

Computer Security Capstone

Chapter 2: Cryptographic Tools

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Outline

- Confidentiality with Symmetric Encryption
- Message Authentication and Hash Function
- Public-Key Encryption
- Digital Signatures and Key Management
- Random and Pseudorandom Numbers

Confidentiality with Symmetric Encryption

- Symmetric encryption

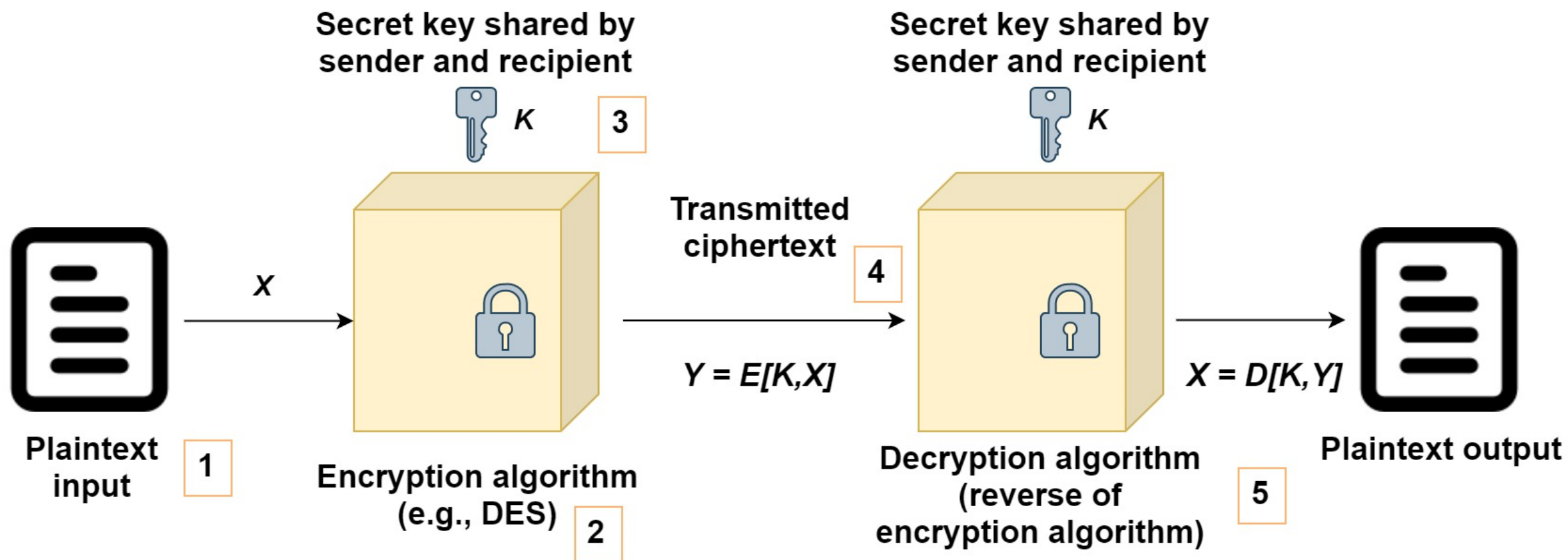
- Providing confidentiality for transmitted or stored data
- Conventional encryption or single-key encryption

- Two requirements for its secure use

- A strong encryption algorithm
 - Opponent: Unable to decrypt ciphertext or discover the key, (given pairs of ciphertexts and plaintexts, as well as the algorithm)
- Secure key distribution and maintenance

Simplified Model of Symmetric Encryption

- Five major components



Attacking Symmetric Encryption

Cryptanalytic Attacks

- Exploit
 - ▣ Nature of the algorithm
 - ▣ General characteristics of the plaintext
 - ▣ Sample plaintext-ciphertext pairs
- Deduce a specific plaintext or the key

Brute-Force Attacks

- Exploit
 - ▣ Knowledge about the expected plaintext
- Try all possible keys on some ciphertexts
 - ▣ Until an intelligible translation into plaintext is obtained
 - ▣ On average half of all possible keys must be tried to achieve success

Data Encryption Standard (DES)

- Adopted in 1977 by the NIST (FIPS PUB 46)

- Most widely used encryption scheme: aka Data Encryption Algorithm (DEA)
- 64-bit plaintext blocks and a 56-bit key → 64-bit ciphertext blocks

- Security concerns

- Algorithm: characteristics may be exploited?
 - Most-studied encryption algo: numerous attempts to find weakness, but no fatal one yet
- Key length: too short
 - 56 bits → $2^{56} = 7.2 \times 10^{16}$ possible keys (Inadequate for today's processor speed)

FIPS (Federal Information Processing Standard): describes document processing, encryption algorithms and other technology standards for use within non-military government agencies and by government contractors and vendors who work with the agencies.

Brute-Force Attacks against DES

- On average, half the key space has to be searched
 - One DES encryption per micro second → more than 1000 years (3.6×10^{16} keys)
- In July 1998, EFF broke a DES encryption
 - DES cracker: less than \$250,000, less than three days
 - <http://cs-exhibitions.uni-klu.ac.at/index.php?id=263>
- Encryption speeds advance
 - Seagate Technology [SEAG08]
 - Multicore computers (2008): 10^9 per second
 - EE Times [AROR12]
 - Contemporary supercomputer (2012): 10^{13} per second → break DES within 1 hour



DES cracker
circuit board:
key discovery
in 56 hours

EFF (Electronic Frontier Foundation): the leading nonprofit organization defending civil liberties in the digital world

Triple DES (3DES)

- Part of the DES in 1999: FIPS PUB 46-3

- Repeat basic DES algorithm 3 times using either 2 or 3 unique keys
- A key size of 112 or 168 bits

- Two attractions

- 168-bit key length: overcomes brute-force attack of DES
- Underlying encryption algorithm is the same as in DES

- Two drawbacks

- Sluggish algorithm/software: not efficient software code and three times as DES
- Uses a 64-bit block size: not efficient and not secure

Advanced Encryption Standard (AES)

- AES: now widely available in commercial products
 - ❑ A replacement for 3DES
 - ❑ 3DES was not reasonable for long term use
- NIST called for proposals for a new AES in 1997
 - ❑ Security strength: equal to or better than 3DES
 - ❑ Significantly improved efficiency
 - ❑ Symmetric block cipher
 - ❑ 128-bit data and 128/192/256 bit keys
- Selected Rijndael Algorithm in Nov. 2001: published as FIPS 197

Symmetric Block Encryption Algorithms

- Block ciphers: most commonly used symmetric encryption
 - ▣ Fixed-size blocks of plaintext → blocks of ciphertext of equal size

	DES	Triple DES	AES
Plaintext block size (bits)	64	64	128
Ciphertext block size (bits)	64	64	128
Key size (bits)	56	112/168	128/192/256

DES = Data Encryption Standard

AES = Advanced Encryption Standard

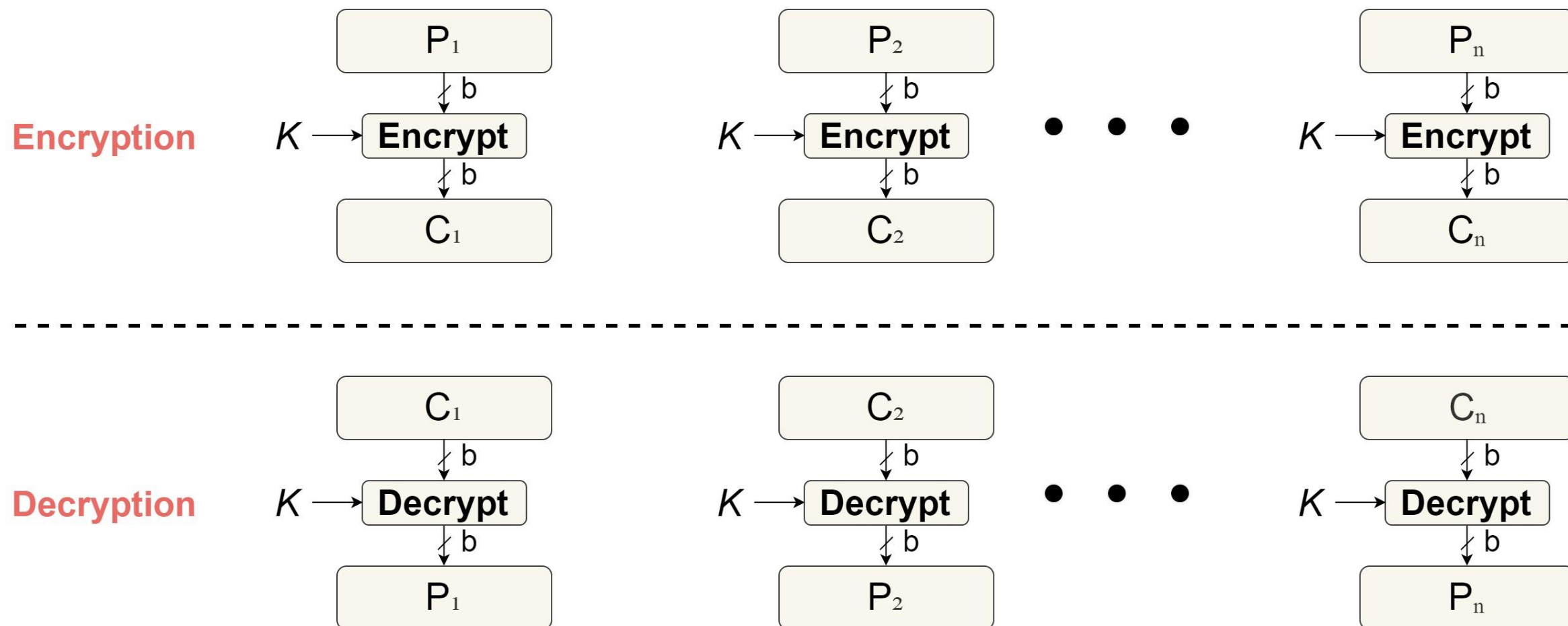
Average Time Required for Exhaustive Key Search

Key size (bits)	Cipher	Number of Alternative Keys	Time Required at 10^9 decryptions/s	Time Required at 10^{13} decryptions/s
56	DES	$2^{56} \approx 7.2 \times 10^{16}$	2^{55} ns = 1.125 years	1 hour
128	AES	$2^{128} \approx 3.4 \times 10^{38}$	2^{127} ns = 5.3×10^{21} years	5.3×10^{17} years
168	Triple DES	$2^{168} \approx 3.7 \times 10^{50}$	2^{167} ns = 5.8×10^{33} years	5.8×10^{29} years
192	AES	$2^{192} \approx 6.3 \times 10^{57}$	2^{191} ns = 9.8×10^{40} years	9.8×10^{36} years
256	AES	$2^{256} \approx 1.2 \times 10^{77}$	2^{255} ns = 1.8×10^{60} years	1.8×10^{56} years

Practical Security Issues

- How to apply the symmetric encryption to a unit of data larger than a single 64-bit or 128-bit block?
 - e.g., E-mail messages, network packets, and database records
- Simplest approach: electronic codebook (ECB)
 - Multiple block encryption
 - Each block of plaintext is encrypted using the same key
- Issue: Cryptanalysts may exploit regularities in the plaintext

Illustration of the ECB Mode

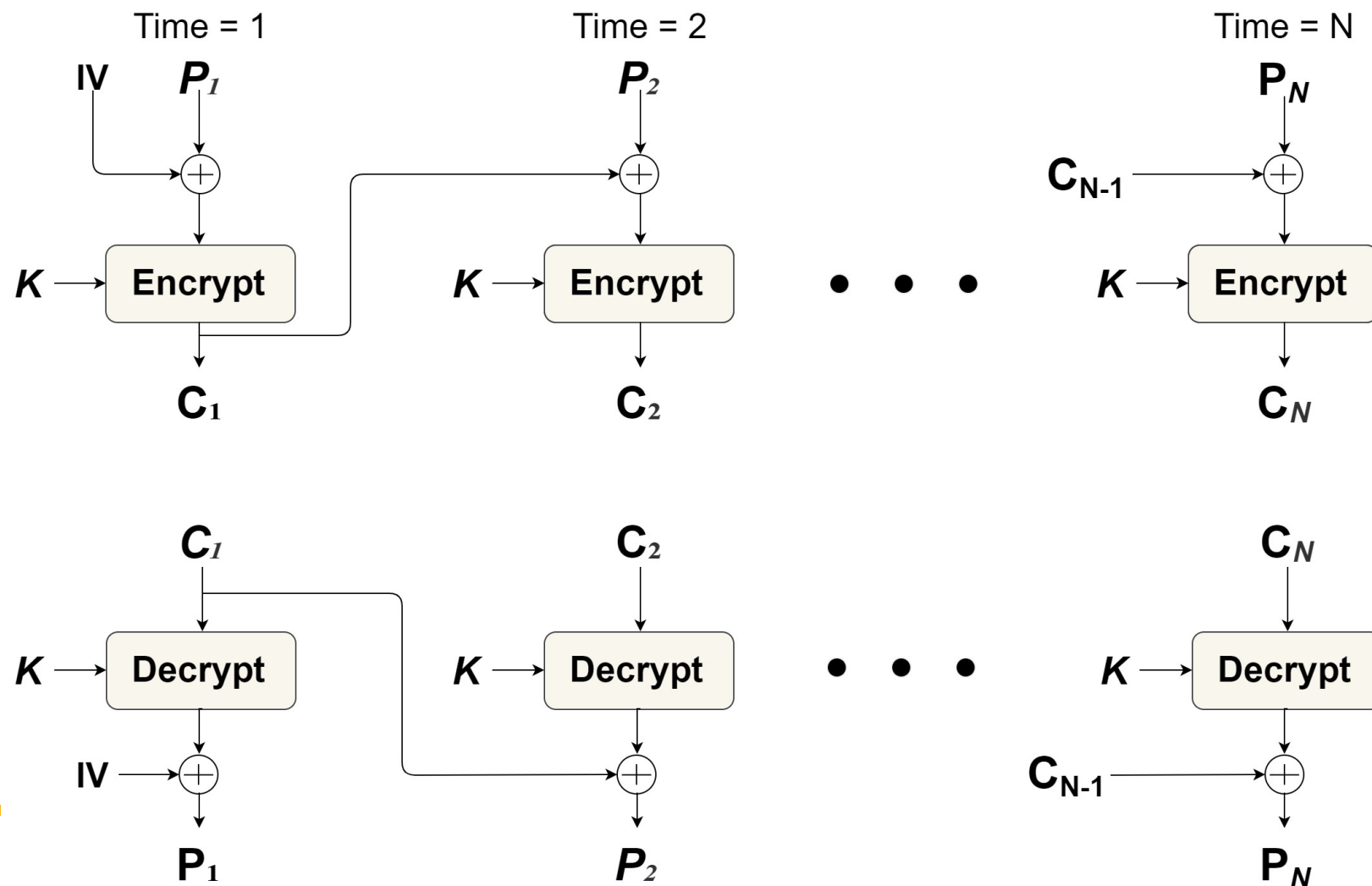


How to address the ECB issue (regularities)?

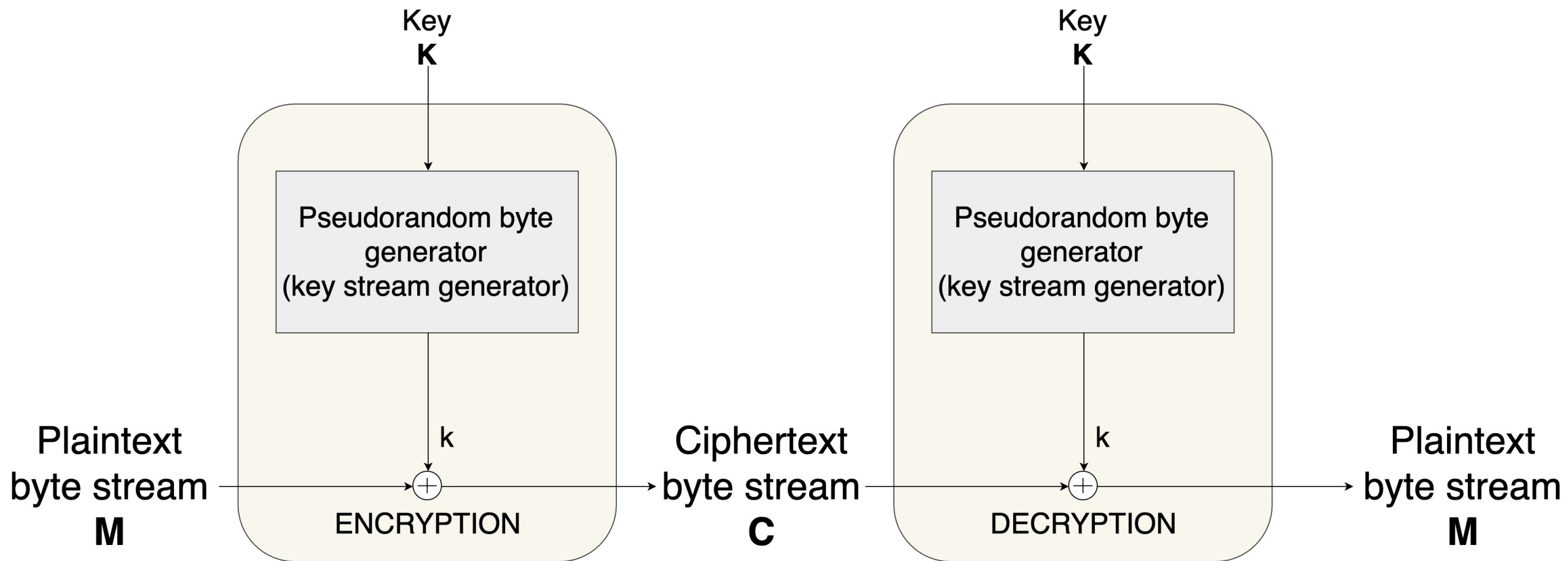
Modes of Operation

- Five modes of operation defined by NIST
 - ▣ ECB, CBC, etc.
- CBC can overcome the weakness of ECB

Cipher Block Chaining (CBC) Mode



The Other Cipher Method: Stream Cipher



Block & Stream Ciphers

Block Cipher

- Processing one block at a time
- Each input block → an output block
- Pro: Can reuse keys
 - ▣ More common
- Apps: file transfer, e-mail, and database

Stream Cipher

- Processing input elements continuously
- One element at a time
 - ▣ Typically: one byte; one bit or larger units are also allowed
- Pro: almost always faster (XOR) and use far less code
- Apps: data stream over a communication channel or a browser/Web link

The keystream must not be reused!!

Consider: $(A \oplus K) \oplus (B \oplus K) = A \oplus B$

Message Authentication and Hash Functions

- Message (data) authentication

- Communicating parties can verify that received/stored messages are authentic
- Against falsification of data and transactions

- Two major aspects to verify

- Message contents: not altered
- Message source: authentic

- Another aspect

- Message timeliness/sequence: not artificially delayed or replayed

Can We Use Symmetric Encryption?

- Authentic source

- Only the sender and the receiver share the key

- No altered contents

- An error-detection code

- Proper message timeliness

- A sequence number or a timestamp

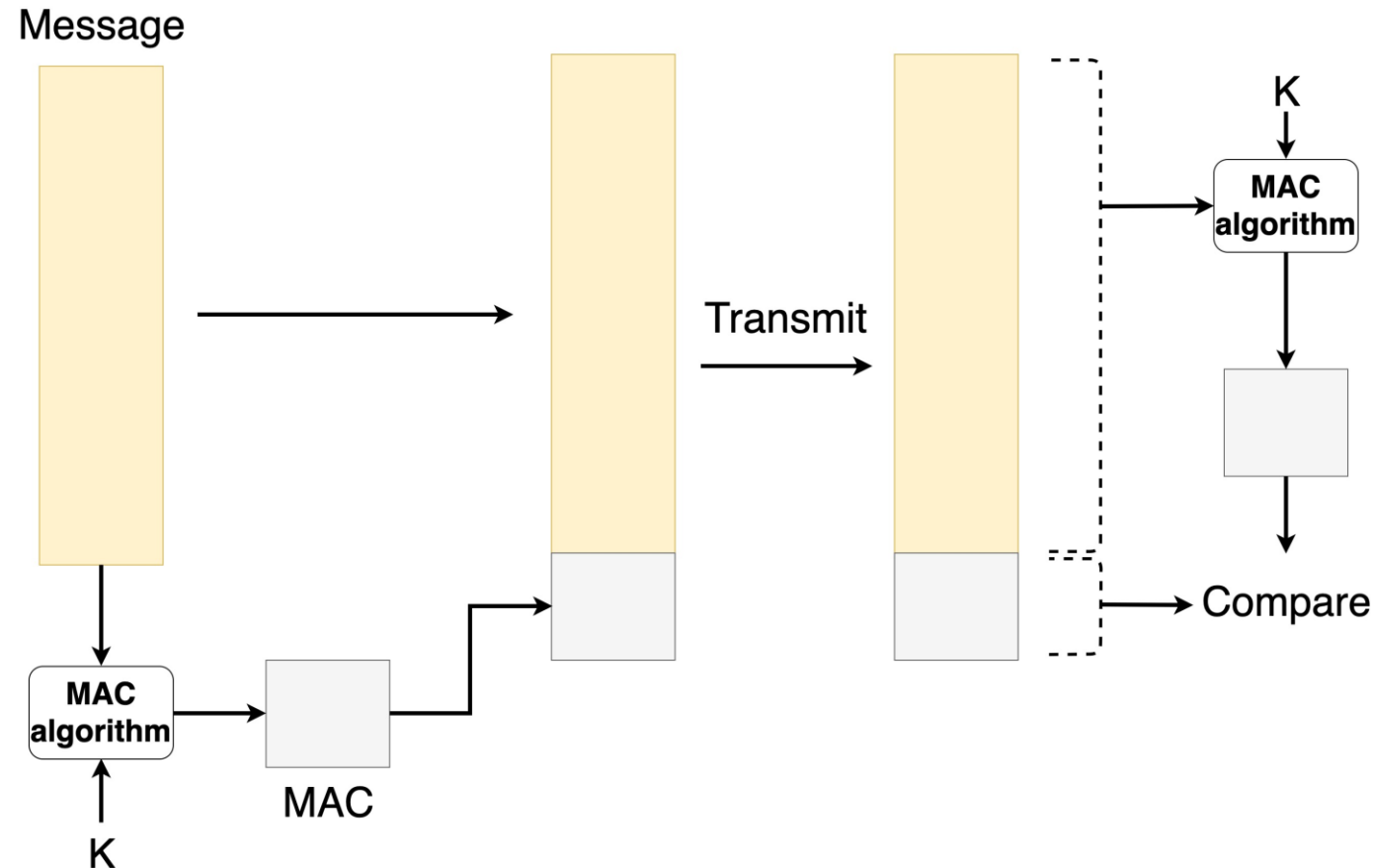
It seems proper, but it is not a suitable tool for data authentication. Why?

Message Authentication w/o Encryption

- Message authentication: a separate function from message encryption
 - e.g., a broadcast message to many destinations, not able to decrypt all incoming messages, wasteful of processor resources
- How? an auth. tag is generated and appended to each message
 - Message Authentication Code (MAC)
 - One-Way Hash Function

MAC

- Use a secret key to generate a small block of data
- Assumption: two communicating parties share the secret key
- MAC algorithms: NIST recommends DES
 - ❑ MAC: last 16- or 32-bit code of the encrypted message



What are the drawbacks?

MAC (Cont.)

- Drawbacks

- ❑ Encryption software is quite slow
- ❑ Encryption hardware costs are non-negligible
- ❑ Encryption hardware is optimized toward large data sizes

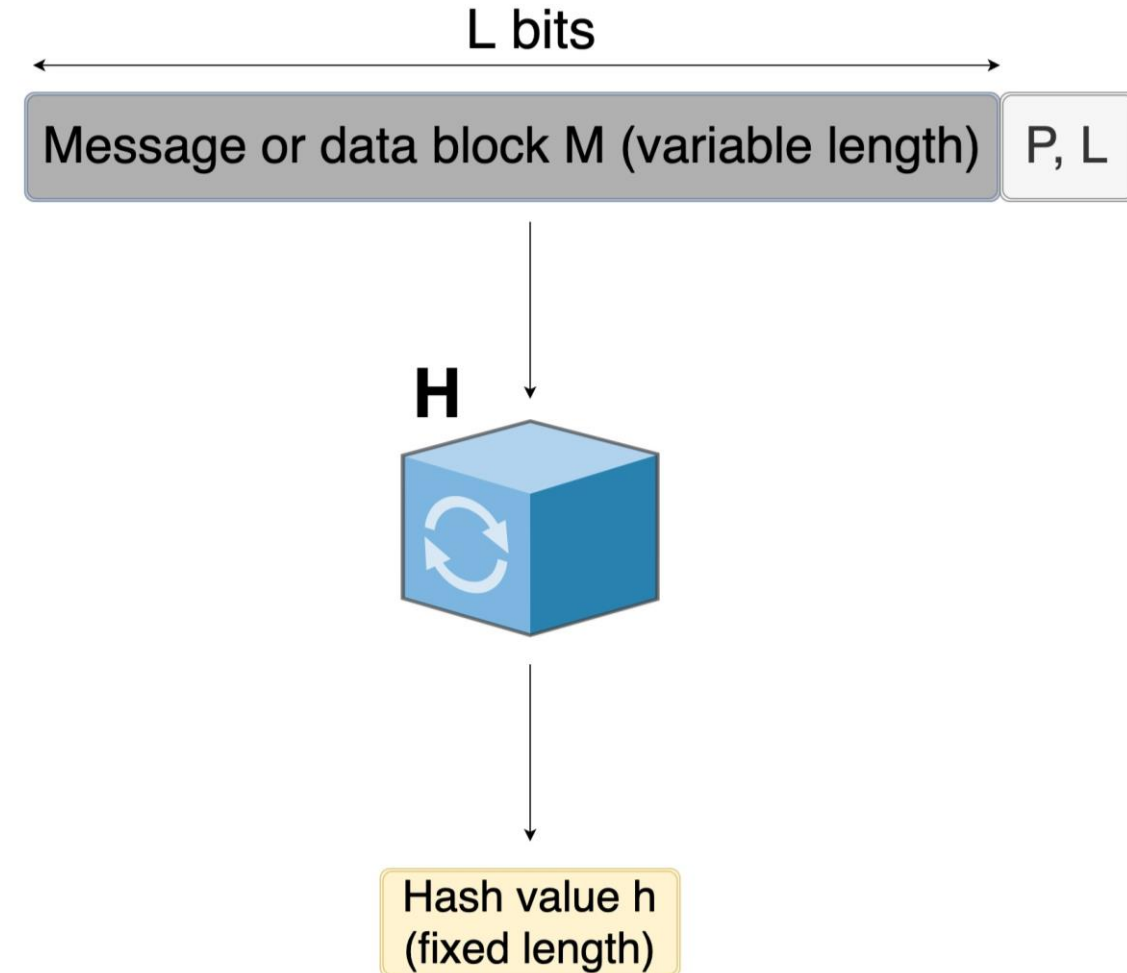
Does message authentication really need encryption of the message?

- No. Why?

- ❑ Authentication need not be reversible
- ❑ Only need a way to generate a tag which can verify messages

One-way Hash Function

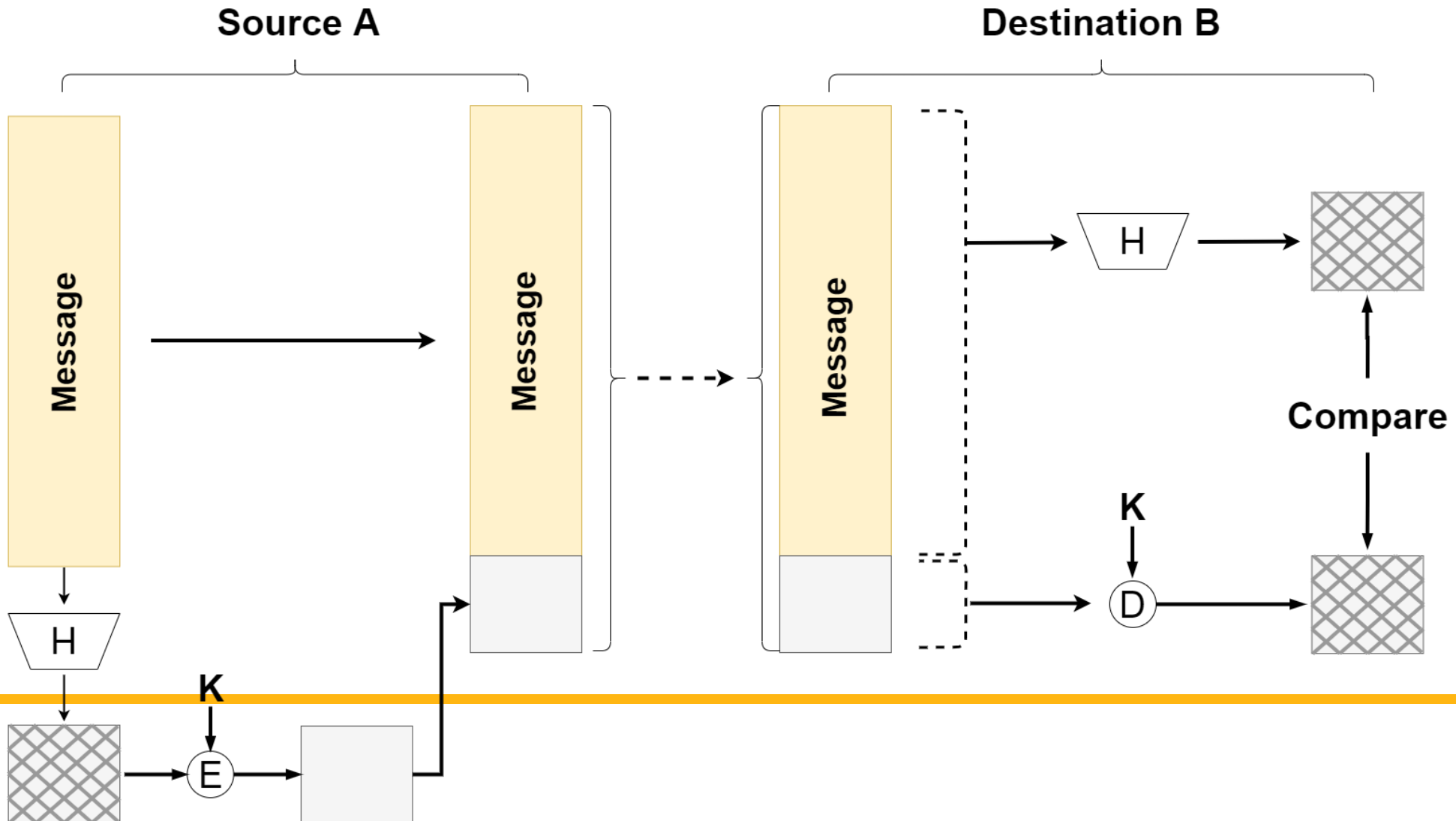
- A variable-size message $M \rightarrow$
Tag: a fixed-size message digest $H(M)$
- Unlike MAC
 - does not take a secret key as input



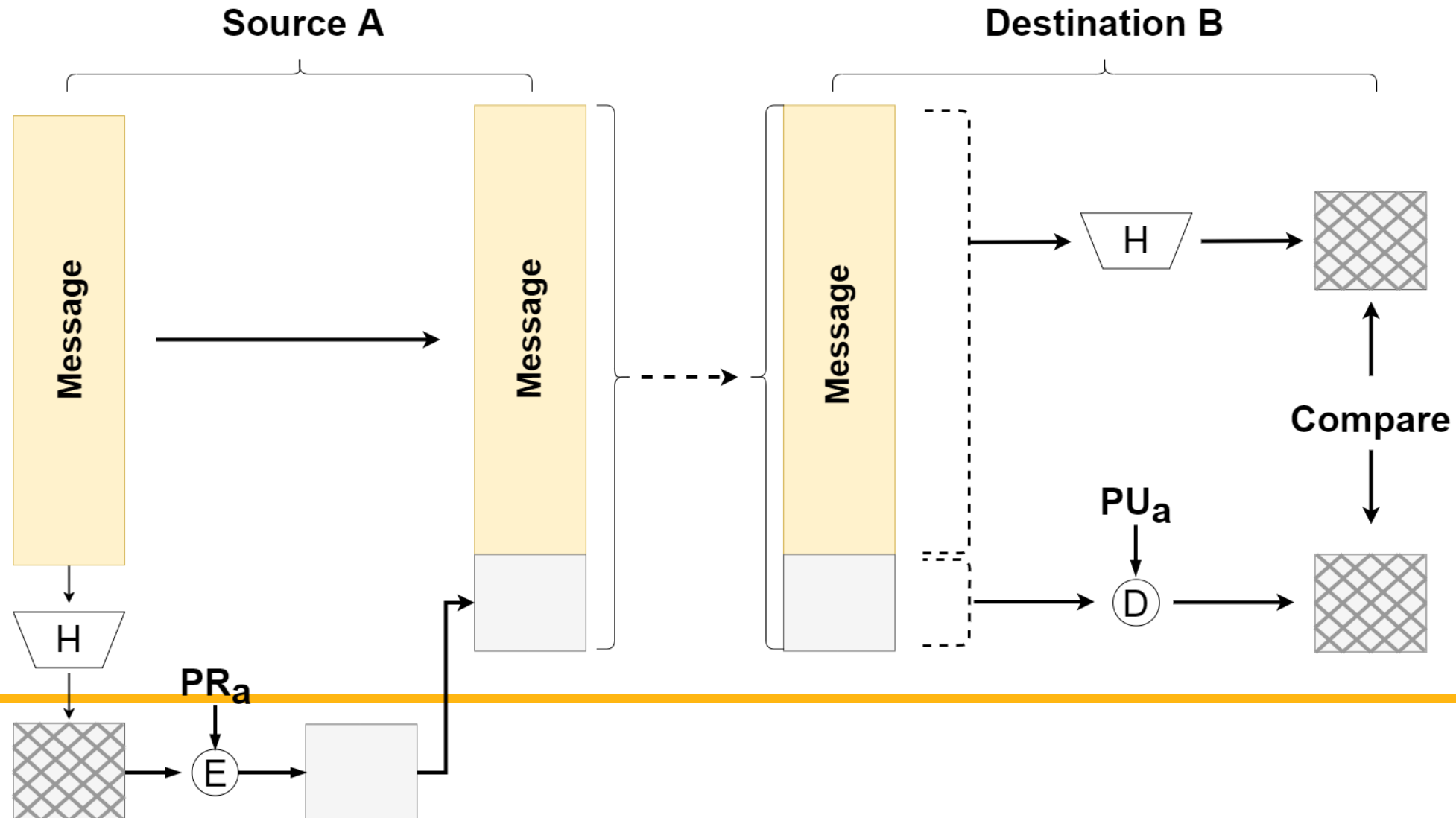
P, L = padding plus length field

The hash value ensures only unaltered contents.
How about authentic source?

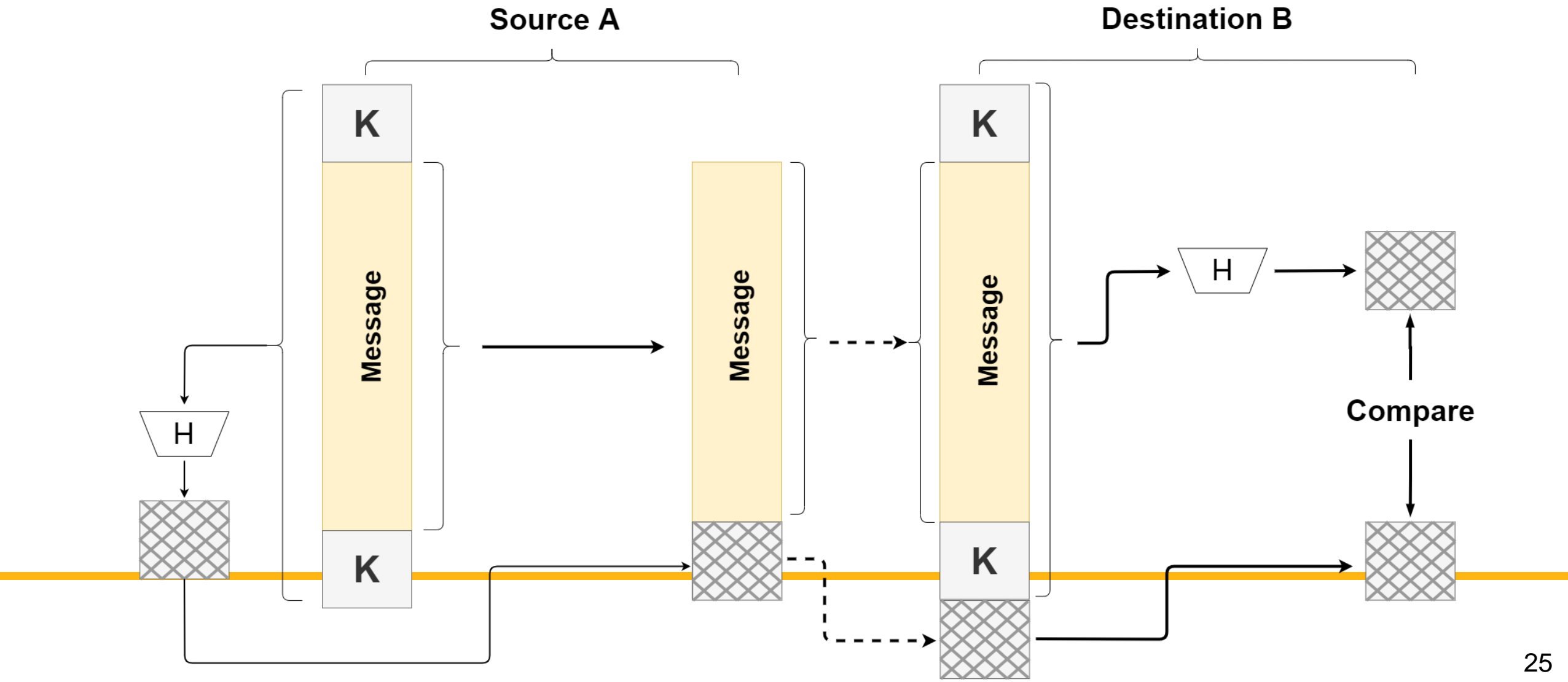
Hash Function w/ Symmetric Encryption



Hash Function w/ Public-key Encryption



Hash Function w/o Encryption: Keyed Hash MAC



Secure Hash Functions

- A hash function H must have the following properties

- H can be applied to a block of data of any size
- H produces a fixed-length output
- $H(x)$ is relatively easy to compute for any given x
 - Making both hardware and software implementations practical
- One-way (pre-image resistant)
 - For any given code h , it is computationally infeasible to find x such that $H(x) = h$
- Second pre-image (weak collision) resistant
 - For any given block x , it is computationally infeasible to find $y \neq x$ with $H(y) = H(x)$
- Collision (strong collision) resistant
 - It is computationally infeasible to find any pair (x, y) such that $H(x) = H(y)$

Security of Hash Functions

- Two attack approaches

- Cryptanalysis: exploits logical weaknesses in algorithms → algorithm dependent
- Brute-force → solely on the length of the hash code

- Security strength against brute-force attacks

- Preimage resistant: 2^n
- Second preimage resistant: 2^n
- Collision resistant: $2^{n/2}$
 - Based on that a birthday attack on a message digest of size n produces a collision
 - 128-bit MD5 [VANO94]: 24 days
 - 160-bit SHA: > 4000 years

Security of Hash Functions (Cont.)

- Most widely used hash algorithm: Secure Hash Algorithm (SHA)
 - ❑ Developed by NIST, and published in FIPS 180, 1993
 - ❑ SHA-1 160-bit (1995)
 - ❑ SHA-256, SHA-384, SHA-512 (2002)
- Applications of hash functions
 - ❑ Message authentication
 - ❑ Digital signatures (later in this Chapter)
 - ❑ Passwords (Chapter 3)
 - ❑ Intrusion detection (Chapter 8)

Public-Key Encryption

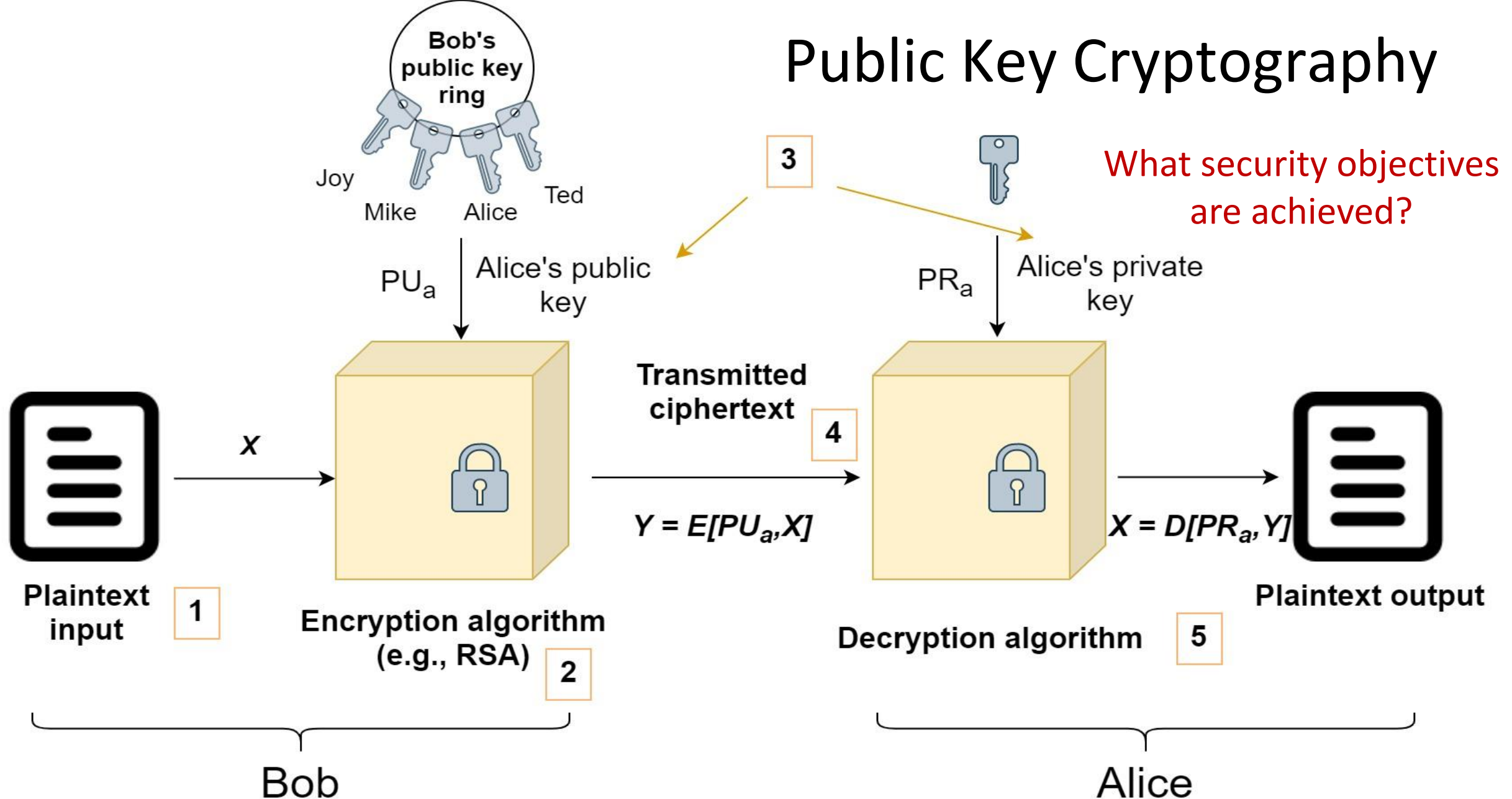
- Public-key cryptography

- Proposed by Diffie and Hellman in 1976
- Asymmetric algorithm
 - Two separate keys: public and private keys
- Based on mathematical functions
 - Different from symmetric algorithms: simple operations on bit patterns

- Three common *misconceptions* for public-key encryption

- More secure than symmetric ones
- A general-purpose technique that has made symmetric ones obsolete
- Key distribution is trivial

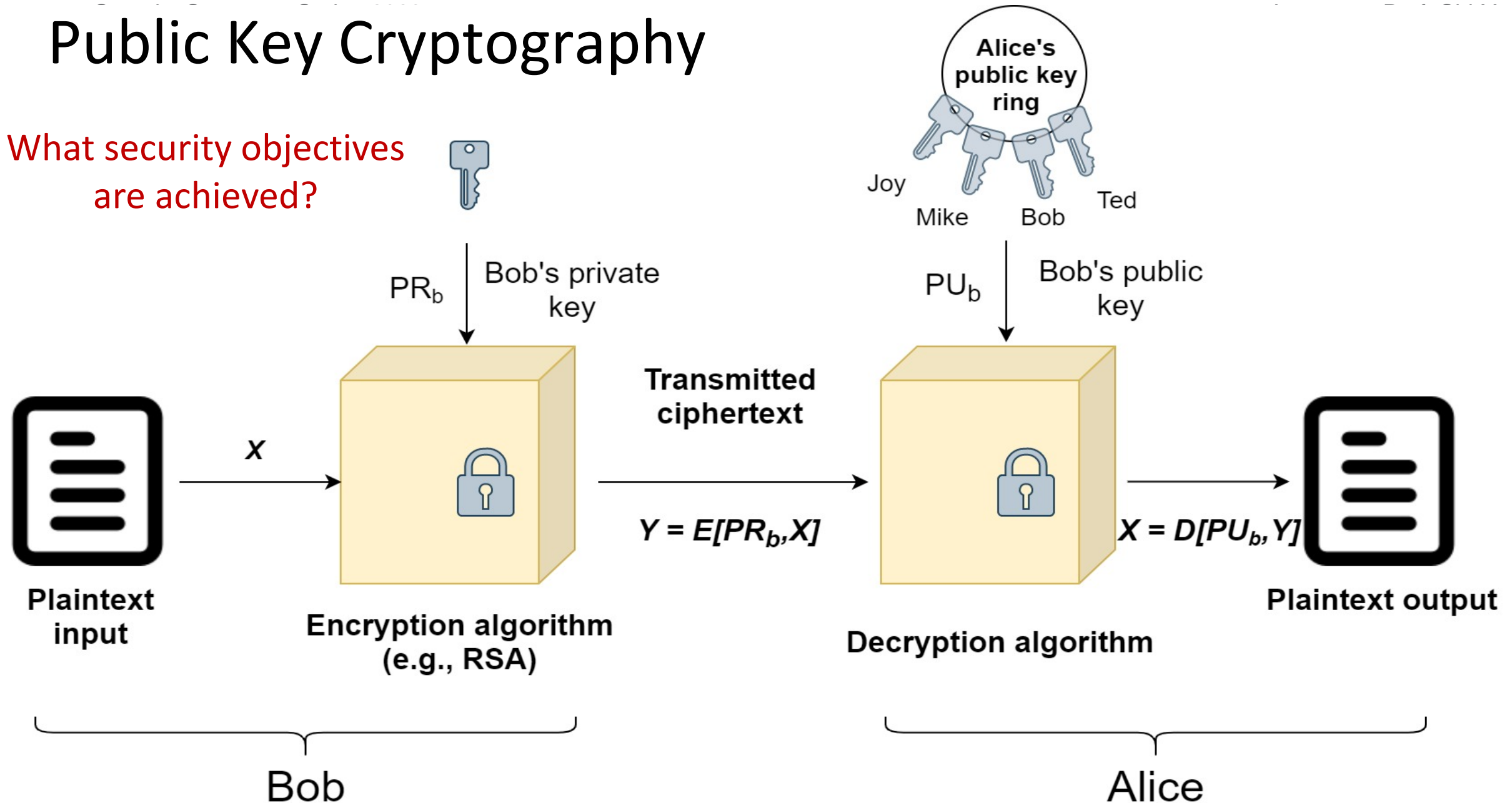
Public Key Cryptography



(a) Encryption with public key

Public Key Cryptography

What security objectives are achieved?



(b) Encryption with private key

Requirements for Public-Key Cryptosystems

- Computationally easy
 - ❑ Create key pairs
 - ❑ for sender knowing public key to encrypt messages
 - ❑ for receiver knowing private key to decrypt ciphertext
- Computationally infeasible for opponent knowing public key
 - ❑ Determine private key
 - ❑ Recover original message, which is encrypted by public key
- Either of private and public keys can be used for encryption

Encryption Algorithms and Applications

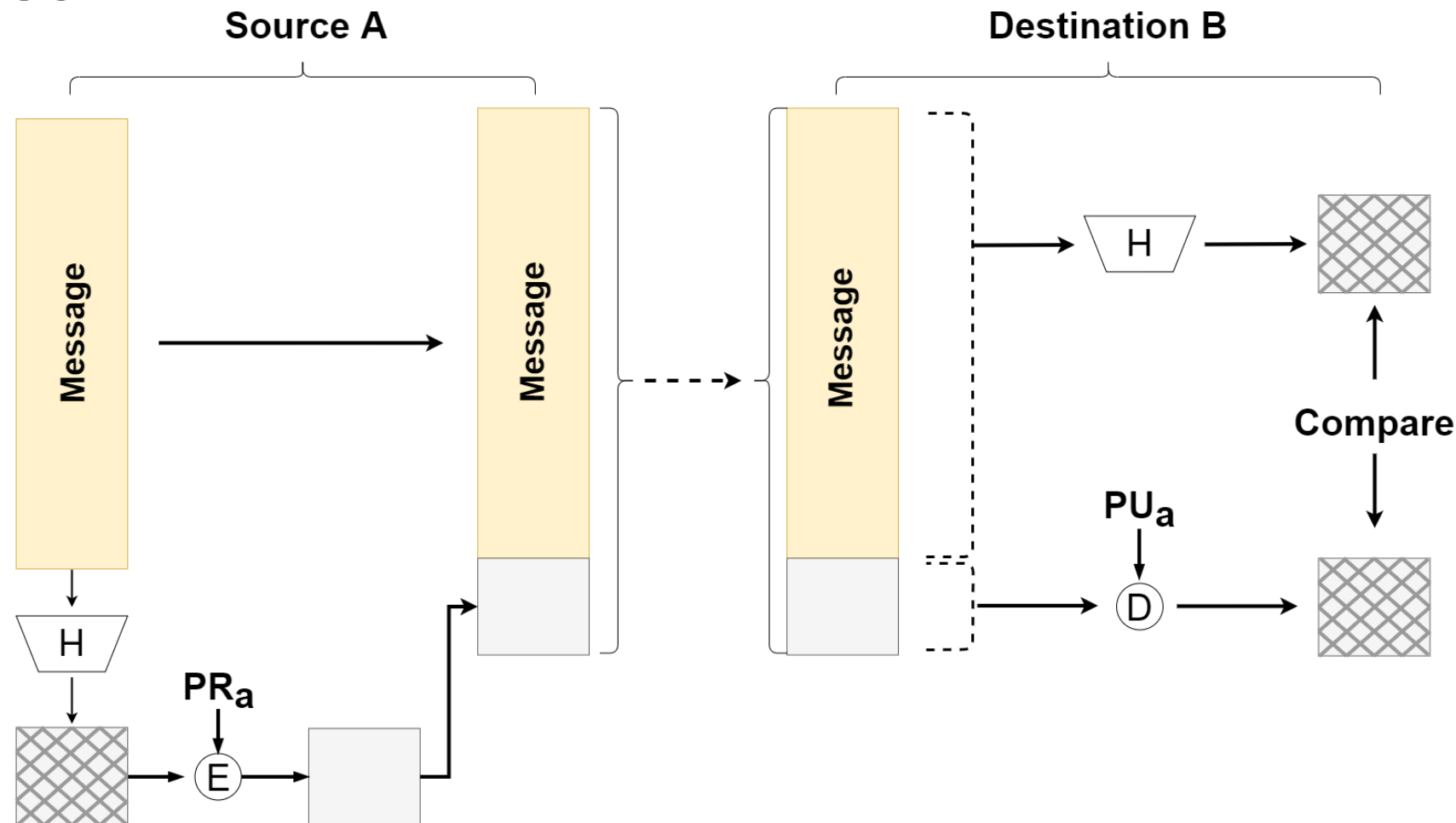
Algorithm	Digital Signature	Symmetric Key Distribution	Encryption of Secret Keys
RSA	Yes	Yes	Yes
Diffie-Hellman	No	Yes	No
DSS	Yes	No	No
Elliptic Curve	Yes	Yes	Yes

Digital Signatures and Key Management

- Digital signatures
- Public-key certificates: secure distribution of public keys
- Symmetric key exchange using public-key encryption
- Digital envelopes: distribution of secret keys

Digital Signatures

- Encrypts hash code with private key
- What security objectives can be achieved?
- Do they provide confidentiality?



Public-Key Certificates

- Public key distribution

- Any person can release his or her public key
- But anyone can forge such a public announcement

- Solution: public-key certificate

- Certificate: a public key + a user ID of the key owner
- The whole block signed by a trusted third party, CA (Certificate Authority)
 - CA has to be trusted by the user community (e.g., government)
- Certificate also includes CA information and validity period

- X.509 standard: universally accepted certificates

- Used in most network security apps: IPSec, TLS, SSH, S/MIME

Survey of CAs

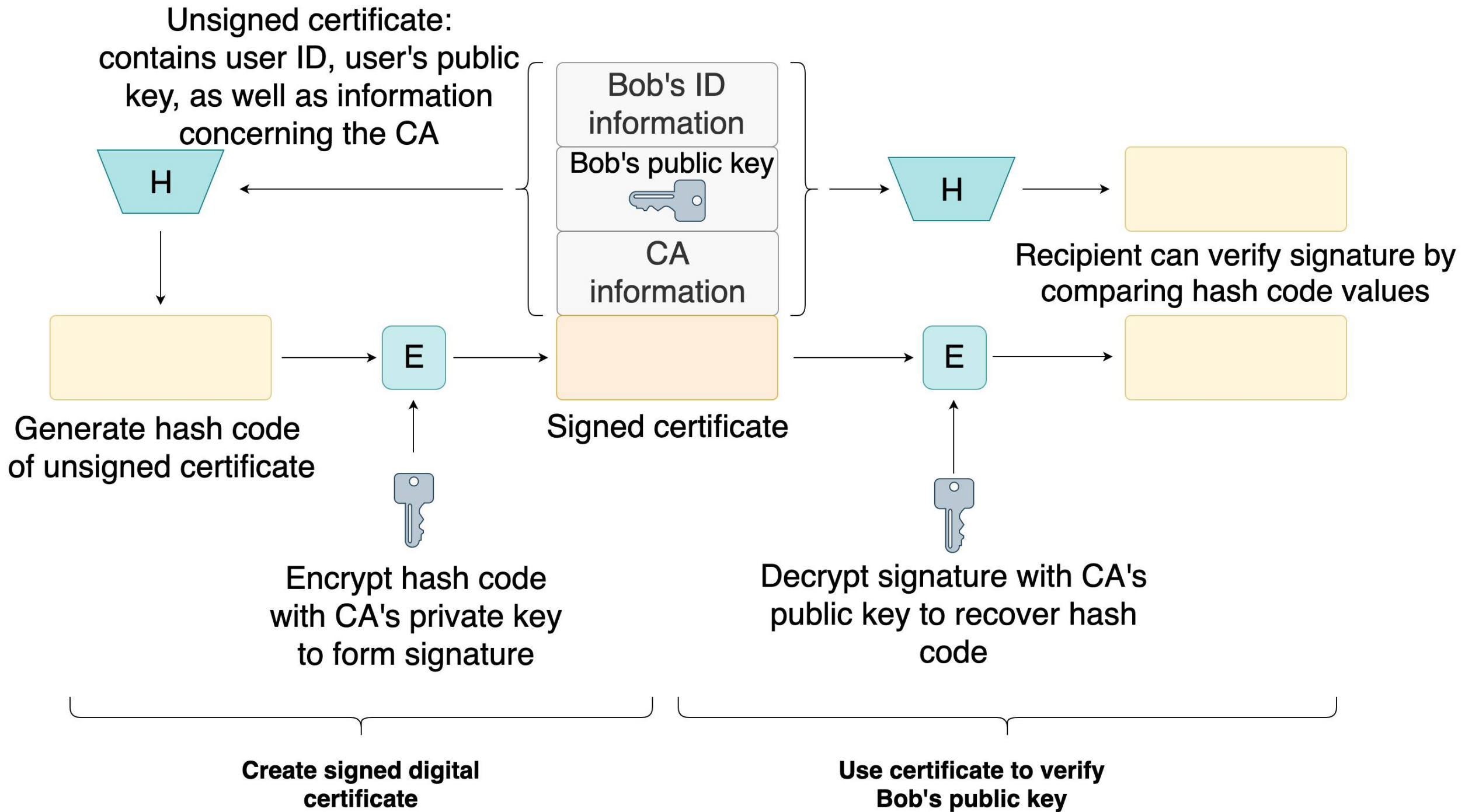
2017

Rank	Issuer	Usage	Market share
1	<u>Comodo</u>	16.7%	38.4%
2	<u>IdenTrust</u>	13.9%	32.0%
3	<u>Symantec</u>	5.6%	12.9%
4	<u>GoDaddy</u>	3.3%	7.5%
5	<u>GlobalSign</u>	1.9%	4.5%
6	<u>DigiCert</u>	1.0%	2.2%
7	Certum	0.3%	0.7%
8	<u>Entrust</u>	0.2%	0.4%
9	<u>Secom</u>	0.1%	0.3%
10	Actalis	0.1%	0.3%
11	<u>Trustwave</u>	0.1%	0.2%
12	<u>Let's Encrypt</u>	0.1%	0.2%

2018

Rank	Issuer	Usage	Market share
1	<u>IdenTrust</u>	20.4%	39.7%
2	<u>Comodo</u>	17.9%	34.9%
3	<u>DigiCert</u>	6.3%	12.3%
4	<u>GoDaddy</u>	3.7%	7.2%
5	<u>GlobalSign</u>	1.8%	3.5%
6	Certum	0.4%	0.7%
7	Actalis	0.2%	0.3%
8	<u>Entrust</u>	0.2%	0.3%
9	<u>Secom</u>	0.1%	0.3%
10	<u>Let's Encrypt</u>	0.1%	0.2%
11	<u>Trustwave</u>	0.1%	0.1%
12	WISeKey Group	< 0.1%	0.1%

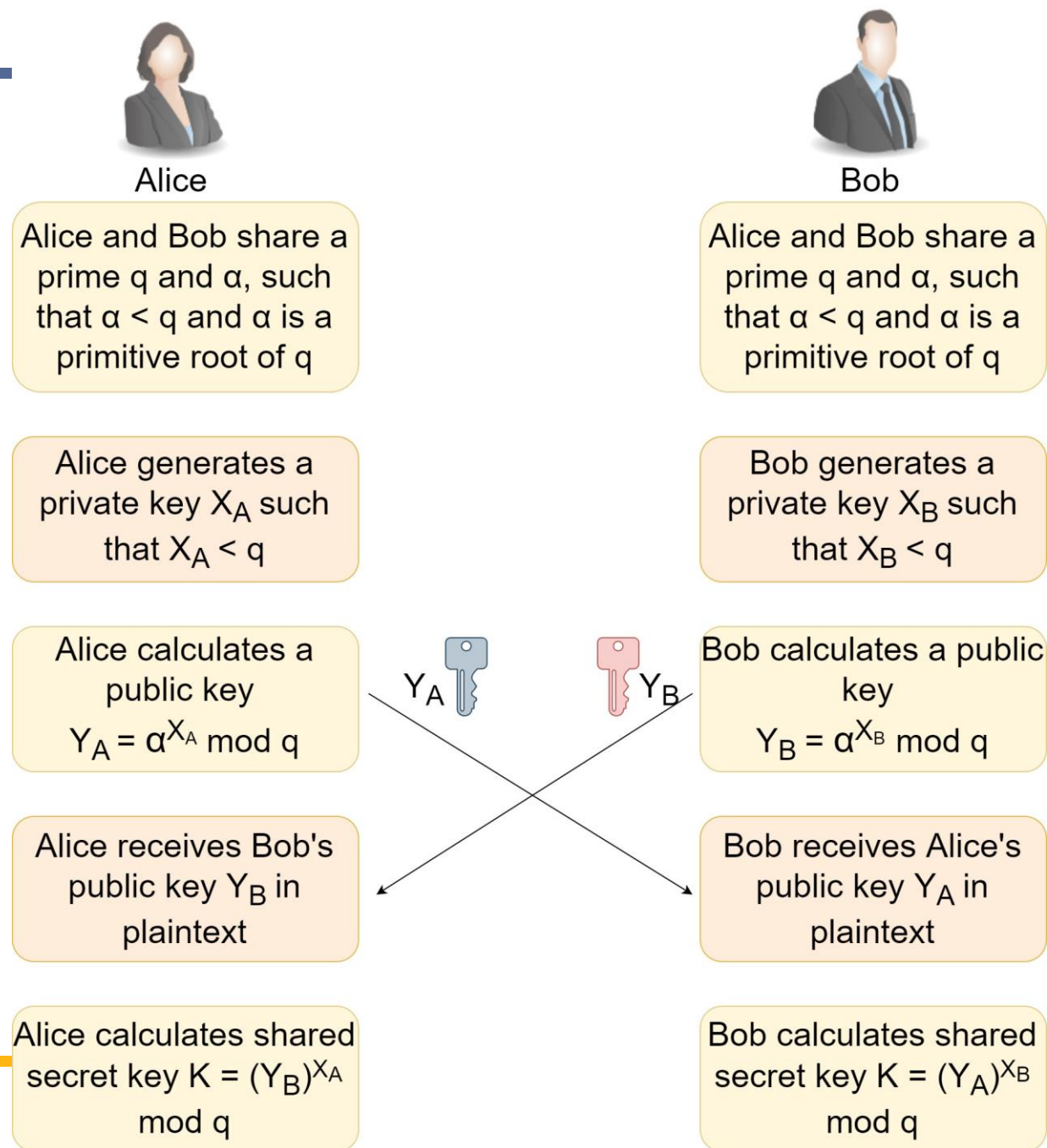
Reference: Surveys from w3techs.com



Symmetric Key Exchange

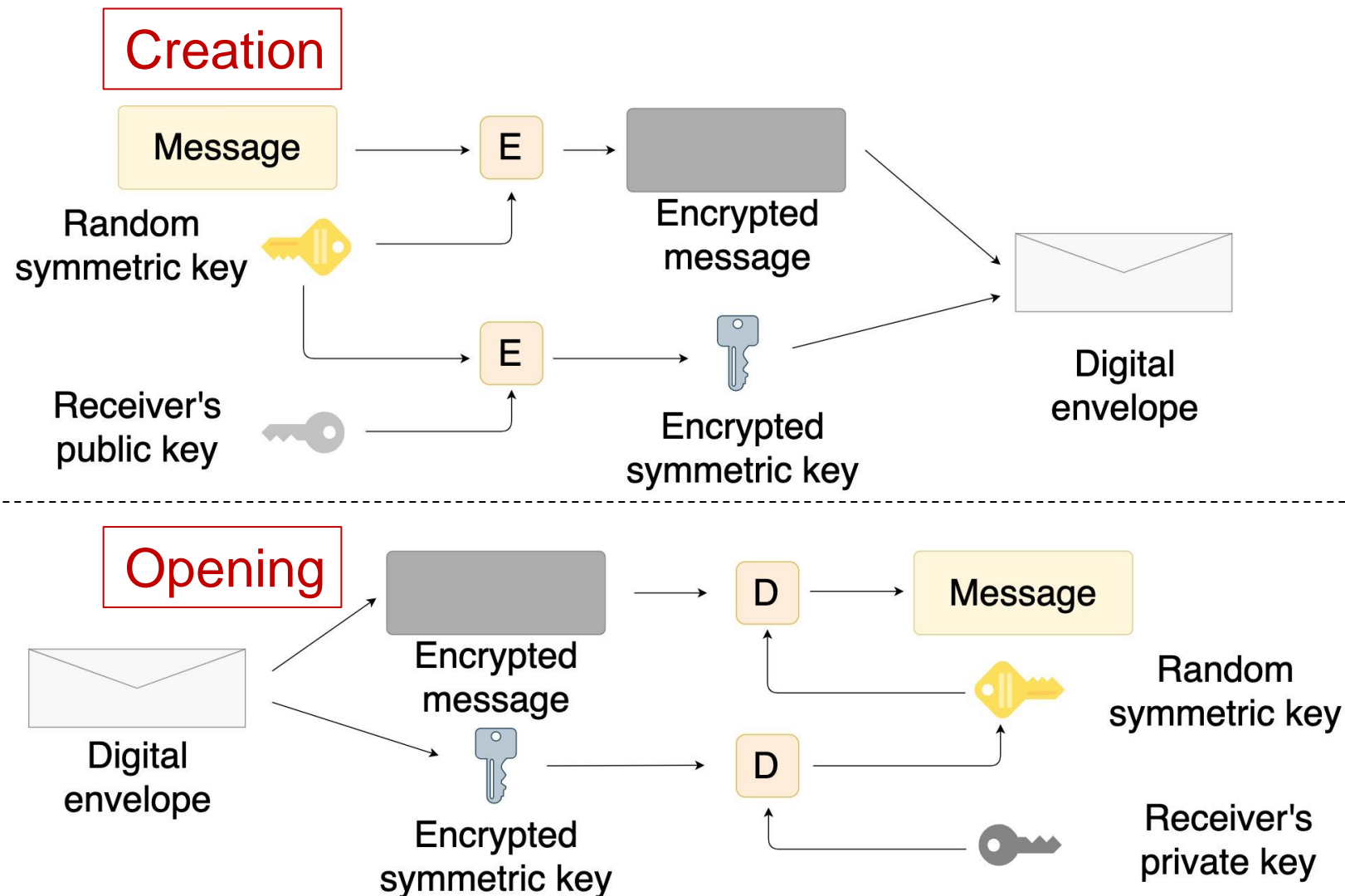
● Diffie-Hellman key exchange

- ❑ Major drawback: no authentication of two communicating partners
- ❑ Many variations proposed to overcome this problem



Digital Envelopes

- Protect a message without having the same secret key between sender/receiver
 - Equate to the same thing as a sealed envelope containing a letter
- Can also be used to deliver a symmetric key only



Random Numbers

- Used for generation of

- Keys for public-key algorithms, stream keys for symmetric stream cipher, symmetric keys, etc.

- Two distinct requirements

- Randomness

- Uniform distribution: frequency of each number should be roughly the same
 - Independence: no one value can be inferred from the others

- Unpredictability

- Not so much that the sequence of numbers be statistically random
 - But, the successive members of the sequence are unpredictable

Random vs. Pseudorandom

- Pseudorandom: its use is widely accepted
 - ❑ Apps typically use deterministic algorithms
 - ❑ Sequences: not statistically random, but can satisfy statistical randomness tests
- Random: true random number generator (TRNG)
 - ❑ Use a nondeterministic source to produce randomness
 - ❑ Most operate by measuring unpredictable natural processes
 - E.g., radiation, gas discharge, leaky capacitors, thermal noise
 - ❑ Increasing provided on modern chips

Questions?