U1M4.LW.Access and Join Methods Part 1 Report Mastykina Elizaweta

1. Table access full scan

1.1. Task 1: Full Scans and the High-water Mark and Block reading

Creation of the table t2:

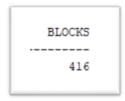
```
CREATE TABLE t2 AS
SELECT TRUNC( rownum / 100 ) id, rpad( rownum,100 ) t_pad
FROM dual
CONNECT BY rownum < 100000;</pre>
```

Creation of the index:

```
CREATE INDEX t2_idx1 ON t2
  ( id );
```

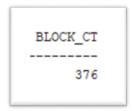
Allocated Blocks Count:

```
SELECT BLOCKS FROM user_segments WHERE segment_name =
'T2';
```



Used Blocks Count:

```
SELECT COUNT(Distinct
(dbms rowid.rowid block number(rowid))) block ct FROM t2;
```



Execution of the full table scan:

```
SET autotrace ON;
 SELECT COUNT ( * )
 FROM t2 ;
 SET autotrace OFF;
Statistics
_____
            1 CPU used by this session
             3 CPU used when call started
             3 DB time
            33 Requests to/from client
            33 SQL*Net roundtrips to/from client
             2 buffer is not pinned count
           523 bytes received via SQL*Net from client
          60463 bytes sent via SQL*Net to client
             2 calls to get snapshot scn: kcmgss
             6 calls to kcmgcs
           380 consistent gets
           380 consistent gets from cache
           380 consistent gets pin
           380 consistent gets pin (fastpath)
             2 cursor authentications
             2 execute count
       12451840 logical read bytes from cache
           376 no work - consistent read gets
            42 non-idle wait count
             2 opened cursors cumulative
             1 opened cursors current
             2 parse count (total)
           380 session logical reads
             1 sorts (memory)
          1581 sorts (rows)
           376 table scan blocks gotten
          99999 table scan disk non-IMC rows gotten
          99999 table scan rows gotten
             1 table scans (short tables)
```

Autotrace Disabled

34 user calls

```
COUNT(*)

99999

PLAN_TABLE_OUTPUT

SQL_ID 2tn5lubvpkphz, child number 0

SELECT COUNT(*) FROM t2

Plan hash value: 3321871023

| Id | Operation | Name | E-Rows |

| 0 | SELECT STATEMENT | | | |

| 1 | SORT AGGREGATE | | 1 |

PLAN_TABLE_OUTPUT

| 2 | TABLE ACCESS FULL | T2 | 99999 |
```

Delete All Rows from table and repeat previous steps:

```
DELETE FROM t2;

SELECT COUNT(Distinct
(dbms_rowid.rowid_block_number(rowid))) block_ct FROM t2;

SET autotrace ON;

SELECT COUNT( * )
 FROM t2;

SET autotrace OFF;
```

```
BLOCK CT
Autotrace Enabled
Shows the execution plan as well as statistics of the statement.
COUNT (*)
PLAN TABLE OUTPUT
SQL_ID 2tn5lubvpkphz, child number 0
-----
SELECT COUNT ( * ) FROM t2
Plan hash value: 3321871023
               | Name | E-Rows |
| Id | Operation
_____
1 |
PLAN TABLE OUTPUT
 2 | TABLE ACCESS FULL| T2 | 99999 |
Statistics
______
          1 CPU used by this session
          1 CPU used when call started
          6 Requests to/from client
         380 consistent gets
         380 consistent gets from cache
         380 consistent gets pin
         380 consistent gets pin (fastpath)
           7 non-idle wait count
          2 opened cursors cumulative
          1 opened cursors current
          1 pinned cursors current
         380 session logical reads
          7 user calls
```

Autotrace Disabled

Insert 1 row and execute a full table scan and noted the consistent gets (logical block reads):

```
INSERT INTO t2
 ( ID, T_PAD )
 VALUES
 ( 1,'1');
COMMIT;
 SELECT COUNT (Distinct
(dbms rowid.rowid block number(rowid))) block ct FROM t2;
 SET autotrace ON;
 SELECT COUNT ( * )
  FROM t2 ;
 SET autotrace OFF;
 COUNT (*)
PLAN_TABLE_OUTPUT
______
SQL_ID 2tn5lubvpkphz, child number 0
 SELECT COUNT ( * ) FROM t2
Plan hash value: 3321871023
| Id | Operation | Name | E-Rows |
PLAN TABLE OUTPUT
______
| 2 | TABLE ACCESS FULL| T2 | 99999 |
```

1 CPU used by this session 1 CPU used when call started 1 DB time 6 Requests to/from client 380 consistent gets 380 consistent gets from cache 380 consistent gets pin 380 consistent gets pin 380 consistent gets pin 280 consistent gets pin (fastpath) 7 non-idle wait count 2 opened cursors cumulative 1 opened cursors current 1 pinned cursors current 380 session logical reads 7 user calls

Autotrace Disabled

Truncate Table and execute a full table scan and noted the consistent gets (logical block reads):

```
TRUNCATE TABLE t2;
SELECT COUNT(Distinct
(dbms_rowid.rowid_block_number(rowid))) block_ct FROM t2;

SET autotrace ON;

SELECT COUNT( * )
  FROM t2;

SET autotrace OFF;
```

```
BLOCK CT
Autotrace Enabled
Shows the execution plan as well as statistics of the statement.
 COUNT (*)
PLAN TABLE OUTPUT
______
SQL_ID 2tn5lubvpkphz, child number 0
 SELECT COUNT( * ) FROM t2
Plan hash value: 3321871023
| Id | Operation | Name | E-Rows |
_____
1 |
PLAN_TABLE_OUTPUT
| 2 | TABLE ACCESS FULL| T2 | 99999 |
Statistics
______
          1 DB time
          6 Requests to/from client
          1 consistent gets
          1 consistent gets from cache
          1 consistent gets pin
          1 consistent gets pin (fastpath)
          1 enqueue releases
           1 enqueue requests
          7 non-idle wait count
          2 opened cursors cumulative
          1 opened cursors current
          1 pinned cursors current
           1 recursive calls
           1 session logical reads
          7 user calls
```

Table T2 truncated.

Autotrace Disabled

Nº	Count of Blocks	Count of Used Blocks	Count of Rows	Consistent gets	Description
1-3 (After creation)	416	376	99999	380	The HVM remains the same, all blocks up to the HWM are read in and scanned, even if they are empty
4-5 (After deletion)	416	0	0	380	
6-7 (After insertion)	416	1	1	380	
8-9 (After truncation)	6	0	0	1	The space was deallocated and the HWM was reset

This access method implies

iterate over all rows of the table, excluding those that do not satisfy where predicate. It is applied either when

predicate conditions are absent in the index, or when the index does not exist.

2. Index Scan types

2.1. Task 2: Index Clustering factor parameter

Created table t2:

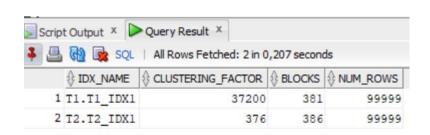
```
CREATE TABLE t2 AS
SELECT TRUNC( rownum / 100 ) id, rpad( rownum, 100 ) t_pad
FROM dual
CONNECT BY rownum < 100000;</pre>
```

Created index:

```
CREATE INDEX t2_idx1 ON t2
  ( id );
```

Created table t1:

```
CREATE TABLE t1 AS
SELECT MOD( rownum, 100 ) id, rpad( rownum, 100 ) t pad
 FROM dual
CONNECT BY rownum < 100000;
Created index:
    CREATE INDEX t1 idx1 ON t1 ( id );
Calculate statistic for both tables:
EXEC
       dbms stats.gather table stats( USER,'t1',method opt=>'FOR
A T.T.
COLUMNS SIZE 1', CASCADE=>TRUE );
EXEC
      dbms stats.gather table stats( USER,'t2',method opt=>'FOR
ALL
COLUMNS SIZE 1', CASCADE=>TRUE );
Computed Index Clustering factor:
SELECT t.table name | | '.' | | i.index name idx name,
 i.clustering factor,
 t.blocks,
 t.num rows
FROM user indexes i, user tables t
```



WHERE i.table name = t.table name

AND t.table name IN('T1', 'T2');

The clustering factor is a number that represents the degree to which data is randomly distributed in a table as compared to the indexed column.

The table rows are synchronized with the index when the clustering factor is close to the number of data blocks and the column value is not row-ordered when the clustering factor approaches the number of rows in the table.

We have different clustering for T1 and T2 tables because they have different structure of column ID. Our index is based on column ID and we can see that in T2 rows with the same id are next to each other, when in T1 they are distributed.

2.2. Task 3: Index Unique Scan

Created index:

o rows selected

```
CREATE UNIQUE INDEX udx_t1 ON t1( t_pad );

SELECT t1.* FROM t1 where t1.t_pad = '1';

NDEX UDX_T1 created.
```

Output is empty because in select statement '1' is string but t_pad type is int

The index structure goes from the root to leaf block to a single entry, gets a rowid, which is used to access the table data block containing one row. The TABLE ACCESS BY INDEX ROWID step in the plan specifies access to a table data block.

2.3. Task 4: Index Range Scan

```
SELECT t2.* FROM t2 where t2.id = '1';
```

```
■ SQL | All Rows Fetched: 100 in 0,045 seconds

I 100

1 1100

2 1101

3 1102

4 1103

5 1104

6 1105

7 1106

8 1107

9 1108

10 1109
```

```
90 1 189

91 1 190

92 1 191

93 1 192

94 1 193

95 1 194

96 1 195

97 1 196

98 1 197

99 1 198
```

A range scan goes through the index structure from the root block to the first leaf block containing an entry that matches the given condition. From this starting point, the rowid is retrieved from the index entry and the table data block (TABLE ACCESS BY INDEX ROWID) is retrieved. After the first row is retrieved, the leaf block of the index is accessed again and the next entry is read to obtain the next row ID. This exchange of data between leaf index blocks and data blocks continues until all matching index entries have been read.

2.4. Task 5: Index Skip Scan

Creation of the table employees:

```
CREATE TABLE employees AS

SELECT *
```

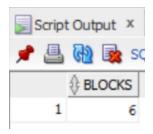
```
FROM scott.emp;
```

Creation of the index:

```
CREATE INDEX idx_emp01 ON employees
  ( empno, ename, job );
```

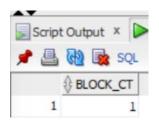
Listed number of allocated blocks and how many blocks contain data:

SELECT BLOCKS FROM user segments WHERE segment name = 'EMP';



SELECT COUNT(DISTINCT (dbms_rowid.rowid_block_number(rowid)))
block_ct

FROM emp ;

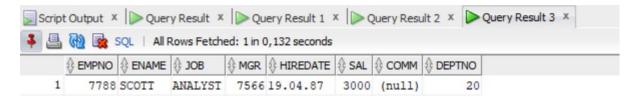


Get trace and statistic of explain plan

SELECT /*+INDEX_SS(emp idx_emp01)*/ emp.* FROM employees
emp where ename = 'SCOTT';



SELECT /*+FULL*/ emp.* FROM employees emp WHERE ename =
'SCOTT';



Nº	Count of Blocks	Count of Used Blocks	Count of Rows	Consistent gets	Description
Full scan	6	1	12	8	Oracle performed SELECT (using the hint /* +FULL(emp)*/) and used the Full Scan algorithm
Index skip scan	6	1	12	2	Oracle performed SELECT (using the hint /*+INDEX_SS(emp idx_emp01)*/) and used the Index Scip Scan algorithm
Without hint	6	1	12	5	Oracle performed SELECT (without using hints) and used the most appropriate Full Scan

The optimizer will tend not to use the index since our predicate did not involve the column X—it might have to inspect each and every index entry in this case (we'll discuss an index skip scan shortly where this is not true). It will typically opt for a full table scan of T instead. That does not preclude the index from being used. If the query was SELECT X,Y FROM T WHERE Y = 5, the optimizer would notice that it did not have to go to the table to get either X or Y (they are in the index) and may very well opt for a fast full scan of the index itself, as the index is typically much smaller than the underlying table.