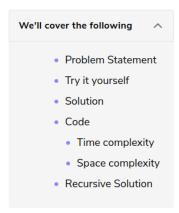




Unique Generalized Abbreviations (hard)



Problem Statement

Given a word, write a function to generate all of its unique generalized abbreviations.

Generalized abbreviation of a word can be generated by replacing each substring of the word by the count of characters in the substring. Take the example of "ab" which has four substrings: "", "a", "b", and "ab". After replacing these substrings in the actual word by the count of characters we get all the generalized abbreviations: "ab", "1b", "a1", and "2".

Example 1:

```
Input: "BAT"
Output: "BAT", "BA1", "B1T", "B2", "1AT", "1A1", "2T", "3"
```

Example 2:

```
Input: "code"
Output: "code", "cod1", "co1e", "co2", "c1de", "c1d1", "c2e", "c3", "1ode", "1od1", "1o1e", "1o
2",
   "2de", "2d1", "3e", "4"
```

Try it yourself

Try solving this question here:

```
# Topo: Write your code here
return result

def main():
    print("Generalized abbreviation are: " +
    | str(generate generalized abbreviation("BAT")))
    print("Generalized abbreviation are: " +
    | str(generate generalized abbreviation("Code")))
    | str(generate generalized abbreviation("code")))

Run

Run

Save Reset (3)
```

Solution

This problem follows the Subsets pattern and can be mapped to Balanced Parentheses. We can follow a similar BFS approach.

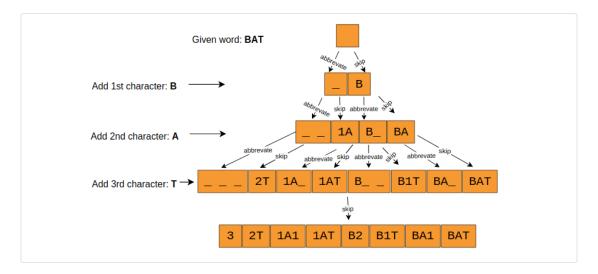
Let's take Example-1 mentioned above to generate all unique generalized abbreviations. Following a BFS approach, we will abbreviate one character at a time. At each step we have two options:

- · Abbreviate the current character, or
- Add the current character to the output and skip abbreviation.

Following these two rules, let's abbreviate BAT:

- 1. Start with an empty word: ""
- 2. At every step, we will take all the combinations from the previous step and apply the two abbreviation rules to the next character.
- 3. Take the empty word from the previous step and add the first character to it. We can either abbreviate the character or add it (by skipping abbreviation). This gives us two new words: ___, _B.
- 4. In the next iteration, let's add the second character. Applying the two rules on _ will give us _ _ and _ _ and _ _ _ and _ _ B gives us _ _ and _ BA .
- 5. The next iteration will give us: _ _ _ , 2T , 1A_ , 1AT , B _ _ , B1T , BA_ , BAT
- 6. The final iteration will give us: 3, 2T, 1A1, 1AT, B2, B1T, BA1, BAT

Here is the visual representation of this algorithm:



Code

Here is what our algorithm will look like:

```
from collections import deque

class AbbreviatedWord:

def __init__(self, str, start, count):
    self.str = str
    self.start = start
    self.count = count

def generate_generalized_abbreviation(word):
    wordLen = len(word)
    result = []
    queue = deque()
    queue.append(AbbreviatedWord(list(), 0, 0))
    while queue:
```

```
abWord = queue.popleft()
        if abWord.start == wordLen:
         if abWord.count != 0:
           abWord.str.append(str(abWord.count))
         result.append(''.join(abWord.str))
         queue.append(AbbreviatedWord(list(abWord.str),
                                      abWord.start + 1, abWord.count + 1))
         if abWord.count != 0:
           abWord.str.append(str(abWord.count))
         newWord = list(abWord.str)
         newWord.append(word[abWord.start])
         queue.append(AbbreviatedWord(newWord, abWord.start + 1, 0))
     return result
     print("Generalized abbreviation are: " +
           str(generate_generalized_abbreviation("BAT")))
           str(generate_generalized_abbreviation("code")))
   main()
Run
                                                                                              Reset []
```

Time complexity

Since we had two options for each character, we will have a maximum of 2^N combinations. If you see the visual representation of Example-1 closely you will realize that it is equivalent to a binary tree, where each node has two children. This means that we will have 2^N leaf nodes and $2^N - 1$ intermediate nodes, so the total number of elements pushed to the queue will be $2^N + 2^N - 1$, which is asymptotically equivalent to $O(2^N)$. While processing each element, we do need to concatenate the current string with a character. This operation will take O(N), so the overall time complexity of our algorithm will be $O(N*2^N)$.

Space complexity

All the additional space used by our algorithm is for the output list. Since we can't have more than $O(2^N)$ combinations, the space complexity of our algorithm is $O(N*2^N)$.

Recursive Solution

Here is the recursive algorithm following a similar approach:

```
🚣 Java
           Python3
                        ⓒ C++
                                    Js JS
        generate_generalized_abbreviation(word):
      generate abbreviation recursive(word, list(), 0, 0, result)
      return result
    def generate abbreviation recursive(word, abWord, start, count, result):
      if start == len(word):
        if count != 0:
          abWord.append(str(count))
        result.append(''.join(abWord))
        generate abbreviation recursive(
          word, list(abWord), start + 1, count + 1, result)
         abWord.append(str(count))
        newWord = list(abWord)
        newWord.append(word[start])
```

