

What are algorithms?

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Formally, given the specification of a problem

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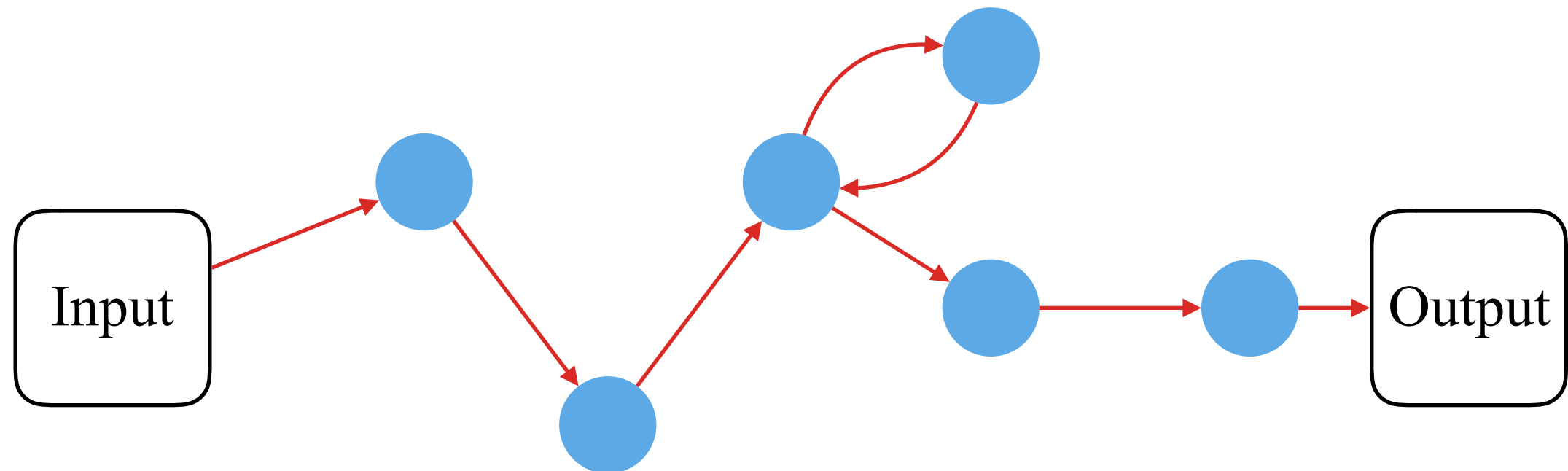
Input

A diagram illustrating the flow of an algorithm. It consists of two rounded rectangular boxes. The first box on the left is labeled 'Input'. A horizontal arrow points from the right side of the 'Input' box to the left side of the 'Output' box. The second box on the right is labeled 'Output'.

Output

# What is an algorithm?

Formally, given the specification of a problem



an algorithm is computational procedures that take some values as **input** and produce some values as **output**.

# What is an algorithm?

Here is an informal example:

Input



Output



# What is an algorithm?

Here is an informal example:

Algorithm 1

foreach egg {



}

Input



Output



# What is an algorithm?

Here is an informal example:

Algorithm 1

foreach egg {



}

Algorithm 2

foreach egg {



}

Input



Output



# Selection Sort



# The Champion Problem

Input: an array  $A$  of  $n$  integers.

Output: an index  $k$  so that  $A[k]$  is the minimum value in  $A$ .

A problem instance (an instance)



return value (ret):



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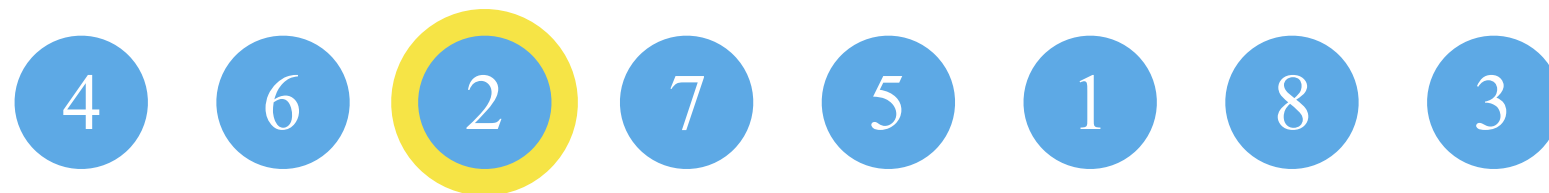


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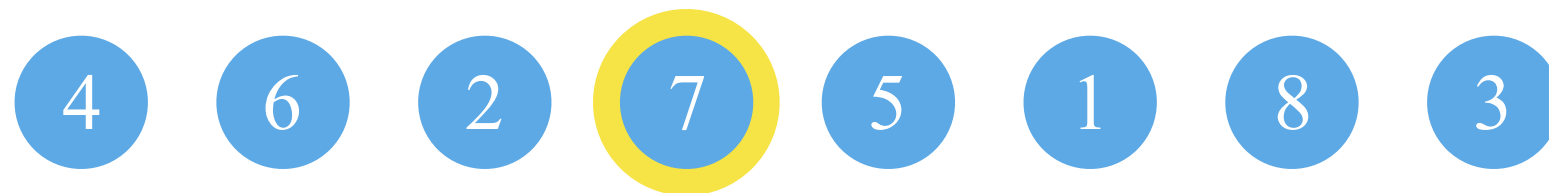


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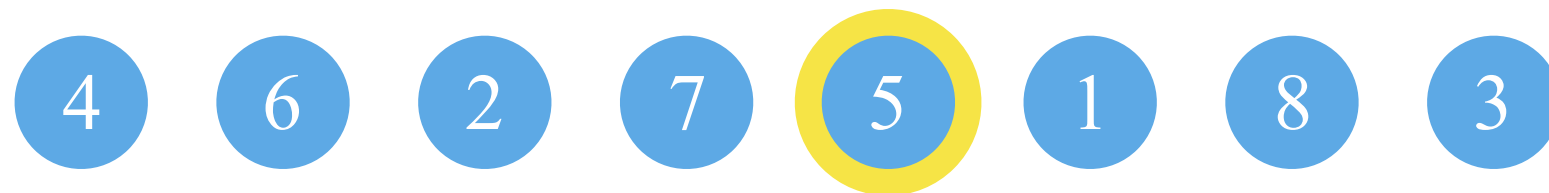


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return value (ret):



# C++ Code

```
int champion(int *s, int n){ // return -1 for empty input

    int ret = 0; // 1 assignment

    for(int i=0; i<n; ++i){ // incur 2n comparisons, ≤ n-1 assignments,
        if(s[i] < s[ret]){ // and n increments
            ret = i;
        }
    }
    return ret;
} // a constant number of operations for the overhead of function call
```

--- total running time ---

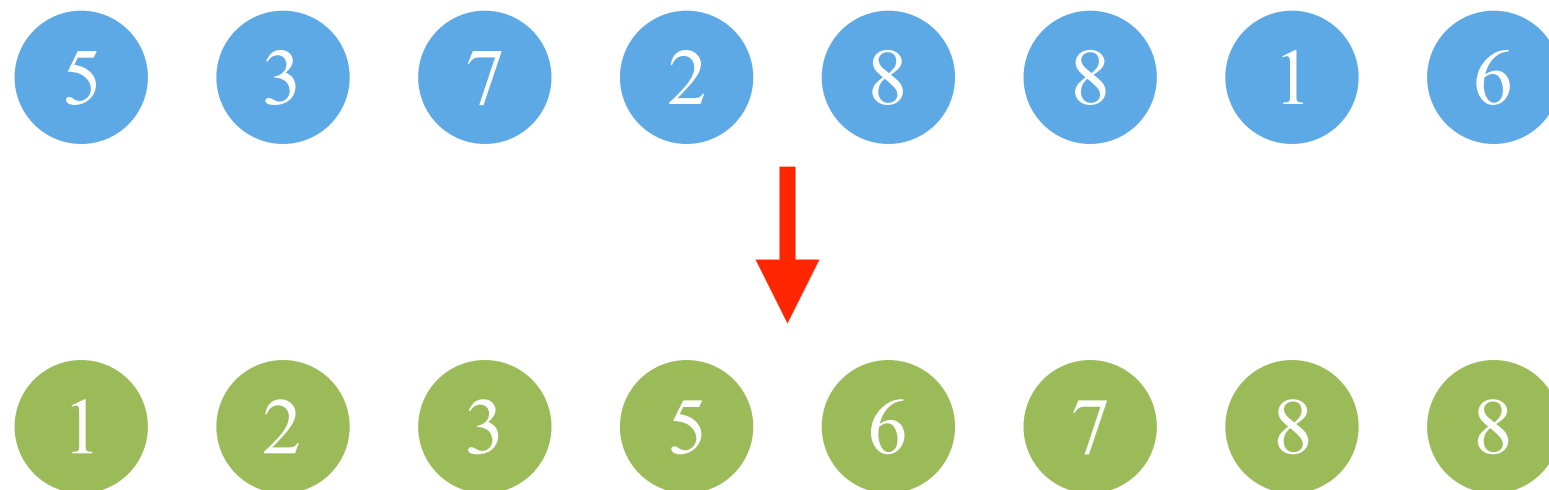
champion() uses at most  $4n + C$  operations for some constant C.

# Sorting Problem

Input: an array  $A$  of  $n$  integers.

Output: the same array with the  $n$  integers ordered nondecrementally.

An instance



# Selection Sort

Input: an array  $A$  of  $n$  integers.

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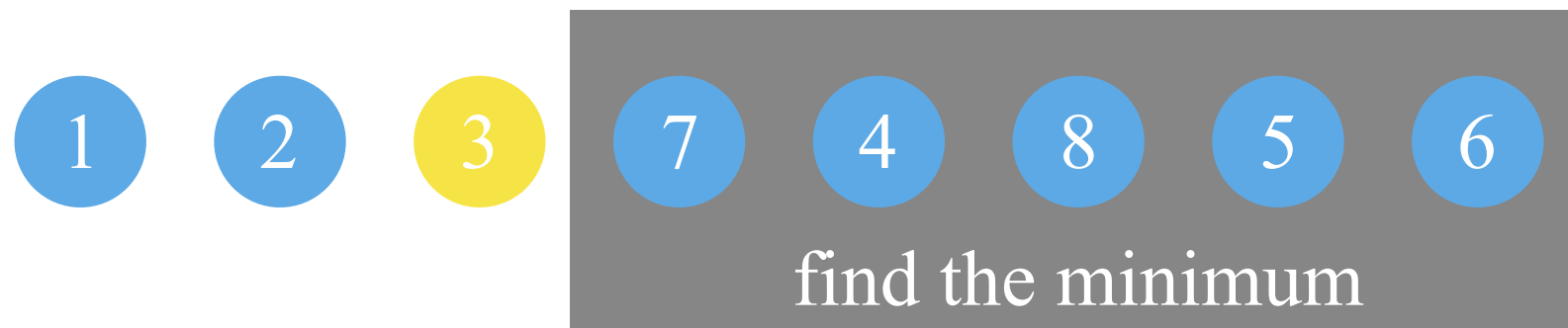


# Selection Sort

Input: an array  $A$  of  $n$  integers.

Output: the same array with the  $n$  integers ordered nondecrementally.

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# C++ Code

```
void selection_sort(int *s, int n){  
    for(int i=0; i<n; ++i){  
        int k = champion(s+i, n-i);  
        int swap = s[i]; s[i] = s[k]; s[k] = swap;  
    }  
}
```

--- about the highlight ---

It is called *reduction*. Reducing one problem X to another problem Y means to devise an algorithm for X using an algorithm for Y as a building block.

selection\_sort() uses at most  $n(4n+C+3)$  operations for some constant C.

# C++ Code

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void selection_sort(int *s, int n){  
    for(int i=0; i<n; ++i){  
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Why does the count of operations matter?

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selection\_sort() uses at most  $n(4n+C+3)$  operations for some constant  $C$ .

Why does the count of operations matter?

A: We can use it to estimate the running time of the program.  $10^8$  operations takes roughly 1 second. Hence, sorting  $10^4$  integers by selection sort takes roughly 4 seconds.

# Insertion Sort

# Insert $x$ into a sorted array $A$ while retaining $A$ sorted

Input: a sorted array  $A$  of  $n$  integers and an integer  $x$ .

Output: a sorted array that comprises all elements in  $A$  and  $x$ .

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# C++ Code

```
void insert(int *s, int n, int x){ // array s has length  $\geq n+1$ 
0:  bool placed = false; // whether x has been placed in s
1:  for(int i=n-1; i>=0 && !placed; --i){
2:      if(s[i] > x){
3:          s[i+1] = s[i];
4:      }else{
5:          s[i+1] = x; placed = true;
6:      }
7:  }
8:  if(!placed) s[0] = x;
9:}
```

-----

Line 1 comprises 1 assignment,  $n$  comparisons, and  $n$  decrements.

Line 2 comprises  $n$  comparisons and  $n$  dereference.

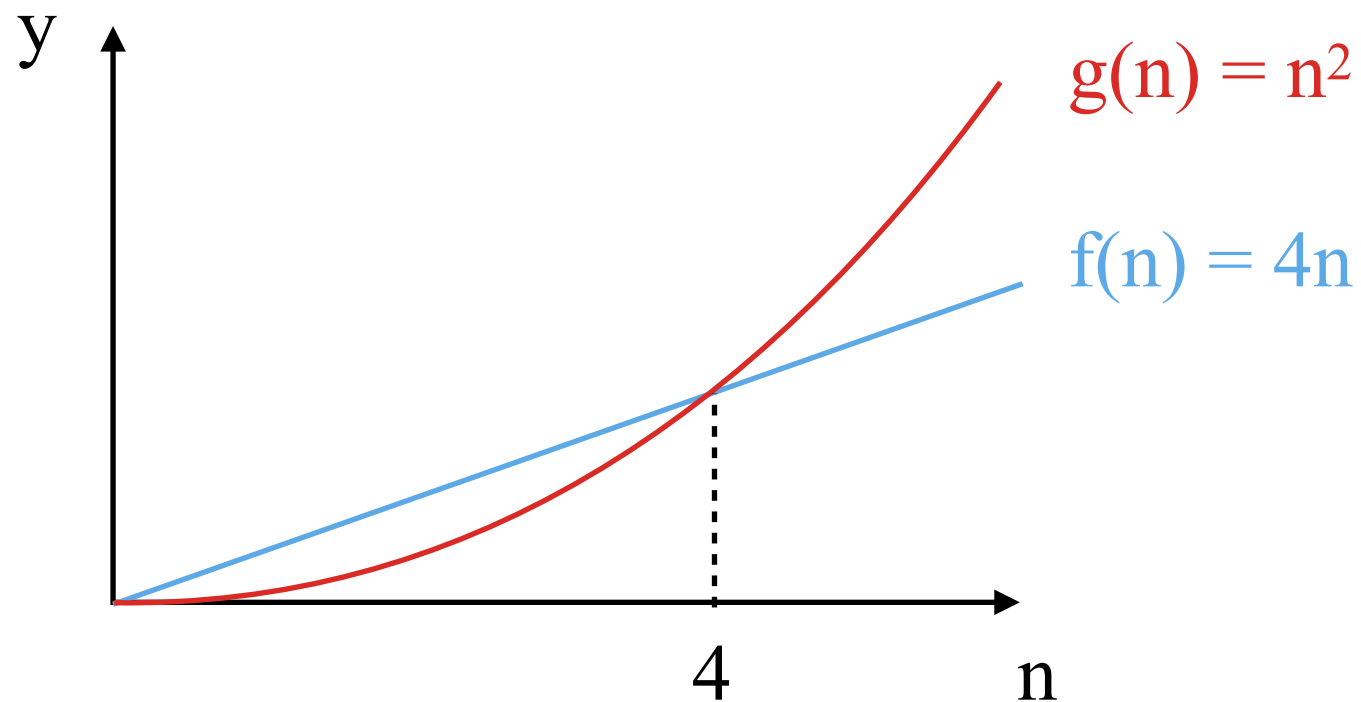
Line 3 comprises  $n$  assignments,  $n$  additions, and  $2n$  dereferences ...

It is cumbersome (and error-prone) to count the exact operations that an algorithm uses.

# Asymptotic Notation: O-Notation

$O(g(n))$  is pronounced as big-Oh of  $g$  of  $n$ .

$f(n) = O(g(n))$  means that  $f(n) \leq C \cdot g(n)$  for every  $n \geq n_0$  for some constants  $C$  and  $n_0$ .



-----

We can write  $4n = O(n)$  by setting  $(C, n_0) = (4, 1)$  or  $4n = O(n^2)$  by setting  $(C, n_0) = (1, 4)$ .

# Asymptotic Notation: O-Notation

$O(g(n))$  is pronounced as big-Oh of  $g$  of  $n$ .

More formally,  $f(n) = O(g(n))$  means that  $f(n)$  is a function contained in the set of functions  $\{h(n) : \text{there exists positive constants } n_0 \text{ and } C \text{ so that } C \cdot g(n) \geq h(n) \text{ for every } n \geq n_0\}$ .

Because it is cumbersome to determine the constant  $C$  and we simply need to estimate the running time, we usually use asymptotic notation to denote the time complexity of an algorithm.

--- Example ---

1. Selection sort runs in  $O(n^2)$  time, so it can sort  $10^4$  integers in seconds.
2. Insert  $x$  into a sorted array runs in  $O(n)$ , so in seconds one can complete  $10^4$  insertions.

# Insertion Sort

Input: an array  $A$  of  $n$  integers.

Output: the same array with the  $n$  integers ordered nondecrementally.

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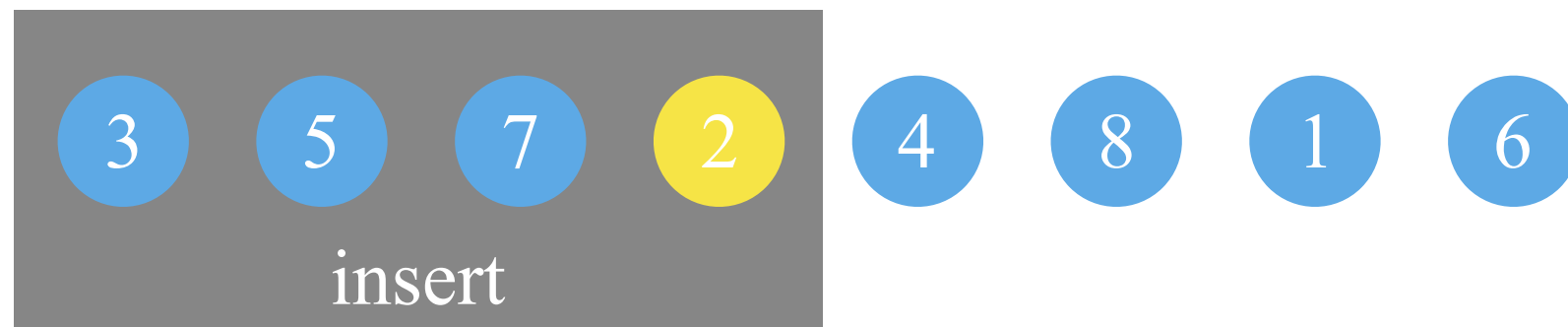


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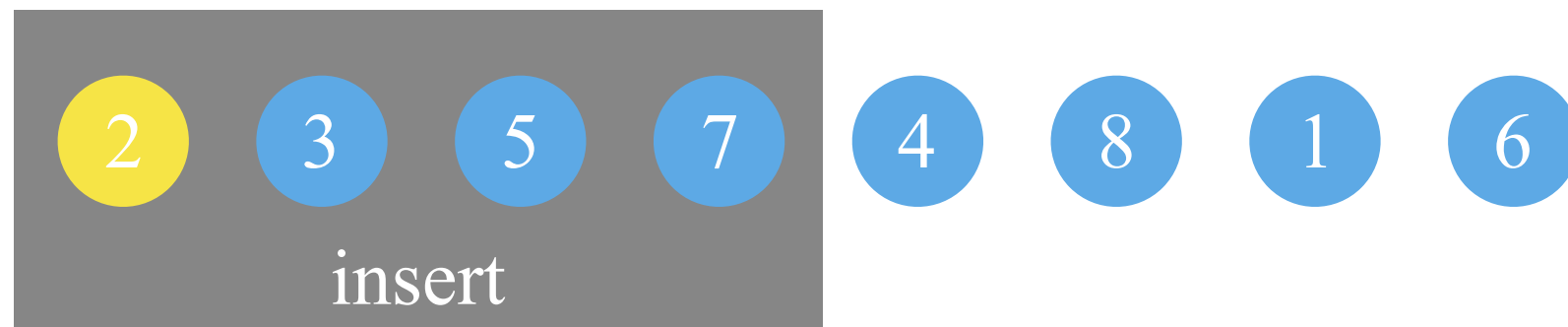


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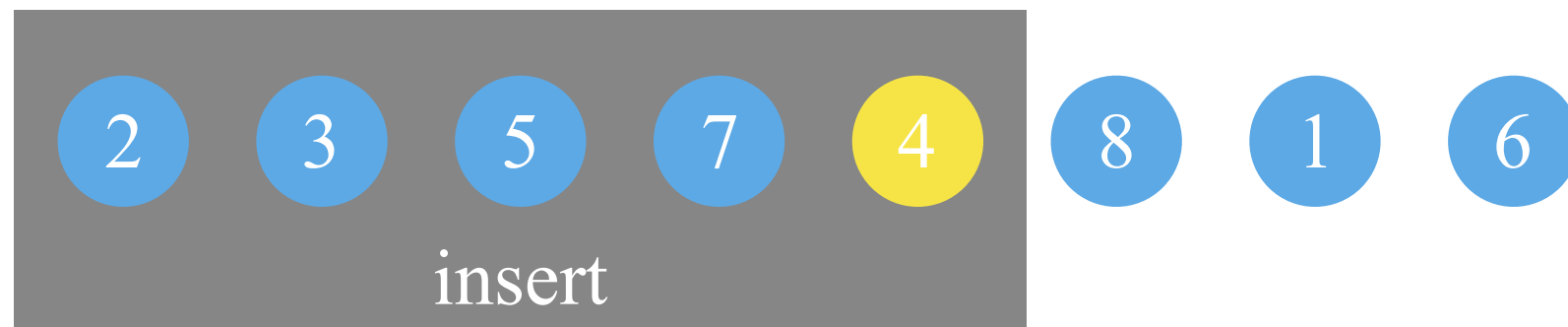


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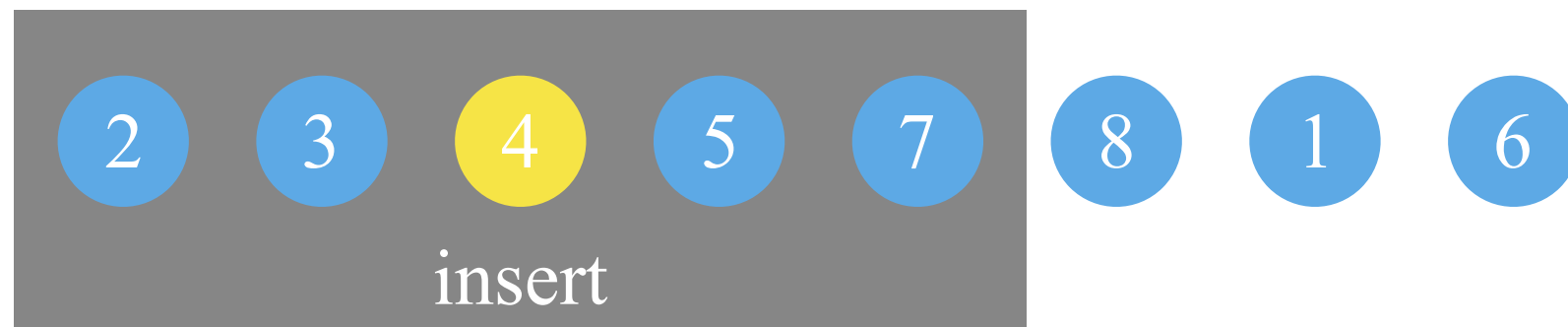


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# C++ Code

```
void insertion_sort(int *s, int n){  
    for(int i=1; i<n; ++i){  
        insert(s, i, s[i]);  
    }  
}
```

--- about the highlight ---

Again, we use a reduction here.

The running time is  $O(n) \cdot O(n) = O(n^2)$ . Why does this equality hold?

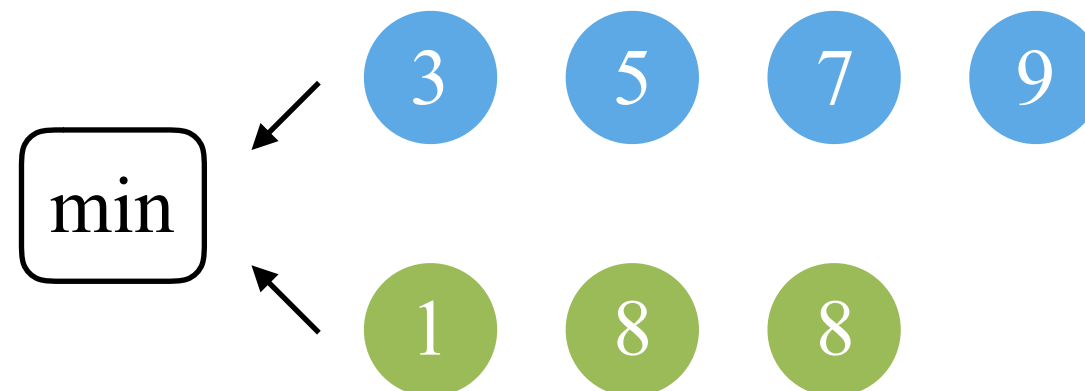
# Merge Sort

# Merge two sorted arrays into one

Input: two sorted arrays A and B of integers.

Output: a sorted array that comprises all elements in A and B.

An instance

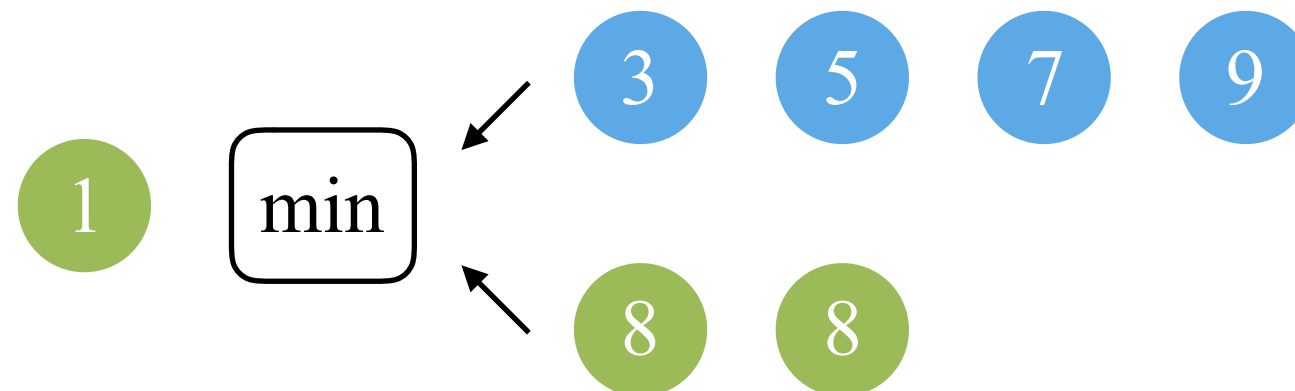


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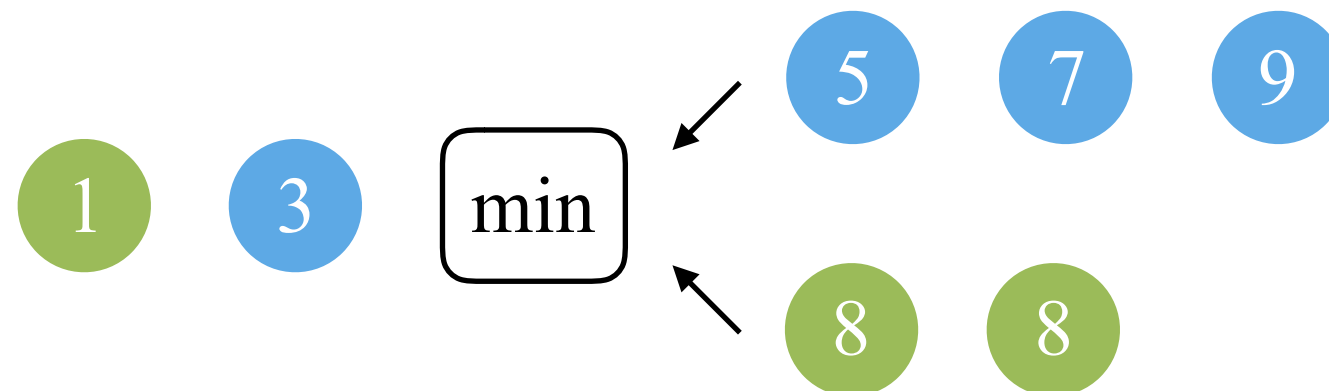


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# C++ Code

```
int* merge(int *s, int n, int *r, int m){  
  
    int *ret = new int [n+m];  
    int i = 0, j = 0, k = 0;  
  
    while( i < n || j < m ){  
        if(i < n && j < m){ // when both arrays are not empty  
            ret[k++] = ((s[i] < r[j]) ? s[i++] : r[j++]);  
        }else{  
            ret[k++] = ((i < n) ? s[i++] : r[j++]);  
        }  
    }  
}  
-----
```

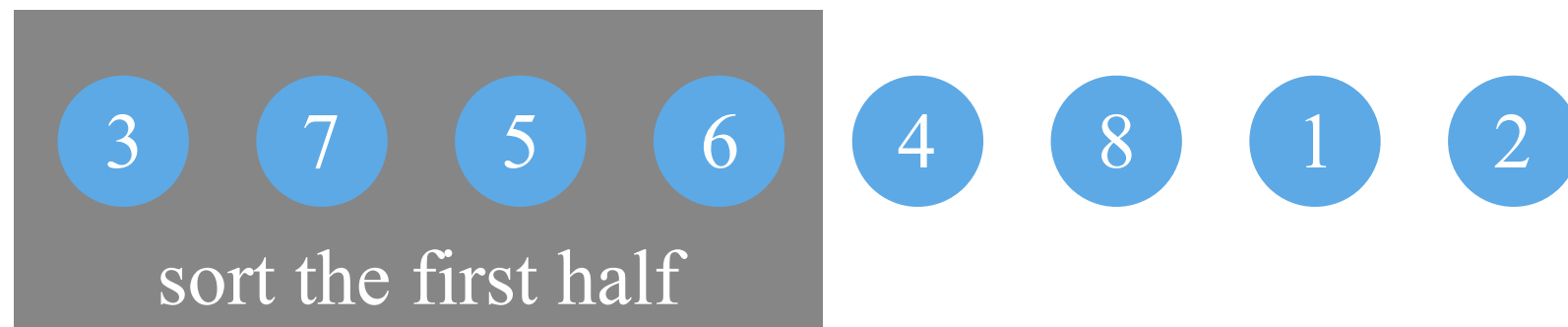
Merging two sorted arrays takes  $O(n+m)$  time.

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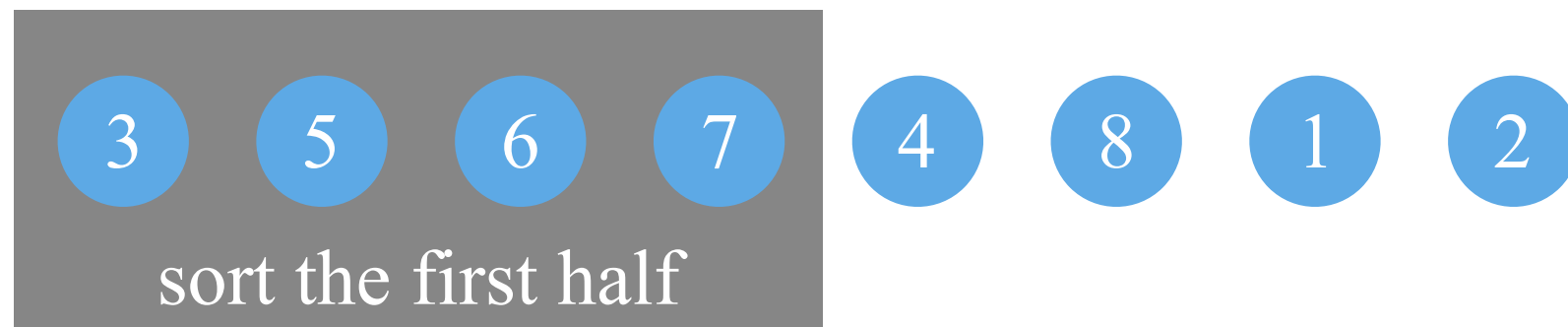


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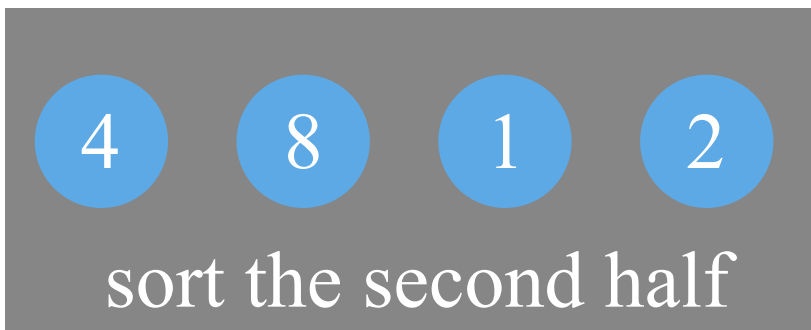


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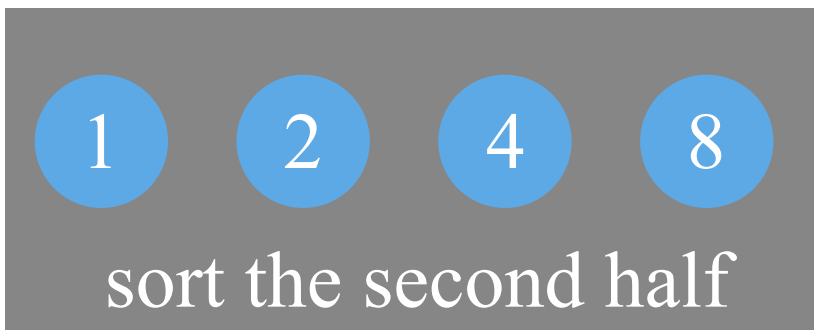


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# C++ Code

```
void merge_sort(int *s, int n){  
  
    if(n == 1) return;  
  
    int k = n/2;  
    merge_sort(s, k);  
    merge_sort(s+k, n-k);  
  
    int *r = merge(s, k, s+k, n-k);  
    memcpy(s, r, sizeof(int)*n);  
}
```

--- about the highlight ---

A reduction from a problem to itself is called *recursion*. A recursion **usually** requires that the instance size decreases monotonically. Why?



# C++ Code

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    if(n == 1) return;  
  
    int k = n/2;  
    merge_sort(s, k);  
    merge_sort(s+k, n-k);  
  
    int *r = merge(s, k, s+k, n-k);  
    memcpy(s, r, sizeof(int)*n);  
}
```

-----

Merge sort needs at most  $T(n) = T(\lfloor n/2 \rfloor) + T(\lceil n/2 \rceil) + c_1 n$  operations where  $T(1) = c_2$ .

# Recursion-Tree Method

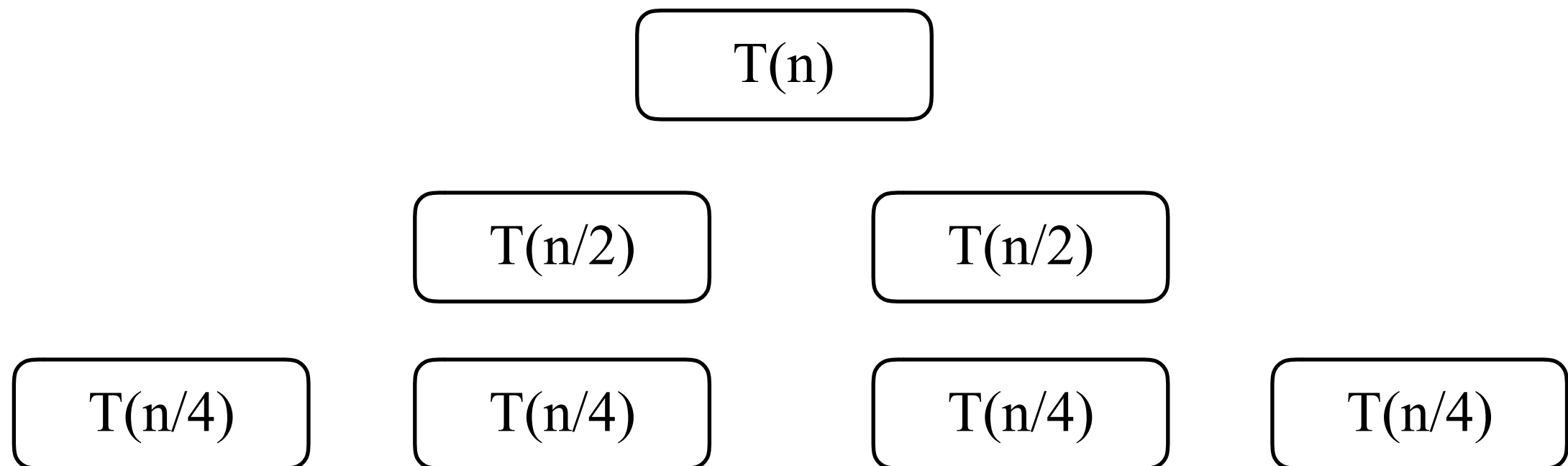
Merge sort needs at most  $T(n) = T(\lfloor n/2 \rfloor) + T(\lceil n/2 \rceil) + c_1 n$  operations where  $T(1) = c_2$ . We have seen how to verify the guess  $T(n) = O(n \log n)$ .  
**How to come up with a guess?**

We simply need a guess, so we may drop the floor and the ceiling functions, and ignore the constants. We get:

$$T(n) = \begin{cases} 2T(n/2) + n & \text{if } n > 1 \\ 1 & \text{otherwise} \end{cases}$$

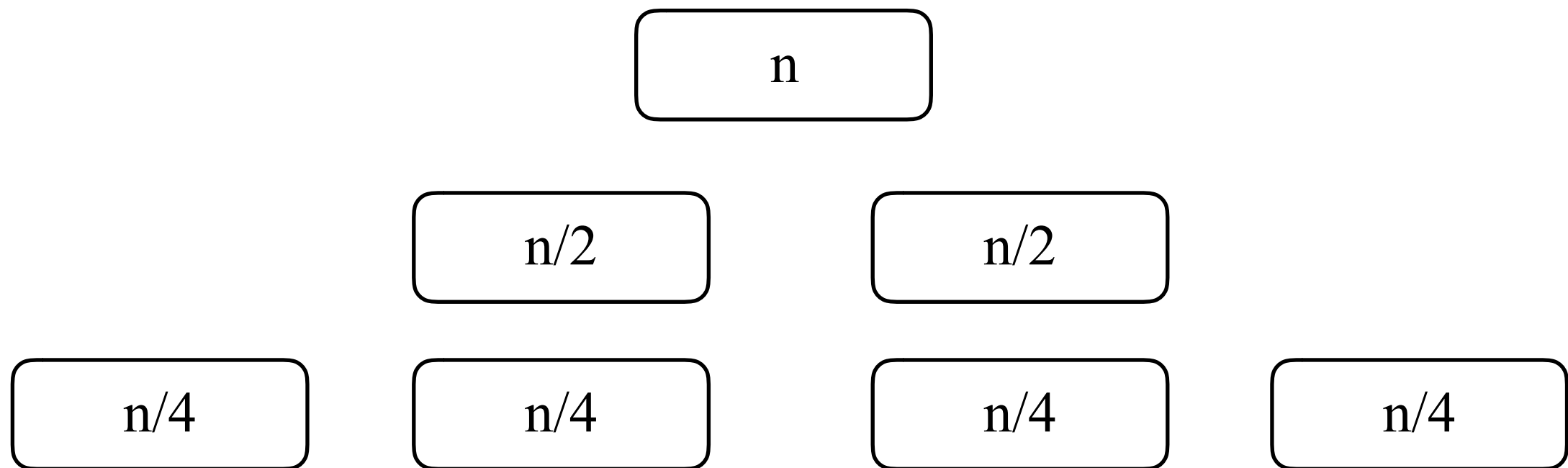
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There are  $\log_2 n$  layers, and for each layer the sum of cost is  $n$ . Consequently, the total cost is  $O(n \log n)$ .