

Tasks Scheduling (medium)

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Problem Statement

There are 'N' tasks, labeled from '0' to 'N-1'. Each task can have some prerequisite tasks which need to be completed before it can be scheduled. Given the number of tasks and a list of prerequisite pairs, find out if it is possible to schedule all the tasks.

Example 1:

```
Input: Tasks=3, Prerequisites=[0, 1], [1, 2]
Output: true
Explanation: To execute task '1', task '0' needs to finish first. Similarly, task '1' needs to finish before '2' can be scheduled. A possible scheduling of tasks is: [0, 1, 2]
```

Example 2:


```
Input: Tasks=3, Prerequisites=[0, 1], [1, 2], [2, 0]
Output: false
Explanation: The tasks have cyclic dependency, therefore they cannot be scheduled.
```


Example 3:


```
Input: Tasks=6, Prerequisites=[2, 5], [0, 5], [0, 4], [1, 4], [3, 2], [1, 3]
Output: true
Explanation: A possible scheduling of tasks is: [0 1 4 3 2 5]
```


Try it yourself

Try solving this question here:

 Java

 Python3

 JS

 C++

```
1 def is_scheduling_possible(tasks, prerequisites):
2     # TODO: Write your code here
3     return False
4
5
6 def main():
7     print("Is scheduling possible: " +
8         str(is_scheduling_possible(3, [[0, 1], [1, 2]])))
9     print("Is scheduling possible: " +
10         str(is_scheduling_possible(3, [[0, 1], [1, 2], [2, 0]])))
11     print("Is scheduling possible: " +
12         str(is_scheduling_possible(6, [[0, 4], [1, 4], [3, 2], [1, 3]])))
13
14 main()
15
```

Run

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Reset



Solution

This problem is asking us to find out if it is possible to find a topological ordering of the given tasks. The tasks are equivalent to the vertices and the prerequisites are the edges.

We can use a similar algorithm as described in [Topological Sort](#) to find the topological ordering of the tasks. If the ordering does not include all the tasks, we will conclude that some tasks have cyclic dependencies.

Code

Here is what our algorithm will look like (only the highlighted lines have changed):

Java Python3 C++ JS

```
1 from collections import deque
2
3
4 def is_scheduling_possible(tasks, prerequisites):
5     sortedOrder = []
6     if tasks <= 0:
7         return False
8
9     # a. Initialize the graph
10    inDegree = {i: 0 for i in range(tasks)} # count of incoming edges
11    graph = {i: [] for i in range(tasks)} # adjacency list graph
12
13    # b. Build the graph
14    for prerequisite in prerequisites:
15        parent, child = prerequisite[0], prerequisite[1]
16        graph[parent].append(child) # put the child into it's parent's list
17        inDegree[child] += 1 # increment child's inDegree
18
19    # c. Find all sources i.e., all vertices with 0 in-degrees
20    sources = deque()
21    for key in inDegree:
22        if inDegree[key] == 0:
23            sources.append(key)
24
25    # d. For each source, add it to the sortedOrder and subtract one from all of its children's in-degrees
26    # if a child's in-degree becomes zero, add it to the sources queue
27    while sources:
28        vertex = sources.popleft()
29        sortedOrder.append(vertex)
30        for child in graph[vertex]: # get the node's children to decrement their in-degrees
31            inDegree[child] -= 1
32            if inDegree[child] == 0:
33                sources.append(child)
34
35    # if sortedOrder doesn't contain all tasks, there is a cyclic dependency between tasks, therefore, we
36    # will not be able to schedule all tasks
37    return len(sortedOrder) == tasks
38
39
40 def main():
41     print("Is scheduling possible: " +
42           str(is_scheduling_possible(3, [[0, 1], [1, 2]])))
43     print("Is scheduling possible: " +
44           str(is_scheduling_possible(3, [[0, 1], [1, 2], [2, 0]])))
45     print("Is scheduling possible: " +
46           str(is_scheduling_possible(6, [[0, 4], [1, 4], [3, 2], [1, 3]])))
47
48 main()
49
```

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Time complexity

In step 'd', each task can become a source only once and each edge (prerequisite) will be accessed and removed once. Therefore, the time complexity of the above algorithm will be $O(V + E)$, where 'V' is the total number of tasks and 'E' is the total number of prerequisites.

number of tasks and E is the total number of prerequisites.

Space complexity

The space complexity will be $O(V + E)$, since we are storing all of the prerequisites for each task in an adjacency list.

Similar Problems

Course Schedule: There are 'N' courses, labeled from '0' to 'N-1'. Each course can have some prerequisite courses which need to be completed before it can be taken. Given the number of courses and a list of prerequisite pairs, find if it is possible for a student to take all the courses.

Solution: This problem is exactly similar to our parent problem. In this problem, we have courses instead of tasks.

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Topological Sort (medium)

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Tasks Scheduling Order (medium)

✓ Completed

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