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“본 강의 동영상 및 자료는 대한민국 저작권법을 준수합니다. 본 강의 동영상 및 자료는 상명대학교 재학생들의 수업목적으로 제작·배포되는 것이므로, 수업목적으로 내려받은 강의 동영상 및 자료는 수업목적 이외에 다른 용도로 사용할 수 없으며, 다른 장소 및 타인에게 복제, 전송하여 공유할 수 없습니다. 이를 위반해서 발생하는 모든 법적 책임은 행위 주체인 본인에게 있습니다.”

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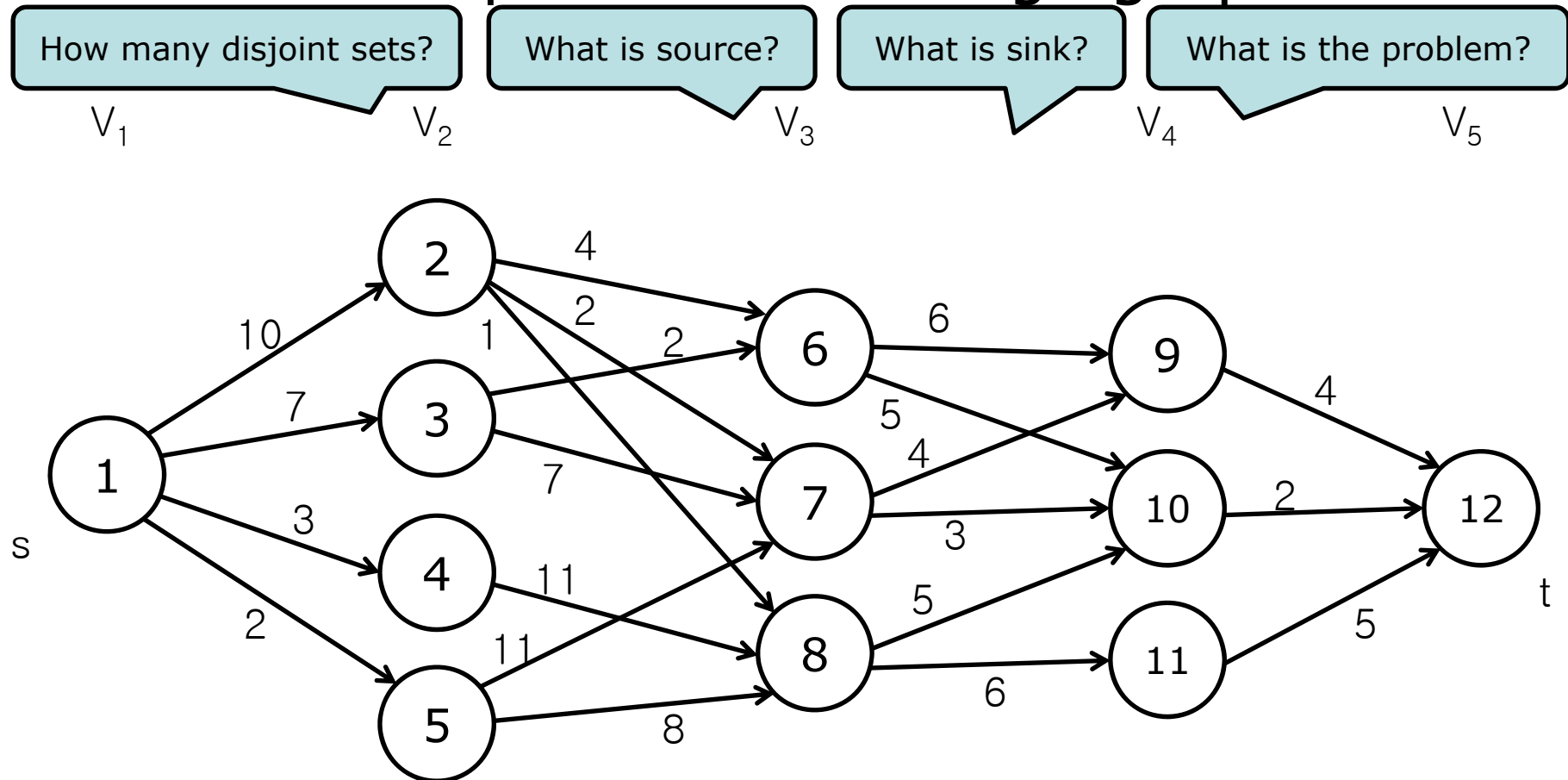
## 6.3 Multistage graph

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- Definition of multistage graph
    - A directed graph
      - The vertices are partitioned into  $k \geq 2$  disjoint sets  $V_i$ ,  $1 \leq i \leq k$ .
      - If  $\langle u, v \rangle$  is an edge in  $E$ , then  $u \in V_i$  and  $v \in V_{i+1}$ , for some  $i$ .
      - $|V_1| = |V_k| = 1$ .
      - If  $s \in V_1$ , then  $s$  is the source
      - If  $t \in V_k$ , then  $t$  is the sink
    - Let  $c(i, j)$  be the cost of edge  $\langle i, j \rangle$ .
      - The cost of a path from  $s$  to  $t$  is the sum of the costs of the edges on the path.
      - The multistage graph problem is to find a minimum cost path from  $s$  to  $t$ .
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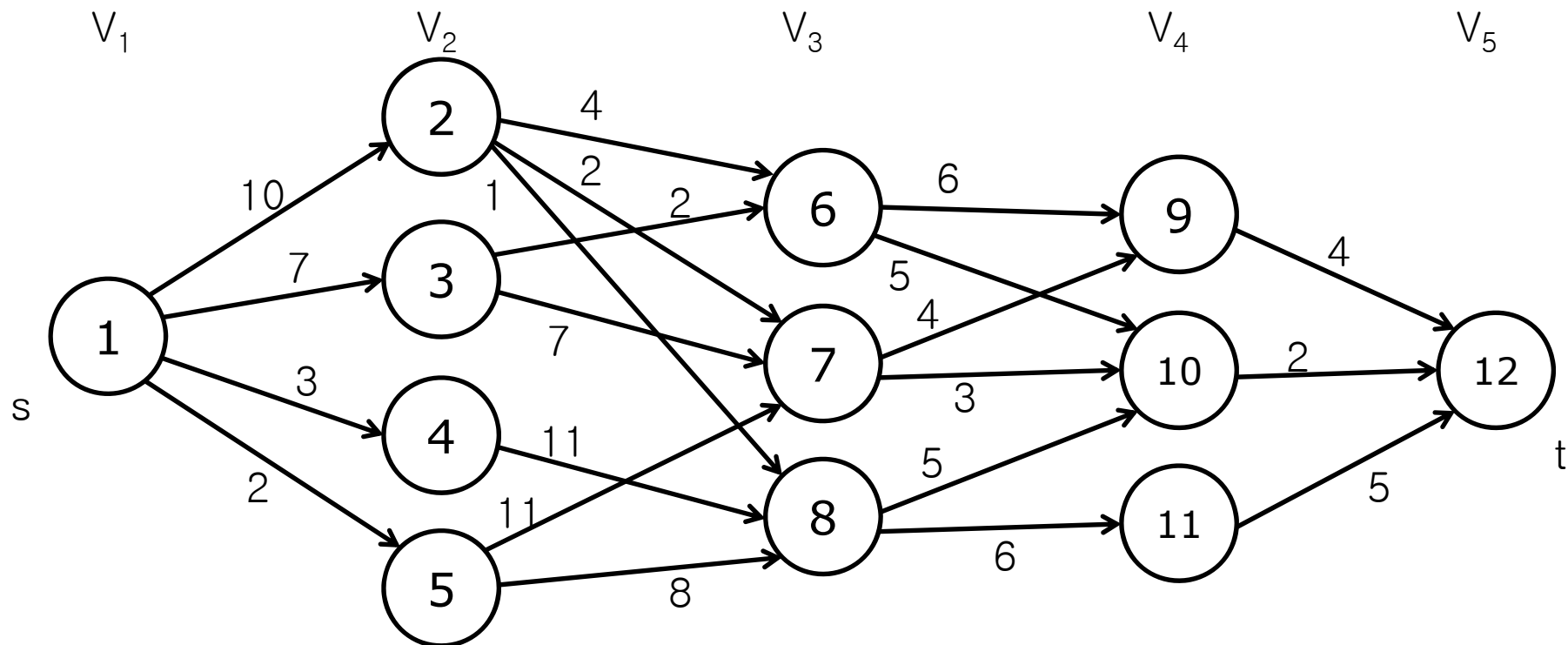
## 6.3 Multistage graph

- Definition of multistage graph
  - An example of a multistage graph



## 6.3 Multistage graph

- Problem
  - Find an optimal path from source to sink



## 6.3 Multistage graph

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- Strategy
  - $P(i, j)$ 
    - minimum cost path from vertex  $j$  in  $V_i$  to sink
  - $COST(i, j)$ 
    - the cost of  $P(i, j)$

Algorithm

$$COST(i, j) = \min_{l \in V_{i+1}} \{c(j, l) + COST(i+1, l)\}$$

- What is the problem to solve?
    - $P(1, 1)$
    - $COST(1, 1)$
-

## 6.3 Multistage graph

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- Algorithm

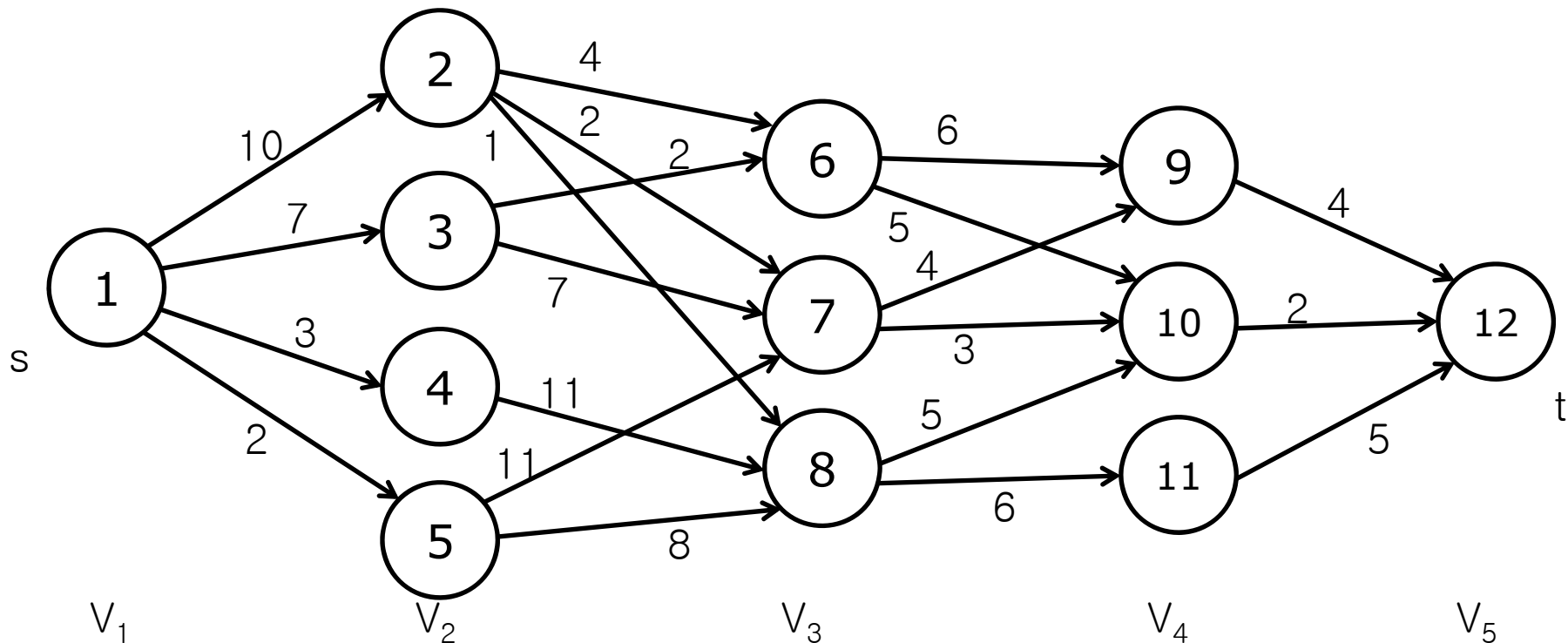
```
void MULTIGRAPH ( Graph G)
{
    int **COST, **D;

    estimate COST(k - 1);
    for ( i = k - 2; i >= 1; i-- ) {
        for all vertices j at stage i
            COST(i, j)  $\leftarrow$  min(cost(j, 1) + COST(i+1, 1));
            D(i, j)  $\leftarrow$  1;
    }
}
```

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## 6.3 Multistage graph

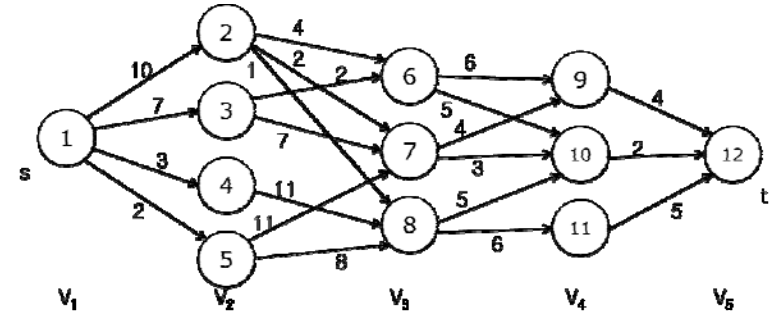
- Example (At  $V_1$ )
  - $\text{COST}(1, 1) =$





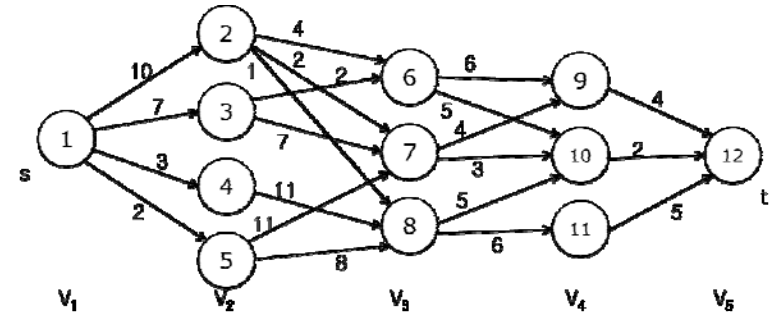
## 6.3 Multistage graph

- Example (At  $V_4$ )
  - $\text{COST}(4, 9) =$
  - $\text{COST}(4, 10) =$
  - $\text{COST}(4, 11) =$



## 6.3 Multistage graph

- Example (At  $V_4$ )
  - $\text{COST}(4, 9) = 4$
  - $\text{COST}(4, 10) = 2$
  - $\text{COST}(4, 11) = 5$



## 6.3 Multistage graph

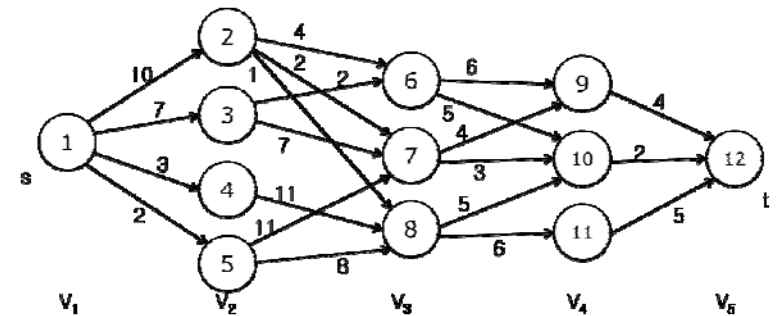
- Example (At  $V_3$ )

- $COST(3, 6) =$

- $COST(3, 7) =$

- $COST(3, 8) =$

$$COST(i, j) = \min_{l \in V_{i+1}} \{c(j, l) + COST(i+1, l)\}$$



## 6.3 Multistage graph

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- Example (At  $V_3$ )
    - $\text{COST}(3, 6) = \min\{6 + \text{COST}(4, 9),$   
 $5 + \text{COST}(4, 10)\}$
    - $\text{COST}(3, 7) = \min\{4 + \text{COST}(4, 9),$   
 $3 + \text{COST}(4, 10)\}$
    - $\text{COST}(3, 8) = \min\{5 + \text{COST}(4, 10),$   
 $6 + \text{COST}(4, 11)\}$
-

## 6.3 Multistage graph

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- Example (At  $V_3$ )
    - $\text{COST}(3, 6) = \min\{6 + \text{COST}(4, 9),$   
 $\quad\quad\quad 5 + \text{COST}(4, 10)\}$   
 $= \min\{6 + 4, 5 + 2\} = 7$
    - $\text{COST}(3, 7) = \min\{4 + \text{COST}(4, 9),$   
 $\quad\quad\quad 3 + \text{COST}(4, 10)\}$   
 $= \min\{4 + 4, 3 + 2\} = 5$
    - $\text{COST}(3, 8) = \min\{5 + \text{COST}(4, 10),$   
 $\quad\quad\quad 6 + \text{COST}(4, 11)\}$   
 $= \min\{5 + 2, 6 + 5\} = 7$
-

## 6.3 Multistage graph

- Example (At  $V_2$ )

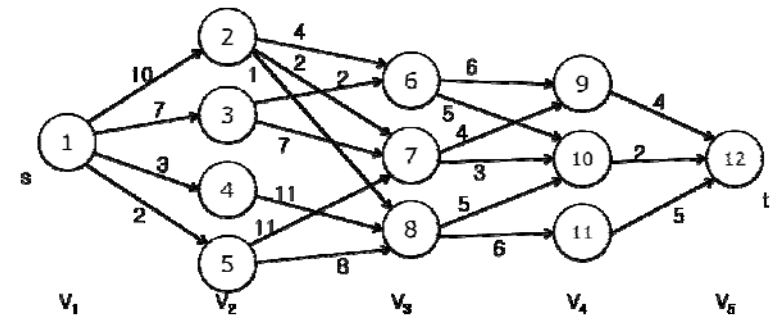
- $COST(2, 2) =$

- $COST(2, 3) =$

- $COST(2, 4) =$

- $COST(2, 5) =$

$$COST(i, j) = \min_{l \in V_{i+1}} \{c(j, l) + COST(i+1, l)\}$$



## 6.3 Multistage graph

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- Example (At  $V_2$ )
  - $COST(2, 2) = \min\{4 + COST(3, 6),$   
 $2 + COST(3, 7),$   
 $1 + COST(3, 8)\}$
  - $COST(2, 3) = \min\{2 + COST(3, 6),$   
 $7 + COST(3, 7)\}$
  - $COST(2, 4) = \min\{11 + COST(3, 8)\}$
  - $COST(2, 5) = \min\{11 + COST(3, 7),$   
 $8 + COST(3, 8)\}$

$$COST(i, j) = \min_{l \in V_{i+1}} \{c(j, l) + COST(i+1, l)\}$$

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## 6.3 Multistage graph

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- Example (At  $V_2$ )

- $\text{COST}(2, 2) = \min\{4 + \text{COST}(3, 6),$   
 $2 + \text{COST}(3, 7),$   
 $1 + \text{COST}(3, 8)\}$   
 $= \min\{4 + 7, 2 + 5, 1 + 7\} = 7$
  - $\text{COST}(2, 3) = \min\{2 + \text{COST}(3, 6),$   
 $7 + \text{COST}(3, 7)\}$   
 $= \min\{2 + 7, 7 + 5\} = 9$
  - $\text{COST}(2, 4) = \min\{11 + \text{COST}(3, 8)\} = 11 + 7 = 18$
  - $\text{COST}(2, 5) = \min\{11 + \text{COST}(3, 7),$   
 $8 + \text{COST}(3, 8)\}$   
 $= \min\{11 + 5, 8 + 7\} = 15$
-

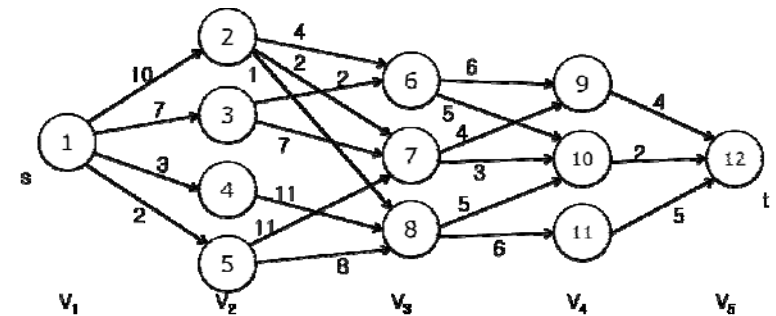


## 6.3 Multistage graph

- Example (At  $V_1$ )
  - $COST(1, 1) =$

–  $P(1, 1) =$

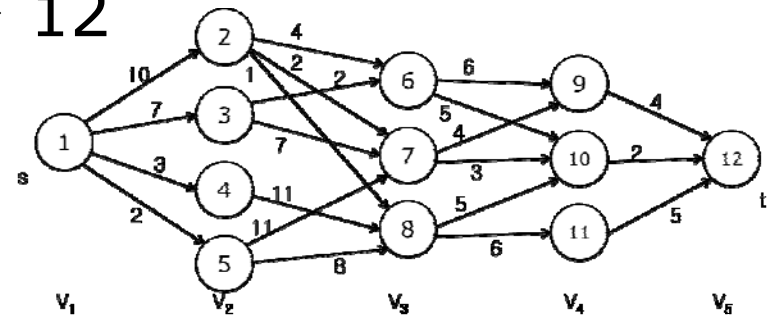
$$COST(i, j) = \min_{l \in V_{i+1}} \{c(j, l) + COST(i+1, l)\}$$



## 6.3 Multistage graph

- Example (At  $V_1$ )
  - $\text{COST}(1, 1) = \min\{10 + \text{COST}(2, 2),$   
 $7 + \text{COST}(2, 3),$   
 $9 + \text{COST}(2, 4),$   
 $2 + \text{COST}(2, 5)\}$

–  $P(1, 1) = 1 \rightarrow 3 \rightarrow 6 \rightarrow 10 \rightarrow 12$

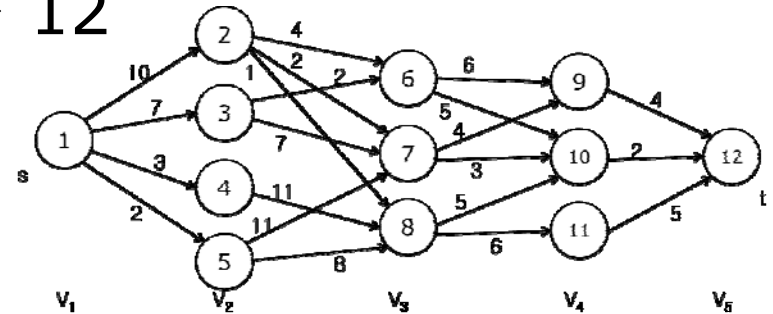


## 6.3 Multistage graph

- Example (At  $V_1$ )

$$\begin{aligned} - \text{COST}(1, 1) &= \min\{10 + \text{COST}(2, 2), \\ &\quad 7 + \text{COST}(2, 3), \\ &\quad 3 + \text{COST}(2, 4), \\ &\quad 2 + \text{COST}(2, 5)\} \\ &= \min\{10 + 7, 7 + 9, 3 + 18, 2 + 15\} \\ &= 16 \end{aligned}$$

$$- P(1, 1) = 1 \rightarrow 3 \rightarrow 6 \rightarrow 10 \rightarrow 12$$



## 6.3 Multistage graph

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- Application: Resource allocation problem
    - Allocating  $n$  units of resources to  $r$  projects
    - $N(i, j)$ 
      - The net profit of allocating  $j$  units of resources to project  $i$ .
    - How to maximize total net profit?
    - Formulate this problem using multistage graph
-

## 6.3 Multistage graph

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- Application: Resource allocation problem
    - Multistage graph
      - $(r+1)$  stage graph problem
      - Stage  $i$  represents project  $i$ .
      - $(n+1)$  vertices  $V(i, j)$  are associated with stage  $i$ .
      - Stage 1 and  $r+1$  has one vertex:  $V(1, 0) = s$  and  $V(r+1, n) = t$ .
      - Vertex  $V(i, j)$  represents the stage in which a total of  $j$  units of resources have been allocated to projects  $1, 2, \dots, i-1$ .
      - The edges are of the form  $\langle V(i, j), V(i+1, l) \rangle$ .
        - The cost of the edges is  $N(i, l-j)$ .
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## 6.3 Multistage graph

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- Example
  - 4 units to 3 teams

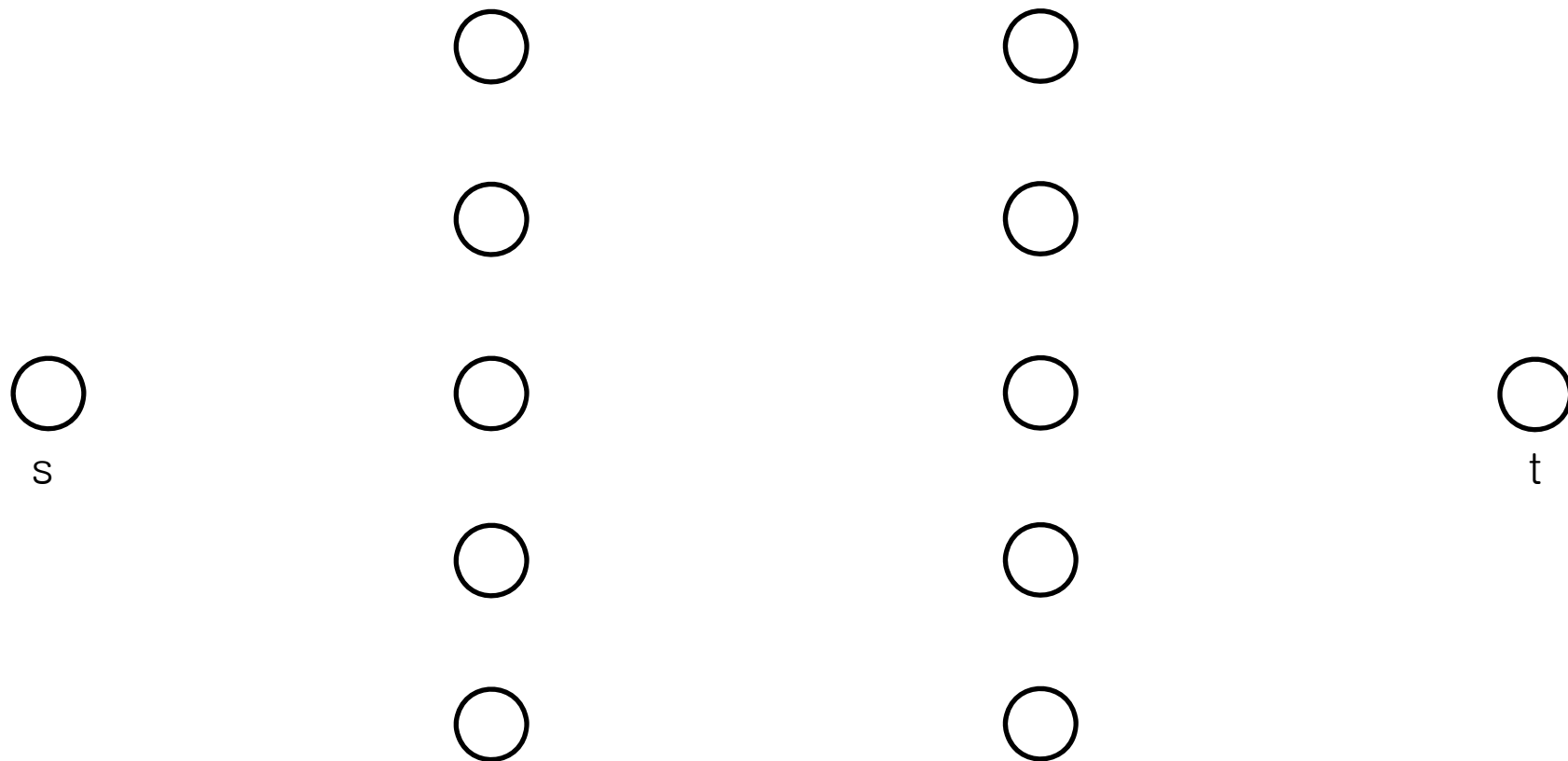
	1 <sup>st</sup> team	2 <sup>nd</sup> team	3 <sup>rd</sup> team
N=0	0	5	3
N=1	10	15	5
N=2	20	20	10
N=3	25	22	15
N=4	25	25	25

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## 6.3 Multistage graph

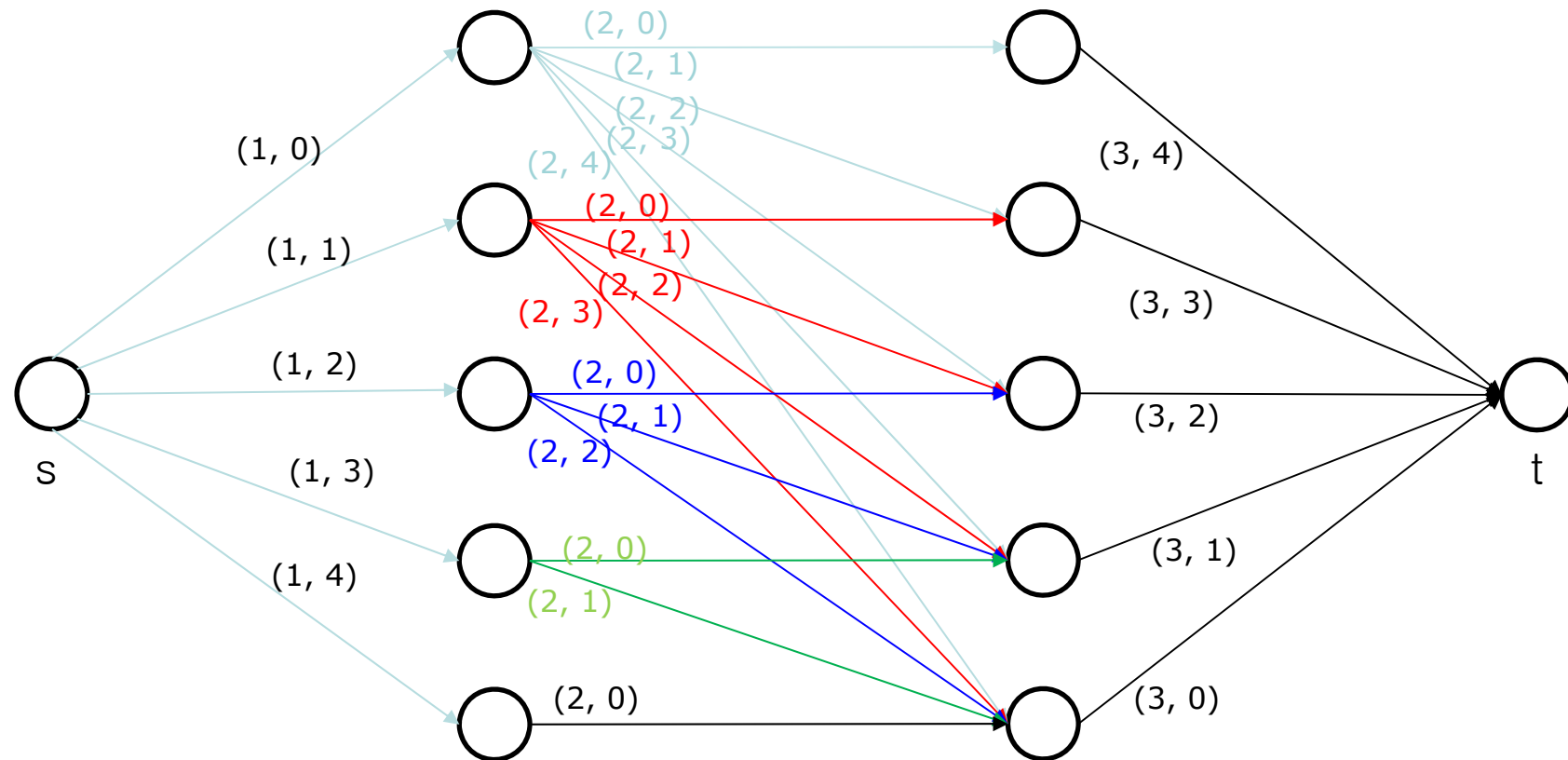
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- Application: Resource allocation problem
  - Multistage graph of 4 units to 3 teams



## 6.3 Multistage graph

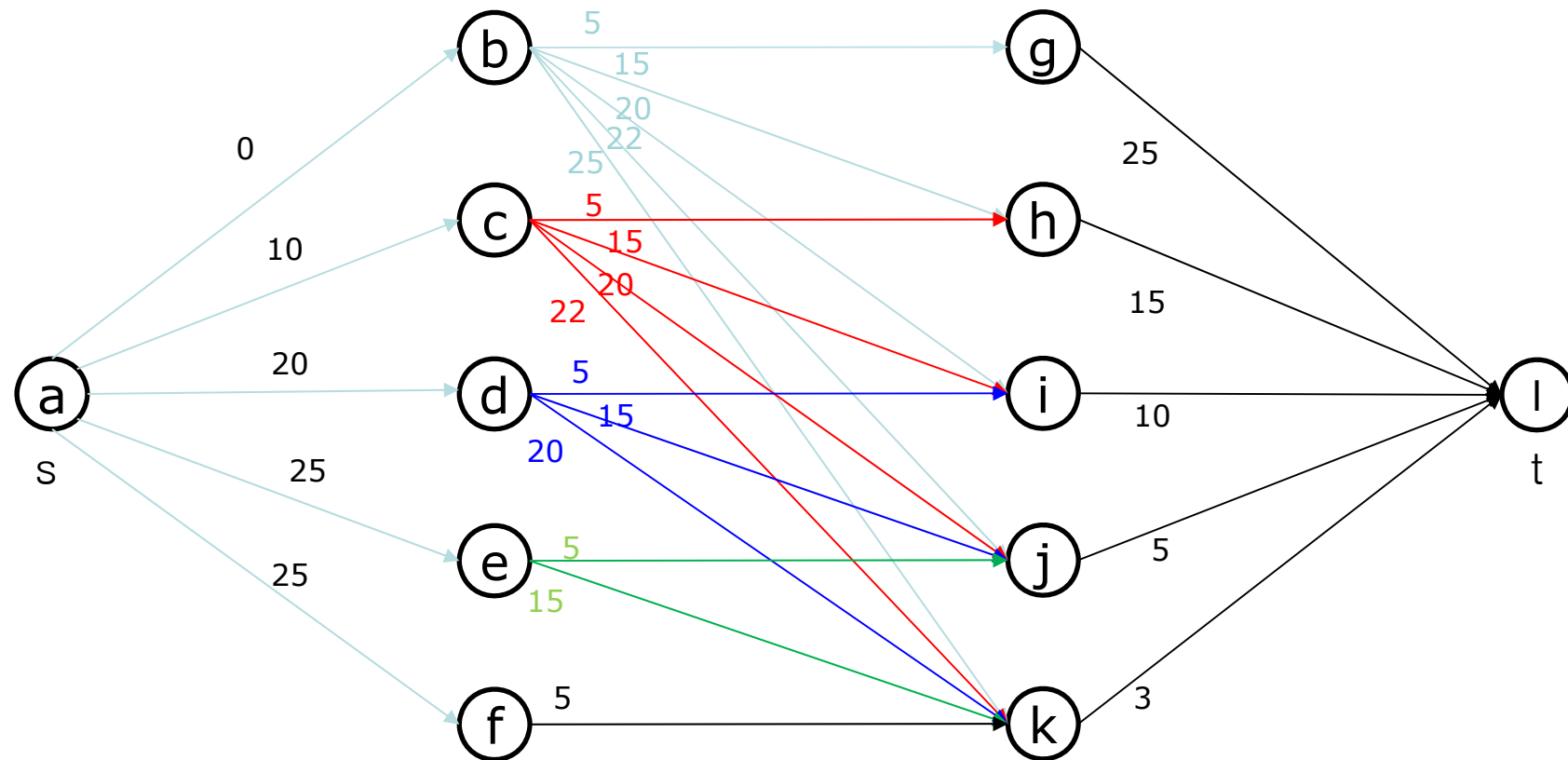
- Application: Resource allocation problem
  - Multistage graph of 4 units to 3 teams





## 6.3 Multistage graph

- Application: Resource allocation problem
  - Multistage graph of 4 units to 3 teams



## 6.3 Multistage graph

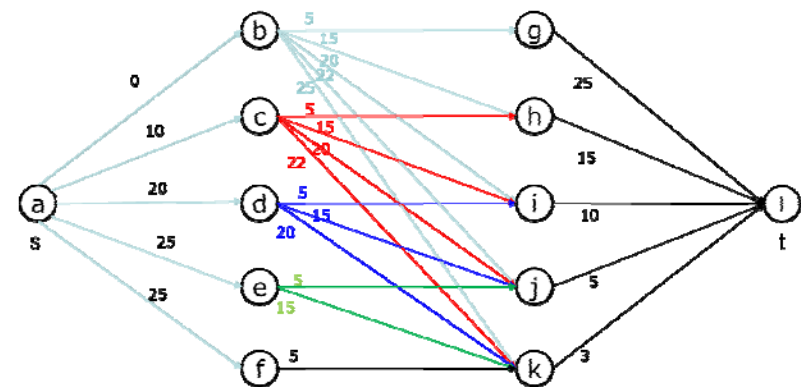
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- Application: Resource allocation problem
    - Multistage graph of 4 units to 3 teams
    - $\text{Cost}(1, a) = ?$
-

## 6.3 Multistage graph

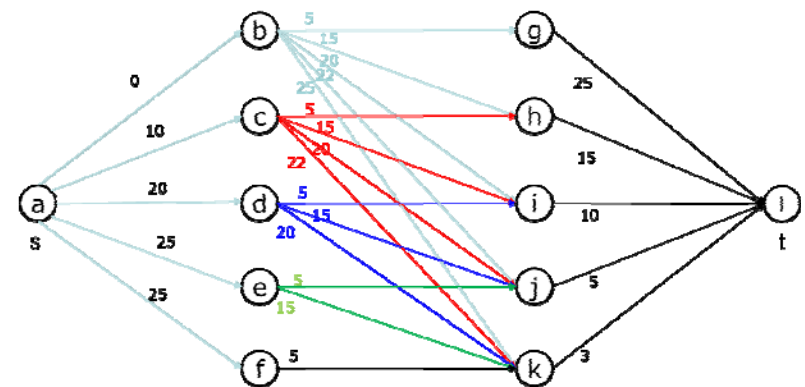
- Application: Resource allocation problem
  - Multistage graph of 4 units to 3 teams

- $\text{Cost}(3, g) =$
- $\text{Cost}(3, h) =$
- $\text{Cost}(3, i) =$
- $\text{Cost}(3, j) =$
- $\text{Cost}(3, k) =$



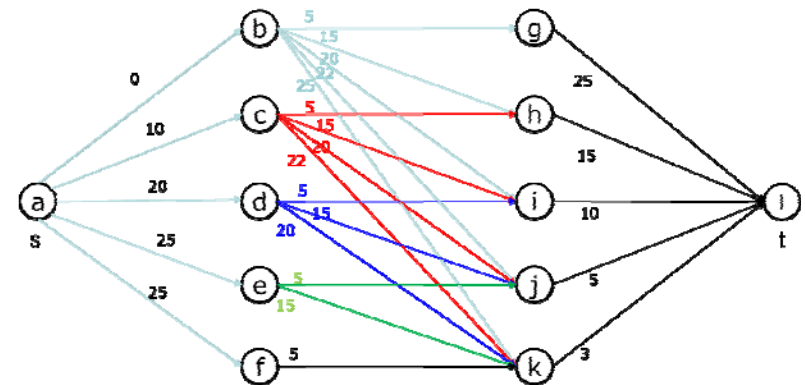
## 6.3 Multistage graph

- Application: Resource allocation problem
  - Multistage graph of 4 units to 3 teams
  - $\text{Cost}(2, b) =$
  - $\text{Cost}(2, c) =$
  - $\text{Cost}(2, d) =$
  - $\text{Cost}(2, e) =$
  - $\text{Cost}(2, f) =$



## 6.3 Multistage graph

- Application: Resource allocation problem
  - Multistage graph of 4 units to 3 teams
  - Cost (1, a) =



# All about Dynamic Programming

Type	Stepwise approach	Recursive structure	Approach	Time	Specialty
0/1 KNAP	Choosing objects	$\text{KNAP}(1, n, M) = \max \{ \text{KNAP}(2, n, M), \text{KNAP}(2, n, M - w_1) + p_1 \}$	Forward	$O(2^n)$	Smart Degenerate Case
Weighted Interval Scheduling	Selecting intervals	$\text{OPT}(j) = \max \{ \text{OPT}(p(j)) + v_j, \text{OPT}(j-1) \}$	Backward	$O(n)$	Auxiliary memory
Multistage Graph	Determining vertex in the next stage	$\text{COST}(i, j) = \min_{l \in V_{i+1}} \{ c(j, l) + \text{COST}(i+1, l) \}$	Backward	$O(m)$	Resource allocation problem
All Pairs Shortest Path					

## 6.3 Multistage graph

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- 다음 중 옳지 않은 문장을 모두 고르시오.
  - (a) multistage graph는 acyclic graph이다.
  - (b) multistage graph에서 계층의 수는 vertex의 수에 비례한다.
  - (c) multistage graph 문제는 Dijkstra 알고리즘으로 해결할 수 있다.
  - (d) multistage graph 문제는 backward approach이다.
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