Seminar 3 - Exerciții rezolvate

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1

Cuprins

Teorema lui Master pentru "subtract-and-conquer"

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T(n) = a \cdot T(n-b) + f(n), a, b \in \mathbb{N}^*
T(i) \in \Theta(1), \forall i \in \mathbb{N}^*, i < b^a
f(n) \in O(n^p)
aPentru simplitatea formulelor am folosit și cazul de bază pentru <math>n = 0.
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Notăm $h = \left| \frac{n}{h} \right|$.

$$T(n) = a \cdot T(n - b) \qquad +f(n)$$

$$a \cdot T(n - b) = a^{2} \cdot T(n - 2 \cdot b) \qquad +a \cdot f(n - b)$$

$$a^{2} \cdot T(n - 2 \cdot b) = a^{3} \cdot T(n - 3 \cdot b) \qquad +a^{2} \cdot f(n - 2b)$$
...
$$a^{k} \cdot T(n - k \cdot b) = a^{k+1} \cdot T(n - (k+1) \cdot b) \qquad +a^{k} \cdot f(n - k \cdot b)$$
...
$$a^{h-1} \cdot T(n - (h - 1) \cdot b) = a^{h} \cdot \underbrace{T(n - h \cdot b)}_{\Theta(1)} \qquad +a^{h-1} \cdot f(n - (h - 1)b)$$

$$\Rightarrow T(n) = a^{\lfloor \frac{n}{b} \rfloor} \cdot \Theta(1) + \sum_{k=0}^{h-1} a^{k} \cdot f(n - k \cdot b)$$

$$Q(n) = \sum_{k=0}^{h-1} a^{k} \cdot f(n - k \cdot b) = \sum_{k=0}^{h-1} a^{k} O((n - k \cdot b)^{p}) = \sum_{k=0}^{h-1} a^{k} O(n^{p})$$

$$\sum_{k=0}^{1} a^{i} = \frac{a^{j} - 1}{a - 1}$$

$$\sum_{k=0}^{h-1} a^{k} = \begin{cases} O(1) & a < 1 \text{ deoarece } \lim_{n \to \infty} a^{\lfloor \frac{n}{b} \rfloor} = 0 \\ O(n^{p}) & a > 1 \end{cases}$$

$$Q(n) = \begin{cases} O(n^{p}) & a < 1 \\ O(a^{\lfloor \frac{n}{b} \rfloor} \cdot n^{p}) & a > 1 \end{cases}$$

$$a^{\lfloor \frac{n}{b} \rfloor} \cdot \Theta(1) = \begin{cases} O(1) & a < 1 \\ O(n) & a = 1 \\ O(a^{\lfloor \frac{n}{b} \rfloor} \cdot n^{p}) & a > 1 \end{cases}$$

$$T(n) = \begin{cases} O(n^p) & a < 1 \\ O(n^{p+1}) & a = 1 \\ O(a^{\lfloor \frac{n}{b} \rfloor} \cdot n^p) & a > 1 \end{cases}$$