

Conflicting Appointments (medium)

We'll cover the following ^

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Problem Statement

Given an array of intervals representing 'N' appointments, find out if a person can **attend all the appointments**.

Example 1:

```
Appointments: [[1,4], [2,5], [7,9]]
Output: false
Explanation: Since [1,4] and [2,5] overlap, a person cannot attend both of these appointments.
```

Example 2:


```
Appointments: [[6,7], [2,4], [8,12]]
Output: true
Explanation: None of the appointments overlap, therefore a person can attend all of them.
```


Example 3:


```
Appointments: [[4,5], [2,3], [3,6]]
Output: false
Explanation: Since [4,5] and [3,6] overlap, a person cannot attend both of these appointments.
```


Try it yourself

Try solving this question here:

 Java

 Python3

 JS


 C++

```
1 def can_attend_all_appointments(intervals):
2     # TODO: Write your code here
3     return False
4
5
6 def main():
7     print("Can attend all appointments: " + str(can_attend_all_appointments([[1, 4], [2, 5], [7, 9]])))
8     print("Can attend all appointments: " + str(can_attend_all_appointments([[6, 7], [2, 4], [8, 12]])))
9     print("Can attend all appointments: " + str(can_attend_all_appointments([[4, 5], [2, 3], [3, 6]])))
10
11
12 main()
13
```

Run

Save

Reset



Solution

The problem follows the [Merge Intervals](#) pattern. We can sort all the intervals by start time and then check if any two intervals overlap. A person will not be able to attend all appointments if any two appointments overlap.

Code

Here is what our algorithm will look like:

Java Python3 C++ JS

```
1 def can_attend_all_appointments(intervals):
2     intervals.sort(key=lambda x: x[0])
3     start, end = 0, 1
4     for i in range(1, len(intervals)):
5         if intervals[i][start] < intervals[i-1][end]:
6             # please note the comparison above, it is "<" and not "<="
7             # while merging we needed "<=" comparison, as we will be merging the two
8             # intervals having condition "intervals[i][start] == intervals[i - 1][end]" but
9             # such intervals don't represent conflicting appointments as one starts right
10            # after the other
11            return False
12    return True
13
14
15 def main():
16     print("Can attend all appointments: " + str(can_attend_all_appointments([[1, 4], [2, 5], [7, 9]])))
17     print("Can attend all appointments: " + str(can_attend_all_appointments([[6, 7], [2, 4], [8, 12]])))
18     print("Can attend all appointments: " + str(can_attend_all_appointments([[4, 5], [2, 3], [3, 6]])))
19
20
21 main()
22
```

Run Save Reset

Time complexity

The time complexity of the above algorithm is $O(N * \log N)$, where 'N' is the total number of appointments. Though we are iterating the intervals only once, our algorithm will take $O(N * \log N)$ since we need to sort them in the beginning.

Space complexity

The space complexity of the above algorithm will be $O(N)$, which we need for sorting. For Java, `Arrays.sort()` uses [Timsort](#), which needs $O(N)$ space.

Similar Problems

Problem 1: Given a list of appointments, find all the conflicting appointments.

Example:

```
Appointments: [[4,5], [2,3], [3,6], [5,7], [7,8]]
Output:
[4,5] and [3,6] conflict.
[3,6] and [5,7] conflict.
```

