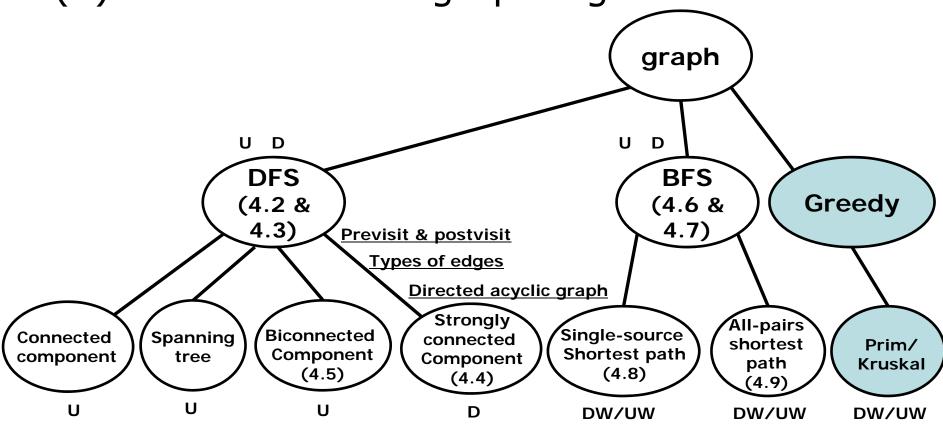
```
Tree Prim( Vertex V, Edge E )
{
    Vertex *U;
    vertex u, v;
    T - { };
    U = { A };
    while (U != V) {
        (u, v) = lowest cost edge with u in U and v in (V - U);
        T += (u, v);
        U |= v;
    }
    return T;
}
```

#### • Ex)





(5) Classification of graph algorithms



# All about graph

Туре	Purpose	Operations	Performance
DFS	Traverse all vertices	Visiting all vertices & visiting all edges	O(n) + O(m)
SCC	Finding SCC	DFS on G <sup>R</sup> and G	O(DFS)
BFS	Traverse all vertices	Visiting all vertices & visiting all edges	O(n) + O(m)
Dijkstra	Single source shortest path	Visiting all edges & managing queue	O(n)*get_min + O(m)*modify $\rightarrow$ O(n <sup>2</sup> ) or O(n log n)
Floyd	All pairs shortest path	Incrementing k	O(n³)
Kruskal (Greedy)	Minimum-cost spanning tree	Sorting edge & selecting (cycle checking)	$O(m log m) + O(n^2)$ $\rightarrow O(m log m)$
Prim (Greedy)	Minimum-cost spanning tree	Managing heap of edges	O(m log m)
MultiStage (Dynamic)			

다음 설명 중 옳은 것을 모두 고르시오.

- (a) Prim 알고리즘은 cycle check를 하지 않기 때문에 Kruskal 알고리즘보다 성능이 더 좋다
- (b) Prim 알고리즘은 우선순위 큐를 이용해서 edge를 관리한다
- (c) Prim 알고리즘은 n > m인 모든 graph에 대해서 NULL을 return한다
- (d) Prim 알고리즘에서 while-loop은 O(m)번 수행된다