

Amdahl's Law

The necessity to parallelize

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- Algorithms can easily exceed any reasonable amount of runtime.
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- Problems
 - There are task that can be solved only sequentially (eg. reading from disk).
 - There are task which requires a lot of effort and work to parallelize.
- We expect - or at least hope - that parallelization shortens runtime significantly.



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- Amdahl's law now can be formulated as the following:

$$Q_{\text{speed up}}(N) = \frac{1}{S + \frac{P}{N}} = \frac{1}{(1 - P) + \frac{P}{N}}, \quad (2)$$

Where N is the number of parallel threads, and the runtime S was expressed in the denominator using equation (1).



Sample outputs of method I.

