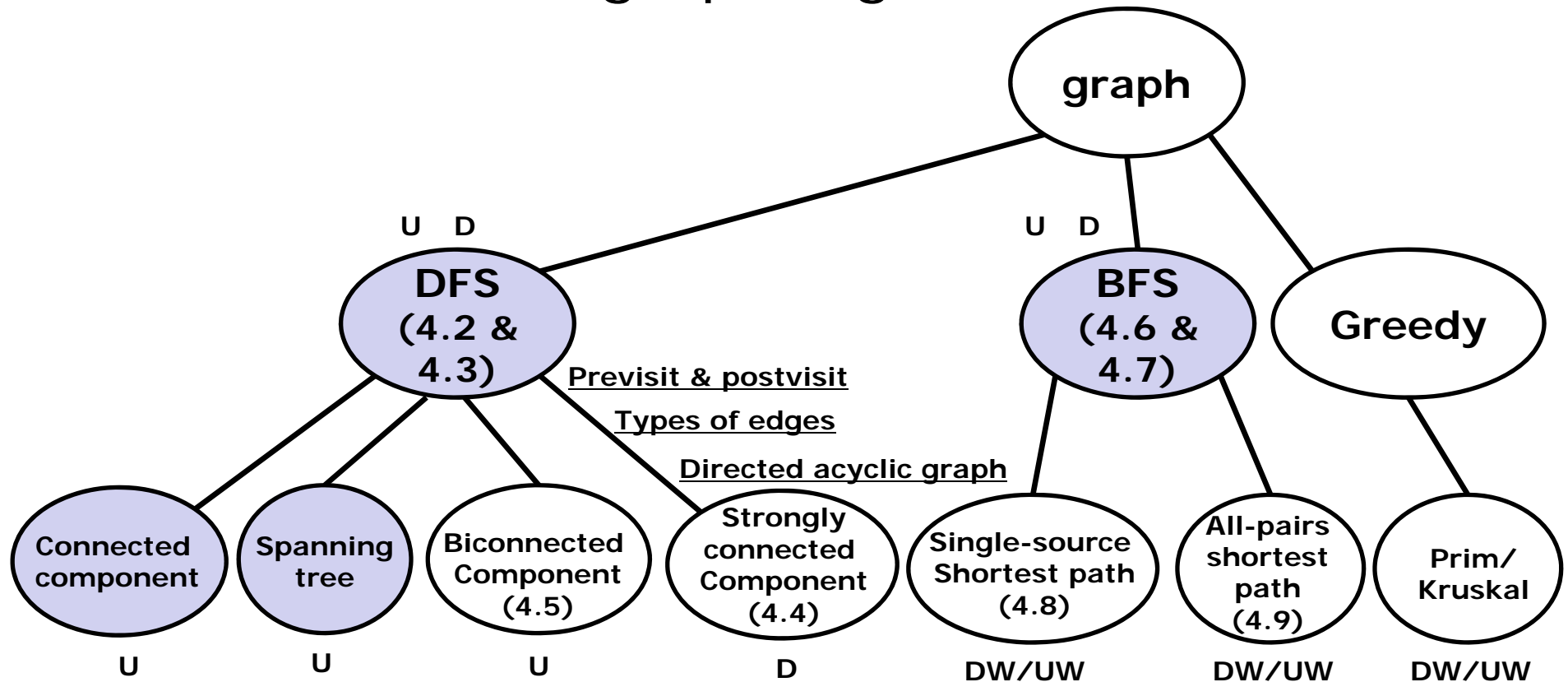

“본 강의 동영상 및 자료는 대한민국 저작권법을 준수합니다. 본 강의 동영상 및 자료는 상명대학교 재학생들의 수업목적으로 제작·배포되는 것이므로, 수업목적으로 내려받은 강의 동영상 및 자료는 수업목적 이외에 다른 용도로 사용할 수 없으며, 다른 장소 및 타인에게 복제, 전송하여 공유할 수 없습니다. 이를 위반해서 발생하는 모든 법적 책임은 행위 주체인 본인에게 있습니다.”

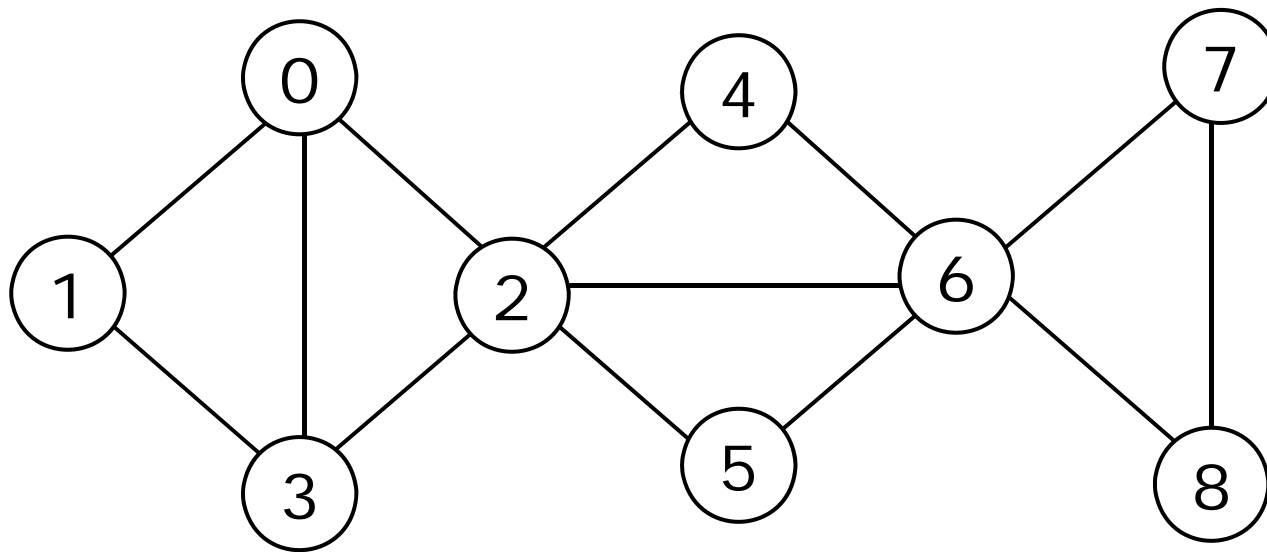
4.5 Biconnected Components

Classification of graph algorithms



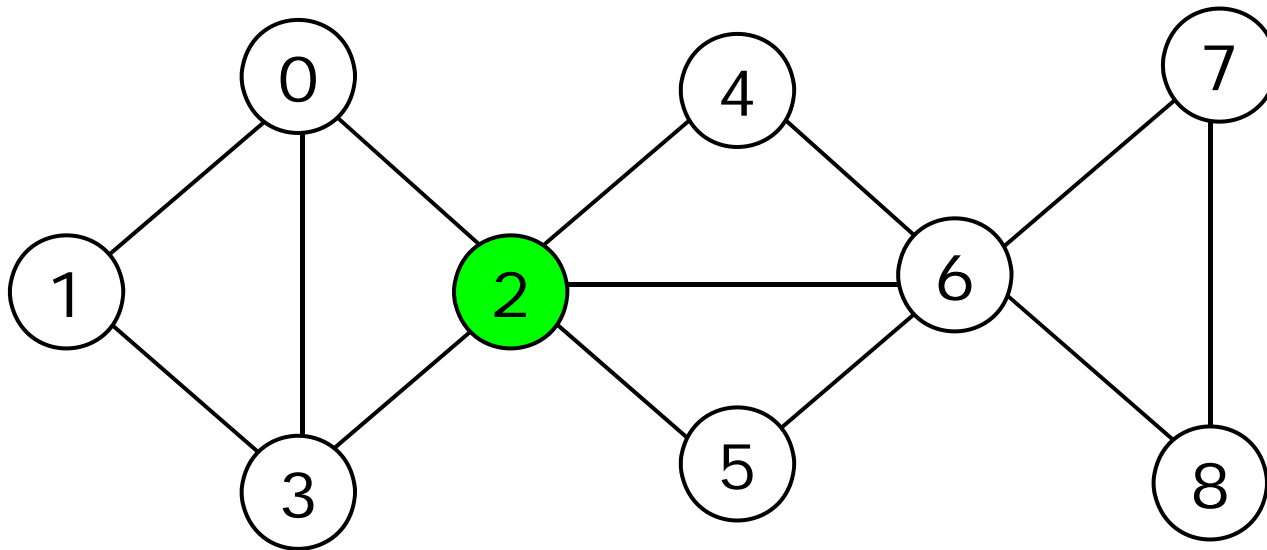
4.5 Biconnected Components

- Articulation point
 - A vertex v of G such that the deletion of v , together with all edges incident to v , produces a graph G' that has at least two connected components



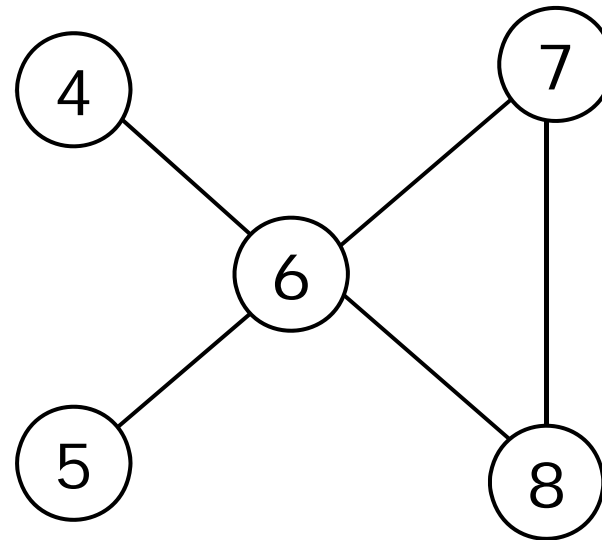
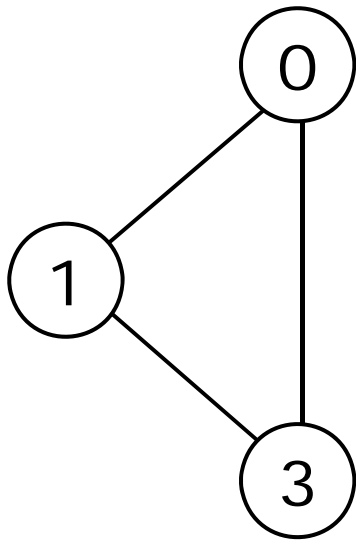
4.5 Biconnected Components

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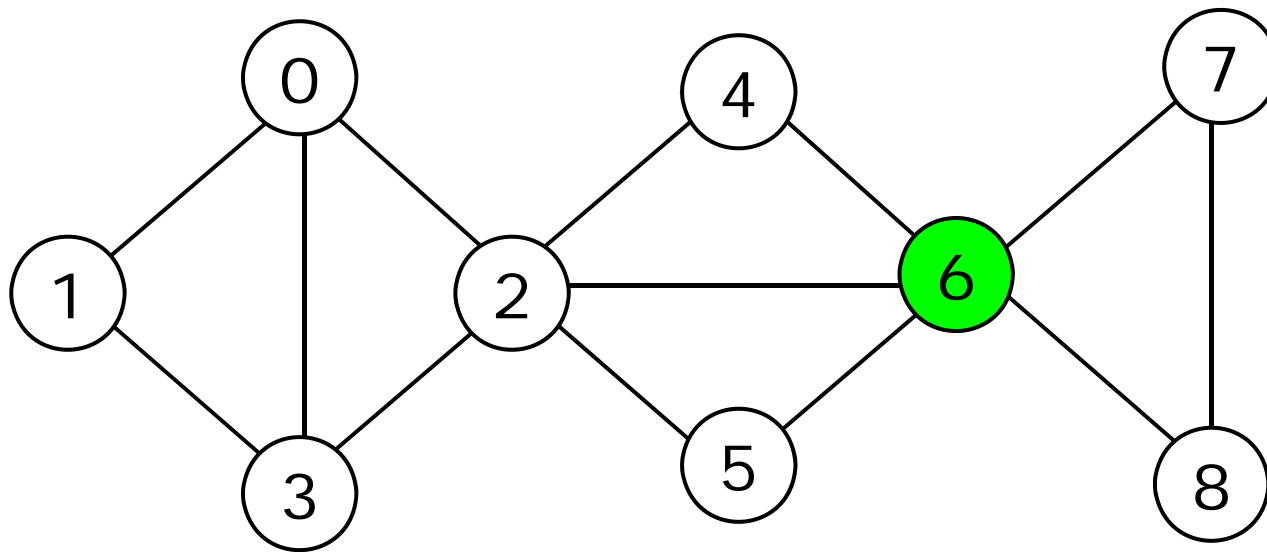
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- Articulation point
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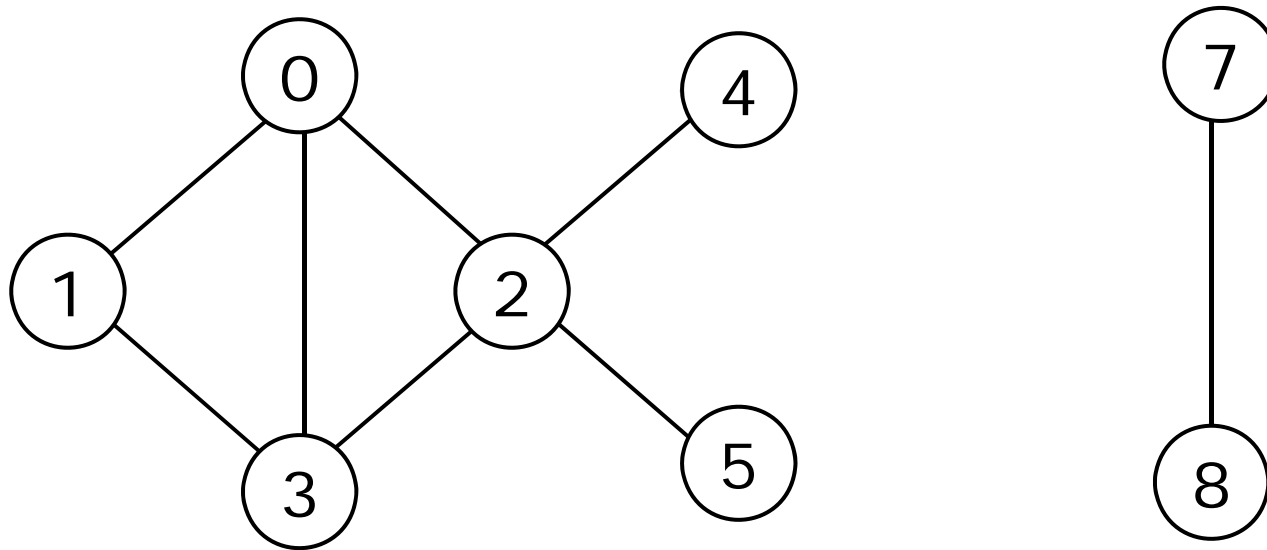
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- Articulation point
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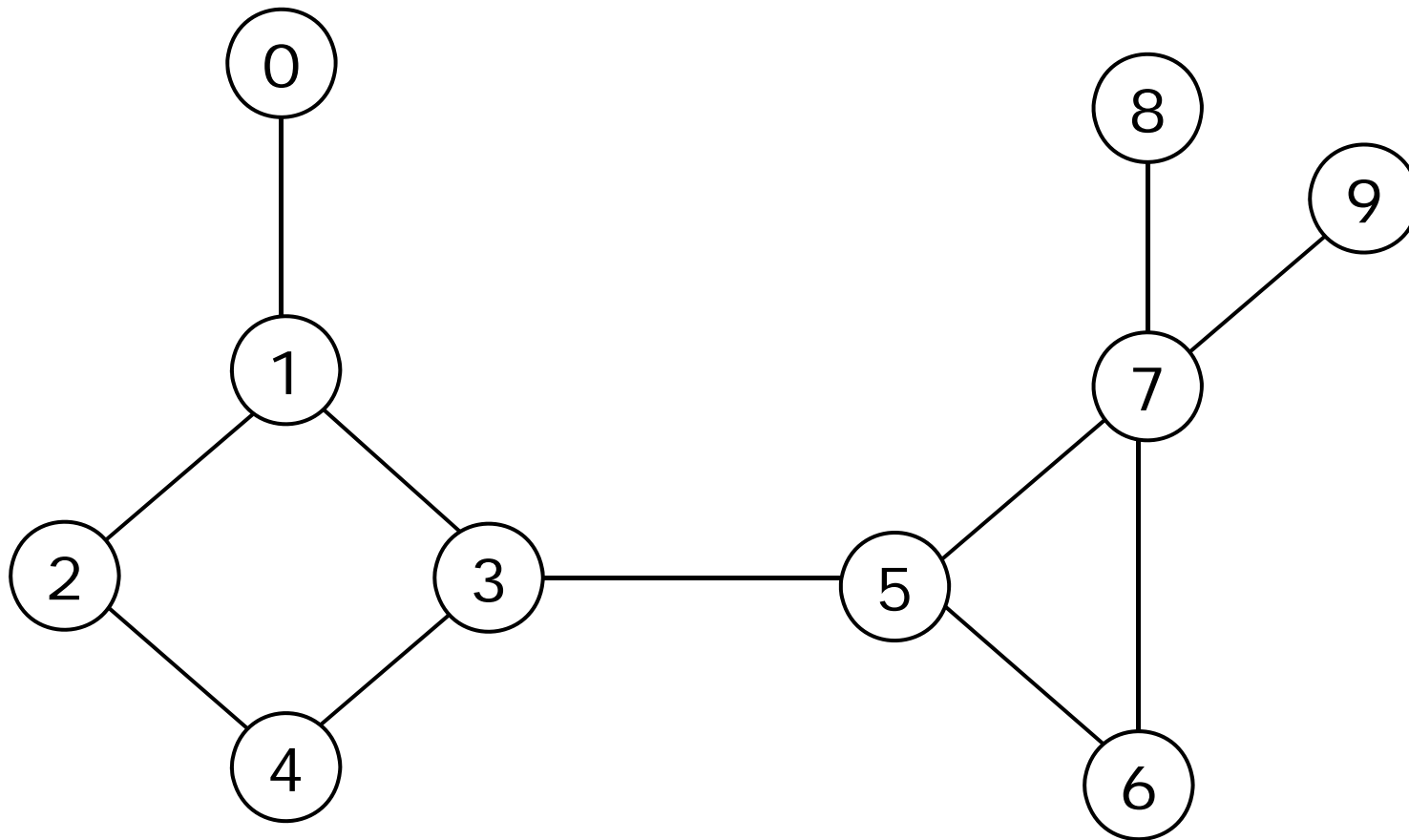
4.5 Biconnected Components

- Articulation point
 - A vertex v of G such that the deletion of v , together with all edges incident to v , produces a graph G' that has at least two connected components



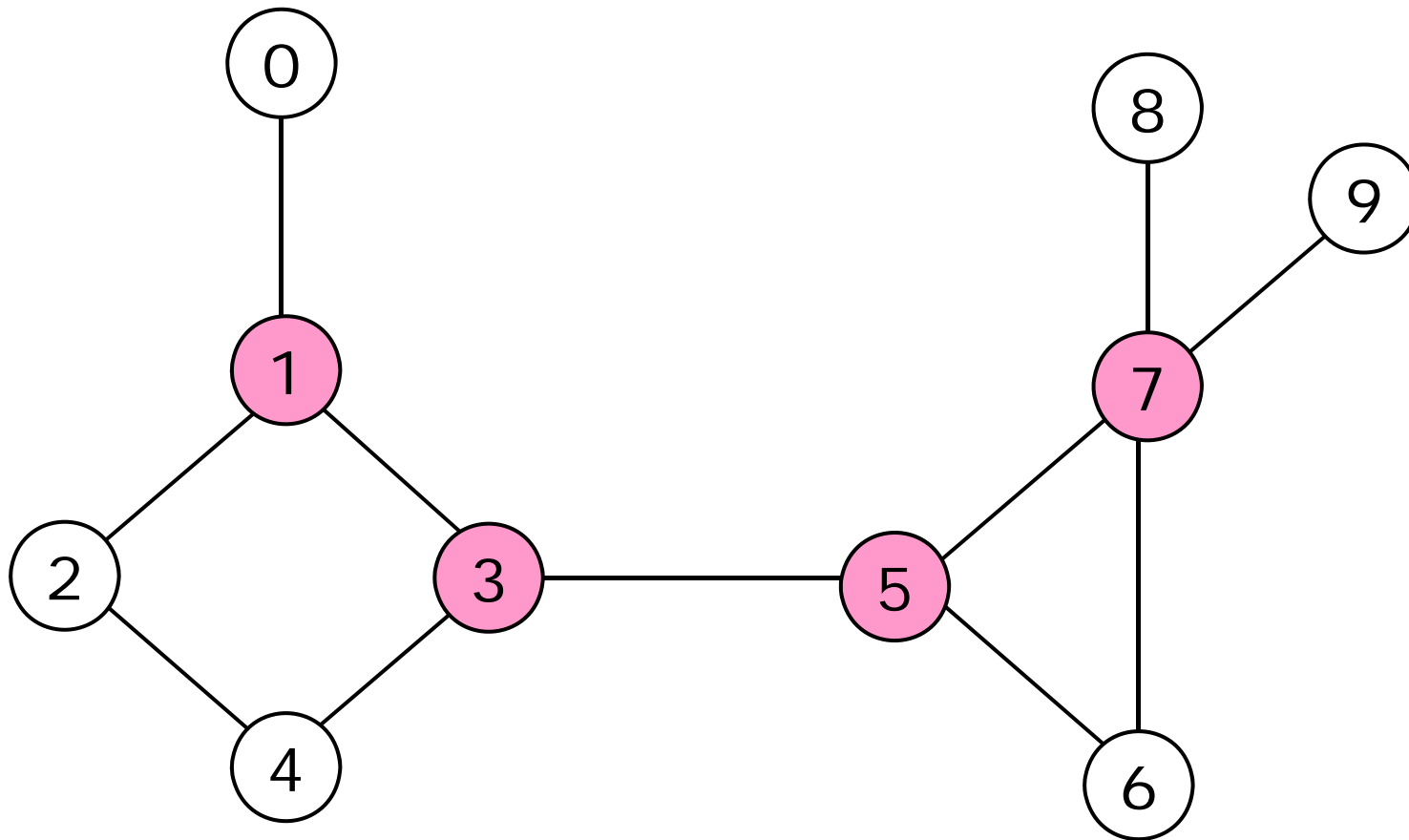
4.5 Biconnected Components

- Ex) What is the articulation point on this graph?



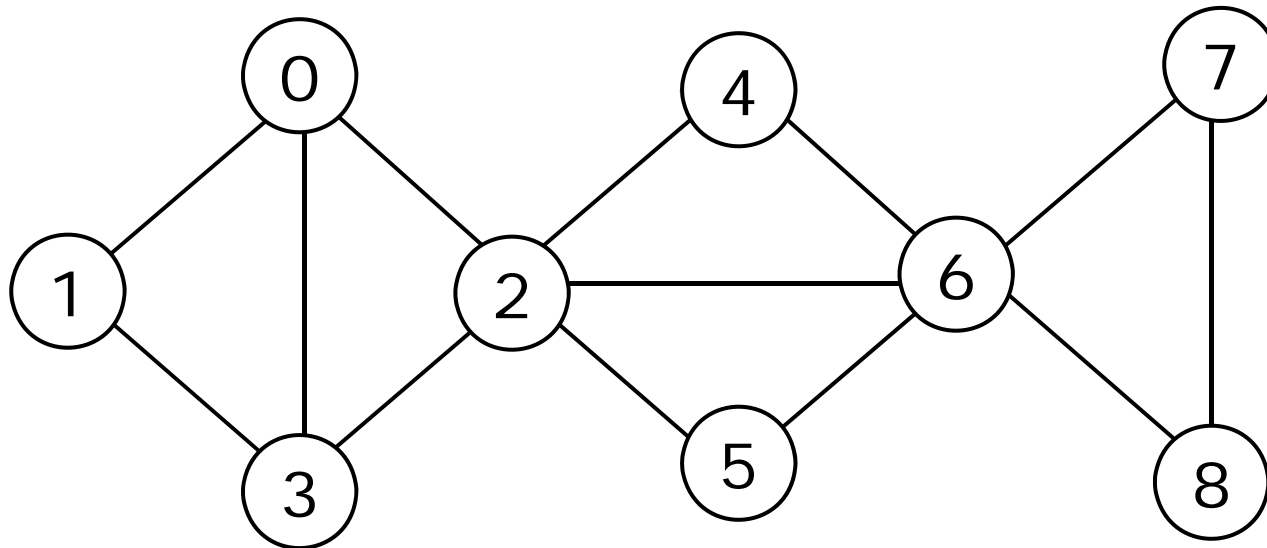
4.5 Biconnected Components

- Ex) What is the articulation point on this graph?



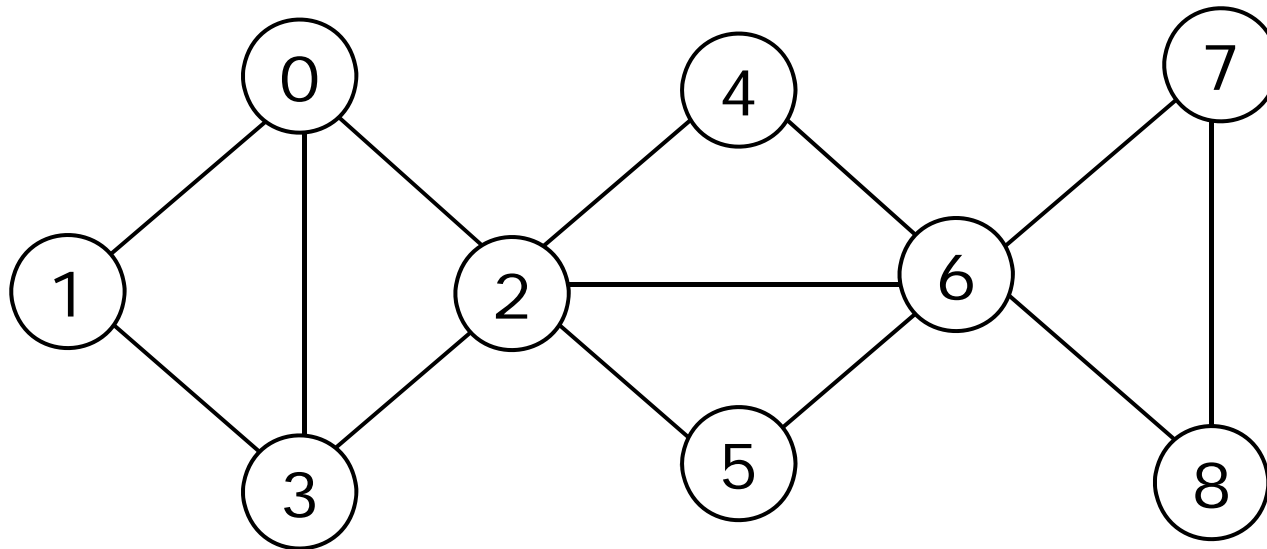
4.5 Biconnected Components

- Biconnected graph
 - A connected graph that has no articulation points



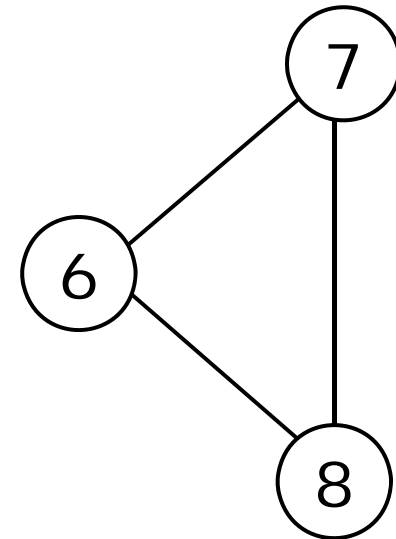
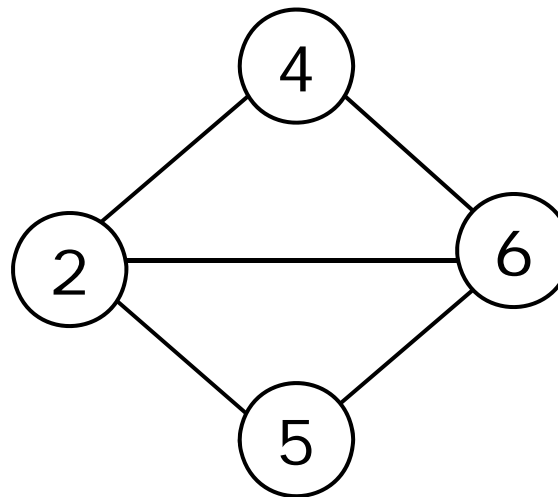
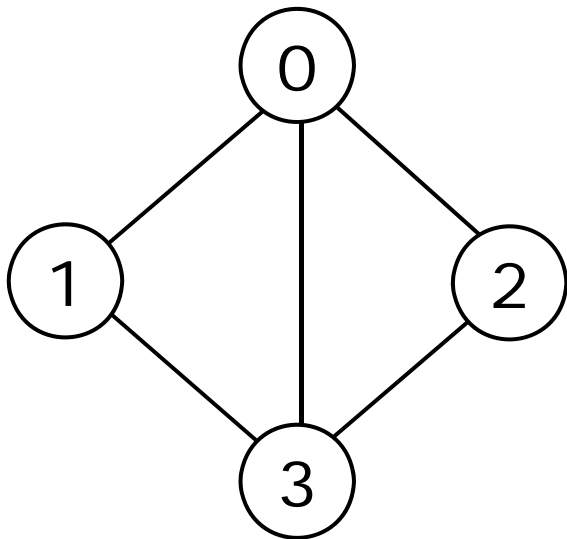
4.5 Biconnected Components

- Biconnected graph
 - A connected graph that has no articulation points
- Biconnected component
 - A maximal biconnected subgraph



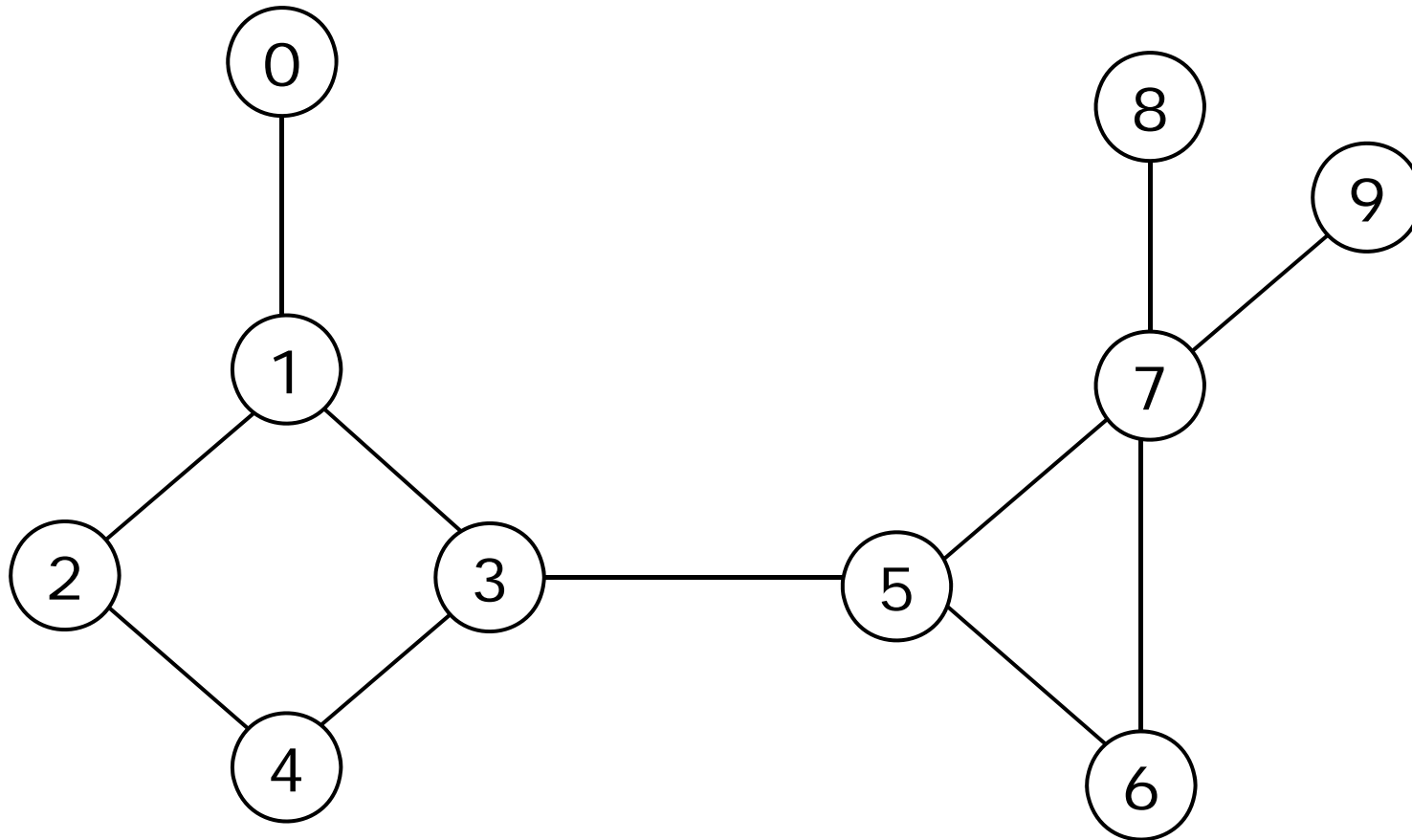
4.5 Biconnected Components

- Biconnected graph
 - A connected graph that has no articulation points
- Biconnected component
 - A maximal biconnected subgraph



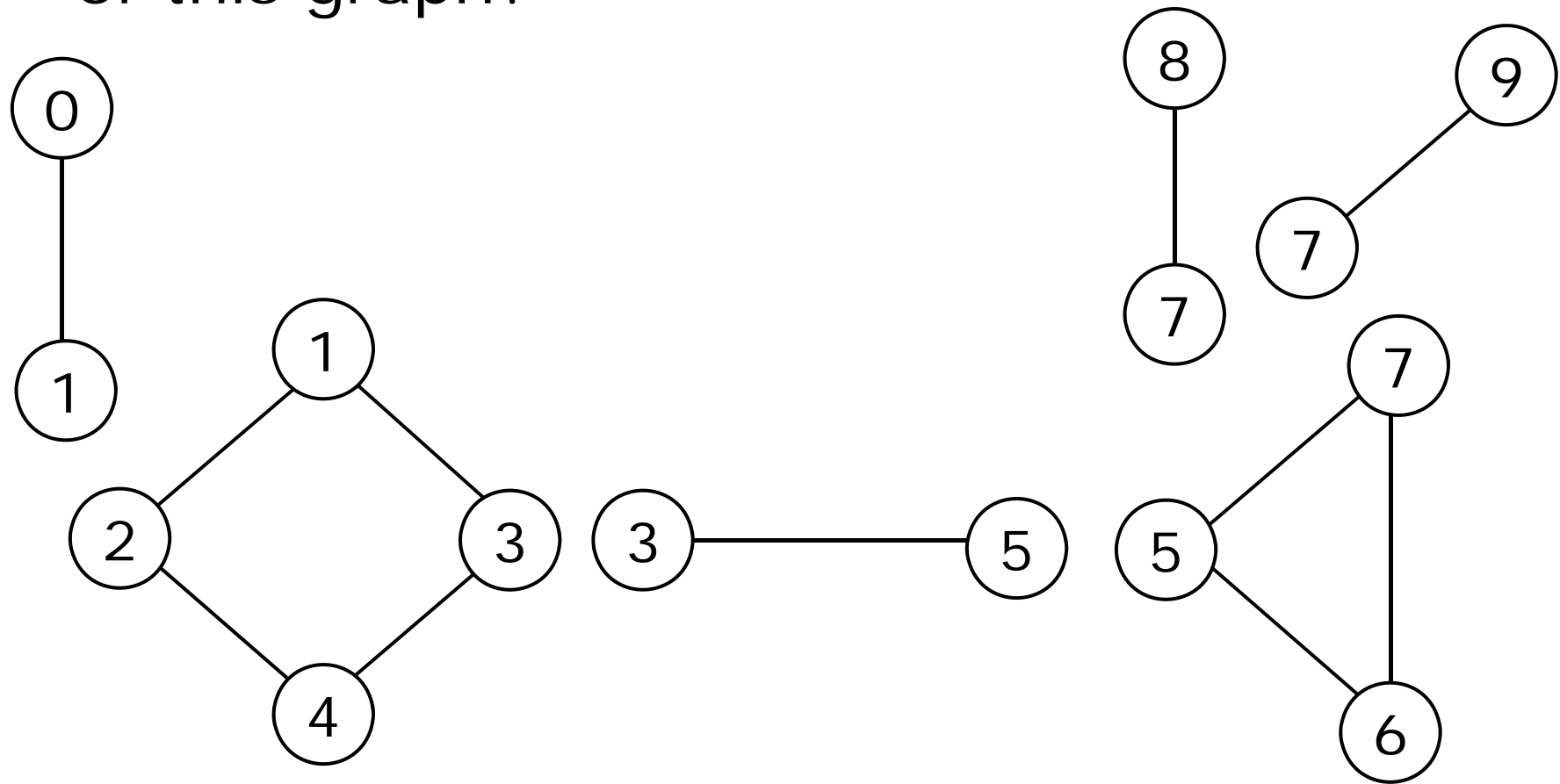
4.5 Biconnected Components

- Ex) What are the biconnected components of this graph?



4.5 Biconnected Components

- Ex) What are the biconnected components of this graph?

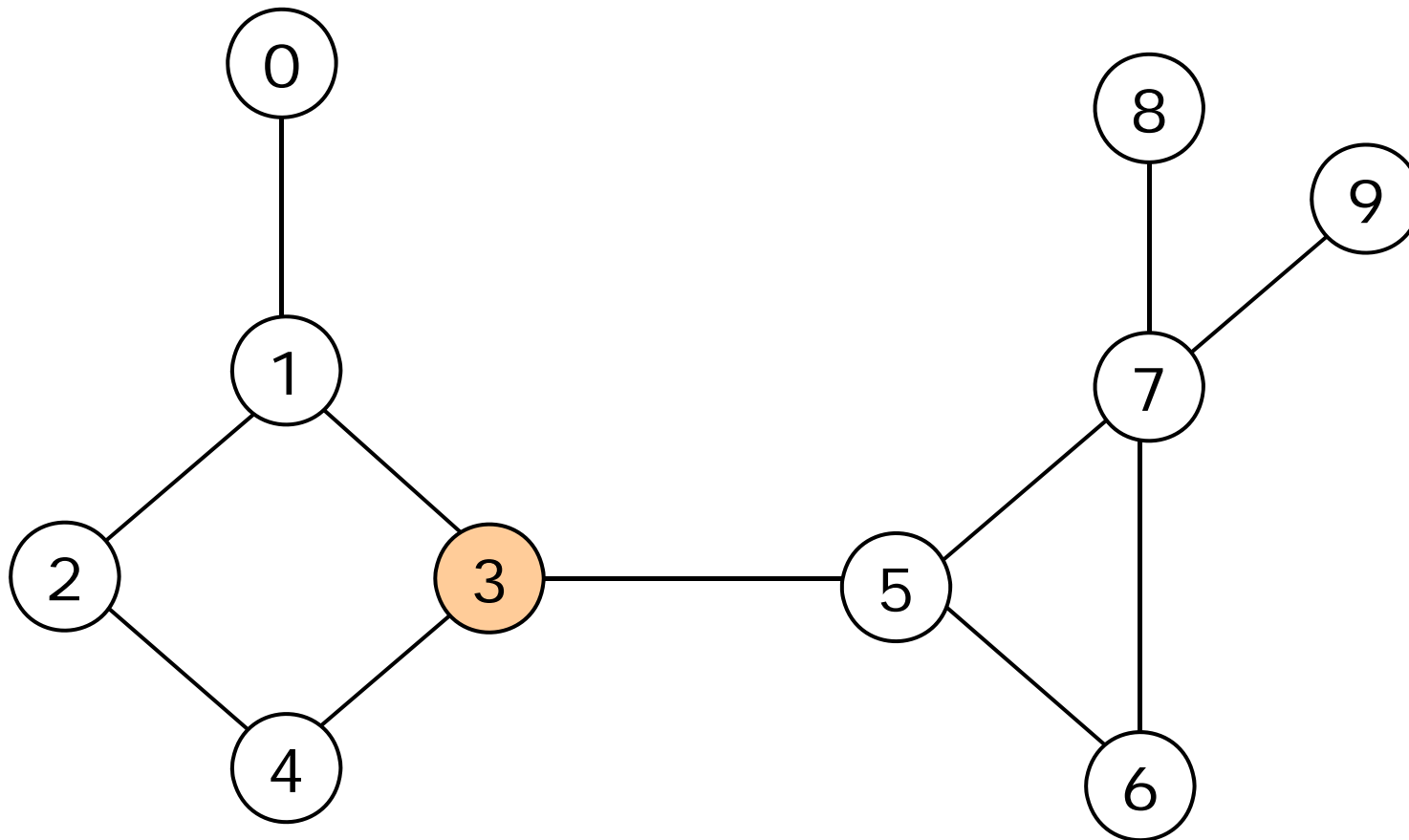


4.5 Biconnected Components

- How to find biconnected components of a graph
 - Use dfs ()
 - dfn (depth-first number) of a vertex
 - The sequence in which the vertices are visited during depth-first search
-

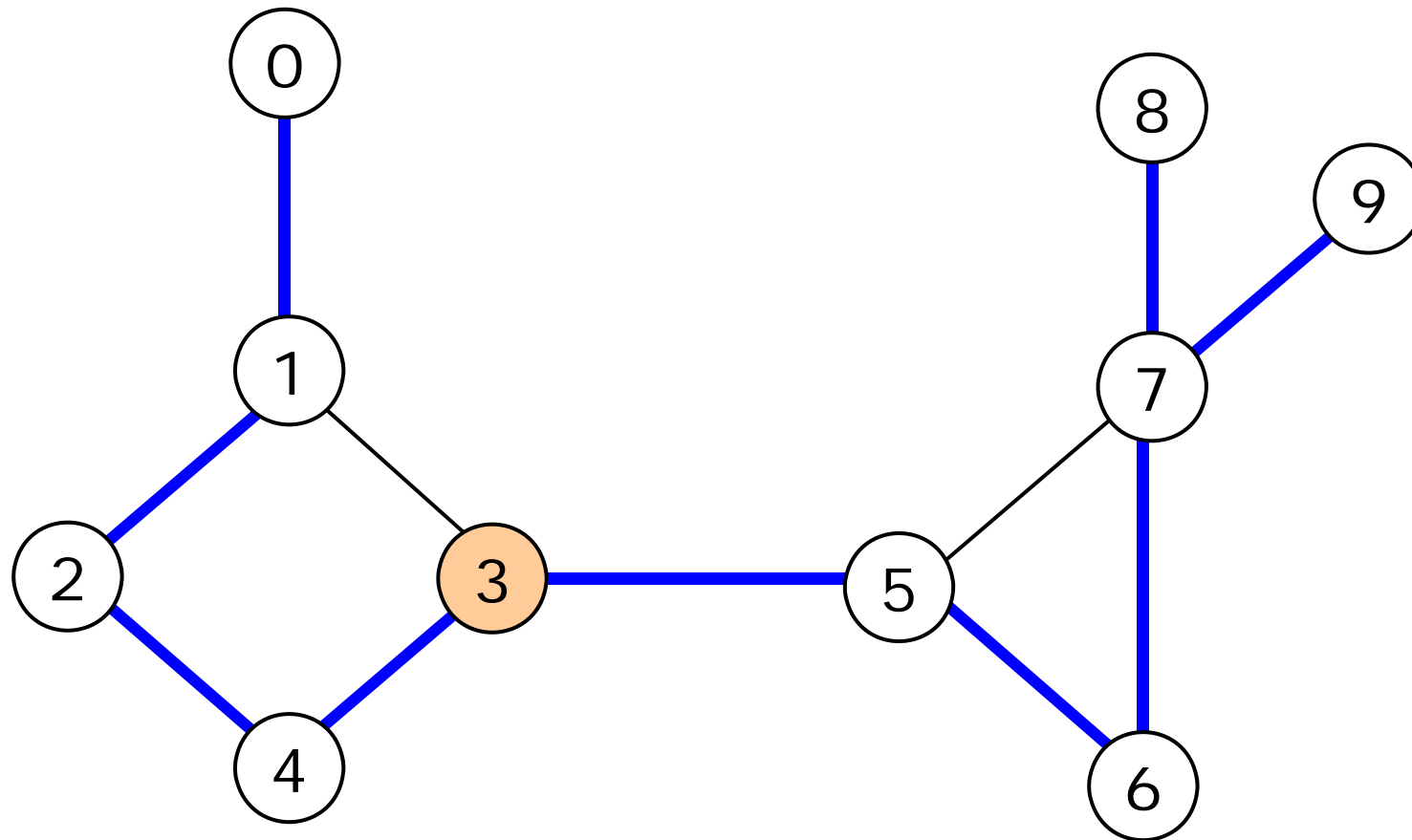
4.5 Biconnected Components

- Ex) What is dfn, if 3 is a root of depth-first spanning tree?



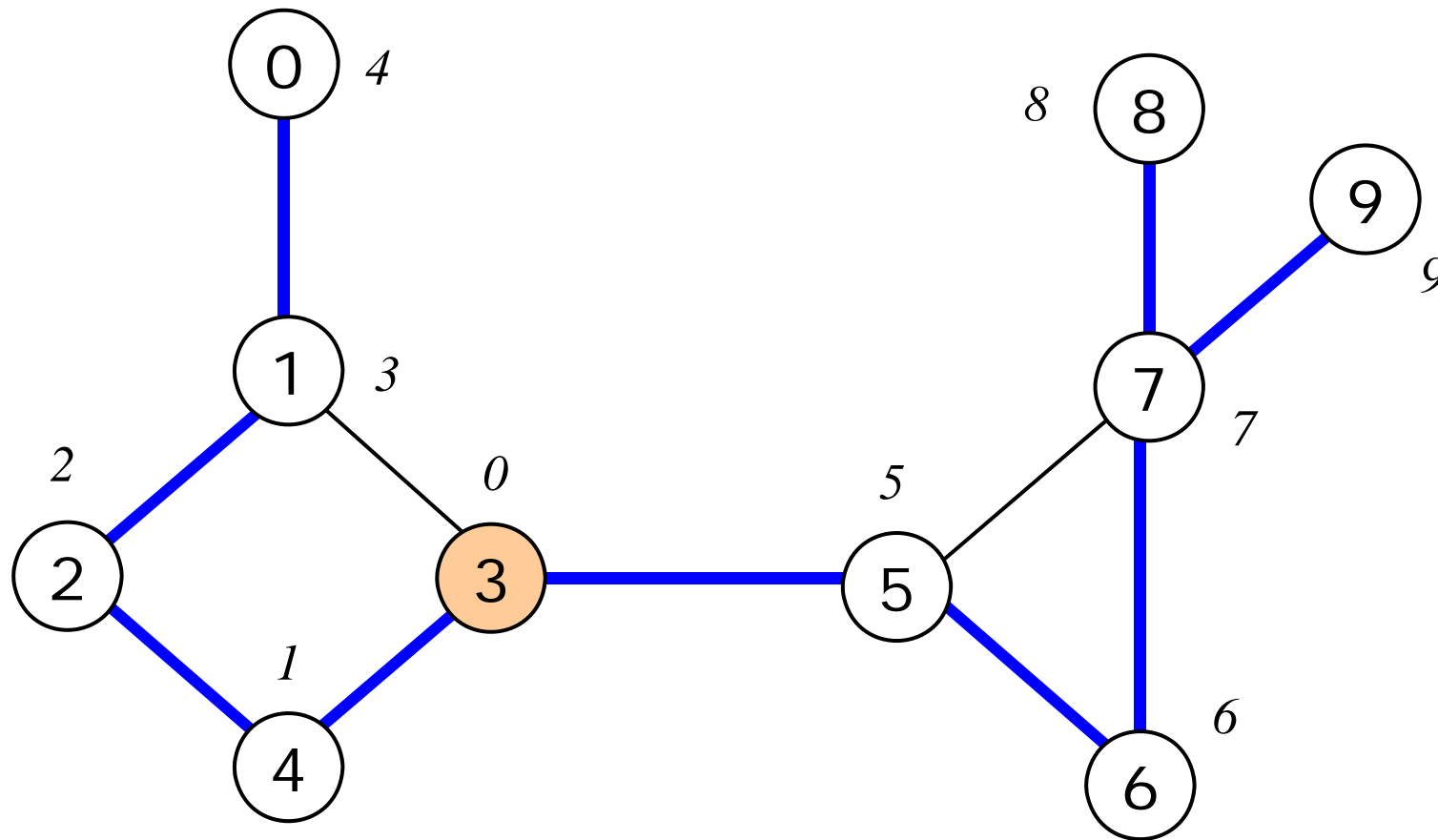
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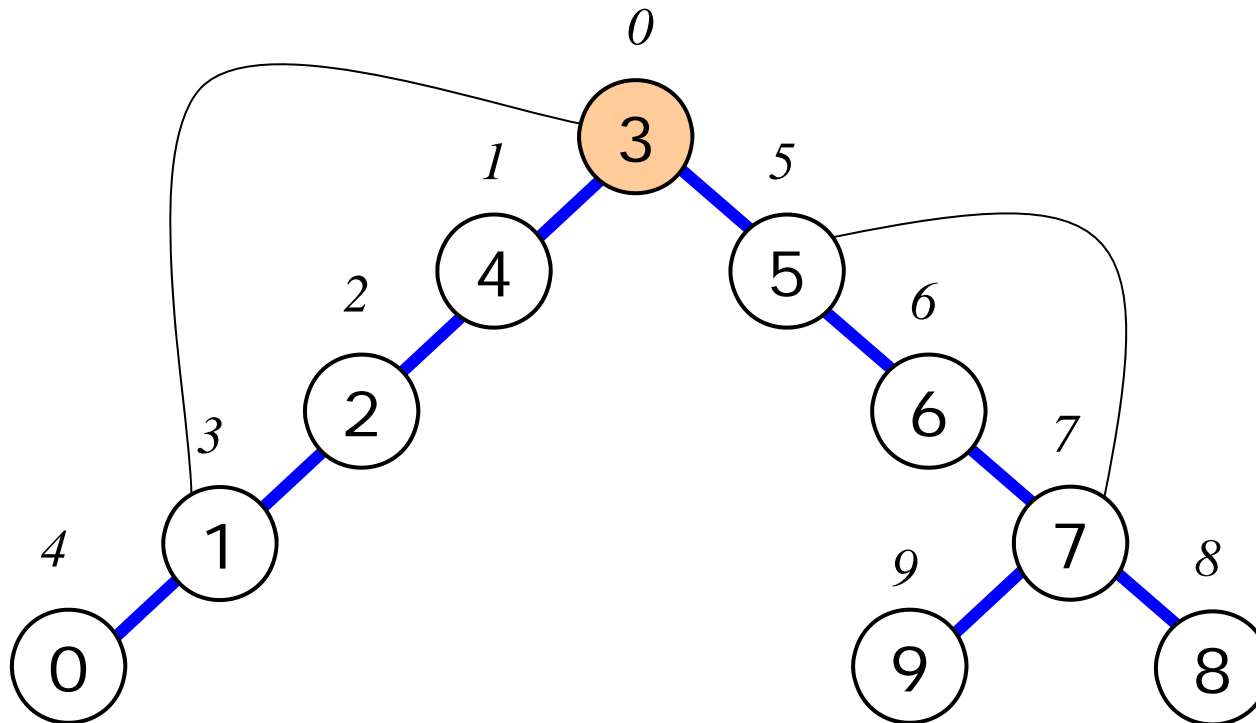
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4.5 Biconnected Components

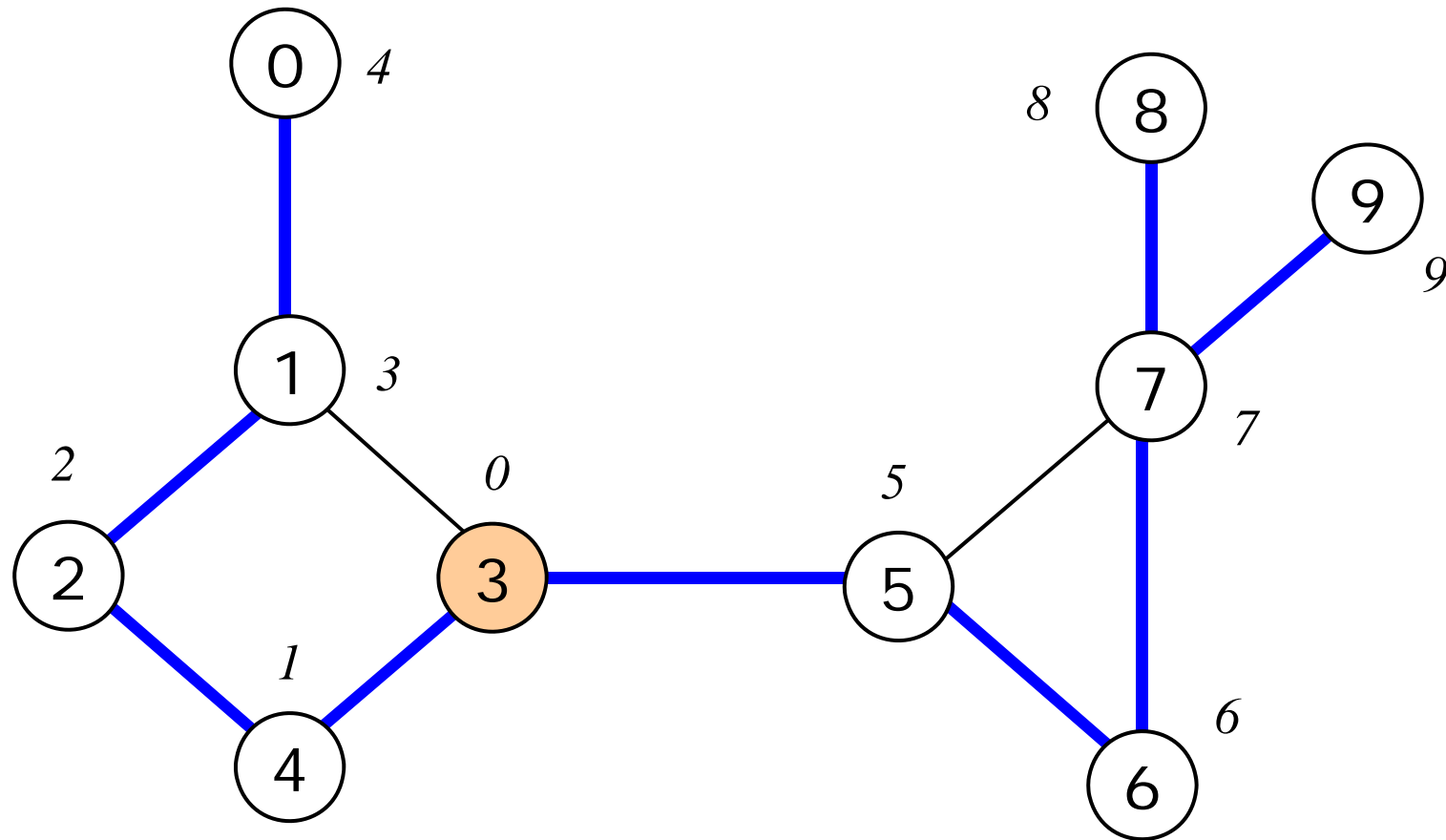
- Property
 - u is an ancestor of v in the depth-first spanning tree $\rightarrow \text{dfn}(u) < \text{dfn}(v)$

4.5 Biconnected Components

- Back edge
 - Edges in G = edges in the spanning tree + nontree edges
 - Back edge:
 - A nontree edge (u, v) , if u is an ancestor of v or vice versa
 - In depth-first spanning tree, all the nontree edges are back edges
-

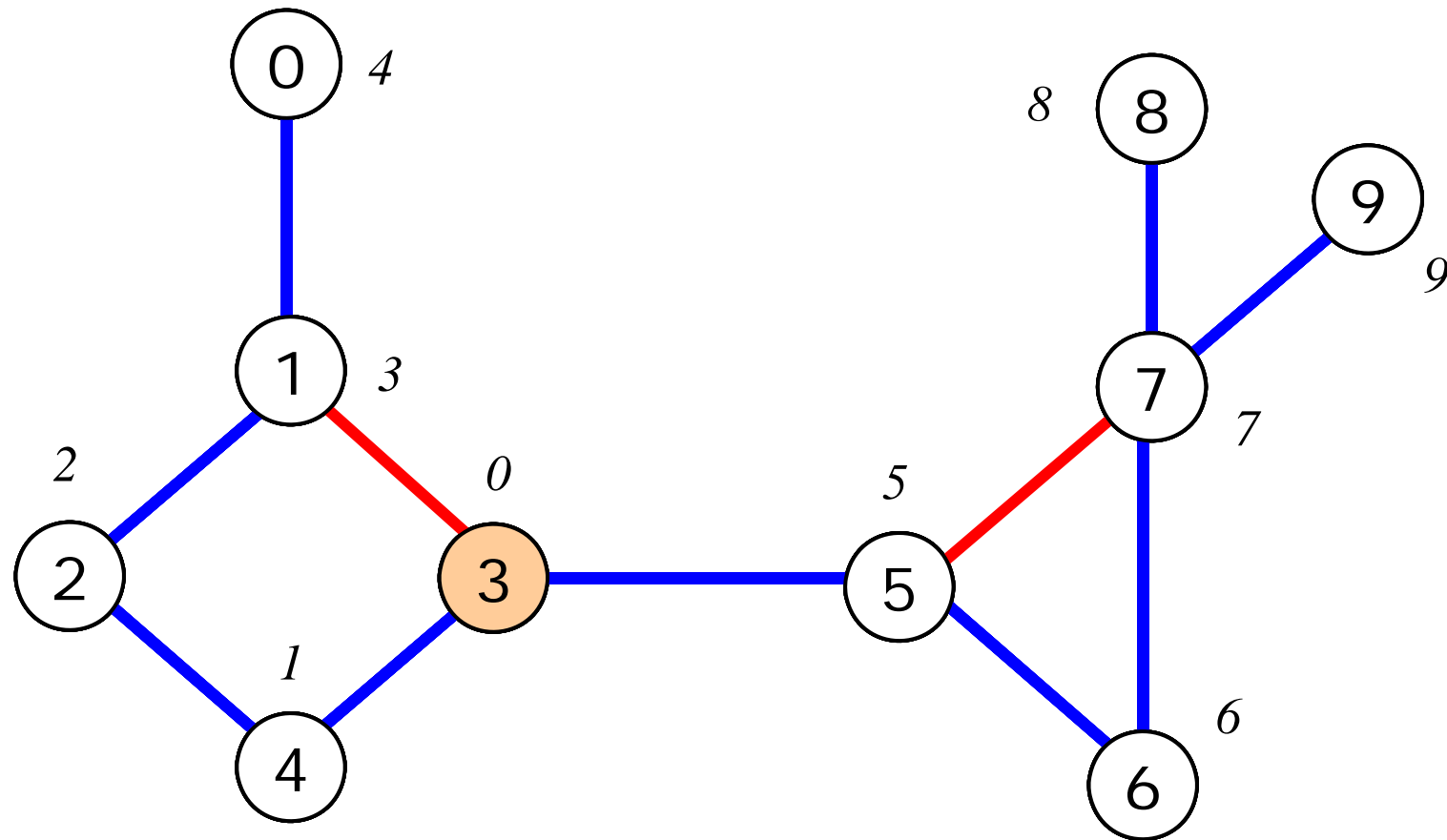
4.5 Biconnected Components

- Ex) What are the back edges of this graph?



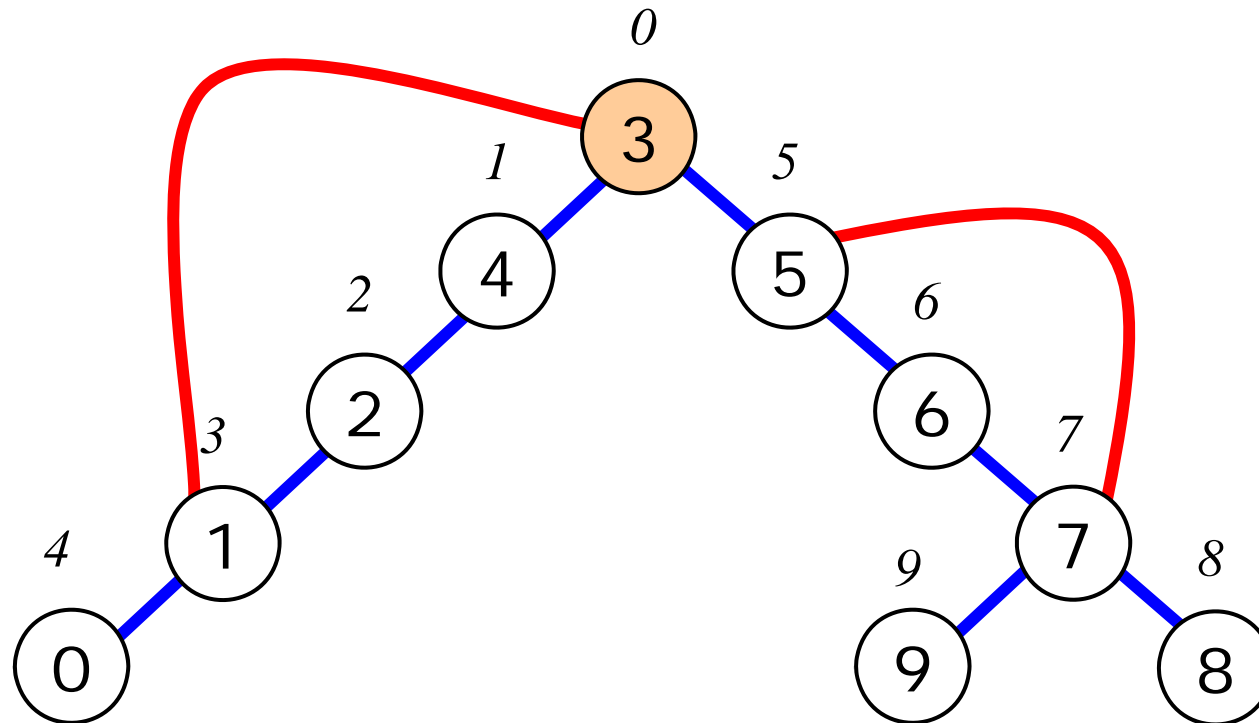
4.5 Biconnected Components

- Ex) What are the back edges of this graph?



4.5 Biconnected Components

- Ex) What are the back edges of this graph?

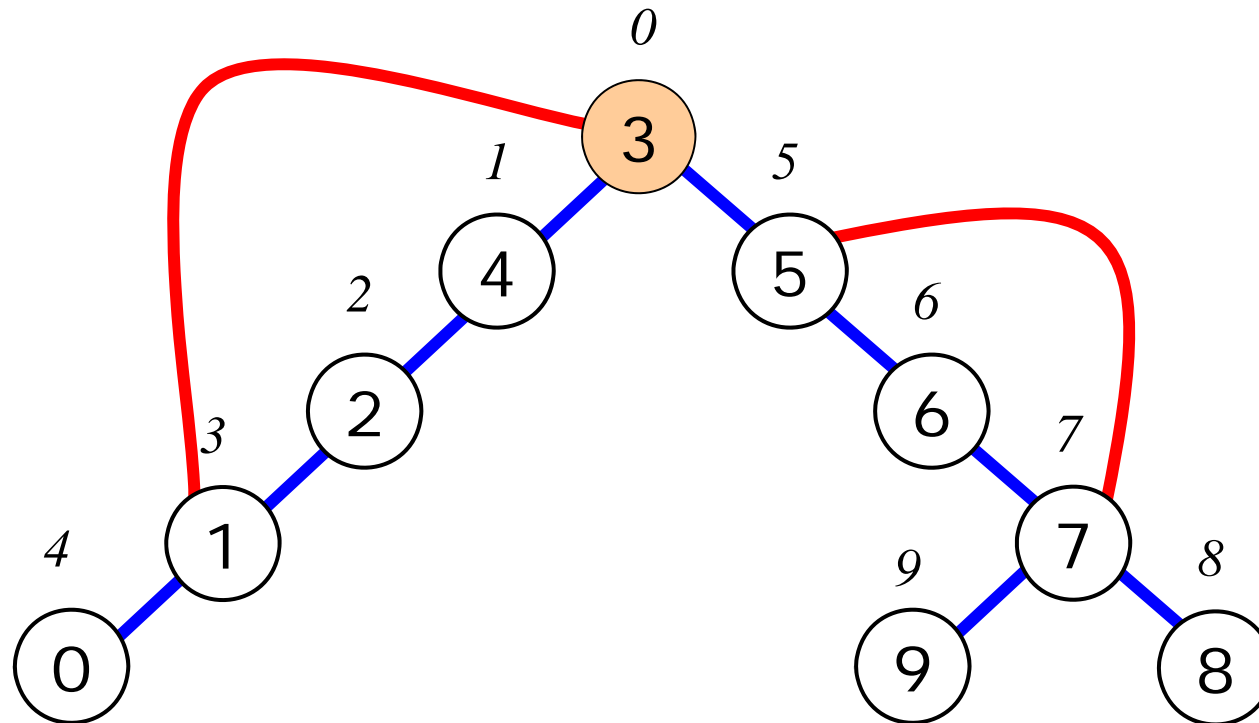


4.5 Biconnected Components

- Articulation points (1)
 - A root of a depth-first spanning tree is an articulation point, if it has at least two childs

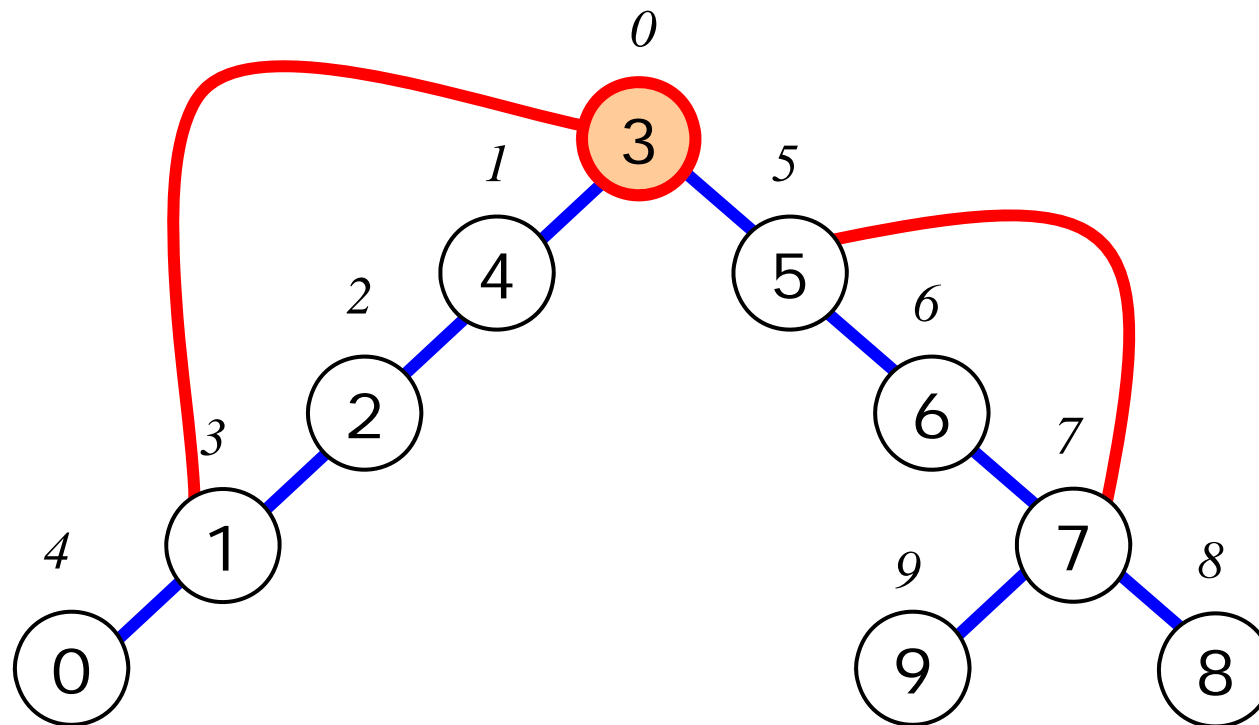
4.5 Biconnected Components

- Articulation points
 - Ex) What are the articulation points?



4.5 Biconnected Components

- Articulation points
 - Ex)

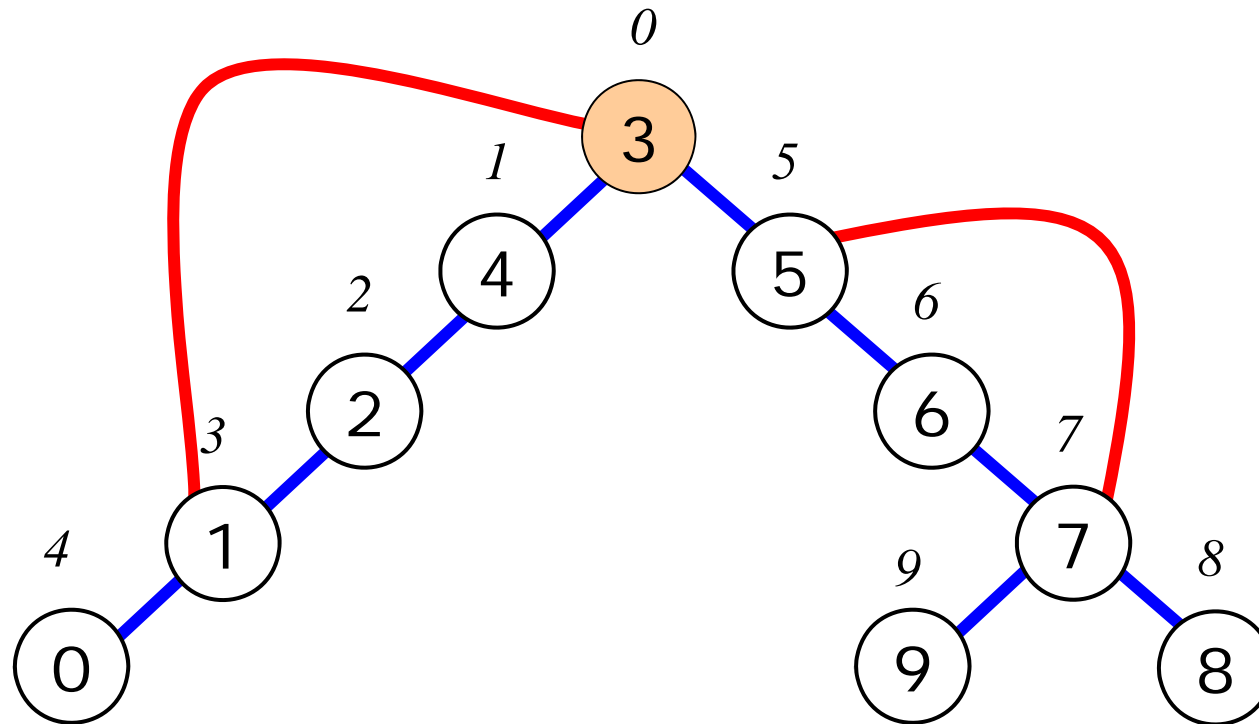


4.5 Biconnected Components

- Articulation points (2)
 - A vertex u , if it has at least one child w such that a path (w , descendants of w , and a single back edge, ancestor of u) does not exist

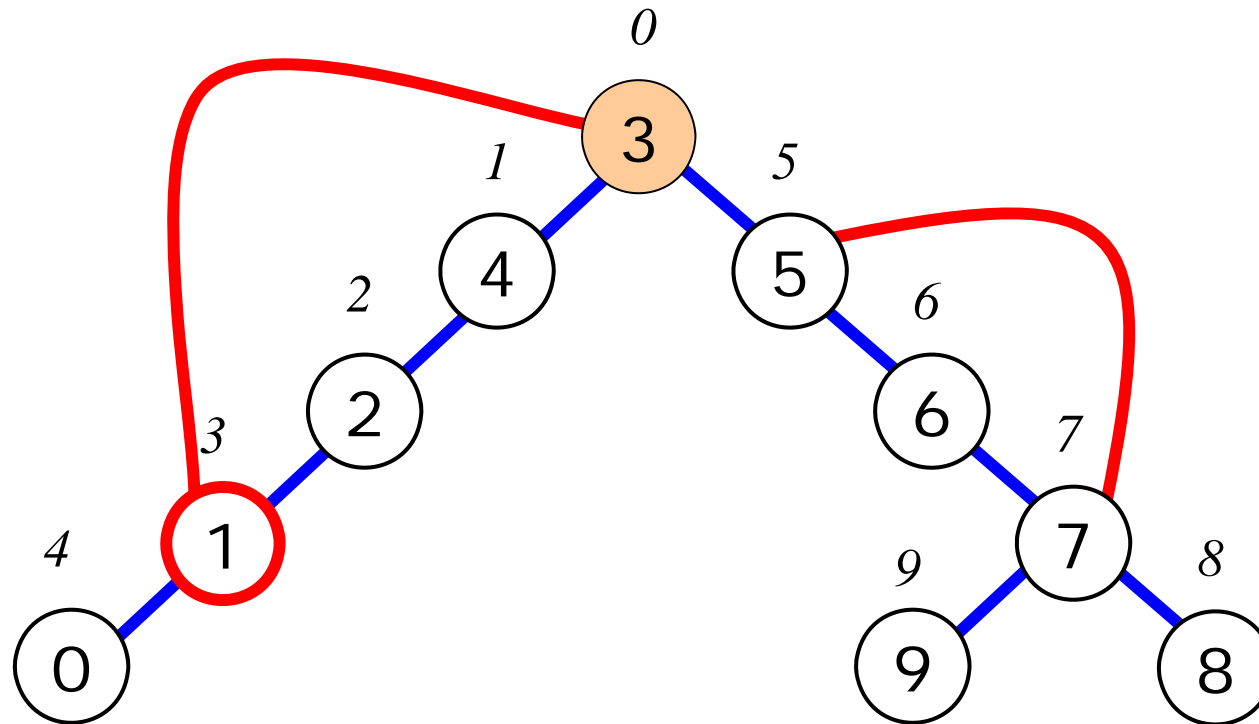
4.5 Biconnected Components

- Articulation points
 - Ex) What are the articulation points?



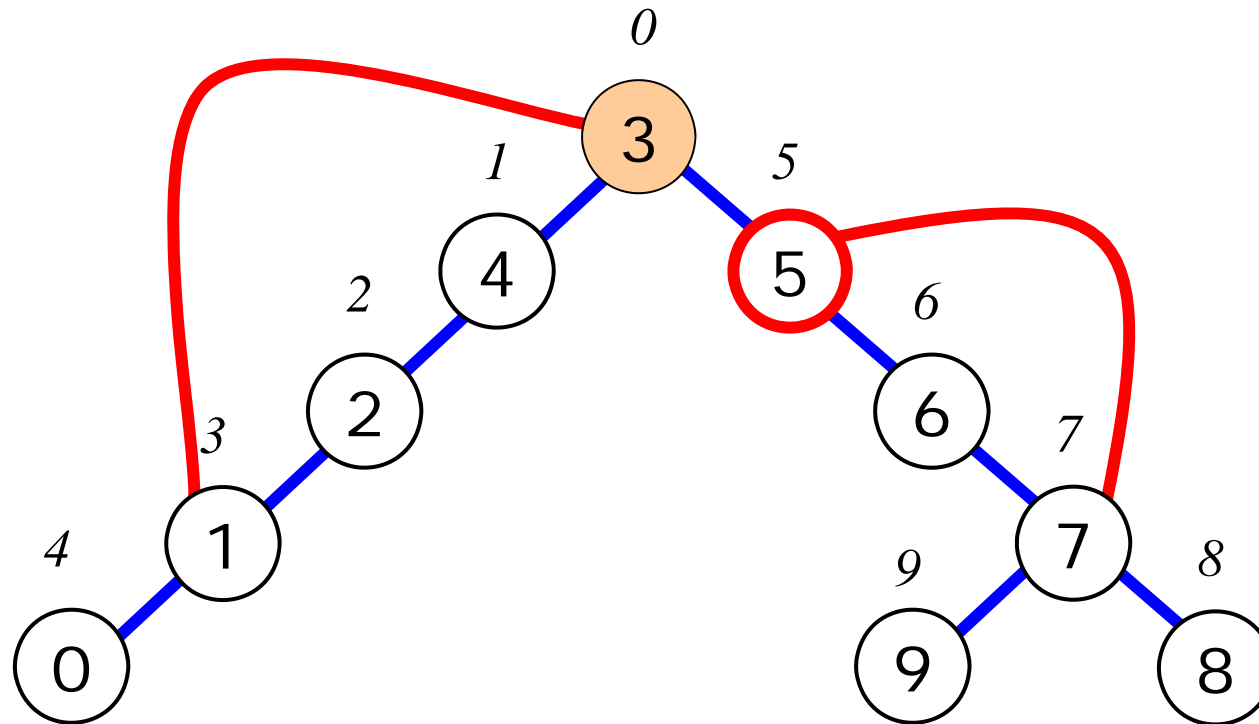
4.5 Biconnected Components

- Articulation points
 - Ex) What are the articulation points?



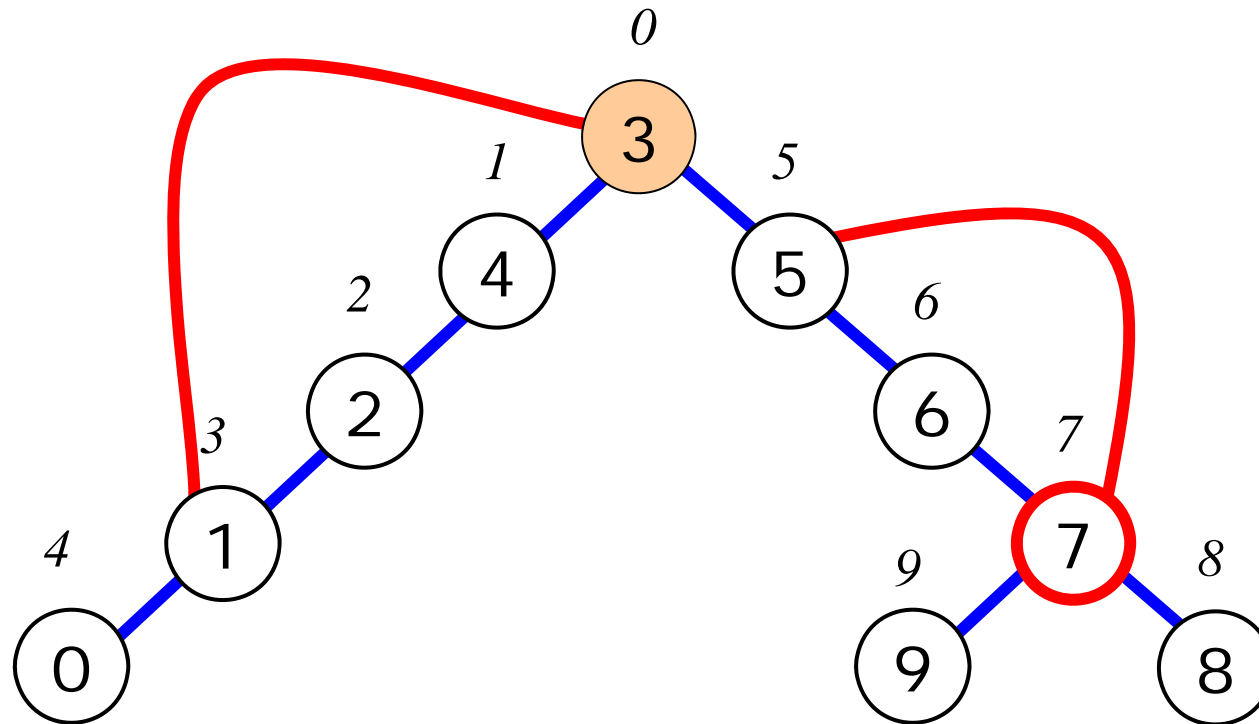
4.5 Biconnected Components

- Articulation points
 - Ex) What are the articulation points?



4.5 Biconnected Components

- Articulation points
 - Ex) What are the articulation points?



4.5 Biconnected Components

- How to find articulation points?
 - Define a new value low for each vertex u , such as $low(u)$
 - $low(u)$
 - The lowest depth first number that we can reach from u using a path of descendants followed by at most one back edge

$$low(u) = \min \{ dfn(u), \\ \min \{ low(w) \mid w \text{ is a child of } u \}, \\ \min \{ dfn(v) \mid (u, v) \text{ is a back edge} \} \}$$

4.5 Biconnected Components

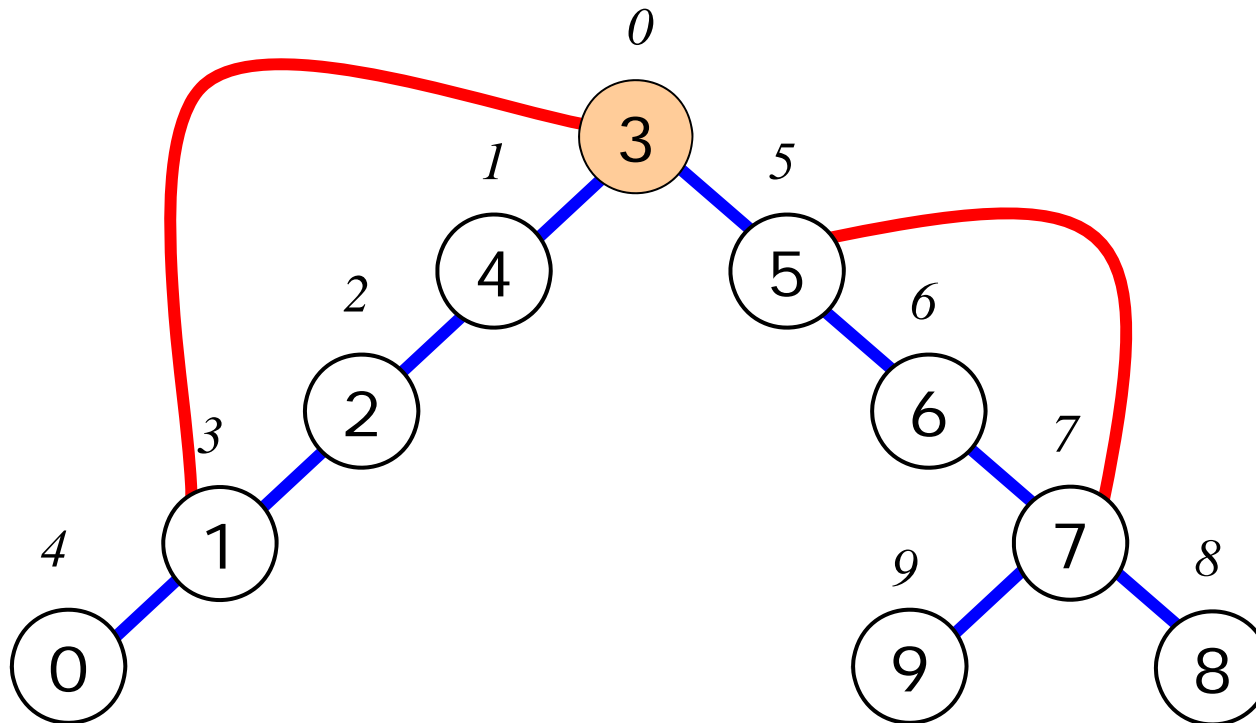
- $\text{low}(u)$

- Ex) What are $\text{low}(u)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u, v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

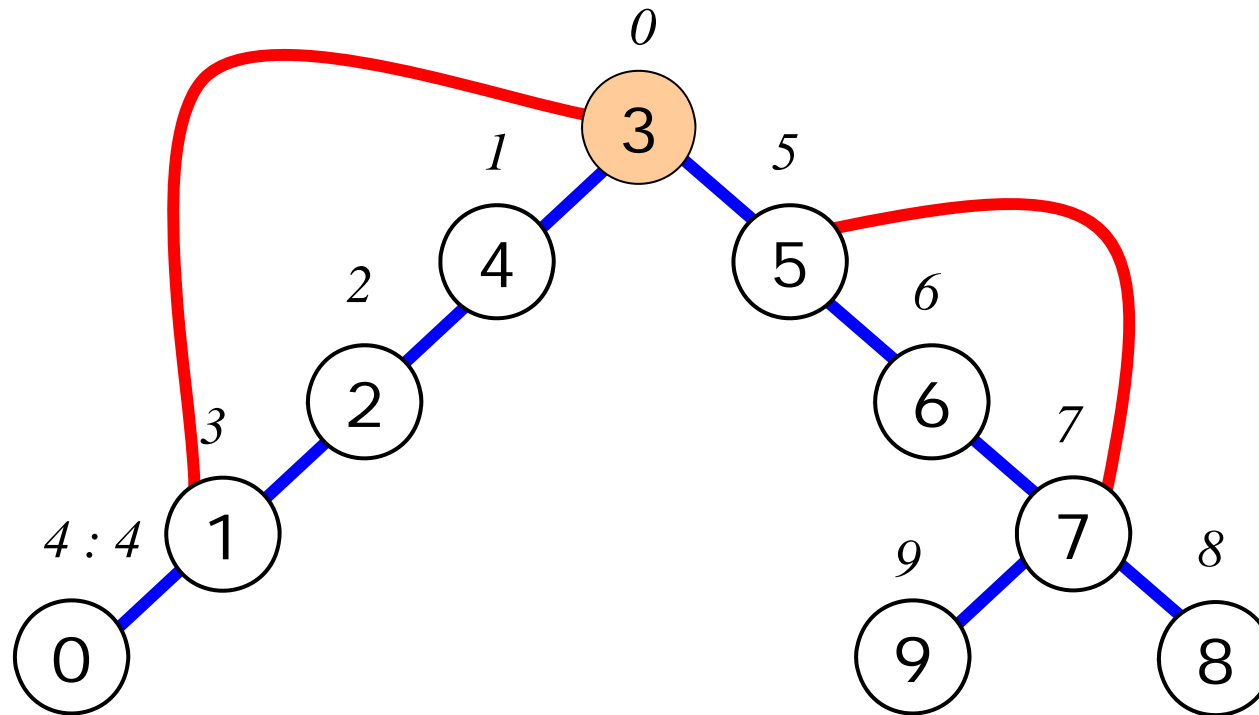
- $\text{low}(u)$

- Ex) What are $\text{low}(0)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u, v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

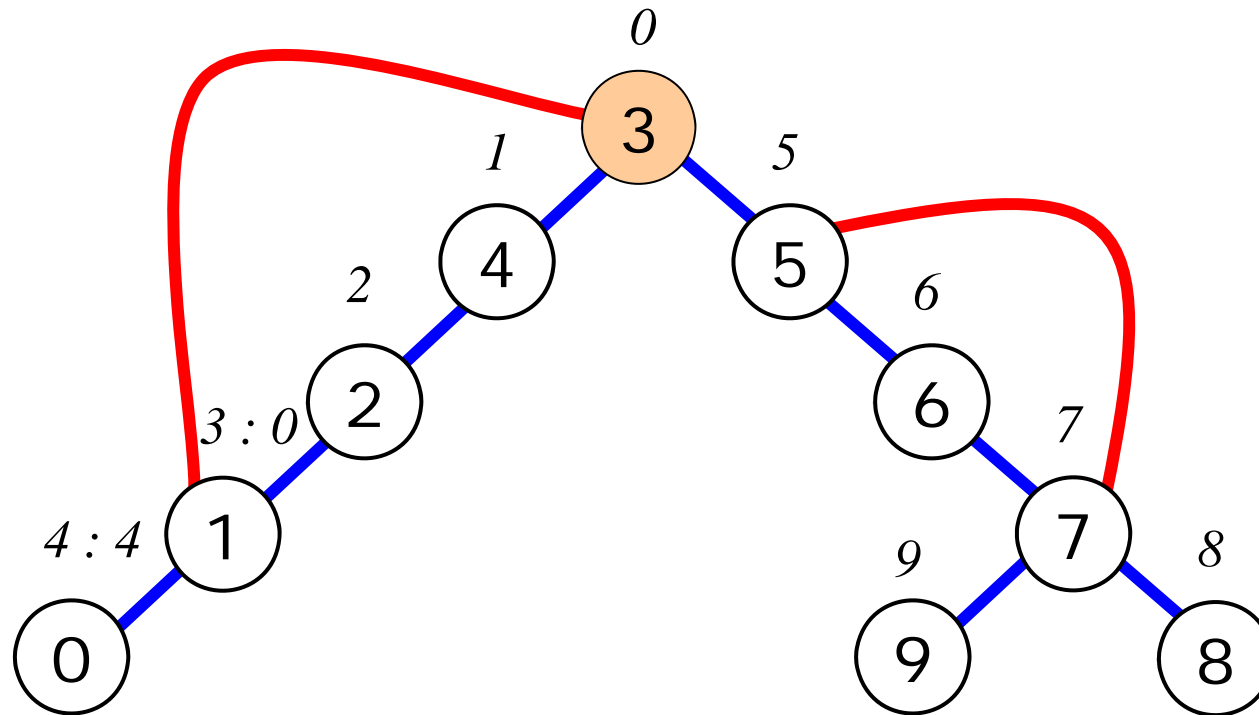
- $low(u)$

- Ex) What are $low(1)$?

$$low(u) = \min\{dfn(u),$$

$$\min\{low(w) \mid w \text{ is a child of } u\},$$

$$\min\{dfn(v) \mid (u,v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

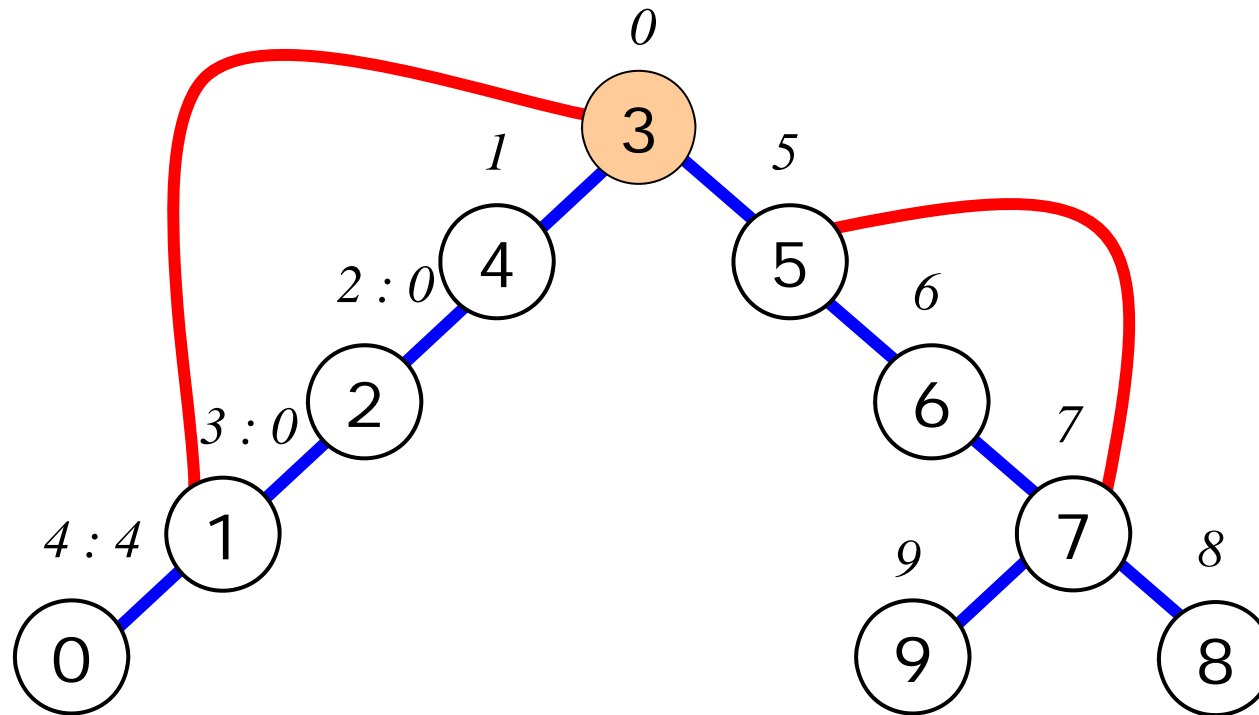
- $\text{low}(u)$

- Ex) What are $\text{low}(2)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u,v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

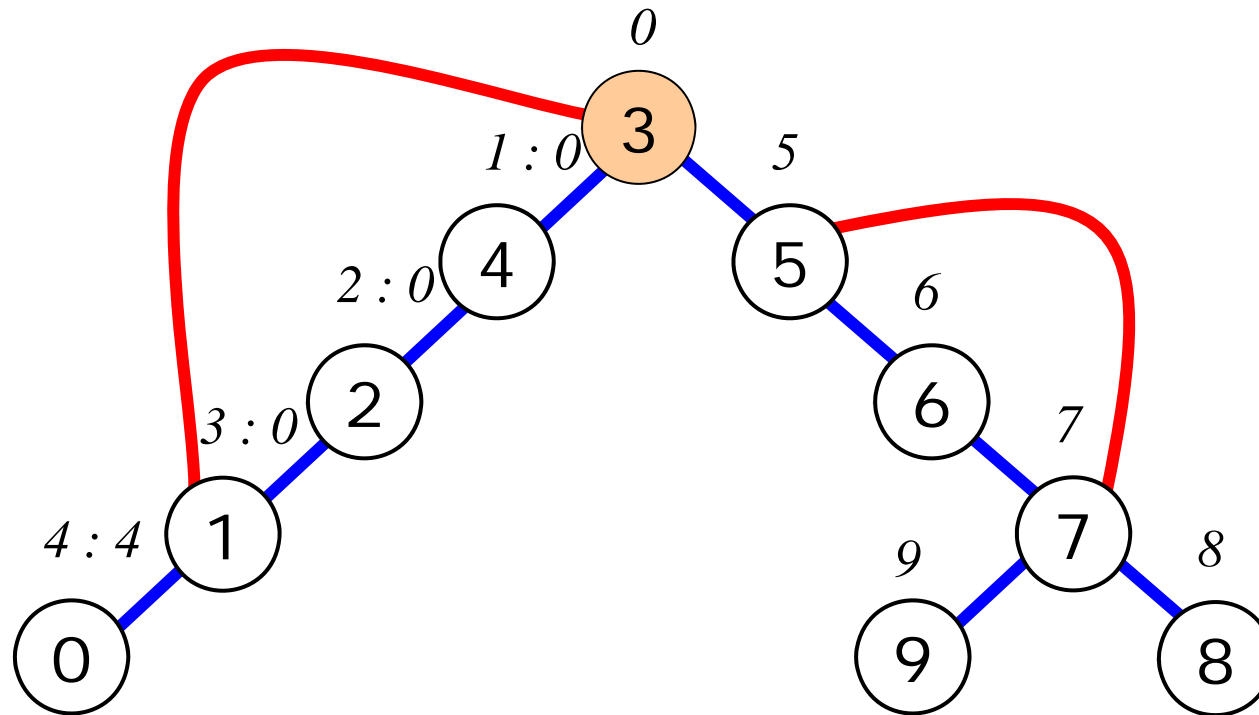
- $\text{low}(u)$

- Ex) What are $\text{low}(4)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u,v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

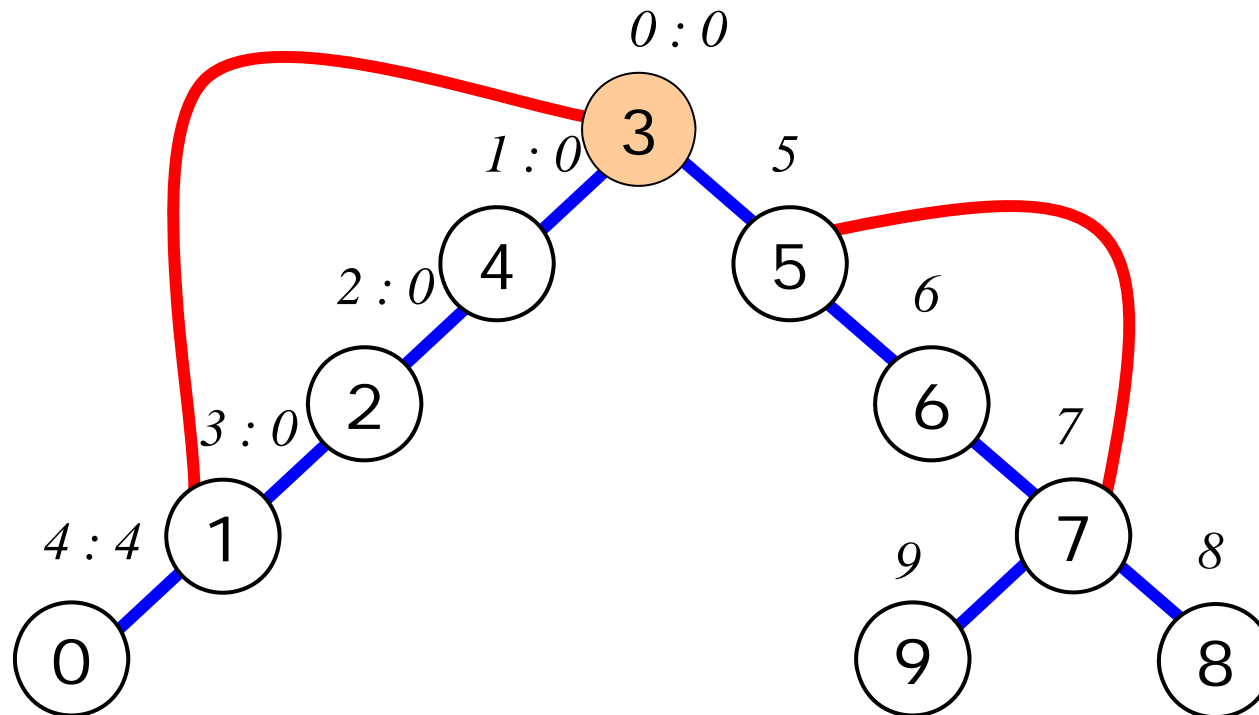
- $\text{low}(u)$

- Ex) What are $\text{low}(3)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u, v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

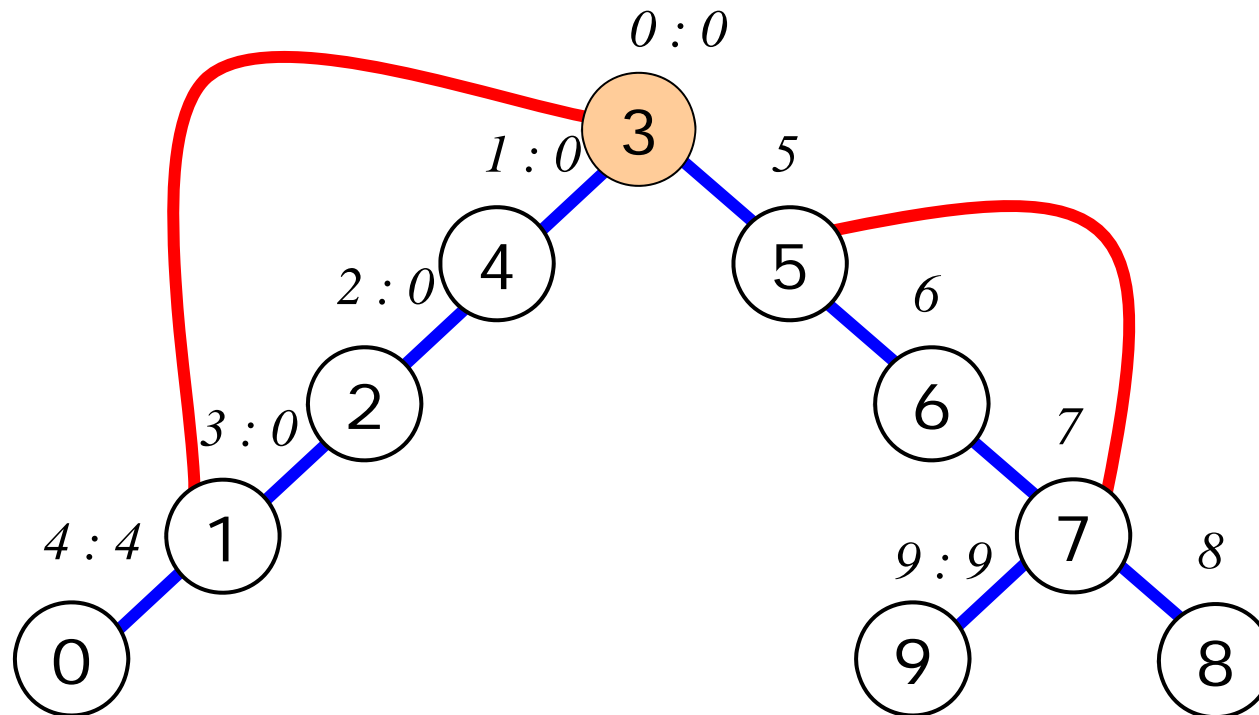
- $\text{low}(u)$

- Ex) What are $\text{low}(9)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u,v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

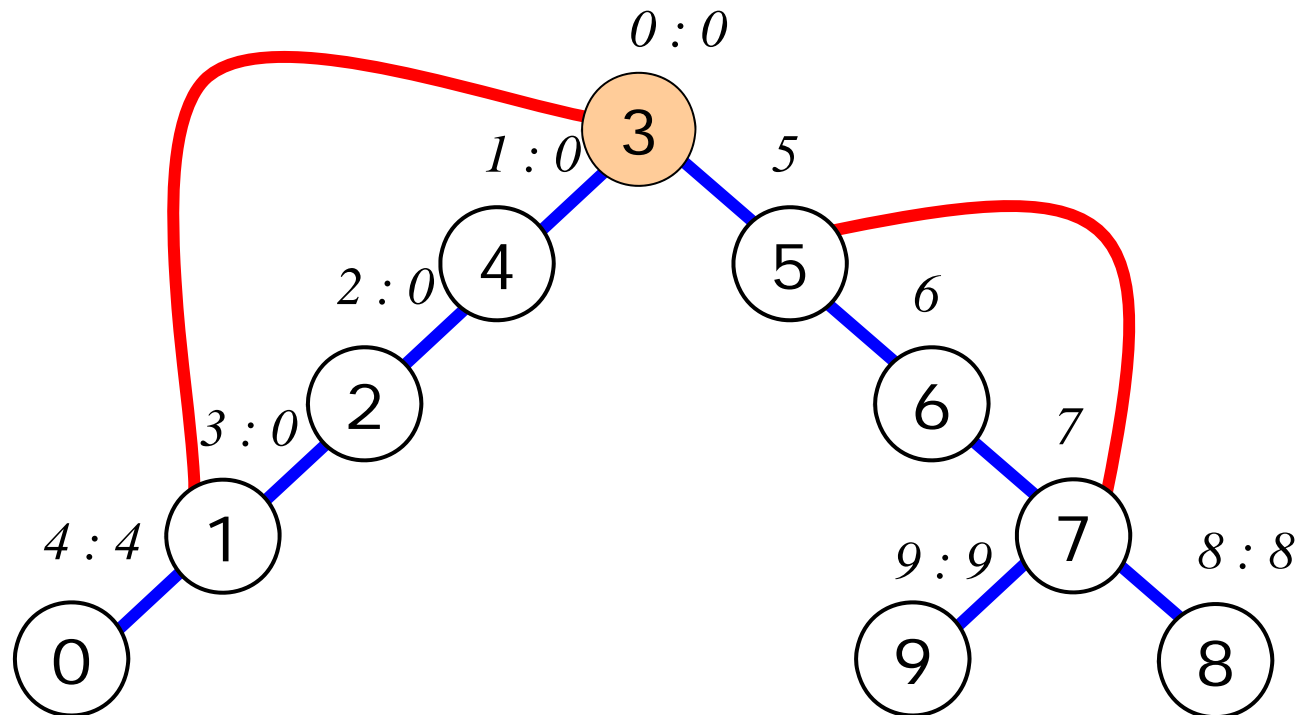
- $low(u)$

- Ex) What are $low(8)$?

$$low(u) = \min\{dfn(u),$$

$$\min\{low(w) \mid w \text{ is a child of } u\},$$

$$\min\{dfn(v) \mid (u,v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

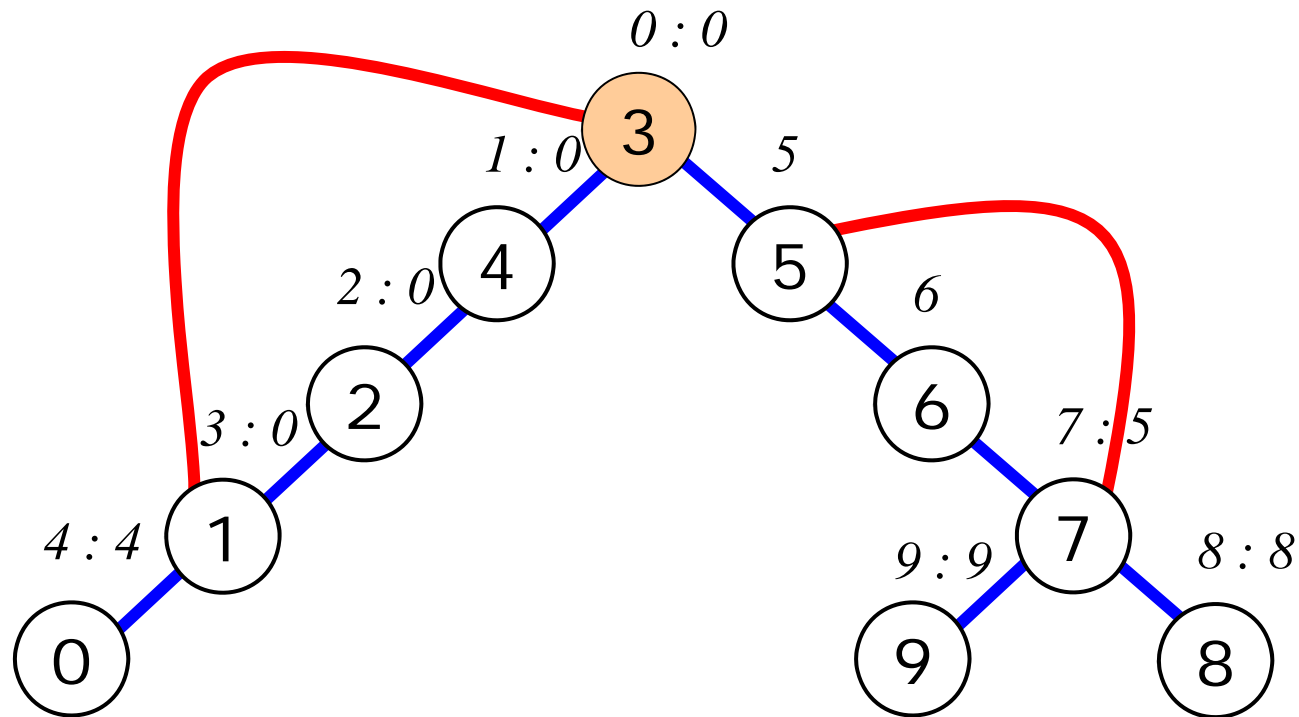
- $\text{low}(u)$

- Ex) What are $\text{low}(7)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u, v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

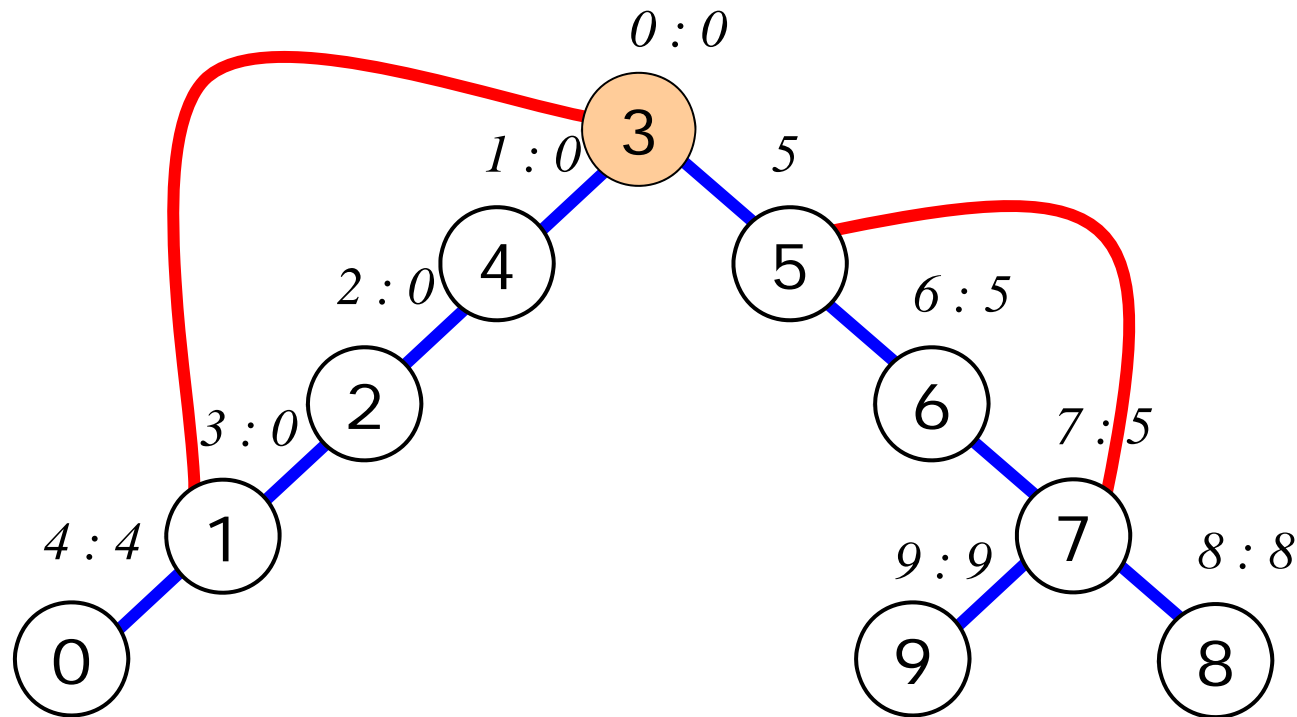
- $\text{low}(u)$

- Ex) What are $\text{low}(6)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u, v) \text{ is a back edge}\}\}$$



4.5 Biconnected Components

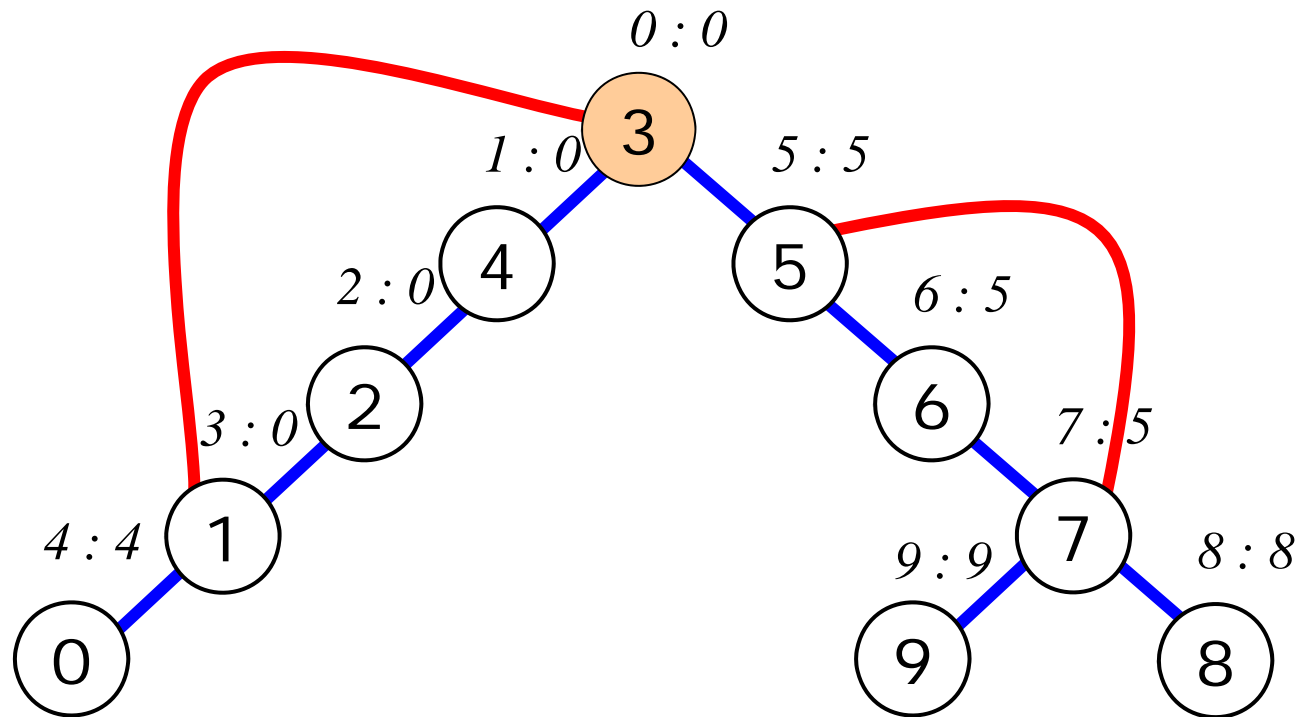
- $\text{low}(u)$

- Ex) What are $\text{low}(5)$?

$$\text{low}(u) = \min\{\text{dfn}(u),$$

$$\min\{\text{low}(w) \mid w \text{ is a child of } u\},$$

$$\min\{\text{dfn}(v) \mid (u,v) \text{ is a back edge}\}\}$$



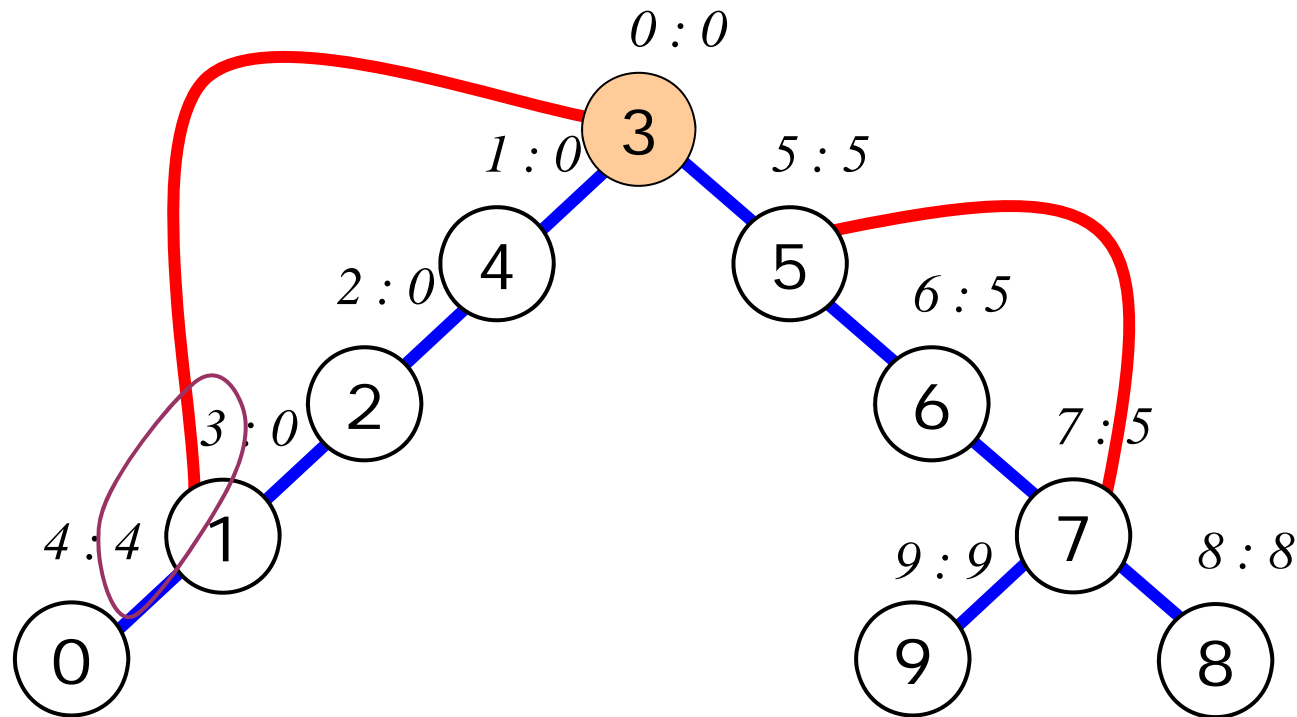
4.5 Biconnected Components

- Articulation points
 - u is an articulation point,
 - if u is either the root of the spanning tree with two or more childs,
 - or u is not a root and has a child w such that $\text{low}(w) \geq \text{dfn}(u)$

	0	1	2	3	4	5	6	7	8	9
<i>dfn</i>	4	3	2	0	1	5	6	7	8	9
<i>low</i>	4	0	0	0	0	5	5	5	8	9

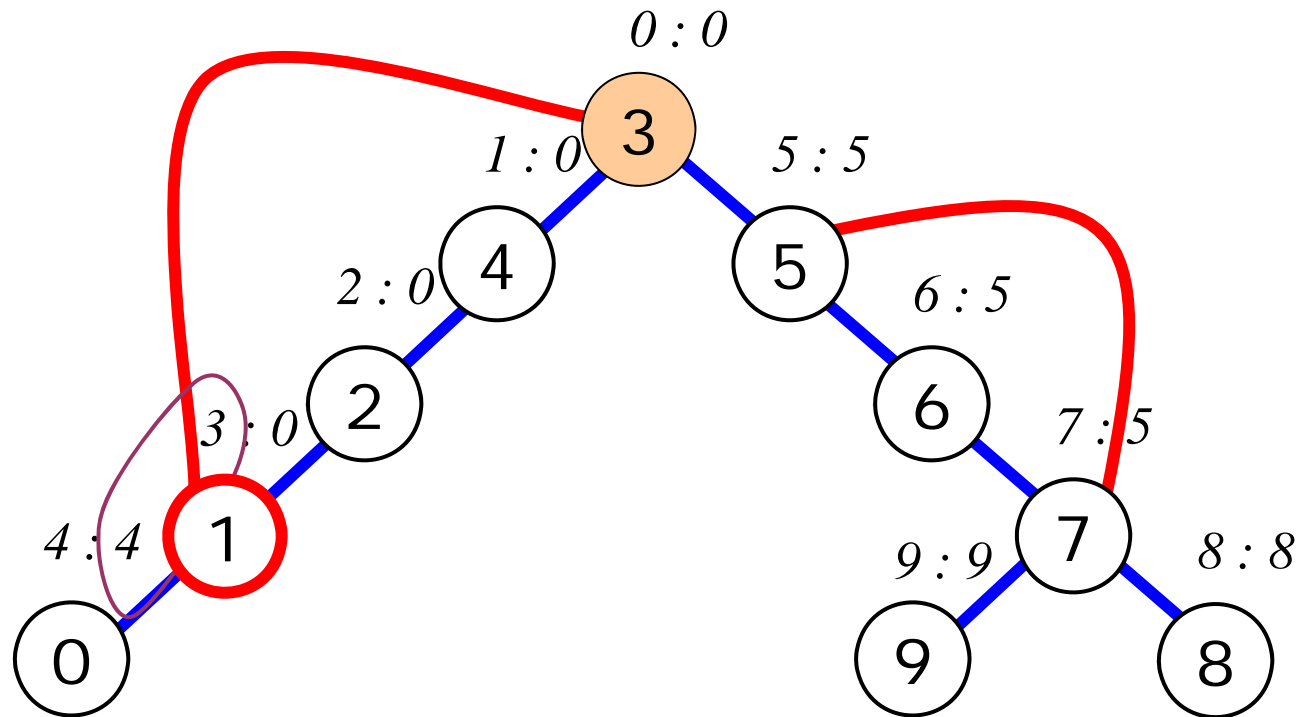
4.5 Biconnected Components

- Determine articulation points
 - At 1, $\text{low}(w) \geq \text{dfn}(u)$?



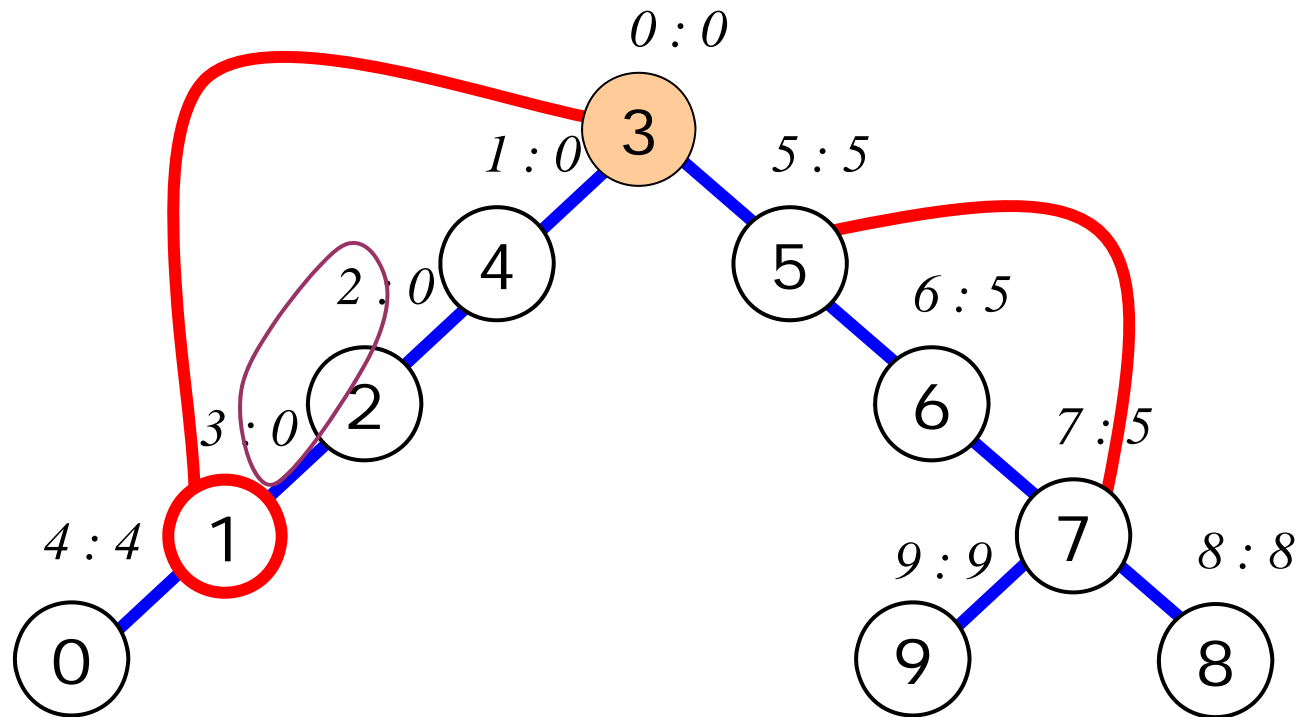
4.5 Biconnected Components

- Determine articulation points
 - At 1, $\text{low}(w) \geq \text{dfn}(u)$?



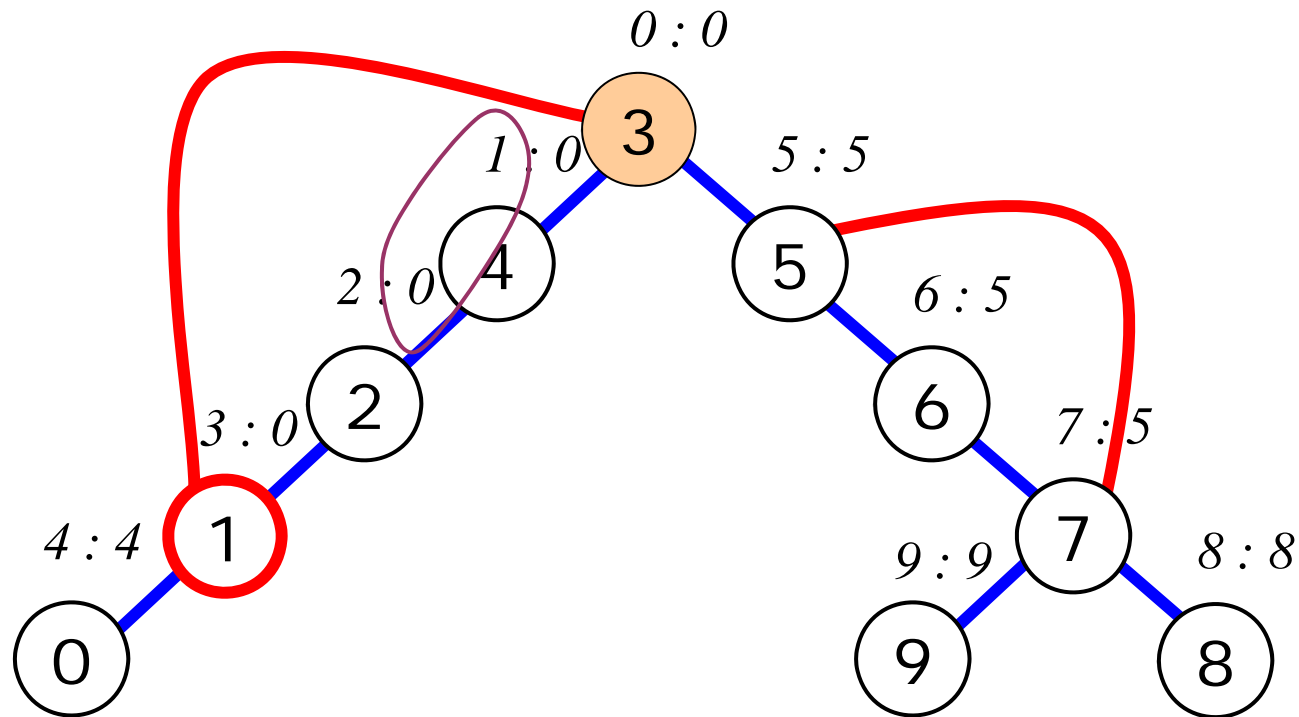
4.5 Biconnected Components

- Determine articulation points
 - At 2, $\text{low}(w) \geq \text{dfn}(u)$?



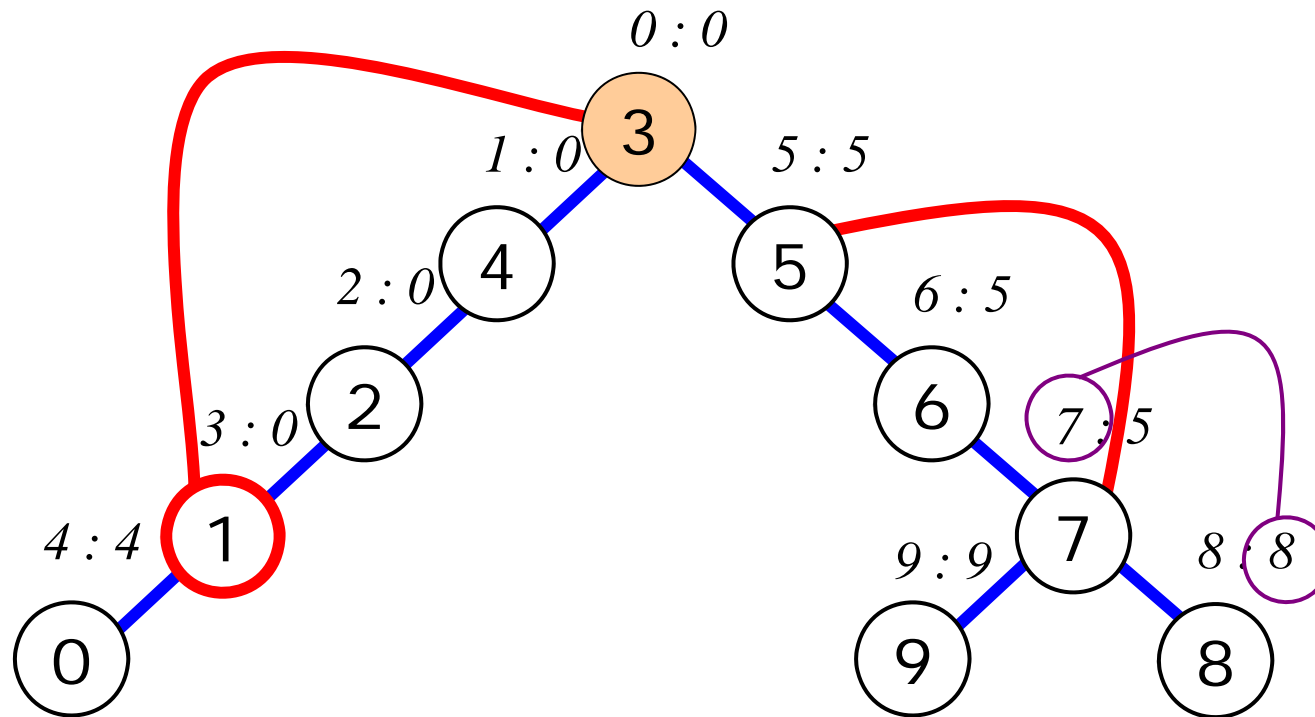
4.5 Biconnected Components

- Determine articulation points
 - At 4, $\text{low}(w) \geq \text{dfn}(u)$?



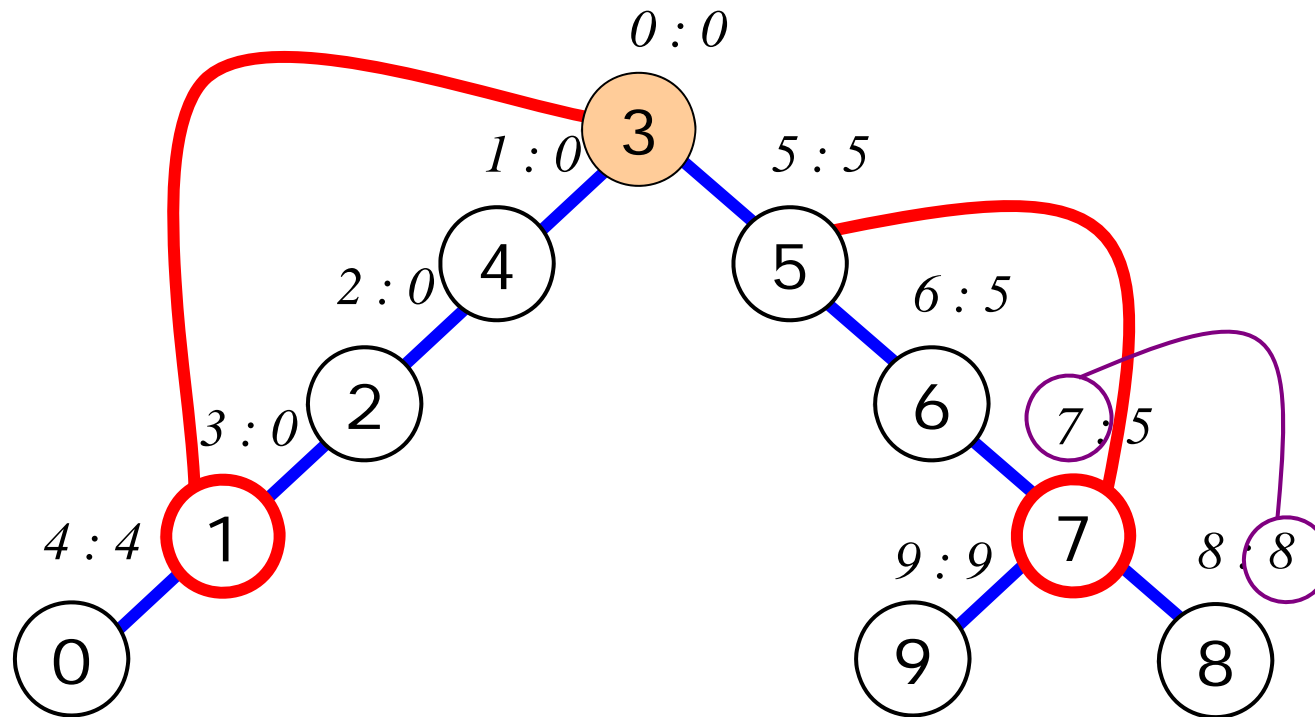
4.5 Biconnected Components

- Determine articulation points
 - At 7, $\text{low}(w) \geq \text{dfn}(u)$?



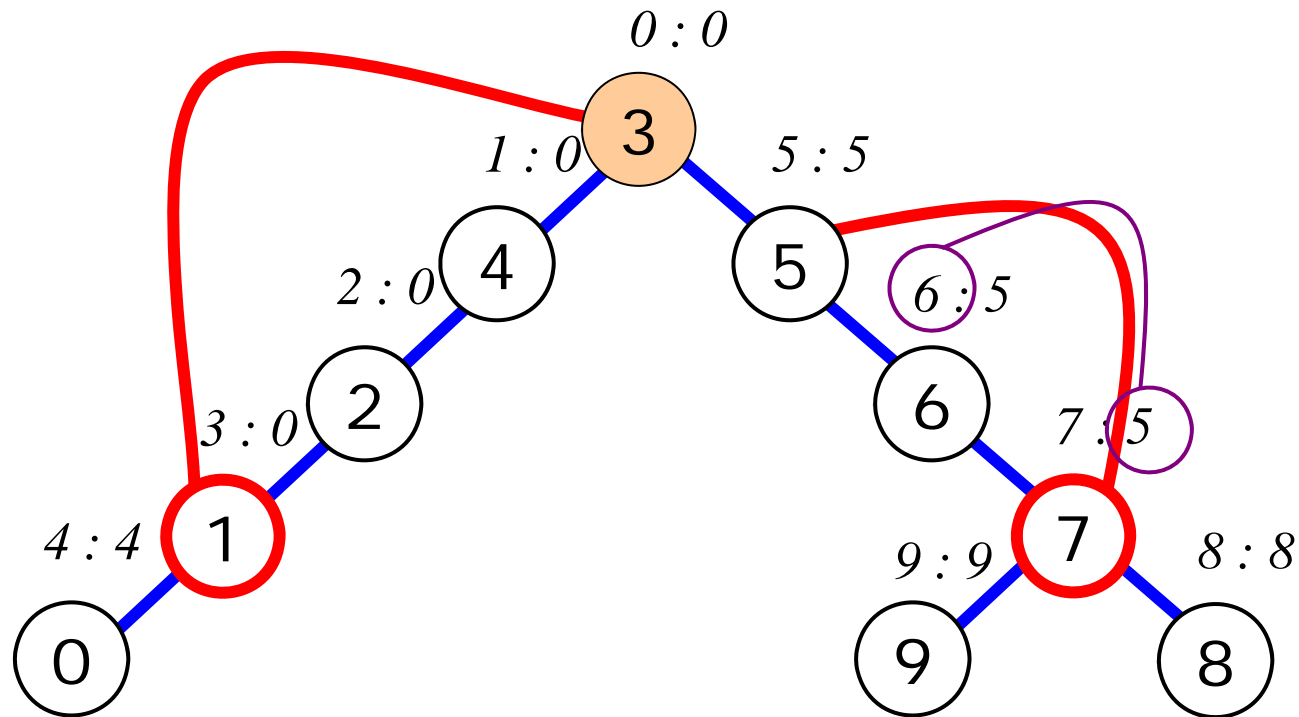
4.5 Biconnected Components

- Determine articulation points
 - At 7, $\text{low}(w) \geq \text{dfn}(u)$?



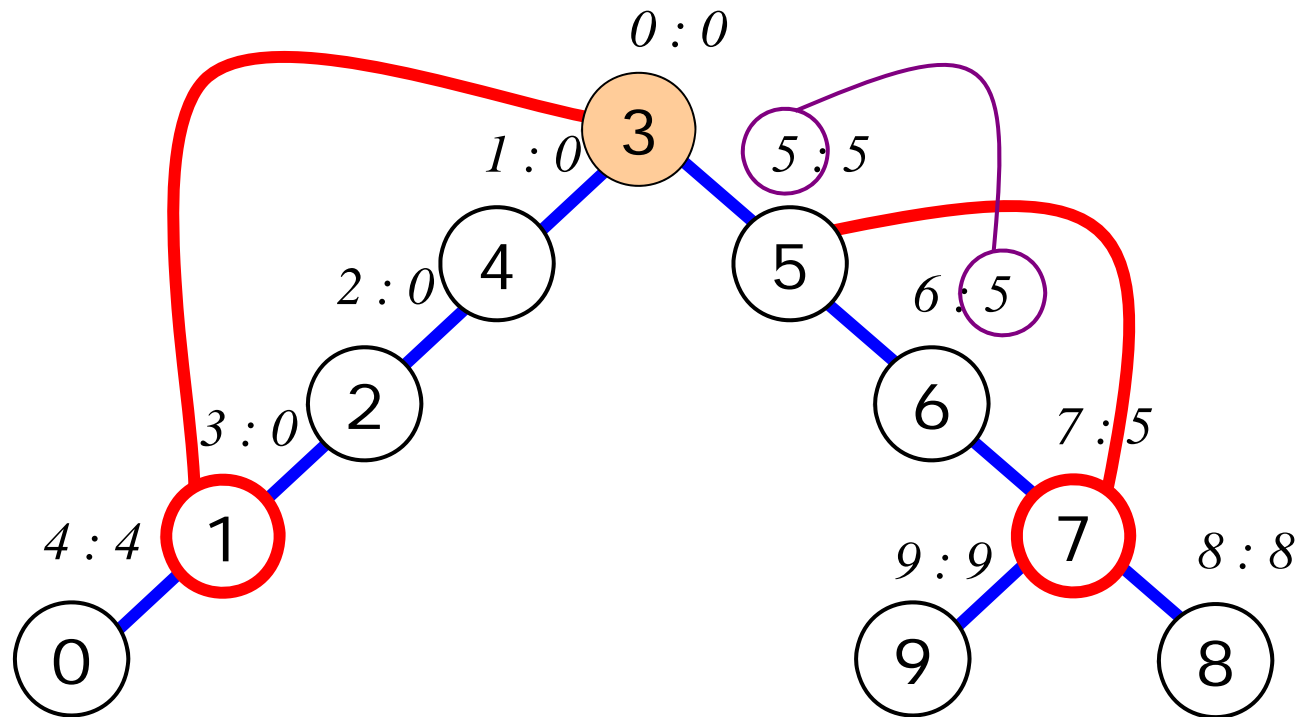
4.5 Biconnected Components

- Determine articulation points
 - At 6, $\text{low}(w) \geq \text{dfn}(u)$?



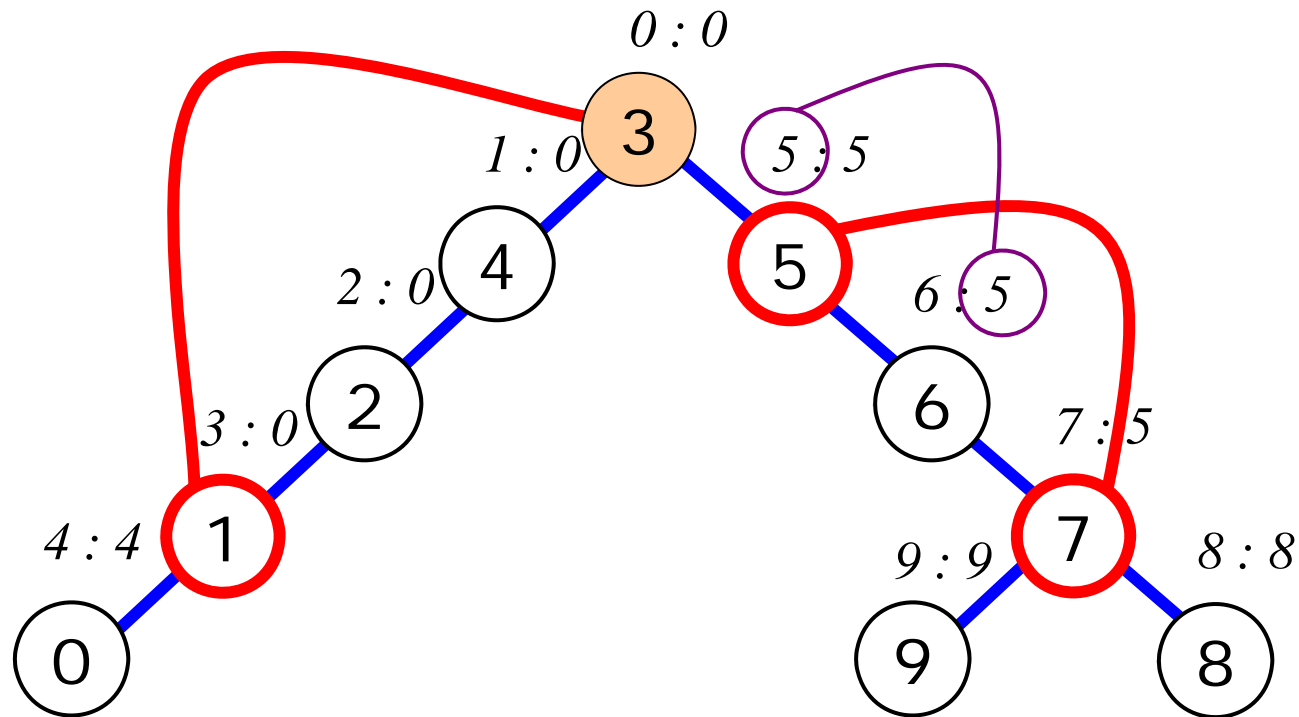
4.5 Biconnected Components

- Determine articulation points
 - At 5, $\text{low}(w) \geq \text{dfn}(u)$?



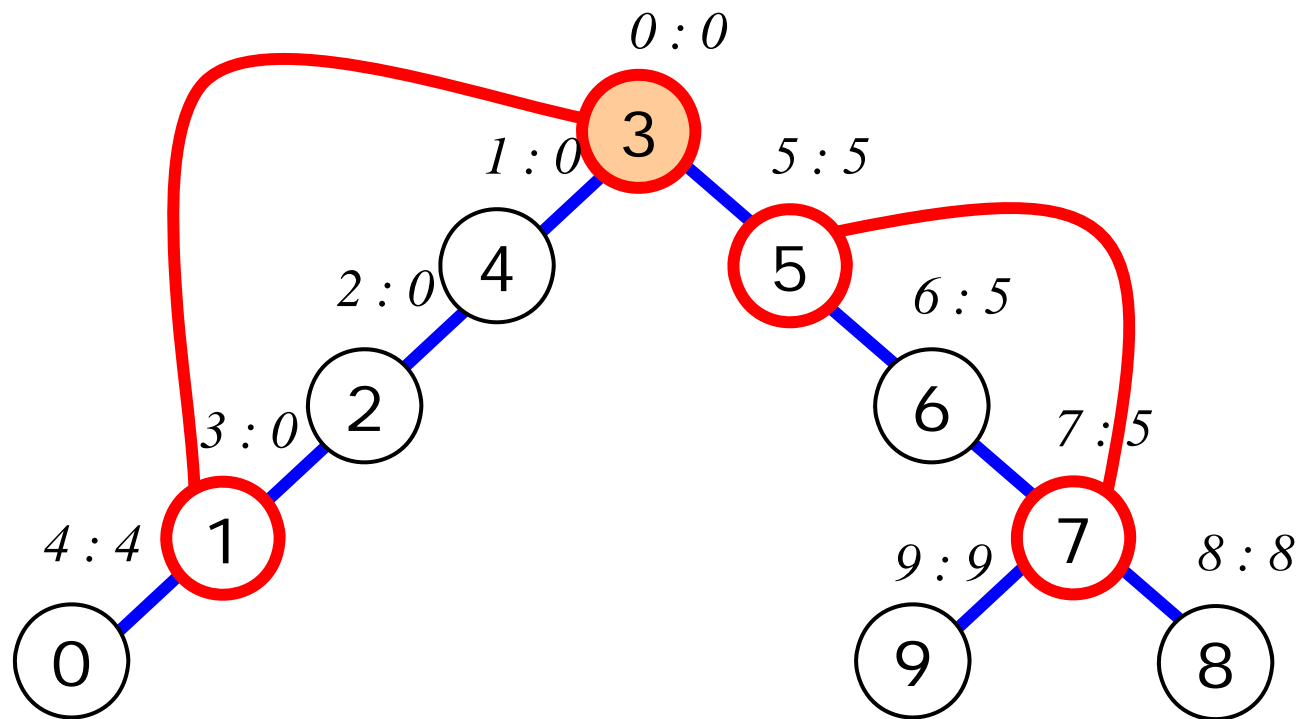
4.5 Biconnected Components

- Determine articulation points
 - At 5, $\text{low}(w) \geq \text{dfn}(u)$?



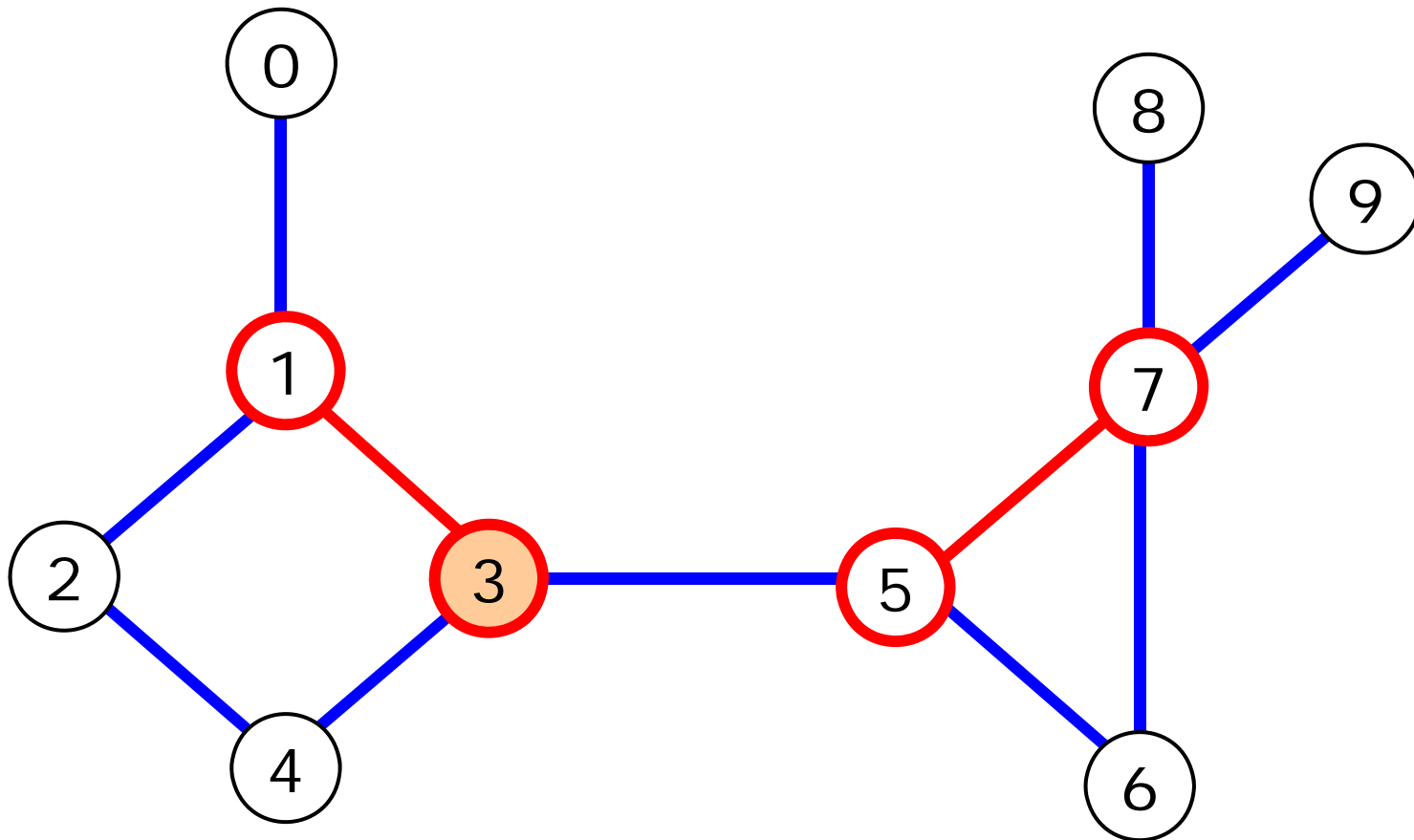
4.5 Biconnected Components

- Total articulation points
 - $\{u \mid \text{low}(w) \geq \text{dfn}(u)\} + \text{root}$



4.5 Biconnected Components

- Total articulation points
 - $\{u \mid \text{low}(w) \geq \text{dfn}(u)\} + \text{root}$



4.5 Biconnected Components

- DFS (1) → computing dfn
 - Build an array for dfn & low, instead of visit
 - Initialize dfn[v] = -1 (visit[v] = FALSE)
 - Call dfs (u, v) instead of dfs (u) (v: parent of u)
 - Declare num and initialize to 0

```
int visit[MAX_VERTEX]; // FALSE로 초기화
void dfs ( u )
{
    visit[u] = TRUE;

    for ( w = graph[u]; w; w = w->link )
        if ( !visit[w] )
            dfs ( w );
}
```

```
int dfn[MAX_VERTEX]; // -1로 초기화
int low[MAX_VERTEX]; // -1로 초기화
int num = 0;
void dfs1 ( u, v ) // v는 u의 부모
{
    dfn[u] = num++;

    for ( w = graph[u]; w; w = w->link )
        if ( dfn[w] < 0 )
            dfs1 ( w, u );
}
```

4.5 Biconnected Components

- DFS (2) → computing low

```
void dfs2 ( u, v ) // v는 u의 부모
{
    dfn[u] = low[u] = num++;

    for ( w = graph[u]; w; w = w->link )
        if ( dfn[w] < 0 ) {
            dfs2 ( w, u );
            low[u] = min (low[u], low[w]);
        }
        else if ( w != v )
            low[u] = min (low[u], low[w]);
}
```

4.5 Biconnected Components

- DFS (3) → computing articulation point

```
void dfs3 ( u, v ) // v는 u의 부모
{
    dfn[u] = low[u] = num++;

    for ( w = graph[u]; w; w = w->link )
        if ( dfn[w] < 0 ) {
            dfs3 ( w, u );
            low[u] = min (low[u], low[w]);
            if ( low[w] >= dfn[u] )
                print ("u: articulation point");
        }
    else if ( w != v )
        low[u] = min (low[u], low[w]);
}
```

4.5 Biconnected Components

- DFS (4) → biconnected component

```
void dfs4 ( u, v ) // v는 u의 부모
{
    dfn[u] = low[u] = num++;

    for ( w = graph[u]; w; w = w->link )
        if ( dfn[w] < 0 ) {
            push (u, w);
            dfs4 ( w, u );
            low[u] = min (low[u], low[w]);
            if ( low[w] >= dfn[u] ) {
                print("new bicon:");
                do {
                    pop (&x, &y);
                    print ( <x, y> );
                } while ( x != u || y != w );
            }
        }
    else if ( w != v )
        low[u] = min (low[u], low[w]);
}
}
```

4.5 Biconnected Components

- Time complexity
 - Time complexity of DFS $\rightarrow O(n + m)$
 - Time complexity of BCC $\rightarrow O(n + m)$

All about graph

Type	Purpose	Operations	Performance
DFS	Traverse all vertices	Visiting all vertices & visiting all edges	$O(n) + O(m)$
SCC	Finding SCC	DFS on G^R and G	$O(\text{DFS})$
BCC	Finding BCC	DFN & Low \rightarrow Articulation Point \rightarrow BCC in DFS	$O(\text{DFS})$
BFS			
Dijkstra			
Floyd			
Kruskal (Greedy)			
Prim (Greedy)			
MultiStage (Dynamic)			