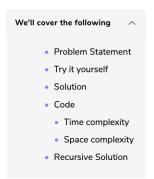


Permutations (medium)



Problem Statement

Given a set of distinct numbers, find all of its permutations.

Permutation is defined as the re-arranging of the elements of the set. For example, {1, 2, 3} has the following six permutations:

```
1. {1, 2, 3}
2. {1, 3, 2}
3. {2, 1, 3}
4. {2, 3, 1}
5. {3, 1, 2}
```

6. {3, 2, 1}

If a set has 'n' distinct elements it will have n! permutations.

Example 1:

```
Input: [1,3,5]
Output: [1,3,5], [1,5,3], [3,1,5], [3,5,1], [5,1,3], [5,3,1]
```

Try it yourself

Try solving this question here:

```
⊗ C++
   class Permutations {
     public static List<List<Integer>> findPermutations(int[] nums) {
       List<List<Integer>> result = new ArrayList<>();
       return result;
     public static void main(String[] args) {
      List<List<Integer>> result = Permutations.findPermutations(new int[] { 1, 3, 5 });
       System.out.print("Here are all the permutations: " + result);
15 }
Run
                                                                                     Save Reset []
```

Solution

This problem follows the Subsets pattern and we can follow a similar Breadth First Search (BFS) approach. However, unlike Subsets, every permutation must contain all the numbers.

Let's take the example-1 mentioned above to generate all the permutations. Following a BFS approach, we will consider one number at a time:

1. If the given set is empty then we have only an empty permutation set: []

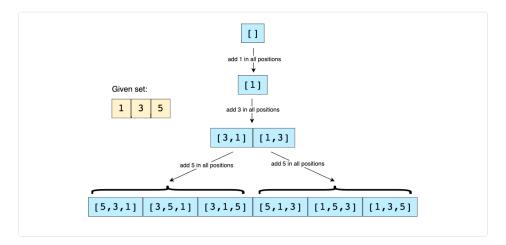
- 2. Let's add the first element (1), the permutations will be: [1]
- 3. Let's add the second element (3), the permutations will be: [3,1], [1,3]
- 4. Let's add the third element (5), the permutations will be: [5,3,1], [3,5,1], [3,1,5], [5,1,3], [1,5,3], [1,3,5]

Let's analyze the permutations in the 3rd and 4th step. How can we generate permutations in the 4th step from the permutations of the 3rd step?

If we look closely, we will realize that when we add a new number (5), we take each permutation of the previous step and insert the new number in every position to generate the new permutations. For example, inserting '5' in different positions of [3,1] will give us the following permutations:

- 1. Inserting '5' before '3': [5,3,1]
- 2. Inserting '5' between '3' and '1': [3,5,1]
- 3. Inserting '5' after '1': [3,1,5]

Here is the visual representation of this algorithm:



Code

Here is what our algorithm will look like:

```
🏂 Java
                        G C++
                                    Js JS
    class Permutations {
      public static List<List<Integer>> findPermutations(int[] nums) {
        List<List<Integer>> result = new ArrayList<>();
        Queue<List<Integer>>> permutations = new LinkedList<>();
        permutations.add(new ArrayList<>());
         for (int currentNumber : nums) {
           int n = permutations.size();
            List<Integer> oldPermutation = permutations.poll();
             for (int j = 0; j <= oldPermutation.size(); j++) {</pre>
              List<Integer> newPermutation = new ArrayList<Integer>(oldPermutation);
              newPermutation.add(j, currentNumber);
              if (newPermutation.size() == nums.length)
                result.add(newPermutation);
                permutations.add(newPermutation);
         return result:
             static void main(String[] args)
Run
                                                                                        Save Reset []
```

Time complexity

We know that there are a total of N! permutations of a set with 'N' numbers. In the algorithm above, we are iterating through all of these permutations with the help of the two 'for' loops. In each iteration, we go through all the current permutations to insert a new number in them on line 17 (line 23 for C++ solution). To insert a number into a permutation of size 'N' will take O(N), which makes the overall time complexity of our algorithm O(N*N!).

Space complexity

All the additional space used by our algorithm is for the result list and the queue to store the intermediate permutations. If you see closely, at any time, we don't have more than N! permutations between the result list and the queue. Therefore the overall space complexity to store N! permutations each containing N elements will be O(N*N!).

Recursive Solution

Here is the recursive algorithm following a similar approach:

