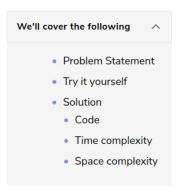




Subarrays with Product Less than a Target (medium)



Problem Statement

Given an array with positive numbers and a target number, find all of its contiguous subarrays whose **product is less than the target number**.

Example 1:

```
Input: [2, 5, 3, 10], target=30
Output: [2], [5], [2, 5], [3], [5, 3], [10]
Explanation: There are six contiguous subarrays whose product is less than the target.
```

Example 2:

```
Input: [8, 2, 6, 5], target=50
Output: [8], [2], [8, 2], [6], [2, 6], [5], [6, 5]
Explanation: There are seven contiguous subarrays whose product is less than the target.
```

Try it yourself

Try solving this question here:



Solution

This problem follows the **Sliding Window** and the **Two Pointers** pattern and shares similarities with Triplets with Smaller Sum with two differences:

- 1. In this problem, the input array is not sorted.
- 2. Instead of finding triplets with sum less than a target, we need to find all subarrays having a product less than the target.

The implementation will be quite similar to Triplets with Smaller Sum.

Code 7

Here is what our algorithm will look like:

```
Js JS
           Python3
                         ⊚ C++
👙 Java
      rom collections import deque
    def find subarrays(arr, target):
      result = []
      product = 1
      for right in range(len(arr)):
        product *= arr[right]
        while (product >= target and left < len(arr)):</pre>
          left += 1
        temp list = deque()
        for i in range(right, left-1, -1):
          temp_list.appendleft(arr[i])
          result.append(list(temp list))
   def main():
      print(find_subarrays([2, 5, 3, 10], 30))
      print(find subarrays([8, 2, 6, 5], 50))
Run
                                                                                                            ::3
```

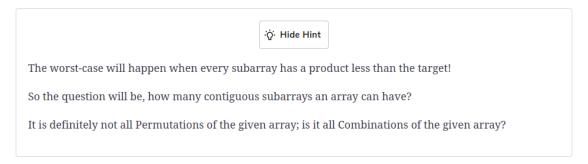
Time complexity

The main for-loop managing the sliding window takes O(N) but creating subarrays can take up to $O(N^2)$ in the worst case. Therefore overall, our algorithm will take $O(N^3)$.

Space complexity

Ignoring the space required for the output list, the algorithm runs in O(N) space which is used for the temp list.

Can you try estimating how much space will be required for the output list?



It is not all the Combinations of all elements of the array!

For an array with distinct elements, finding all of its contiguous subarrays is like finding the number of ways to choose two indices, \mathbf{i} and \mathbf{j} , in the array such that $\mathbf{i} \leftarrow \mathbf{j}$.

If there are a total of n elements in the array, here is how we can count all the contiguous subarrays:

- When i = 0, j can have any value from 0 to n-1, giving a total of n choices.
- When i = 1, j can have any value from 1 to n-1, giving a total of n-1 choices.
- Similarly, when i = 2, j can have n-2 choices.

...

• When i - n-1 i can only have only 1 choice

- which I - HI, J can only have only I choice.

Let's combine all the choices:

Which gives us a total of: n*(n+1)/2

So, at most, we need space for $O(n^2)$ output lists. At worst, each subarray can take O(n) space, so overall, our algorithm's space complexity will be $O(n^3)$.

