

## Triplet Sum Close to Target (medium)

### We'll cover the following ^

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## Problem Statement #

Given an array of unsorted numbers and a target number, find a **triplet in the array whose sum is as close to the target number as possible**, return the sum of the triplet. If there are more than one such triplet, return the sum of the triplet with the smallest sum.

Example 1:

```
Input: [-2, 0, 1, 2], target=2
Output: 1
Explanation: The triplet [-2, 1, 2] has the closest sum to the target.
```

Example 2:

```
Input: [-3, -1, 1, 2], target=1
Output: 0
Explanation: The triplet [-3, 1, 2] has the closest sum to the target.
```

Example 3:

```
Input: [1, 0, 1, 1], target=100
Output: 3
Explanation: The triplet [1, 1, 1] has the closest sum to the target.
```

## Try it yourself #

Try solving this question here:

Java

Python3

JS

C++

```
1 def triplet_sum_close_to_target(arr, target_sum):
2     # TODO: Write your code here
3     return -1
4
```

Test

Save

Reset

## Solution #

This problem follows the **Two Pointers** pattern and is quite similar to [Triplet Sum to Zero](#).

We can follow a similar approach to iterate through the array, taking one number at a time. At every step, we will save the difference between the triplet and the target number, so that in the end, we can return the triplet with the closest sum.

## Code #

Here is what our algorithm will look like:

Java Python3 C++ JS

```
1 import math
2
3
4 def triplet_sum_close_to_target(arr, target_sum):
5     arr.sort()
6     smallest_difference = math.inf
7     for i in range(len(arr)-2):
8         left = i + 1
9         right = len(arr) - 1
10        while (left < right):
11            target_diff = target_sum - arr[i] - arr[left] - arr[right]
12            if target_diff == 0: # we've found a triplet with an exact sum
13                return target_sum - target_diff # return sum of all the numbers
14
15            # the second part of the following 'if' is to handle the smallest sum when we have more than one so
16            if abs(target_diff) < abs(smallest_difference) or (abs(target_diff) == abs(smallest_difference) and
17                smallest_difference = target_diff # save the closest and the biggest difference
18
19            if target_diff > 0:
20                left += 1 # we need a triplet with a bigger sum
21            else:
22                right -= 1 # we need a triplet with a smaller sum
23
24        return target_sum - smallest_difference
25
26
27 def main():
28     print(triplet_sum_close_to_target([-2, 0, 1, 2], 2))
29     print(triplet_sum_close_to_target([-3, -1, 1, 2], 1))
30     print(triplet_sum_close_to_target([1, 0, 1, 1], 100))
31
32
33 main()
34
```

Run Save Reset

## Time complexity #

Sorting the array will take  $O(N * \log N)$ . Overall, the function will take  $O(N * \log N + N^2)$ , which is asymptotically equivalent to  $O(N^2)$ .

## Space complexity #

The above algorithm's space complexity will be  $O(N)$ , which is required for sorting.

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Triplet Sum to Zero (medium)Triplets with Smaller Sum (medium)

✓ Completed