

# Solutions to Extremal Combinatorics

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## **Abstract**

The following document contains solutions to selected problems in the book ‘Extremal Combinatorics’ by Stasys Jukna. Some problems which are either trivial or require nothing more than brute force/proof mirroring have been omitted.

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# 1 Counting

## 1.1 Review

- The binomial theorem and binomial coefficients provide a way to count objects.
- Exact values of binomial coefficients are hard to compute, but they can be roughly estimated. Consider the following inequalities:

1.  $1 + t < e^t$  for  $t \neq 0$
2.  $1 - t > e^{-t-t^2/2}$  for  $0 < t < 1$

they can be used to prove

$$\left(\frac{n}{k}\right)^k \leq \binom{n}{k} \quad \sum_{i=1}^k \binom{n}{i} \leq \left(\frac{en}{k}\right)^k.$$

- The Stirling approximation is another approximation which shows

$$n! \approx \left(\frac{n}{e}\right)^n \sqrt{2\pi n} e^{\alpha_n}$$

with  $1/(12n+1) < \alpha_n < 1/12n$ . Similarly, an asymptotic approximation for the  $k$ -factorial is

$$(n)_k = n^k e^{-k^2/2n - k^3/6n^2 + o(1)}$$

which is valid for  $k = o(n^{3/4})$ .

- The stars and bars technique allows you to partition a set into ordered partitions.
- Double counting can be done using incidence matrices: in this technique you count the rows and the columns of a zero-one matrix, which should yield the same sums.
- The number of vertices in a graph with odd degree is even.
- Jensen's Inequality: a function is convex if

$$f(\lambda a + (1 - \lambda)b) \leq \lambda f(a) + (1 - \lambda)f(b)$$

for  $0 \leq \lambda \leq 1$ . Suppose that each  $\lambda_i$  is in  $[0, 1]$ , and  $\sum \lambda_i = 1$ . If  $f$  is convex then

$$f\left(\sum_{i=1}^n \lambda_i x_i\right) \leq \sum_{i=1}^n \lambda_i f(x_i).$$

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- Inclusion-exclusion provides a way to sum up all the elements in some sets that intersect.
- Derangements are permutations which don't fix any points.

### 1.2 Problems

**In-text.** Prove:

$$\sum_{x \in Y} d(x) = \sum_{A \in \mathcal{F}} |Y \cap A| \text{ for any } Y \subseteq X. \quad (1.1)$$

$$\sum_{x \in X} d(x)^2 = \sum_{x \in \mathcal{F}} \sum_{x \in A} d(x) = \sum_{A \in \mathcal{F}} \sum_{B \in \mathcal{F}} |A \cap B|. \quad (1.2)$$

**Answer.** We can proceed by a counting argument. For the first part, consider the incidence matrix  $M = (m_{x,a})$  of  $\mathcal{F}$ . Adding the rows belonging to the elements of  $x \in Y$  gives us the left hand side. If we 'slice' the matrix in this way, however (removing all the rows belonging to elements  $x \notin Y$ ) then summing via the columns we get the sum of only those elements which are both in  $A$  and in  $Y$ , i.e.  $|A \cap Y|$ .

For the second part, it is enough to notice that if we replace each entry 1 in the incidence matrix with  $d(x)$ , then adding along the rows gives  $\sum_{x \in X} d(x)^2$  while adding along the columns gives the term in the central equality (we add  $d(x)$  for every  $x \in A$ , and then we sum over each of the  $A$ s). The final equality follows from noticing that

$$\begin{aligned} \sum_{x \in \mathcal{F}} \sum_{x \in A} d(x) &= \sum_{A' \in \mathcal{F}} \left( \sum_{x \in A} |A \cap A'| \right) \\ &= \sum_{A \in \mathcal{F}} \sum_{B \in \mathcal{F}} |A \cap B|. \end{aligned}$$

where the second equality follows from substituting the first part.

**Question 1.1** In how many ways can we distribute  $k$  balls to  $n$  boxes so that each box has at most one ball?

**Answer.** Depends on whether  $k \leq n$  or  $k > n$ . If  $k \leq n$ , we can choose  $\binom{n}{k}$  boxes, and the balls can be ordered in  $k!$  ways. it follows that the number of ways is  $\binom{n}{k} k! = (n)_k$ .

If  $k > n$ , then there are  $\binom{k}{n}$  ways we can choose the balls, and we can order them in  $n!$  ways. This gives us the answer of  $(k)_n$ .

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**Question 1.2** Show that for every  $k$  the product of any  $k$  consecutive natural numbers is divisible by  $k!$ .

**Answer.** Let the  $k$  numbers be  $n + 1, \dots, n + k$ . Consider the number  $\binom{n+k}{k}$ . This is the number of ways you can choose  $k$  numbers from  $n + k$  numbers, and is an integer. Expanding, we get

$$\binom{n+k}{k} = \frac{n!}{k!(n-k)!}$$

which clearly shows that  $k!$  divides  $(n+1)(n+2)\dots(n+k)$ .

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**Question 1.3** Show that the number of pairs  $(A, B)$  of distinct subsets of  $\{1, \dots, n\}$  with  $A \subset B$  is  $3^n - 2^n$ .

**Answer.** We can proceed as follows. Select a subset  $S$  of size  $k$ , and then select a subset of  $S$ . The former can be done in  $\binom{n}{k}$  ways, while the latter can be done in  $2^k - 1$  ways. This gives us the sum

$$\sum_{k=0}^n \binom{n}{k} (2^k - 1)$$

The binomial theorem tells us that this is equal to  $(2+1)^n - (1+1)^n = 3^n - 2^n$ .

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**Question 1.4** Show that

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}.$$

**Answer.** We will count the number of ways to choose  $k$  balls from  $n$  balls. Just choosing the  $k$  balls is the left hand side. Alternately, we can choose one ball first, and then choose  $k - 1$  balls from the remaining  $n - 1$  balls. However, this leads to a  $k$ -recounting, since each  $k$  size subset is selected  $k$  times (once when each element is the ‘fixed’ element). We are done.

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**Question 1.6** There is a set of  $2n$  people:  $n$  male and  $n$  female. A good party is a set with the same number of male and female. How many possibilities are there to build such a good party?

**Answer.** For each  $k$  there are  $k$  ways to choose men and  $k$  ways to choose women, so the total number of parties of size  $k$  are  $\binom{n}{k}^2$ . Adding over all the  $k$ 's we get

$$\sum_{i=0}^n \binom{n}{i}^2.$$

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**Question 1.7** Use Proposition 1.3 to show that

$$\sum_{i=0}^r \binom{n+i-1}{i} = \binom{n+r}{r}$$

**Answer.** We expand the RHS using Pascal's identity, then recursively expand one of the terms.

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**Question 1.8** Let  $0 \leq a \leq m \leq n$  be integers. Show that

$$\sum_{i=m}^n \binom{i}{a} = \binom{n+1}{a+1} - \binom{m}{a+1}.$$

**Answer.** Same trick as the previous question, expand  $\binom{n+1}{a+1}$ .

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**Question 1.9** Prove the Cauchy-Vandermonde identity:

$$\binom{p+q}{k} = \sum_{i=0}^k \binom{p}{i} \binom{q}{k-i}.$$

**Answer.** We count twice. The left is selecting  $k$  items from  $p+q$  items. Another way we can count this is to select  $i$  items from  $p$  items and  $k-i$  items from  $q$  items, then add over every  $i$ .

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**Question 1.10** Show that

$$\sum_{k=0}^n \binom{n}{k}^2 = \binom{2n}{n}.$$

**Answer.** Start with the previous problem, then set  $p = q = k = n$ . It follows that

$$\binom{2n}{n} = \sum_{k=0}^n \binom{n}{k} \binom{n}{n-k}.$$

However, we know that to choose  $k$  things, we can choose  $k$  or choose  $n - k$ , so the latter terms are equal.

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**Question 1.11** Prove the following analogy of the binomial theorem for factorials:

$$(x + y)_n = \sum_{k=0}^n \binom{n}{k} (x)_k (y)_{n-k}.$$

**Answer.** Immediately follows from the same consideration as the typical proof of the binomial theorem.

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**Question 1.28** This took me way too fucking long. Here's how you do it:

First, let  $|E|$  be the number of edges in the graph. Note that  $|B|D \geq |E| \geq |A|d$  by definition, since the *minimum* number of edges leaving  $A$  is  $|A|d$ , and the *maximum* number of edges leaving  $B$  is  $|B|D$ . Together with the fact given in the question this gives  $|A|d = |B|D$ . Furthermore, because the number of edges is sandwiched between these two, quantities, we know that  $E = |A|d$ . This tells us that every vertex in  $A$  *must* have degree  $d$ . Now take  $B_0$  to be the set such that every neighbor has  $\alpha D/2$  vertices to  $A_0$ .

It follows that the number of edges leaving  $A_0$  is  $|A_0|d = \alpha|A|d = \alpha|B|D$ . Of these, some go to  $B_0$ , while others go outside of  $B_0$ . It follows that  $|E_{A_0}| = |E_{A_0 \rightarrow B_0}| + |E_{A_0 \rightarrow B \setminus B_0}|$ . We know that  $|E_{A_0 \rightarrow B_0}| \leq \alpha|B_0|D/2$ . Suppose that  $|E_{A_0 \rightarrow B \setminus B_0}| \geq |E_{A_0 \rightarrow B_0}|$ , which implies  $|E_{A_0}| = \alpha|B|D \leq \alpha|B_0|D/2 + \alpha|B_0|D/2 = \alpha|B_0|D$ , which is preposterous since  $|B_0| \leq |B|$ . This proves (iii), that  $|E_{A_0 \rightarrow B \setminus B_0}| \leq |E_{A_0 \rightarrow B_0}|$ , which implies that more than half of the edges leaving  $A_0$  go to  $B_0$ .

The first part now easily follows. The number of edges going from  $A_0$  to  $B_0$  is exactly the number of edges going from  $B_0$  to  $A_0$ , which we know is greater than

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$\alpha|B|D/2$  (which is half the number of edges leaving  $A_0$ ). The edges leaving  $B_0$  include this number *and more*. This means that  $|B_0|D \geq \alpha|B|D/2$ , which gives us (i), ie.  $|B_0| \geq \alpha|B|/2$ .

Very annoying problem. It's not conceptually difficult but there's a number of very shitty moving parts. Happy I solved it, finally.

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**Question 1.29** This was also mildly annoying, but I got it after a few minutes of serious thought. You apply Jensen's inequality to the parameters  $f(x) = x^{t/s}$ ,  $\lambda_i = 1/n$ , and  $x_i = a_i^s$ . The inequality follows from the fact that  $f$  is convex iff  $t \geq s$ .

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**Question 1.37** This is not interesting, but is used later in the book. Really the only thing here is that this is in some sense 'partial' Inclusion-Exclusion: it's written in a misleading way, but it's actually just inclusion-exclusion with some of the terms chopped off.

In the odd cases, the opposite will happen: it'll be greater for odd  $k$  and lesser for even  $k$ .



## 2 Advanced Counting

### 2.1 Review

Some filling-in-the-details here, in order to make sure that I'm completely understanding what's happening.

**Page 29, Prop. 2.8.** The final inequality follows from the averaging principle, but it's not as immediate as meets the eye. We basically use the fact that the elements of  $\mathcal{F}$  form a partition of the vertices of  $G$ , and the fact that every single one of them is a clique. In particular, for any  $X, Y \in \mathcal{F}$ , the edges from  $X$  to  $Y$  are not *complete*, and hence  $X \cup Y$  is not a clique (also because of the independence of  $\alpha$ , none of the subsets are cliques either). This means that this is a complete characterization of *all the cliques* in the graph. We then get that the average size of a clique is

$$\frac{1}{|\mathcal{F}|} \sum_{X_i \in \mathcal{F}} |X_i| = \frac{n}{|\mathcal{F}|}$$

and it follows from the averaging principle there exists some clique with size greater than or equal to the average clique, and hence

$$\omega(G) \geq \frac{n}{|\mathcal{F}|}.$$

### 2.2 Problems

**Question 2.1** Start with the sets  $A_i$  and construct new sets  $A_{i,j} = A_i \cap A_j$  and  $A'_i = A_i \setminus (\cap_{j \neq i} A_j)$ . For each  $i$ ,  $A_i = A'_i \cup (\cup_{j \neq i} [A_i \cap A_j])$ . Then by the union bound we have

$$|A_i| \leq |A'_i| + \sum_{j \neq i} |A_i \cap A_j|$$

and it follows that

$$\sum_{i=1}^m |A_i| \leq \sum_{i=1}^m |A'_i| + \sum_i \sum_{j \neq i} |A_i \cap A_j|$$

Since  $|A_i \cap A_j| \leq t$ , we have the latter half of the RHS to be  $\leq t \binom{m}{2}$ . For the former, we notice that  $A'_i \cap A'_j = \emptyset$  by definition. Then  $\sum_{i=1}^m |A'_i| \leq n$ . We have

$$\sum_{i=1}^m |A_i| \leq n + t \binom{m}{2}.$$