

Chapter 9: Virtual-Memory Management

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Chapter 9: Virtual Memory

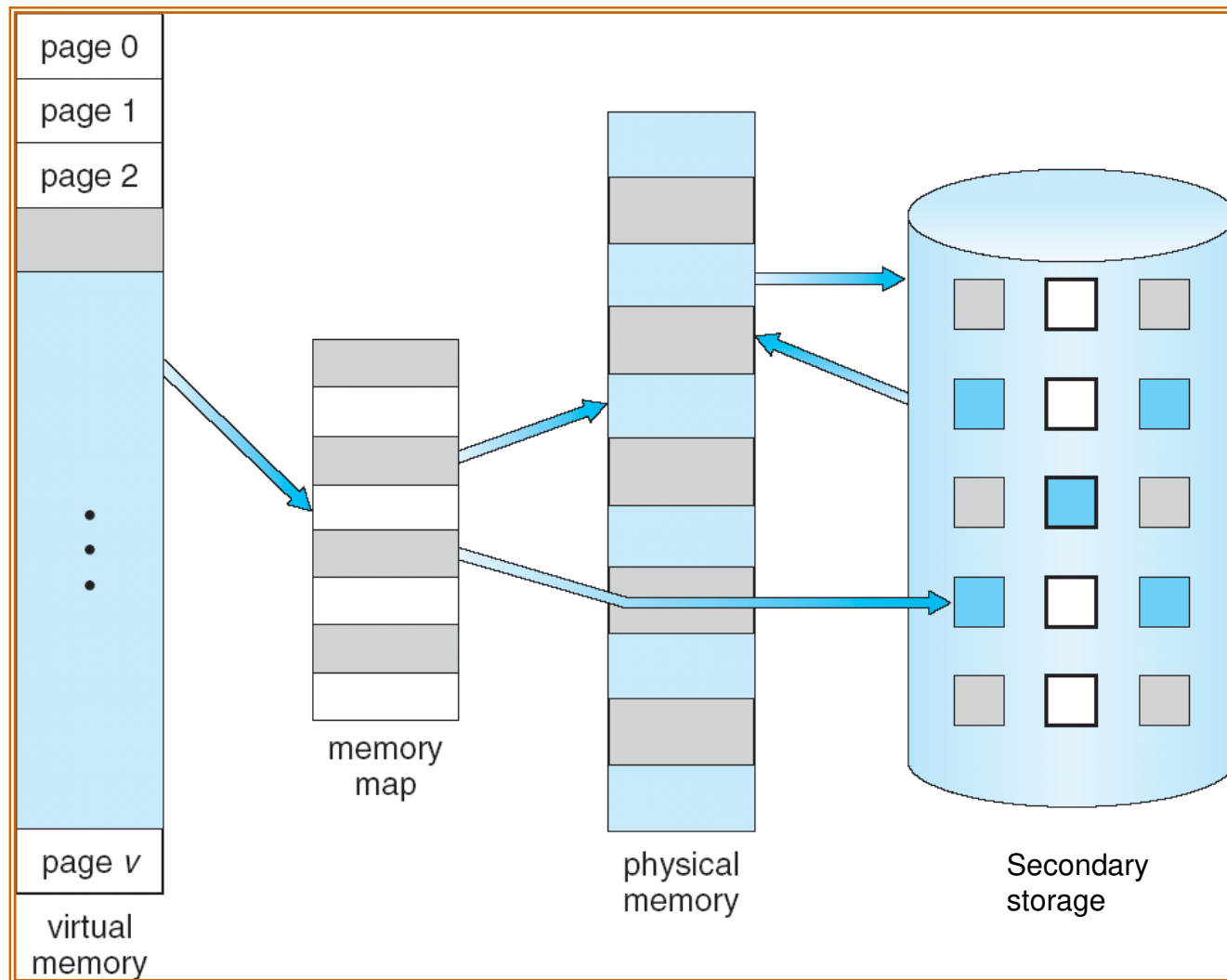
- Demand Paging
- Page Replacement
- Swapping
- Memory-Mapped Files
- Copy-on-Write
- Shared memory
- Performance Issues
- Operating System Examples

DEMAND PAGING

Virtual Memory based on Demand Paging

- Recap: Virtual memory – separation of user logical memory from physical memory
- Virtual memory based on demand paging
 - Only **part of the program** needs to be in memory for execution
 - Logical address space can therefore be **much larger** than physical address space
 - Allows for more efficient process creation
 - Improves the degree of multiprogramming
 - Allows address spaces to be shared by several processes

Virtual Memory That is Larger Than Physical Memory



Demand Paging

- An essential mechanism to implement virtual memory; requiring a (large) secondary storage as the backup of memory pages
- Bringing a page to memory only when it is needed
 - Less memory needed
 - Faster response
 - More users/processes
 - Less I/O needed
- Page is needed \Rightarrow reference to it (load or store)
 - In-memory \Rightarrow normal reference
 - not-in-memory \Rightarrow bring to memory

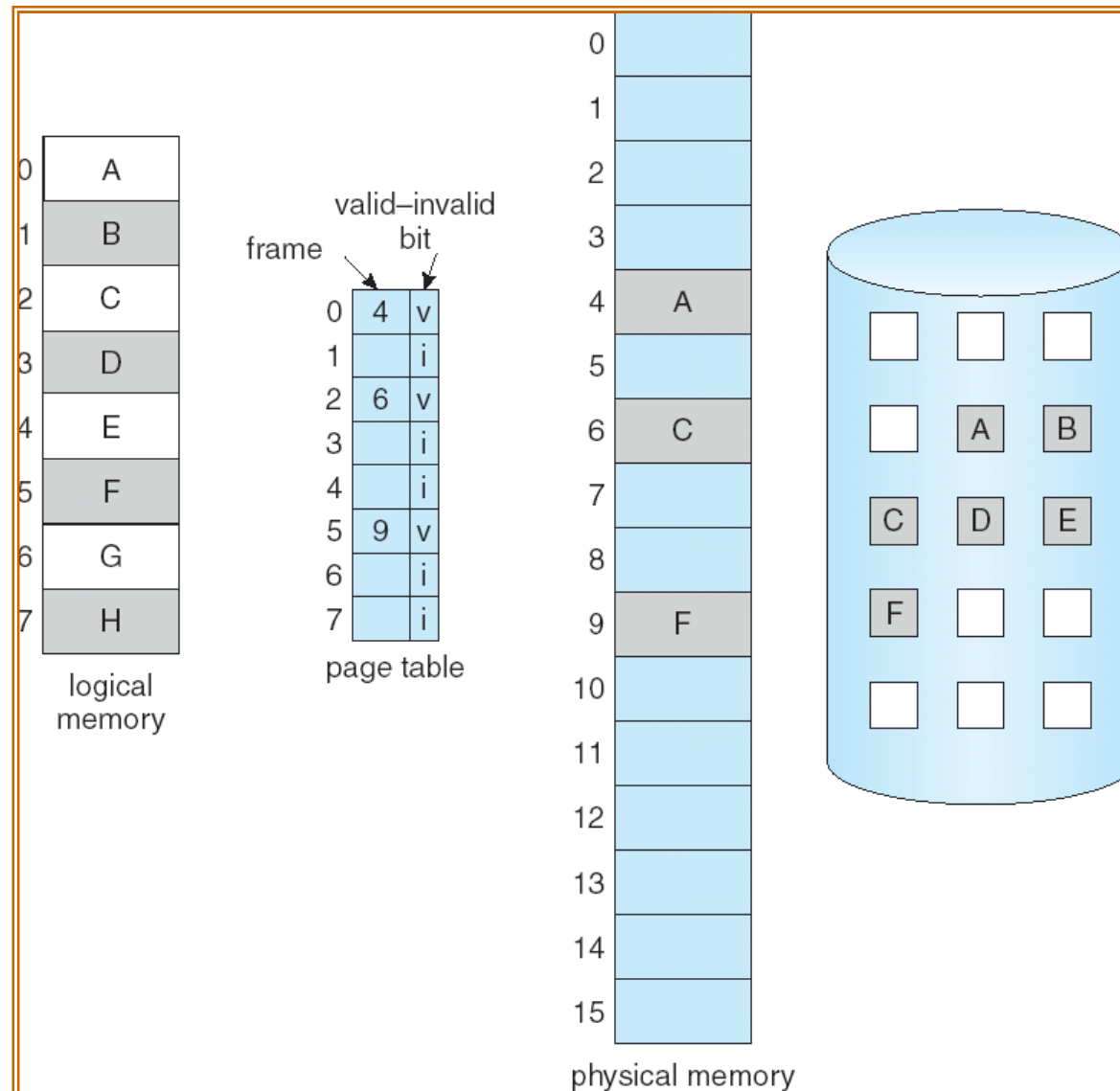
Valid-Invalid Bit of Page Table Entry

- With each page table entry a valid–invalid bit is associated
(1 \Rightarrow in-memory, 0 \Rightarrow not-in-memory)
- Initially valid–invalid bit is set to 0 on all entries
- During address translation, if valid–invalid bit in page table entry is 0 \Rightarrow **page fault**

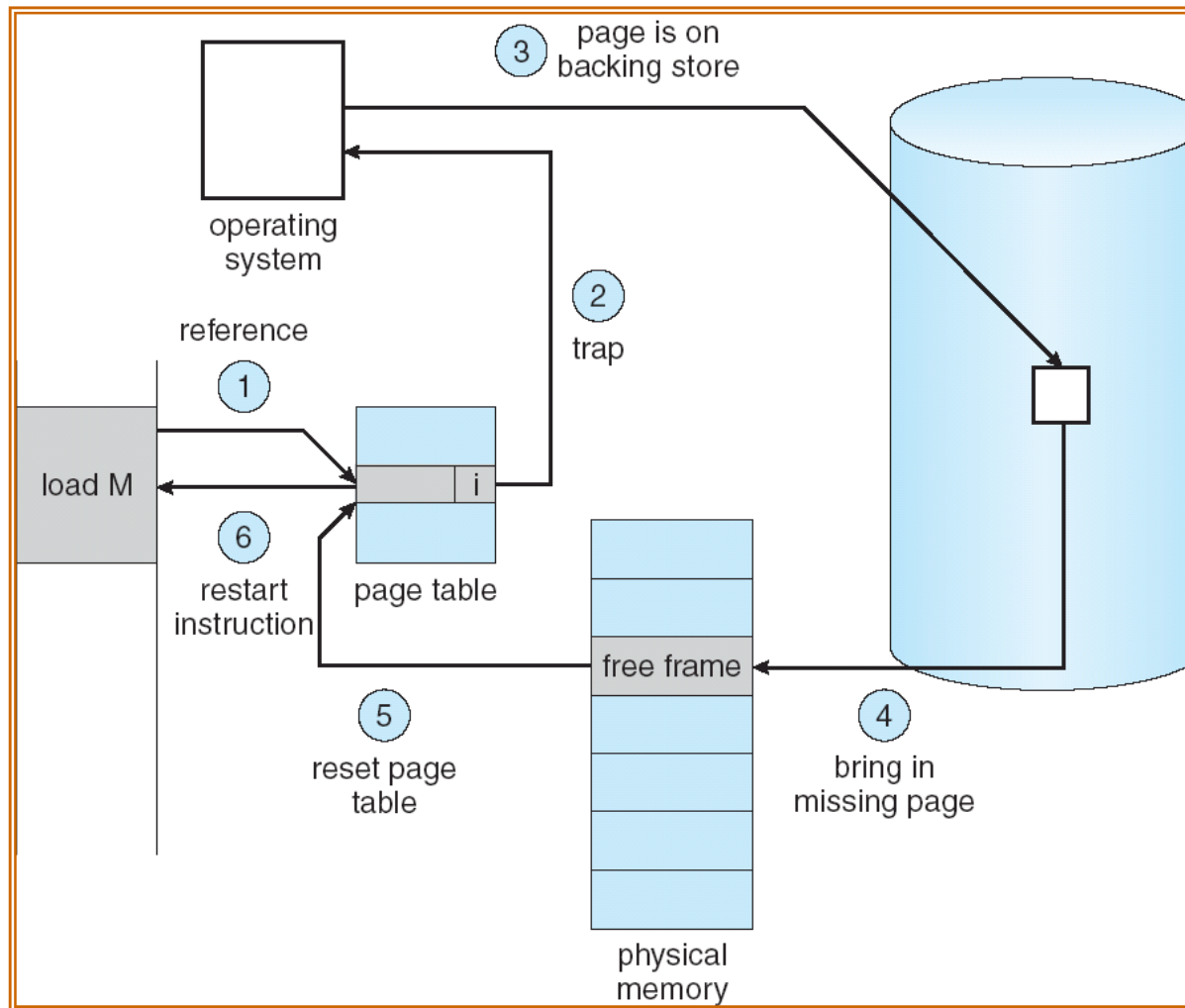
Frame #	valid- invalid bit
	1
	1
	1
	1
	0
⋮	
	0
	0

page table

Page Table When Some Pages Are Not in Main Memory



Steps of Handling a Page Fault



1. Trap to the operating system.
2. Save the user registers and process state.
3. Determine that the interrupt was a page fault.
4. Check that the page reference was legal and determine the location of the page on the disk.
5. Issue a read from the disk to a free frame:
 - a. Wait in a queue for this device until the read request is serviced.
 - b. Wait for the device seek and/or latency time.
 - c. Begin the transfer of the page to a free frame.
6. While waiting, allocate the CPU to some other user (CPU scheduling, optional).
7. Receive an interrupt from the disk I/O subsystem (I/O completed).
8. Save the registers and process state for the other user (if step 6 is executed).
9. Determine that the interrupt was from the disk.
10. Correct the page table and other tables to show that the desired page is now in memory.
11. Wait for the CPU to be allocated to this process again.
12. Restore the user registers, process state, and new page table, and then resume the interrupted instruction.

- Up to how many page faults will be caused by the following instruction?

MOV EAX, DWORD PTR [EDX]

- Total 2
 - The instruction itself once and the access of memory location [EDX] the other one
 - Instructions and data do not span across page boundaries

Performance of Demand Paging

- Page Fault Rate $0 \leq p \leq 1.0$
 - if $p = 0$ no page faults
 - if $p = 1$, every reference is a fault

- Effective Access Time (EAT)

$$\begin{aligned} \text{EAT} = & (1 - p) \times \text{memory access} \\ & + p (\text{page fault overhead} \\ & + [\text{swap page out}] \\ & + \text{swap page in} \\ & + \text{restart overhead}) \end{aligned}$$

Suppose that the page-fault service time is 8 ms, including disk I/O and all necessary memory access

A memory access takes 200ns (TLB hit+TLBmiss)

The page-fault ratio is p

The EAT is

$$(1-p)*200+p*8ms \\ =200+7999800p$$

The RHS term dominates the EAT! p should be as low as possible!!

If $p=1/1000$, $EAT = 200+7999 \sim 8.2\mu s$, 40 times slower!!

If the expected slowdown is no larger than 10% compared to 200ns, then

$$220 \geq 200+7999800p$$

$$20 \geq 7999800p$$

$P \leq 0.0000025$ in other words, no more than 1 page fault should happen out of 399,990 memory access.

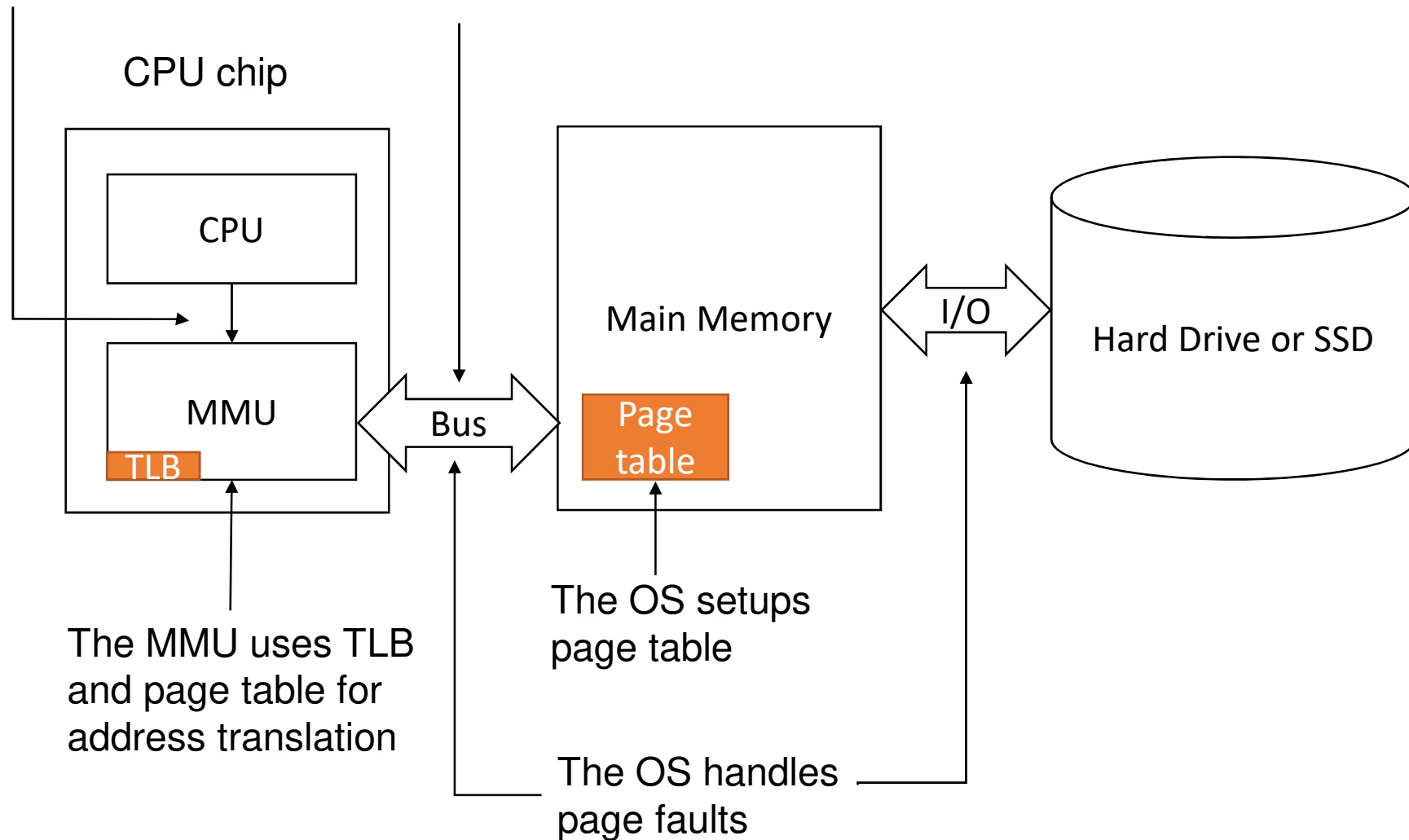
****Consider a demand-paging system with a paging disk that has an average access and transfer time of 10 milliseconds. Addresses are translated through a page table in main memory, with an access time of 4 microsecond per memory access. Thus, each memory reference through the page table takes two accesses. To improve this time, we have added an associative memory that reduces access time to one memory reference, if the page-table entry is in the associative memory.**

Assume that 87.5% of the accesses are in the associative memory and that, 20% of the remain 12.5% cause page faults. What is the effective memory access time?

- If TLB hit
 - won't be a page fault
 - TLB access
- If TLB miss
 - If not a page fault
 - TLB access + page table access + TLB update
 - If a page fault
 - The page is not in memory
 - TLB access + page table access + page fault handling + TLB update

The CPU sends
virtual address

The MMU generates
physical addresses

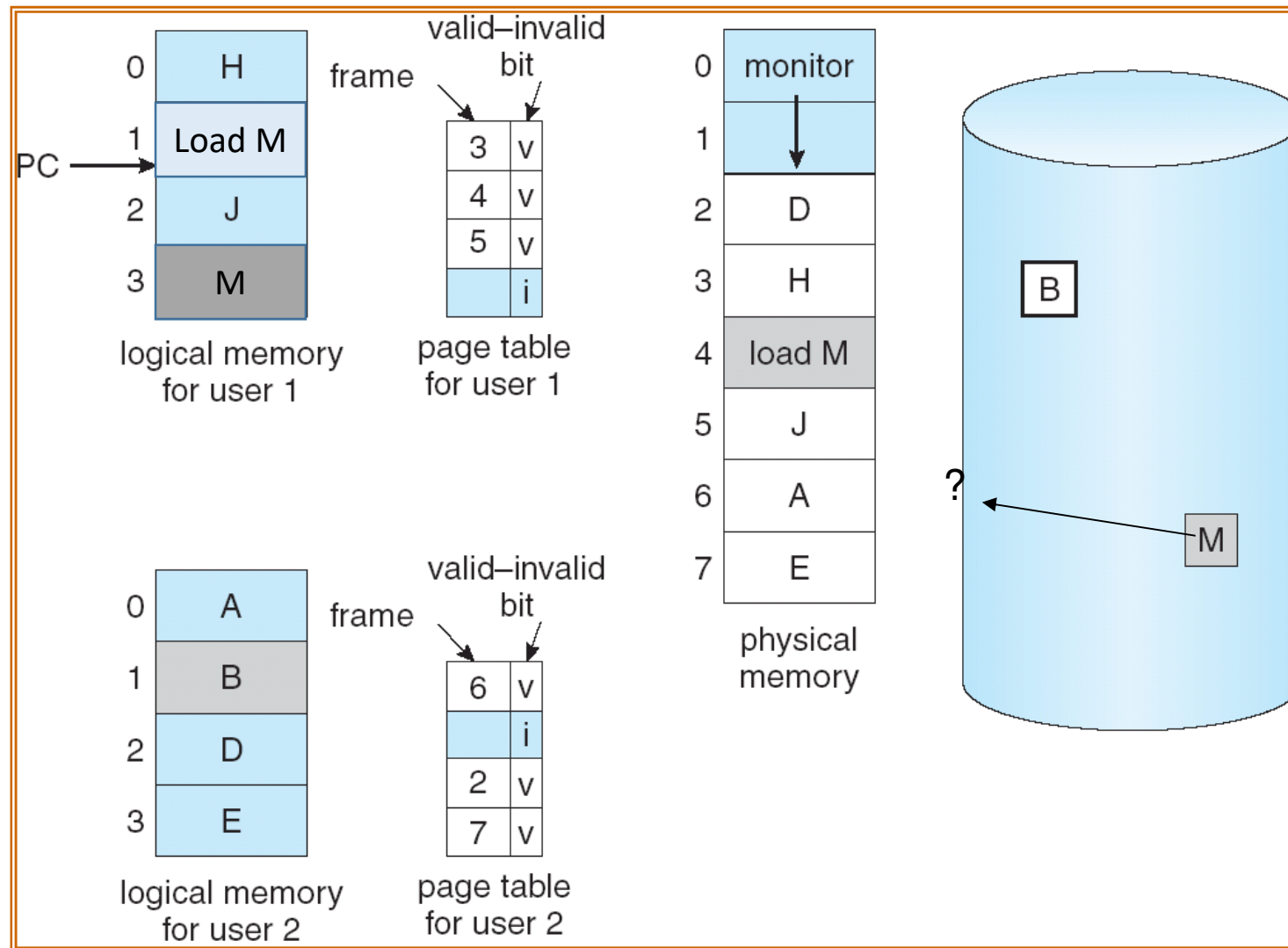


PAGE REPLACEMENT

Page Replacement

- Prevent over-allocation of memory by modifying page-fault service routine to include page replacement
- Page replacement completes separation between logical memory and physical memory – large virtual memory can be provided on a smaller physical memory

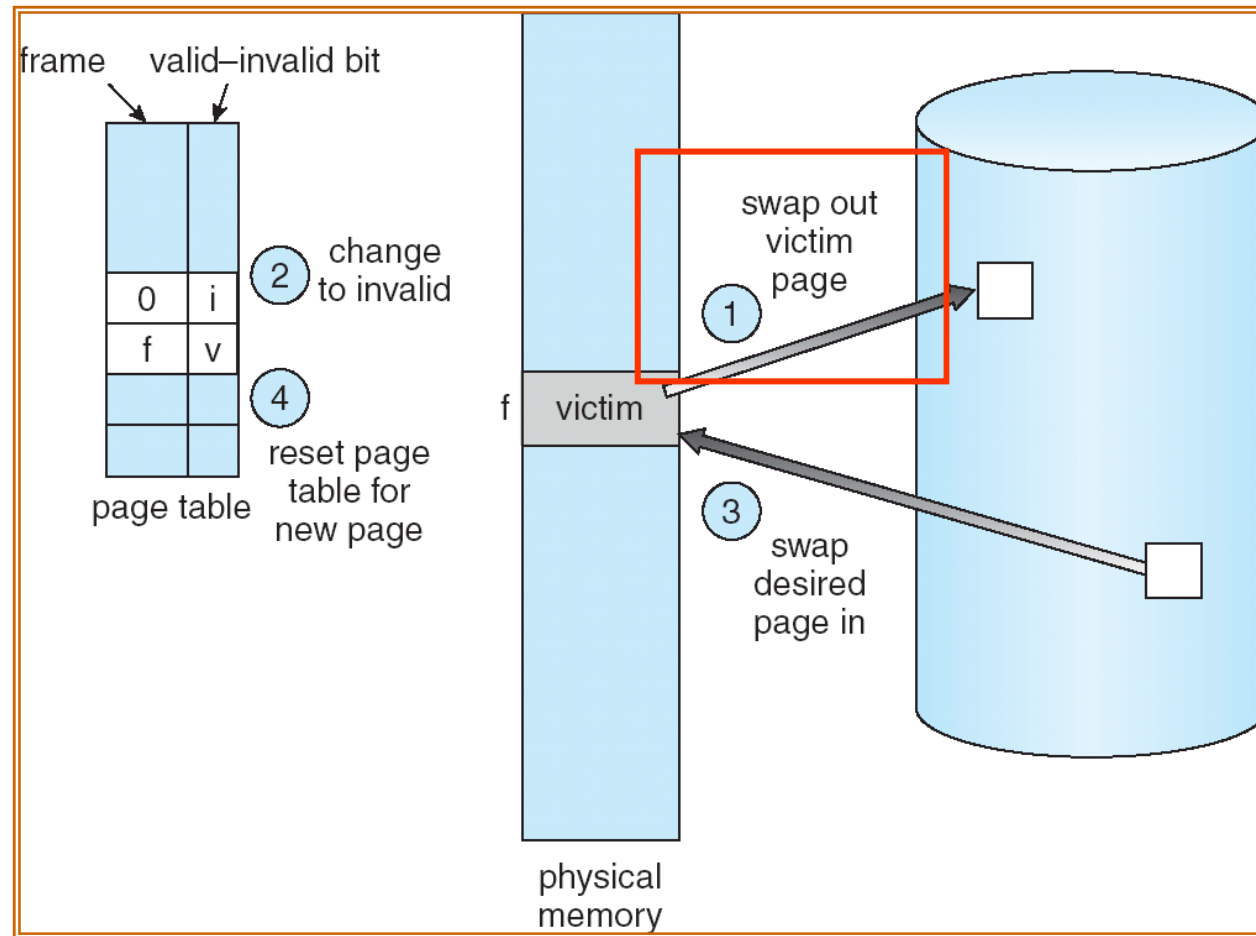
Need For Page Replacement



Basic Page Replacement

- Find the location of the desired page on disk
- Find a free frame:
 - If there is a free frame, use it
 - If there is no free frame, use a page replacement algorithm to select a **victim** frame
- Read the desired page into the free frame
- Update the page and frame tables
- Restart the process

Page Replacement



Page Replacement

- To replace a page, the victim page must be **written back** to the backing store. It may double the time of page-fault handling
- Use modify (dirty) bit to reduce overhead of page transfers – only modified pages are written to disk

Page Replacement Policy

- Replace a page that is unlikely to be used in the near future to reduce the overhead of reading a page from disk

Reference string

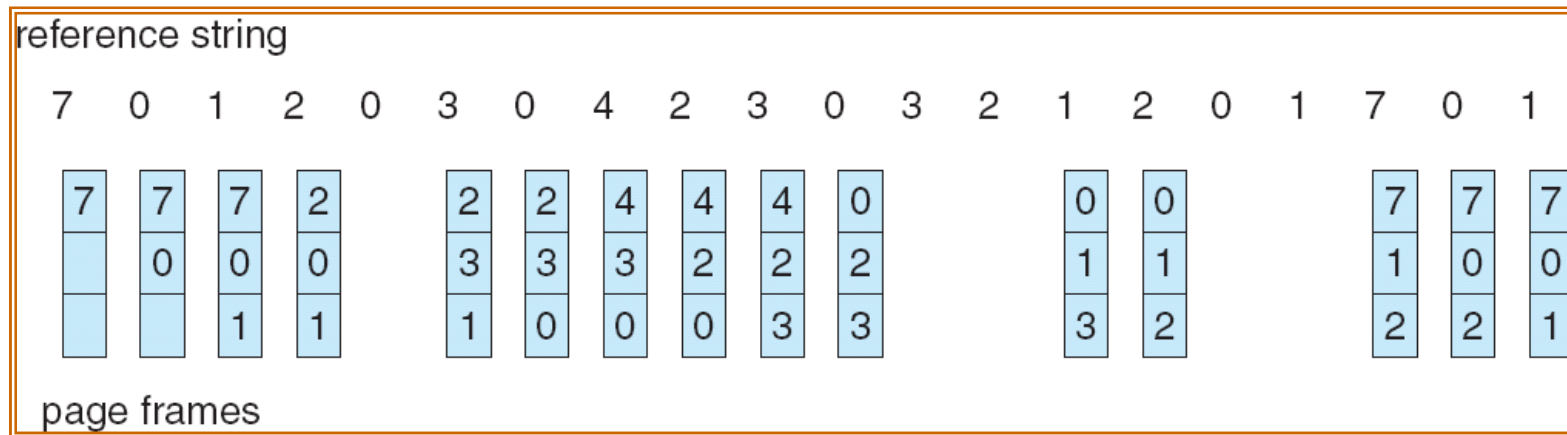
0100, 0432, 0101, 0612, 0102, 0103, 0104, 0101, 0611, 0102, 0103,
0104, 0101, 0610, 0102, 0103, 0104, 0101, 0609, 0102, 0105



Page size=100B

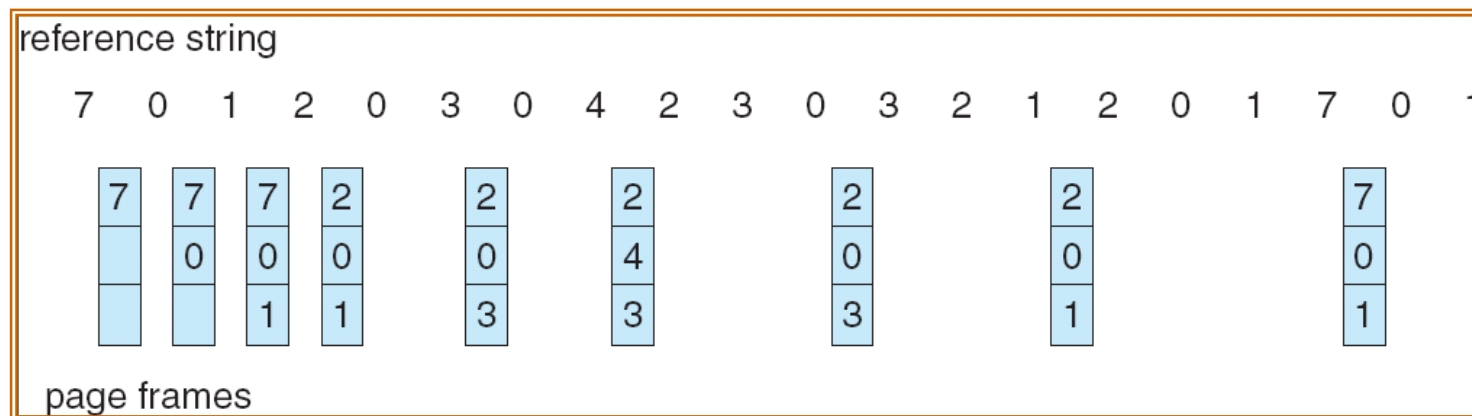
1, 4, 1, 6, 1, 6, 1, 6, 1, 6, 1

FIFO Page Replacement



15 page faults

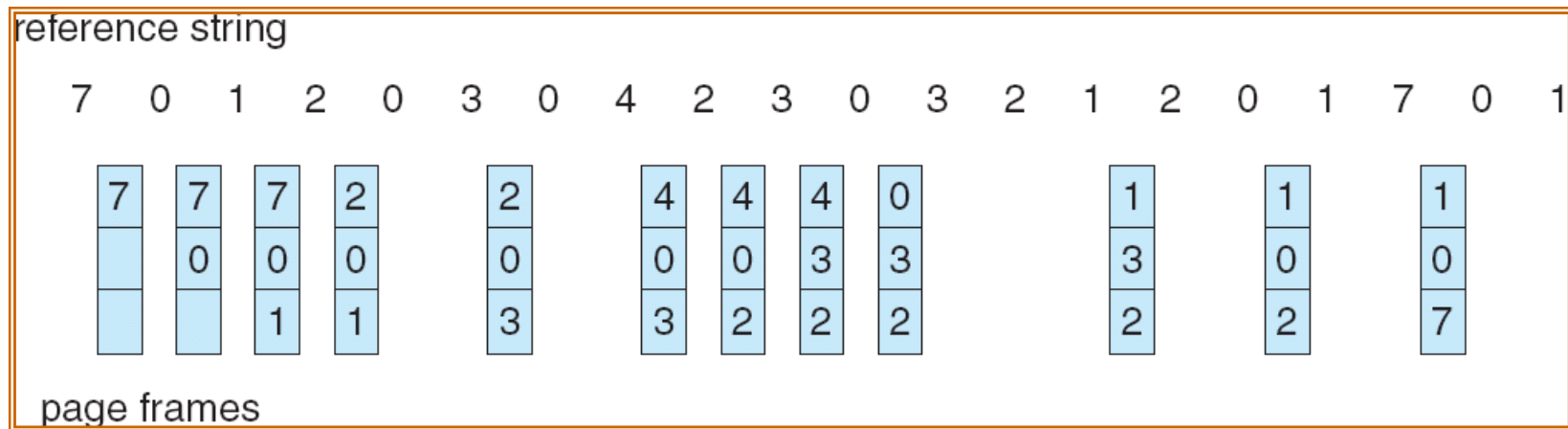
Optimal Page Replacement (OPT)



9 page faults.

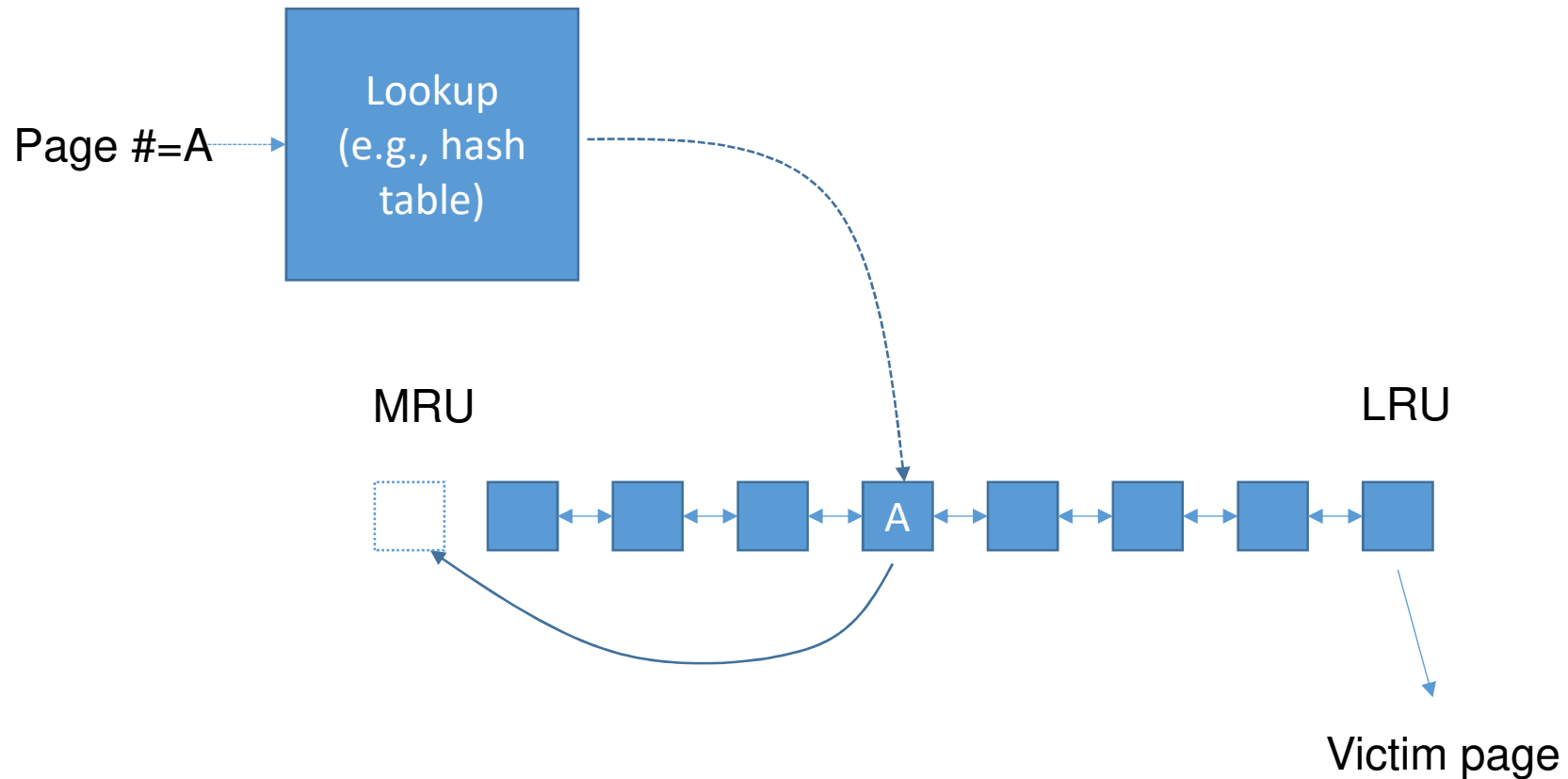
OPT is not applicable to real systems, however.

Least-Recently Used (LRU)

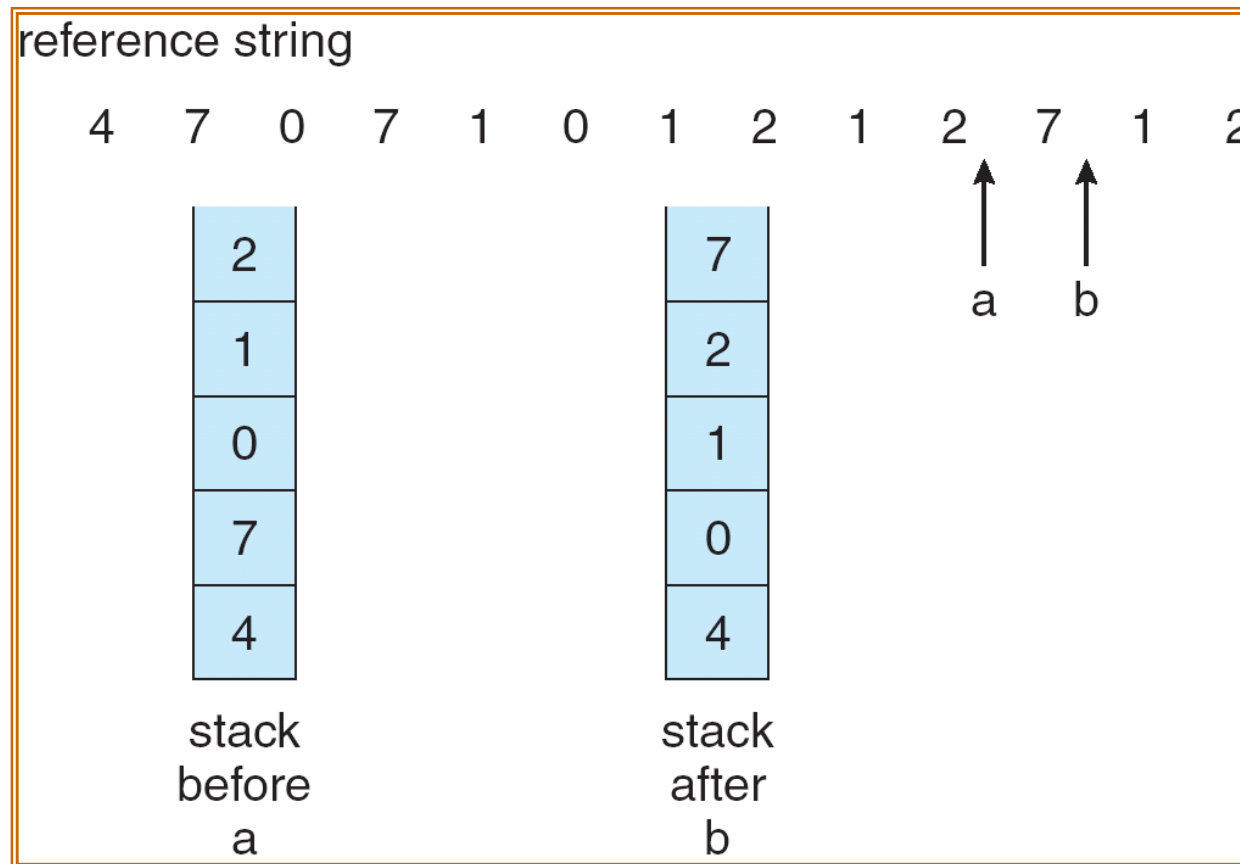


12 page faults.

An LRU Implementation



The “Stack” Property of LRU

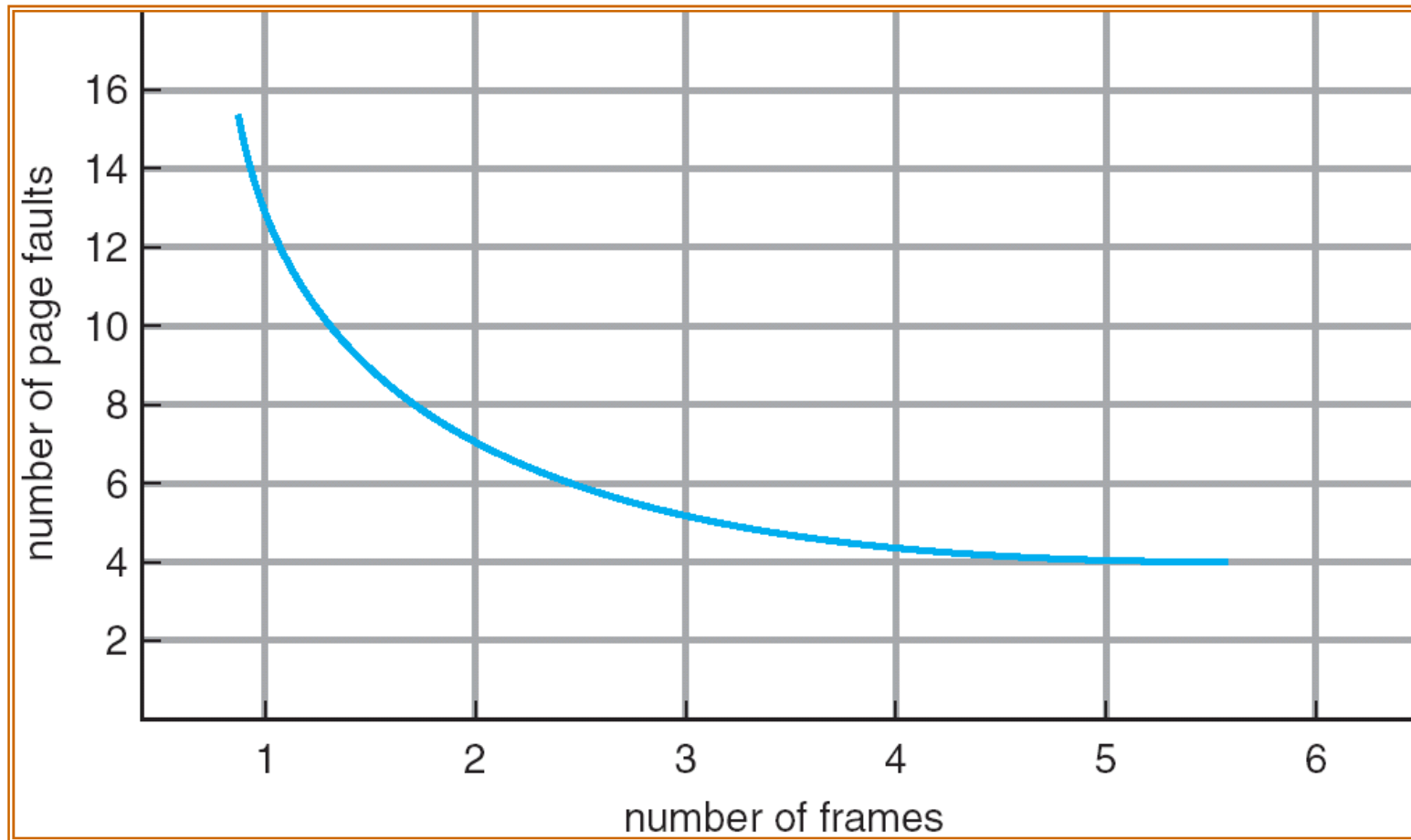


See what happens if there are four frames only.

The “Stack” Property of LRU

- The page set of LRU with K pages is a superset of that of LRU with $K-1$ pages; can be proved by math. induction
- This property suggests that LRU performance is always better as more frames are available
- Also useful in cache simulation. Simulating K pages get results of $K, K-1, K-2, \dots$

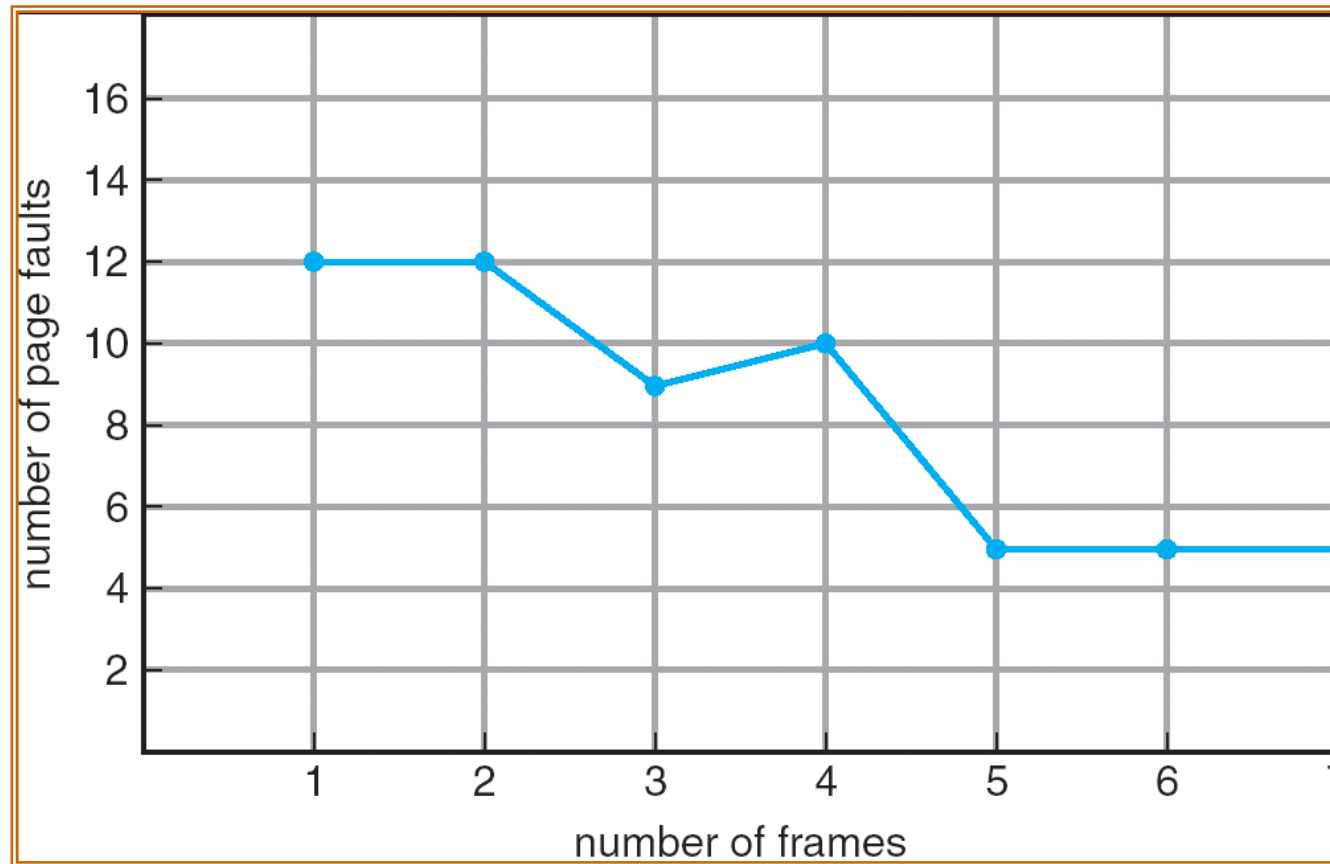
Trends of Page Fault # vs. Frame



The margin quickly saturates beyond a sweet spot

Belady's Anomaly of FIFO

1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5



FIFO does not have the stack property
Adding more frames unexpectedly defrags the performance

Reproducing Belady's Anomaly

1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5

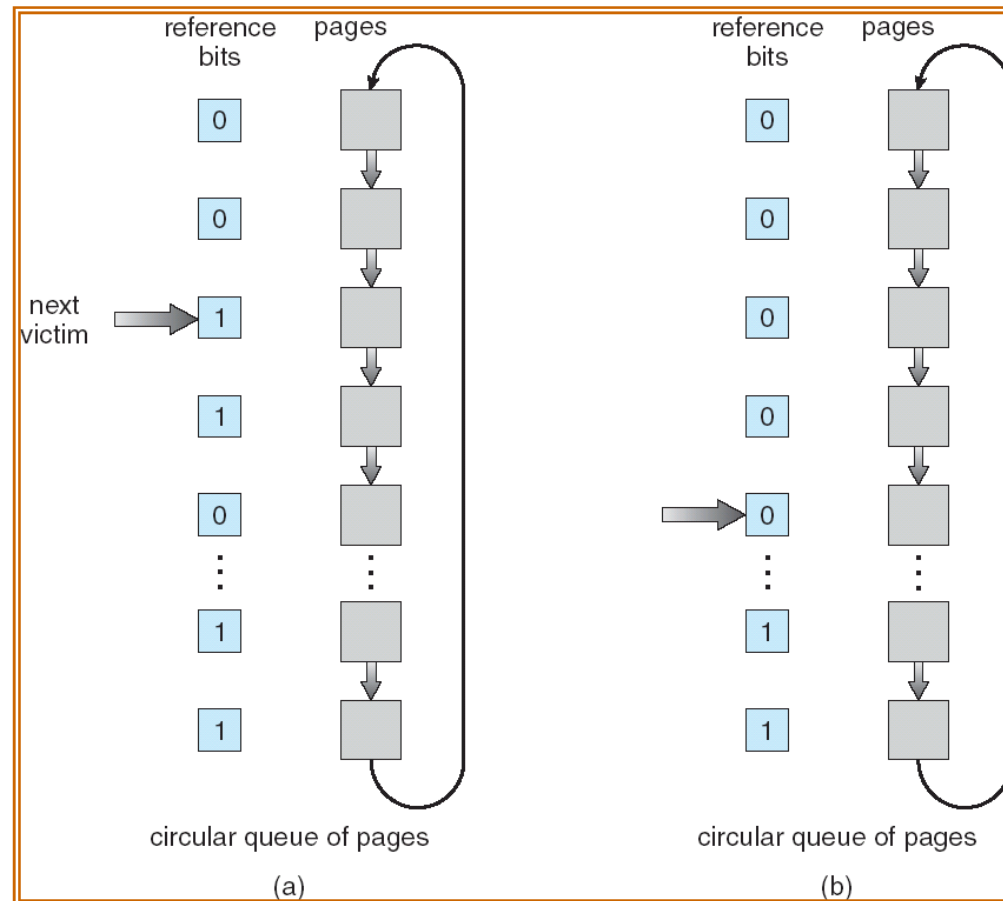
Belady's Anomaly

- FIFO suffers from Belady's anomaly
- LRU and OPT do not
- LRU and OPT are “stack algorithms”, i.e., the page set with n frames is a subset of that with $n+1$ frames

LRU Approximation Algorithms

- Reference bit
 - With each page associate a bit, initially = 0
 - When page is referenced bit set to 1
 - Replace the one which is 0 (if one exists). We do not know the order, however
- Second chance
 - Need reference bit
 - Clock replacement
 - If page to be replaced (in clock order) has reference bit = 1 then:
 - set reference bit 0
 - leave page in memory (i.e., give a second chance)
 - replace next page (in clock order), subject to same rules

Second-Chance (clock) Page-Replacement Algorithm



It degrades into the FIFO policy in the worst case

Enhanced Second-Chance Algorithm

(ref, dirty)

1. (0, 0) neither recently used nor modified—best page to replace
 2. (0, 1) not recently used but modified—not quite as good, because the page will need to be written out before replacement
 3. (1, 0) recently used but clean—probably will be used again soon
 4. (1, 1) recently used and modified—probably will be used again soon, and the page will be need to be written out to disk before it can be replaced
- The first page encountered at the lowest nonempty class is replaced
 - This method considers to reduce I/O costs.

Counting Algorithms

- Keep a counter of the number of references that have been made to each page
- **LFU** Algorithm: replaces page with smallest count
 - Periodically bit shift for counters (aging)
- **MFU** Algorithm: based on the argument that the page with the smallest count was probably just brought in and has yet to be used
 - Preserving pages newly brought in

Recency vs. Frequency

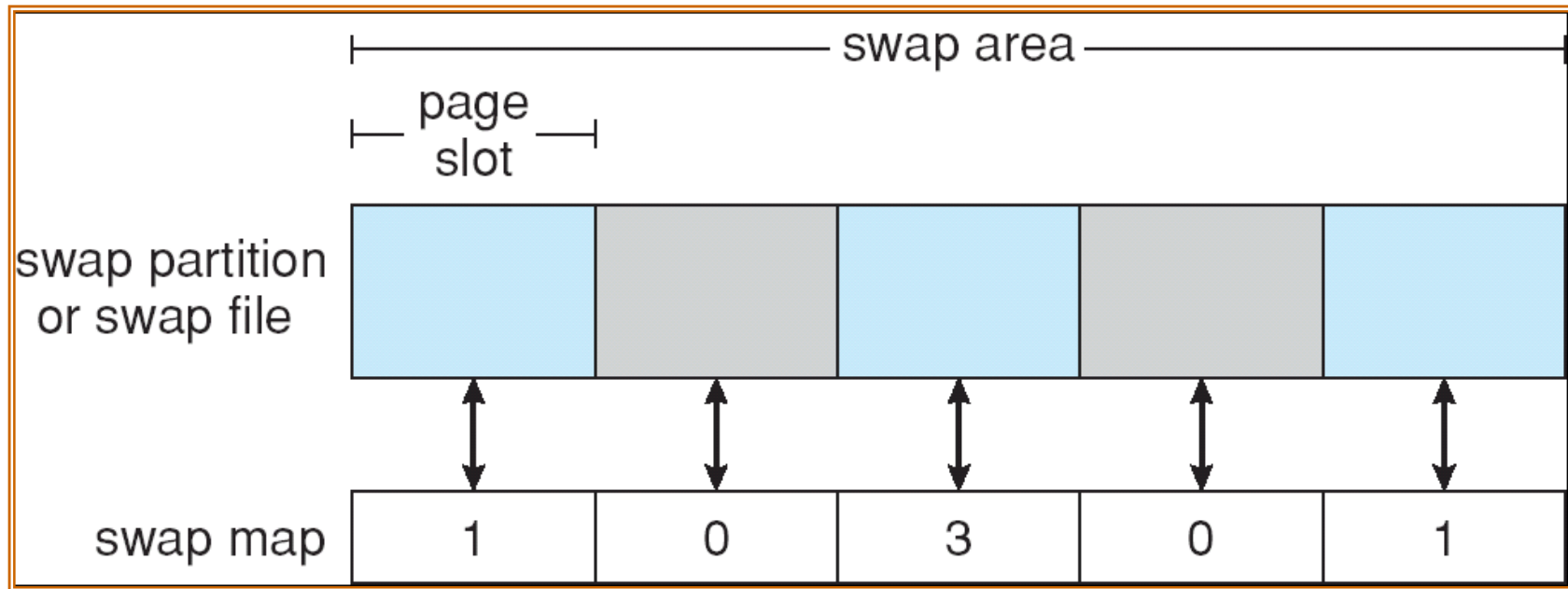
- LRU
 - Replace the least recently used page
 - Fast response to locality change
 - Sequential reference will wash away all cached pages
- LFU
 - Replace the least frequently used page
 - Resistant to sequential reference
 - Need time to warm up and cool down a page

Swap Space Management

Swap-Space Management

- Swap-space — Virtual memory uses disk space as an extension of main memory
- Swap-space can be carved out of the normal file system (Windows/Linux) or can be in a separate disk partition (Linux)
- Swap-space management
 - Kernel uses swap maps to track swap-space use
 - 4.3BSD allocates swap space when process starts; holds text segment (the program) and data segment
 - Solaris 2 allocates swap space only when a page is forced out of physical memory, not when the virtual memory page is first created

The Data Structures for Swapping on Linux Systems



Why there are entries > 1 ?

Applications of Demand Paging

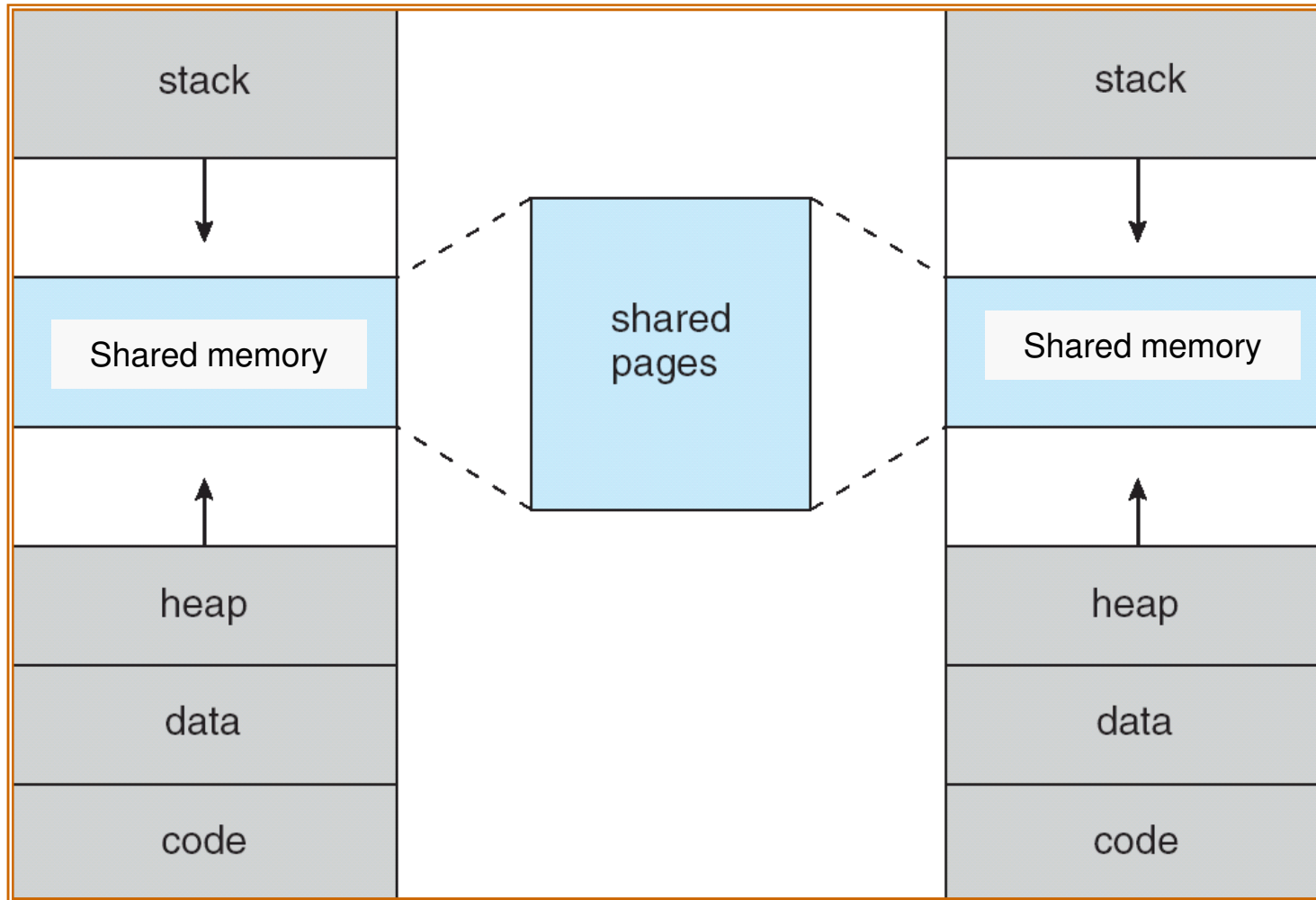
OS Features based on Demand Paging

- Page Sharing
 - Two processes share a memory region
 - Shared memory can be mapped to a file
- Copy-on-write
 - A fast implementation of memory duplication between processes, e.g., fast fork()
 - Pages are shared until being modified
- Memory-mapped file
 - Memory region backed by a regular file, not the swap space

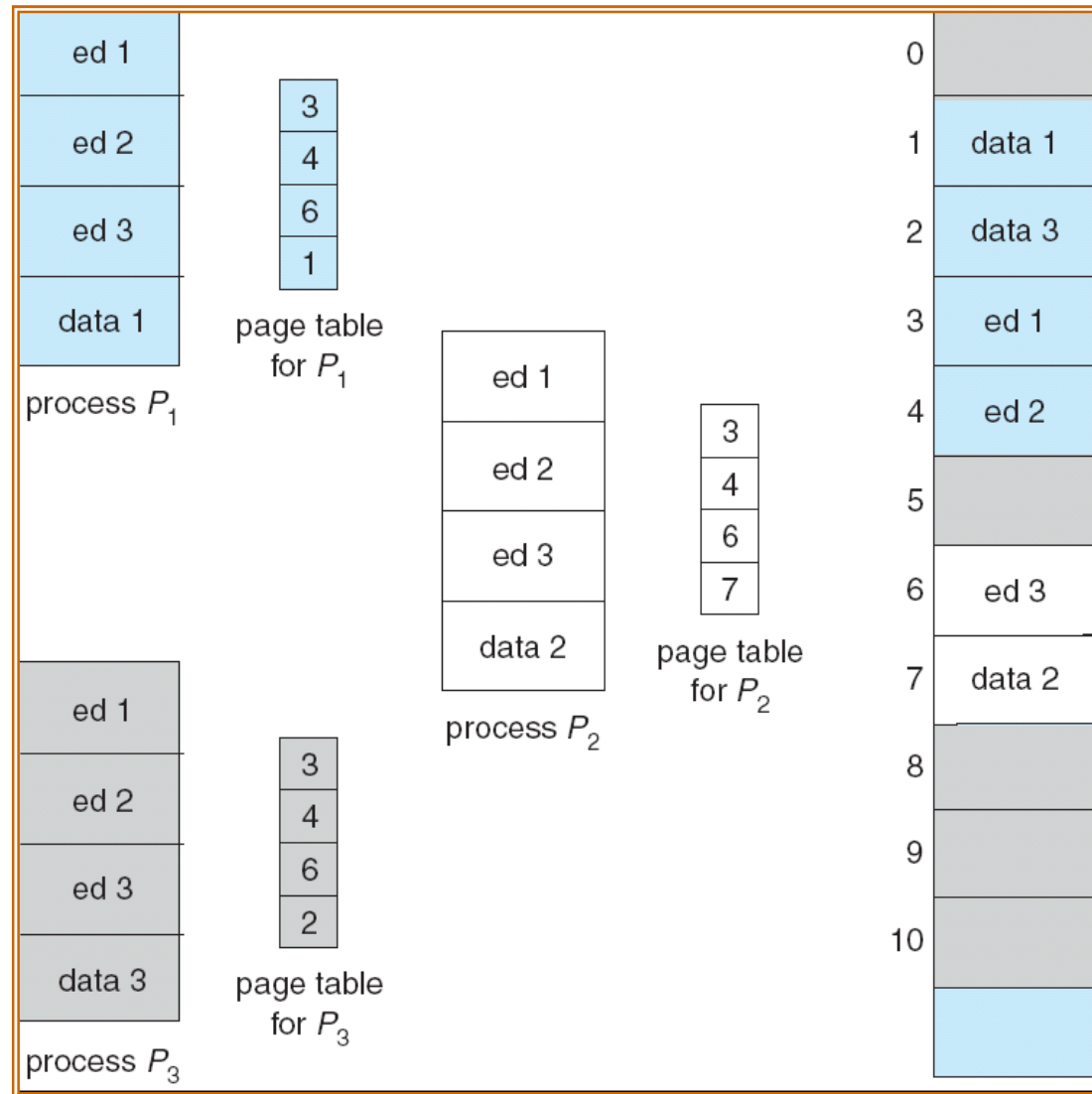
Page sharing

- Shared code
 - Dynamic linking-loading libraries (e.g., libc)
 - Kernel code
 - Handled by OS
- Shared memory
 - Creating a piece of shared memory: `shmget()`
 - Mapping a piece of shared memory to process address space: `shmat()`

Shared Memory

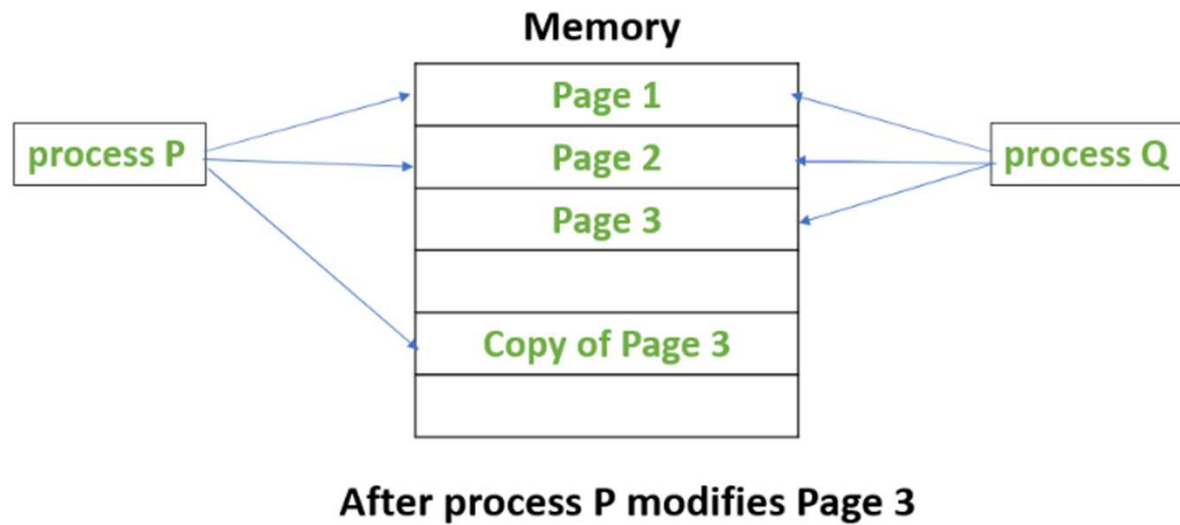
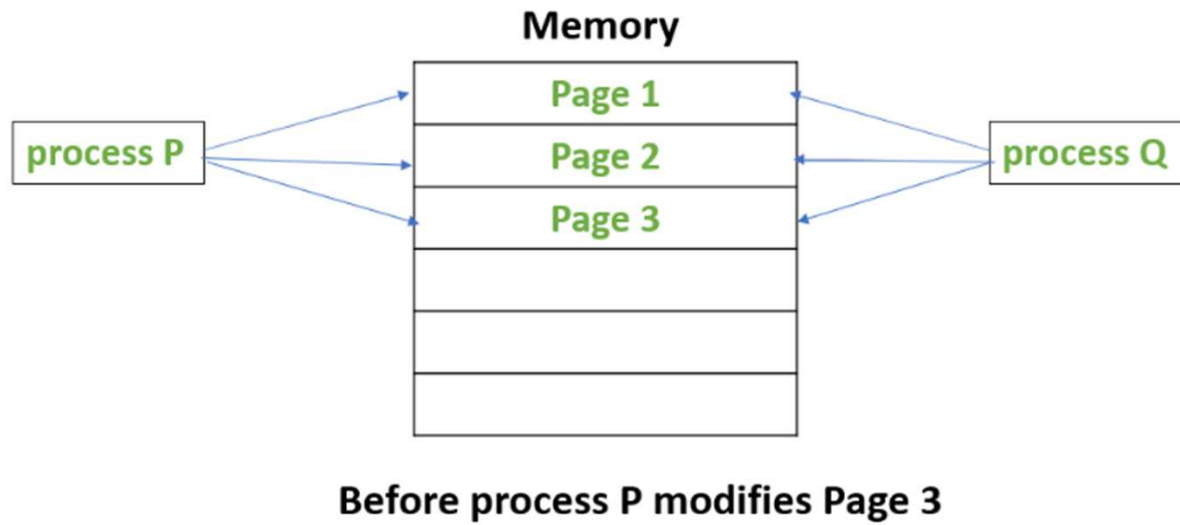


Shared Code



Copy-on-Write

- An extension to shared page; for fast duplication of memory
- Copy-on-Write (COW) allows both parent and child processes to initially share the same pages in memory
 - If either process modifies a shared page, a copy of the page is then created
- A RO/RW bit for each page is necessary to trap writes to shared page; write to the page cause an exception/trap
- COW allows more efficient process creation as only modified pages are copied



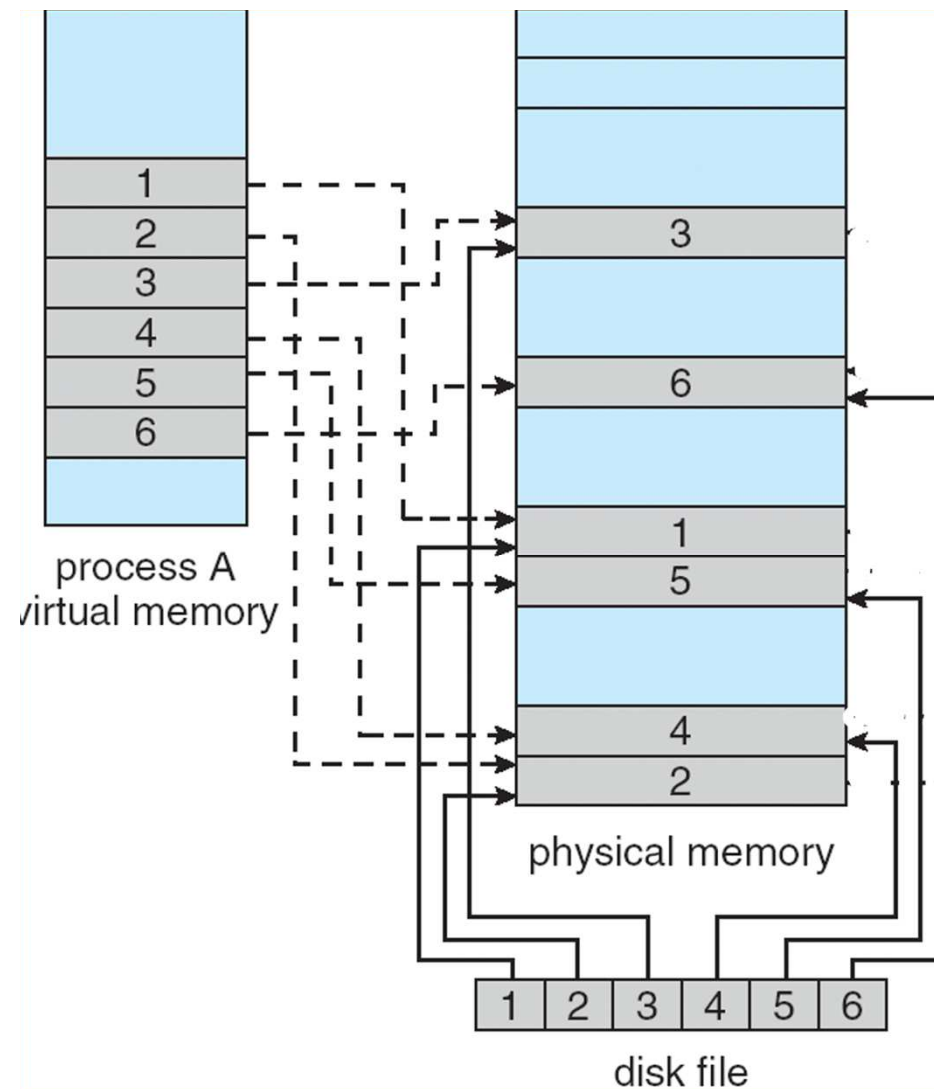
fork() and COW

- fork() uses copy-on-write
- vfork() does not use copy-on-write. It suspends the parent process and let the child process use the memory pages of the parent
 - The child uses on top of the parent's stack
 - The child is supposed to call exec() immediately
 - Child's leftovers on the stack is cleaned up upon exec()
- vfork() is for CPUs without a MMU. Otherwise, fork() with COW is efficient enough

Memory-Mapped Files

- A segment of virtual memory that is linearly mapped to a disk file
- Reading the memory segment triggers page faults, which brings a page-sized portion of the file through demand paging
- Writing the memory segments makes pages dirty; dirty pages are flushed to disk file in background or explicitly

Memory Mapped Files

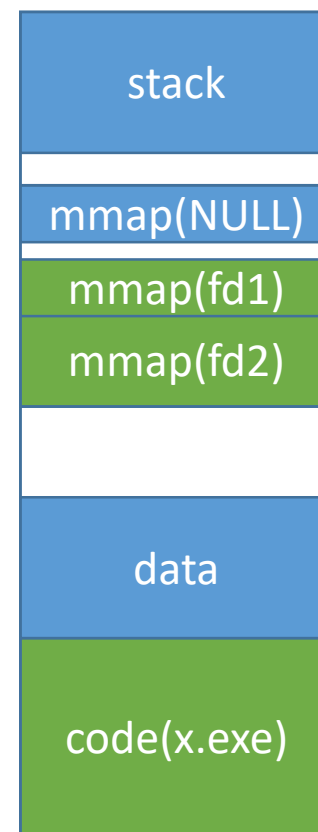


Memory-Mapped Files

- Available to user program through the `mmap()` system call
- Simplifies file access by treating file I/O through memory rather than `read()` `write()` system calls
 - Simple byte-level memory operations
 - No need to enter the kernel
- Shared memory can be mapped to a file as a memory-mapped file
- If the mapped file is `null`, pages go to the swap space
 - Application \rightarrow `malloc()` , glibc \rightarrow `mmap(null)`, system call

Mapping of Pages: File and Anonymous

- **Anonymous** pages belong to memory segments of processes, such as data, stack, and `mmap(null)` regions
- **File** pages belong to memory regions that are mapped to a (named) file
- Anonymous pages are backed by the swap device; file pages are backed by files
- They share the same mechanism for paging in and paging out but with different swap-out destinations



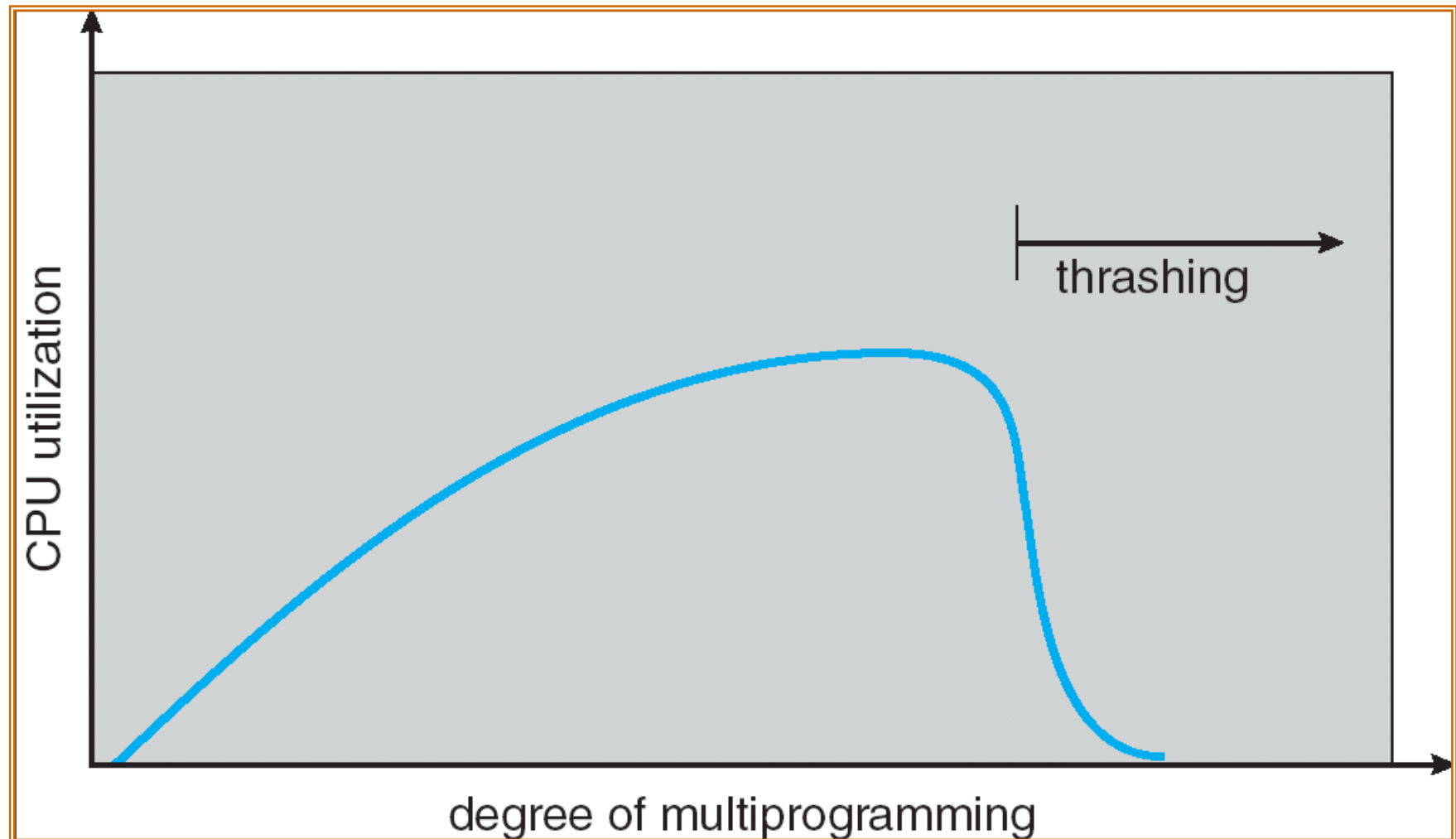
- What type of pages the following memory regions are associated with?
- Code from executable files
- Process stacks
- Process heaps
- Memory shared between processes

PERFORMANCE ISSUES

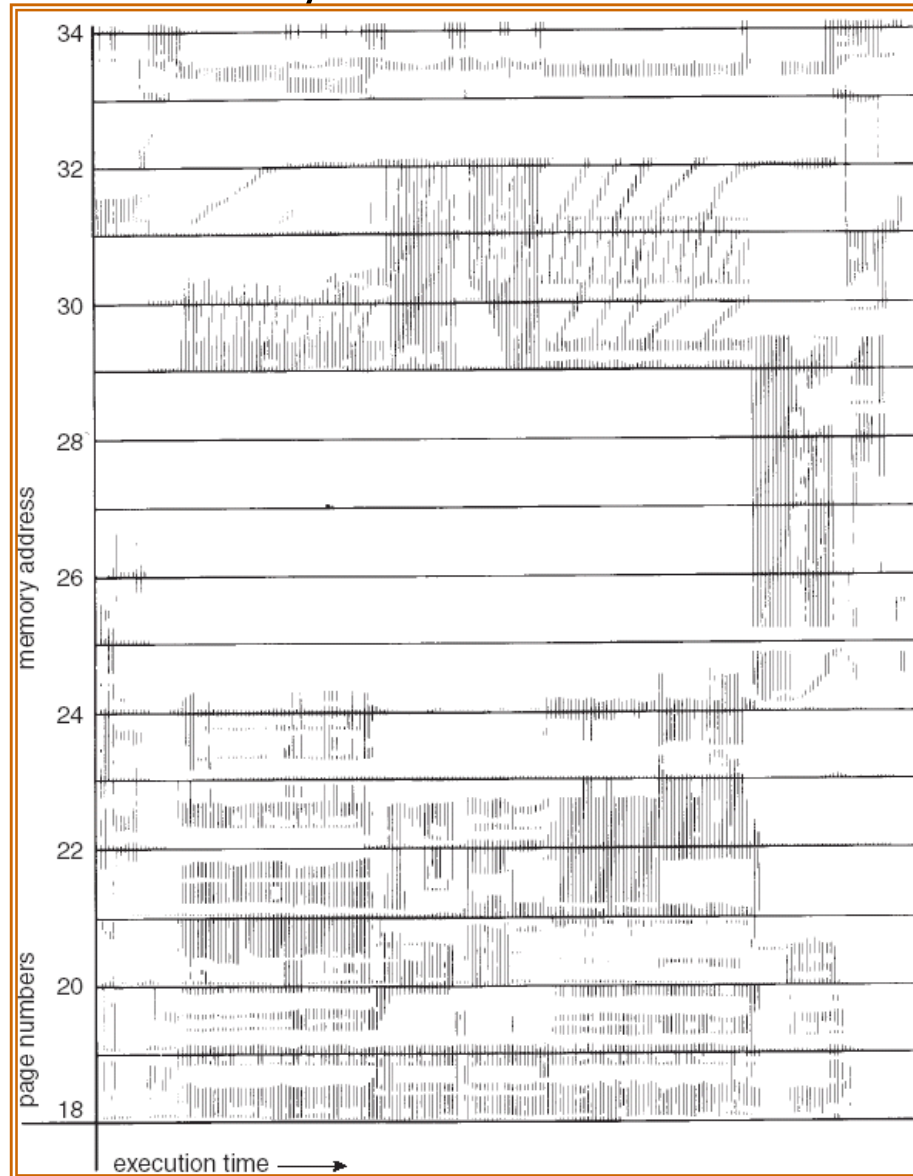
Thrashing

- If a process does not have “enough” pages, the page-fault rate will be very high. This leads to:
 - a low CPU utilization
 - operating system thinks that it needs to increase the degree of multiprogramming
 - another process is added to the system
- **Thrashing** \equiv a state that the kernel is busy swapping pages in and out

Thrashing (Cont.)



Locality In A Memory-Reference Pattern



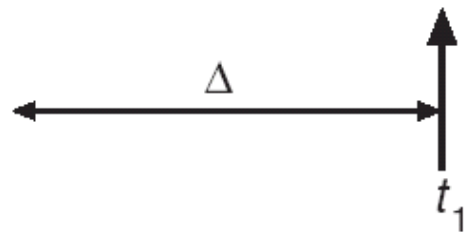
Working-Set Model

- $\Delta \equiv$ working-set window \equiv a fixed number of page references. Example: 10,000 references
- WSS_i (working set of Process P_i) = a set of pages referenced in the most recent Δ
 - if Δ too small will not encompass the entire locality
 - if Δ too large will unnecessarily encompass old localities
- $D = \sum WSS_i \equiv$ total demand frames; m is the total number of frames
- if $D > m \Rightarrow$ Thrashing: De-allocate pages from processes

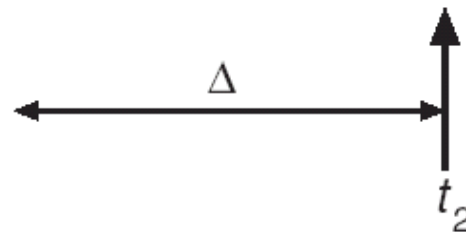
Working-set model

page reference table

. . . 2 6 1 5 7 7 7 7 5 1 6 2 3 4 1 2 3 4 4 4 4 3 4 3 4 4 4 4 1 3 2 3 4 4 4 3 4 4 4 . . .



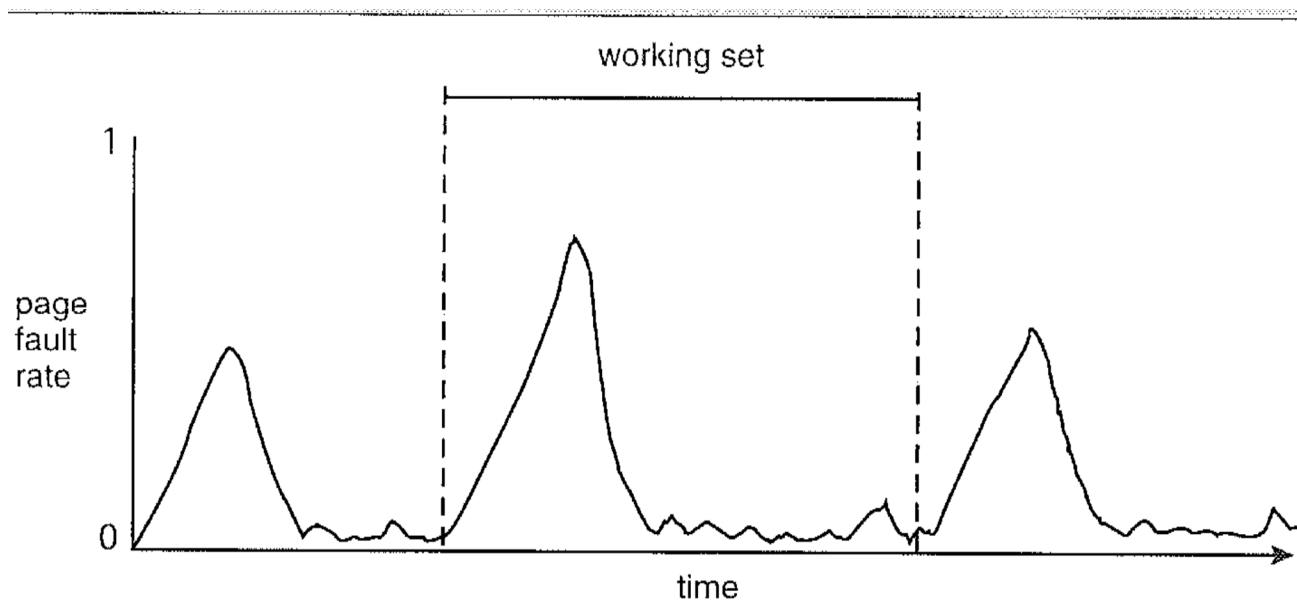
$$WS(t_1) = \{1, 2, 5, 6, 7\}$$



$$WS(t_2) = \{3, 4\}$$

Remark: Migration of Working Sets

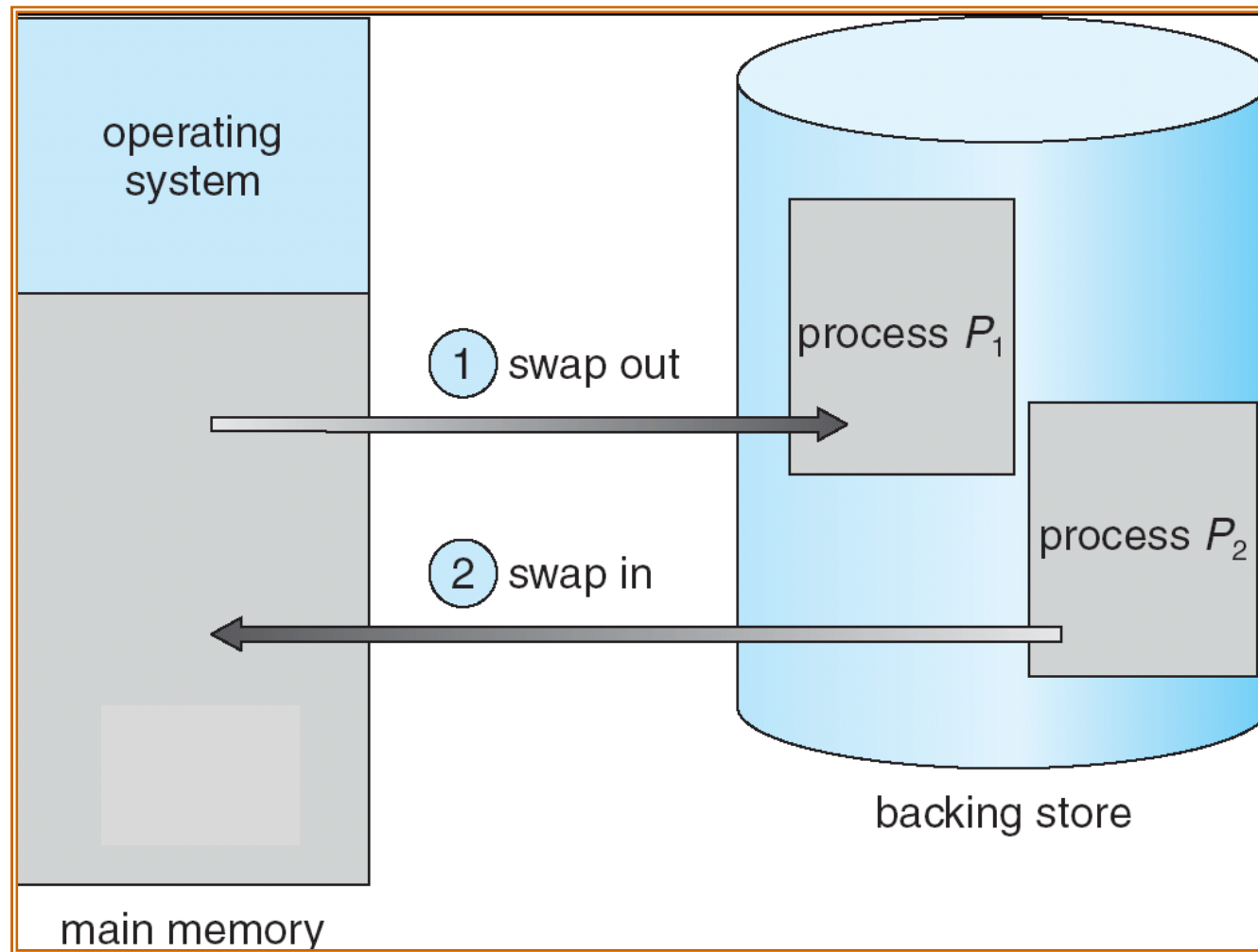
- Working set is a time-variant, when a process migrates from a working set to another, the page fault rate also increases (but not thrashing)



Thrashing Management

- Swapper
 - AKA Midterm scheduler; process-based swapping
 - A process can be swapped temporarily out of memory to a backing store, and then brought back into memory for continued execution
 - Memory of a swapped-out process are released; useful to thrashing management

Schematic View of Process Swapping



Thrashing Management

- Modern UNIX OSes, however, do not implement process swapping and use finer-granularity page swapping instead
- The kernel cannot reclaim free pages quickly enough when the system is under thrashing
- Process termination: Linux uses Out-of-Memory (OOM) Killer as the last resort

Background Dirty Page Flushing

- Reclaiming a clean page is very fast (just discard it), but reclaiming a dirty page involves a slow I/O (page write back)
- The OS flushes dirty pages in **background**
 - Making dirty pages clean
 - Detaching the writing latency from page replacement
- `kswapd` and `pdflush` in Linux

Pre-paging (Pre-fetching)

- On a page fault, read a number of pages ahead of the requested page
- Pros 1: Less disk head movement
 - Sequential disk access is highly efficient
- Pros 2: Fewer page faults
 - Spatial locality
 - Fetch a few sequential pages ahead
 - 128 KB=32 pages in Linux
- Cons: If pre-paged pages are not used, I/O and memory were wasted

Program Structure

- Program structure
 - `Int[128,128] data;`
 - Each row is stored in one page
 - Program 1

```
for (j = 0; j < 128; j++)  
  for (i = 0; i < 128; i++)  
    data[i,j] = 0;
```

128 x 128 = 16,384 page faults

- Program 2

```
for (i = 0; i < 128; i++)  
  for (j = 0; j < 128; j++)  
    data[i,j] = 0;
```

128 page faults

Page size=128 integers
Suppose that we have < 128 frames...

TLB Reach & Large Pages

- TLB Reach = (TLB Size) X (Page Size)
 - The amount of memory accessible without page table lookup
- Ideally, the working set of each process is stored in the TLB. Otherwise there is a high degree of TLB miss
- Increase the TLB size
 - Impractical !! Too expensive
- Increase the Page Size
 - May have negative performance impacts: internal fragmentation, more I/O traffic, etc.
- Provide Multiple Page Sizes
 - For a good balance between I/O traffic volume and TLB hit ratio
 - IA-64 supports 4 KB and 4 MB pages

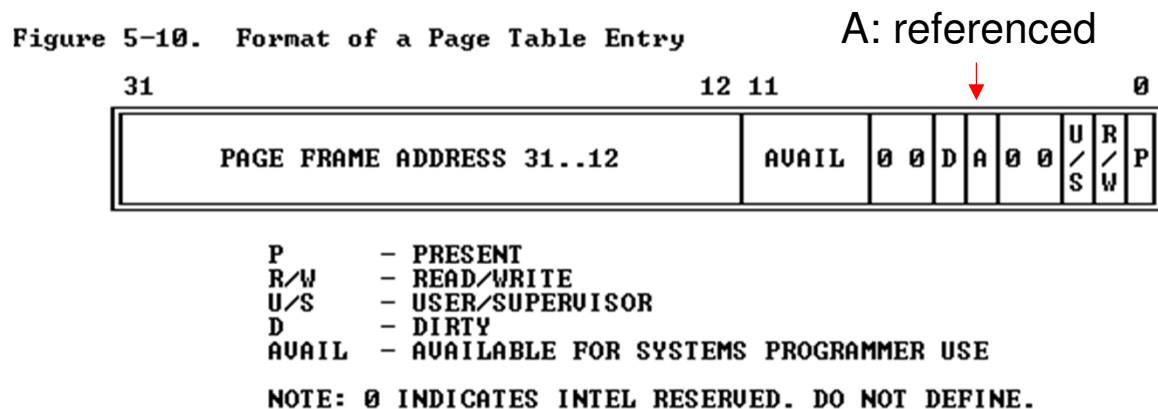
Page Size Tradeoff

- Large pages
 - Small page table (good)
 - Large TLB reach (good)
 - May bring unused data into memory (bad)
- Small pages
 - Large page table (bad)
 - Small TLB reach (bad)
 - Less unused data in memory (good)

Operating System Examples

Status Bits Associated with a Page (x86)

- Valid (Present) bit: does the page present in main memory?
- RW bit: is the page read-only or read-write?
- Reference bit: is the page referenced?
- Dirty bit: is a page modified?



Linux Page Reclaiming

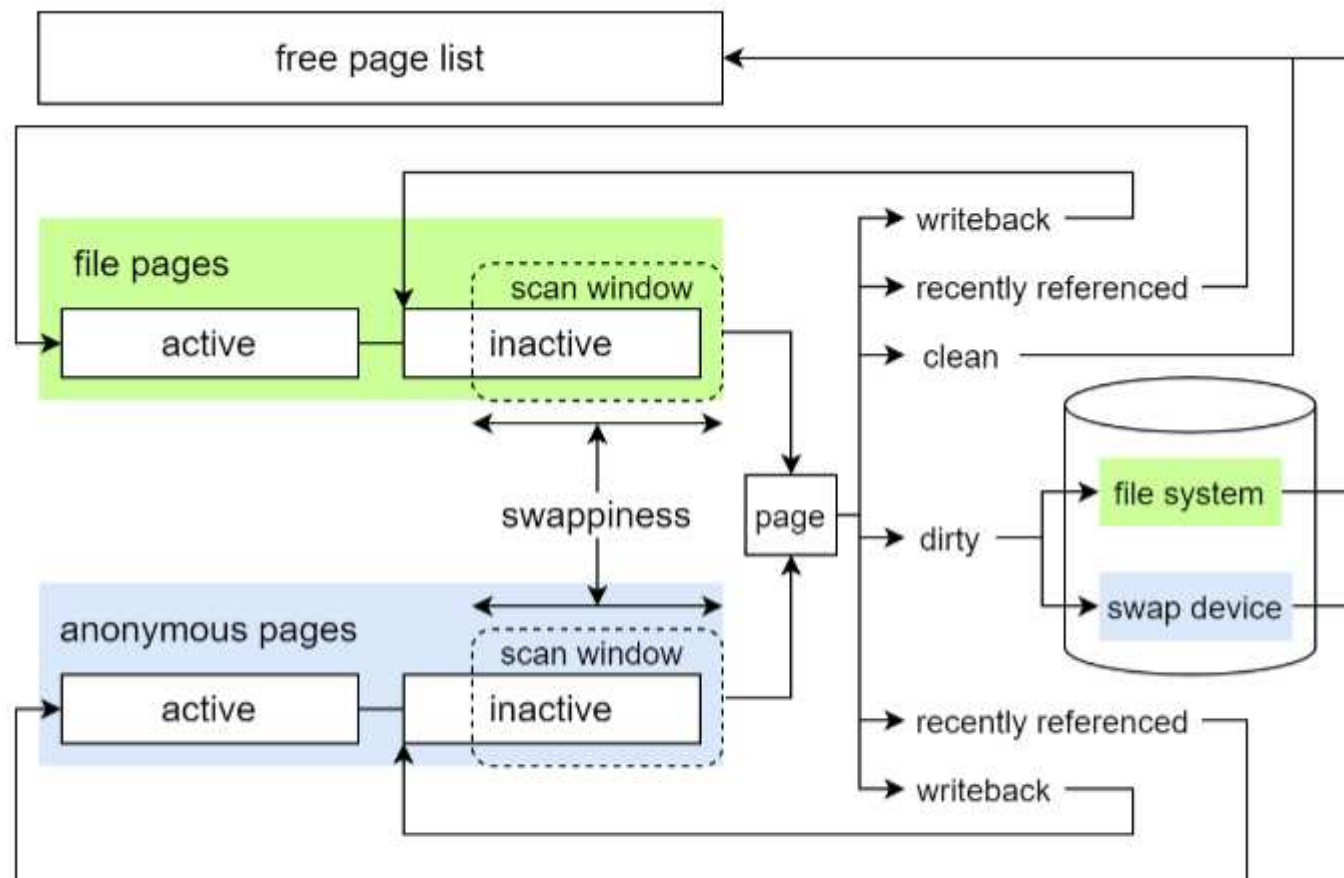


Figure A.1: An overview of page scanning for reclaiming free memory.

-王詠萱，壓縮交換:不以快閃記憶體壽命為代價之安卓裝置使用體驗改善方法，陽明交大碩士論文，2021

-Yong-Xuan Wang, Chung-Hsuan Tsai and Li-Pin Chang, "Killing Processes or Killing Flash? Escaping from the Dilemma Using Lightweight, Compression-Aware Swap for Mobile Devices," International Conference on Embedded Software (EMSOFT 2021)

Check file pages and anon pages of process

- Run “ps” to find out the ID of a process of yours
- `cat /proc/[pid]/status | grep Rss`
 - In general, a process use more file pages than anon pages; that's why swappiness is set to 60 by default
($\text{reclaim prob.} = (200 - 60) / 200$)

```
[lpchang@linux1 3253000]$ ps
  PID TTY          TIME CMD
3253000 pts/52    00:00:00 tcsh
3255286 pts/52    00:00:00 ps
[lpchang@linux1 3253000]$ cat /proc/3253000/status |grep Rss
RssAnon:                960 kB
RssFile:                 3172 kB
RssShmem:                 0 kB
[lpchang@linux1 3253000]$
```

End of Chapter 9