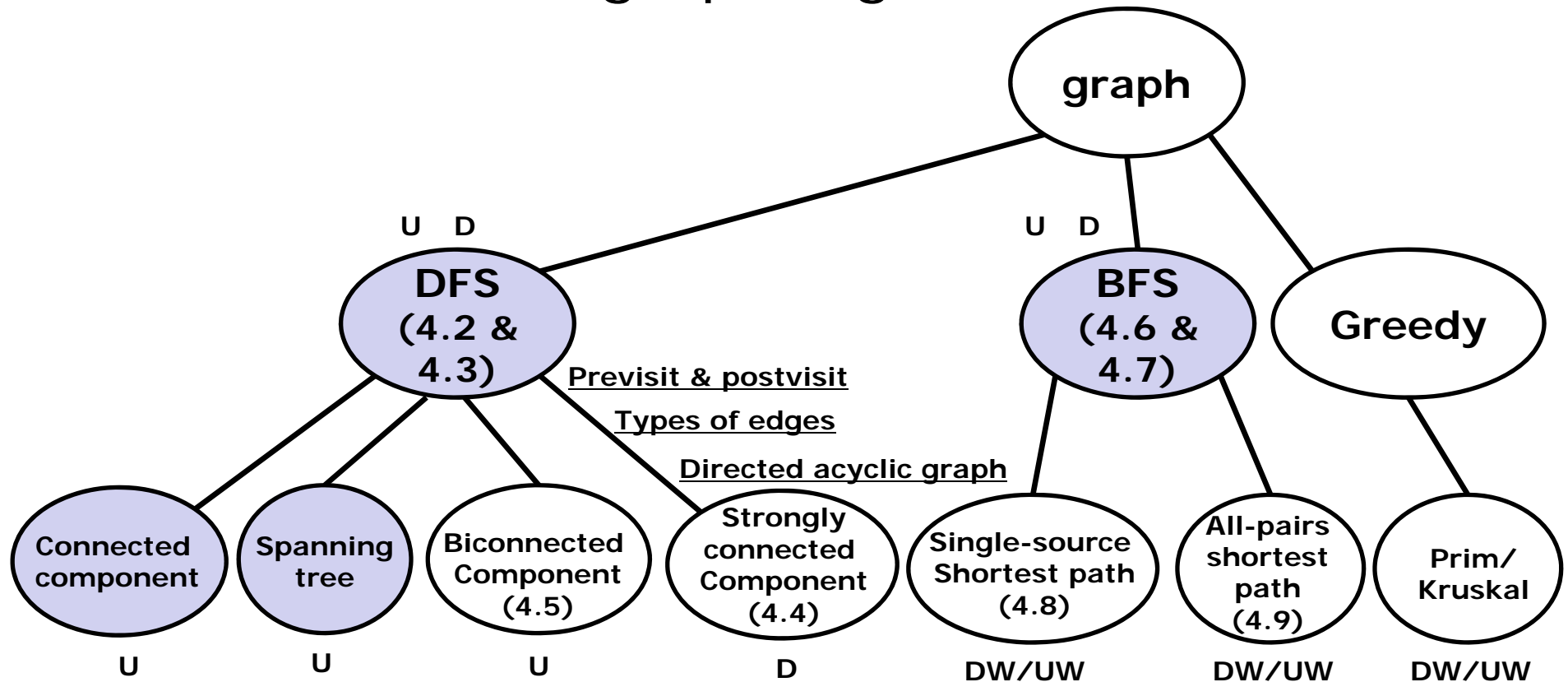

“본 강의 동영상 및 자료는 대한민국 저작권법을 준수합니다. 본 강의 동영상 및 자료는 상명대학교 재학생들의 수업목적으로 제작·배포되는 것이므로, 수업목적으로 내려받은 강의 동영상 및 자료는 수업목적 이외에 다른 용도로 사용할 수 없으며, 다른 장소 및 타인에게 복제, 전송하여 공유할 수 없습니다. 이를 위반해서 발생하는 모든 법적 책임은 행위 주체인 본인에게 있습니다.”

4.6 Distance

Classification of graph algorithms

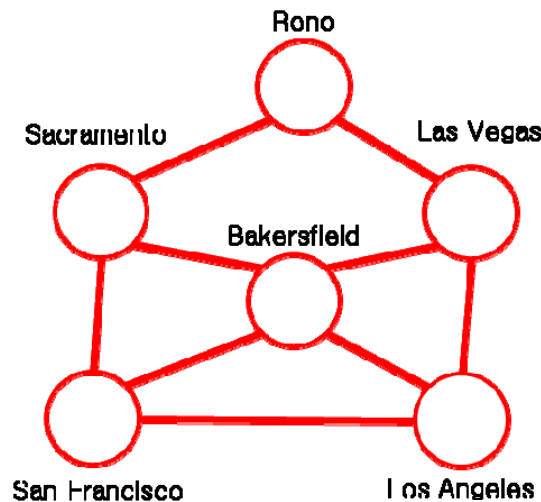


4.6 Distance

- Why graph?
 - Ex) 6 cities in American west coast
 - {Sacramento, San Francisco, Reno, Bakersfield, Los Angeles, Las Vegas}
 - In San Francisco, we can go to Sacramento, Bakersfield & Los Angeles
 - In Sacramento, we can go to San Francisco, Bakersfield & Reno
 - In Reno, we can go to Sacramento & Las Vegas
 - In Bakersfield, we can go to San Francisco, Reno, Los Angeles & Las Vegas
 - In Los Angeles, we can go to San Francisco, Bakersfield & Las Vegas
 - In Las Vegas, we can go to Reno, Bakersfield & Los Angeles

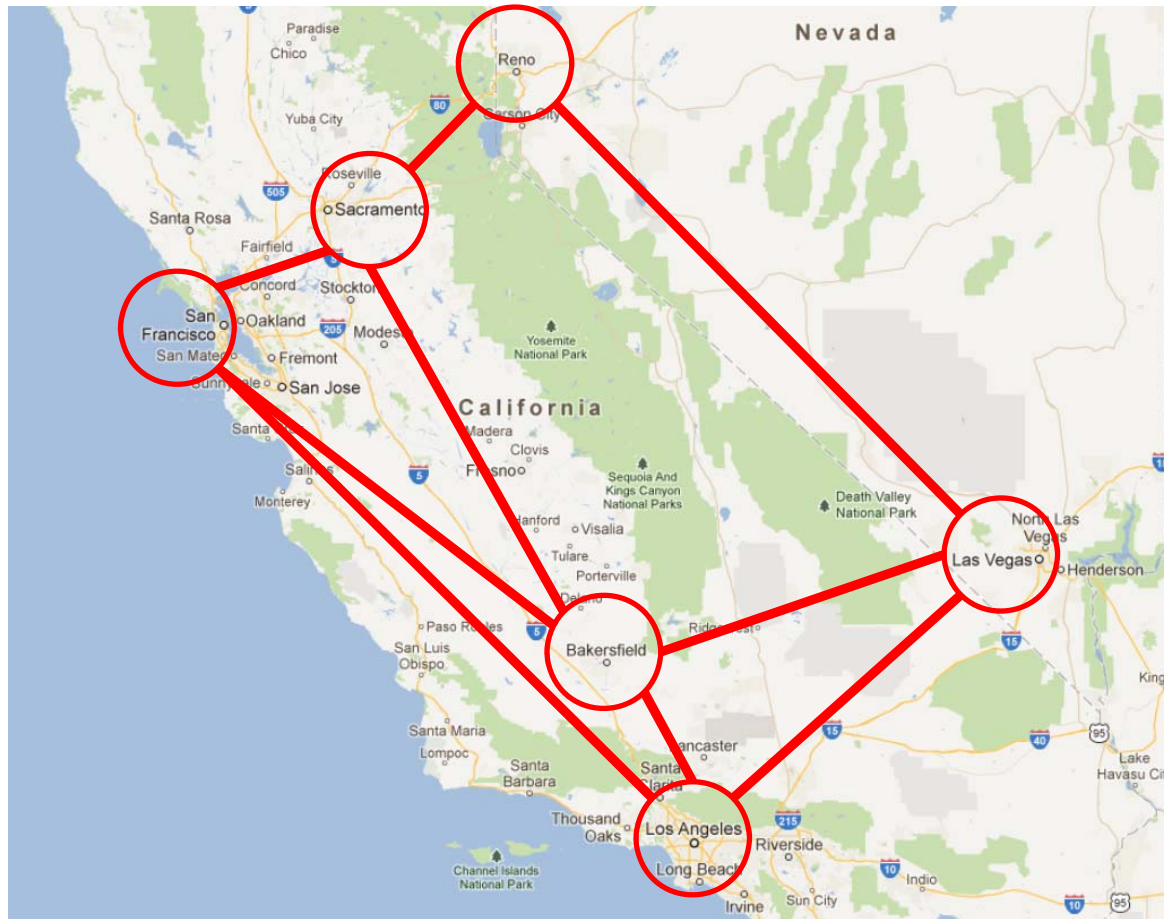
4.6 Distance

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 - In Los Angeles, we can go to San Francisco, Bakersfield & Las Vegas
 - In Las Vegas, we can go to Reno, Bakersfield & Los Angeles



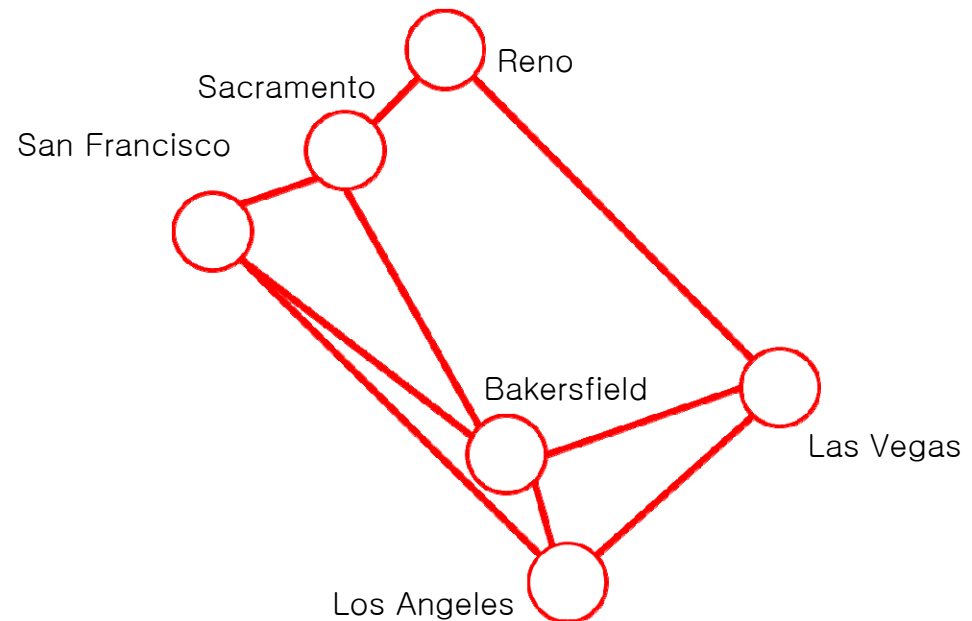
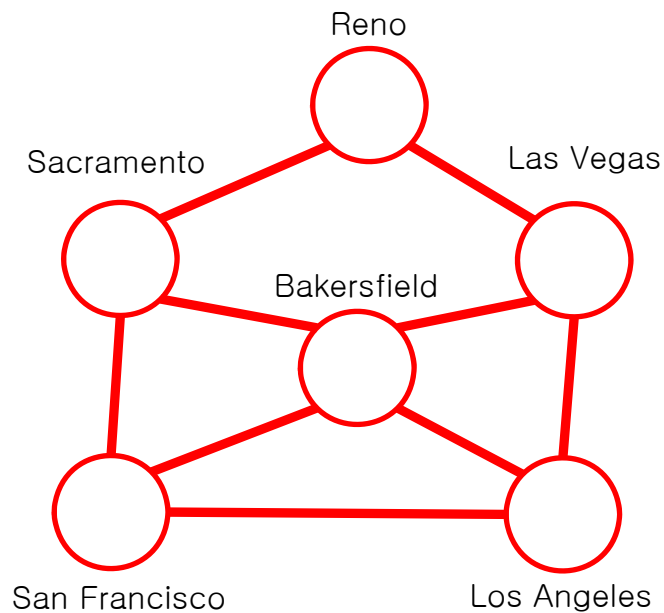
4.6 Distance

- In real world,
 - Distances between two nodes are not identical



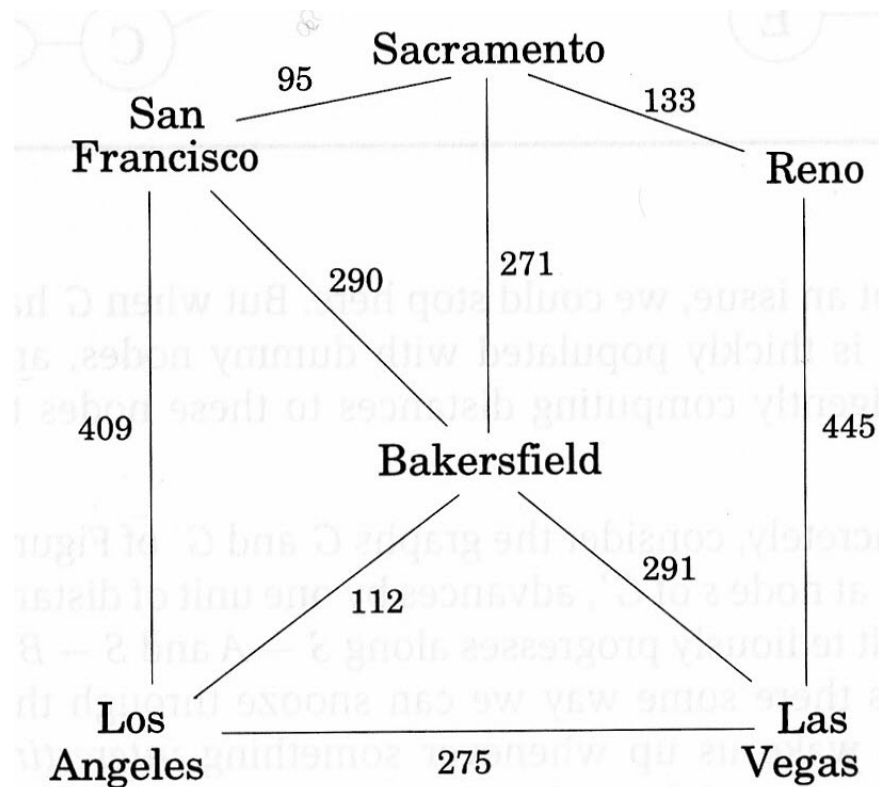
4.6 Distance

- Why weighted graph?
 - Edges on a graph have different weights (distance)
 - Distance
 - The length of the shortest path between them



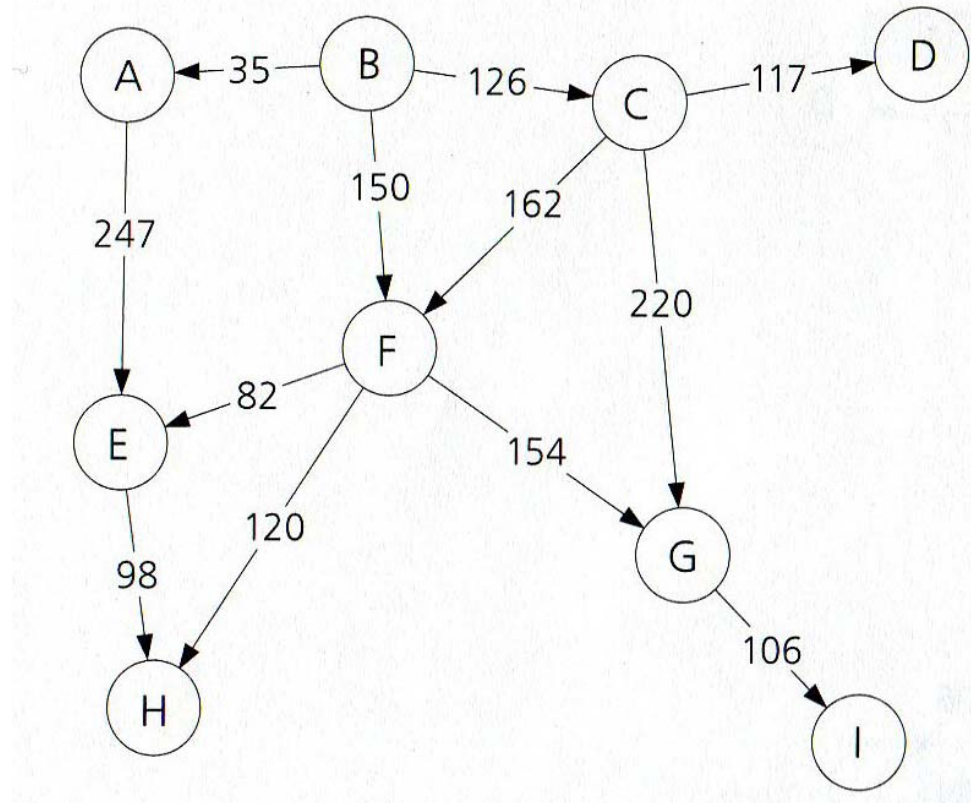
4.6 Distance

- Weighted graph
 - Distance between two nodes
 - The length of the shortest path between them



4.6 Distance

- Finding shortest path between two nodes in a graph
 - Ex) What is the shortest path from B to G?



4.7 Breadth-first search

- Basic strategy
 - In visiting a vertex v ,
 - Mark all the adjacent vertices as visited
 - Add the vertices to the queue
 - Use a queue
- Compare to Depth-first search
 - Visit connected nodes as far as possible
 - Implemented using a recursive call
 - Use a stack

4.7 Breadth-first search

- Recall: Queue
 - Property of Queue
 - First-In First-Out
 - Important points of Queue
 - Front
 - Rear
 - Operations of Queue
 - Insert ()
 - Remove ()

4.7 Breadth-first search

- Algorithm

```
procedure bfs ( G, s )
Input: Graph  $G = (V, E)$ ,  $s$  = start vertex
Output: For all vertices  $u$  reachable from  $s$ ,  $\text{dist}[u]$  is
set to the distance from  $s$  to  $u$ 

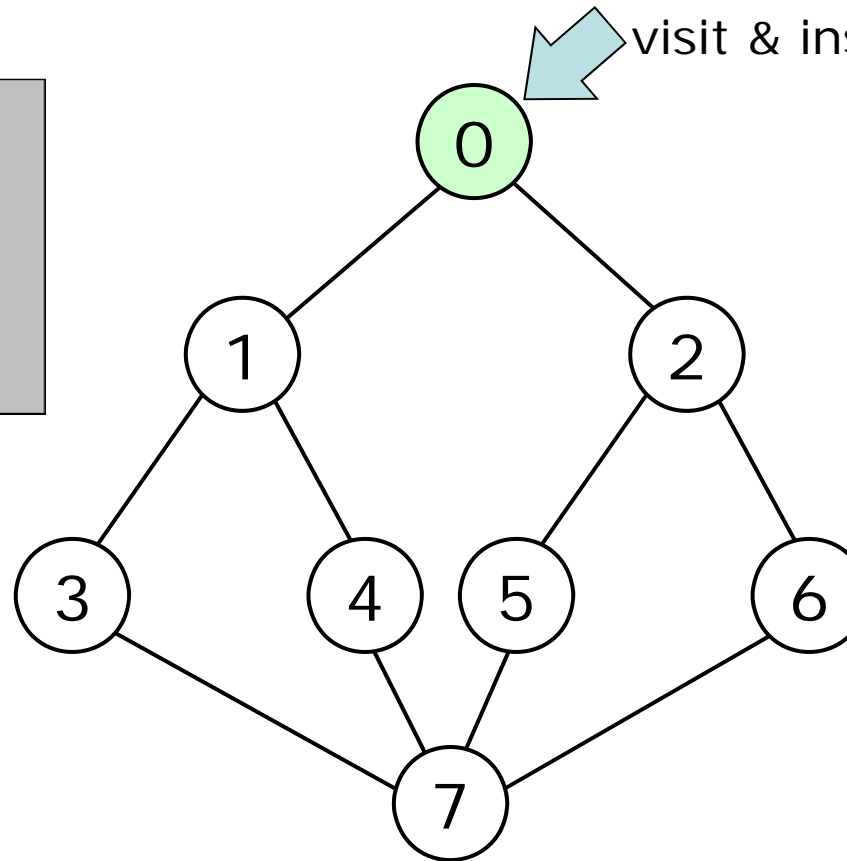
For all  $u \in V$ 
     $\text{dist}[u] = \infty$ ;

 $\text{dist}[s] = 0$ ;
 $Q.\text{insert} ( s )$ ;
while (  $!Q.\text{is\_empty} ( )$  ) {
     $u = Q.\text{remove} ( )$ ;
    for all edges  $(u, v) \in E$ 
        if (  $\text{dist}[v] == \infty$  )
             $Q.\text{insert} ( v )$ ;
             $\text{dist}[v] = \text{dist}[u] + 1$ ;
}
```

4.7 Breadth-first search

– Example

```
dist[s] = 0;
Q.insert ( s );
while ( !Q.is_empty ( ) ) {
    u = Q.remove ( );
    for all edges (u, v) ∈ E
        if ( dist[v] == ∞ )
            Q.insert ( v );
            dist[v] = dist[u] + 1;
}
```



0	0
1	∞
2	∞
3	∞
4	∞
5	∞
6	∞
7	∞

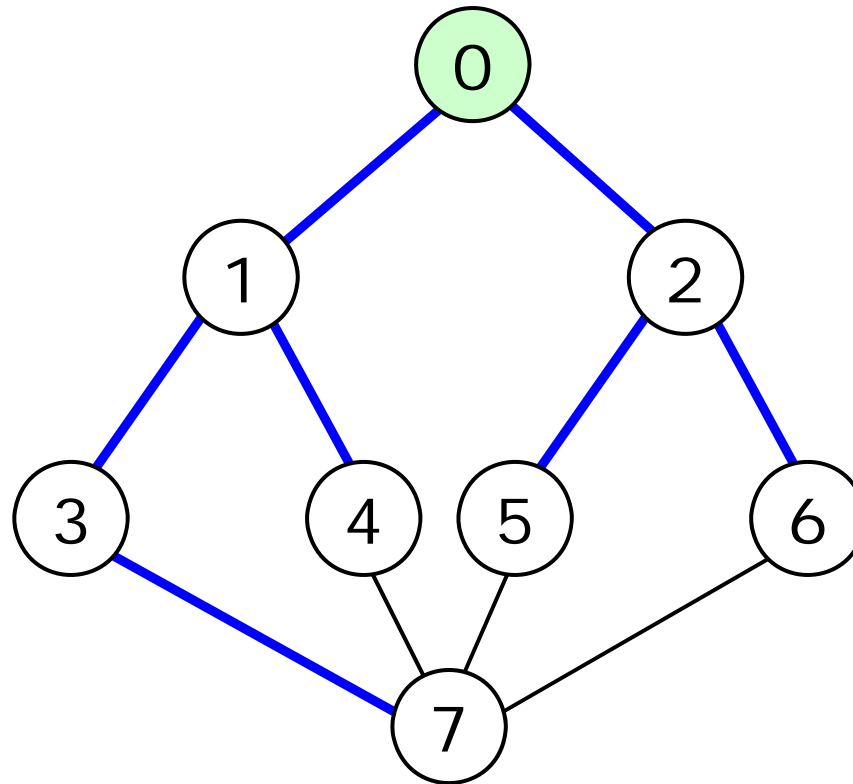
dist[u]

queue

0							
---	--	--	--	--	--	--	--

4.7 Breadth-first search

– Example

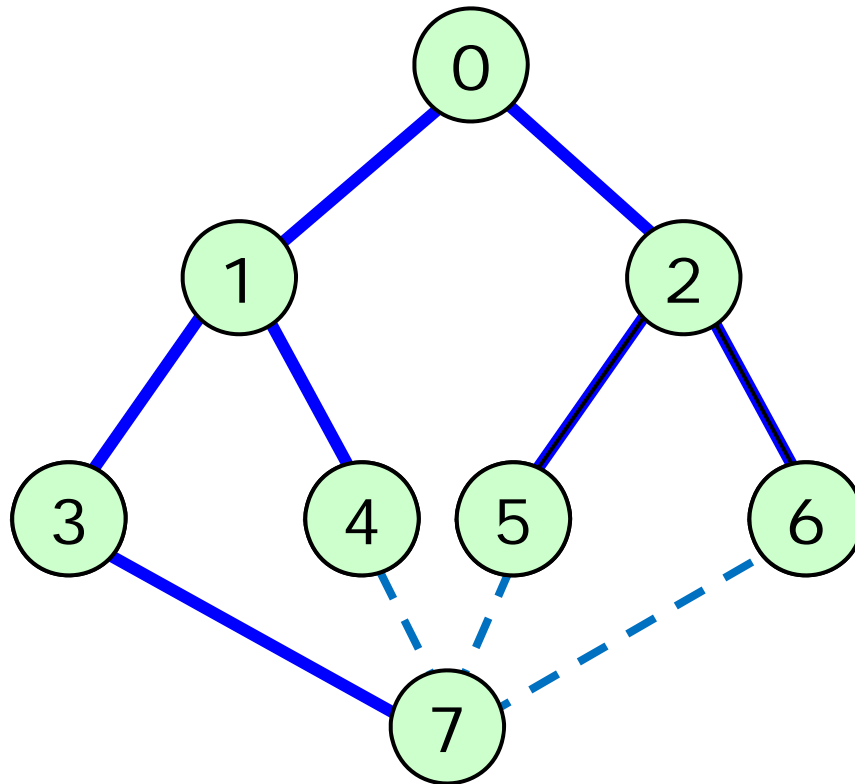


0	0
1	1
2	1
3	2
4	2
5	2
6	2
7	3

dist[u]

4.7 Breadth-first search

– Breadth-first spanning tree



4.7 Breadth-first search

– Practical implementation using STL

```
#include <stdio.h>
#include <vector>
#include <queue>
#include <algorithm>
using namespace std;

vector<int> edge[10001];
int dist[10001];

/*visit 변수 초기화 함수*/
void initDist(){
    for (int i = 0; i < 10001; i++){
        dist[i] = inf;
    }
}
```

4.7 Breadth-first search

– Practical implementation using STL

```
int main()
{
    //dist 변수 초기화 후 bfs
    initDist();
    for ( i = 1; i <= n; i++ ) {
        if ( dist[i] == inf )
            bfs(i);
    }
    return 0;
}
```


4.7 Breadth-first search

```
void bfs(int s)
{
    queue<int> q;
    int now, k, next;

    q.push(s);
    dist[s] = 0;
    while (!q.empty()){
        now = q.front();
        q.pop();
        printf("%d ", now);
        for (k = 0; k < edge[now].size(); k++){
            next = edge[now][k];
            if (dist[next] == inf){
                dist[next] = dist[now] + 1;
                q.push(next);
            }
        }
    }
}
```

4.7 Breadth-first search

- 시간 복잡도 = $O(n + m)$

```
dist[s] = 0;
Q.insert ( s );
while ( !Q.is_empty ( ) ) {
    u = Q.remove ( );
    for all edges (u, v) ∈ E
        if ( dist[v] == ∞ )
            Q.insert ( v );
            dist[v] = dist[u] + 1;
}
```

$O(m)$

u에 연결된 모든 edge에 대해서

모든 u에 대해서

```
initVisit();
for ( i = 1; i <= n; i++ ) {
    if ( dist[i] == inf )
        bfs(i);
}
```

$O(n)$

4.7 Breadth-first search

- bfs에 대한 설명 중 잘못된 것을 모두 고르시오.
 - (a) 그래프에서 두 vertex 사이의 거리를 edge의 weight로 표현할 수 있다.
 - (b) edge의 weight는 두 vertex 사이의 거리를 나타내는 값이기 때문에 음수가 올 수 없다.
 - (c) bfs는 stack을 이용하고 dfs는 queue를 이용한다.
 - (d) bfs의 시간복잡도는 dfs의 시간복잡도와 같다.