

## Solution Review: Problem Challenge 4

### We'll cover the following ^

- Words Concatenation (hard)
- Solution
- Code
  - Time Complexity
  - Space Complexity

## Words Concatenation (hard) #

Given a string and a list of words, find all the starting indices of substrings in the given string that are a **concatenation of all the given words** exactly once **without any overlapping** of words. It is given that all words are of the same length.

### Example 1:

```
Input: String="catfoxcat", Words=["cat", "fox"]
Output: [0, 3]
Explanation: The two substring containing both the words are "catfox" & "foxcat".
```

### Example 2:

```
Input: String="catcatfoxfox", Words=["cat", "fox"]
Output: [3]
Explanation: The only substring containing both the words is "catfox".
```

## Solution #

This problem follows the **Sliding Window** pattern and has a lot of similarities with [Maximum Sum Subarray of Size K](#). We will keep track of all the words in a **HashMap** and try to match them in the given string. Here are the set of steps for our algorithm:

1. Keep the frequency of every word in a **HashMap**.
2. Starting from every index in the string, try to match all the words.
3. In each iteration, keep track of all the words that we have already seen in another **HashMap**.
4. If a word is not found or has a higher frequency than required, we can move on to the next character in the string.
5. Store the index if we have found all the words.

## Code #

Here is what our algorithm will look like:

Java Python3 C++ JS

```
1 def find_word_concatenation(str1, words):
2     if len(words) == 0 or len(words[0]) == 0:
3         return []
4
5     word_frequency = {}
6
7     for word in words:
8         if word not in word_frequency:
9             word_frequency[word] = 0
```

```

10 | word_frequency[word] += 1
11 |
12 | result_indices = []
13 | words_count = len(words)
14 | word_length = len(words[0])
15 |
16 | for i in range((len(str1) - words_count * word_length)+1):
17 |     words_seen = {}
18 |     for j in range(0, words_count):
19 |         next_word_index = i + j * word_length
20 |         # Get the next word from the string
21 |         word = str1[next_word_index: next_word_index + word_length]
22 |         if word not in word_frequency: # Break if we don't need this word
23 |             break
24 |
25 |         # Add the word to the 'words_seen' map
26 |         if word not in words_seen:
27 |             words_seen[word] = 0
28 |         words_seen[word] += 1
29 |
30 |         # No need to process further if the word has higher frequency than required
31 |         if words_seen[word] > word_frequency.get(word, 0):
32 |             break
33 |
34 |         if j + 1 == words_count: # Store index if we have found all the words
35 |             result_indices.append(i)
36 |
37 | return result_indices
38 |
39 |
40 | def main():
41 |     print(find_word_concatenation("catfoxcat", ["cat", "fox"]))
42 |     print(find_word_concatenation("catcatfoxfox", ["cat", "fox"]))
43 |
44 |
45 | main()
46 |

```

Run

Save

Reset



## Time Complexity #

The time complexity of the above algorithm will be  $O(N * M * Len)$  where 'N' is the number of characters in the given string, 'M' is the total number of words, and 'Len' is the length of a word.

## Space Complexity #

The space complexity of the algorithm is  $O(M)$  since at most, we will be storing all the words in the two **HashMaps**. In the worst case, we also need  $O(N)$  space for the resulting list. So, the overall space complexity of the algorithm will be  $O(M + N)$ .

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✓ Completed

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