

Level Order Successor (easy)

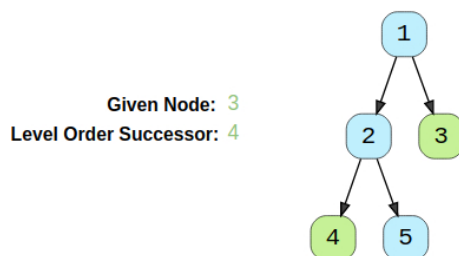
We'll cover the following ^

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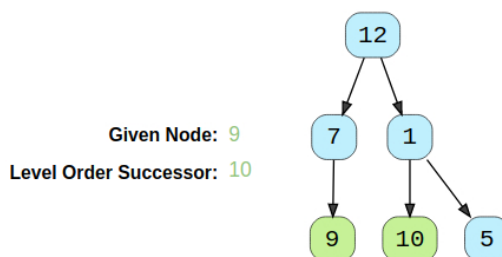
Problem Statement

Given a binary tree and a node, find the level order successor of the given node in the tree. The level order successor is the node that appears right after the given node in the level order traversal.

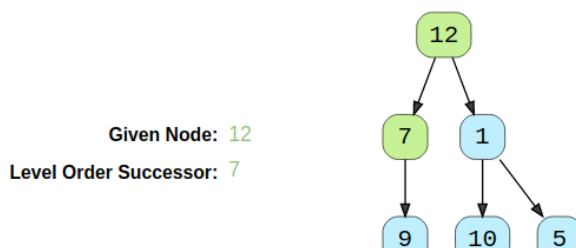
Example 1:



Example 2:



Example 3:



Try it yourself

Try solving this question here:

Java Python3 JS C++

```
1 from collections import deque
2
3
4 class TreeNode:
5     def __init__(self, val):
6         self.val = val
7         self.left, self.right = None, None
8
9
10 def find_successor(root, key):
11     # TODO: Write your code here
12     return None
13
14
15 def main():
16     root = TreeNode(12)
17     root.left = TreeNode(7)
18     root.right = TreeNode(1)
19     root.left.left = TreeNode(9)
20     root.right.left = TreeNode(10)
21     root.right.right = TreeNode(5)
22     result = find_successor(root, 12)
23     if result:
24         print(result.val)
25     result = find_successor(root, 9)
26     if result:
27         print(result.val)
28
29
30 main()
31
```

Run Save Reset

Solution

This problem follows the [Binary Tree Level Order Traversal](#) pattern. We can follow the same **BFS** approach. The only difference will be that we will not keep track of all the levels. Instead we will keep inserting child nodes to the queue. As soon as we find the given node, we will return the next node from the queue as the level order successor.

Code

Here is what our algorithm will look like; most of the changes are in the highlighted lines:

Java Python3 C++ JS

```
1 from collections import deque
2
3
4 class TreeNode:
5     def __init__(self, val):
6         self.val = val
7         self.left, self.right = None, None
8
9
10 def find_successor(root, key):
11     if root is None:
12         return None
13
14     queue = deque()
15     queue.append(root)
16     while queue:
17         currentNode = queue.popleft()
18         # insert the children of current node in the queue
19         if currentNode.left:
20             queue.append(currentNode.left)
21         if currentNode.right:
22             queue.append(currentNode.right)
23
24     return queue[0] if queue else None
25
26
27 def main():
28     root = TreeNode(12)
29     root.left = TreeNode(7)
30     root.right = TreeNode(1)
31     root.left.left = TreeNode(9)
32     root.right.left = TreeNode(10)
33     root.right.right = TreeNode(5)
34     result = find_successor(root, 12)
35     if result:
36         print(result.val)
37     result = find_successor(root, 9)
38     if result:
39         print(result.val)
40
41
42 main()
43
```

```
20     queue.append(currentNode.left)
21     if currentNode.right:
22         queue.append(currentNode.right)
23
24     # break if we have found the key
25     if currentNode.val == key:
26         break
27
28     return queue[0] if queue else None
29
30
31 def main():
32     root = TreeNode(12)
33     root.left = TreeNode(7)
34     root.right = TreeNode(1)
35     root.left.left = TreeNode(9)
36     root.right.left = TreeNode(10)
37     root.right.right = TreeNode(5)
38     result = find_successor(root, 12)
39     if result:
40         print(result.val)
41     result = find_successor(root, 9)
42     if result:
43         print(result.val)
44
45
46 main()
47
```

Run

Save

Reset

Time complexity

The time complexity of the above algorithm is $O(N)$, where 'N' is the total number of nodes in the tree. This is due to the fact that we traverse each node once.

Space complexity

The space complexity of the above algorithm will be $O(N)$ which is required for the queue. Since we can have a maximum of $N/2$ nodes at any level (this could happen only at the lowest level), therefore we will need $O(N)$ space to store them in the queue.

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✓ Completed

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