



**LUDDY**

SCHOOL OF INFORMATICS, COMPUTING, AND ENGINEERING

Intro

Generic types  $\Delta$

Illustrative language

Not a fan of ~~the~~ <sup>how they handle</sup> polymorphic functions  $\text{fbb}$

$\Delta t. \lambda x \in t. x$  has type  $\Delta t. t \rightarrow t$

Notational Preliminaries

$[f/x/x']$  denotes  $\lambda y \in S. \text{if } y=x \text{ then } x' \text{ else } f(y)$

$f \in S \Rightarrow S' \quad x \in S \quad x' \in S'$

$S \Rightarrow S' \equiv S, S'$

Syntax  $\prod_{x \in S} F(x)$  all functions  $f$  whose domain is  $S$  and  $\forall x \in S, f(x) \in F(x)$

$w_1 / w_2$  replace all  $t$  in  $w_1$  with  $w_2$

$R_{QW}$   $Q(x)$  induces type  $w$   
I don't understand (2a)

Not sure I fully understand  $R_{QW}$

Semantics

$B \in W \Rightarrow D^T \Rightarrow D \approx \approx$

type to (type to domain)  
to domain





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$$M_{QW} \in R_{QW} \Rightarrow \prod_{\bar{D} \in \bar{D}'} (Env_Q(\bar{D}) \rightarrow B[W](\bar{D}))$$

$$Env_Q(\bar{D}) = \prod_{x \in V} B[\theta(x)](\bar{D})$$

### Representations

$$rep(D, D') = \{ \langle \phi, \psi \rangle \mid \phi \in D \rightarrow D', \psi \in D' \rightarrow D, \psi \circ \phi \sqsubseteq I_D, \phi \circ \psi \sqsubseteq I_{D'} \}$$

$$\rho = \langle \phi, \psi \rangle \in rep(D, D')$$

$x$  represents  $x'$  according to  $\rho$ , iff

$$x \sqsubseteq \psi(x') \quad \text{or} \quad \phi(x) \sqsubseteq x'$$

### Representation Thm

Sure...

### Full semantics

Category ☹️ or what did we get ourselves into

### Syntactic Manipulations

Bro created Ackermann's function ☹️