

Solution Review: Problem Challenge 1

We'll cover the following

- Quadruple Sum to Target (medium)
- Solution
 - Code
 - Time complexity
 - Space complexity

Quadruple Sum to Target (medium)

Given an array of unsorted numbers and a target number, find all **unique quadruplets** in it, whose **sum is equal to the target number**.

Example 1:

```
Input: [4, 1, 2, -1, 1, -3], target=1
Output: [-3, -1, 1, 4], [-3, 1, 1, 2]
Explanation: Both the quadruplets add up to the target.
```

Example 2:

```
Input: [2, 0, -1, 1, -2, 2], target=2
Output: [-2, 0, 2, 2], [-1, 0, 1, 2]
Explanation: Both the quadruplets add up to the target.
```

Solution

This problem follows the **Two Pointers** pattern and shares similarities with [Triplet Sum to Zero](#).

We can follow a similar approach to iterate through the array, taking one number at a time. At every step during the iteration, we will search for the quadruplets similar to [Triplet Sum to Zero](#) whose sum is equal to the given target.

Code

Here is what our algorithm will look like:

```
Java Python3 C++ JS
1 def search_quadruplets(arr, target):
2     arr.sort()
3     quadruplets = []
4     for i in range(0, len(arr)-3):
5         # skip same element to avoid duplicate quadruplets
6         if i > 0 and arr[i] == arr[i - 1]:
7             continue
8         for j in range(i + 1, len(arr)-2):
9             # skip same element to avoid duplicate quadruplets
10            if j > i + 1 and arr[j] == arr[j - 1]:
11                continue
12            search_pairs(arr, target, i, j, quadruplets)
13    return quadruplets
14
15
16 def search_pairs(arr, target_sum, first, second, quadruplets):
17     left = second + 1
18     right = len(arr) - 1
19     while (left < right):
20         quad_sum = arr[first] + arr[second] + arr[left] + arr[right]
21         if quad_sum == target_sum: # found the quadruplet
```

```
21     if quad_sum == target_sum: # found the quadruplet
22         quadruplets.append(
23             [arr[first], arr[second], arr[left], arr[right]])
24         left += 1
25         right -= 1
26     while (left < right and arr[left] == arr[left - 1]):
27         left += 1 # skip same element to avoid duplicate quadruplets
28     while (left < right and arr[right] == arr[right + 1]):
29         right -= 1 # skip same element to avoid duplicate quadruplets
30     elif quad_sum < target_sum:
31         left += 1 # we need a pair with a bigger sum
32     else:
33         right -= 1 # we need a pair with a smaller sum
34
35
36 def main():
37     print(search_quadruplets([4, 1, 2, -1, 1, -3], 1))
38     print(search_quadruplets([2, 0, -1, 1, -2, 2], 2))
39
40
41 main()
42
```

Run

Save

Reset

Time complexity

Sorting the array will take $O(N * \log N)$. Overall `searchQuadruplets()` will take $O(N * \log N + N^3)$, which is asymptotically equivalent to $O(N^3)$.

Space complexity

The space complexity of the above algorithm will be $O(N)$ which is required for sorting.

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Problem Challenge 2

✓ Completed

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