Computer Security Capstone

Chapter 22: Internet Security Protocols and Standards

Chi-Yu Li (2023 Spring)
Computer Science Department
National Yang Ming Chiao Tung University

Outline

- Secure E-mail
- DomainKeys Identified Mail
- Secure Sockets Layer (SSL) and Transport Layer Security (TLS)
- HTTPS (HTTP over SSL)
- IPv4 and IPv6 Security

Secure E-mail

MIME

- An Internet mail format: an extension to the old RFC 822 specification
 - RFC 822 defines a simple heading with To, From, Subject
 - Assumes ASCII text format
- Provides new header fields
 - Defining information about the message body

S/MIME

- Secure/Multipurpose Internet Mail Extension
- Security enhancement to the MIME
 Internet e-mail format
 - Based on technology from RSA Data Security
- Provides the ability to sign and/or encrypt e-mail messages

MIME and S/MIME Message Examples

```
From: John Doe <example@example.com>
MTME-Version: 1.0
Content-Type: multipart/mixed;
                                           MIME
        boundary="XXXXboundary text"
This is a multipart message in MIME format.
--XXXXboundary text
Content-Type: text/plain
this is the body text
--XXXXboundary text
Content-Type: text/plain;
Content-Disposition: attachment;
        filename="test.txt"
this is the attachment text
--XXXXboundary text--
```

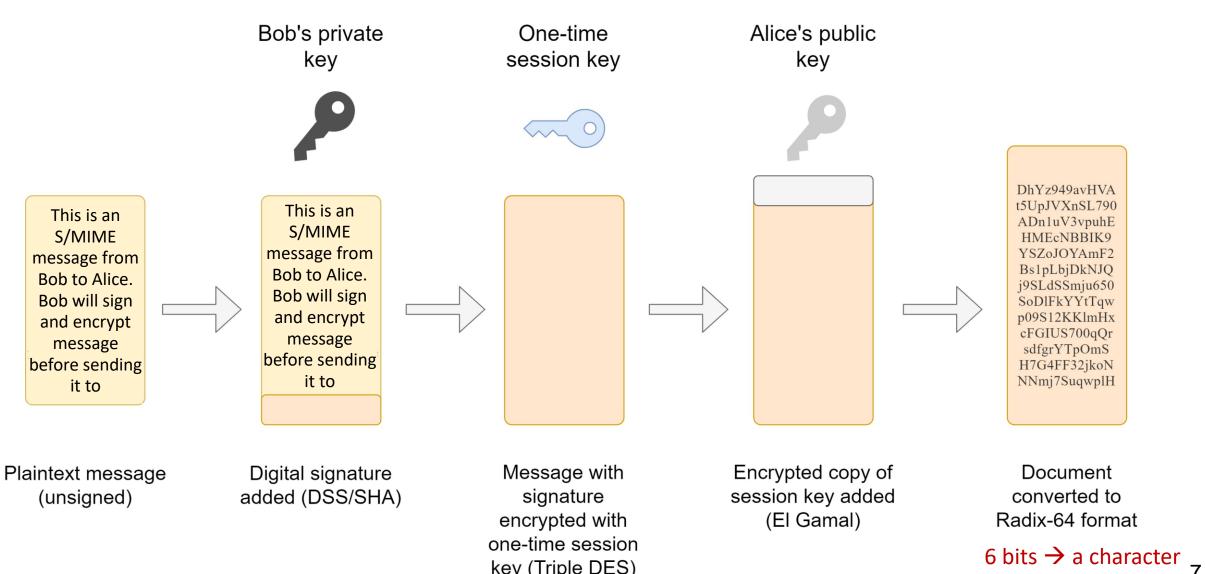
```
Content-Type: multipart/signed;
protocol="application/pkcs7-signature"; micalg="shal";
boundary="---Alexis"
                                                   S/MIME
-----Alexis
Content-Type: text/plain
This is an example of a clear signed message.
-----Alexis
Content-Type: application/pkcs7-signature; name="smime.p7s"
Content-Transfer-Encoding: base64
Content-Disposition: attachment; filename="smime.p7s"
MIHmBgkqhkiG9w0BBwKggdgwgdUCAQMxCTAHBgUrDgMCGjALBgkqhkiG9w0BBwEx
gbcwgbQCAQOAFFJz90NbzoEOUQOJxop1+azsEdH9MAcGBSsOAwIaMA0GCSqGSIb3
DQEBAQUABIGAmO1mvk4maqMgb+aYomDTO5fPnRK/0p8w4ExJnai+MknjHmo5VX99
J0zxf+myPuBn1ajZKxoDIRSHpeMp8oToAKUgWNQm7b/yAqRmtCQvmaeI7rnMULVU
cfQfsI0Hz/ujJ02wUUMjFKmLgo7Q/C7++JTJTad9+SADRcNnVEWGjMg=
-----Alexis--
```

Computer Security Capstone, Spr	Туре	Subtype	Description
MIME Content Types	Text	Plain	Unformatted text; may be ASCII or ISO 8859.
		Enriched	Provides greater format flexibility.
	Multipart	Mixed	The different parts are independent but are to be transmitted together. They should be presented to the receiver in the order that they appear in the mail message.
		Parallel	Differs from Mixed only in that no order is defined for delivering the parts to the receiver.
		Alternative	The different parts are alternative versions of the same information. They are ordered in increasing faithfulness to the original, and the recipient's mail system should display the "best" version to the user.
		Digest	Similar to Mixed, but the default type/subtype of each part is message/rfc822.
	Message	Rfc822	The body is itself an encapsulated message that conforms to RFC 822.
		Partial	Used to allow fragmentation of large mail items, in a way that is transparent to the recipient.
		External-body	Contains a pointer to an object that exists elsewhere.
	Image	Jpeg	The image is in JPEG format, JFIF encoding.
		gif	The image is in GIF format.
	Video	Mpeg	MPEG format.
	Audio	Basic	Single-channel 8-bit ISDN mu-law encoding at a sample rate of 8 kHz.
	Application	PostScript	Adobe Postscript.
		octet-stream	General binary data consisting of 8-bit bytes.

S/MIME Content Types

Туре	Subtype	smime Parameter	Description
Multipart	Signed		A clear-signed message in two parts: one is the message and the other is the signature.
Application	pkcs7-mime	signedData	A signed S/MIME entity.
	pkcs7-mime	envelopedData	An encrypted S/MIME entity.
	pkcs7-mime	degenerate signedData	An entity containing only public-key certificates.
	pkcs7-mime	CompressedData	A compressed S/MIME entity.
	pkcs7- signature	signedData	The content type of the signature subpart of a multipart/signed message.

Typical S/MIME Process for Creating an S/MIME Message



key (Triple DES)

S/MIME Functions

Enveloped data

Encrypted content and associated keys

Signed data

Encoded
message +
signed
digest

Clear-signed data

Cleartext
message +
encoded
signed
digest

Signed and enveloped data

Nesting of signed and encrypted entities

Enveloped Data Using Public-key Infrastructure

- Default algorithms used for encrypting S/MIME messages are AES and RSA
 - $\square M \rightarrow AES(M) \rightarrow AES(M) + RSA(K)$
 - M: message
 - K: a pseudorandom secret key
 - Use recipient's public RSA key to encrypt K
- Public-key certificate: basic tool permits widespread use of S/MIME
 - □ Using certificates that conform to the international standard X.509v3

Signed and Clear-signed Data

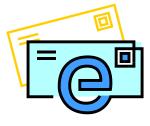
- Signed data → Base64(content + sig)
- Clear-signed data → content + Base64(sig)
- Default algorithms used for signing messages
 - □ DSS(SHA-1(Message), DSS-private-key)
 - □ DSS: Digital Signature Standard; SHA: Secure Hash Algorithm
- Alternative
 - □ RSA(SHA-1/MD5(Message), RSA-private-key)
- Radix-64/Base64: mapping content/sig into printable ASCII characters

Outline

- Secure E-mail
- DomainKeys Identified Mail
- Secure Sockets Layer (SSL) and Transport Layer Security (TLS)
- HTTPS (HTTP over SSL)
- IPv4 and IPv6 Security

DomainKeys Identified Mail (DKIM)

- A specification of cryptographically signing e-mail messages
 - ☐ A signing domain can claim responsibility for a message
 - □ RFC 4871 standard: DKIM signatures
- Has been widely adopted by a range of e-mail providers
 - □ e.g., gmail, yahoo, and many ISPs
- Why DKIM?



Why DKIM?

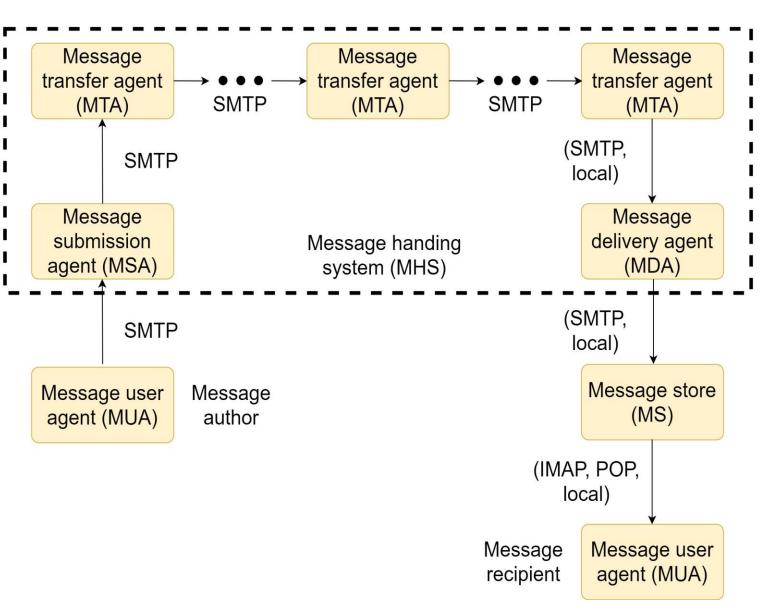
An email authentication technique that is transparent to the end user

- Reasons
 - □ S/MIME depends on both the sending and receiving users employing S/MIME
 - But, the bulk of incoming mail does not use S/MIME
 - □ S/MIME signs only the message content
 - Header information concerning origin can be compromised
 - ☐ Transparent to the user; (s)he need take no action
 - Applied to all mails from cooperating domains
 - □ Preventing forgers from masquerading as good senders

Internet Mail Architecture

User world: MUA

- Transfer world: MHS
 - □ Composed of MTAs



Simple Example of DKIM Deployment

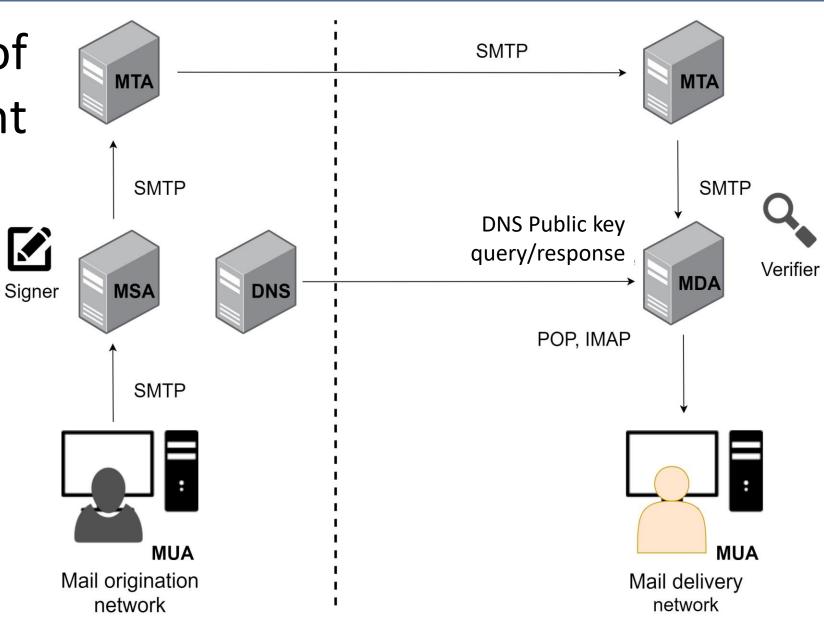
DNS = domain name system

MDA = mail delivery agent

MSA = mail submission agent

MTA = message transfer agent

MUA = message user agent



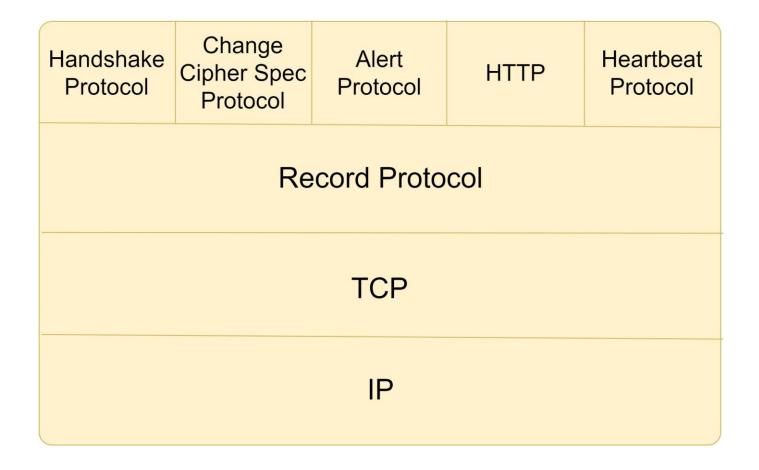
Outline

- Secure E-mail
- DomainKeys Identified Mail
- Secure Sockets Layer (SSL) and Transport Layer Security (TLS)
- HTTPS (HTTP over SSL)
- IPv4 and IPv6 Security

SSL and TLS

- Secure Socket Layer (SSL)
 - ☐ One of the most widely used security services
 - SSL3.0, SSL2.0 (deprecated in 2015)
- Transport Layer Security (TLS)
 - Becoming Internet standard RFC4346
 - ☐ General-purpose service: as a set of protocols that rely on TCP
 - □ TLS1.3, TLS1.2, TLS1.1 (new standards)
 - ☐ Two implementation choices
 - As part of the underlying protocol suite: transparent to apps
 - Be embedded in specific packages: e.g., most browsers come equipped with SSL

SSL/TLS Protocol Stack



TLS Concepts

TLS Connection

- A transport (in the OSI layering model definition) that provides a suitable type of service
- Peer-to-peer relationships
- Transient
- Every connection is associated with one session

TLS Session

- An association between a client and a server
- Created by the handshake protocol
- Define a set of cryptographic security parameters
- Used to avoid the expensive negotiation of new security parameters for each connection

TLS Record Protocol Operation

Two services

□ Confidentiality

■ Message integrity

Application data

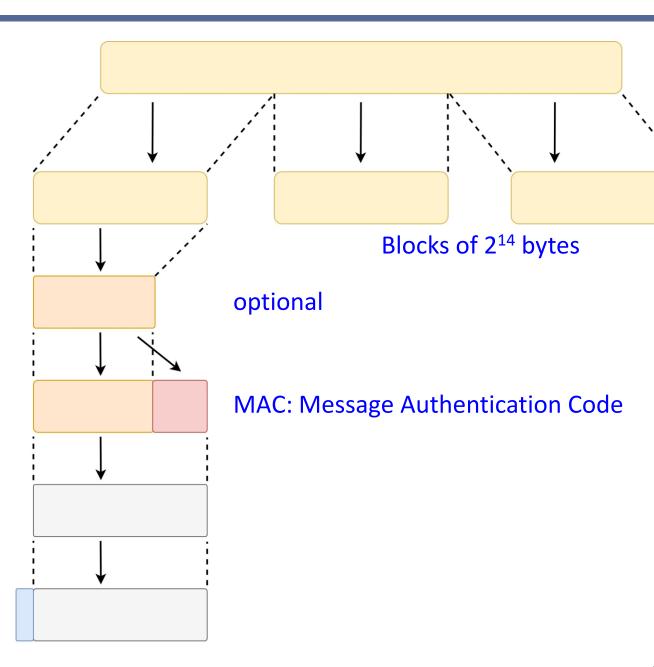
Fragment

Compress

Add MAC

Encrypt

Append TLS record header



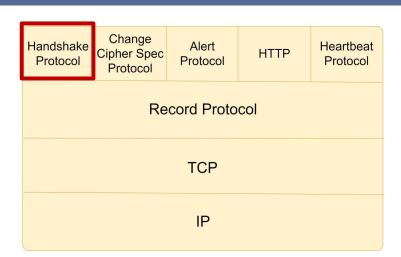
Handshake Protocol



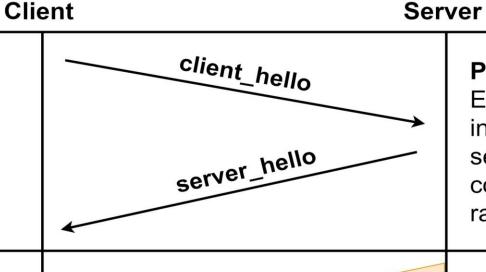
- Most complex part of TLS
- Used before any application data are transmitted
- Allows server and client to:



- Comprises a series of messages exchanged by client and server
- Exchange has four phases

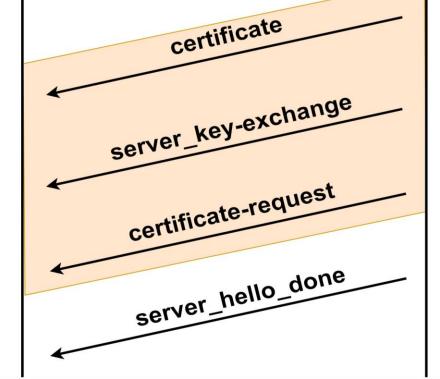


Handshake Protocol Action I



Phase 1

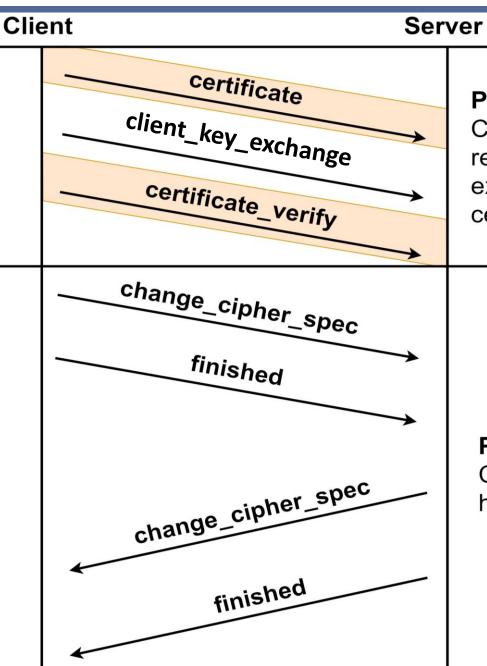
Establish security capabilities, including protocol version, session ID, cipher suite, compression method, and initial random number.



Phase 2

Server may send certificate, key exchange, and request certificate. Server signals end of hello message phase.

Handshake Protocol Action II



Phase 3

Client sends certificate if requested. Client sends key exchange. Client may send certificate verification.

Phase 4

Change cipher suite and finish handshake protocol.

ClientHello (RFC)

```
CipherSuite TLS RSA WITH NULL MD5
                                                      = { 0x00,0x01 };
CipherSuite TLS_RSA_WITH_NULL_SHA
                                                        { 0x00,0x02 };
CipherSuite TLS RSA WITH NULL SHA256
                                                        { 0x00,0x3B };
CipherSuite TLS_RSA_WITH_RC4_128_MD5
                                                          0x00,0x04 }:
CipherSuite TLS RSA WITH RC4 128 SHA
                                                        { 0x00,0x05 }:
CipherSuite TLS RSA WITH 3DES EDE CBC SHA
                                                      = { 0x00,0x0A };
CipherSuite TLS RSA WITH AES 128 CBC SHA
                                                        { 0x00,0x2F };
CipherSuite TLS_RSA_WITH_AES_256_CBC_SHA
                                                        { 0x00,0x35 };
CipherSuite TLS RSA WITH AES 128 CBC SHA256
                                                      = \{ 0 \times 00, 0 \times 3C \};
CipherSuite TLS RSA WITH AES 256 CBC SHA256
                                                      = \{ 0 \times 00.0 \times 3D \}:
```

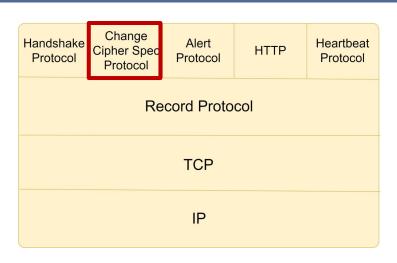
- struct {
 - ProtocolVersion client_version;
 - □ Random random;
 - SessionID session_id;
 - □ CipherSuite cipher suites <2..2^16-2>;
 - □ CompressionMethod compression_methods <1..2^8-1>;} ClientHello;

ServerHello (RFC)

struct {
 ProtocolVersion server_version;
 Random random;
 SessionID session_id;
 CipherSuite cipher_suite;
 CompressionMethod compression_method;} ServerHello;

Change Cipher Spec Protocol

- One of four TLS specific protocols that use the TLS Record Protocol
 - ☐ the simplest

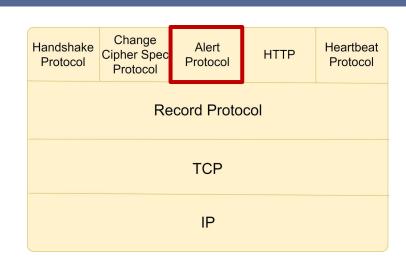


Consists of a single message: a single byte with the value 1

- Purpose: cause pending state to be copied into the current state
 - ☐ Hence, updating the cipher suite in use

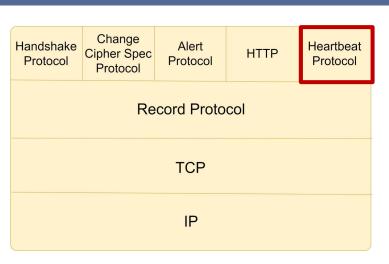
Alert Protocol

- Two bytes
 - ☐ First byte: warning (1), fatal (2)
 - For fatal, TLS implementation **terminates** the TLS connection.
 - For other TLS connections using the same TLS session, they may continue, but no new TLS connections may be established
 - ☐ Second byte: specific alert code (RFC5246-appendix)
 - close_notify(0), unexpected_message(10), etc.



Heartbeat Protocol

- A periodic signal generated by hardware or software
 - □ Indicates normal operation or synchronizes other parts of a system
 - ☐ Typically used to monitor the availability of a protocol entity
 - ☐ Timer: 1s to 60s (RFC6347)
 - ☐ Its Use is established during Phase 1 of the Handshake Protocol

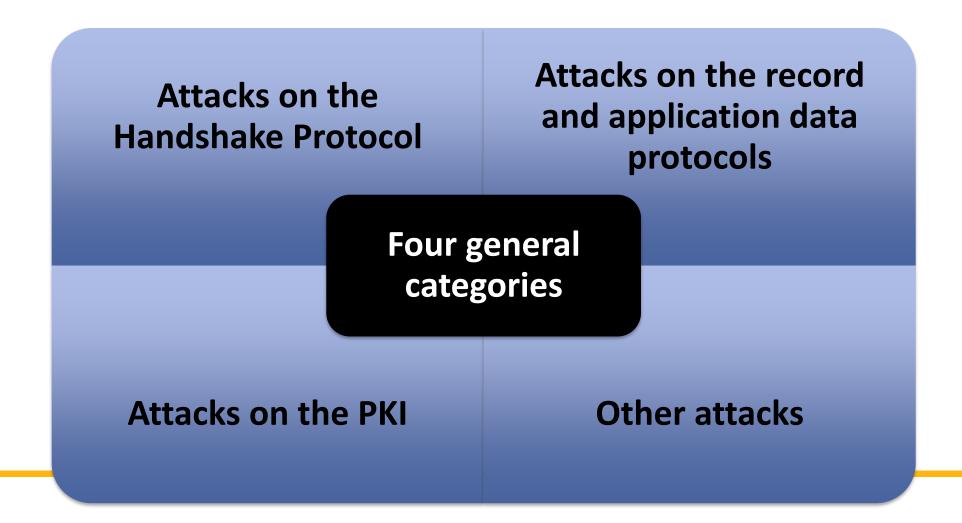


Serves two purposes:

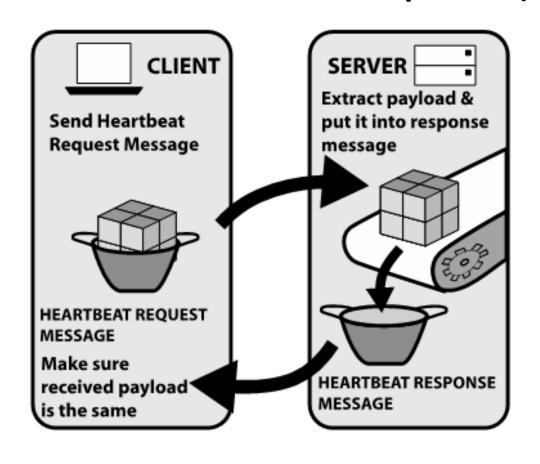
- Assures the sender that the recipient is still alive
- ☐ Generates activity across the connection during idle periods

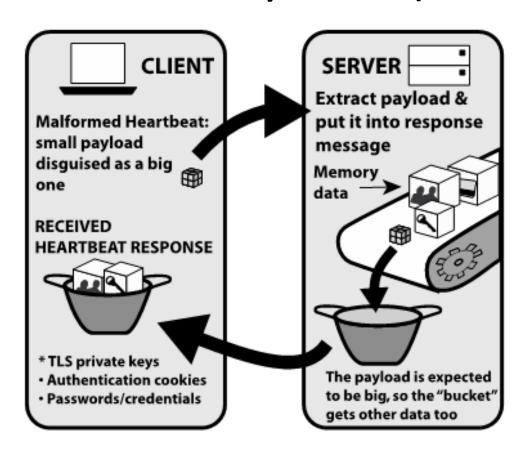
```
struct {
   HeartbeatMessageType type;
   uint16 payload_length;
   opaque payload[HeartbeatMessage.payload_length];
   opaque padding[padding_length];
} HeartbeatMessage;
```

SSL/TLS Attacks



The Heartbleed Exploit (Source: BAE Systems)





(a) How TLS Heartbeat Protocol works

(b) How TLS Heartbleed exploit works

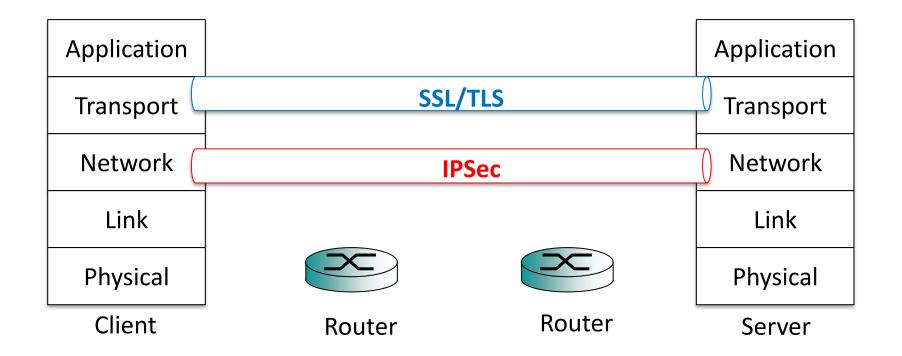
Outline

- Secure E-mail
- DomainKeys Identified Mail
- Secure Sockets Layer (SSL) and Transport Layer Security (TLS)
- HTTPS (HTTP over SSL)
- IPv4 and IPv6 Security

HTTPS (HTTP over SSL/TLS)

- Combination of HTTP and SSL
 - ☐ HTTP over TLS (RFC2818)
 - □ Secure communication between a Web browser and a Web server
- Built into all modern Web browsers
 - □ URL addresses begin with https://
- Connection initiation and closure
 - ☐ HTTP client act as TLS client
 - □ TLS handshake → HTTP request
 - ☐ Three levels: HTTP, TLS session, TCP

IPSec v.s. SSL/TLS



Applications of IPSec

- Secure branch office connectivity over the Internet
- Secure remote access over the Internet

- Establishing extranet and intranet connectivity with partners
- Enhancing electronic commerce security

Benefits of IPSec

- Strong security applied to all traffic at a firewall or router
- Resistant to bypass at a firewall

- Transparent to apps
 - □ No need to change software
- Transparent to end users
 - No need to train users on security mechanisms

Benefits of IPSec (Cont.)

- Routing apps: prevent attackers from disrupting communications or diverting some traffic
 - □ A router advertisement from an authorized router
 - A neighbor advertisement from an authorized router
 - ☐ A redirect message from the router to which the initial packet was sent
 - A routing update is nor forged

The Scope of IPSec

- Two main functions
 - Encapsulating Security Payload (ESP): a combined authentication/encryption function
 - □ A key exchange function: Internet Key Exchange standard (IKEv2)
- VPN: both authentication and encryption are generally desired
- Authentication Header (AH): authentication-only function (deprecated)

Security Associations

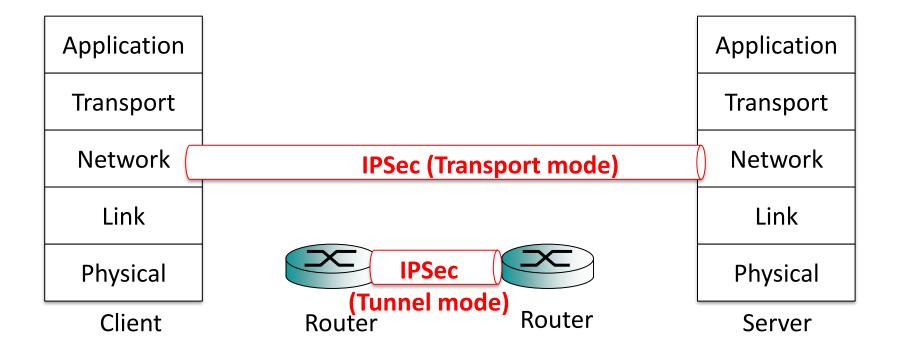
- A key concept of IPSec
 - One-way relationship between a sender and a receiver
 - Two-way secure exchange: two SAs are required
- Uniquely identified by three parameters
 - Security parameter index (SPI)
 - □ IP destination address
 - □ Protocol identifier: AH or ESP

Security Associations (Cont.)

- Characterized by the following parameters
 - □ Sequence number counter: 32-bit
 - □ Sequence counter overflow: A flag → whether overflow → an auditable event
 - ☐ Anti-replay window: defining a sliding window
 - □ AH information
 - Algorithm, keys, key lifetimes, etc.
 - **□** ESP information
 - Algorithm, keys, init values, key lifetimes, etc.
 - ☐ Lifetime of this security association
 - ☐ IPSec protocol mode: tunnel or transport
 - Path MTU

Two IPSec Operation Modes

Transport and Tunnel modes



Encapsulating Security Payload

 Providing authentication and confidentiality services Bit: 24 31 **Security Parameters Index (SPI) Authentication Coverage Sequence Number Transport-level segment (transport mode)** IP packet (tunnel mode) Sonfidentiality Payload Data (variable) Soverage Encrypted data Padding(0 - 255 bytes) Data integrity **Next Header Pad Length** protection **Authentication Data (variable)** 41

Transport and Tunnel Modes

Transport Mode

- Protection: the payload of an IP packet
- Typically used for end-to-end communication between two hosts
- ESP protects the IP payload but not the IP header

Tunnel Mode

- Protection: the entire IP packet
- Entire original packet travels through a tunnel from one point to another
- Used when one or both ends of a security association are a security gateway
- Hosts on networks behind firewalls may engage in secure communications without implementing IPSec

IPSec: AH + ESP

IP AH only

ID Hoodoro	ders AH	TCP/UDP	TCP/UDP
IP Headers		Headers	Payload

- IP AH + ESP
 - □ Transport mode

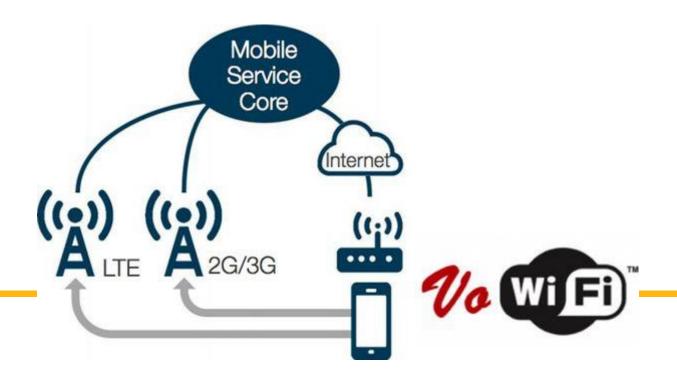
IP Headers	ESP Headers	Encrypted TCP/UDP Headers	Encrypted TCP/UDP Payload	ESP Trailer	ESP Auth (AH)
------------	----------------	---------------------------------	---------------------------------	-------------	------------------

□ Tunnel mode

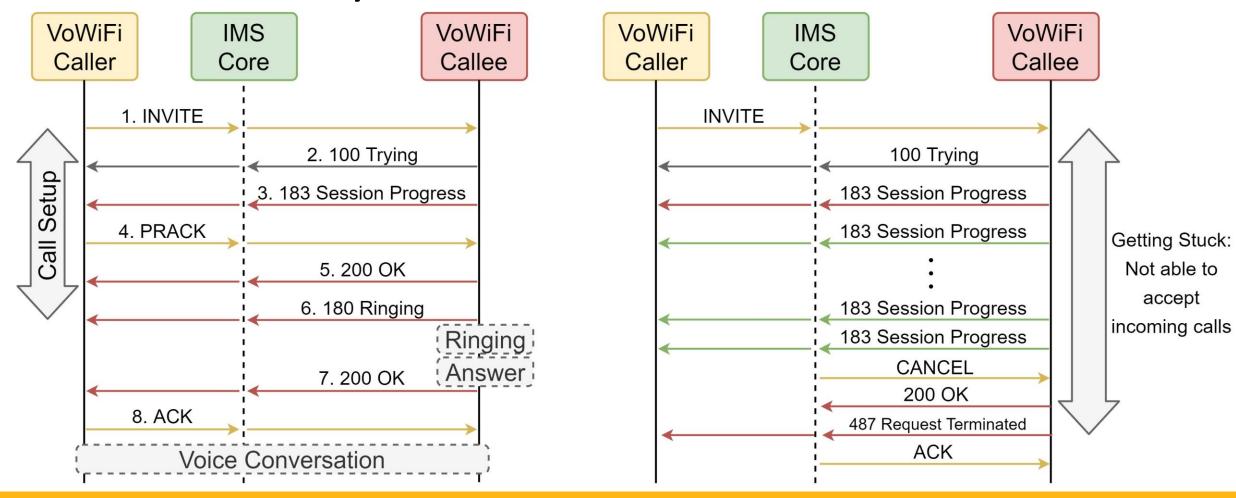
|--|

Case Study: VoWi-Fi Session Hijacking

- VoWi-Fi (Voice over Wi-Fi): telephony calls over Wi-Fi
 - □ Complements VolTE (Vocie over LTE) for poor-signal areas
 - ☐ Supported by a system called IMS (IP Multimedia Subsystem)

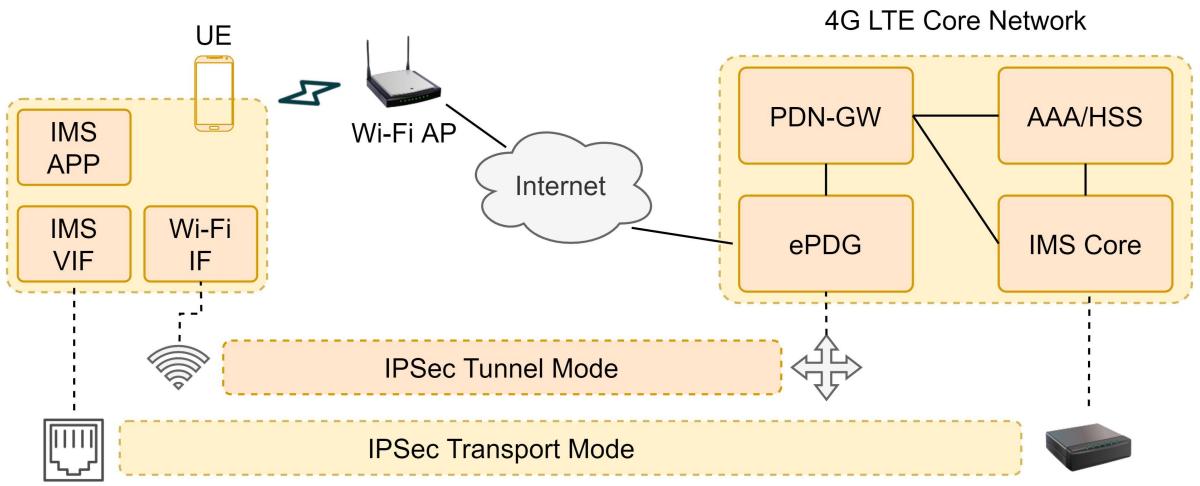


A Vulnerability at the Call State Machine



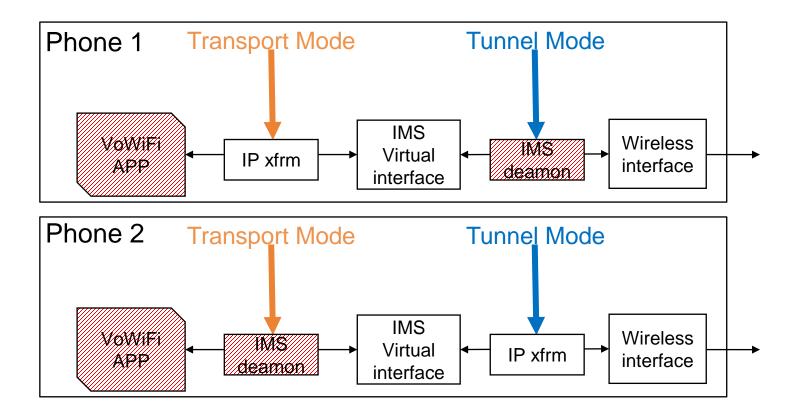
VoWi-Fi SIP Sessions Protected by IPSec

• How to hijack the session to exploit the vulnerability?



47

IPSec Implementation Analysis



 XFRM ("transform"): The Linux kernel's IP framework for transforming packets (such as encrypting payloads)

□ Support IPSec

Example: XFRM IPSec SAs

Phone 1

```
dreamlte:/ # ip xfrm state
src 10.159.204.230 dst 100.64.45.207
        proto esp spi 0x00017502 regid 512 mode transport
        replay-window 4
        auth-trunc hmac(sha1) 0x35a5dccbf4cfa5f2c54999b6216d4cec 96
        enc ecb(cipher null)
        anti-replay context: seq 0x5, oseq 0x0, bitmap 0x0000001f
        sel src 0.0.0.0/0 dst 0.0.0.0/0
src 100.64.45.207 dst 10.159.204.230
        proto esp spi 0x77b513cc regid 511 mode transport
        replay-window 4
        auth-trunc hmac(sha1) 0x35a5dccbf4cfa5f2c54999b6216d4cec 96
        enc ecb(cipher null)
        anti-replay context: seq 0x0, oseq 0x0, bitmap 0x00000000
        sel src 0.0.0.0/0 dst 0.0.0.0/0
src 10.159.204.230 dst 100.64.45.207
        proto esp spi 0x00017501 regid 510 mode transport
        replay-window 4
        auth-trunc hmac(sha1) 0x35a5dccbf4cfa5f2c54999b6216d4cec 96
        enc ecb(cipher null)
        anti-replay context: seq 0x24, oseq 0x0, bitmap 0xffffffff
        sel src 0.0.0.0/0 dst 0.0.0.0/0
src 100.64.45.207 dst 10.159.204.230
        proto esp spi 0x80cfc9f0 regid 509 mode transport
        replay-window 4
        auth-trunc hmac(sha1) 0x35a5dccbf4cfa5f2c54999b6216d4cec 96
        enc ecb(cipher null)
        anti-replay context: seq 0x0, oseq 0x33, bitmap 0x00000000
        sel src 0.0.0.0/0 dst 0.0.0.0/0
```

Example: XFRM IPSec SAs

Phone 2

```
ASUS_Z01KDA:/ # ip xfrm state
src 223.22.236.1 dst 192.168.0.121
    proto esp spi 0x1f35c02b reqid 0 mode tunnel
    replay-window 0 flag af-unspec
    auth-trunc hmac(shal) 0xc9d7d3aa2f57e7b0cf2216d795478a04d7efa010 96
    enc cbc(des3_ede) 0x43098b1d8fae1344671ba8e61e955c0fc5ee8706c6ac8cbb
    encap type espinudp sport 4500 dport 42987 addr 0.0.0.0
    anti-replay context: seq 0x0, oseq 0x0, bitmap 0x00000000
src 192.168.0.121 dst 223.22.236.1
    proto esp spi 0xef011233 reqid 0 mode tunnel
    replay-window 0 flag af-unspec
    auth-trunc hmac(shal) 0x02027e6811161a69e9b16088d89c5dbd5cd58583 96
    enc cbc(des3_ede) 0x961028b4d0a30ebeadaa031918dd5f623204a30d3f3fbd52
    encap type espinudp sport 42987 dport 4500 addr 0.0.0.0
    anti-replay context: seq 0x0, oseq 0x65, bitmap 0x00000000
```

IPSec SAs Negotiation

```
Source
                 Destination
                                  Protocol
                                          Length
                                                  Info
                                              557 Request: REGISTER sip:ims.mnc092.
100.72.58.140
                 10.159.204.229
                                  SIP
10.159.204.229 100.72.58.140
                                              894 Status: 401 Unauthorized 11030230
                                 SIP

▼ Session Initiation Protocol (REGISTER)
   Request-Line: REGISTER sip:ims.mnc092.mcc466.3gppnetwork.org SIP/2.0

▼ Message Header

      Via: SIP/2.0/TCP 100.72.58.140:5060; branch=z9hG4bK-524287-1---0fe6d123f0b
        Require: sec-agree
         [truncated]Security-Client: ipsec-3gpp;prot=esp;mod=trans;spi-c=58645;sp.
Source
                Destination
                                 Protocol
                                         Length
                                             557 Request: REGISTER sip:ims.mnc092.m
100.72.58.140
                 10.159.204.229
                                  SIP
                                             894 Status: 401 Unauthorized 110302303
10.159.204.229 100.72.58.140
                                  SIP

▼ Session Initiation Protocol (401)
   > Status-Line: SIP/2.0 401 Unauthorized 11030230325

▼ Message Header

     > Security-Server: ipsec-3gpp; alg=hmac-sha-1-96; prot=esp; mod=trans; ealg=null
```

IPSec Transport Packet

```
Destination
Source
                             Protocol
                                    Length
                                          Info
1224 Request: INVITE sip:+
                                                                    023292
                                       464 Status: 100 Trying
10.156.204.230 100.64.54.225 SIP
                                       212 Status: 183 Session Progress
10.156.204.230 100.64.54.225 SIP/SDP
                                      1096 Request: PRACK sip:sgc_c@10.156.20
100.64.54.225 10.156.204.230 SIP
                                       720 Status: 200 OK
10.156.204.230 100.64.54.225 SIP
                                       840 Status: 180 Ringing
10.156.204.230 100.64.54.225 SIP
 Encapsulating Security Payload
    ESP SPI: 0xc0448453 (3225715795)
    ESP Sequence: 9
    ESP Pad Length: 1
    Next header: TCP (0x06)
    Authentication Data: 273cc5e6fa6ef23184580d6a

▼ Session Initiation Protocol (INVITE)
  > Request-Line: INVITE sip:+
23292; phone-context=ims.mnc
.mcc
.3g
  > Message Header

▼ Message Body

    Session Description Protocol
        Session Description Protocol Version (v): 0
       > Owner/Creator, Session Id (o): 24465697594557 1535544072365129 0 IN
        Session Name (s): FORGED SIP
```

How to Fabricate a IPSec-protected SIP Message?

- SIP message content
 - □ caller and callee numbers
 - ☐ Static info: SIP server/proxy addresses and supported features, etc.
 - □ Dynamic info: tag, call id, branch id, etc.
- ESP integrity protection
 - □ Obtain the IPSec security parameters from the SAD (Security Association DB)
 - SPI and hash algorithm/keys
 - ☐ Keep track of the sequence number of IPSec session
 - Apply default ESP padding algorithm
 - ☐ Generate authentication data using the hash algorithm/keys

Questions?